

Trees (III)

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Note: The most of the information of these slides was extracted and adapted from Weiss's book, "*Data Structures and Algorithm Analysis in Java*". They are provided for COP3530 students only. Not to be published or publicly distributed without permission by the publisher.



COP-3530 - Data Structures



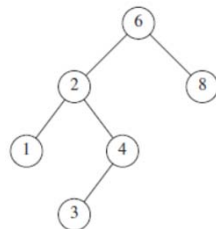
Module #3: Trees (part III)

Outline:

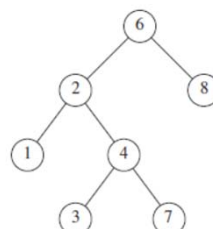
- Binary Search Trees (BST).
- Main operations on BST:
 - Insertion/Add.
 - Remove/Del.
 - Find.
 - Complexity analysis.
- Java implementation of BST.
- Build a BST.
- An application of BST.

Binary Search Trees (BST)

- For simplicity, let assume that each node in the tree stores an integer and also assume that all the items are distinct.
- **Definition:** A **binary tree** is a **binary search tree** is that for every node, X , in the tree:
 - I. the values of all the items in its **left subtree** are **smaller than the item in X** , and,
 - II. the values of all the items in its **right subtree** are **larger than the item in X** .
- **Examples:**



Binary Search Tree



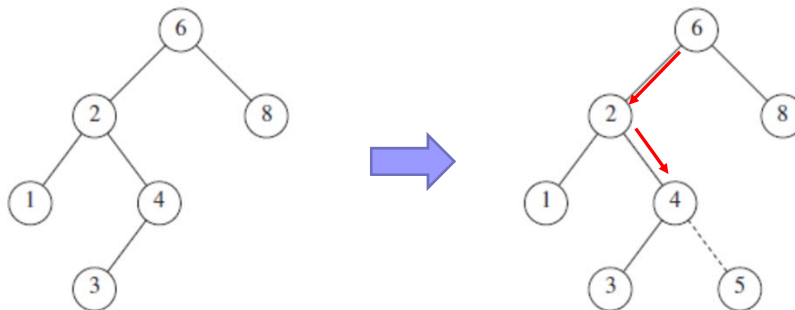
Binary Tree

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Binary Search Trees. Insertion

- To **insert/add** x into **BST** T proceed down the tree as you would with a contains.
 - (i) If x is found, do nothing (or "update" something).
 - (ii) Otherwise, insert x at the last spot on the path traversed.

- **Example:** insert(5)



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Binary Search Trees. Insertion Java Code

```
/**
 * Internal method to insert into a subtree.
 * @param x the item to insert.
 * @param t the node that roots the subtree.
 * @return the new root of the subtree.
 */
private BinaryNode<AnyType> insert( AnyType x, BinaryNode<AnyType> t )
{
    if( t == null )
        return new BinaryNode<>( x, null, null );
    int compareResult = x.compareTo( t.element );
    if( compareResult < 0 ) //goto left subtree
        t.left = insert( x, t.left );
    else if( compareResult > 0 ) //goto right subtree
        t.right = insert( x, t.right );
    else
        ; // Duplicate; do nothing
    return t;
}
```

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Binary Search Trees. Remove

- As is common with many data structures, the hardest operation is deletion.

(i) If the node is a leaf, it can be deleted immediately.

(ii) If the node has one child, the node can be deleted after its parent adjusts a link to bypass the node

□ **Example:** Case (ii): remove(4)

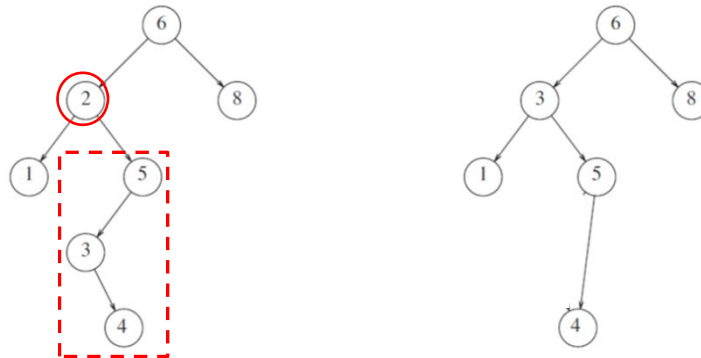


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Binary Search Trees. Remove

- (iii) If the node has two children.
 - Replace the data of this node with the smallest data of the right subtree (which is easily found)
 - Recursively delete that node (which is now empty).

- **Example:** Case (iii): remove(2)

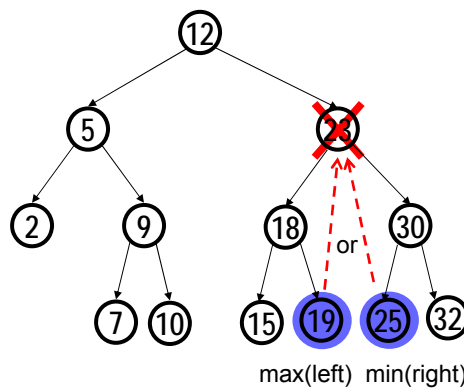


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Binary Search Trees. Remove

- (iii) If the node has two children.
 - replace the data of this node with the smallest data of the right subtree (which is easily found)
 - Recursively delete that node (which is now empty).

- **Example (alternative):** Case (iii): remove(23)



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Binary Search Trees. Remove Java Code



```
/**
 * Internal method to remove from a subtree.
 * @param x the item to remove.
 * @param t the node that roots the subtree.
 * @return the new root of the subtree.
 */
private BinaryNode<AnyType> remove( AnyType x, BinaryNode<AnyType> t )
{
    if( t == null )
        return t; // Item not found; do nothing
    int compareResult = x.compareTo( t.element );
    if( compareResult < 0 )
        t.left = remove( x, t.left );
    else if( compareResult > 0 )
        t.right = remove( x, t.right );
    else if( t.left != null && t.right != null ) // Two children
    {
        t.element = findMin( t.right ).element;
        t.right = remove( t.element, t.right );
    }
    else
        t = ( t.left != null ) ? t.left : t.right;
    return t;
}
```

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Binary Search Trees. Java Code



- See **Java** Code implementation of **Binary Search Tree** in:

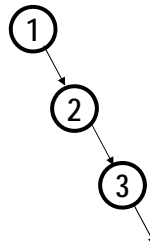
http://users.cis.fiu.edu/~weiss/cop3530_sum09/Day13.java

(Prof. Mark Weiss Implementation)

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Insert & Remove. Worst case running time.

- For the **BST** – **Find**, **Insert** and **Remove** are $O(N)$ (**worst case**)
- **Example:**



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Building a BST (buildBST)

- Suppose we need to build a **BST** from the following ordered sequence of integers:

1, 2, 3, 4, 5, 6, 7, 8, 9

- How to insert this sequence into an **empty BST**?

- **First option:** insert the elements in a given order.

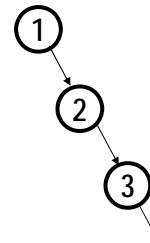
What's the tree?

Not a good tree!

we want to avoid height = N (max height)

What the running time?

$O(N^2)$



- **Second option:** insert the elements in the reverse order?

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Building a BST (buildBST)

- Suppose we need to build a **BST** from the following ordered sequence of integers:

1, 2, 3, 4, 5, 6, 7, 8, 9

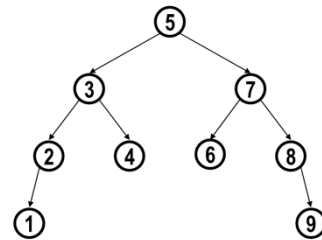
- How to insert this sequence into an **empty BST**?
 - **Third option:** re-arrange elements: median first, then left median, right median, etc. (5, 3, 7, 2, 1, 4, 8, 6, 9)

What's the tree?

Good tree! (balanced tree)

What the running time?

$O(N \log N)$ for building the **BST!!!**



- **Important:** The expected depth of any node is $O(\log N)$

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Applications of Binary Trees (BT) and BST

- A **BST** is a prominent data structure used in many systems programming applications for representing and managing dynamic sets.
- The **Average case** complexity of **Find**, **Insert**, and **Remove** operations is $O(\log N)$, where N is the number of nodes in the tree.
- **Examples:**
 - A **Binary Search Partition** is used in almost every 3D video game to determine what objects need to be rendered.
 - **BST** is used in Unix kernels for managing a set of virtual memory areas.
 - **BST** may also be used to solve some of database problems, for example, indexing.
 - . . .

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