

§ 3. Algorithm of Resolution

Now we are ready to prove

$$\text{Ord II in dim } \leq n-1 \xrightarrow{T_1} \text{Ord I in dim } n \xrightarrow{T_2} \text{Ord II in dim } n.$$

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T_1 : We start from $(X, I, E) \quad E = (E^1, \dots, E^s)$, $\max \text{ord } I \leq m$
 If $\max \text{ord } I < m$, the process is trivial. We assume $\max \text{ord } I = m$

Step 1: Construct $T_{I_1}: (X_{r_1}, I_{r_1}, E_{r_1}) \rightarrow \dots \rightarrow (X, I, E)$ s.t.

$$\text{Supp } T_{I_1}^{-1} E \cap \text{Supp } (I_{r_1}, m) = \emptyset.$$

Consider the equivalent ideal $C(I)$ now

1.1. Let Z_0 be the union of all irr comp of E' contained in $\text{Supp}(I, m)$.
 Blow up Z_0 , we have (for irr comp E'^k in E' for example) $T_{I_0}: X \rightarrow X$
 $\max \text{ord}_{E'^k} I = m \Rightarrow \max_{T_{I_0}^{-1} E'^k} \text{ord } I \leq m - m = 0$.
 $\Rightarrow T_{I_0}^{-1} E'^k$ and $\text{Supp}(I_{r_1}, m)$ are disjoint.

1.2. Now, set $S = E'$, $E_S = (E - E')|_S = (0, E^2|_S, \dots, E^s|_S)$

consider $(S, I|_S, m, E|_S)$ apply Ord II in dim $\leq n-1$, we get
 $B\Omega_1: (S_{E_1}, I_{E_1}|_{S_{E_1}}, m, E|_S)$ lift $(X_{r_{E_1}}, I_{r_{E_1}}, m, E_S) \rightarrow \dots \rightarrow (\dots)$
 s.t. $\text{Supp } (I_{r_{E_1}}, m) \cap T_{E_1}^{-1} S = \emptyset$

Inductively, we have $T_{I_1}: (\dots) \rightarrow \dots (\dots)$ s.t. $\text{Supp } T_{I_1}^{-1} E \cap \text{Supp } (I_{r_1}, m) = \emptyset$
 Functionality follows from previous remark.

Step 2: Start from $J = H(C(I_{r_1}))$, $T = X_{r_1}$, $F = (T_{E_1}^{-1} E, E^{s+1}, \dots, E^{s+r_1})$
 (Y, J, m, F).

$\forall y \in \text{Supp } (J, m)$, $\exists U_y$ h.t. $T \subset U_y$. Since in step 1, all blow-up is snc,
 we can take h s.t. H_h snc with F ($T_{E_1}^{-1} E$ away from $\text{Supp}(I, m)$).
 Locally consider $(H_h, J|_{H_h}, m, F|_{H_h})$

Apply Ord II in dim $< n$ and lift it and globalized it to

$$T_{I_2}: (X_{r_2}, I_{r_2}, E_{r_2}) \rightarrow (X_{r_1}, I_{r_1}, E_{r_1})$$

$$\text{s.t. } \text{Supp } (I_{r_2}, m) = \emptyset !$$

(Note, all blow up seq is also for $C(I)$ and I , $T H(C(I)) = C(I)$)

Functionality follows from previous rem.

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□.

T2: We start with a marked triple (X, I, m, E) $E = (E^1, \dots, E^s)$. 2.

Step 0:

We may write $I = N(I) M(I)$, $M(I) = \mathcal{O}_X(-\sum_{i=1}^s E_i)$ and $\text{Supp } N(I)$ does not contain any of E^i !

Rem: if $E = \emptyset$, $M(I) = \emptyset$ and $I = N(I)$.

Step 1: Write $(N(I), 1) + (I, m) = (J, s)$ here s a number,

Write $m_J = \max \text{ord } J$, Run $\text{Ord } I$ to (X, J, E) with m_J , we get (X_1, J_1, E_1) s.t. $\max \text{ord } J_1 = m_{J_1} < m_J$.

(Note, $\text{Supp } (J, m_J)_K \subseteq \text{Supp } (J, s)_K = \text{Supp } (N(I), 1)_K \cap \text{Supp } (I, m)_K$)

Inductively we get

(X_r, J_r, E_r) s.t. $\max \text{ord } J_r < s$.

this implies $\emptyset = \text{Supp } (J, s)_r = \text{Supp } (N(I), 1)_r \cap \text{Supp } (I, m)_r$
 $\text{Supp } (N(I)_r)$.

To sum up, we have $(X_r, I_r, m, E_r) \rightarrow \dots \rightarrow (X, I, m, E)$

s.t. $N(I)_r \cap \text{Supp } (I, m)_r = \emptyset$.

Rem: $N(I)_r$ $N(I_r)$ differs by some exceptional comps $E^{k (k > s)}$
 $M(I)_r$ $M(I_r)$ and is contained in $M(I_r)$.

Step 2: $I = M(I) = \mathcal{O}_X(-\sum a_j E^j)$ $E = (E^1, E^2, \dots, E^s, \dots)$.

2.1 Sub { $E^1, E^2, \dots, E^s, \dots$ } has a lexicographic order.

($x \dots$).

2.2. $\forall x \in X$, set $p(x) = (\{E^{j_1}, \dots, E^{j_k}\})$ the maximal subset (in above order) satisfying

(1) E^{j_i} pass x $\forall 1 \leq i \leq k$.

(2) $\sum a_{j_i} \geq m$ (3) $a_{j_1} + \dots + \hat{a}_{j_i} + \dots + a_{j_k} < m$.

$D_{p(x)} = \bigcap_{i=1}^k E^{j_i}$, and it is the focus that is a maximal component of

$\text{Supp } (I, m)$.

$v = (\max \text{ord } I, \text{member of maximal comp of } \text{Supp } (I, m) \text{ attain maxord})$
 $= (m, n)$

Each time, we blow up $D_{p(x)}$ $\Rightarrow D_{p(x)} = \{x_{ji} = \dots = x_{jk} = 0\}$.

$$\begin{aligned} \forall x \in D_{p(x)}, \quad & \text{since } \sum_{i=1}^k q_{ji} - m < a_{ji} \quad \forall |s| \leq k, \\ I_x = \prod_t x_t^{a_t} & \\ \text{and } I_x = \prod_{\substack{t=s_i \\ \text{in word}}} x_t^{a_t} \cdot x_{j_i}^{\sum a_{ji} - m} & < q_{ji}^{\sum a_{ji} + \hat{a}_{ji} + \dots + a_{jk} - m + a_{ji}} \end{aligned}$$

$$\Rightarrow \text{ord}_y T_x^{-1}(I) < \text{ord}_x I \quad \forall \pi(y)=x.$$

$\Rightarrow \nu$ decrease strictly in the lexicographic order \Rightarrow the procedure terminates with $\nu = (< m, \#)$. $\Rightarrow \max \text{ord } I_r < m$ eventually.

Functionality follows from the sm invariant property of $p(x)$ and ν . \square .

OCE: Commute resp to closed embedding or $\text{Ord } II \underset{Is}{\rightarrow} E \neq \emptyset$

$$T: S \rightarrow X$$

Now, $T^* Is$ in X we have locally

$T^* Is$ is adding some smooth element $\{u\}$.

$$\text{Ord}(T^* Is, 1) = 1 \quad T(Is) = T^* Is \leftarrow \text{sm element.}$$

Recall T_1 : We restrict to sm by \mathfrak{m} in T to do induction, so for this case, we just restrict $T^* Is$ to $\{u\}$ and exactly get I , the procedure commute with closed embedding. \square