

# Programming Language Concepts

Programming Language Theory

# Topics

- Memory Management
  - Static Management
  - Dynamic Management w/ Stack
  - **Dynamic Management w/ Heap**
  - Scope Rule Implementation

# Heap

- Why we need ***Heap*** for memory management?
  - We already have stack, and it seems natural to manage memory for procedures.
- Some languages have statements which allow ***explicit memory allocation***.

# Explicit Memory Allocation

- With explicit memory allocation, there is no guarantee of LIFO.
- In the example, pointer variable **p** is the first one allocated, and also the first one deallocated.
- If we use the stack, we can't deallocate **p** before **q**, since **q** is at the top.

```
int *p, *q;  
p = malloc(sizeof(int));  
q = malloc(sizeof(int));  
*p = 1;  
*q = 2;  
free(p);  
free(q);
```

**Allocation**

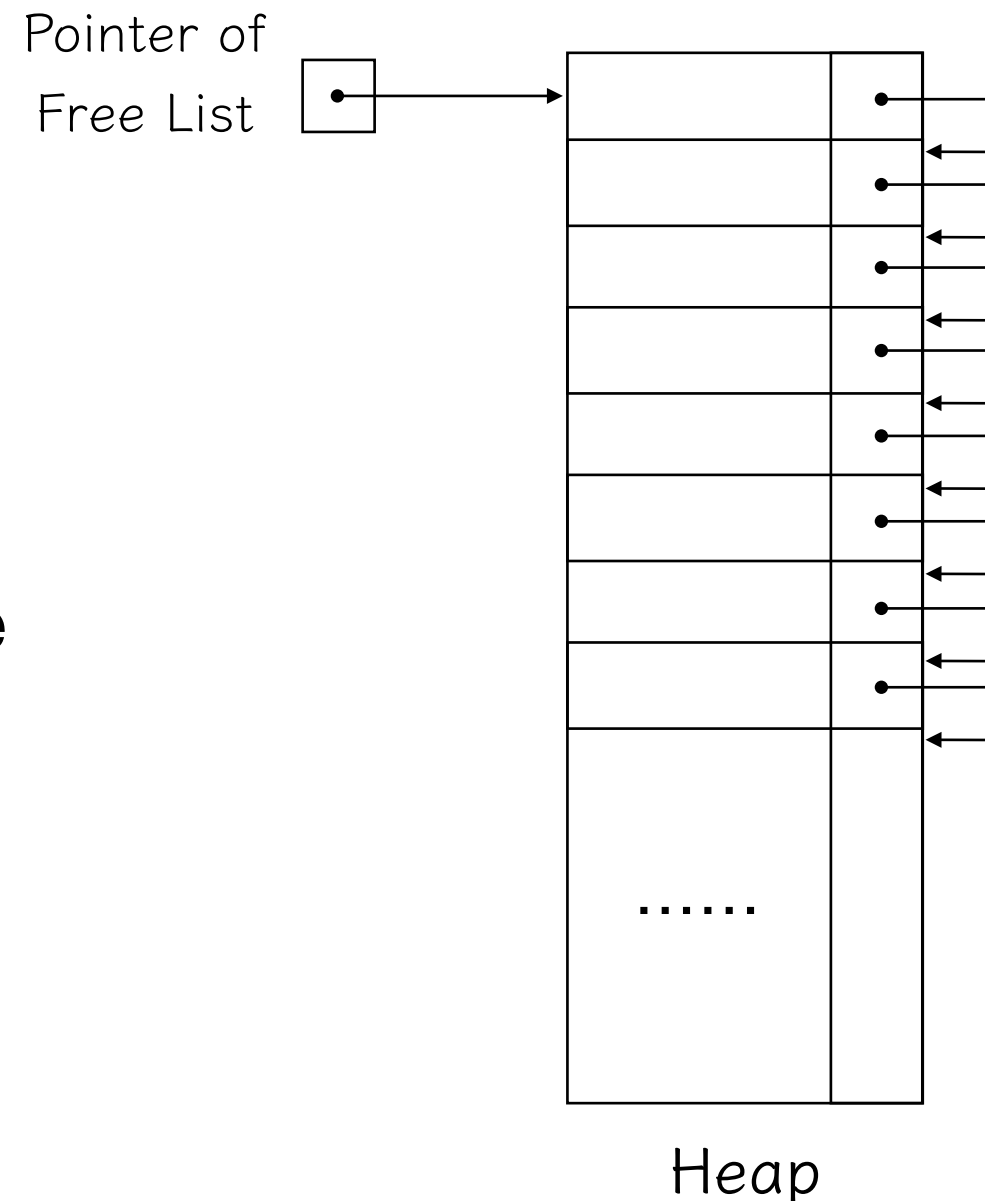
**Deallocation**

# Heap Management

- Heap management methods fall into two main categories,
- based on whether the memory blocks are considered,
  - *Fixed Length Blocks*
  - *Variable Length Blocks*

# Fixed Length Blocks

- Divide the heap to multiple fixed length blocks.
- Using a free list to maintain the list of free blocks.
- For each request, the first free block will be allocated.
- The pointer of the free list points to the first block of the list.



# Fixed Length Blocks

- When there is a request, the first block is assigned and the block is removed from the free list.
- When a block is freed (or deallocated), the block is back to the free list.



# Fixed Length Blocks

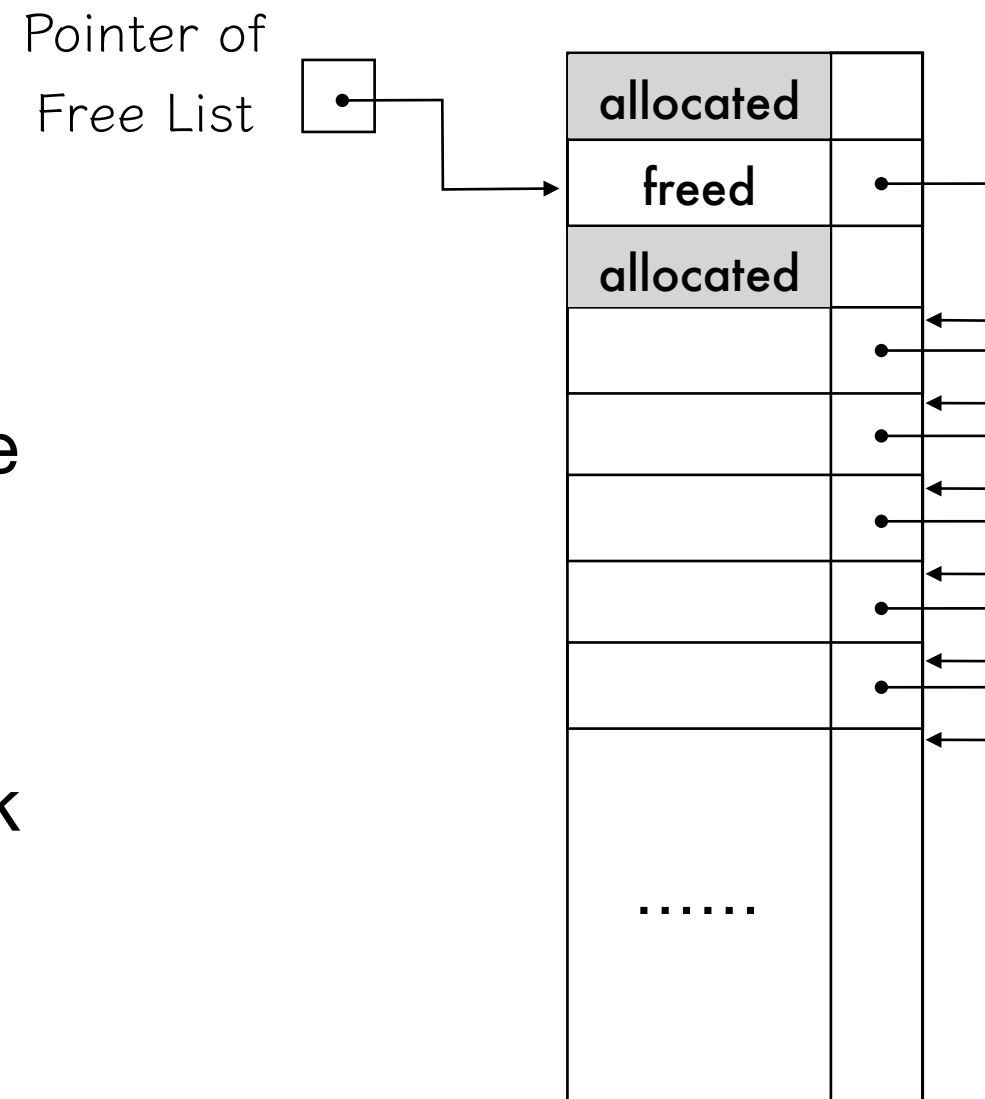
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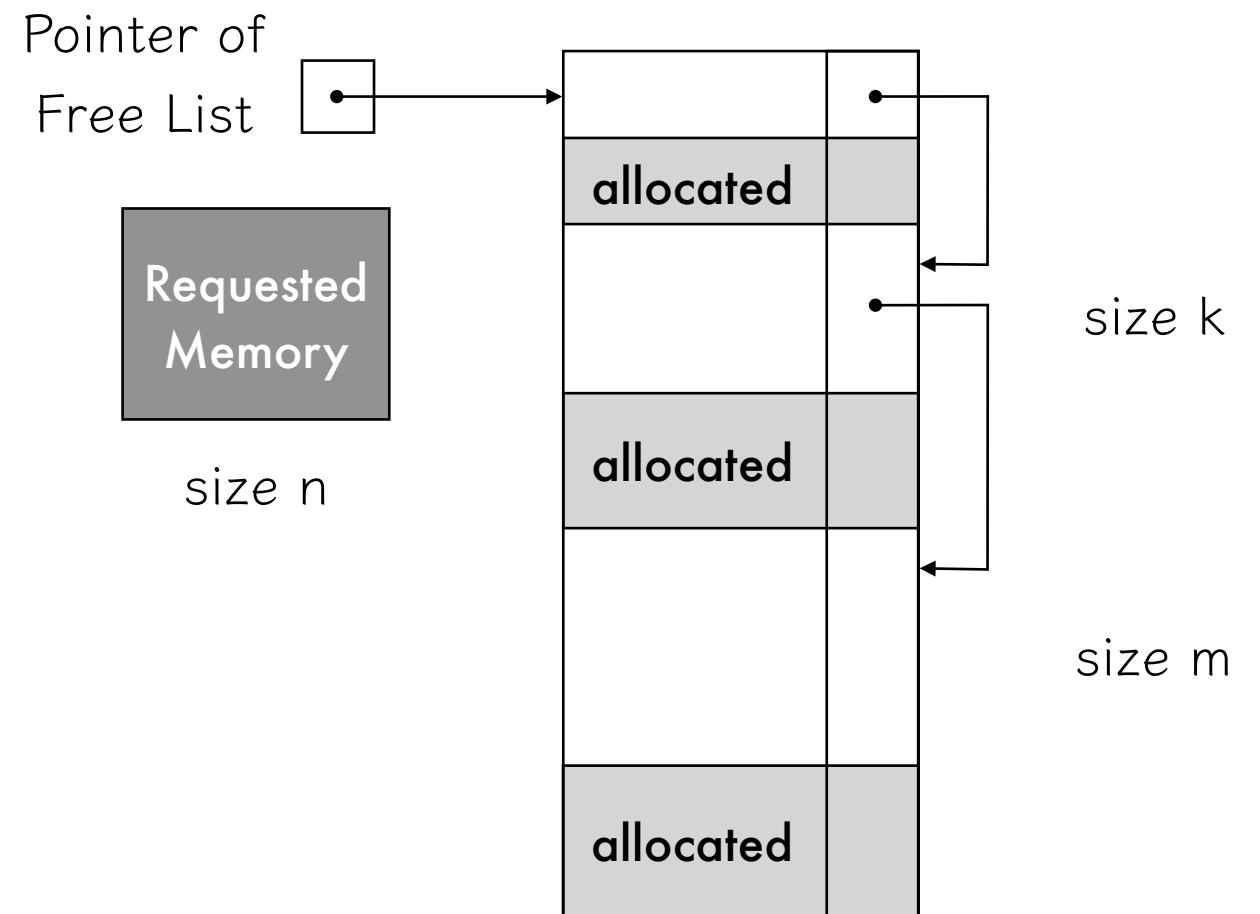
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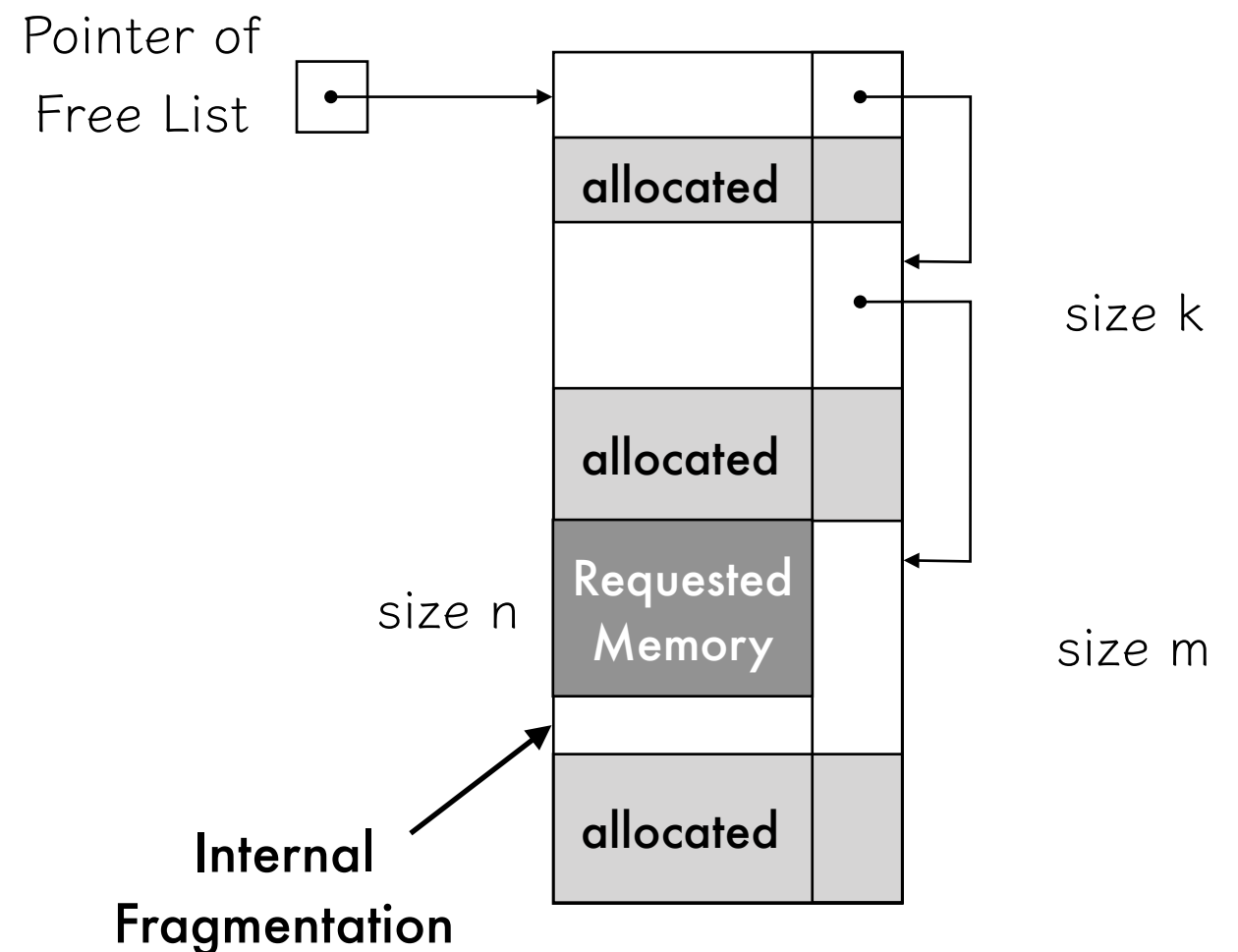
# Variable Length Blocks

- Similar to fixed length blocks, it maintains a free list for available blocks.
- The size of blocks can be different.
- When a request for memory of size  $n$ , it allocates a free block fits to this size.
  - e.g.)  $n > k$  and  $n < m$ .
- The third free block is allocated.



# Fragmentation

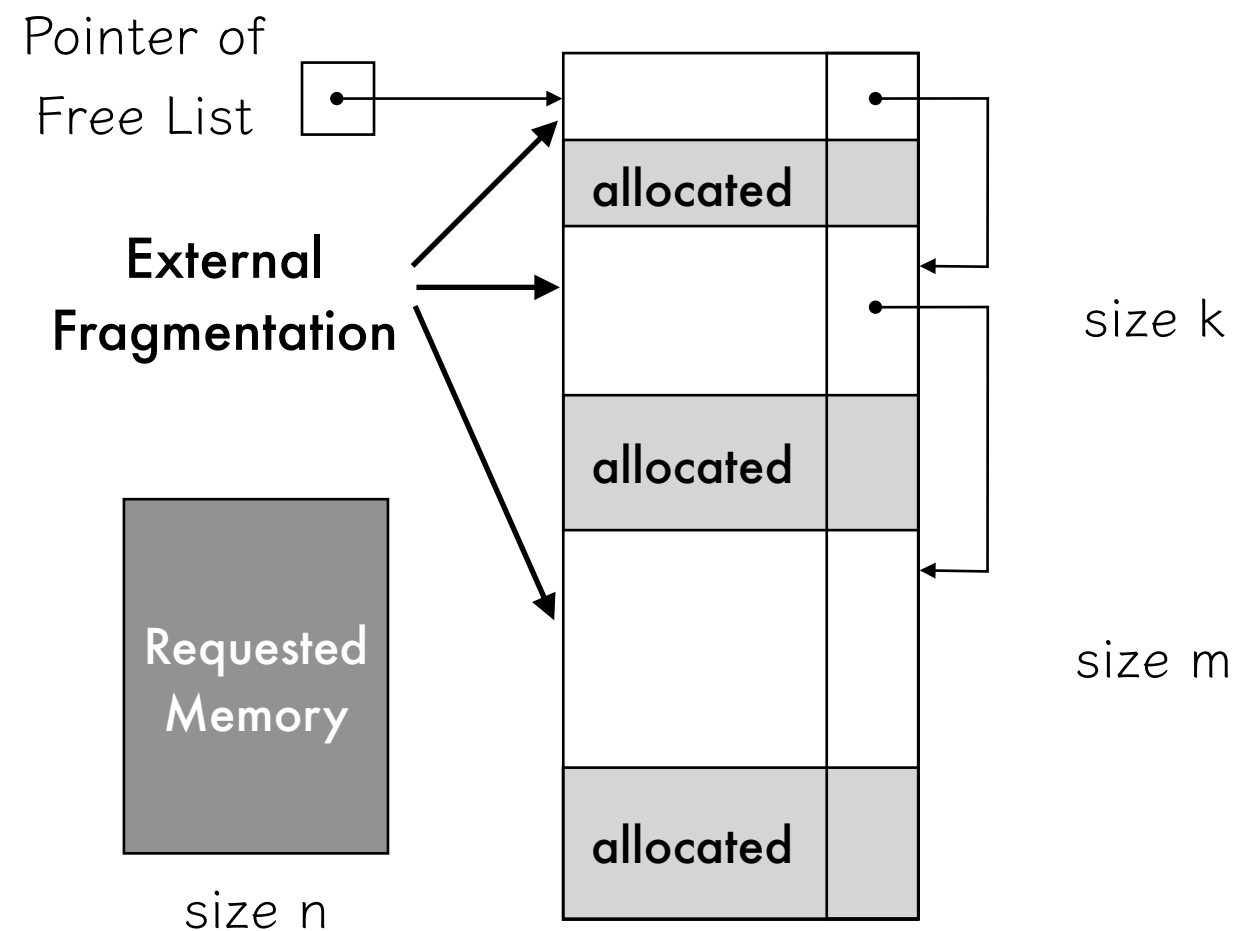
- Variable length method causes fragmentation.
- Due to fragmentation, memory space is wasted, or it reduces the performance of programs.
- **Internal Fragmentation:**  
Allocated block size is greater than the requested size.
  - $m > n$ , then  $d = m - n$  is wasted.



# Fragmentation

- ***External Fragmentation***

- Due to the scattered free blocks, requested memory cannot be allocated,
- even if there exists enough space.
- $m + k > n$ , but they are not consecutive.



# Using Single Free List

- When there is a request for memory allocation of size  $n$ ,
- Directly use the free list.
  - First Fit: allocate the first block bigger than size  $n$ .
  - Best Fit: allocate the size  $k \geq n$  block which has the minimum  $d = k - n$ .
- Free Memory Compaction.
  - When the end of the heap is reached, move all active blocks to the end.

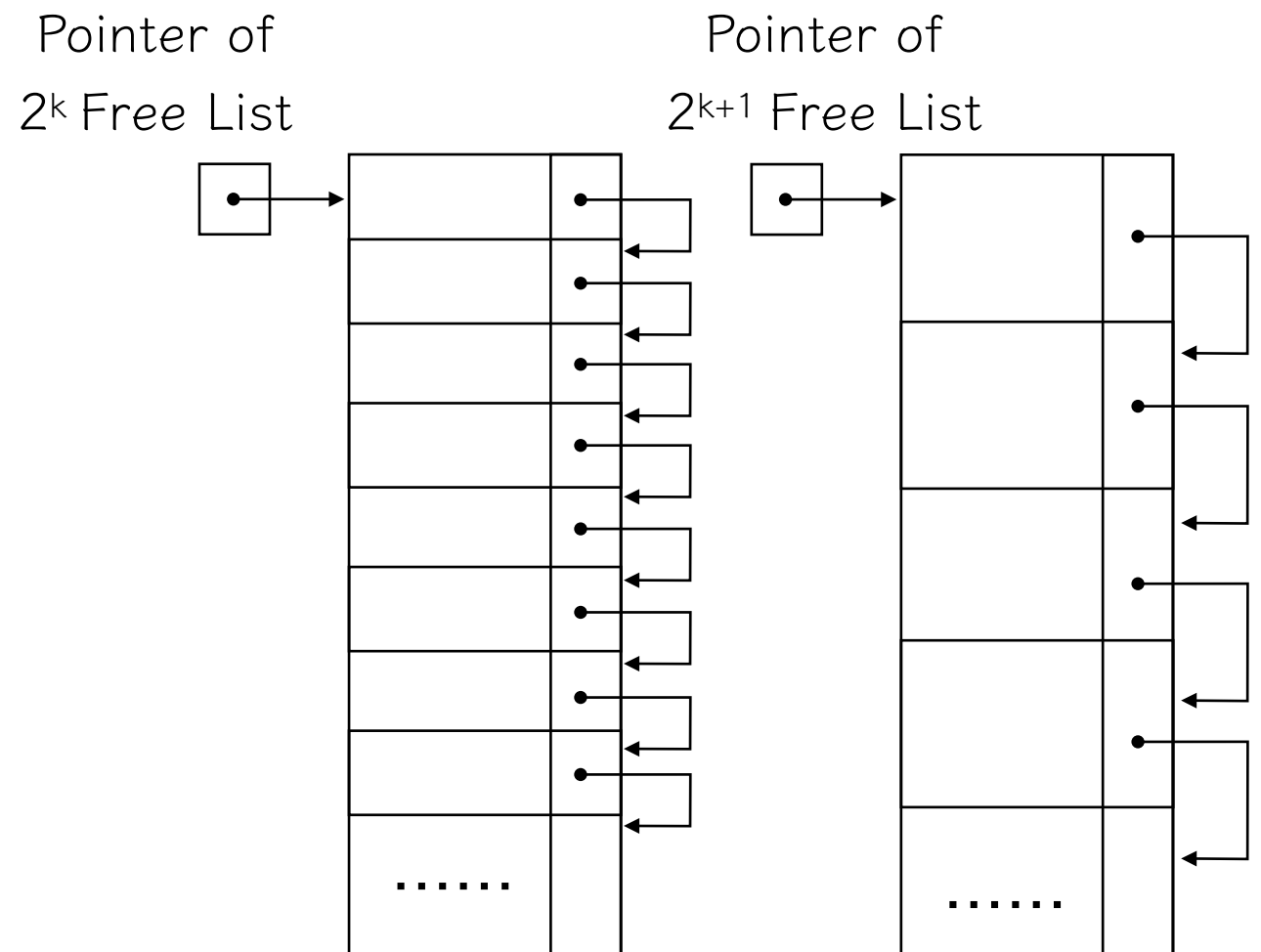
# Multiple Free Lists

- Buddy System
  - Have multiple free lists with size power of 2 (i.e.  $2^n$ ).
  - For size  $n$  request, find a block from the free list of  $2^k \geq n$  blocks.
  - If there is no available block, then search  $2^{k+1}$  free list next.
- Fibonacci Heap
  - Instead of  $2^n$  free lists, use Fibonacci numbers as block sizes in free lists.
  - $\text{Fib}(n) = \text{Fib}(n-1) + \text{Fib}(n-2)$

# Multiple Free Lists

- **Buddy System**

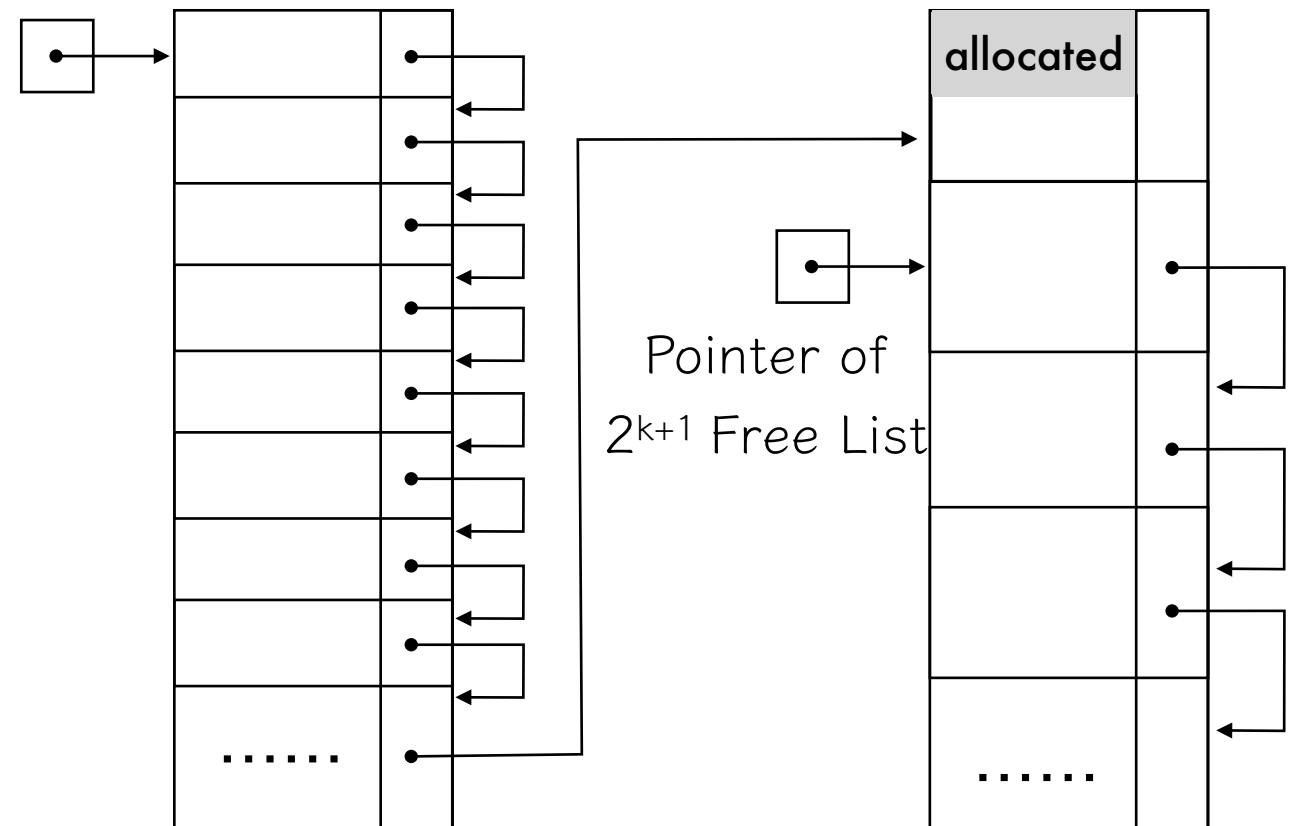
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# Multiple Free Lists

- When a free block is found in  $2^{k+1}$  free list,
  - Split this block into two  $2^k$  blocks.
  - Allocate one of them, and connect the other to  $2^k$  free list.

Pointer of  
 $2^k$  Free List

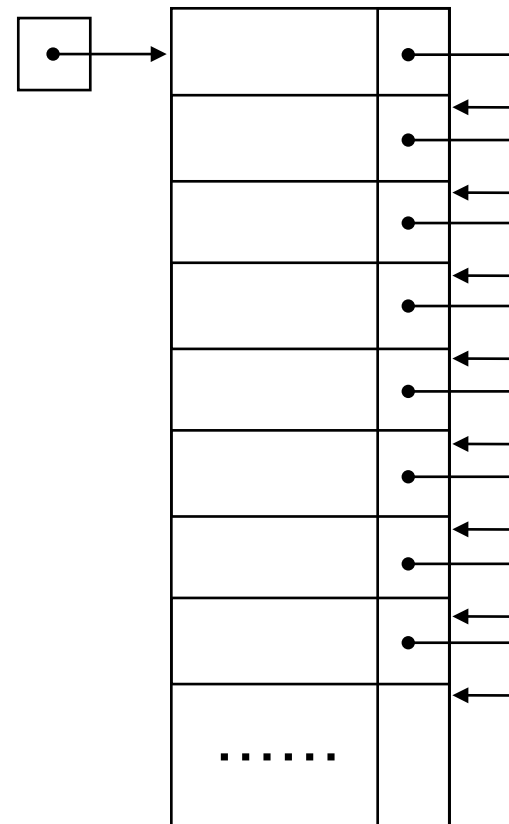




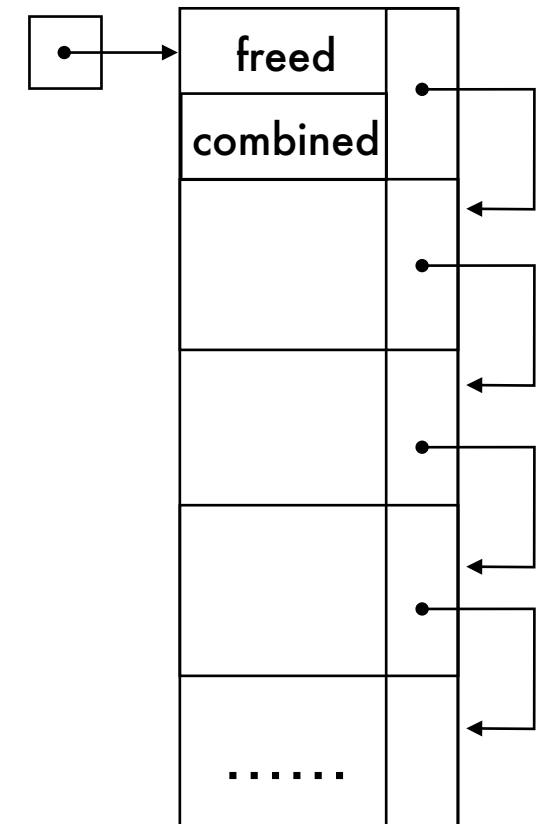
# Multiple Free Lists

- Next time the allocated block is freed,
- Find its *buddy* which is resulted by the split, and check it is also free.
- Combine them and attach it to  $2^{k+1}$  free list again.

Pointer of  
 $2^k$  Free List



Pointer of  
 $2^{k+1}$  Free List



# Summary

- Dynamic Memory Management w/ Heap
  - Differences between Fixed/Variable Length Blocks.
  - Fragmentation
  - Heap Management w/ Single / Multiple Free Lists