

Lecture 16 — October 21

Lecturer: Sanghavi

Scribe: Jimmy Lin and Taewan Kim

16.1 Recap

In the last lecture, dual of *Semi-Definite Programming (SDP)* problems with linear objective function are derived as

$$\begin{aligned} \min \quad & -\langle G, Z \rangle \\ \text{s.t.} \quad & \langle F_i, Z \rangle = c_i, \quad \forall i \\ & z \succeq 0 \end{aligned} \tag{16.1}$$

Afterwards, it is shown how to formulate following problem as linear SDP optimization:

- find a matrix with largest eigenvalue,
- find sum of r -largest eigenvalues of a given matrix and
- find sum of singular values of a symmetric but not PSD matrix.

Then study of SDP is then extended to non-linear objective, which is called *Log Determinant Optimization*. The general form is as follows:

$$\begin{aligned} \min_x \quad & c^T x - \log \det G(x) \\ \text{s.t.} \quad & G(x) \succeq 0 \\ & F(x) \geq 0 \end{aligned} \tag{16.2}$$

where $-\log \det G(x)$ is proven to be convex function.

Also, the following problems are formulated as Log Determinant Optimization problem:

- find the minimal-volume ellipsoid that contains all given points,
- find the maximum-volume ellipsoid enclosed within a given polyhedron,
- find the most likely parameters that generates a given set of samples
- find the variance matrix of gaussian channel with maximum capacity.

In this lecture, we will discuss *Proximal Gradient Algorithm*. Section 16.2 illustrates motivations for the Proximal Gradient Algorithm. Section 16.3 will provide introduction and definition of the Proximal Operator. Section 16.4 gives details of the Proximal Gradient Algorithm. Section 16.5 is to touch convergence analysis of the Proximal Gradient Algorithm.

16.2 Motivations

As has been seen before, in order to get to $\|x - x^*\| \leq \epsilon$, we have complexity

- $\mathcal{O}(\frac{1}{\epsilon})$ for gradient descent method with *smooth* objective $f(\cdot)$
- $\mathcal{O}(\log(\frac{1}{\epsilon}))$ for gradient descent method with *strongly convex* objective $f(\cdot)$
- $\mathcal{O}(\frac{1}{\epsilon^2})$ for subgradient descent method with *non-smooth* objective $f(\cdot)$

16.2.1 Iterative Shrinkage-Thresholding Algorithm (ISTA)

Consider the unconstrained optimization problem with l_1 regularization

$$\min_x \quad \underbrace{\frac{1}{2} \|y - Ax\|_2^2}_{\text{smooth/"nice"}} + \underbrace{\lambda \|x\|_1}_{\text{not smooth/"nice", but "special"}} \quad (16.3)$$

Say if $A = I$, then

$$\begin{aligned} \min_x \quad & \frac{1}{2} \|y - x\|_2^2 + \lambda \|x\|_1 \\ = \min_x \quad & \sum_i \left\{ \frac{1}{2} (y_i - x_i)^2 + \lambda |x_i| \right\} \end{aligned} \quad (16.4)$$

Obviously, such transformation is justifiable since the original problem is separable. Now we derive closed-form solution for each "separated" problem:

$$x_i - y_i + \lambda \quad \text{if } x_i > 0 \quad (16.5)$$

$$x_i - y_i - \lambda \quad \text{if } x_i < 0 \quad (16.6)$$

$$-y_i + r\lambda, r \in (-1, 1) \quad \text{if } x_i = 0 \quad (16.7)$$

Suppose $y_i > 0$, then it obvious that $x_i \geq 0$.

- if $x_i > 0$, then $\exists x_i = y_i - \lambda$.
- if $x_i = 0$, then

Similarly, if $y_i < 0$, then exists $x_i = y_i + \lambda$ if $x_i < 0$.

16.2.2 Motivation 2

Consider another optimization problem

$$\begin{aligned} \min_x \quad & \underbrace{g(x)}_{\text{"nice"}} \\ \text{s.t.} \quad & x \in Q \end{aligned} \quad (16.8)$$

where Q is a convex set and is easy to project onto, e.g. $Q : \{x \mid \|x\|_\infty \leq 1\}$.
The above optimization problem equates to

$$\min_x g(x) + I_Q(x) \quad (16.9)$$

where

$$I_Q(x) = \begin{cases} 0 & \text{if } x \in Q \\ \infty & \text{if } x \notin Q \end{cases} \quad (16.10)$$

Hence, we have

$$[P_Q(y)]_i = \begin{cases} y_i & \text{if } |y_i| \leq 1 \\ \text{sign}(y_i) & \text{otherwise} \end{cases} \quad (16.11)$$

16.3 Proximal Operator

Proximal gradient is an algorithm for unconstrained problems with a cost function which can be expressed sum of two functions.

$$\text{minimize } f(x) = g(x) + h(x) \quad (16.12)$$

$g(x)$: 'nice', (ex) convex, L -lipschitz gradient

$h(x)$: 'special', (ex) convex, not smooth, **prox** is easy to calculate

In the above equation, we used the term 'nice' and 'special' to express the characteristic of each function g and h . Specific conditions on g and h will be more clear in the theorem which will be provided later. And the following is a definition of prox function (proximal operator) which clarifies the meaning of **prox** in the condition of h .

Definition 1. (Prox function/ Proximal operator)

The prox function (or proximal operator) of a function $h(\cdot)$ is defined as:

$$\text{prox}_{th}(v) = \arg \min_x \left(h(x) + \frac{1}{2t} \|x - v\|_2^2 \right) \quad (16.13)$$

One natural example of proximal operator is projection on set Q . Suppose you perform a gradient descent and require a solution to be in a certain set Q . So, in each step, x_+ should satisfy the condition $x_+ \in Q$, and this can be done by using projection P_Q as in Figure 16.1.

$$x^{(k+1)} = P_Q(x^{(k)} - t\nabla g(x^{(k)})) \quad (16.14)$$

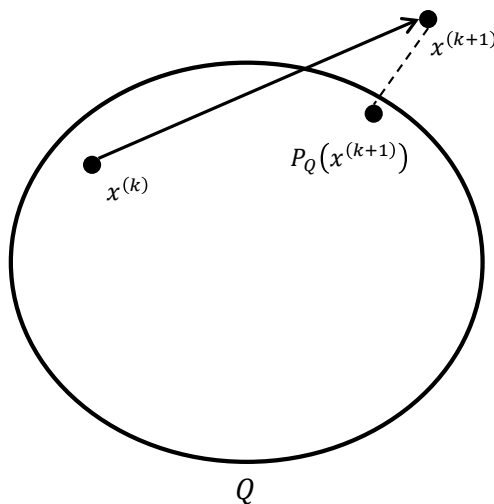


Figure 16.1. Projection on Q for gradient step

And this projection operator is a specific version of proximal operator by using $h(x) = I_Q(x)$ where $I_Q(\cdot)$ is a indicator function of Q .

$$\begin{aligned}
 \mathbf{prox}_h(x) &= \arg \min_u \left(h(u) + \frac{1}{2} \|u - x\|_2^2 \right) \\
 &= \arg \min_u \left(I_Q(u) + \frac{1}{2} \|u - x\|_2^2 \right) \\
 &= \arg \min_{u \in Q} \left(\frac{1}{2} \|u - x\|_2^2 \right) \\
 &= P_Q(x)
 \end{aligned}$$

where indicator function I_Q is defined as,

$$I_Q(x) = \begin{cases} 0 & \text{if } x \in Q \\ \infty & \text{if } x \notin Q \end{cases}$$

Third equality comes from the fact that $u \notin Q$ gives infeasible solution x for minimizing the function. And the last equality is the general definition of projection based on the Euclidean distance.

16.4 Proximal Gradient Algorithm

Now we can introduce a proximal gradient algorithm which uses two black boxes, one outputs $\nabla g(x)$ and the other outputs $\mathbf{prox}_{th}(v)$ (Figure 16.2), for the task of minimizing $f(x) = g(x) + h(x)$.

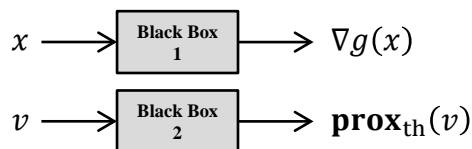


Figure 16.2. Using two black boxes for proximal gradient

- Proximal Gradient Algorithm

Proximal gradient algorithm is defined as the following update rule.

$$x_+ \leftarrow \mathbf{prox}_{th}(x - t\nabla g(x)) \quad (16.15)$$

By using the definition of proximal operator, update of proximal gradient can be expressed as follows:

$$\begin{aligned}
 x_+ &\leftarrow \mathbf{prox}_{th}(x - t\nabla g(x)) \\
 &= \arg \min_u \left(h(u) + \frac{1}{2t} \|u - x + t\nabla g(x)\|_2^2 \right) \\
 &= \arg \min_u \left(h(u) + \langle \nabla g(x), u - x \rangle + \frac{1}{2t} \|u - x\|_2^2 + \frac{t}{2} \|\nabla g(x)\|_2^2 \right) \\
 &= \arg \min_u \left(h(u) + \langle \nabla g(x), u - x \rangle + \frac{1}{2t} \|u - x\|_2^2 + g(x) \right) \quad (16.16)
 \end{aligned}$$

Last equality comes from the fact that $\|\nabla g(x)\|_2^2$ and $g(x)$ do not depend on u . In equation (16.4), an important fact to point out is that rear part has a form of quadratic approximation of $g(u)$ around x .

$$g(x) + \langle \nabla g(x), u - x \rangle + \frac{1}{2t} \|u - x\|_2^2 \quad (16.17)$$

16.5 Convergence Analysis

The convergence of proximal gradient algorithm requires $O(1/\epsilon)$ of iteration. Following theorem specifies the condition and time complexity.

Theorem 16.1. *If g is convex and g has L -lipschitz gradient, i.e. $\|\nabla g(x) - \nabla g(y)\|_2 \leq L\|x - y\|_2$, using fixed size $t < 1/L$ on proximal gradient algorithm gives $O(1/\epsilon)$ of convergence to the optimal point x^* , where $x^* = \arg \min_x (g(x) + h(x))$.*

Before providing a proof of the theorem, let's define a function $G_t(x)$ which satisfies the following condition. Role of the $G_t(x)$ function is to remove the proximal operator on the

whole equation and make it similar to the general update as before.

$$\begin{aligned} x_+ &\leftarrow x - tG_t(x) = \mathbf{prox}_{th}(x - t\nabla g(x)) \\ \Leftrightarrow G_t(x) &\triangleq \frac{1}{t}(x - \mathbf{prox}_{th}(x - t\nabla g(x))) \end{aligned}$$

Claim 1. $G_t(x) - \nabla g(x) \in \partial h(x - tG_t(x))$

Lemma 16.2. For any point z , $f(x_+) \leq f(z) + \langle G_t(x), x - z \rangle - \frac{t}{2}\|G_t(x)\|_2^2$

Proof of Claim 1 and Lemma 16.2 is provided after the proof of Theorem 16.1.

Proof (Theorem 16.1): Putting $z = x$ in the Lemma 16.2 gives,

$$f(x_+) \leq f(x) - \frac{t}{2}\|G_t(x)\|_2^2$$

So, this gives the decreasing property $f(x_+) \leq f(x)$ since step size satisfies $t \geq 0$. And putting $z = x^*$ gives,

$$\begin{aligned} f(x_+) &\leq f^* - \frac{t}{2}\|G_t(x)\|_2^2 + \langle G_t(x), x - x^* \rangle \\ \Leftrightarrow f(x_+) - f^* &\leq \frac{1}{2t} [\|x - x^*\|_2^2 - \|x - x^* - tG_t(x)\|_2^2] \\ &= \frac{1}{2t} [\|x - x^*\|_2^2 - \|x_+ - x^*\|_2^2] \end{aligned}$$

Adding up over T iterations gives

$$\begin{aligned} \sum_{k=1}^T (f(x^{(k)}) - f^*) &\leq \frac{1}{2t} [\|x^{(0)} - x^*\|_2^2 - \|x^{(T)} - x^*\|_2^2] \\ &\leq \frac{1}{2t}\|x^{(0)} - x^*\|_2^2 \end{aligned}$$

From the fact that $f(x_+) \leq f(x)$ as shown above, $f(x^{(k)}) - f^* \geq f(x^{(T)}) - f^*$ for $k = 1, \dots, T$. And this gives,

$$f(x^{(T)}) - f^* \leq \frac{1}{2tT}\|x^{(0)} - x^*\|_2^2$$

which concludes the proof of $O(1/\epsilon)$ convergence, i.e. it requires $k = O(1/\epsilon)$ of iteration number to have $f(x^{(k)}) - f^* \leq \epsilon$.

□

Proof (Lemma 16.2): From the L-lipschitz continuous of gradient on g ,

$$g(y) \leq g(x) + \langle \nabla g(x), y - x \rangle + \frac{L}{2} \|y - x\|_2^2$$

So, for $y = x_+ = x - tG_t(x)$

$$g(x_+) \leq g(x) - \langle \nabla g(x), tG_t(x) \rangle + \frac{L}{2} t^2 \|G_t(x)\|_2^2 \quad (16.18)$$

And recall the condition of subgradient. If $a \in \partial h(x_+)$ then,

$$h(z) \geq h(x_+) + \langle a, z - x_+ \rangle$$

From the Claim 1, $G_t(x) - \nabla g(x) \in \partial h(x - tG_t(x))$. By setting $a = G_t(x) - \nabla g(x)$ gives,

$$\begin{aligned} h(z) &\geq h(x_+) + \langle G_t(x) - \nabla g(x), z - x_+ \rangle \\ \Leftrightarrow h(x_+) &\leq h(z) - \langle G_t(x) - \nabla g(x), z - x_+ \rangle \end{aligned} \quad (16.19)$$

Now using $f(x_+) = g(x_+) + h(x_+)$, and adding (16.18) and (16.19):

$$\begin{aligned} f(x_+) &\leq g(x) - \langle \nabla g(x), tG_t(x) \rangle + \frac{L}{2} t^2 \|G_t(x)\|_2^2 \\ &\quad + h(z) - \langle G_t(x) - \nabla g(x), z - x_+ \rangle \\ &\leq g(z) - \langle \nabla g(x), z - x \rangle - \langle \nabla g(x), tG_t(x) \rangle + \frac{L}{2} t^2 \|G_t(x)\|_2^2 \\ &\quad + h(z) - \langle G_t(x) - \nabla g(x), z - x_+ \rangle \\ &\quad (\because g(x) \leq g(z) - \langle \nabla g(x), z - x \rangle \text{ from convexity of } g) \\ &= f(z) - \langle \nabla g(x), z - x + tG_t(x) - (z - x_+) \rangle \\ &\quad + \frac{L}{2} t^2 \|G_t(x)\|_2^2 - \langle G_t(x), z - x_+ \rangle \\ &= f(z) + \frac{L}{2} t^2 \|G_t(x)\|_2^2 - \langle G_t(x), z - x + tG_t(x) \rangle \\ &\leq f(z) + \langle G_t(x), x - z \rangle + \frac{t}{2} \|G_t(x)\|_2^2 - t \|G_t(x)\|_2^2 \quad (\because t < 1/L) \\ &= f(z) + \langle G_t(x), x - z \rangle - \frac{t}{2} \|G_t(x)\|_2^2 \end{aligned}$$

This concludes the proof of Lemma 16.2. □

Remaining part is the proof of Claim 1, i.e. $G_t(x) - \nabla g(x) \in \partial h(x - tG_t(x))$.

Proof (Claim 1): To show that $G_t(x) - \nabla g(x)$ is a subgradient of $h(x - tG_t(x)) = h(x_+)$, let's show that the following property of proximal operator is true.

$$u = \mathbf{prox}_{th}(x) \Leftrightarrow \frac{1}{t}(x - u) \in \partial h(u) \quad (16.20)$$

Proof of the property (16.20) is based on the definition of proximal operator, which means that $\mathbf{prox}_{th}(x)$ is a minimizer.

$$\begin{aligned}
& \frac{1}{t}(x - u) \in \partial h(u) \\
& \Leftrightarrow h(u) + \left\langle \frac{1}{t}(x - u), y - u \right\rangle \leq h(y), \quad \forall y \\
& \Leftrightarrow h(u) + \left\langle \frac{1}{t}(x - u), y - u \right\rangle - \frac{1}{2t}\|y - u\|_2^2 \leq h(y) - \frac{1}{2t}\|y - u\|_2^2 \leq h(y), \quad \forall y \\
& \Leftrightarrow h(u) + \left\langle \frac{1}{t}(x - u) + \frac{1}{2t}(u - y), y - u \right\rangle \leq h(y), \quad \forall y \\
& \Leftrightarrow h(u) + \frac{1}{2t}\langle 2(x - u) + (u - y), (x - u) - (x - y) \rangle \leq h(y), \quad \forall y \\
& \Leftrightarrow h(u) + \frac{1}{2t}\langle (x - u) + (x - y), (x - u) - (x - y) \rangle \leq h(y), \quad \forall y \\
& \Leftrightarrow h(u) + \frac{1}{2t}(\|u - x\|_2^2 - \|y - x\|_2^2) \leq h(y), \quad \forall y \\
& \Leftrightarrow h(u) + \frac{1}{2t}\|u - x\|_2^2 \leq h(y) + \frac{1}{2t}\|y - x\|_2^2, \quad \forall y \\
& \Leftrightarrow u = \arg \min_z \left(h(z) + \frac{1}{2t}\|z - x\|_2^2 \right) \\
& \Leftrightarrow u = \mathbf{prox}_{th}(x)
\end{aligned}$$

Now put $u = x_+ = \mathbf{prox}_{th}(x - \nabla g(x))$ in equation (16.20).

$$\begin{aligned}
& \Rightarrow \frac{1}{t}(x - \nabla g(x) - x_+) \in \partial h(x_+) \\
& \Rightarrow \frac{1}{t}(x - \nabla g(x) - x + tG_t(x)) \in \partial h(x_+) \\
& \Rightarrow G_t(x) - \nabla g(x) \in \partial h(x_+)
\end{aligned}$$

□