

Large Scale Optimization: Lecture 16

Fall 2014

The University of Texas at Austin

Recap

- Dual of *Semi-Definite Programming (SDP)* problems with linear objective function are derived as

$$\begin{aligned} \min \quad & -\langle G, Z \rangle \\ \text{s.t.} \quad & \langle F_i, Z \rangle = c_i, \quad \forall i \\ & Z \succeq 0 \end{aligned} \tag{1}$$

- Application of linear SDP optimization:
 - Find a matrix with largest eigenvalue,
 - Find sum of r -largest eigenvalues of a given matrix and
 - Find sum of singular values of a symmetric but not PSD matrix.

Recap

- Study of SDP is then extended to non-linear objective, called *Log Determinant Optimization*. The general form is as follows:

$$\begin{aligned} \min_x \quad & c^T x - \log \det G(x) \\ \text{s.t.} \quad & G(x) \succeq 0 \\ & F(x) \geq 0 \end{aligned} \tag{2}$$

where $-\log \det G(x)$ is proven to be convex function.

Recap

- Following problems are formulated as Log Determinant Optimization problem:
 - Find the minimal-volume ellipsoid that contains all given points,
 - Find the maximum-volume ellipsoid enclosed within a given polyhedron,
 - Find the most likely parameters that generates a given set of samples
 - Find the variance matrix of gaussian channel with maximum capacity.

Motivations for Proximal Gradient Descent

- Motivation 1: Iterative Shrinkage-Thresholding Algorithm (ISTA)
- Motivation 2

Iterative Shrinkage-Thresholding Algorithm (ISTA)

Consider the unconstrained optimization problem with l_1 regularization

$$\min_x \quad \underbrace{\frac{1}{2} \|y - Ax\|_2^2}_{\text{smooth/"nice"}} + \underbrace{\lambda \|x\|_1}_{\text{not smooth/"nice", but "special"}} \quad (3)$$

If $A = I$, then

$$\begin{aligned} & \min_x \quad \frac{1}{2} \|y - x\|_2^2 + \lambda \|x\|_1 \\ &= \min_x \quad \sum_i \left\{ \frac{1}{2} (y_i - x_i)^2 + \lambda |x_i| \right\} \end{aligned} \quad (4)$$

Iterative Shrinkage-Thresholding Algorithm (ISTA)

Now we derive closed-form solution for each "separated" problem:

$$\begin{aligned}x_i - y_i + \lambda &= 0 & \text{if } x_i \geq 0 \\x_i - y_i - \lambda &= 0 & \text{if } x_i < 0\end{aligned}\tag{5}$$

Suppose $|y_i| \leq \lambda$, then $\exists r \in (-1, 1)$, such that $-y_i + r\lambda = 0$.

Otherwise, in the case of $|y_i| > \lambda$, then

- if $y_i > 0$, then $\exists x_i = y_i - \lambda$ if $x_i > 0$.
- if $y_i < 0$, then exists $x_i = y_i + \lambda$ if $x_i < 0$.

Motivation 2

Consider another optimization problem

$$\begin{aligned} \min_x \quad & \underbrace{g(x)}_{\text{"nice"}} \\ \text{s.t.} \quad & x \in Q \end{aligned} \tag{6}$$

where Q is a convex set and is easy to project onto, e.g.

$Q : \{x \mid \|x\|_\infty \leq 1\}$.

The above optimization problem equates to

$$\min_x \quad g(x) + I_Q(x) \tag{7}$$

where

$$I_Q(x) = \begin{cases} 0 & \text{if } x \in Q \\ \infty & \text{if } x \notin Q \end{cases} \tag{8}$$

Hence, we have

$$[P_Q(y)]_i = \begin{cases} y_i & \text{if } |y_i| \leq 1 \\ \text{sign}(y_i) & \text{otherwise} \end{cases} \tag{9}$$

Proximal Gradient Algorithm

- Unconstrained problem with sum of two functions:

$$\text{minimize } f(x) = g(x) + h(x) \quad (10)$$

$g(x)$: '*nice*', (ex) convex, L-lipschitz gradient

$h(x)$: '*special*', (ex) convex, not smooth, **prox** is inexpensive

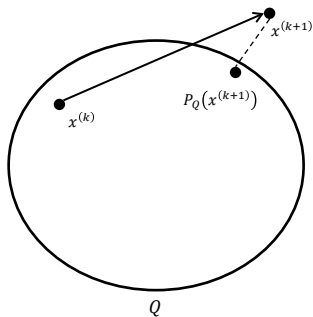
Definition

- **Prox function/ Proximal operator**

The prox function (or proximal operator) of a function $h(\cdot)$ is defined as:

$$\text{prox}_{th}(v) = \arg \min_x \left(h(x) + \frac{1}{2t} \|x - v\|_2^2 \right) \quad (11)$$

Prox function/ Proximal operator



■ Ex) Projection on set Q

$$x^{(k+1)} = P_Q \left(x^{(k)} - t \nabla g(x^{(k)}) \right)$$

$$\begin{aligned} \mathbf{prox}_h(x) &= \arg \min_u \left(I_Q(u) + \frac{1}{2} \|u - x\|_2^2 \right) \\ &= \arg \min_{u \in Q} \left(\frac{1}{2} \|u - x\|_2^2 \right) \\ &= P_Q(x) \end{aligned}$$

$$I_Q(x) = \begin{cases} 0 & \text{if } x \in Q \\ \infty & \text{if } x \notin Q \end{cases}$$

Proximal Gradient Algorithm

- Proximal Gradient Algorithm

$$x_+ \leftarrow \mathbf{prox}_{th}(x - t\nabla g(x))$$

- Uses two black boxes for the task:

$$\text{minimize } f(x) = g(x) + h(x)$$

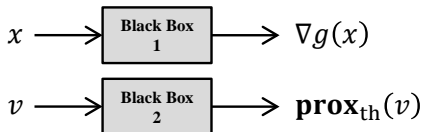


Figure : Using two black boxes for proximal gradient

Proximal Gradient Algorithm

- Form of quadratic approximation of $g(u)$ around x

$$\begin{aligned}x_+ &\leftarrow \mathbf{prox}_{th}(x - t\nabla g(x)) \\&= \arg \min_u \left(h(u) + \frac{1}{2t} \|u - x + t\nabla g(x)\|_2^2 \right) \\&= \arg \min_u \left(h(u) + \langle \nabla g(x), u - x \rangle + \frac{1}{2t} \|u - x\|_2^2 + \frac{t}{2} \|\nabla g(x)\|_2^2 \right) \\&= \arg \min_u \left(h(u) + \langle \nabla g(x), u - x \rangle + \frac{1}{2t} \|u - x\|_2^2 + g(x) \right)\end{aligned}\tag{12}$$

Convergence Analysis of Proximal Gradient

Theorem

If g is convex and g has L -lipschitz gradient, i.e.

$\|\nabla g(x) - \nabla g(y)\|_2 \leq L\|x - y\|_2$, using fixed size $t < 1/L$ on proximal gradient algorithm gives $O(1/\varepsilon)$ of convergence to the optimal point x^ , where $x^* = \arg \min_x (g(x) + h(x))$.*

■ Introducing $G_t(x)$

$$\begin{aligned}x_+ &\leftarrow x - tG_t(x) = \mathbf{prox}_{th}(x - t\nabla g(x)) \\ \Leftrightarrow G_t(x) &\triangleq \frac{1}{t}(x - \mathbf{prox}_{th}(x - t\nabla g(x)))\end{aligned}$$

Convergence Analysis of Proximal Gradient

- Preliminaries for proof

Claim

$$G_t(x) - \nabla g(x) \in \partial h(x - tG_t(x))$$

Lemma

For any point z , $f(x_+) \leq f(z) + \langle G_t(x), x - z \rangle - \frac{t}{2} \|G_t(x)\|_2^2$

Proof of Claim and Lemma: Later

Proof of Theorem

Putting $z = x$ in the Lemma gives,

$$f(x_+) \leq f(x) - \frac{t}{2} \|G_t(x)\|_2^2$$

This gives the decreasing property $f(x_+) \leq f(x)$

And putting $z = x^*$ gives,

$$\begin{aligned} f(x_+) &\leq f^* - \frac{t}{2} \|G_t(x)\|_2^2 + \langle G_t(x), x - x^* \rangle \\ \Leftrightarrow f(x_+) - f^* &\leq \frac{1}{2t} [\|x - x^*\|_2^2 - \|x - x^* - tG_t(x)\|_2^2] \\ &= \frac{1}{2t} [\|x - x^*\|_2^2 - \|x_+ - x^*\|_2^2] \end{aligned}$$

Proof of Theorem

Adding up over T iterations gives

$$\begin{aligned}\sum_{k=1}^T \left(f(x^{(k)}) - f^* \right) &\leq \frac{1}{2t} \left[\|x^{(0)} - x^*\|_2^2 - \|x^{(T)} - x^*\|_2^2 \right] \\ &\leq \frac{1}{2t} \|x^{(0)} - x^*\|_2^2\end{aligned}$$

From the fact that $f(x_+) \leq f(x)$,

$$f(x^{(k)}) - f^* \geq f(x^{(T)}) - f^* \text{ for } k = 1, \dots, T$$

This gives,

$$f(x^{(T)}) - f^* \leq \frac{1}{2tT} \|x^{(0)} - x^*\|_2^2$$

$\Rightarrow O(1/\varepsilon)$ convergence

(requires $k = O(1/\varepsilon)$ of iteration number for $f(x^{(k)}) - f^* \leq \varepsilon$)

Proof of Lemma

From the L -lipschitz continuous of gradient on g ,

$$g(y) \leq g(x) + \langle \nabla g(x), y - x \rangle + \frac{L}{2} \|y - x\|_2^2$$

So, for $y = x_+ = x - tG_t(x)$

$$g(x_+) \leq g(x) - \langle \nabla g(x), tG_t(x) \rangle + \frac{L}{2} t^2 \|G_t(x)\|_2^2 \quad (13)$$

Recall the definition of subgradient,

$$a \in \partial h(x_+) \iff h(z) \geq h(x_+) + \langle a, z - x_+ \rangle$$

By Claim, $G_t(x) - \nabla g(x) \in \partial h(x - tG_t(x))$

$$\begin{aligned} h(z) &\geq h(x_+) + \langle G_t(x) - \nabla g(x), z - x_+ \rangle \\ \Leftrightarrow h(x_+) &\leq h(z) - \langle G_t(x) - \nabla g(x), z - x_+ \rangle \end{aligned} \quad (14)$$

Proof of Lemma

Add (13), (14) and use $f(x_+) = g(x_+) + h(x_+)$

$$\begin{aligned} f(x_+) &\leq g(x) - \langle \nabla g(x), tG_t(x) \rangle + \frac{L}{2} t^2 \|G_t(x)\|_2^2 \\ &\quad + h(z) - \langle G_t(x) - \nabla g(x), z - x_+ \rangle \\ &\leq g(z) - \langle \nabla g(x), z - x \rangle - \langle \nabla g(x), tG_t(x) \rangle + \frac{L}{2} t^2 \|G_t(x)\|_2^2 \\ &\quad + h(z) - \langle G_t(x) - \nabla g(x), z - x_+ \rangle \\ &\quad (\because g(x) \leq g(z) - \langle \nabla g(x), z - x \rangle \text{ from convexity of } g) \\ &= f(z) + \frac{L}{2} t^2 \|G_t(x)\|_2^2 - \langle G_t(x), z - x + tG_t(x) \rangle \\ &\leq f(z) + \langle G_t(x), x - z \rangle + \frac{t}{2} \|G_t(x)\|_2^2 - t \|G_t(x)\|_2^2 \quad (\because t < 1/L) \\ &= f(z) + \langle G_t(x), x - z \rangle - \frac{t}{2} \|G_t(x)\|_2^2 \end{aligned}$$

Proof of Claim

Claim. $G_t(x) - \nabla g(x) \in \partial h(x_+)$

- Property of proximal operator

$$u = \mathbf{prox}_{th}(x) \Leftrightarrow \frac{1}{t}(x - u) \in \partial h(u) \quad (15)$$

(Proof: Use that $\mathbf{prox}_{th}(x)$ is a minimizer)

Now put $u = x_+ = \mathbf{prox}_{th}(x - \nabla g(x))$ in equation (15),

$$\begin{aligned} \Rightarrow \frac{1}{t}(x - \nabla g(x) - x_+) &\in \partial h(x_+) \\ \Rightarrow \frac{1}{t}(x - \nabla g(x) - x + tG_t(x)) &\in \partial h(x_+) \\ \Rightarrow G_t(x) - \nabla g(x) &\in \partial h(x_+) \end{aligned}$$

Next lecture

- Review of oracle complexity of algorithm
- Nesterov's accelerated gradient descent

Thank you!