401 project notes

## Rough explanation of the paper

## Introduction

The setting: solving 2-D Laplace problems using rational function approximations, plus numerical experiments and examples. It turns out using rational function approximations with exponentially clustered points near singularities gets root-exponential convergence.

Their problem:

$$\Delta u(z) = \nabla^2 u(z) = \left(\frac{\partial^2}{\partial x^2} + \frac{\partial^2}{\partial y^2}\right) u(z) = 0, \ z \in \Omega$$
$$u(z) = h(z), \ z \in \Gamma$$

in a domain  $\Omega$  bounded piece-wise smoothly (with corners) by  $\Gamma$ , with specified boundary data h. This sort of problem comes up a lot in physics: electrostatics, fluid dynamics, heat conduction...

The approach:

$$u(z) \approx \text{Re}[r(z)]$$
  
 $r(z) = \sum_{j=1}^{N_1} \frac{a_j}{z - z_j} + \sum_{j=0}^{N_2} b_j (z - z_*)^j$ 

with poles  $z_j$ . The crux of the method is using exponentially clustered sample points on the boundary with h near corners, along with exponentially clustered poles outside the boundary near the corners.

The structure of the paper consists of: theorems establishing root-exponential convergence with rational approximations, then an algorithm using linear least-squares fitting on the boundary to find coefficients  $a_j$  and  $b_j$ .

## Two Theorems

The theorems involve an analytic function f, while the applications involve a harmonic function u. However, and harmonic function u can be the real part of an analytic function, i.e. f = u + iv.

(\*\*\*pictures, define  $A_{\theta}$ )

(1) Let f be a bounded analytic function in the slit disk  $A_{\pi}$  that satisfies  $f(z) = O(|z|^{\delta})$  as  $z \to 0$  for some  $\delta > 0$ , and let  $\theta \in (0, \pi/2)$  be fixed. Then for some  $0 < \rho < 1$  depending on  $\theta$  but not on f, there exist type (n-1,n) rational functions  $\{r_n\}$ ,  $1 \le n < \infty$ , such that

$$||f - r_n||_{\Omega} = O(e^{-C\sqrt{n}})$$

as  $n \to \infty$  for some C > 0, where  $\Omega = \rho A_{\theta}$ . Moreover, each  $r_n$  can be taken to have simple poles only at

$$\beta_j = -e^{-\sigma j/\sqrt{n}}, \ 0 \le j \le n - 1,$$

where  $\sigma > 0$  is arbitrary.

(2) Let  $\Omega$  be a convex polygon with corners  $w_1, \ldots, w_m$ , and let f be an analytic function in  $\Omega$  that is analytic on the interior of each side segment and can be analytically continued to a disk near each  $w_k$  with a slit along the exterior bisector there. Assume f satisfies  $f(z) - f(w_k) = O(|z - w_k|^{\delta})$  as  $z \to w_k$  for each k for some  $\delta > 0$ . There exist degree n rational functions  $\{r_n\}$ ,  $1 \le n < \infty$  such that

$$||f - r_n||_{\Omega} = O(e^{-C\sqrt{n}})$$

as  $n \to \infty$  for some C > 0. Moreover, each  $r_n$  can be taken to have finite poles only at points exponentially clustered along the exterior bisectors at the corners, with arbitrary clustering parameter  $\sigma$ , as long as the number of poles near each  $w_k$  grows at least in proportion to n as  $n \to \infty$ .

Some extensions: These same results hold for  $\Omega$  bounded by analytic arcs meeting at corners. Additionally, the authors believe these results are valid also for non-convex domains and  $\theta < \pi/2$ . Experiments show that placing poles along exterior bisector is also not necessary.

## Algorithm and Examples

Probably examples and pictures first.