


Modern Birkhäuser Classics


Logic for Computer Scientists  
Uwe Schöningh

Prolog Programming for Artificial Intelligence  
Peter Bratko




FLORIDA  
INTERNATIONAL  
UNIVERSITY

# Logic Programming and PROLOG (III)


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Note: The most of the information of these slides was extracted and adapted from Bratko's book, "Prolog Programming for Artificial Intelligence". They are provided for COT-3541 students only. Not to be published or publicly distributed without permission by the publisher.



## PROLOG. Our previous class. Summary



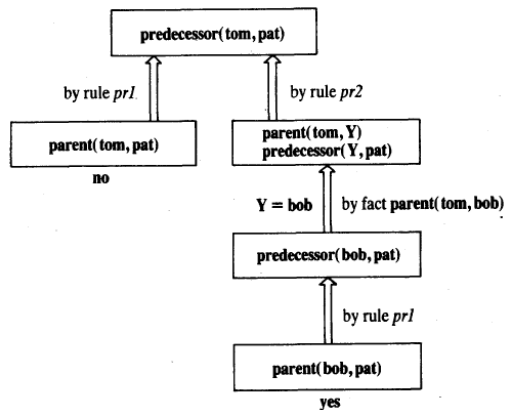
- How **PROLOG** answers questions?
  - **PROLOG** accepts **facts** and **rules** as a set of axioms, and the user's question/query as a conjectured **theorem**; then
  - **PROLOG** tries to prove this theorem - that is, to demonstrate that it can be logically derived from the axioms.
  - The **inference process** is done automatically by the **PROLOG** system and is, in principle, hidden from the user.

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## PROLOG. Our previous class. Summary

How **PROLOG** answers questions?

?- **predecessor**(tom, pat).



```
parent(pam, bob).
parent(tom, bob).
parent(tom, liz).
parent(bob, ann).
parent(bob, pat).
parent(pat, jim).
female(pam).
male(tom).
male(bob).
female(liz).
female(ann).
female(pat).
male(jim).
different(X, Y) :- X \== Y.
offspring(Y, X) :- parent(X, Y).
mother(X, Y) :- parent(X, Y),
                female(X).
grandparent(X, Z) :- parent(X, Y),
                    parent(Y, Z).
sister(X, Y) :- parent(Z, X),
               parent(Z, Y),
               female(X),
               different(X, Y).
pr1 predecessor(X, Z) :- parent(X, Z).
pr2 predecessor(X, Z) :- parent(X, Y),
                        predecessor(Y, Z).
```

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## PROLOG. Our previous class. Summary



### Monkey and Banana



- The problem:
  - There is a **monkey** at the door into a room.
  - In the middle of the room a **banana** is hanging from the ceiling.
  - The monkey is hungry and wants to get the banana, but he cannot stretch high enough from the floor.
  - At the window of the room there is a **box** the monkey may use.

**Can the monkey get the banana?**

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## PROLOG. Our previous class. Summary



### Monkey and Banana (cont...)



- The **monkey** can perform the following **actions**:
  - **Walk** on the floor
  - **Climb** the box
  - **Push** the box around (if it is already at the box)
  - **Grasp** the banana if standing on the box directly under the banana.

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## PROLOG. Our previous class. Summary.



### Monkey and Banana (cont...)



- **Initial State:**
  - Monkey is at the door
  - Monkey is on floor
  - Box is at window
  - Monkey does not have banana
- In **PROLOG**
  - state(atdoor, onfloor, atwindow, hasnot).**
- The **goal of the game** is a situation in which the monkey has the banana; that is, any state in which the last component is **has**:

**state(\_ \_ \_ has).**

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## PROLOG. Our previous class. Summary.



### Monkey and Banana (cont...)



- Move from one state to another  
**move(State1, M, State2)**
  - **State1** is the state before the move.
  - **M** is the move executed and
  - **State2** is the state after the move.
- The move '**grasp**', can be specified in **PROLOG** as:  

```
move(state(middle,onbox,middle,hasnot), % Before move
      grasp, % Move
      state(middle,onbox,middle,has)). % After move
```

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## PROLOG. Our previous class. Summary.



### Monkey and Banana (cont...)



- The other two types of moves, '**push**' and '**climb**', can be similarly specified.
- Example:
  - The move "**push**":  

```
move(state(P1, onfloor, P1, H),
      push(P1, P2),
      state(P2, onfloor, P2, H)).
```
  - The move "**climb**":  

```
move(state(P, onfloor, P, H),
      climb,
      state(P, onbox, P, H)).
```

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## PROLOG. Our previous class. Summary.



### Monkey and Banana (cont...)



- Main question our program will pose:  
**Can the monkey in some initial state get the banana?**
- In terms of PROLOG predicate:  
**canget(State).**  
where the argument State is a state of the monkey world.

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## PROLOG



- Then the **Monkey and Banana PROLOG program** is:

```
move(state(middle,onbox,middle,hasnot),
      grasp,
      state(middle,onbox,middle,has)).

move(state(P,onfloor,P,H),
      climb,
      state(P,onbox,P,H)).

move(state(P1,onfloor,P1,H),
      push(P1,P2),
      state(P2,onfloor,P2,H)).

move(state(P1,onfloor,B,H),
      walk(P1,P2),
      state(P2,onfloor,B,H)).

% change(State): monkey can get banana is State

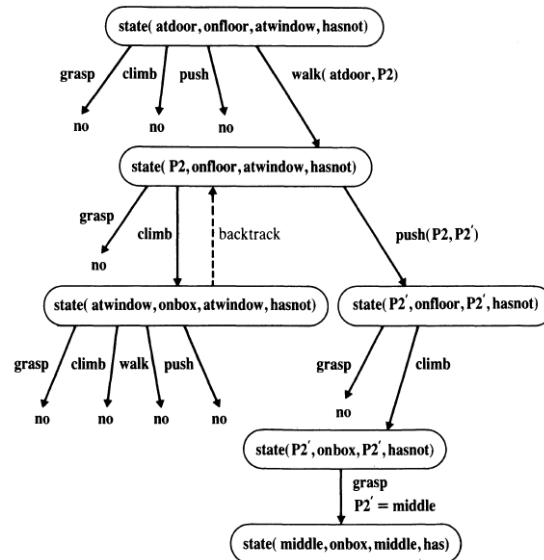
canget(state(_,_,_,has)).           % can 1: Monkey already has it

canget(State1) :-                   % can2: Do some work to get it
    move(State1,Move,State2),      % Do something
    canget(State2).                 % Get it now
```

## PROLOG

### Monkey and Banana (cont...)

- ?- canget(state(atdoor,onfloor,atwindow,hasnot)).



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## PROLOG

### Order of Clauses and Goal

- Consider the following clause:  
 $p \text{ :- } p.$
- This says that "p is true if p is true".
- This is **declarative perfectly** correct but **procedurally is quite inoperable**.
- In fact, such a clause can cause problems to **PROLOG**.
- Consider the question:  
 $? - p.$
- Using this clause, the goal **p** is replaced by the same goal **p**; this will be in turn replaced by **p**, etc.
- **PROLOG** will enter an **infinite loop!!!**

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## PROLOG

### Order of Clauses and Goal (cont...)



- In our **monkey\_and\_banana PROLOG** program we have the following clause order:
  - **Grasp**
  - **Climb**
  - **Push**
  - **Walk**
- Effectively says that the monkey prefers grasping to climbing, climbing to pushing, etc...
- This order of preferences helps the monkey to solve the problem.

**But what could happen if the order was different?**

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## PROLOG

### Order of Clauses and Goal (cont...)



**But what could happen if the order was different?**

- Let assume that the new clause order is:

**Walk – Grasp – Climb – Push**
- Then the execution of our original goal:

**?- canget(state(atdoor, onfloor, atwindow, hasnot)).**

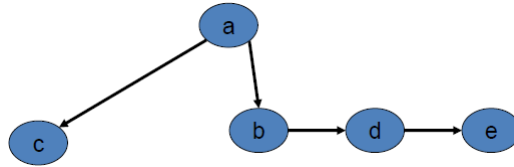
**This results in an infinite loop!**
- As the first move the monkey chooses will always be move, therefore he moves aimlessly around the room.
- **Conclusion:**
  - A program in PROLOG may be declaratively correct, but procedurally incorrect (i.e. Unable to find a solution when a solution actually exists).

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## Exercise (homework):

### Path finding in in a directed acyclic graph:

```
edge(a,b).  
edge(a,c).  
edge(b,d).  
edge(d,e).
```



```
path(X,Y) :- path(X,Z), edge(Z,Y).  
path(X,Y) :- edge(X,Y).
```

- (a) If the above **PROLOG** program is problematic then, what is the correct solution?
- (b) What if the graph is undirected?

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## List processing in PROLOG

### Representation of lists:

- The **list** is a simple **data structure** commonly used in non-numeric programming.

**Example:** A list is a sequence of any number of items, such as:

**ann, tennis, tom, ball**

- This list can be written in **PROLOG** as:

**[ann, tennis, tom, ball]**

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## List processing in PROLOG



How can a list be represented as standard PROLOG object?

- For this we need to consider two cases:
  - the list is **empty** or
  - the list is **non-empty**
- In the case of the empty list we write **[]**
- In the second case, the list can be viewed as consisting of:
  - The first item, called the **head** of the list;
  - The remaining part of the list, called the **tail**.
- For our example list:  
**[ann, tennis, tom, ball]**
- The **head** is **ann** and the **tail** of the list is **[tennis, tom, ball]**

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## List processing in PROLOG



How can a list be represented as standard PROLOG object?

- Let **L = [a,b,c]**
- **PROLOG** provides a notational extension, the **vertical bar**, which separates the **Head** and the **Tail**:  
**L = [a | Tail]**
- We can list any number of elements followed by **'|'** and the list of remaining items.
- So, alternative ways of writing the above list are:  
**[a,b,c] = [a | [b,c]] = [a,b | [c]] = [a,b,c | []]**

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## List processing in PROLOG



To summarize:

- A **list** is a data structure that is either empty or consists of two parts: a **head** and a **tail**. The **tail** itself has to be a list.
- **PROLOG** provides a special notation for lists, thus accepting lists written as:

**[Item1, Item2, ...]**

or

**[Head | Tail ]**

or

**[Item1, Item2, ... | Others]**

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## List processing in PROLOG



**Some operations on Lists. Membership**

- Now we try to implement the **membership** relation as:

**member(X,L)**

where **X** is an object and **L** is a list.

- The goal **member(X,L)** is **True** if **X** occurs in **L**.
- For example:

**member(b,[a,b,c])** is **True**

**member(b,[a,[b,c]])** is **False**

**member([b,c],[a,[b,c]])** is **True**

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## List processing in PROLOG



### Some operations on Lists. Membership (cont...)

- The program for the **membership relation** can be based on the following observation:
  - **X** is a **member** of **L** if either:
    - (1) **X** is the **head** of **L**, or
    - (2) **X** is a member of the **tail** of **L**.
- In **PROLOG** this can be written in two clauses: the first is a simple **fact** and the second is a **rule**:

```
member(X, [X|Tail]).
```

```
member(X, [Head|Tail]) :- member(X, Tail).
```

- One can read the clauses the following way, respectively:
  - **X** is a member of a list whose first element is **X**.
  - **X** is a member of a list whose tail is **Tail** if **X** is a member of **Tail**.

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## List processing in PROLOG



### Some operations on Lists. Membership (cont...)

```
member(X, [X|Tail]).
```

```
member(X, [Head|Tail]) :- member(X, Tail).
```

- **Example:**

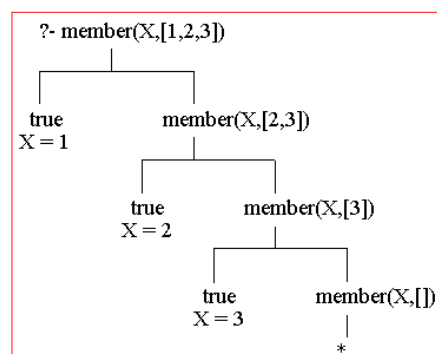
```
?- member(X,[1,2,3]).
```

```
X = 1 ;
```

```
X = 2 ;
```

```
X = 3 ;
```

```
false
```



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## List processing in PROLOG



Some operations on Lists. Membership (cont...)

```
member(X, [X|Tail]).
```

```
member(X, [Head|Tail]) :- member(X, Tail).
```

- **Example:**

```
?- member([3,Y], [[1,a],[2,m],[3,z],[4,v],[3,p]]).
```

```
Y = z ;
```

```
Y = p ;
```

```
false
```

```
?- member(X,[23,45,67,12,222,19,9,6]), Y is X*X, Y < 100.
```

```
X = 9   Y = 81 ;
```

```
X = 6   Y = 36 ;
```

```
false
```

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## List processing in PROLOG



Some operations on Lists. Membership (cont...)

```
member(X, [X|Tail]).
```

```
member(X, [Head|Tail]) :- member(X, Tail).
```

- The definition for the predicate **member** can be written as:

```
member(X,[X|_]).
```

```
member(X,[_|Tail]) :- member(X, Tail).
```

- Remember that '\_' (underscore) designates a "don't-care" variable, usually called **anonymous** variables.
- In general, such variables have names whose first character is the underscore.
- Not having to bind values to anonymous variables saves a little run-space and run-time.

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## List processing in PROLOG



### Some operations on Lists. Concatenation (Append)

- For **concatenation of lists** we will define the relation:

**conc(L1, L2, L3)**

where L1 and L2 are two lists, and L3 is their concatenation.

- For **example**:

**conc([a,b],[c,d],[a,b,c,d])** is **True**

**conc([a,b],[c,d],[a,b,d,c])** is **False**

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## List processing in PROLOG



### Some operations on Lists. Concatenation (cont...)

**conc(L1, L2, L3)**

- For **conc** we have **two cases**, depending on the first argument, L1:

(1) if the 1<sup>st</sup> argument is the empty list then the 2<sup>nd</sup> and the 3<sup>rd</sup> arguments must be the same. This is expressed by the following **PROLOG** fact:

➡ **conc([], L, L).**

(2) If the 1<sup>st</sup> argument of **conc** is non-empty list then it has a **head** and a **tail** and must look like this:

**[X | L1]**

- The result of the concatenation is the list **[X | L3]** where L3 is the concatenation of L1 and L2. In **PROLOG**:

➡ **conc([X | L1], L2, [X | L3]) :- conc(L1,L2,L3).**

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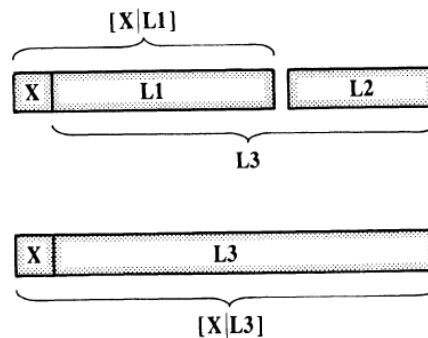
## List processing in PROLOG

Some operations on Lists. Concatenation (cont...)

```
conc([], L, L).
```

```
conc([X | L1], L2, [X | L3]) :- conc(L1, L2, L3).
```

- Illustration of the concatenation of  $[X | L1]$  and some list  $L2$ .



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## List processing in PROLOG

Some operations on Lists. Concatenation (cont...)

```
conc([], L, L).
```

```
conc([X | L1], L2, [X | L3]) :- conc(L1, L2, L3).
```

- Example:

```
?- conc([a,b,c],[1,2,3],L).
```

```
L = [a,b,c,1,2,3]
```

```
?- conc([a,[b,c],d], [a,[],b], L).
```

```
L=[a,[b,c],d,a,[],b]
```

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## List processing in PROLOG

Some operations on Lists. Concatenation (cont...)

```
conc([], L, L).
conc([X | L1], L2, [X | L3]) :- conc(L1, L2, L3).
```

Example:

```
conc([1,2],[3,4],A).
```



```
conc([X|L1],L2,[X|L3]) :- conc(L1,L2,L3).
conc([1|[2]],[3,4],[1|V1]) :- conc([2],[3,4],V1).    A=[1|V1]
```



```
conc([X|L1],L2,[X|L3]) :- conc(L1,L2,L3).
conc([2|[]],[3,4],[2|V2]) :- conc([], [3,4], V2).    V1=[2|V2]
```



```
conc([],L,L).
conc([], [3,4], [3,4]).                                V2 =[3,4]
```

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## List processing in PROLOG

Example: `conc([], L, L).`  
`conc([X | L1], L2, [X | L3]) :- conc(L1, L2, L3).`

```
SWI-Prolog (Multi-threaded, version 6.6.6)
File Edit Settings Run Debug Help
Welcome to SWI-Prolog (Multi-threaded, 32 bits, Version 6.6.6)
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For help, use ?- help(Topic). or ?- apropos(Word).

1 ?-
Warning: c:/users/leslie/documents/fiu/spring_2016/cot3541_spring_2016/prolog/list_
Singleton variables: [Tail]
Warning: c:/users/leslie/documents/fiu/spring_2016/cot3541_spring_2016/prolog/list_
Singleton variables: [Head]
Warning: c:/users/leslie/documents/fiu/spring_2016/cot3541_spring_2016/prolog/list_
Singleton variables: [L1,L2]
Warning: c:/users/leslie/documents/fiu/spring_2016/cot3541_spring_2016/prolog/list_
Singleton variables: [L1,L3]
% c:/Users/leslie/Documents/FIU/Spring_2016/COT3541_Spring_2016/Prolog/list_operati
es
1 ?- trace.
true.
[trace] 1 ?- conc([a,b,c],[d,e,f],X).
Call: (6) conc([a, b, c], [d, e, f], _G532) ? creep
Call: (7) conc([b, c], [d, e, f], _G617) ? creep
Call: (8) conc([c], [d, e, f], _G620) ? creep
Call: (9) conc([], [d, e, f], _G623) ? creep
Exit: (9) conc([], [d, e, f], [d, e, f]) ? creep
Exit: (8) conc([c], [d, e, f], [c, d, e, f]) ? creep
Exit: (7) conc([b, c], [d, e, f], [b, c, d, e, f]) ? creep
Exit: (6) conc([a, b, c], [d, e, f], [a, b, c, d, e, f]) ? creep
X = [a, b, c, d, e, f].
[trace] 2 ?- notrace.
true.
[debug] 3 ?- ■
```

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## List processing in PROLOG

### Some operations on Lists. Concatenation (cont...)

- We can use **conc** in the inverse direction for decomposing a given list into two list.

- Example:**

?- conc(L1,L2,[a,b,c]).

conc([], L, L).

conc([X | L1], L2, [X | L3]) :- conc(L1,L2,L3).

L1=[]  
L2=[a,b,c];

L1=[a]  
L2=[b,c];

L1=[a,b]  
L2=[c];

L1=[a,b,c]  
L2=[];  
false

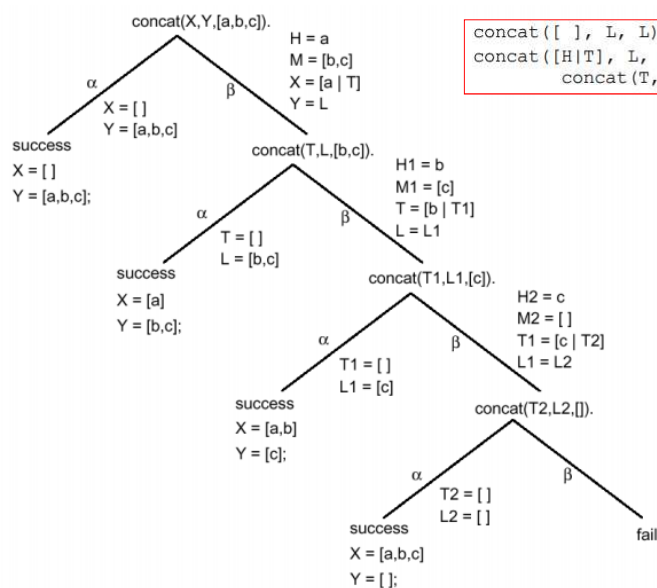
- It is possible to decompose the list **[a,b,c]** in four ways, all of which were found by our program through backtracking.



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## List processing in PROLOG

### Some operations on Lists. Concatenation (cont...)



concat([ ], L, L). % clause a  
concat([H|T], L, [H|M]) :-  
concat(T, L, M). % clause b



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## List processing in PROLOG



### Some operations on Lists. Concatenation (cont...)

- We can use **conc** to look for a certain pattern in a list.
- **Example:** we can find the months "before" and "after" a given month:

```
?- conc(Before, [may | After],  
[jan,feb,mar,apr,may,jun,jul,aug,sep,oct,nov,dec]).
```

Before = [jan,feb,mar,apr]

After = [jun,jul,aug,sep,oct,nov,dec]

- Using **conc**, we can find the immediate month before and the immediate month after May by asking:

```
?- conc(_, [Month1,may,Month2 | _],  
[jan,feb,mar,apr,may,jun,jul,aug,sep,oct,nov,dec]).
```

Month1 = apr

Month2 = jun

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## List processing in PROLOG



### Some operations on Lists. Concatenation (cont...)

```
conc([], L, L).
```

```
conc([X | L1], L2, [X | L3]) :- conc(L1,L2,L3).
```

- Observe that we can use **conc** to delete from some list, L1, everything that follows three successive occurrences of z in L1 together with the three z's.

- **Example:**

```
?- L1 = [a,b,z,z,c,z,z,z,d,e],  
conc(L2,[z,z,z | _], L1).
```

L1 = [a,b,z,z,c,z,z,z,d,e],

L2 = [a,b,z,z,c]

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## List processing in PROLOG

Some operations on Lists. Concatenation (cont...)

```
conc([], L, L).
```

```
conc([X | L1], L2, [X | L3]) :- conc(L1, L2, L3).
```

- Note that using **conc** we can redefine the **membership relation** by the clause:

```
member1(X, L) :- conc(L1, [X | L2], L).
```

- % X is a member of list L, if L can be decomposed into two list  
% and the second one has X as its **head**
- The **member1** clause can be written using **anonymous** variables as:

```
member1(X, L) :- conc(_, [X | _], L).
```



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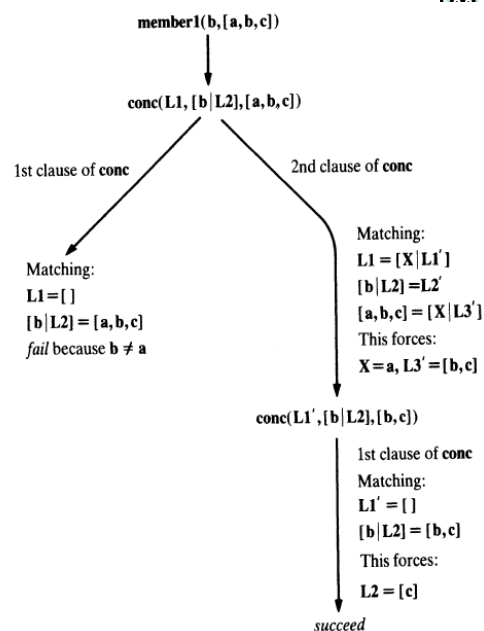
## List processing in PROLOG

- Procedure **member1** finds an item in a given list.

```
conc([], L, L).
```

```
conc([X | L1], L2, [X | L3]) :-  
    conc(L1, L2, L3).
```

```
member1(X, L) :-  
    conc(L1, [X | L2], L).
```



## List processing in PROLOG



### Some operations on Lists. Concatenation (cont...)

- Define a predicate **rotate(X,Y)** where both X and Y are represented by lists, and Y is formed by rotating X to the left by one element.

- Solution:**

- Take the first element off the first list (H) and append it after the tail (i.e. at the end) in the solution (R)

```
rotate([H|T],R) :- conc(T,[H],R).
```

```
conc([], L, L).
```

```
conc([X | L1], L2, [X | L3]) :- conc(L1,L2,L3).
```

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## List processing in PROLOG



### Some operations on Lists. Concatenation (cont...)

- Exercises (Homework):**

- Write a query, using **conc**, to **delete** the last three elements from a list **L** producing another list **L1**.

*Hint: L is the concatenation of L1 and a three-element list.*

- Define the predicate

**last(Item, List)**

so that **Item** is the last element of a list **List** using the **conc** relation/predicate.

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## List processing in PROLOG



### Some operations on Lists. Adding an item to a list

- To **add** an item to a list the easiest way is to put item in front of the list so that it becomes the new **head**.
- If X is the new item and the list is L then the result is:

**[X | L]**

- Using **PROLOG** we can define the **add** procedure as the fact:

**add(X, L, [X | L]).**

- **Example:**

**?- add(a,[1,2,3], L).**

**L = [a,1,2,3]**

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## List processing in PROLOG



### Some operations on Lists. Deleting an item

- **Deleting** an item, X, from a list, L, can be programmed as a relation:

**del(X, L, L1)**

where L1 is equal to the list with the item X removed.

- The **del** relation can be defined similarly to the membership relation. We have, again, two cases:

(1) If X is the **head** of the list then the result, after deletion, is the tail of the list.

(2) If X is in the **tail** then it is deleted from there.

**del(X, [X | Tail], Tail).**

**del(X, [Y | Tail], [ Y | Tail1]) :- del(X, Tail, Tail1).**

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## List processing in PROLOG



### Some operations on Lists. Deleting an item (cont...)

```
del(X, [X | Tail], Tail).  
del(X, [Y | Tail], [Y | Tail1]) :- del(X, Tail, Tail1).
```

- Note that
  - If there are several occurrences of X in the list then **del** will be able to delete anyone of them by backtracking.
  - each alternative execution of del will only delete one occurrence of X, leaving the other untouched.
- Example

```
?- del(a,[a,b,a],L).
```

```
L = [b,a,a];
```

```
L = [a,b,a];
```

```
L = [a,b,a];
```

```
false
```

- Note: **del** will fail if the list does not contain the item to be deleted.

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## List processing in PROLOG



### Some operations on Lists. Deleting an item (cont...)

```
del(X, [X | Tail], Tail).  
del(X, [Y | Tail], [Y | Tail1]) :- del(X, Tail, Tail1).
```

- Observe that **del** can also be used to **add** an item to a list by inserting the new item anywhere in the list.
- **Example:**
  - If we want to insert **a** at any place in the list **[1,2,3]** then then we can do this by asking the query:
  - What's **L** such that after deleting **a** from **L** we obtain **[1,2,3]**?

```
?- del(a,L,[1,2,3]).
```

```
L = [a,1,2,3];
```

```
L = [1,a,2,3];
```

```
L = [1,2,a,3];
```

```
L = [1,2,3,a];
```

```
false
```

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## List processing in PROLOG



### Sublist

- Let us now consider the **sublist** relation.
- This relation has two arguments:
  - a list **L** and
  - a list **S** such that occurs in **L** as its sublist
- So
  - sublist([c,d,e],[a,b,c,d,e,f])** is **True** and
  - sublist([c,e],[a,b,c,d,e,f])** is **False**
- The **sublist** relation can be formulated as:
  - **S** is a **sublist** of **L** if:
    - (1) **L** can be decomposed into two lists, **L1** and **L2**, and
    - (2) **L2** can be decomposed into two lists, **S** and some **L3**
- In **PROLOG** as:  
**sublist(S,L) :- conc(L1,L2,L), conc(S,L3,L2).**

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## List processing in PROLOG



### Sublist (cont...)

**sublist(S,L) :- conc(L1,L2,L), conc(S,L3,L2).**

- Observe that **sublist** relation can also be used, for example, to find all sublists of a given list:

**?- sublist(S,[a,b,c]).**

S = [];  
S = [a];  
S = [a,b]  
S = [a,b,c]  
S = [b];

...

**how many?**

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## List processing in PROLOG



### Permutation

- We will define the **permutation relation** with two lists such that one is the permutation of the other.
- **Example:**  
**?- permutation ([a,b,c],P).**  
P = [a,c,b]  
...
- The intention is to **generate permutations** of a list through **backtracking** using the permutation procedure.
- The procedure for **permutation** can be based on the consideration of two cases, depending on the first list:
  - (1) If the first list is empty then the second list must also be empty.
  - (2) If the first list is not empty then it has the form [X|L], and the permutation can be constructed as: first permuted L obtaining L1, and the insert X at any position into L1.

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## List processing in PROLOG



### Permutation (cont...)

- ...
  - (1) If the first list is empty then the second list must also be empty.
  - (2) If the first list is not empty then it has the form [X|L], and the permutation can be constructed as:
    - first permuted L obtaining L1, and
    - insert X at any position into L1.
- The **PROLOG** clauses for the permutation relation are:  
**permutation([], []).**  
**permutation([X|L],P) :-**  
    **permutation(L,L1),**  
    **insert(X,L1,P).**

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## List processing in PROLOG



### Permutation (cont...)

```
permutation([], []).  
permutation([X | L], P) :-  
    permutation(L, L1),  
    insert(X, L1, P).
```

#### ■ Example:

```
?- permutation([red,blue,green],P).
```

```
P = [red,blue,green];
```

```
P = [red,green,blue];
```

```
P = [blue,red,green];
```

```
P = [blue,green,red];
```

```
P = [green,red,blue];
```

```
P = [green,blue,red];
```

```
false
```