



# An Exploration of a Lightning-Fast Laplace Solver

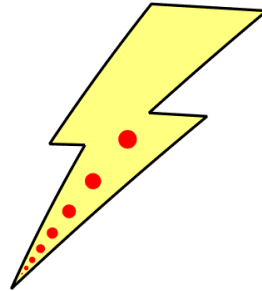
by Jim Vargas

under direction of Dr. Jeff Ovall

Portland State University



This talk is based on...



## *Solving Laplace Problems with Corner Singularities via Rational Functions*

- ...A paper written by Gopal and Trefethen, published in SIAM Journal on Numerical Analysis September 2019
- The Lightning Laplace code, based on the paper, yields accurate approximations quickly (on nice problems)
- <https://epubs.siam.org/doi/pdf/10.1137/19M125947X>
- <https://people.maths.ox.ac.uk/trefethen/lightning.html>



## Here's the problem

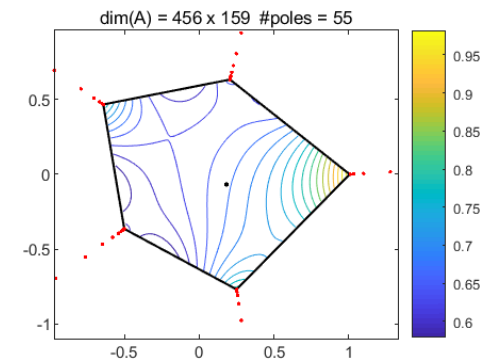
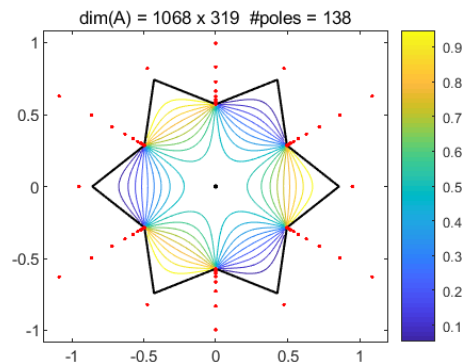
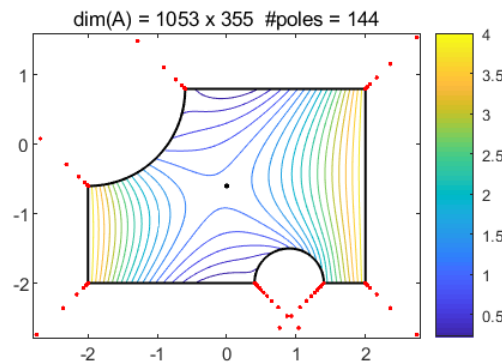
We wish to find a (real) function  $u$  over a domain  $\Omega \subset \mathbb{C}$  which satisfies

$$\Delta u(z) = 0, \quad z \in \Omega \qquad u(z) = h(z), \quad z \in \Gamma.$$

In particular, we want to be able to handle a domain with sharp corners, curves etc.

We will find  $r$ , and approximation of  $u$  ( $u \approx \text{Re}[r]$ ).

$$r(z) = \sum_{j=1}^{N_1} \frac{a_j}{z - z_j} + \sum_{j=0}^{N_2} b_j (z - z_*)^j$$





## Why this problem?

- Problems involving the Laplace operator  $\Delta = \nabla^2$  frequently appear in physical equations:
  - Heat Equation  $\alpha \nabla^2 u = \partial_t u$
  - Schrodinger Equation  $\left[ \frac{-\hbar^2}{2m} \nabla^2 + V \right] \Psi = i\hbar \partial_t \Psi$
  - Wave Equation  $c^2 \nabla^2 u = \partial_t^2 u$
  - And more...
- Functions which satisfy Laplace's Equation have very nice properties, and are called harmonic.



## Some nice properties of functions of interest

- The real and imaginary parts of a holomorphic (and thus also an analytic) function  $f = u + iv$  are harmonic;
- $f$  is also smooth (infinitely differentiable); by extension this applies to  $u$  and  $v$  as well.
- Maximum Principle: a harmonic function on a compact domain attains a max. (and min.) on the boundary.



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On a simply connected domain we can construct a holomorphic function from a harmonic one: given  $u$ , define  $g = u - iu$ . The theory will work with holomorphic functions, which will trickle down to our problem.

If  $r$  approximates  $f$ , having real part  $u$ , the worst we'll do over the whole domain in approximating  $u$  is  $\|u(z) - \operatorname{Re}[r(z)]\|_\infty, z \in \Gamma$ .



## Back to the problem

$$r(z) = \underbrace{\sum_{j=1}^{N_1} \frac{a_j}{z - z_j}}_{\text{"Newman"}} + \underbrace{\sum_{j=0}^{N_2} b_j (z - z_*)^j}_{\text{"Runge"}}$$

- Using the scheme in the paper, we can have root exponentially good approximations for  $u$ . The task at hand is finding the coefficients  $a_j, b_j$ .

$$\|f - r_n\|_{\Omega} = O(e^{-C\sqrt{n}})$$

- The theorems in the paper are based on interpolation, showing existence.
- In the code, the problem is solved via a least squares approach using QR factorization. Code is written in MATLAB.

$$\min_{\substack{\{a_1, \dots, a_{N_1}\} \\ \{b_1, \dots, b_{N_2}\}}} \sum_{j=0}^M |r(y_j) - h(y_j)|^2$$

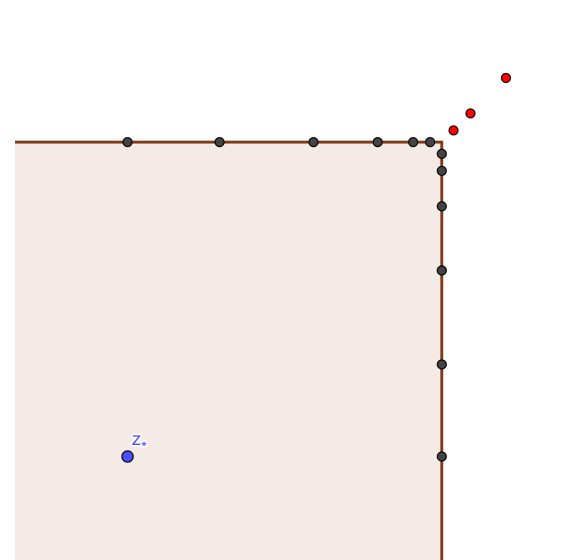


## Describing $r$

$$r(z) = \sum_{j=1}^{N_1} \frac{a_j}{z - z_j} + \sum_{j=0}^{N_2} b_j (z - z_*)^j$$

**Newman Part:** built to handle corners.

- The terms  $z_j$  are poles, exponentially clustered near a corner on the exterior of  $\Omega$  (works for spacing scaled at least  $O(n^{-1/2})$ ).
- "Rational functions are more powerful than polynomials for approximating functions near singularities..."<sup>a</sup>



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<sup>a</sup>Lloyd N. Trefethen. 2013. *Approximation theory and approximation practice*, Society for Industrial and Applied Mathematics.



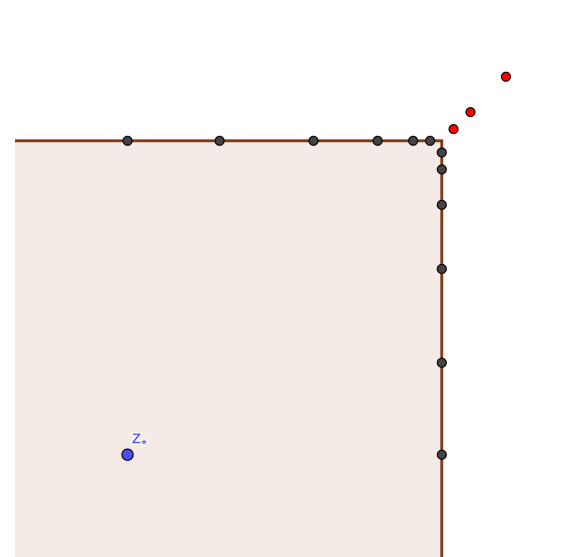


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**The Runge part:** built to handle the interior.

- The term  $z_*$  is an expansion point, near the middle of  $\Omega$ .
- Polynomials can approximate root exponentially well on a nice domain (going back to Runge).





## The function $r$ is harmonic

$$r(z) = \sum_{j=1}^{N_1} \frac{a_j}{z - z_j} + \sum_{j=0}^{N_2} b_j (z - z_*)^j$$

To prove  $r$  is harmonic, consider  $f(z) = 1/z$  and  $g(z) = z^k$ . The function  $f$  can be decomposed as  $f = u + iv$ , where

$$u(x, y) = \frac{x}{x^2 + y^2}$$

$$v(x, y) = \frac{-y}{x^2 + y^2}.$$

Taking derivatives will show that  $u$  and  $v$  satisfy the Cauchy-Riemann equations,  $\partial_x u = \partial_y v$ ,  $\partial_y u = -\partial_x v$ , meaning  $f$  is (holomorphic, and thus) harmonic.



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Writing  $g$  in polar form, then in terms of sines and cosines is enough to see  $g$  is harmonic:

$$g(z) = \rho e^{ik\theta} = \rho [\cos(k\theta) + i \sin(k\theta)].$$

Adding these templates, applying translations and scaling as necessary give us our result.



## An important lemma

**Hermite integral formula** for rational interpolation.

Let  $\Omega$  be a simply connected domain in  $\mathbb{C}$  bounded by a closed curve  $\Gamma$ , and let  $f$  be analytic in that domain and extend continuously to the boundary. Let interpolation points  $\alpha_0, \dots, \alpha_{n-1} \in \Omega$  and poles  $\beta_0, \dots, \beta_{n-1}$  anywhere in the complex plane be given. Let  $r$  be the unique type  $(n-1, n)$  rational function with simple poles at  $\{\beta_j\}$  that interpolate  $f$  at  $\{\alpha_j\}$ . Then for any  $z \in \Omega$ ,

$$f(z) - r(z) = \frac{1}{2\pi i} \int_{\Gamma} \frac{\phi(z)}{\phi(t)} \frac{f(t)}{t - z} dt,$$
$$\phi(z) = \prod_{j=0}^{n-1} (z - \alpha_j) \bigg/ \prod_{j=0}^{n-1} (z - \beta_j).$$



## First Theorem

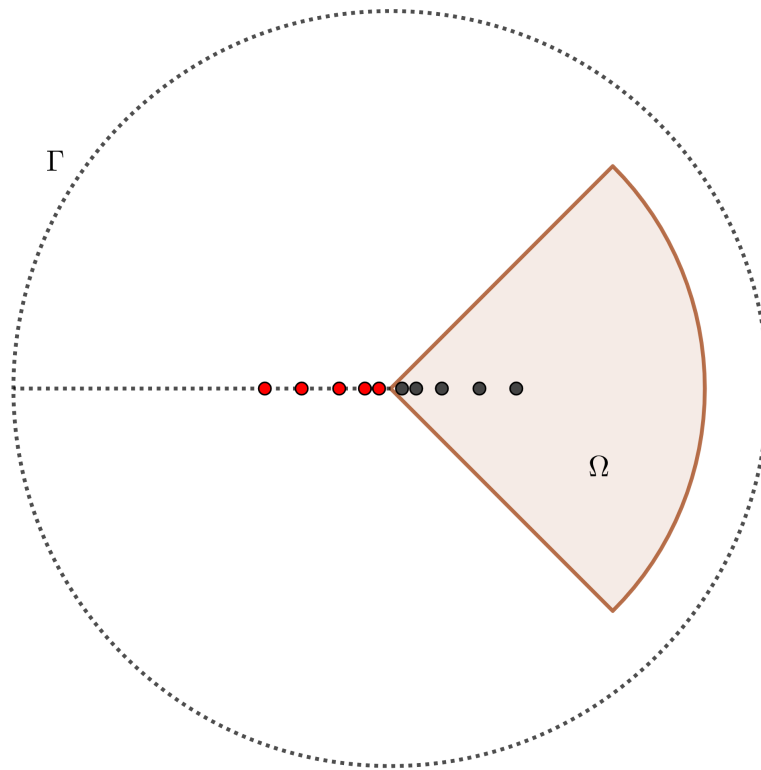
Let  $f$  be a bounded analytic function in the slit disk  $A_\pi$  that satisfies  $f(z) = O(|z|^\delta)$  as  $z \rightarrow 0$  for some  $\delta > 0$ , and let  $\theta \in (0, \pi/2)$  be fixed. Then for some  $0 < \rho < 1$  depending on  $\theta$  but not on  $f$ , there exist type  $(n-1, n)$  rational functions  $\{r_n\}$ ,  $1 \leq n < \infty$ , such that

$$\|f - r_n\|_\Omega = O(e^{-C\sqrt{n}})$$

as  $n \rightarrow \infty$  for some  $C > 0$ , where  $\Omega = \rho A_\theta$ . Moreover, each  $r_n$  can be taken to have simple poles only at

$$\beta_j = -e^{-\sigma j/\sqrt{n}}, \quad 0 \leq j \leq n-1,$$

where  $\sigma > 0$  is arbitrary.



$$A_\theta = \{z \in \mathbb{C} : |z| < 1, \quad |\arg(z)| < \theta\}$$

$$\Omega = \rho A_\theta$$



## Second Theorem

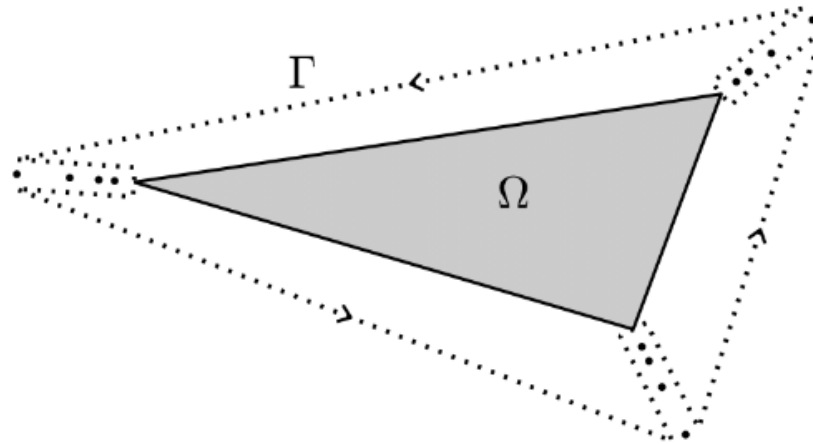
Let  $\Omega$  be a convex polygon with corners  $w_1, \dots, w_m$ , and let  $f$  be an analytic function in  $\Omega$  that is analytic on the interior of each side segment and can be analytically continued to a disk near each  $w_k$  with a slit along the exterior bisector there. Assume  $f$  satisfies  $f(z) - f(w_k) = O(|z - w_k|^\delta)$  as  $z \rightarrow w_k$  for each  $k$  for some  $\delta > 0$ . There exist degree  $n$  rational functions  $\{r_n\}$ ,  $1 \leq n < \infty$  such that

$$\|f - r_n\|_\Omega = O(e^{-C\sqrt{n}})$$

as  $n \rightarrow \infty$  for some  $C > 0$ . Moreover, each  $r_n$  can be taken to have finite poles only at points exponentially clustered along the exterior bisectors at the corners, with arbitrary clustering parameter  $\sigma$ , as long as the number of poles near each  $w_k$  grows at least in proportion to  $n$  as  $n \rightarrow \infty$ .



## Second Theorem: the idea<sup>a</sup>



Split  $f$  into  $2m$  terms, a "Newman" part and a "Runge" part:

$$f = \sum_{k=1}^m f_k + \sum_{k=1}^m g_k.$$

The Runge part can be handled by previously established results, and the Newman part can be handled by applying the first theorem to each corner.

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<sup>a</sup>Image from Gopal, A., & Trefethen, L. N. (2019). *Solving Laplace Problems with Corner Singularities via Rational Functions*. SIAM Journal on Numerical Analysis.





## Some extensions

Numerical experiments show that:

- We can get root exponentially good approximations on non-convex domains;
- We're not limited to sectors and convex polygons, we can have curvy edges.

These theorems apply to a holomorphic function  $f$ , but our problem involves a harmonic  $u$ .

If we assume  $u$  satisfies the corner behavior needed and  $\Omega$  is simply connected, then so will a  $v$ , where we can have an  $f = u + iv$ .



## The Algorithm

1. Define boundary  $\Gamma$ , corners  $w_1, \dots, w_m$ , boundary function  $h$ , tolerance  $\varepsilon$ .
2. For increasing values of  $n$  with  $\sqrt{n}$  approximately evenly spaced;
  - 2a. fix  $N_1 = O(mn)$  poles  $1/(z - z_k)$  clustered outside the corners;
  - 2b. fix  $N_2 + 1 = O(n)$  monomials  $1, (z - z_*)^1, \dots, (z - z_*)^{N_2}$  and set  $N = N_1 + N_2 + 1$ ;
  - 2c. choose  $M \approx 3N$  sample points on a boundary, also clustered near corners;
  - 2d. evaluate at sample points to obtain an  $M \times N$  matrix  $A$  and  $M$ -vector  $b$ ;
  - 2e. solve the least-squares problem  $Ax \approx b$  for the coefficient vector  $x$ ;
  - 2f. exit loop if  $\|Ax - b\|_\infty < \varepsilon$  or if  $N$  is too large or the error is growing.
3. Confirm accuracy by checking the error on a finer boundary mesh.
4. Construct a function to evaluate  $r(z)$  based on computed coefficients  $x$ .

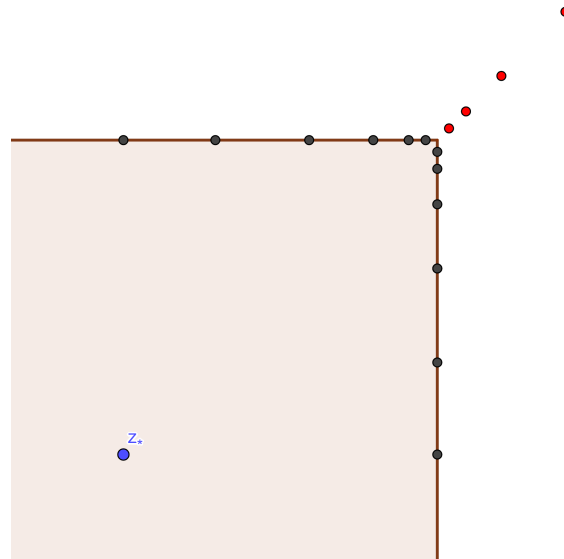


## The code

Code is branded "Lightning Laplace." We enter:

- Corners of a polygonal-ish/curvy domain in  $\mathbb{C}$ ;
- boundary data in the form of a(n) real function handle(s), or scalar values, corresponding to the edges.

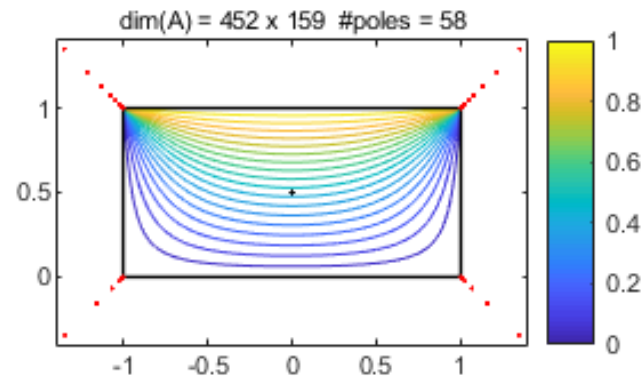
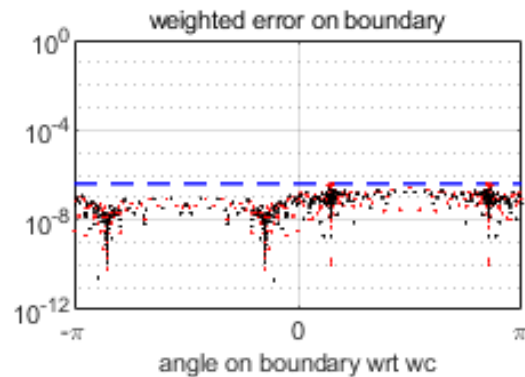
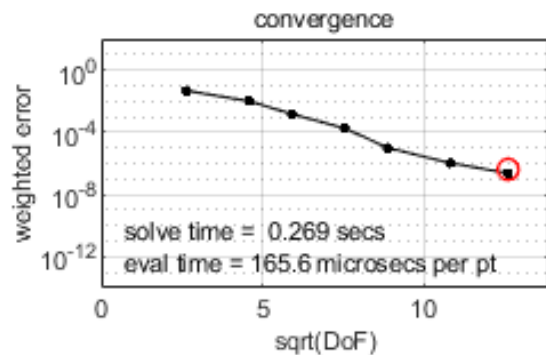
Errors are computed by comparing the procedure with a finer sampling (so not a true error).





```
disp('rectangle with piecewise constant BCs')
P = [-1 1 1+1i -1+1i];
h = [0 0 1 0];
laplace(P,h,'plots','rel');
```

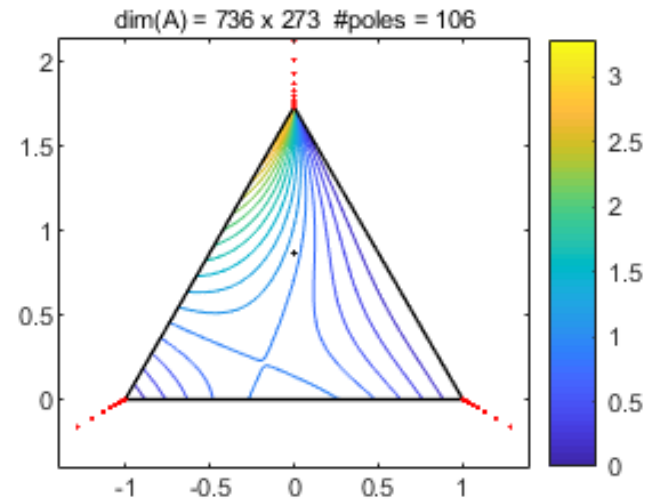
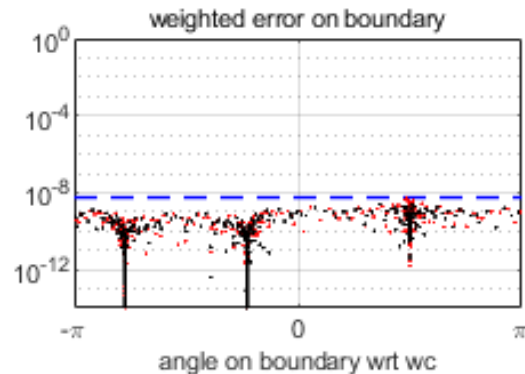
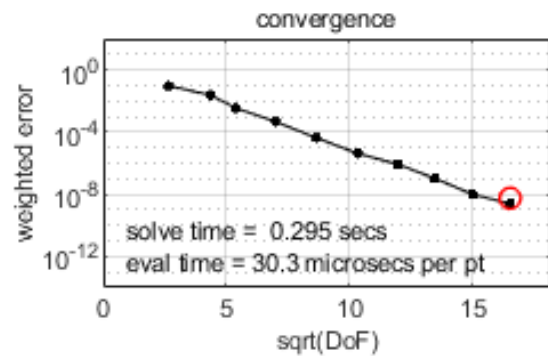
Solve time: 0.269s, Epsilon: 1e-6, dim(A)=452 × 159, #poles: 58

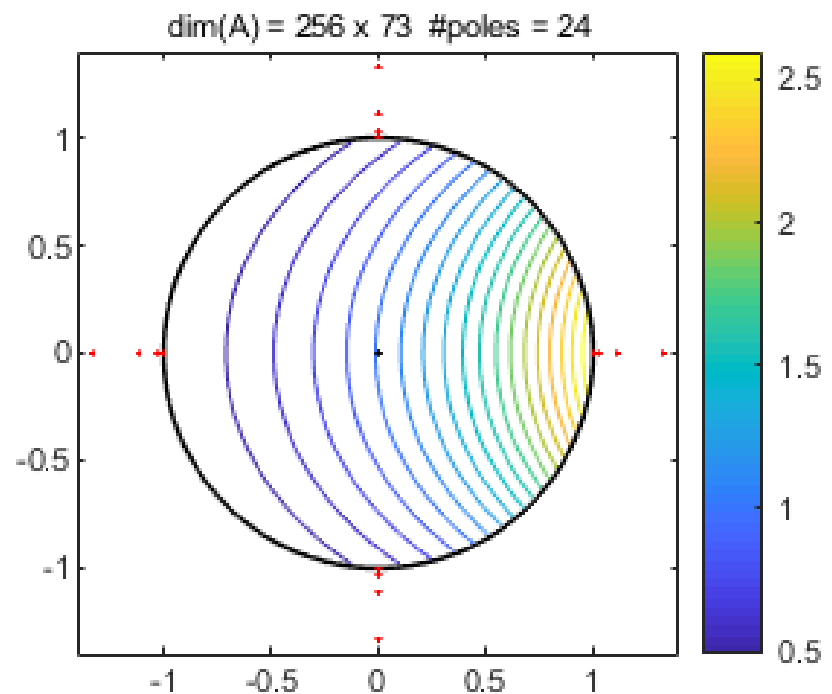
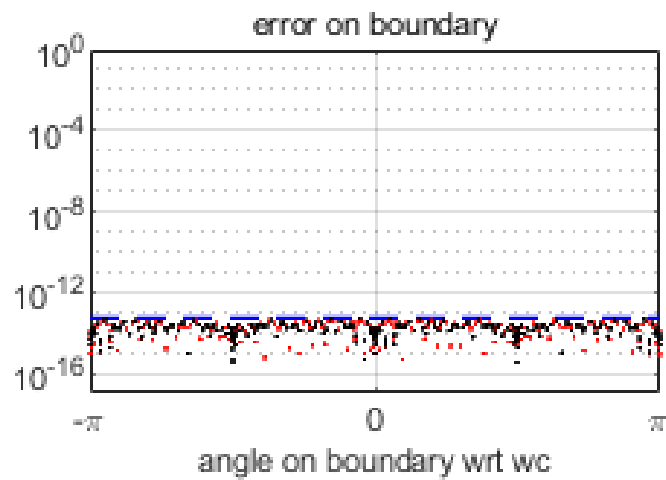
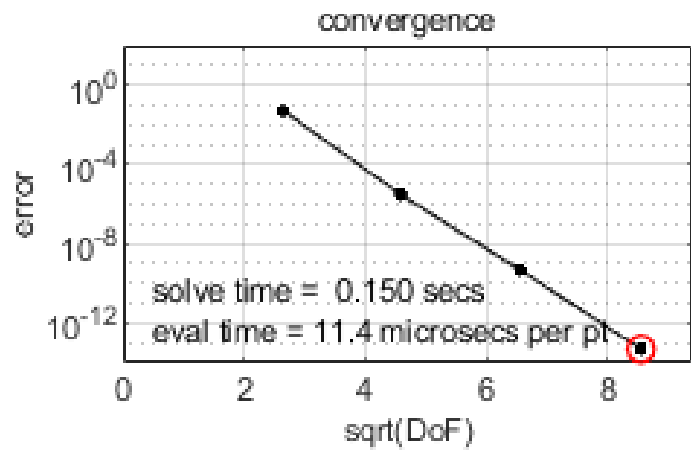


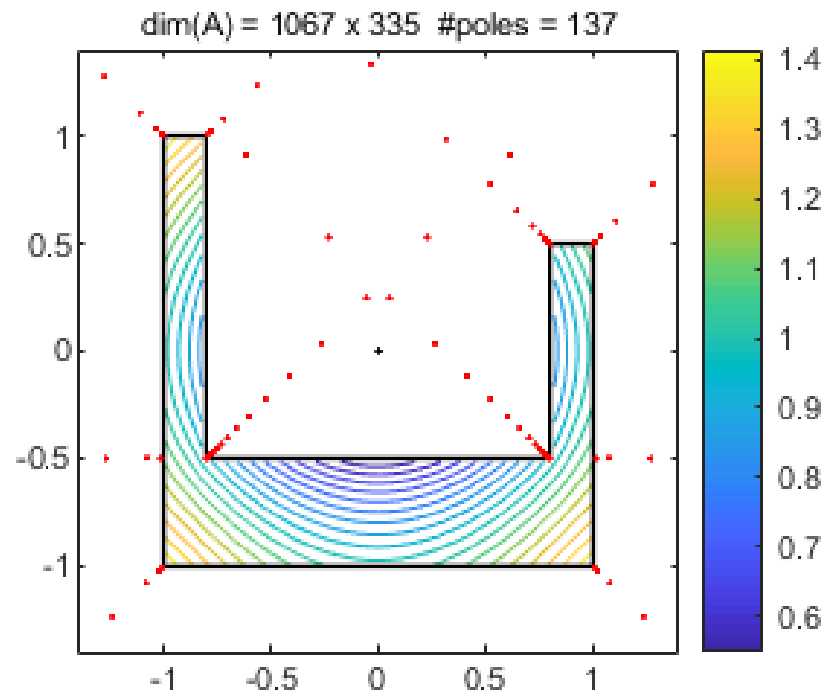
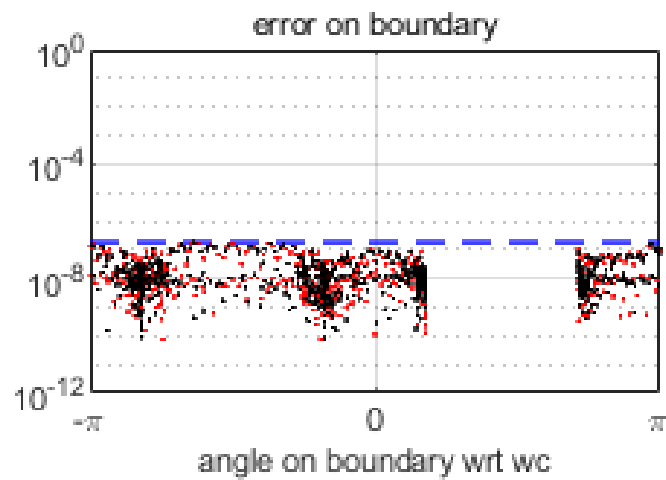
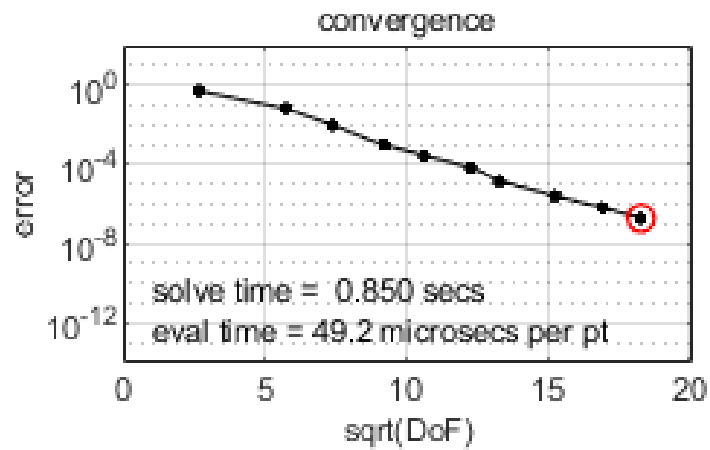


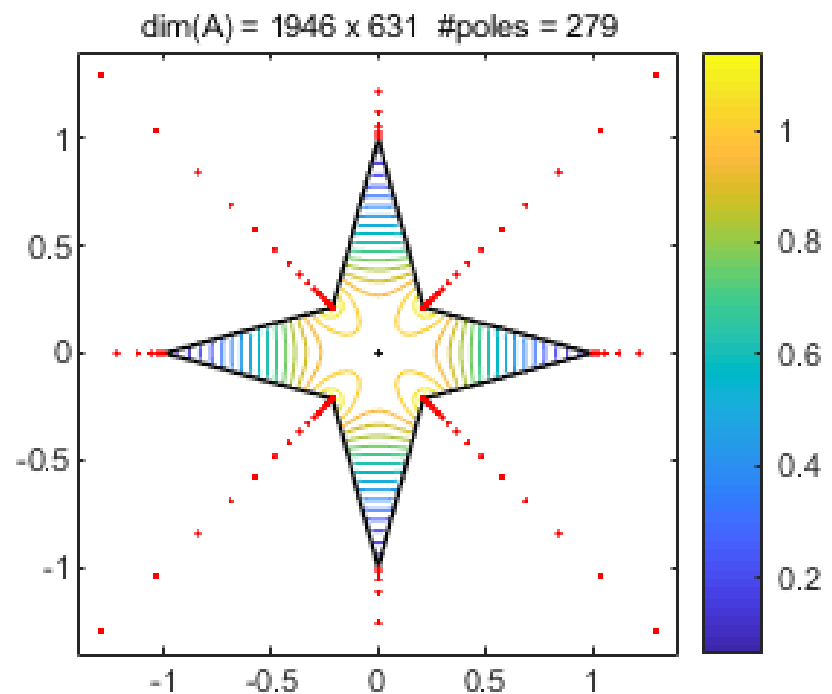
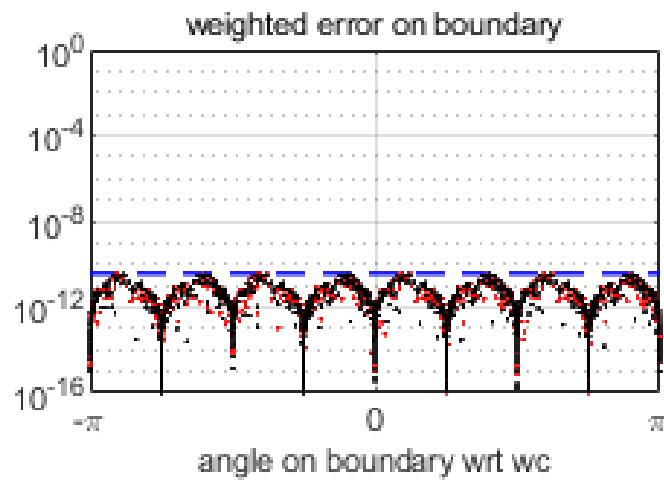
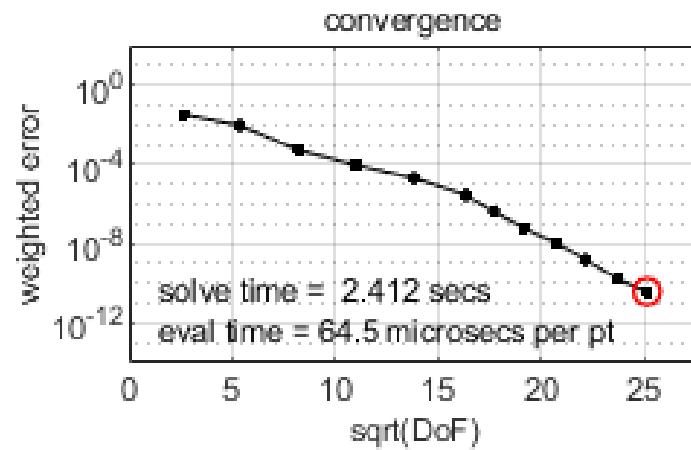
```
disp('equilateral triangle with one non-constant BC')
h = {@(x) cos(pi*x/2), @(z) 0*z, @(z) 2*imag(z)};
P = [-1 1 1i*sqrt(3)];
laplace(P,h,'plots','tol',1e-8);
```

Solve time: 0.295s. Epsilon: 1e-8,  $\dim(A)=736 \times 273$ , #poles: 106

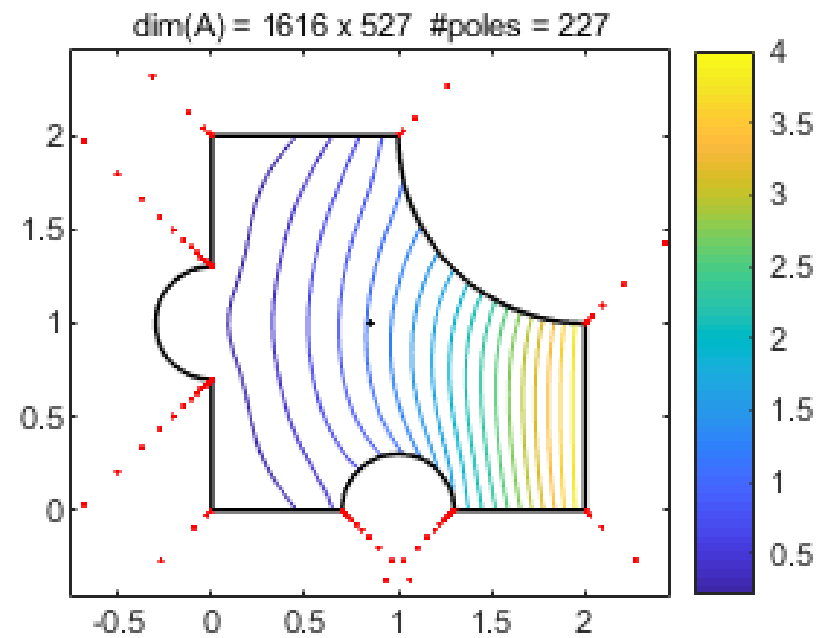
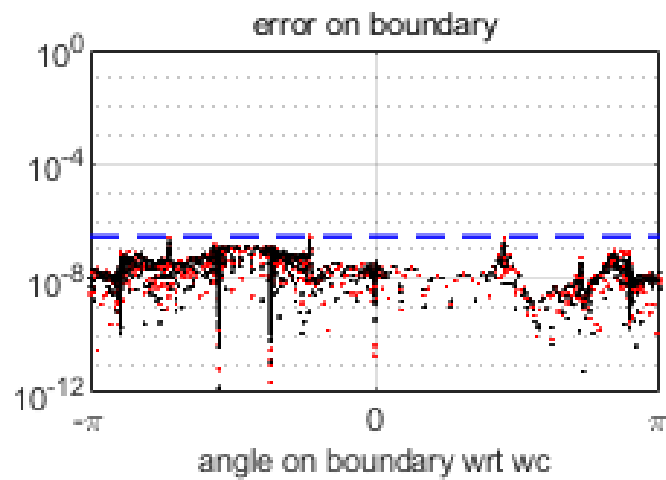
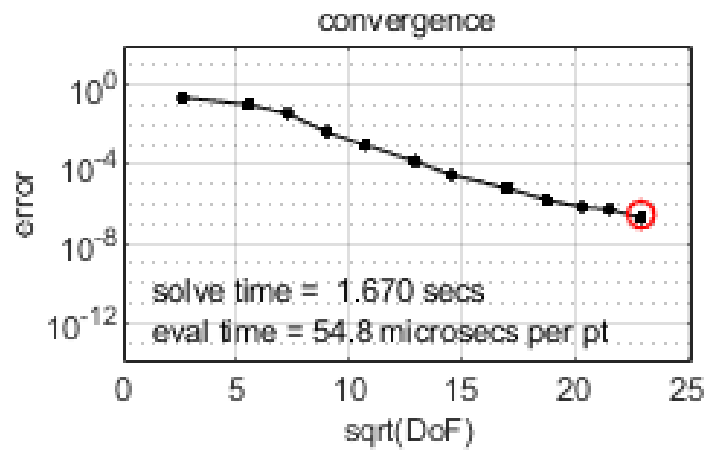


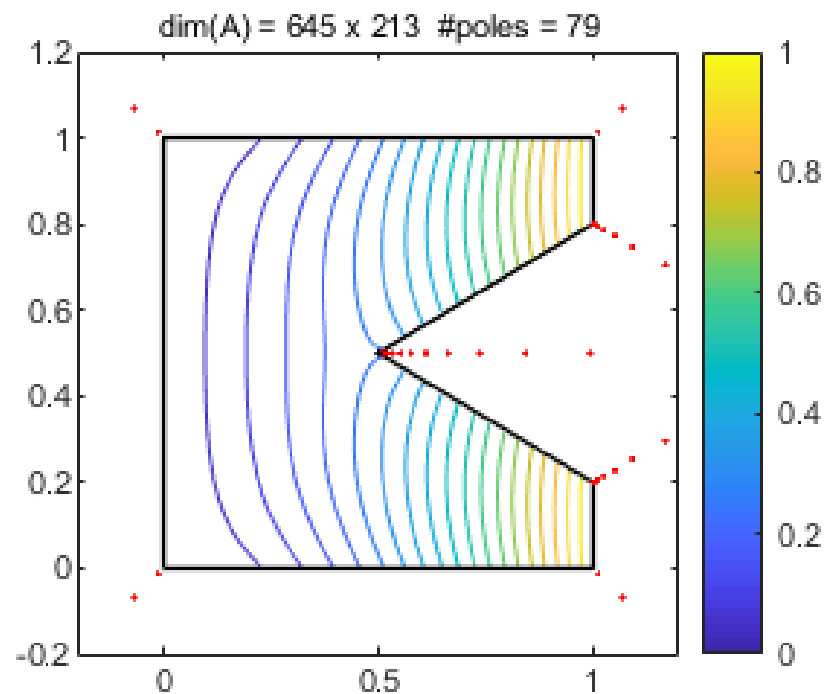
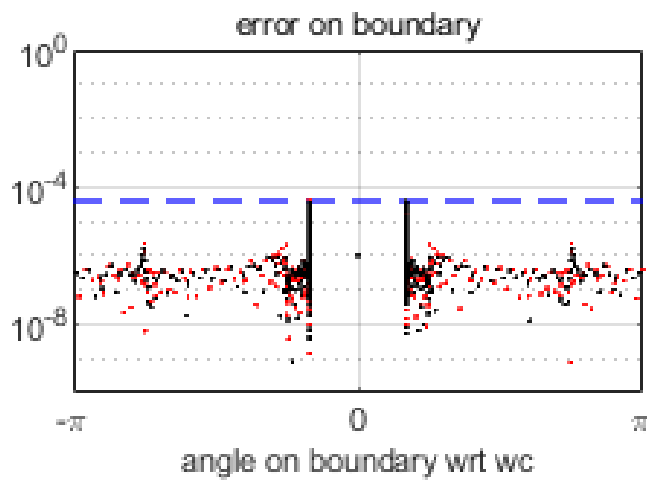
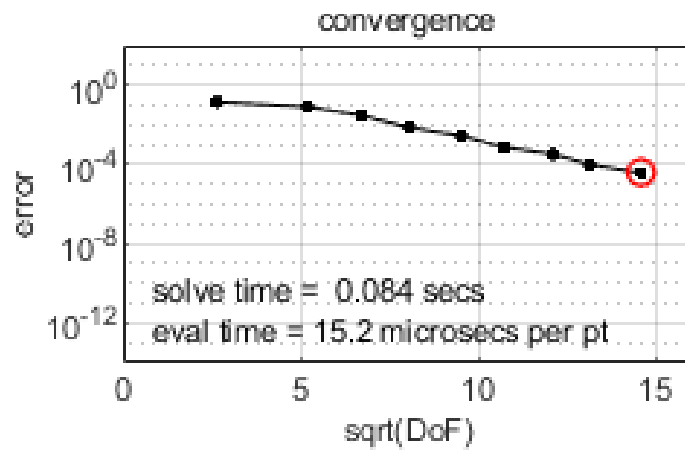


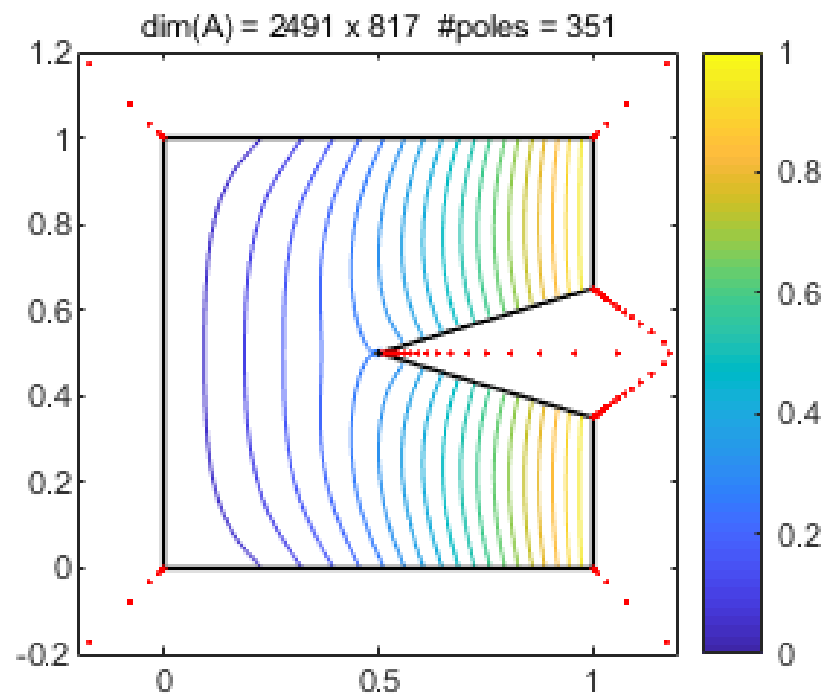
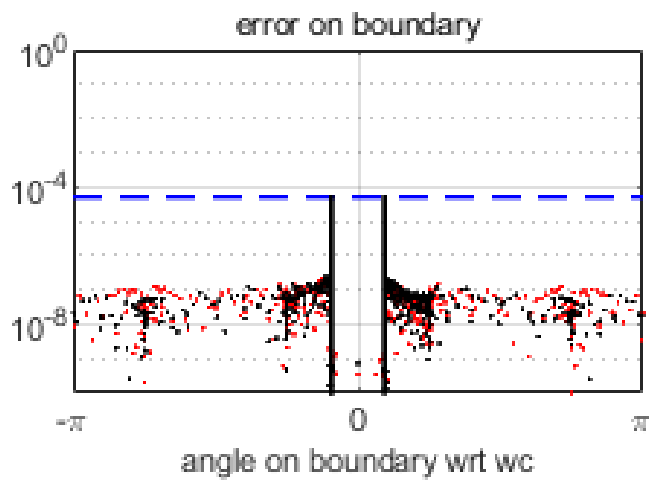
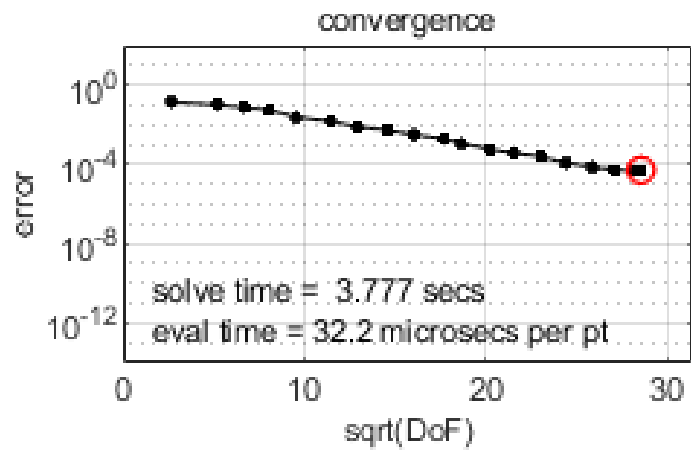


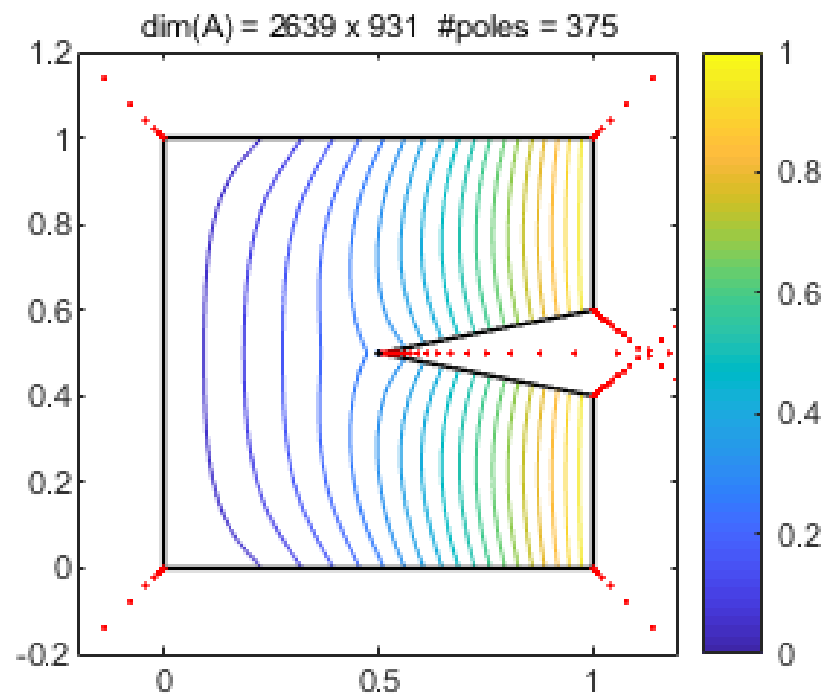
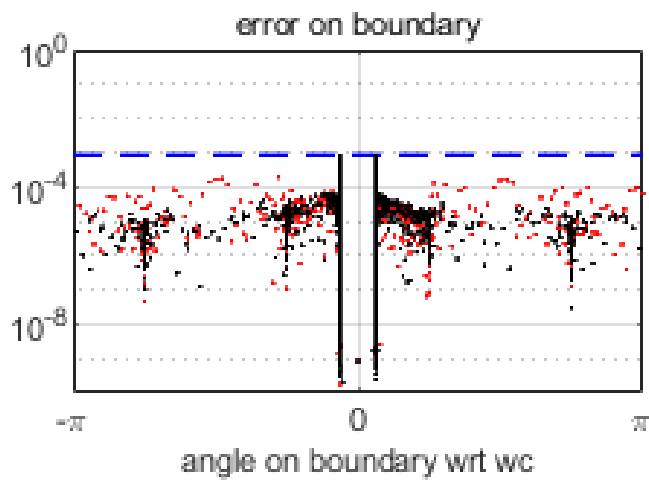
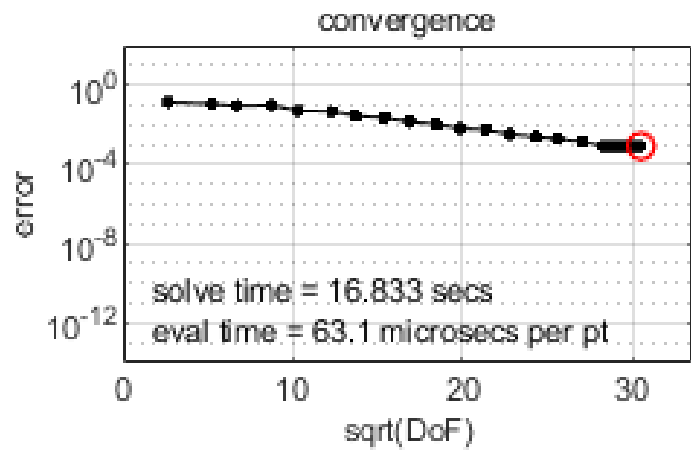


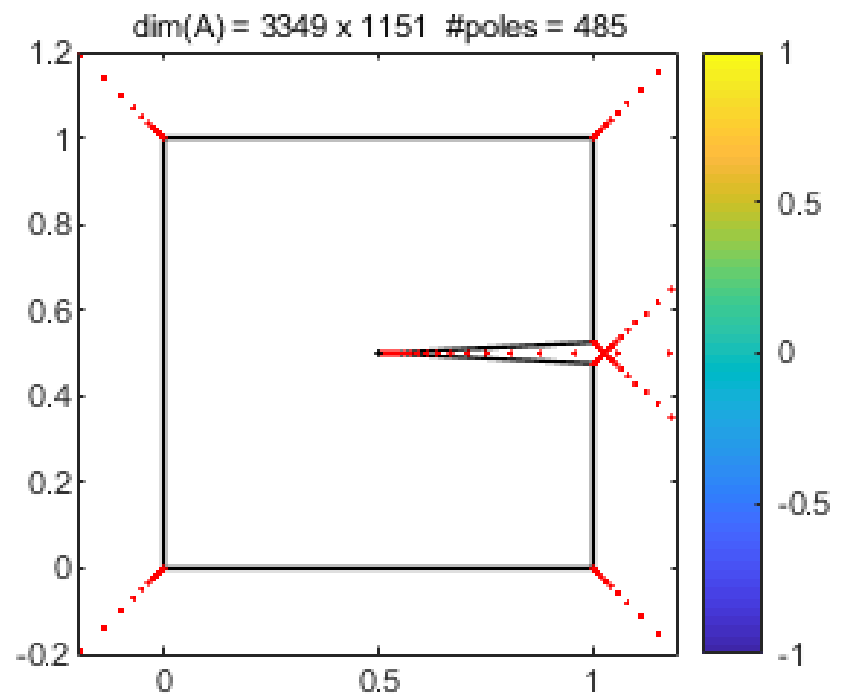
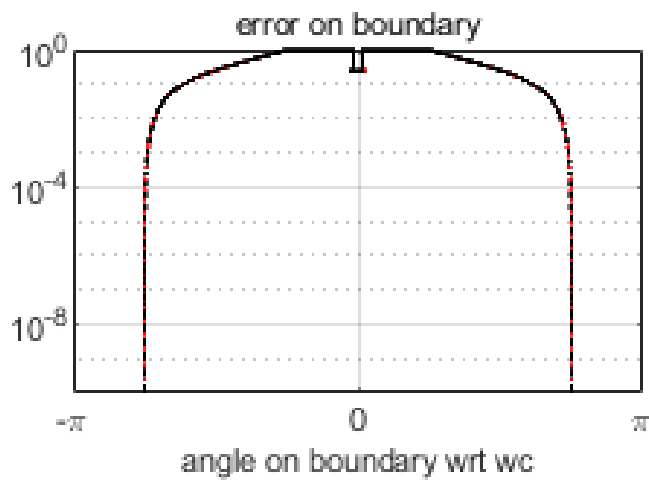
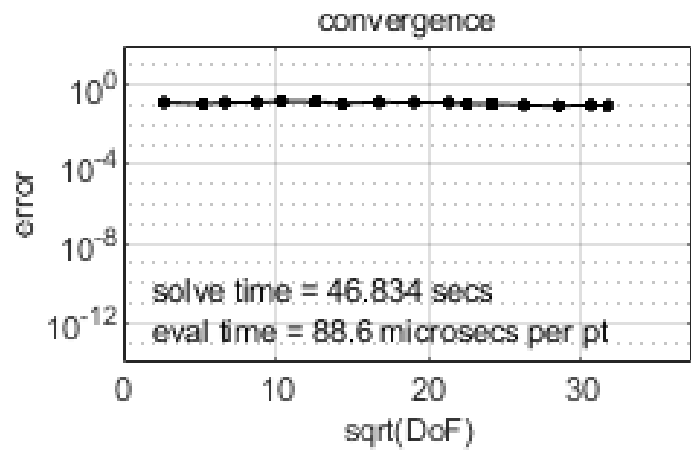


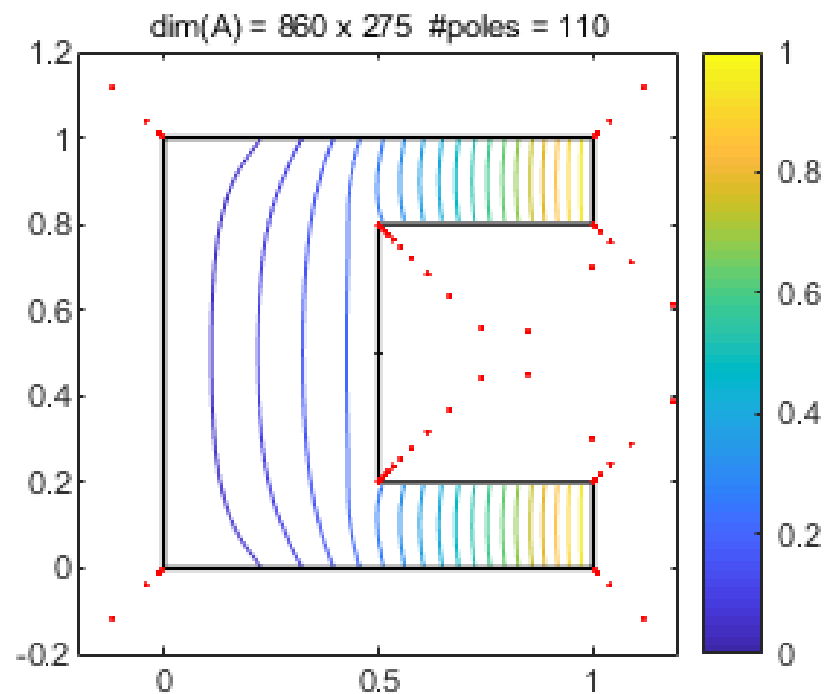
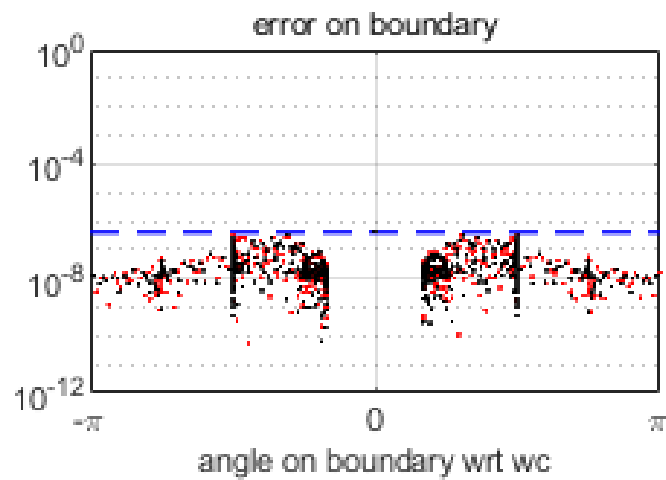
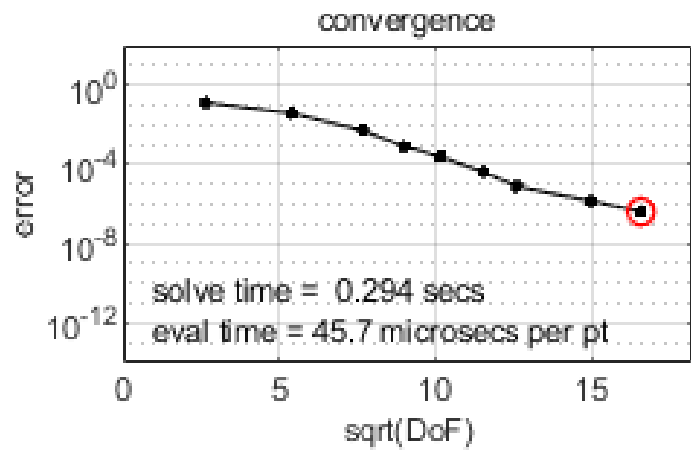


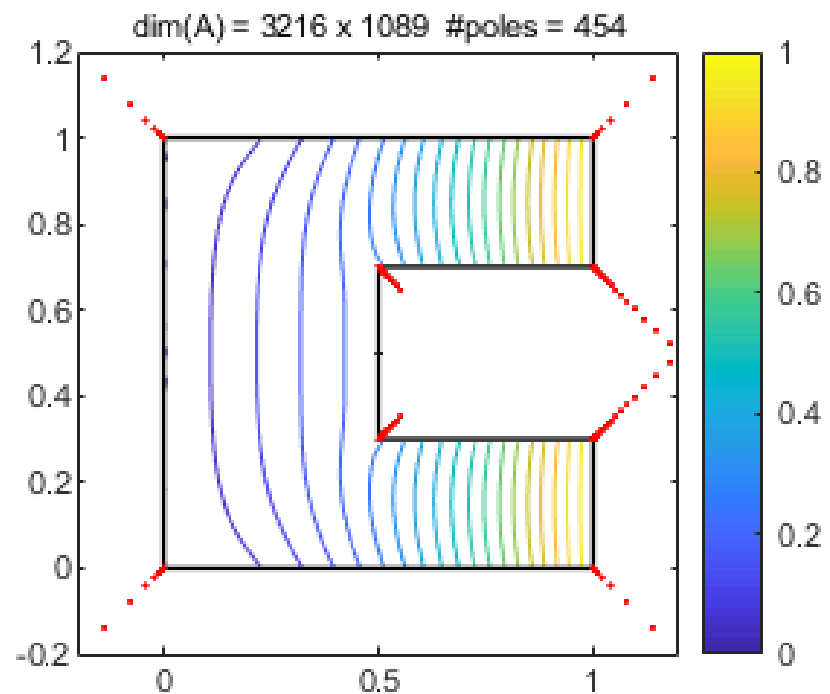
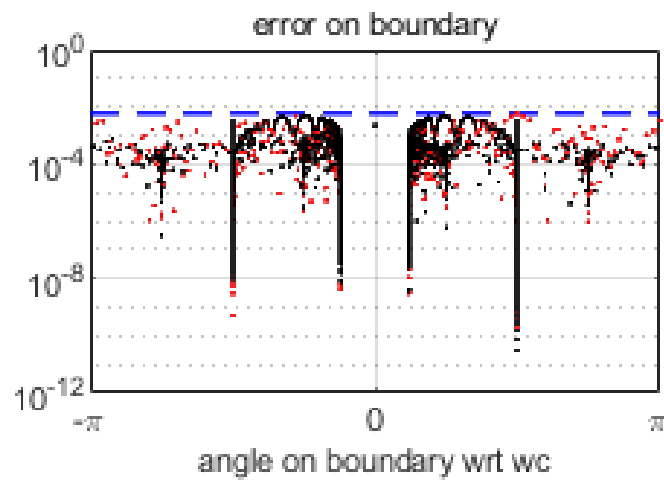
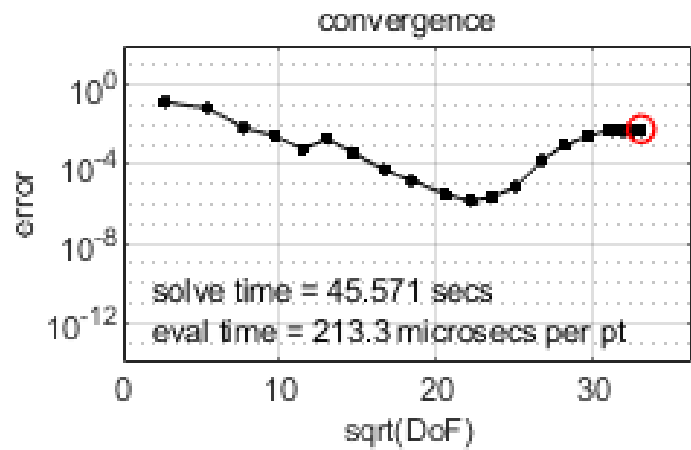


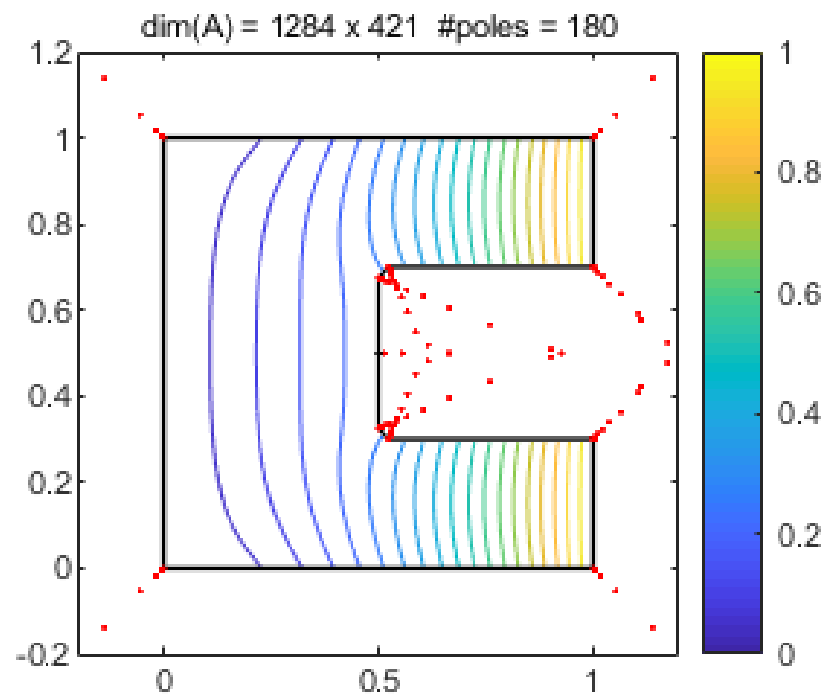
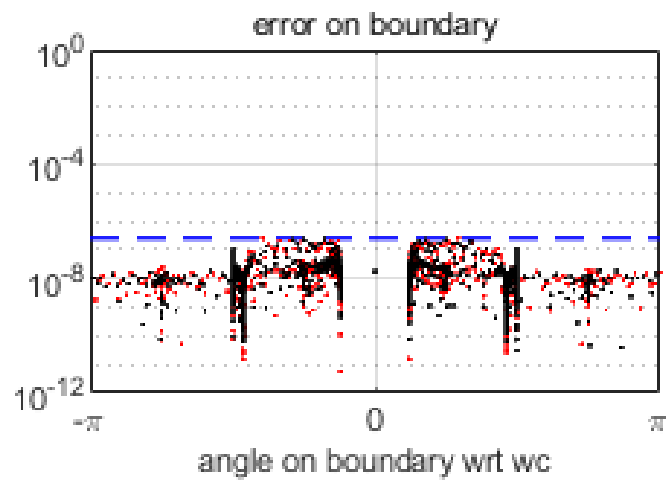
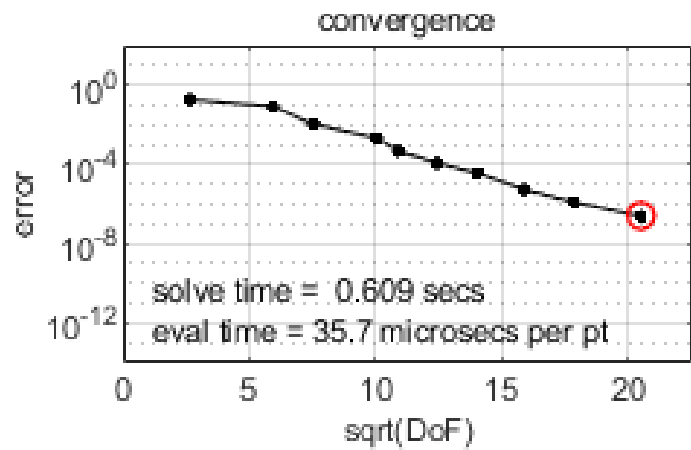




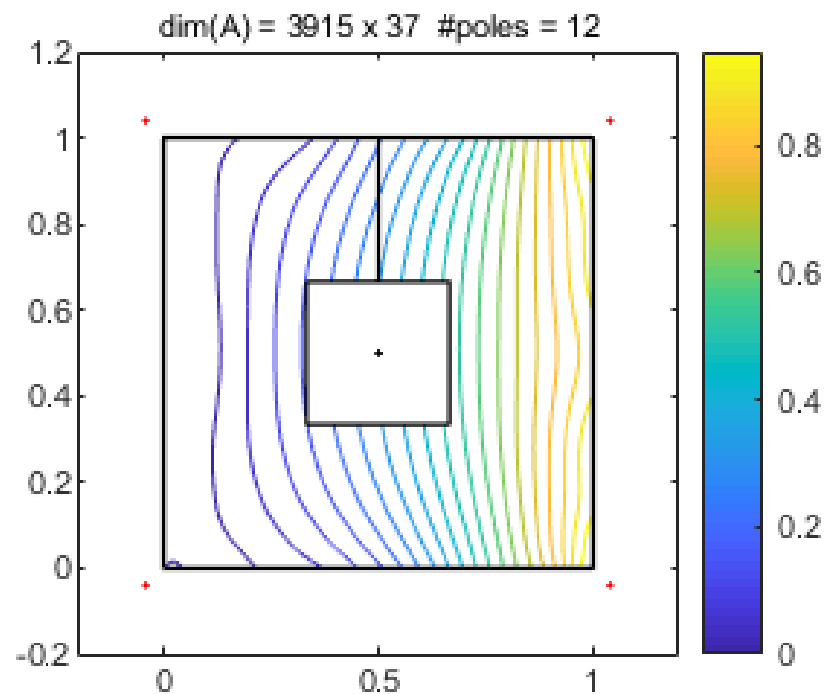
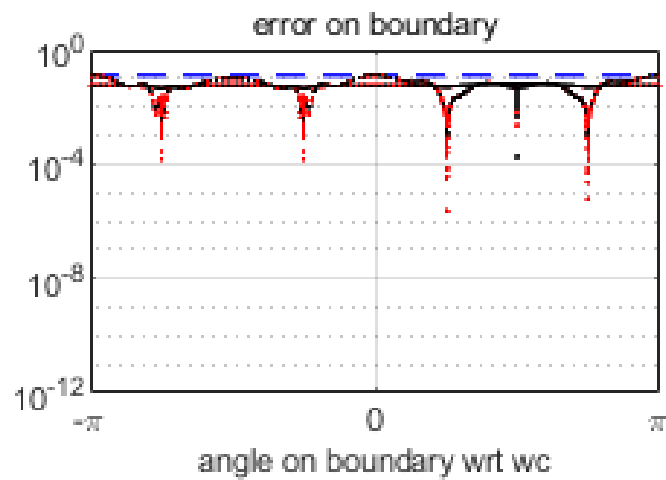
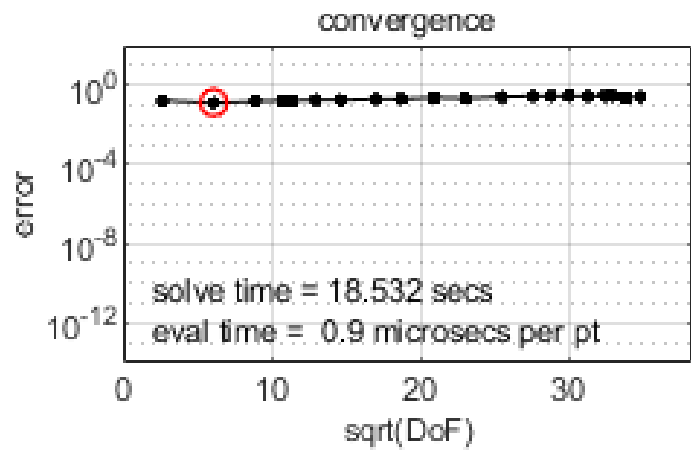


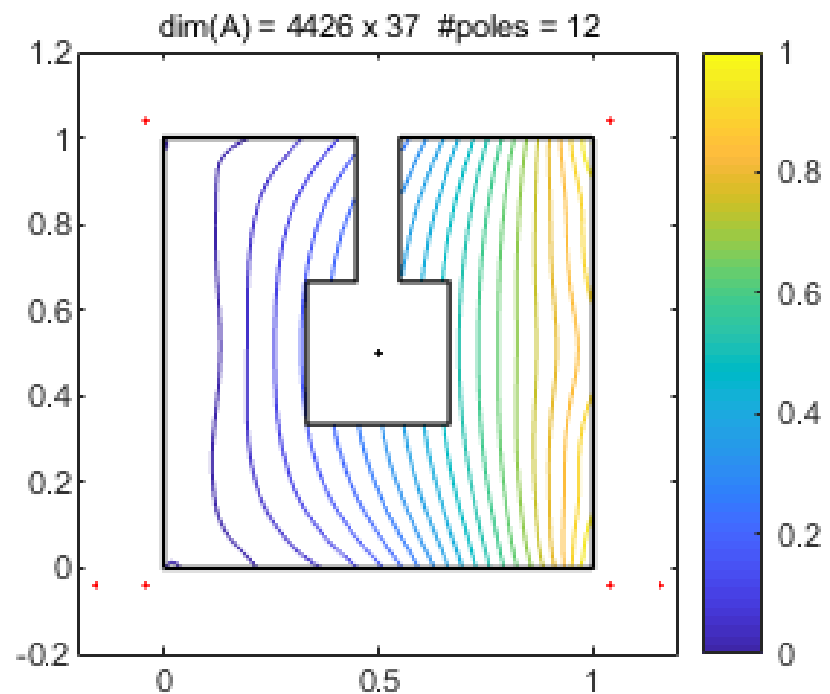
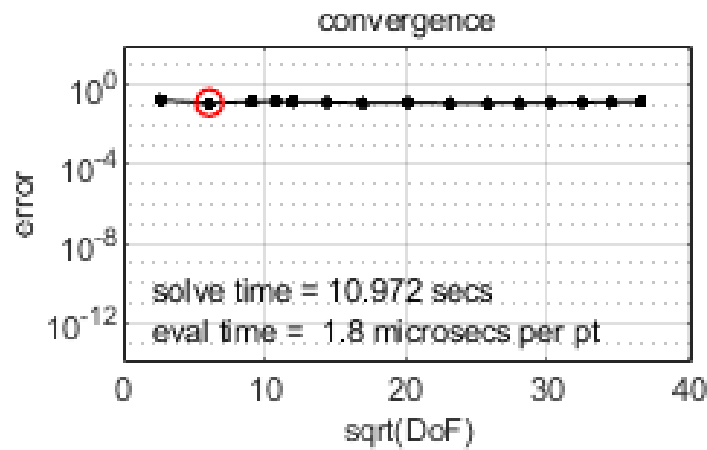


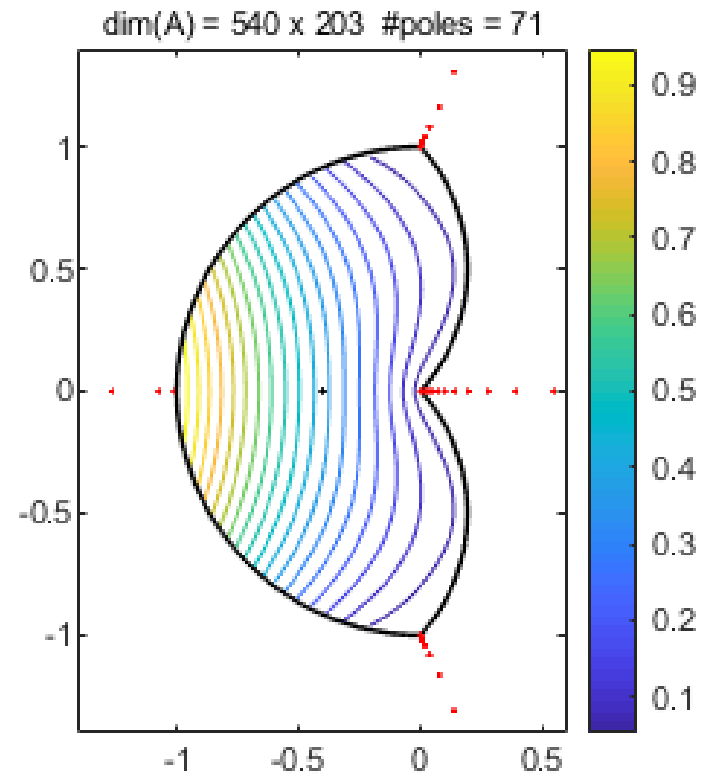
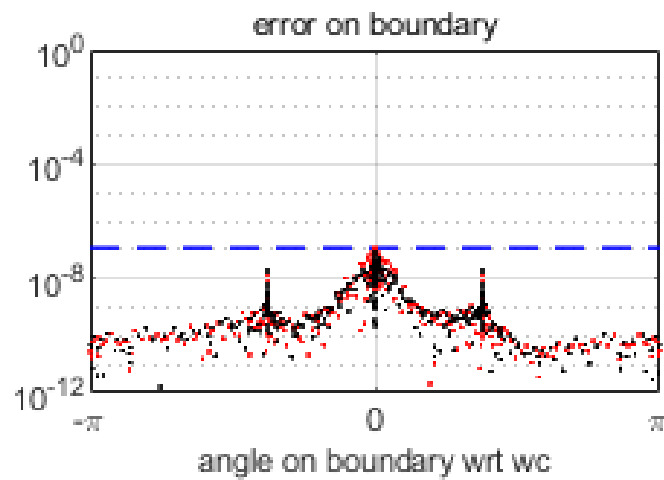
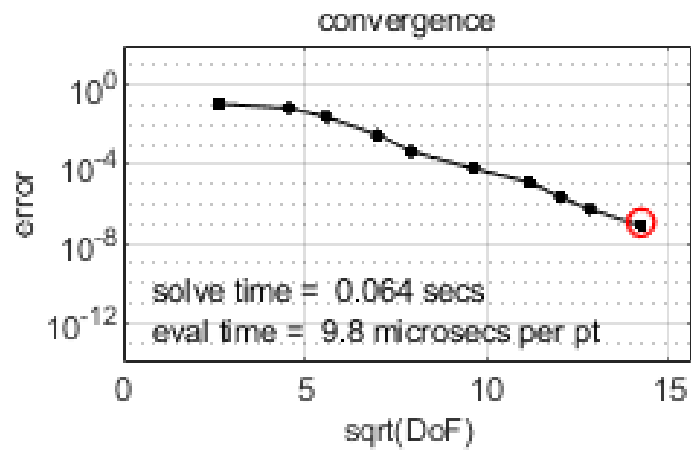


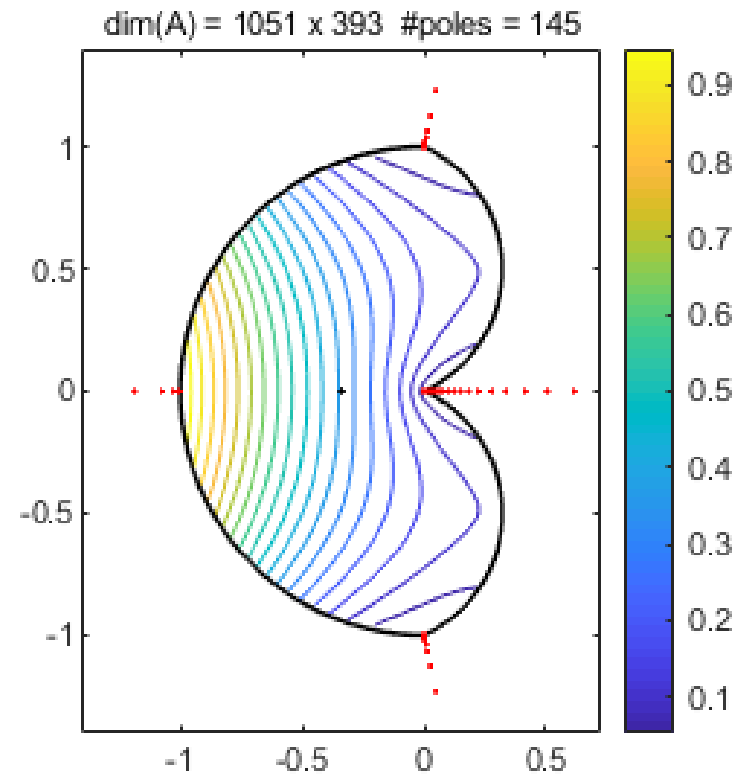
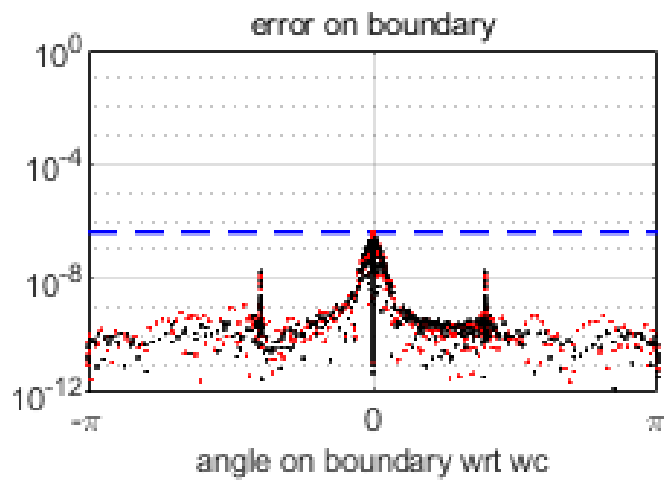
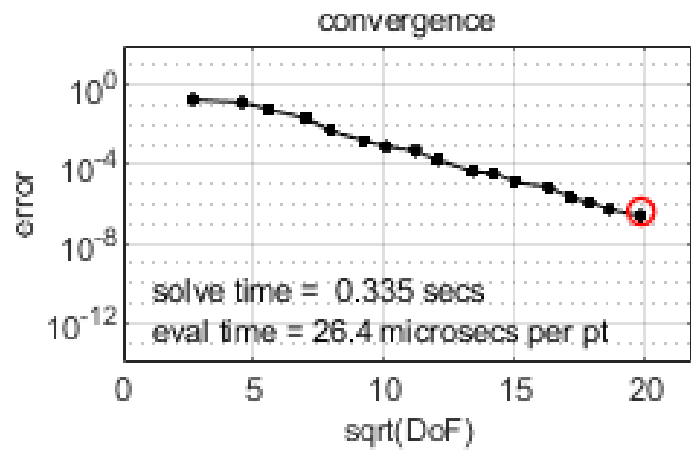


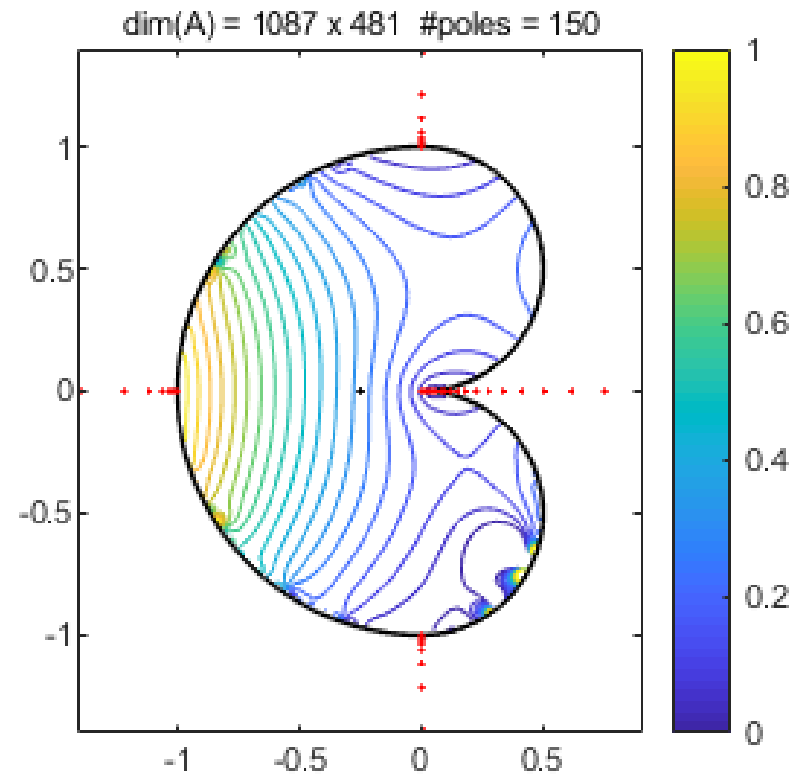
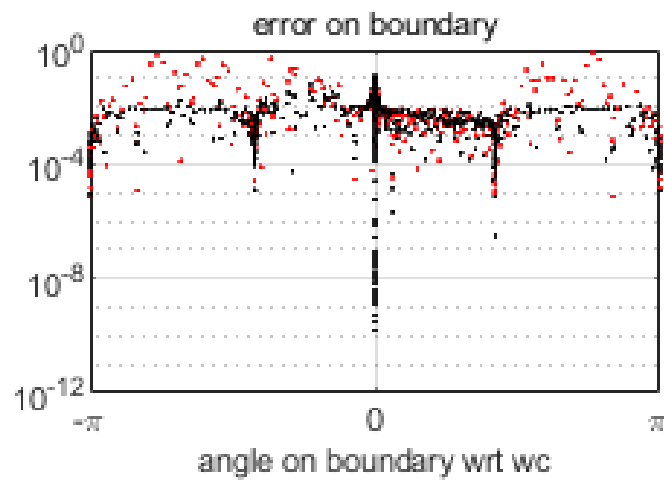
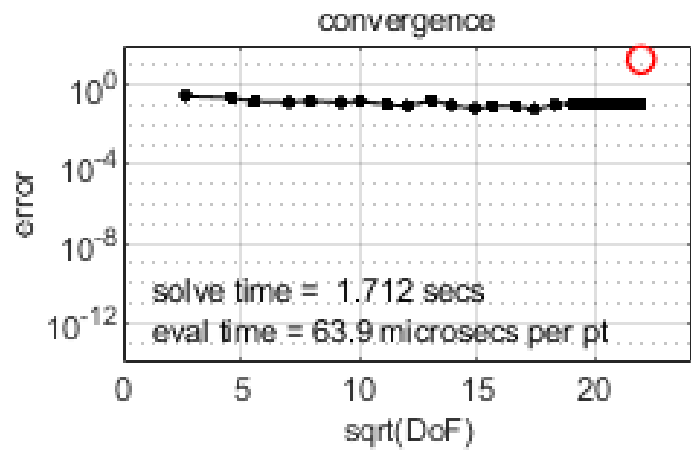


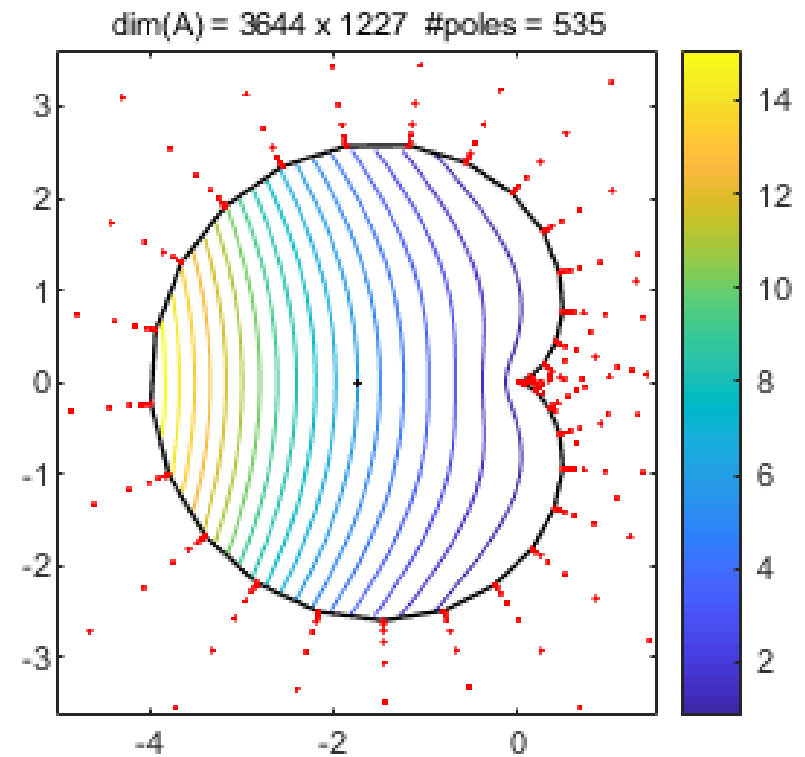
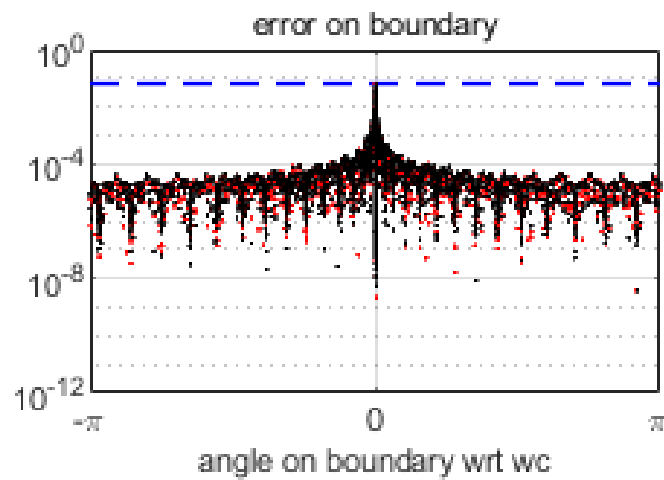
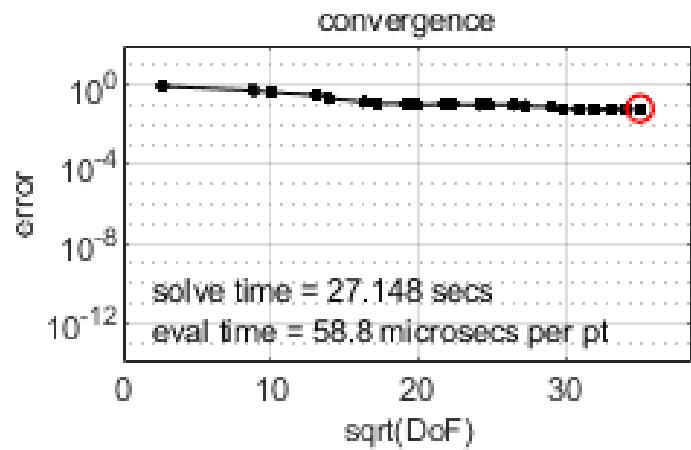














## What else...

Variants:

- Discontinuous boundary conditions: pretty much no problem, except supremum norm convergence to zero. Instead, use a supremum norm weighted by distance to the nearest corner.
- Multiply connected domains: very much a problem.
- Poisson equation  $\Delta u = f$  with boundary conditions: find  $\Delta v = f$  with arbitrary bd. conds. first, then do  $w = u - v$ .
- Faster than root exponential convergence: not possible.
- Domain size, matrix size: Runge part of  $r$  works less optimally far away from  $z_*$ ;  $m$  corners means operation count  $O(m^3 |\log(\varepsilon)|^6)$ .



## What else...

- Authors say Finite Element Methods can't match Lightning Laplace's simplicity and performance.
- Boundary Integral Equations are good when applicable, and is "the most powerful tool currently available."
- Proofs for various more general domains
- Expanding code functionality and usability





The End