

1. 证:

下面证明  $f(n) = \sum_{i=1}^n i^2 = \frac{1}{6}n(n+1)(2n+1)$ ,

当  $n = 1$  时,  $f(1) = 1$ ,

假设  $n = k$  时,  $f(n) = \frac{1}{6}n(n+1)(2n+1)$  成立,

当  $n = k + 1$  时,

$$\begin{aligned} f(n+1) &= f(n) + (n+1)^2 \\ &= \frac{1}{6}n(n+1)(2n+1) + (n+1)^2 \\ &= \frac{1}{6}(n+1)(n+1+1)(2(n+1)+1) \end{aligned}$$

从而  $f(n) = \frac{1}{6}n(n+1)(2n+1)$

故  $f(n) = \Theta(n^3)$

2. 解:

(1)  $T(n) = O(n)$

即证  $\exists c, n_0 > 0, s.t. n > n_0, T(n) \leq cn$

假设  $n < k$  时均成立,

当  $n = k$  时,

$$T(k) = T\left(\frac{k}{2}\right) + k \leq c\frac{k}{2} + k = \left(\frac{c}{2} + 1\right)k \leq ck, \quad c \geq 2,$$

即为  $T(n) = O(n)$

$T(n) = \Omega(n)$

即证  $\exists c, n_0 > 0, s.t. n > n_0, T(n) \geq cn$

假设  $n < k$  时均成立,

当  $n = k$  时,

$$T(k) = T\left(\frac{k}{2}\right) + k \geq c\frac{k}{2} + k = \left(\frac{c}{2} + 1\right)k \geq ck, \quad c \leq 2,$$

即为  $T(n) = \Omega(n)$

(2)  $T(n) = O(n \log^2 n)$

即证  $\exists c, n_0 > 0, s.t. n > n_0, T(n) \leq cn \log^2 n$

假设  $n < k$  时均成立,

当  $n = k$  时,

$$\begin{aligned} T(k) &= 2T\left(\frac{k}{2}\right) + k \log k \leq c \frac{k}{2} \log^2 \frac{k}{2} + k \log k \\ &= ck \log^2 k + (1 - 2c \log 2) n \log n + cn \log^2 2 \\ &\leq ck \log^2 k, \quad c = 1, n_0 \geq 10, \end{aligned}$$

即为  $T(n) = O(n \log^2 n)$

$T(n) = \Omega(n \log^2 n)$

即证  $\exists c, n_0 > 0, s.t. n > n_0, T(n) \geq cn \log^2 n$

假设  $n < k$  时均成立,

当  $n = k$  时,

$$\begin{aligned} T(k) &= 2T\left(\frac{k}{2}\right) + k \log^2 k \geq c \frac{k}{2} \log^2 \frac{k}{2} + k \log k \\ &= ck \log^2 k + (1 - 2c \log 2) n \log n + cn \log^2 2 \\ &\geq ck \log^2 k, \quad c = \frac{1}{2}, n_0 \geq 5, \end{aligned}$$

即为  $T(n) = \Omega(n \log^2 n)$

3. 证:

对于  $\forall x > -1, x > \ln(x+1)$ ,

从而

$$\begin{aligned} 1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n} &\geq \ln\left(\frac{2}{1}\right) + \ln\left(\frac{3}{2}\right) + \dots + \ln\left(\frac{n+1}{n}\right) \\ &= \ln\left(\frac{2}{1} \cdot \frac{3}{2} \cdot \dots \cdot \frac{n+1}{n}\right) \\ &= \ln(n+1) \end{aligned}$$

4. 解:

$T(n) = O(n^2)$

即证  $\exists c, n_0 > 0, s.t. n > n_0, T(n) \leq cn^2$

假设  $n < k$  时均成立,

当  $n = k$  时,

$$T(k) = 2T\left(\frac{k}{2}\right) + k^2 \leq c\left(\frac{k}{2}\right)^2 + k^2 = \left(\frac{c}{4} + 1\right)k^2 \leq ck^2, \quad c \geq \frac{4}{3},$$

即为  $T(n) = O(n^2)$

$$T(n) = \Omega(n^2)$$

即证  $\exists c, n_0 > 0, s.t. n > n_0, T(n) \geq cn^2$

假设  $n < k$  时均成立,

当  $n = k$  时,

$$T(k) = 2T\left(\frac{k}{2}\right) + k^2 \geq c\left(\frac{k}{2}\right)^2 + k^2 = \left(\frac{c}{4} + 1\right)k^2 \geq ck^2, \quad c \leq \frac{4}{3},$$

即为  $T(n) = \Omega(n^2)$

5. 解:

$$T(n) = \Omega(n \log n)$$

即证  $\exists m, n_0 > 0, s.t. n > n_0, T(n) \geq mn \log n$

假设  $n < k$  时均成立,

当  $n = k$  时,

$$\begin{aligned} T(k) &= T\left(\frac{k}{3}\right) + T\left(\frac{2k}{3}\right) + ck \\ &\geq mc\frac{k}{3} \log \frac{k}{3} + mc\frac{2k}{3} \log \frac{2k}{3} + ck \\ &= cmk \log k + \left(\frac{2}{3} \log 2 - \log 3\right) cmk + ck \end{aligned}$$

当  $c \leq 1$ , 取  $m < \frac{1}{3 \log 3 - 2 \log 2}$ , 即有  $T(k) \geq k \log k$ ,

当  $c > 1$ , 取  $m > \frac{1}{3 \log 3 - 2 \log 2}, n_0 = 10$ , 即有  $T(k) \geq k \log k$ ,

即为  $T(n) = \Omega(n \log n)$