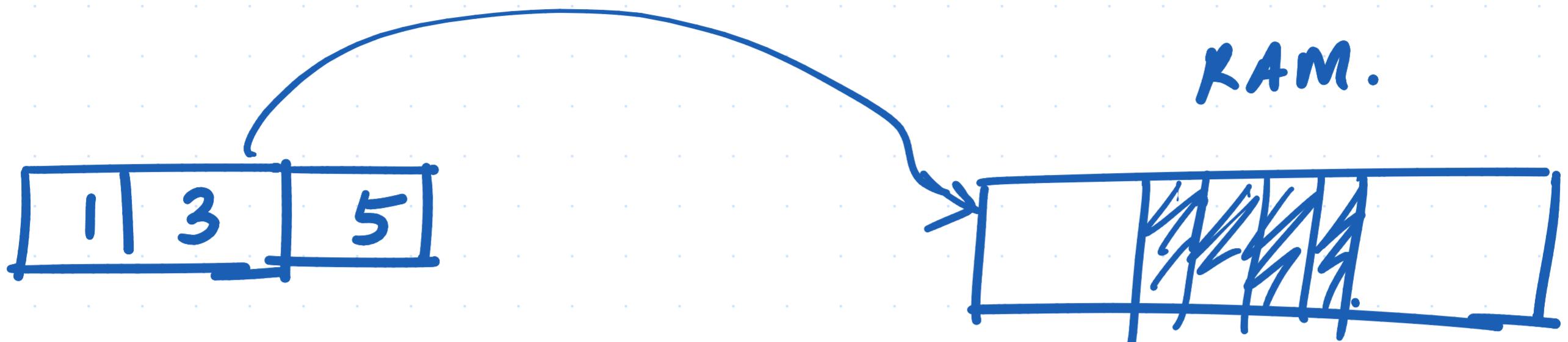


16/08/25

Arrays / Data Structure.



Measured in Bytes

, 8 GB of RAM.

Byte \rightarrow 8 bits.

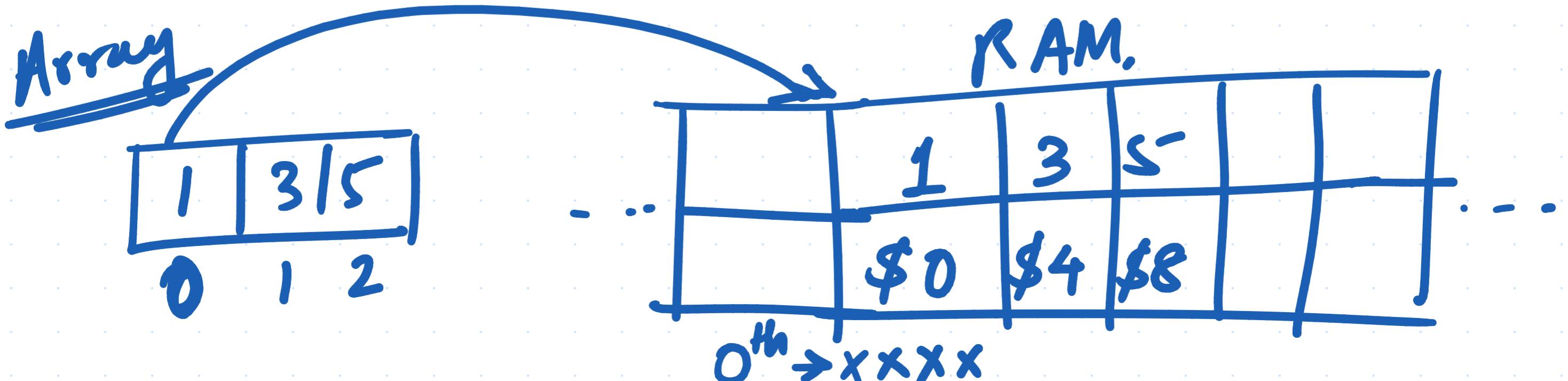
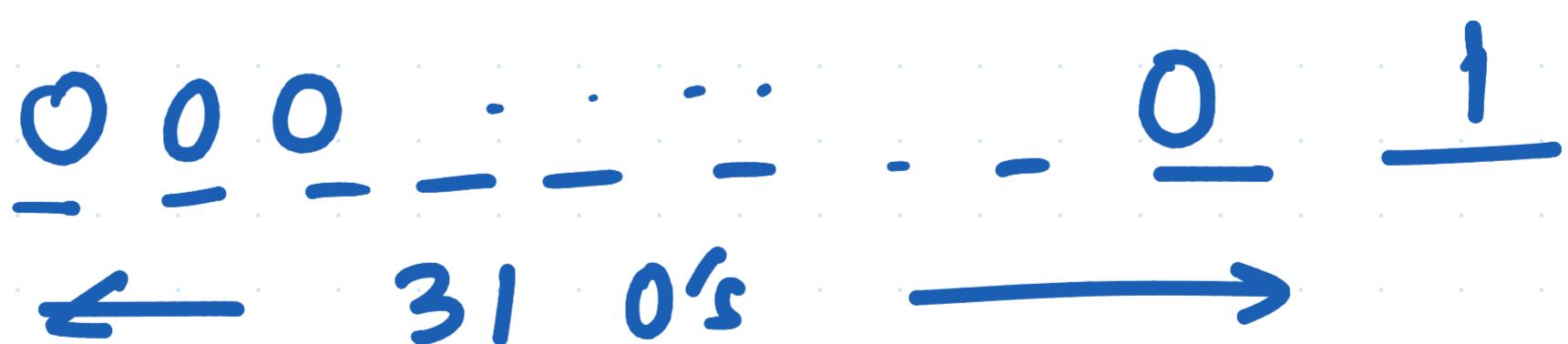
$10^9 \rightarrow$ Bytes

Bit \rightarrow a portion that can store 0/1.

Integer \rightarrow 4 byte.

? Bits? \rightarrow 32 Bits.

1. \rightarrow integer.



$(xxxx + 4) \rightarrow 3$ value.

$(xxxx + 8) \rightarrow 5$ value.

1 byte \rightarrow 1 character

ASCII value representation.

access.



Read $\rightarrow O(1)$. constant time read

Write $\rightarrow O(1)$ constant time write.

Random Access Memory \rightarrow we can access in constant time, any part of the memory given we have the correct location / address for it.

(garbage collector).

C, C++

java:

value	address
1	\$0
3	\$4
5	\$8
7	\$12
6	\$16

C, C++ \rightarrow vectors \rightarrow dynamic. accessed by other / used by others.

int arr[10];

Python → dynamic array → List

value	2	5	3	2				
addr	\$0	\$4	\$8	\$12	\$16	\$20	\$24	

← 8 →

inserting in the middle of the array,
we need to shift elements, so it will
take $O(n)$ time.

Deletion of first & last $\Rightarrow \underline{\underline{O(1)}}$

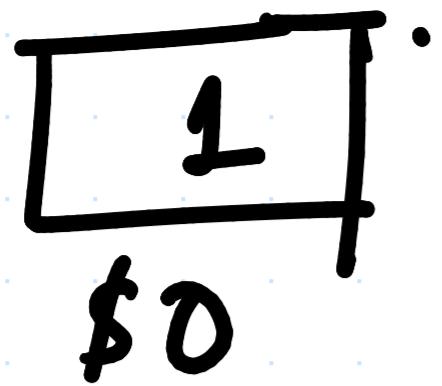
$O(n) \rightarrow$ remove & shift elements of the array

Dynamic array, how is it maintained in
Python or any kind of Prog Lang.

→ No size is mentioned.

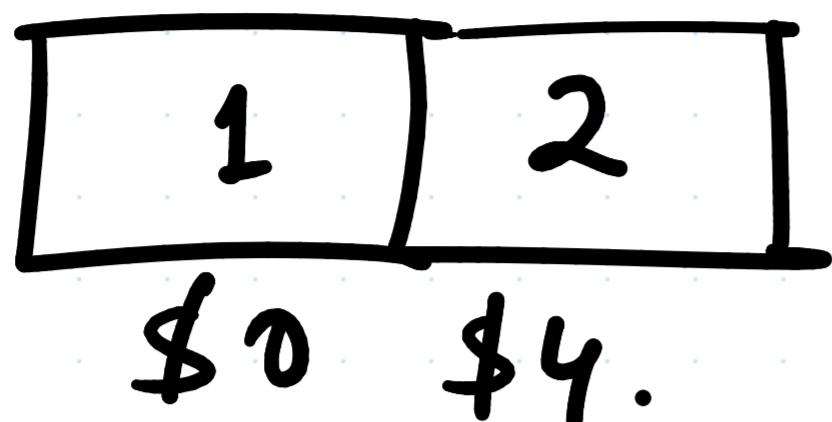
$a = []$; List.

1



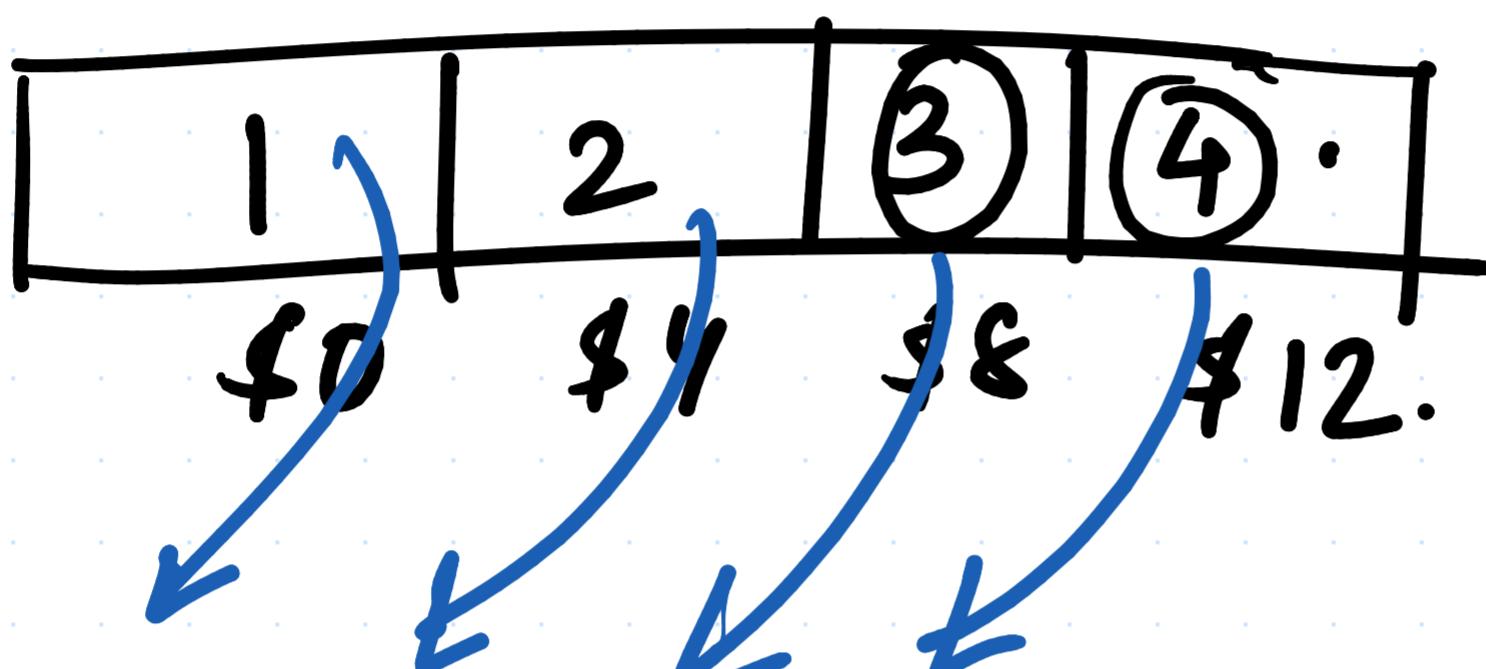
$$2 > 1$$

2

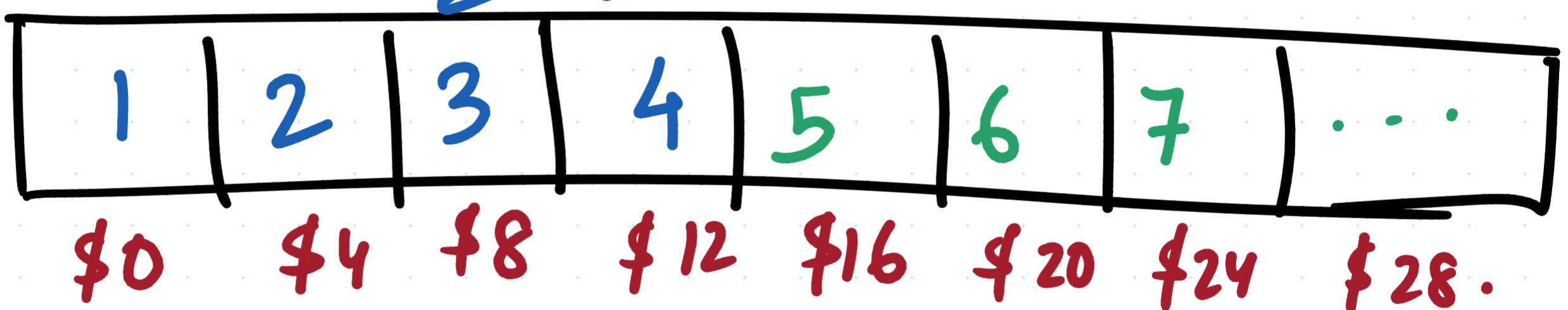


$$4 > 2 + 1$$

4



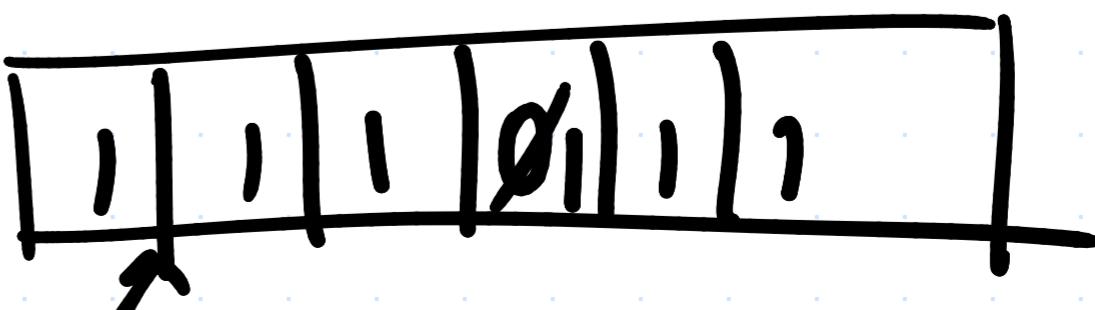
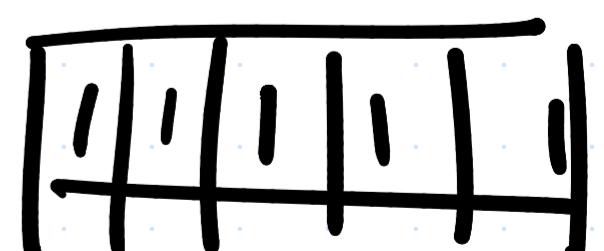
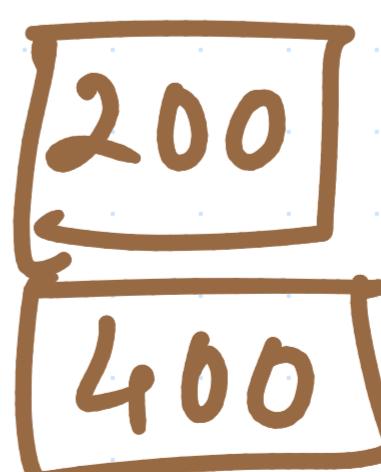
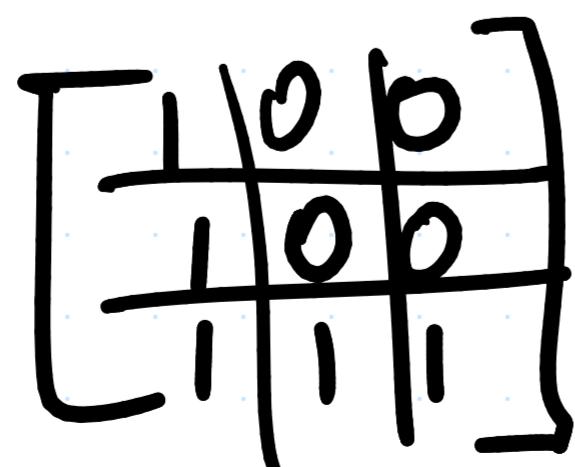
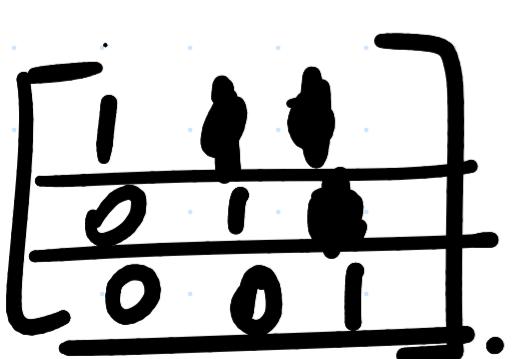
8



Tradeoff b/w time & space.

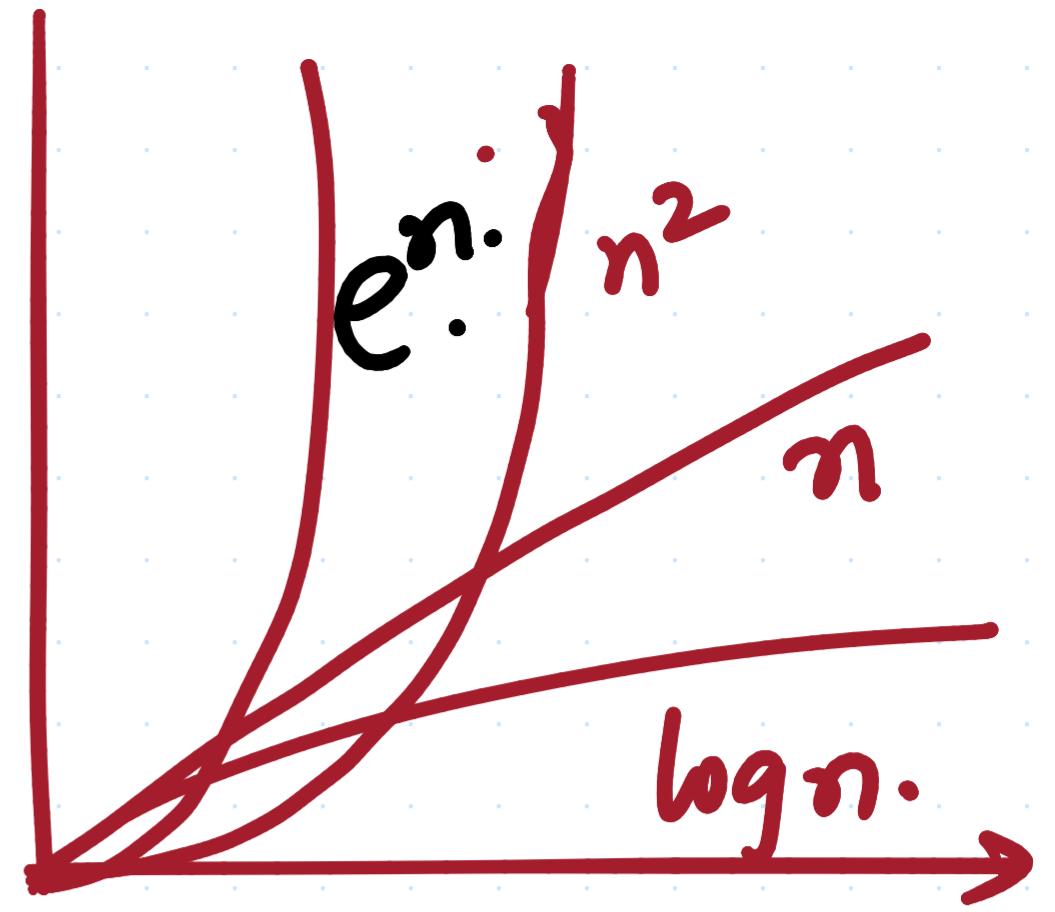


Optimal $\rightarrow \underline{2x}$.



Mostly $O(n)$ & $O(\log n)$

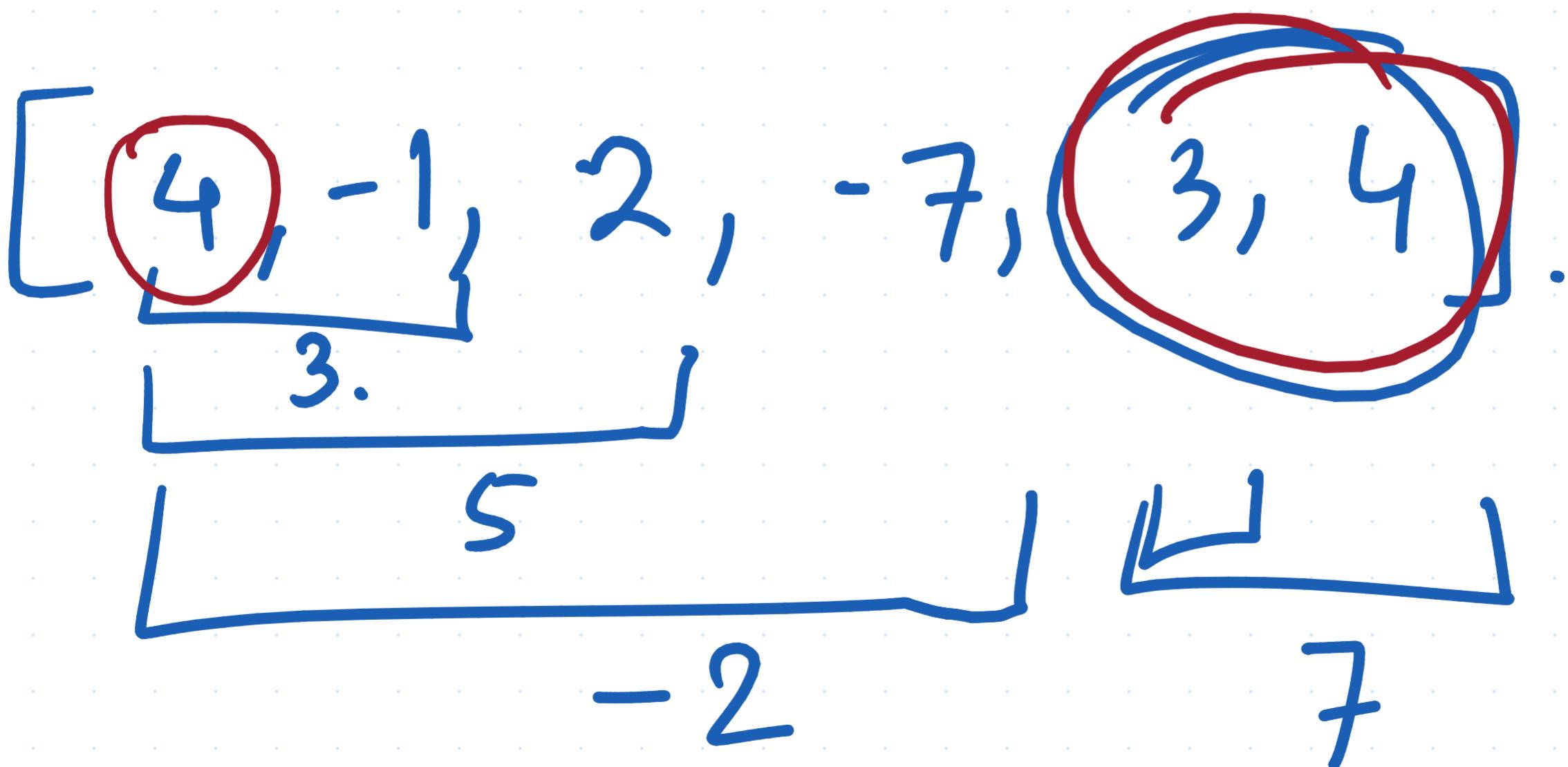
↓
type of Algo.



Brute force $\rightarrow O(n^2)$.

Kadane's Algorithm (2ptr method)

Find a non-empty subarray with the largest sum.

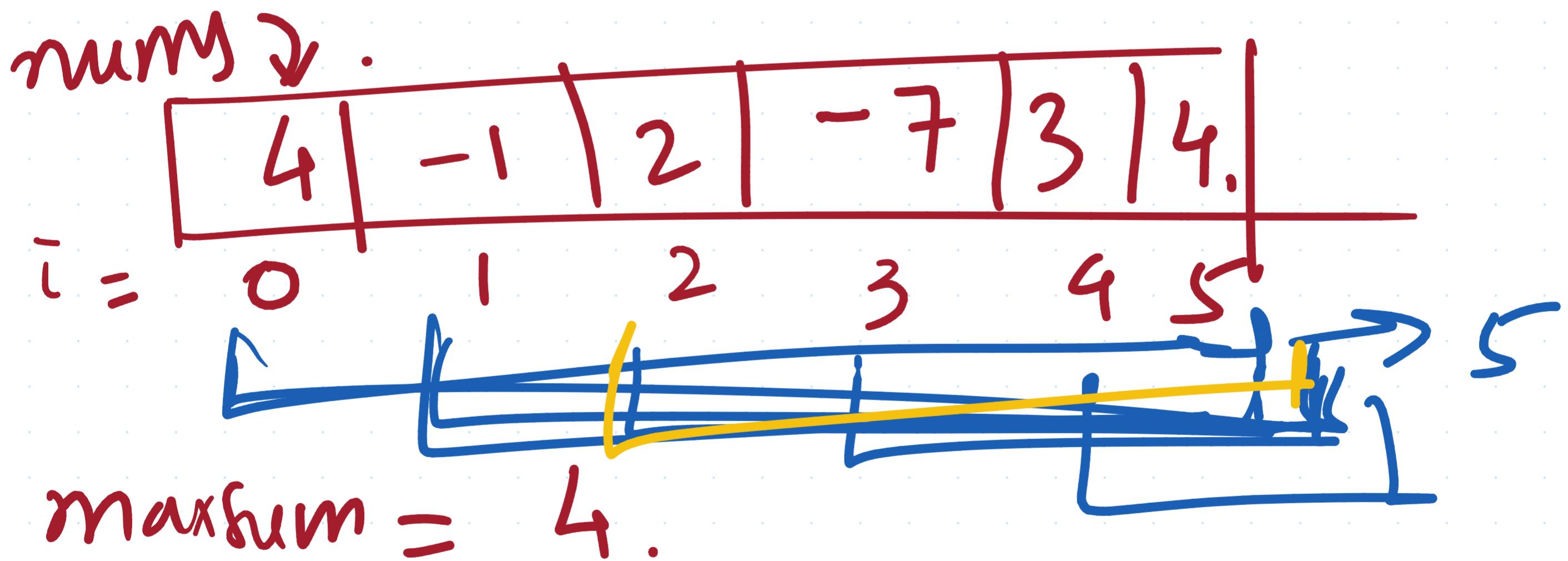


$$O/P \rightarrow 7.$$

```

def bruteForce (nums)
    maxSum = nums [0]
    for i in range (len(nums)):
        curSum = 0
        for j in range (i, len(nums)):
            curSum += nums [j]
            maxSum = max(maxSum, curSum)
    return maxSum.

```



$$\text{curSum} = \cancel{0} \cancel{4} \cancel{3} \cancel{-} \cancel{2} \cancel{5}.$$

$$\text{maxSum} = \underline{\underline{4}}, \underline{\underline{4}}, \underline{\underline{4}}, \underline{\underline{5}}. \underline{\underline{5}} \cdot \underline{\underline{5}}$$

unsum = $\emptyset - X X - b - \beta$

maxm = \$ 8.5 8 7 5



def Kadanes(nums):

maximum = nums[0]

cursum = 0

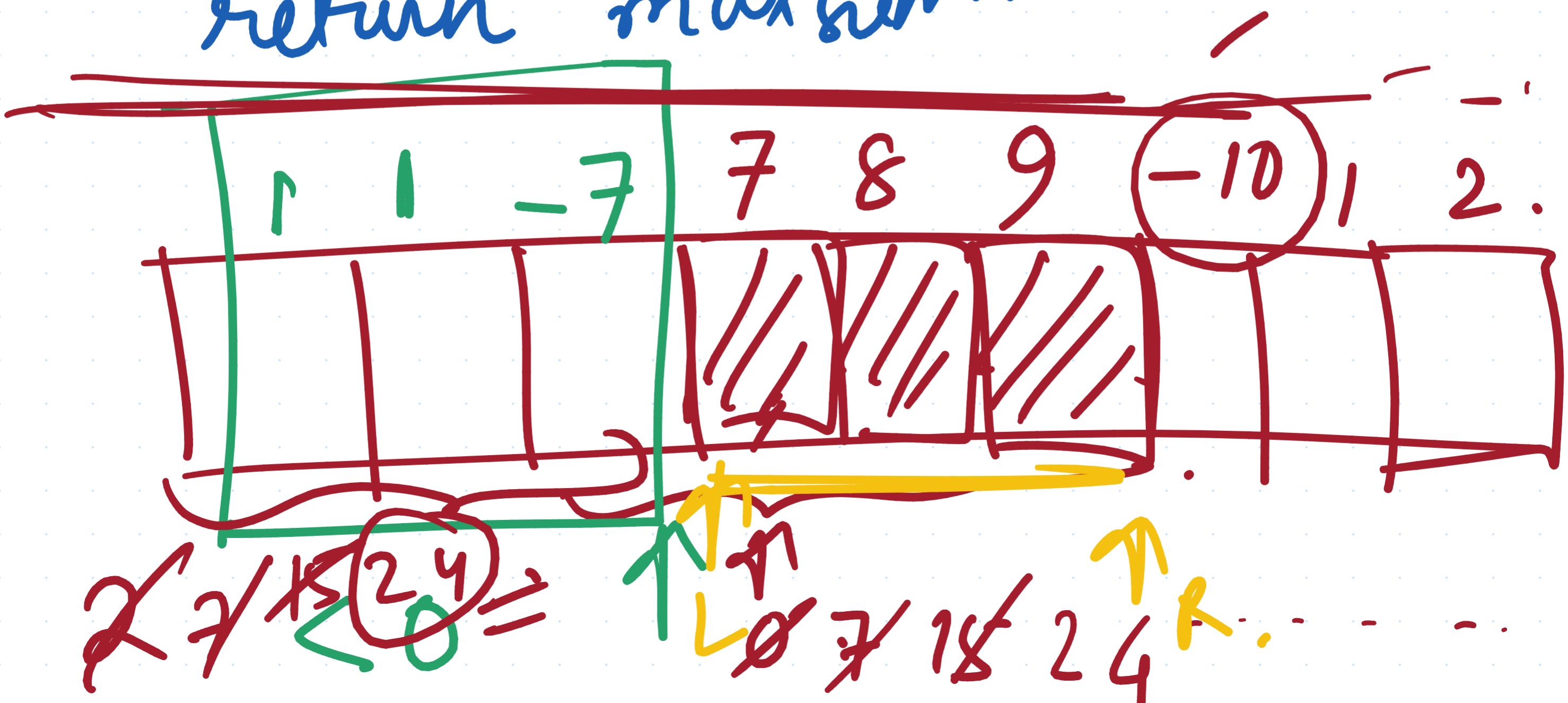
for n in nums: $O(n)$

 · cursum = $\max(\text{cursum}, 0)$

 · cursum += n

 · maximum = $\max(\text{maximum}, \text{cursum})$

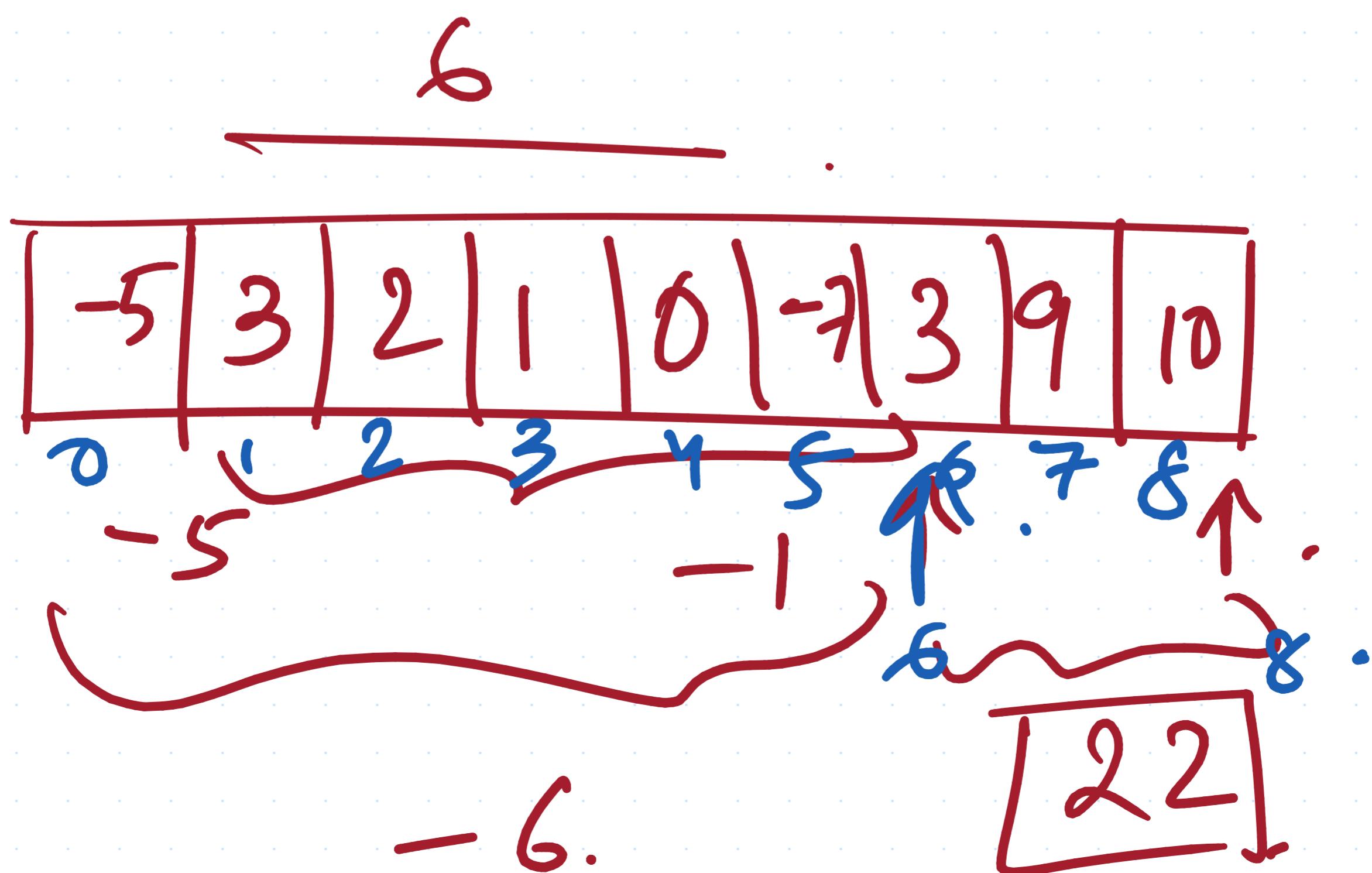
return maximum.



sliding window Version of

Kadane's Algorithm.

Q : Return the left & right index of the max-subarray sum, assuming there is exactly one result.



def slidingWindow(nums):

maxSum = nums[0]

cursum = 0

maxL, maxR = 0, 0

L = 0

for R in range(len(nums)):

, if cursum < 0:

cursum = 0 ; L = R

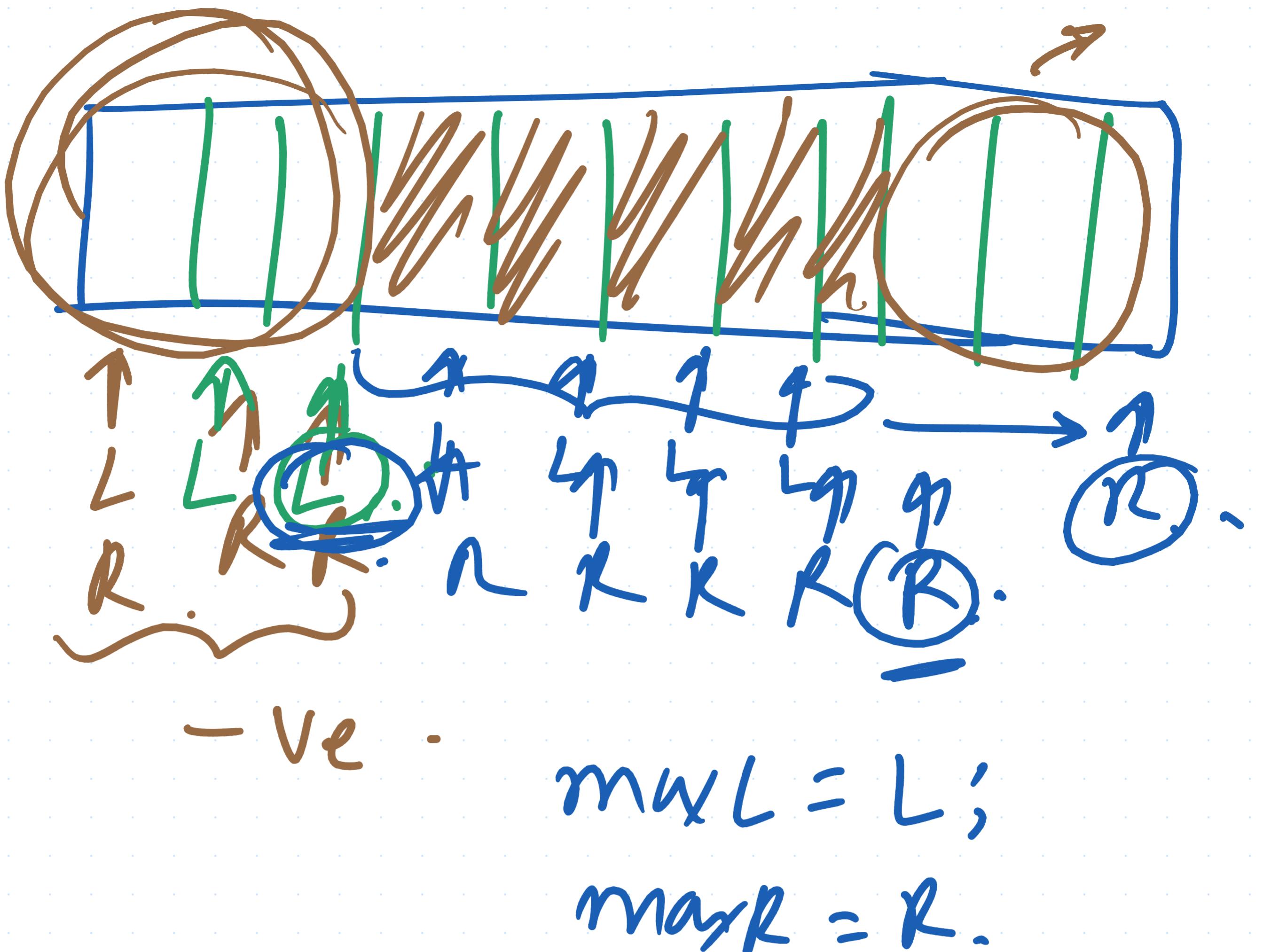
, cursum += nums[R].

, if cursum > maxSum :

, maxSum = cursum

, maxL, maxR = L, R.

return [maxL, maxR].



Sliding Window

Given an array, return true
 if there are two elements
 within a window of size k.
 that are equal.

1	2	3	2	3	3	4	5

window of size. $k = 2$.

1	2	3	2	3	2

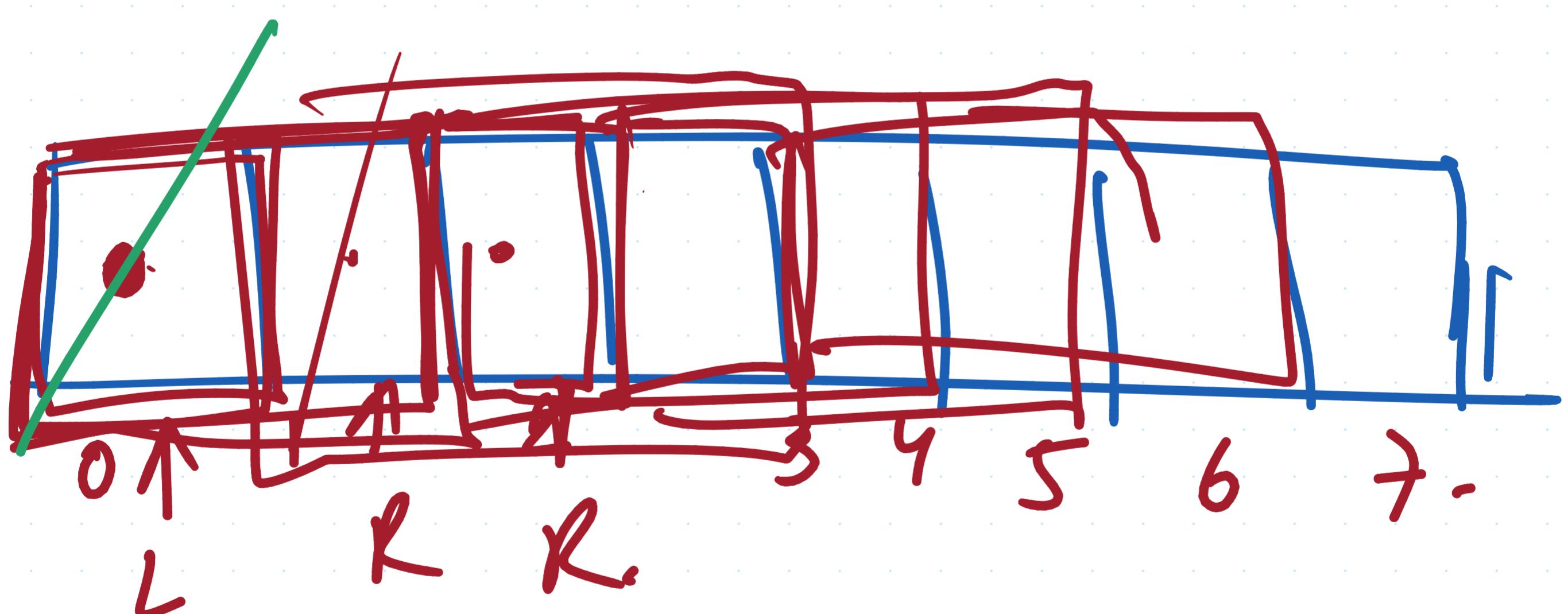
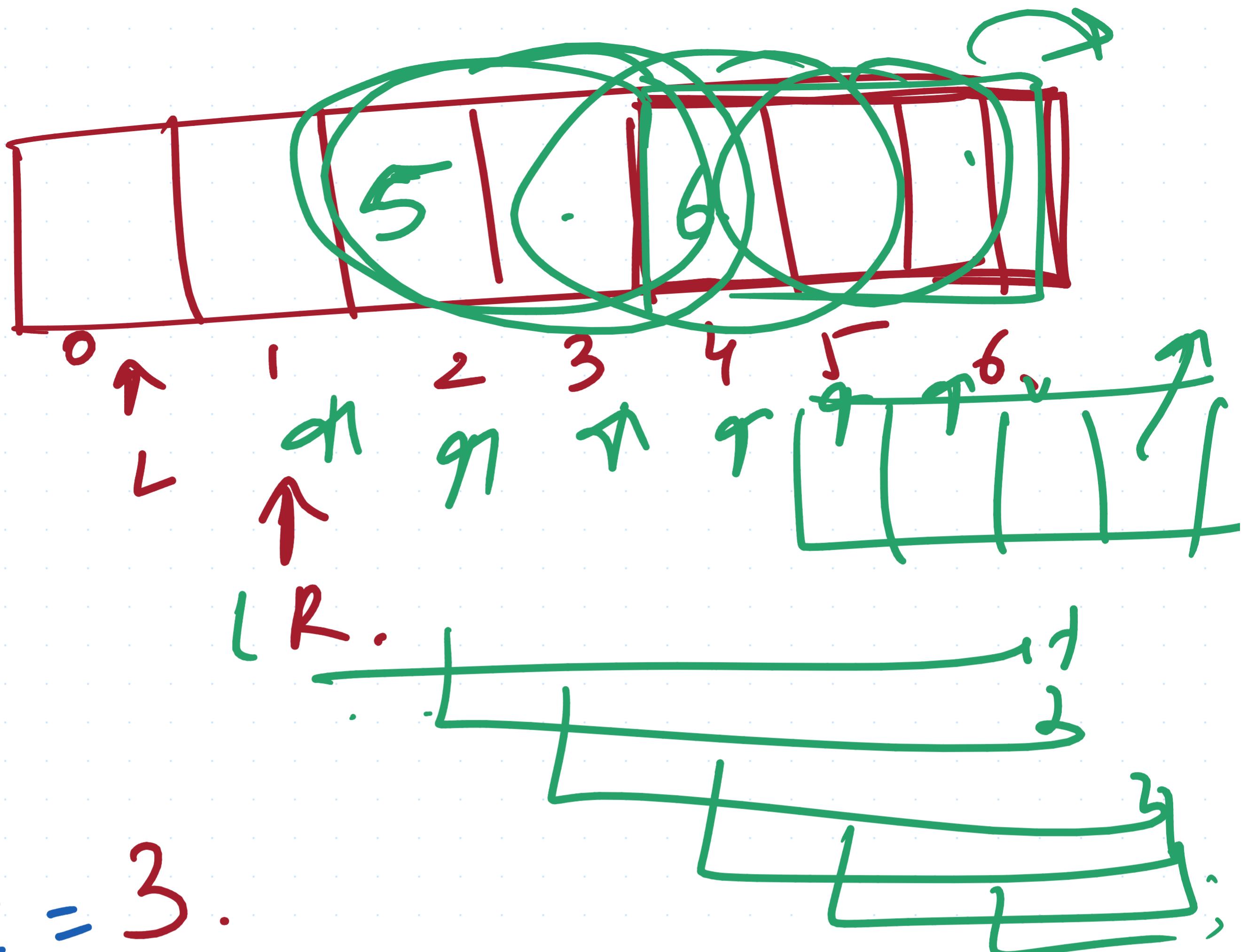
$k = 3$.

2 true.

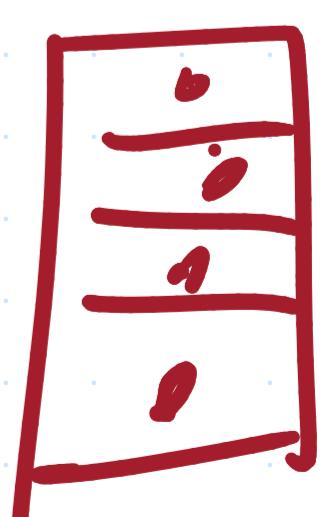
```

def closestBF( nums, k ):
    for L in range( len(nums) ):
        for R in range( L+1, min( len(nums), L+k ) ):
            if nums[L] == nums[R]:
                return True
    return False

```



Sliding window.



$O(1) \rightarrow$ hash set.

$O(n)$ complexity:

def closeDuplicates(nums, k):

window = set()

$l = 0$

for R in range($\text{len}(nums)$):

If $R - l + 1 > k$:

$\underline{\text{window.remove}(nums[l])}$.

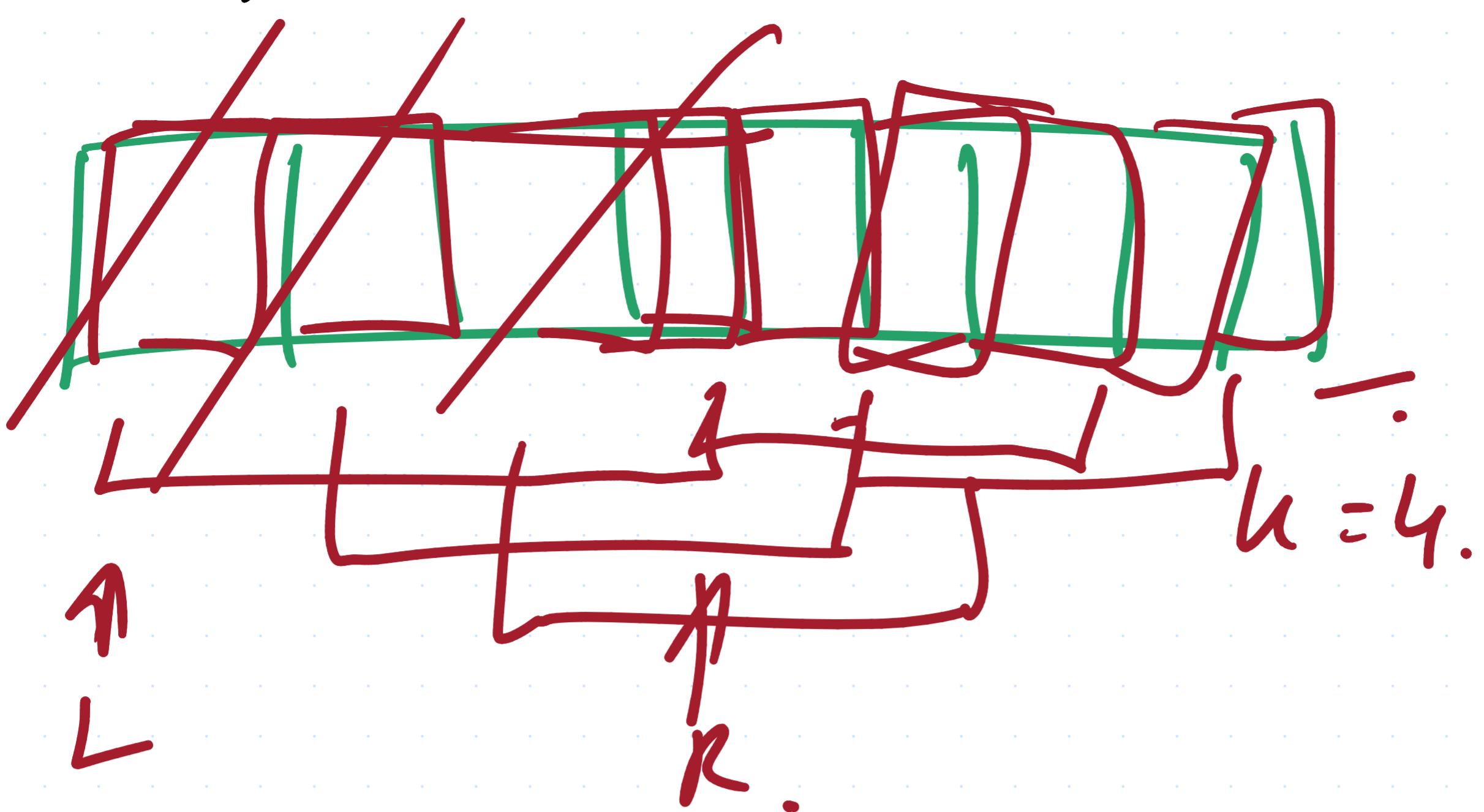
$\underline{l += 1}$

If $nums[k]$ in $\underline{\text{window}}$:

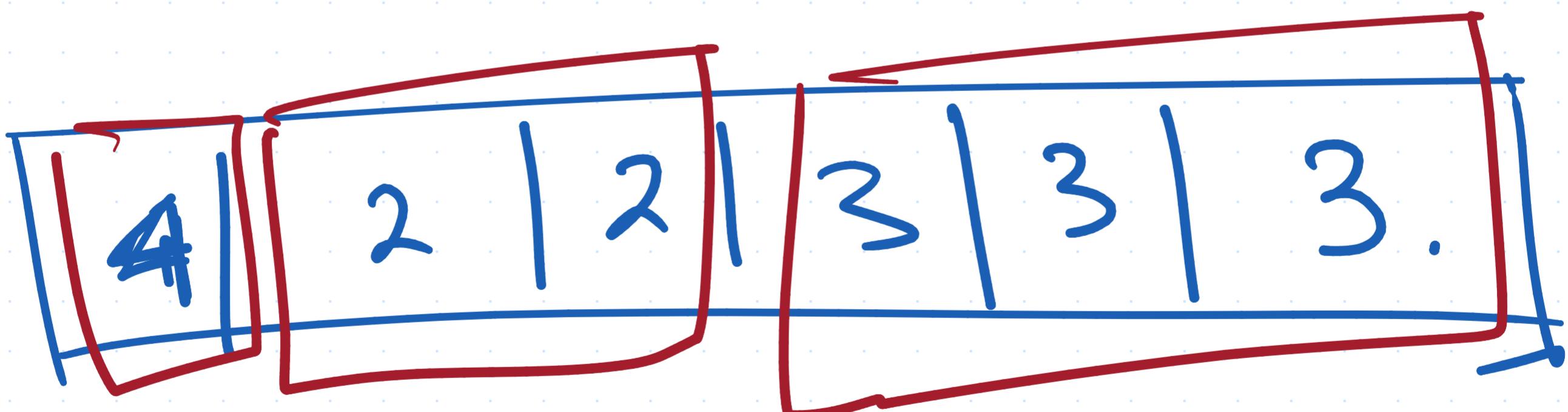
return True

$\underline{\text{window.add}(nums[R])}$

return False.



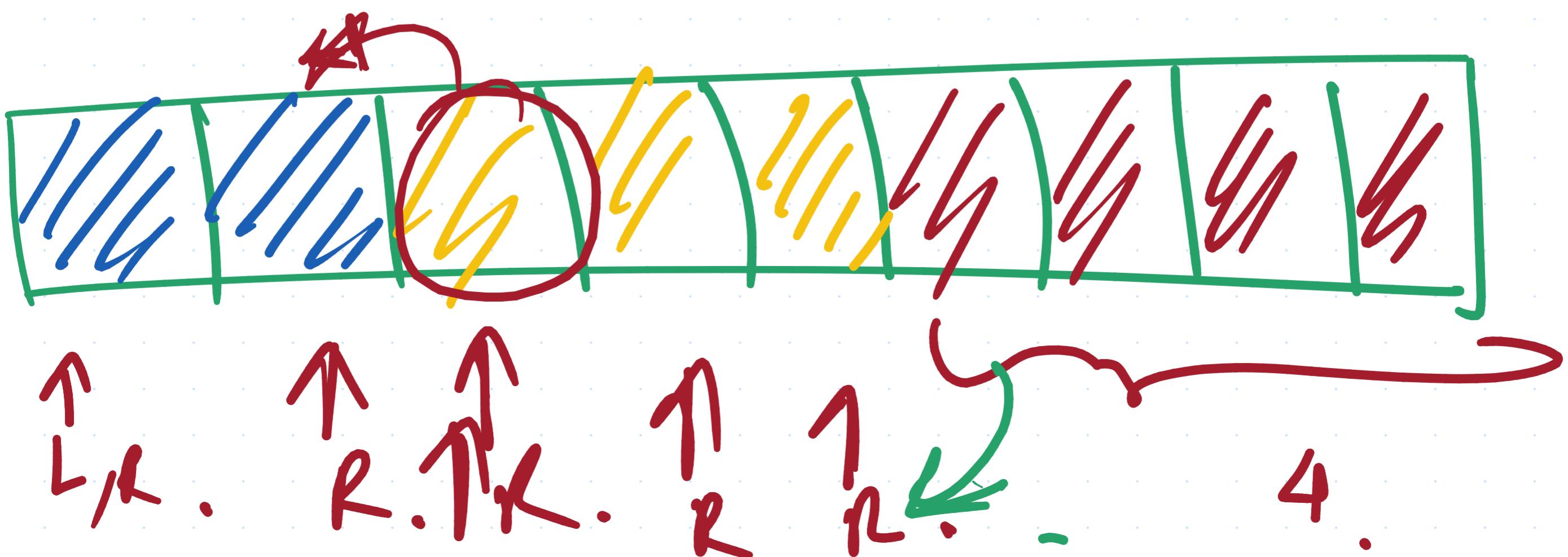
Sliding Window with variable length.



- Find the length of the longest subarray with the same value in each position.

$O(n)$

```
def longestSubarray (nums):
    length = 0; L = 0
    for R in range (len(nums)):
        if nums[L] != nums[R]:
            L = R
        length = max(length, R-L+1)
    return length
```



L

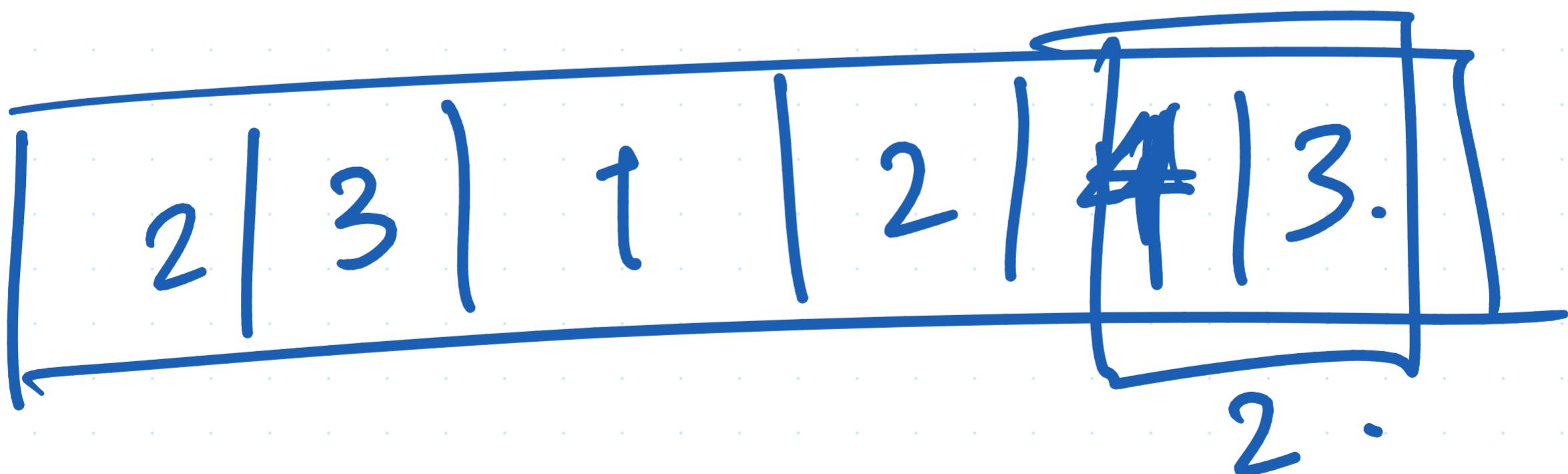
L, R. R.

$$= \text{R} - \text{L} + 1$$

find the min length subarray

where the sum is greater
than or equal to the target.

Assume all values are +ve.



Target = 6.

```
def shortestSubarray(nums, target):
```

```
L, total = 0, 0
```

```
length = float("inf") # infinity
```

```
for R in range(len(nums)):
```

```
    total += nums[R]
```

```
    while total >= target:
```

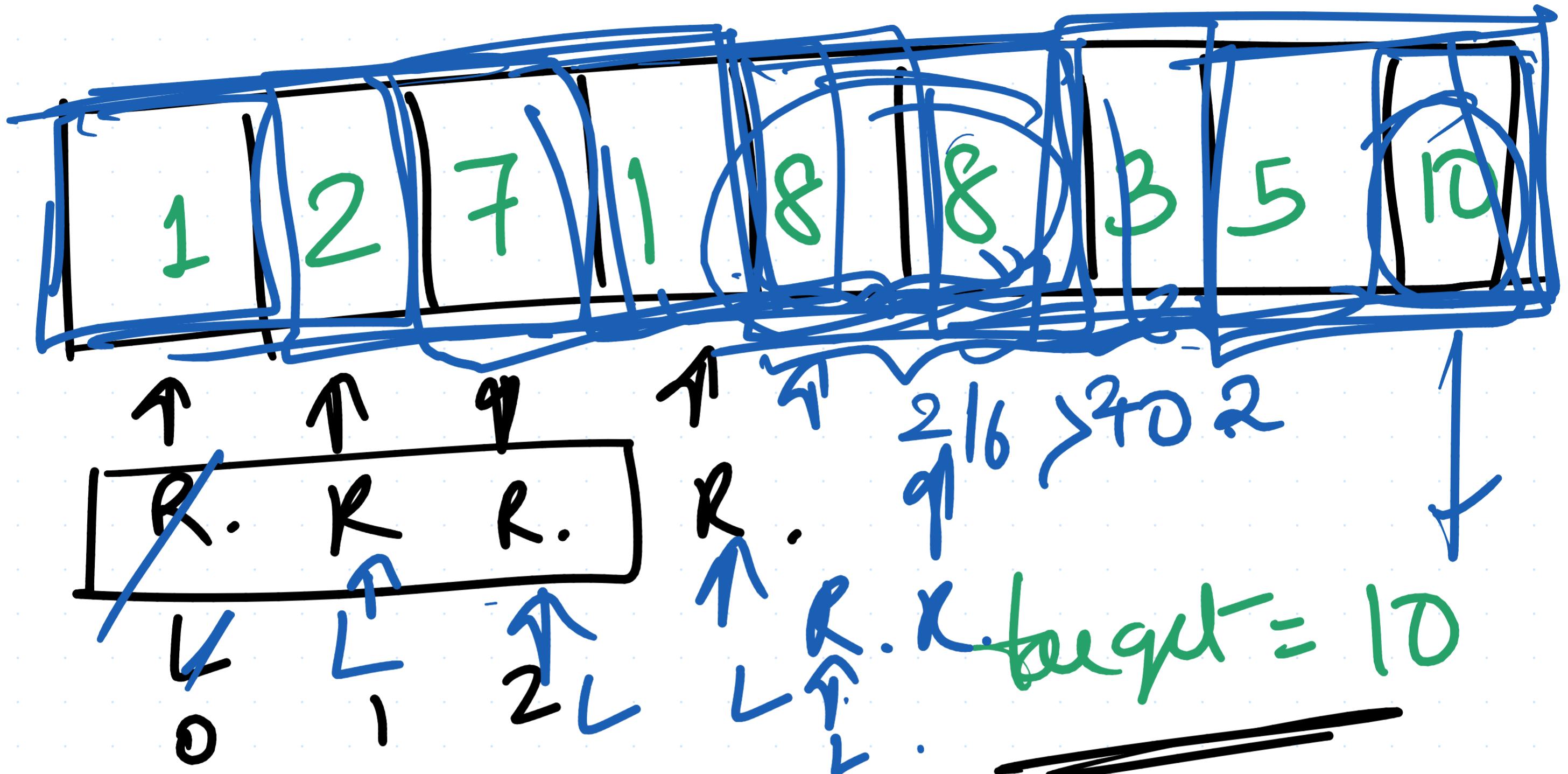
```
        length = min(R - L + 1, length)
```

```
        total -= nums[L]
```

```
        L += 1
```

```
return 0 if length == float("inf")
```

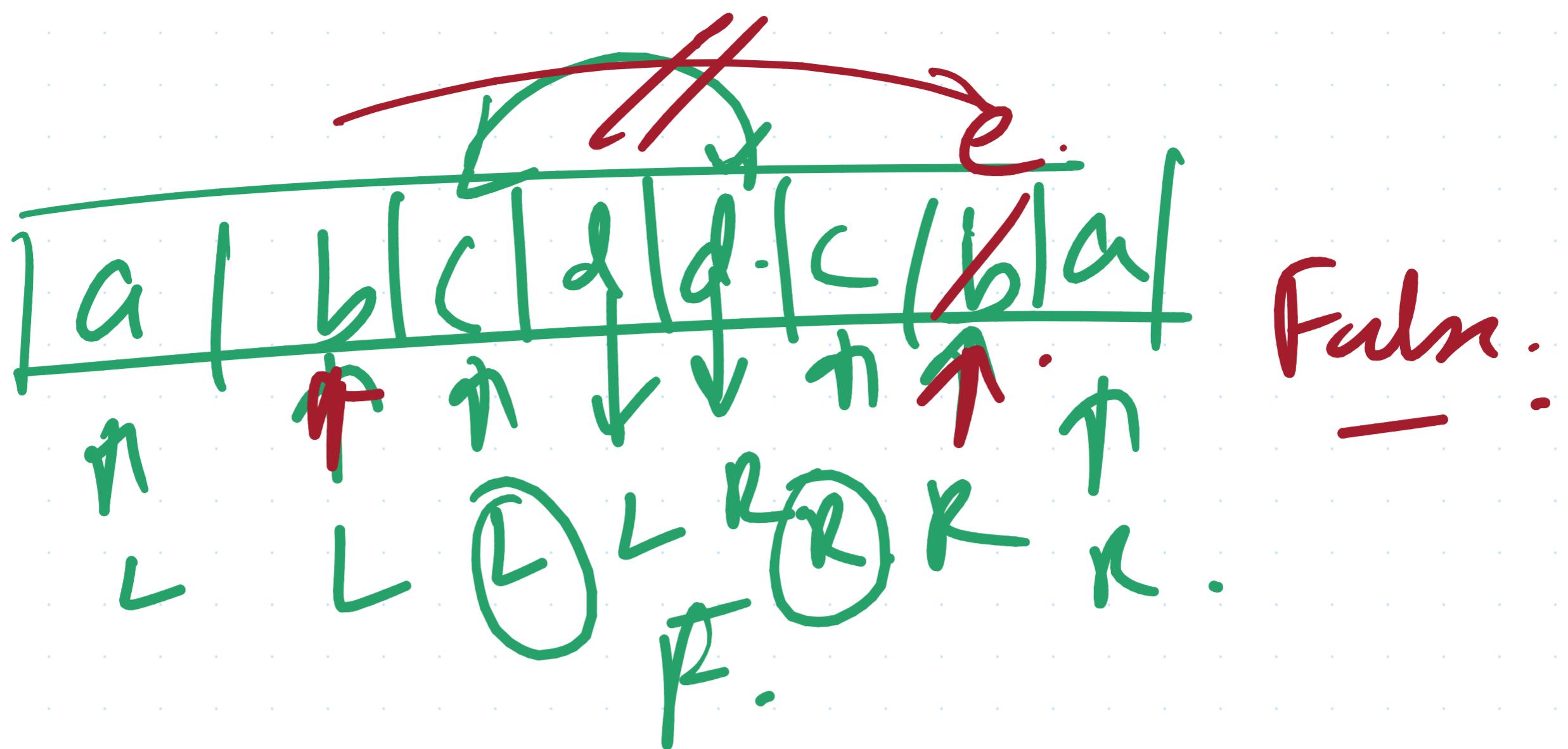
else length.



$$\text{length} = \sum_{i=1}^{10} 11; \quad \underline{\underline{R-L+1}}$$

0/0 → 1.

O(n)



check if an array is palindrome.

def IsPal(word):

L, R = 0, len(word) - 1

while L < R:

} if word[L] != word[R]:

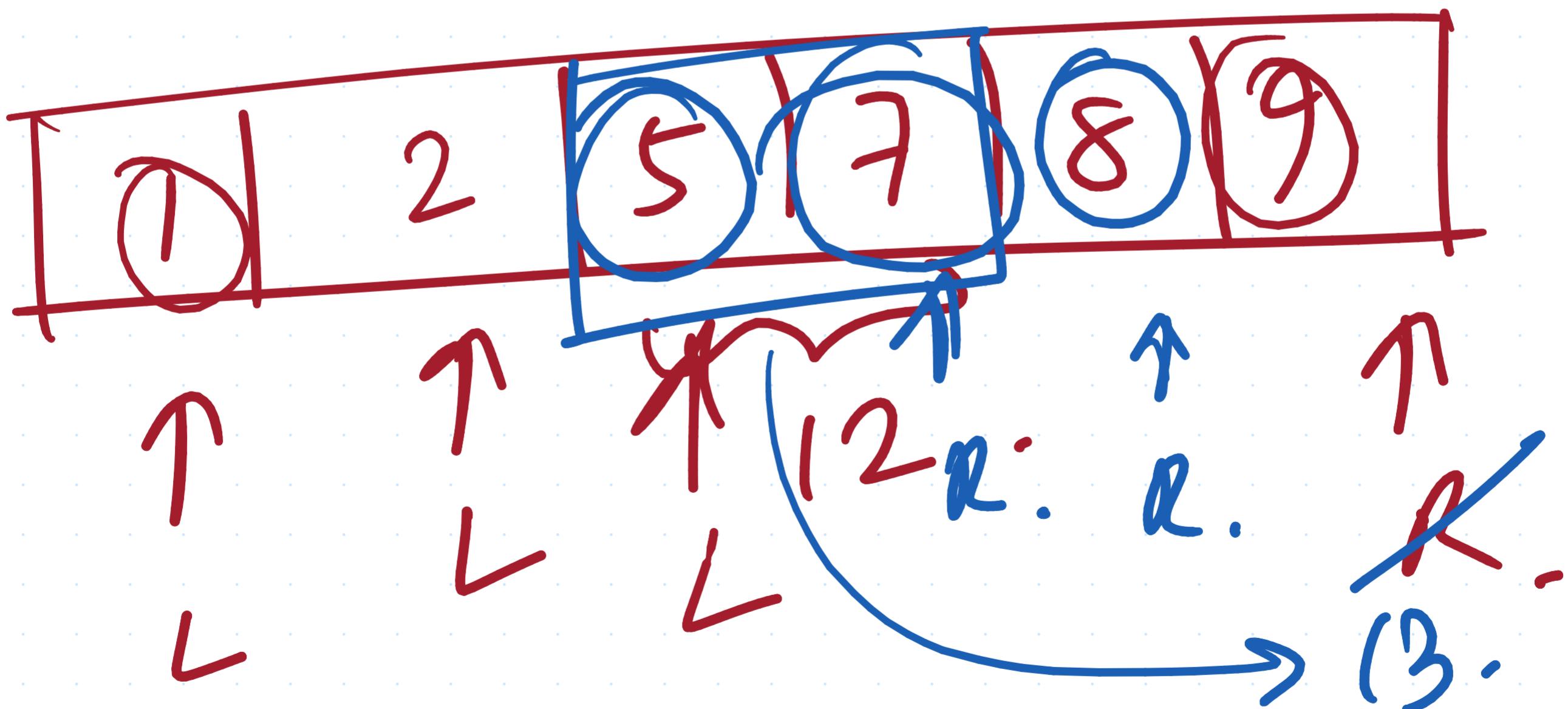
 return False.

} L += 1 ; R -= 1

return True.

Q) Given a sorted np array,
return the two indices of the
two elements which sum up to
the target value. Assume

there is exactly 1 sol.



$$\text{sum} = 10 + 14 \not> 12.$$

(L, R)

fond, true.

def targetsum (nums, target):

L, R = 0, len(nums) - 1

while L < R:

if nums[L] + nums[R] > target:

R -= 1

elif nums[L] + nums[R] <
target:

L += 1

else

return [L, R].