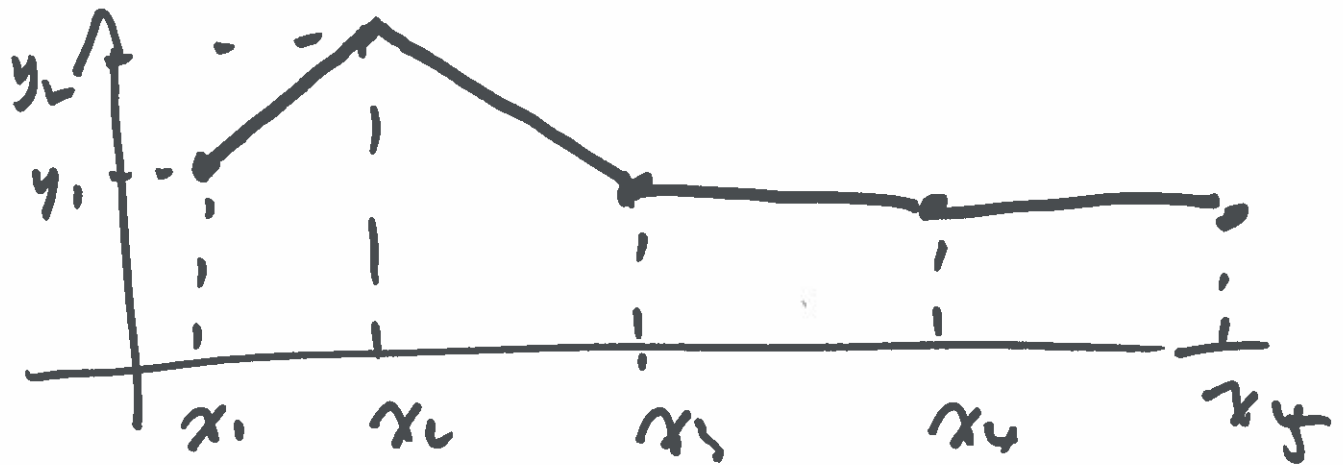


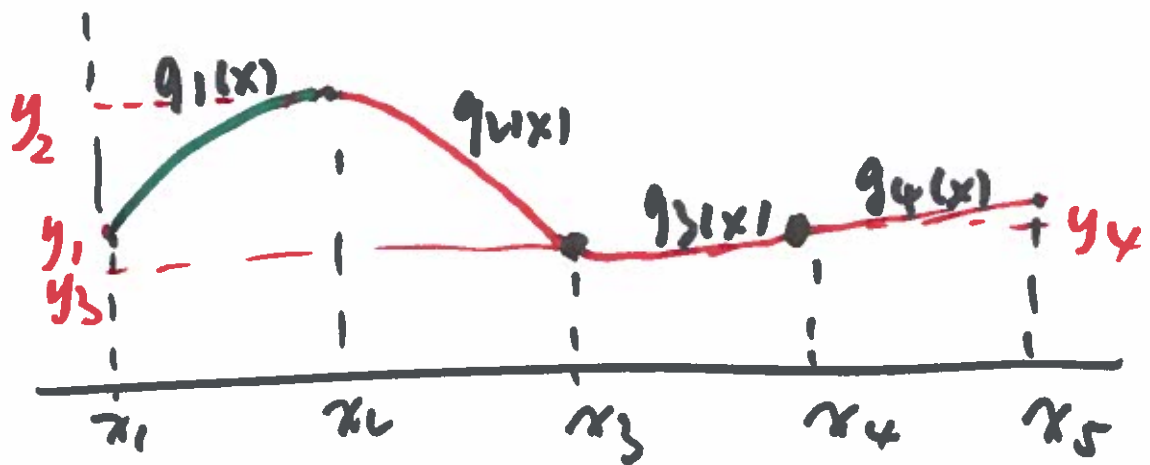
§ 3-2 Interpolation: local

Lecture 15

§ 3-2-1 piecewise linear 3.14.2019



§ 3-2-2 Cubic Splines: better smoothness.



$$g(x) = g_j(x) \quad x \in [x_j, x_{j+1}] \quad (x_j, y_j) \quad j = 1, \dots, n+1$$

property 1:
$$\begin{cases} g_j(x_j) = y_j \\ g_j(x_{j+1}) = y_{j+1} \end{cases}$$

property 2:
$$g_j'(x_j) = g_{j-1}'(x_j) \quad j = 2, \dots, n$$

property 3:
$$g_j''(x_j) = g_{j-1}''(x_j) \quad j = 2, \dots, n$$

$g(x), g'(x), g''(x)$ are all continuous

unique existence

$$g_j(x) = a_j + b_j(x - x_j) + c_j(x - x_j)^2 + d_j(x - x_j)^3$$

\Rightarrow

$$(\tilde{a}_j + \tilde{b}_j x + \tilde{c}_j x^2 + \tilde{d}_j x^3)$$

n intervals $[x_j, x_{j+1}]$ $j=1, \dots, n$

$4n$ unknowns

conditions.

property 1 $\Leftrightarrow 2n$

property 2: $n-1$

property 3: $n-1$

2 conditions short: no unique
nos.

Two more conditions can be imposed at both ends of the data interval

Choice 1: Natural splines

$$g_1''(x_1) = 0, \quad g_n''(x_{n+1}) = 0$$

Choice 2: Clamped splines

$$\begin{aligned} g_1'(x_1) &= \alpha \\ g_n'(x_{n+1}) &= \beta \end{aligned} \quad \begin{array}{l} \alpha, \beta \text{ are} \\ \text{given} \end{array}$$

Choice 3: Not-a-knot spline

$$g_1'''(x_2) = g_2'''(x_2)$$

$$g_{n-1}'''(x_n) = g_n'''(x_n)$$

With any of the choices above,
the cubic spline interpolation
can be uniquely determined.

↪ existence

$$\underline{A} \underline{x} = \underline{b}$$

↑ square and invertible

Example given 3 data points

j	x_j	y_j
1	0	1 (= y_1)
2	$\frac{1}{2}$	-1 (= y_2)
3	1	2 (= y_3)

Find the natural cubic spline interpolation

on $[x_1, x_2] = [0, \frac{1}{2}]$

$$\begin{aligned} g_1(x) &= a_1 + b_1(x-x_1) + c_1(x-x_1)^2 + d_1(x-x_1)^3 \\ &= a_1 + b_1x + c_1x^2 + d_1x^3 \end{aligned}$$

on $[x_2, x_3] = [\frac{1}{2}, 1]$

$$\begin{aligned} g_2(x) &= a_2 + b_2(x-x_2) + c_2(x-x_2)^2 + d_2(x-x_2)^3 \\ &= a_2 + b_2(x-\frac{1}{2}) + c_2(x-\frac{1}{2})^2 + d_2(x-\frac{1}{2})^3 \end{aligned}$$

property

$\Leftrightarrow \begin{cases} a_1 = y_1 = 1 \\ a_2 + \frac{1}{2}b_2 + \frac{1}{4}c_2 + \frac{1}{8}d_2 = -1 \end{cases}$

$$g_1(x_2) = y_2, \quad g_2(x_3) = y_3$$

$$a_2 = y_2 = -1$$

$$a_2 + \frac{1}{2}b_2 + \frac{1}{4}c_2 + \frac{1}{8}d_2 = 2$$

property 2: $g_1'(x_2) = g_2'(x_2)$

$$\Leftrightarrow b_1 + 2c_1x + 3d_1x^2 \Big|_{x=\frac{1}{2}}$$

$$= b_2 + 2c_2(x - \frac{1}{2}) + 3d_2(x - \frac{1}{2})^2 \Big|_{x=\frac{1}{2}}$$

$$\Leftrightarrow b_1 + c_1 + \frac{3}{4}d_1 = b_2$$

property 3: $g_1''(x_2) = g_2''(x_2)$

$$\Leftrightarrow 2c_1 + 6d_1x \Big|_{x=\frac{1}{2}}$$

$$= 2c_2 + 6d_2(x - \frac{1}{2}) \Big|_{x=\frac{1}{2}}$$

$$\Leftrightarrow c_1 + \frac{3}{2}d_1 = c_2$$

Natural spline boundary conditions

$$g_1''(x_1) = 0 \Rightarrow 2c_1 + 6d_1x \Big|_{x=0} = 2c_1 = 0$$

$$g_2''(x_3) = 0 \Rightarrow 2c_2 + 6d_2(x - \frac{1}{2}) \Big|_{x=1} = 2c_2 + 3d_2 = 0$$

$$a_1 = 1, a_2 = -1, c_1 = 0$$

$$d_2 = -10, d_1 = 10, b_1 = -6.5$$

$$\begin{aligned} & b_2 = 1, c_2 = 15 \\ \Rightarrow \left\{ \begin{aligned} g_1(x) &= 1 - 6.5x + 10x^3 \\ g_2(x) &= -1 + (x - \frac{1}{2}) + 15(x - \frac{1}{2})^2 - 10(x - \frac{1}{2})^3 \end{aligned} \right. \quad x \in [0, \frac{1}{2}] \end{aligned}$$

Remark: In general, with more data, all the unknowns can be expressed in terms of $\{c_j\}_{j=1}^n$, and $\underline{c} = \begin{bmatrix} c_1 \\ \vdots \\ c_n \end{bmatrix}$

Satisfies a tr-diagonal system.

$$A\underline{c} = \text{given}$$

Can be solved with $O(n)$ computational complexity

§ 3-2-3 Interpolating errors

'use piecewise linear interpolation as an example'

Assume $\{(x_j, y_j)\}_{j=1}^{n+1}$ are from a given function $f(x)$, namely $y_j = f(x_j)$, we want to bound $|f(x) - g(x)|$, here $g(x)$

is a piecewise linear interpolation

Recall $g(x) = g_j(x)$ on $[x_j, x_{j+1})$

$g_j(x)$ is a polynomial of degree ≤ 1

$g_j(x_j) = y_j$, $g_j(x_{j+1}) = y_{j+1}$

Consider $[x_j, x_{j+1})$, $g_j(x)$ is a global interpolant for (x_j, y_j) and (x_{j+1}, y_{j+1}) based on $f(x)$

Based on the error of the global interpolant.

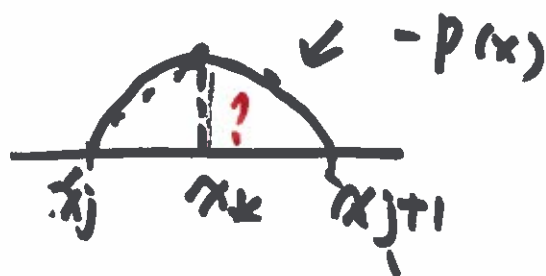
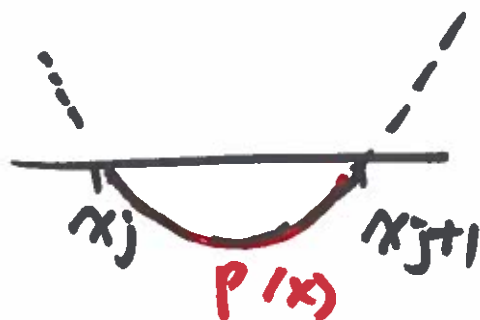
we know ~~as~~ for $x \in [x_j, x_{j+1}]$ $\rightarrow p(x)$

$$f(x) - g_j(x) = \frac{(x-x_j)(x-x_{j+1})}{2} f''(c)$$

c is some number from (x_j, x_{j+1})

$$\Leftrightarrow |f(x) - g_j(x)| \leq \max_{x_j \leq x \leq x_{j+1}} |p(x)| \max_{s \in (x_j, x_{j+1})} |f''(s)|$$

we next take a closer look at $p(x)$



$$(-p(x))' = 0 \Rightarrow x = x_* = \frac{x_j + x_{j+1}}{2}$$

$$(-p(x))''|_{x=x_*} < 0$$

$$-p(x)|_{x=x_*} = \frac{(x_{j+1} - x_j)^2}{8}$$

$$\Rightarrow \left[\begin{array}{l} \text{for } x \in [x_j, x_{j+1}] \\ |f(x) - g_j(x)| \leq \frac{h_j^2}{8} \max_{s \in (x_j, x_{j+1})} |f''(s)| \end{array} \right]$$

Theorem:

Given f on $[a, b]$, f, f', f'' are continuous, consider

$$a = x_1 < x_2 < \dots < x_{n+1} = b$$

and the piecewise linear interpolant $g(x)$ of $f(x)$ $h = \max_{1 \leq j \leq n} |x_{j+1} - x_j|$

then

$$|f(x) - g(x)|$$

$$\leq \frac{1}{8} h^2 \max_{s \in [a, b]} |f''(s)|$$

$$\text{for } x \in [a, b]$$

2nd order accuracy

$$h \rightarrow h/2 \rightarrow h/4$$

$$\text{error} \rightarrow \text{error}/4 \rightarrow \text{error}/16$$

Discussion: for what $f(x)$,

the error in g is zero?

Answer: when $f(x)$ is linear
(polynomial of degree 1)
 $\Leftarrow f''(x) \equiv 0$

§ 3-3 Least squares solution and data fitting

Motivation:

1) $A \in \mathbb{R}^{m \times n}$, $m > n$ $b \in \mathbb{R}^m$

to solve $A\underline{x} = \underline{b}$ for $\underline{x} \in \mathbb{R}^n$

2) To overcome the possible issue of global interpolation (Runge phenomenon)

$$\{(x_j, y_j)\}_{j=1}^{n+1}$$

approximate data by $p_m(x)$, $m < n$

3) Represent or analyze a more scattered data set



Review

matrix - vector multiplication

Given $A = (a_{ij}) \in \mathbb{R}^{m \times n}$

$$\underline{x} = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix} \in \mathbb{R}^n$$

we define $A \underline{x} = \underline{b} = \begin{pmatrix} b_1 \\ \vdots \\ b_m \end{pmatrix} \in \mathbb{R}^m$

$$\text{as } b_1 = \sum_{j=1}^n a_{1j} x_j$$

$$b_2 = \sum_{j=1}^n a_{2j} x_j$$

\vdots

A new view point

$$\underline{b} = \begin{pmatrix} b_1 \\ \vdots \\ b_m \end{pmatrix} = \begin{pmatrix} \sum_{j=1}^n a_{1j} x_j \\ \sum_{j=1}^n a_{2j} x_j \\ \vdots \\ \sum_{j=1}^n a_{mj} x_j \end{pmatrix}$$

$$= \sum_{j=1}^n x_j \begin{pmatrix} a_{1j} \\ a_{2j} \\ \vdots \\ a_{mj} \end{pmatrix}$$

$\leftarrow j^{\text{th}}$
column
of A

Let $\underline{a}_j = \begin{pmatrix} a_{1j} \\ a_{2j} \\ \vdots \\ a_{mj} \end{pmatrix}$ be the j th column of A

then $\underline{b} = A\underline{x} = \sum_{j=1}^n x_j \underline{a}_j$
 $= x_1 \underline{a}_1 + x_2 \underline{a}_2 + \dots + x_n \underline{a}_n$

That is $A\underline{x}$ is a linear combination of column vectors of A , with the coefficients being the entries of \underline{x} .

notation

$$A = [\underline{a}_1, \underline{a}_2, \dots, \underline{a}_n]$$

$$A\underline{x} = (\underline{a}_1, \underline{a}_2, \dots, \underline{a}_n) \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix} \\ = x_1 \underline{a}_1 + x_2 \underline{a}_2 + \dots + x_n \underline{a}_n$$

Definition

Range of A , also denoted as $\text{range}(A)$

$$\text{range}(A) = \{ x_1 \underline{a}_1 + x_2 \underline{a}_2 + \dots + x_n \underline{a}_n : \\ \forall x_i \in \mathbb{R}, i=1, \dots, n \} \\ \subset \mathbb{R}^m$$

linearly independent:

$$\underline{a}_1 \dots \underline{a}_n$$

when none of them is a
linear combination of the
other

$$\text{Or } \sum_{i=1}^n x_i \underline{a}_i = \underline{0} \Leftrightarrow x_1 = x_2 = \dots = x_n = 0$$

(no redundancy)