

# KV-Auditor: Auditing Local Differential Privacy for Correlated Key–Value Estimation

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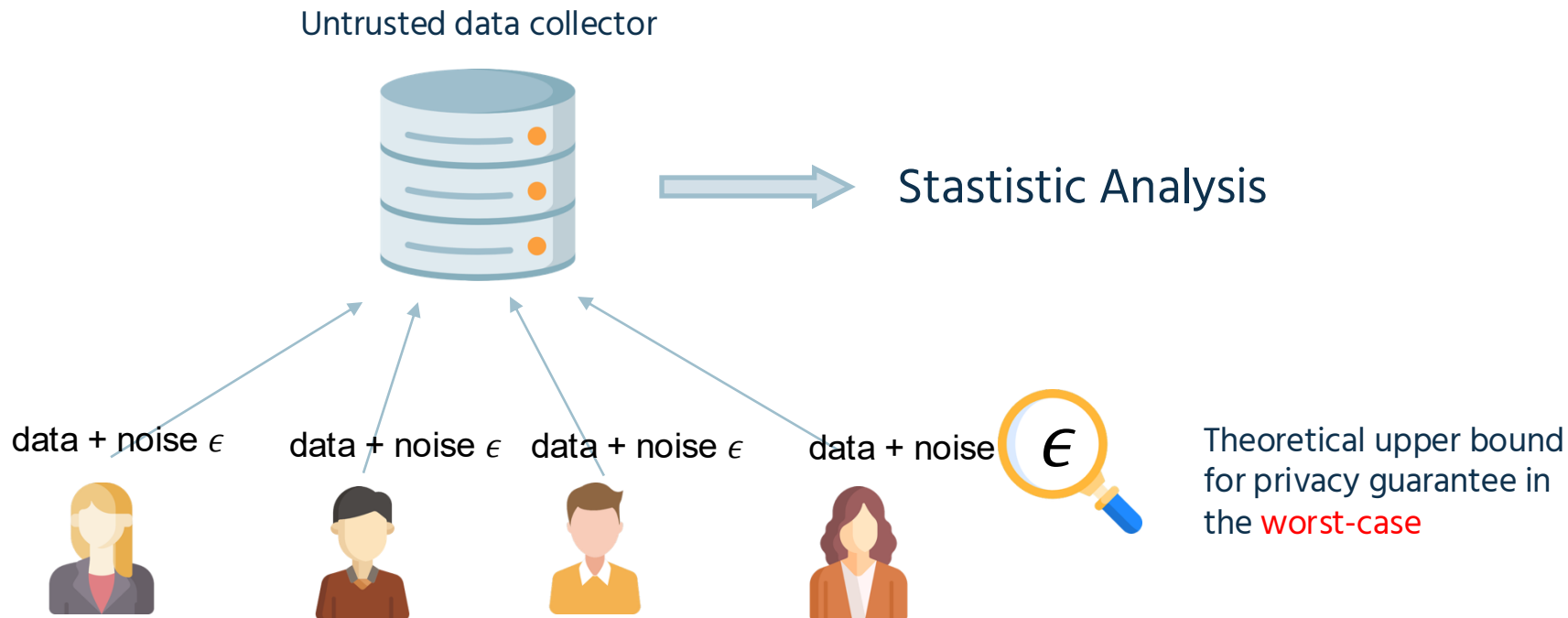


# Privacy Challenges in Data Collection



User data collection introduces privacy risks

# Local Differential Privacy (LDP)



# Limitations of Theoretical $\epsilon$

## Algorithm Designers



- Proof/Implementation errors
- Overly loose bound

$\epsilon$ -LDP

## Users



$\epsilon = 0.5 ?$

- Hard to understand

# Privacy Auditing for Differential Privacy

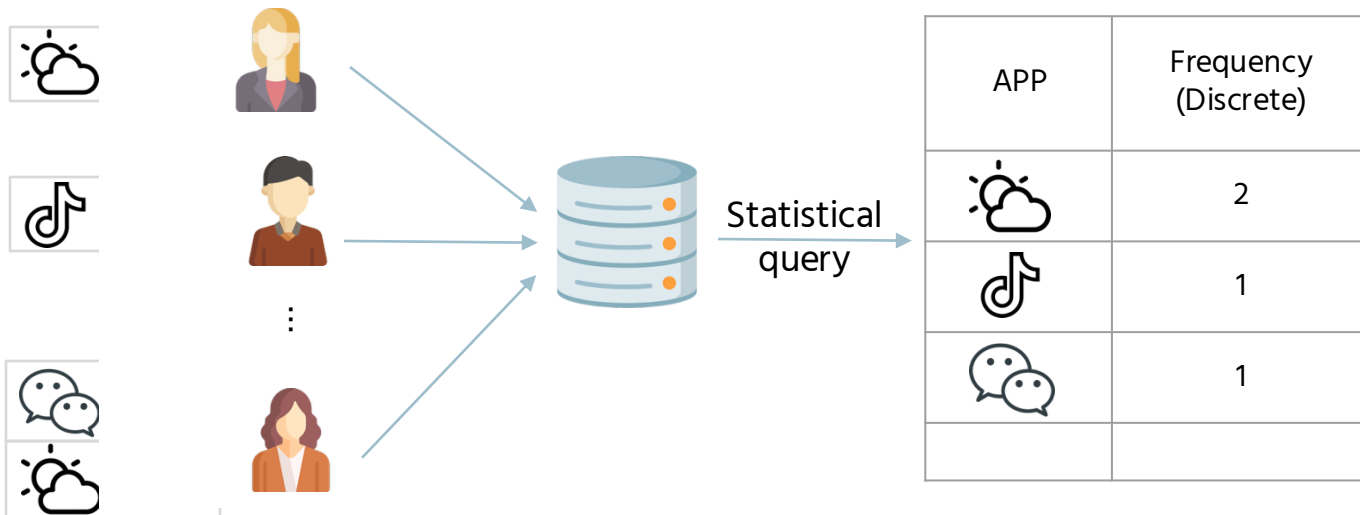


# Privacy Auditing for Differential Privacy

- ✓ Auditing under Centralized Setting
- ✓ Auditing LDP protocols for discrete data (frequency estimation)
- ✗ Auditing LDP protocols for key-value data

## Challenges:

- Continuous values  $\rightarrow$  cannot be enumerated
- Key-value correlation



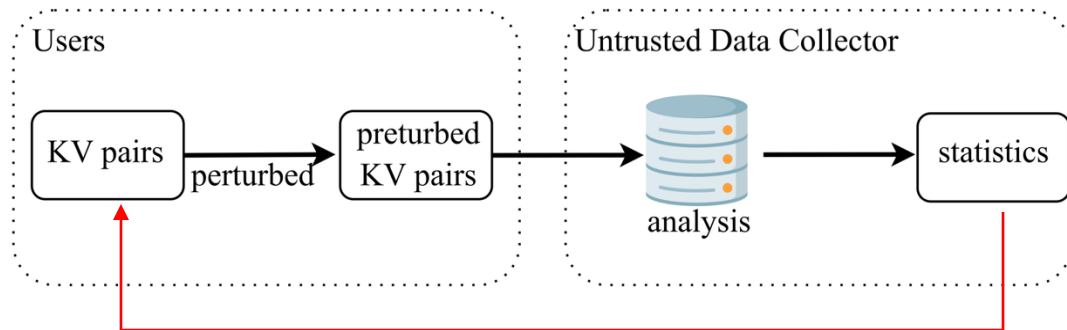
# LDP Protocols for Correlated Key-Value Estimation

## Interactive protocols

- PrivKVM -- multi-round estimation
- PrivKVM\* -- multi-bucket extension

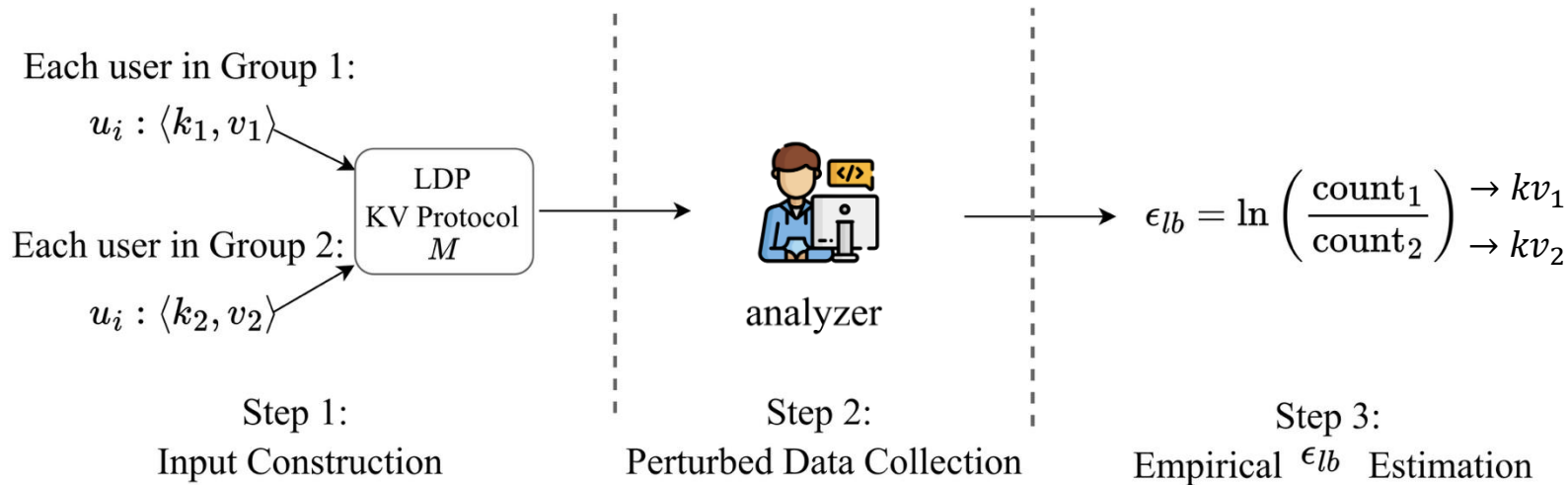
## Non-interactive protocols

- PCKV -- one-shot perturbation with padding & sampling



aggregated statistics returned (interactive only)

# Workflow of KV-Auditor



CPP-UE/CPP-GRR:  $kv_1 = \langle k, 1 \rangle, kv_1 = \langle k, -1 \rangle$   
 PCKV-UE/PCKV-GRR:  $kv_1 = \langle k_1, 1 \rangle, kv_1 = \langle k_2, -1 \rangle$   
 $k_1 \neq k_2$


Calculate with the intersection of the perturbed data




# KV-Auditor for Non-interactive Protocols

- **KV-Auditor for Non-interactive Protocols Perturbed data keeps a key-value structure; format depends on the mechanism.**

UE: perturbed key + perturbed vector (e.g., PCKV-UE, CPP-UE).

 , 0.5h  $\rightarrow \langle k = 1, v = [1, 0, 0, 0] \rangle$

GRR: perturbed key + perturbed integer (e.g., PCKV-GRR, CPP-GRR).

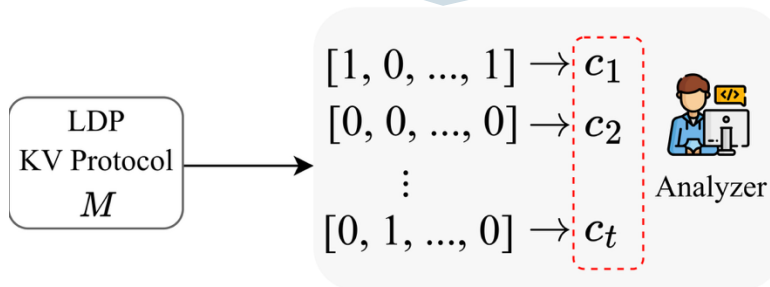
 , 0.5h  $\rightarrow \langle k = 1, v = -1 \rangle$

- **We design two auditors: Horizontal KV-Auditor (HKV-Auditor) and Vertical KV-Auditor (VKV-Auditor) to estimate the empirical lower bound.**

# KV-Auditor for Non-interactive Protocols

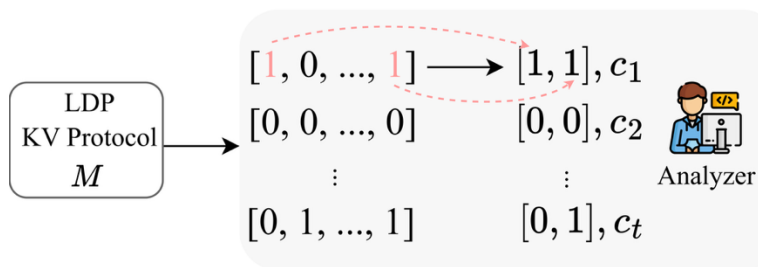
## HKV-Auditor

Treats the perturbed record as a unit



- ✓ Provides tighter  $\epsilon_{lb}$
- ✗ Requires more than  $10^8$  users when the bit length  $b$  is large.

## VKV-Auditor



Collects two bits from each perturbed vector

- ✓ Has shorter auditing time.
- ✗ Ignores bit dependencies.



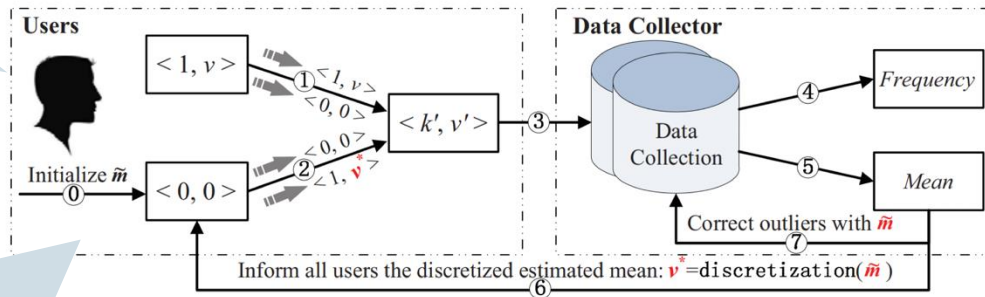
# Challenges of KV-Auditor for Interactive Protocols

- *Does the increment of privacy leakage diminish as the number of iterations increases?*
- *What is the underlying reason for the deceleration in the rate of privacy loss increment?*
- *Can the theoretical upper bound of the allocated budget for each iteration be further tightened?*



# Auditing with KV-Auditor

In auditing, user data fixed  $\rightarrow$  constant privacy leakage



In auditing, mean changes until stable  $\rightarrow$  changing privacy leakage

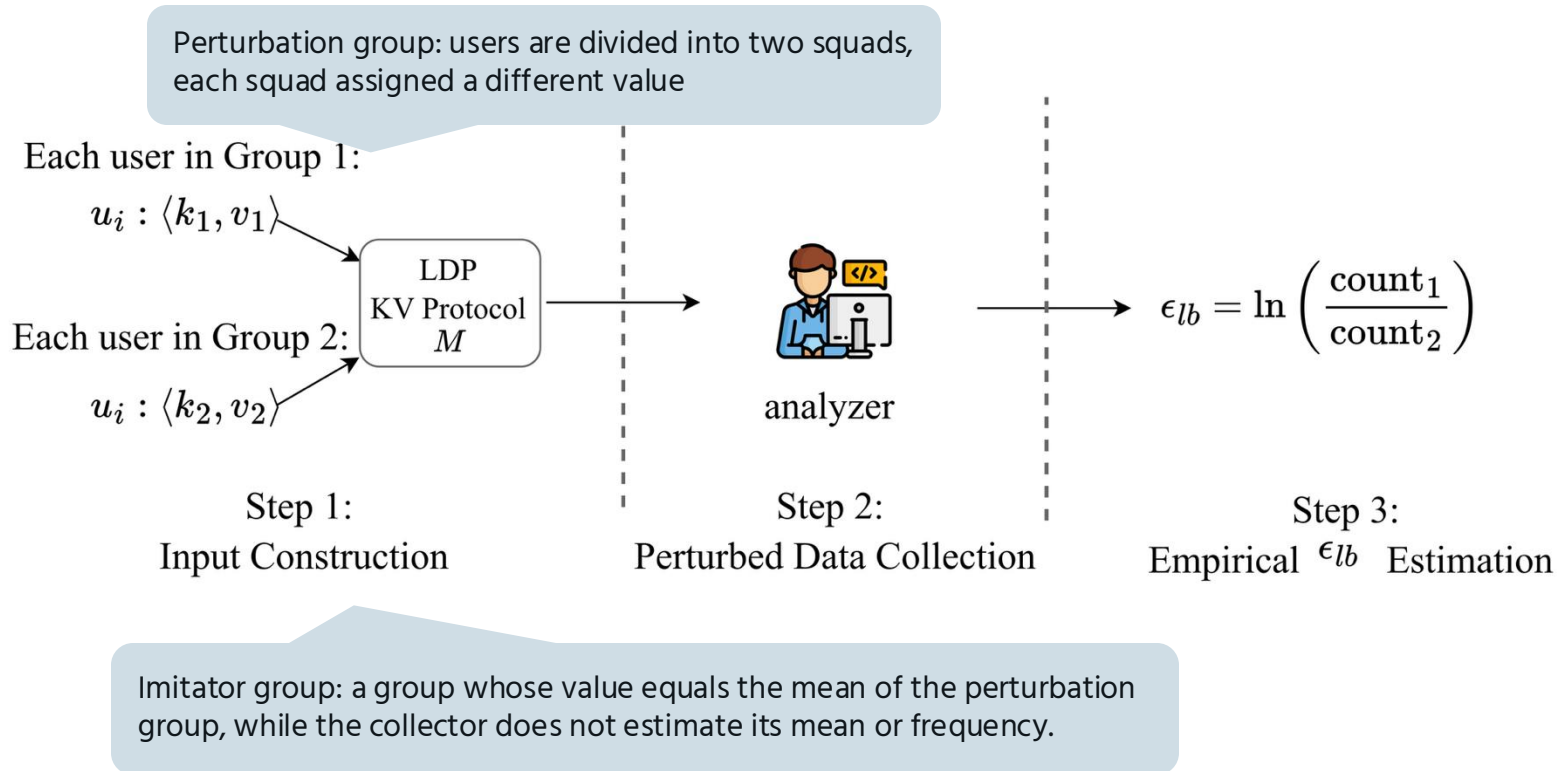
## Stage 1: Distribution Collection

- Perturb each user 10 times, assume all possess KV pair.
- The perturbed distribution is the average over 10 runs

## Stage 2: Distribution Separation

- Perturb once across 10 iterations
- Analyzer separates the mean-perturbed distribution (scaling applied)

# KV-Auditor with Segmentation



# Experiments

- **Audited LDP protocols**

Non-interactive protocols: PCKV-UE, PCKV-GRR

Interactive protocols (PrivKVM / PrivKVM\*): CPP-UE, CPP-GRR; CPP-UE\*, CPP-GRR\*

- **Input Construction**

PCKV-UE, PCKV-GRR:  $kv_1 = \langle k=k_1, v = 1 \rangle$ ,  $kv_2 = \langle k=k_2, v = -1 \rangle$   $k_1 \neq k_2$

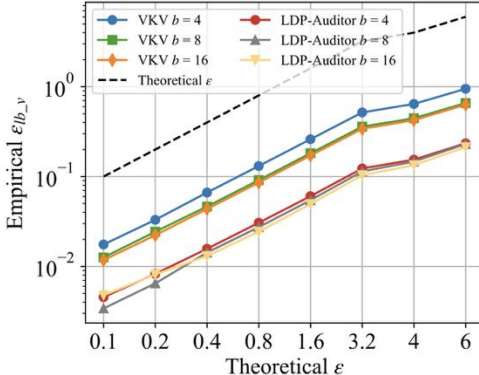
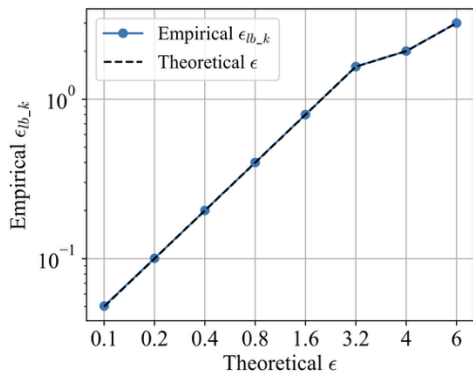
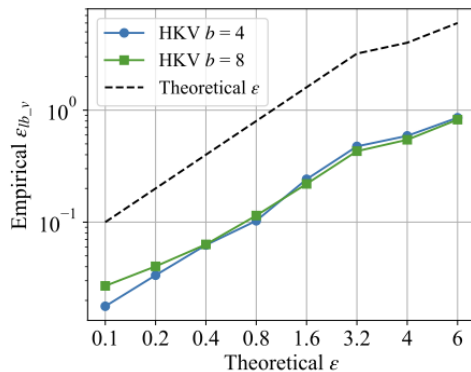
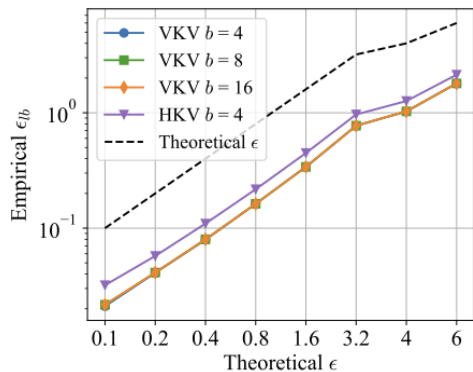
CPP-UE and CPP-GRR:  $kv_1 = \langle k=k, v = 1 \rangle$ ,  $kv_2 = \langle k=k, v = -1 \rangle$ ,

- **Observed Reports by Analyzer(after perturbation):**

UE  $\rightarrow$  binary vector (not one-hot)

GRR  $\rightarrow$  randomized key + randomized value

# Experiments



- For PCKV-GRR, large  $N_{key}$  causes more privacy leakage.
- HKV-Auditor is inaccurate for UE with a large bit length  $b$ , while VKV-Auditor is stable.
- Our KV-Auditor is tighter than the LDP-Auditor.
- The theoretical of keys in CPP is tight.

# Experiments

Theoretical	Empirical $\epsilon_{lb}$				
$\epsilon$	$c = 1$	$c = 2$	$c = 3$	$c = 4$	$c = 5$
0.1	0.0068	0.0092	0.0134	0.0140	0.0139
0.2	0.0076	0.0152	0.0214	0.0234	0.0252
0.4	0.0080	0.0298	0.0366	0.0410	0.0406
0.8	0.0148	0.0460	0.0646	0.0710	0.0752
1.6	0.0538	0.1224	0.1237	0.1257	0.1261
3.2	0.3064	0.4241	0.4248	0.4250	0.4265
4	0.5020	0.5580	0.5569	0.5568	0.5553
6	1.0392	0.6218	0.6343	0.6390	0.6475

In CPP-UE, privacy leakage decreases with iterations due to mean convergence.

Theoretical	Empirical $\epsilon_{lb}$				
$\epsilon$	$c = 1$	$c = 2$	$c = 3$	$c = 4$	$c = 5$
0.1	0.0096	0.0078	0.0082	0.0070	0.0071
0.2	0.0165	0.0112	0.0098	0.0092	0.0085
0.4	0.0270	0.0161	0.0164	0.0153	0.0143
0.8	0.0537	0.0299	0.0275	0.0265	0.0237
1.6	0.1240	0.0558	0.0528	0.0481	0.0495
3.2	0.3090	0.1103	0.1069	0.1078	0.1051
4	0.4219	0.1431	0.1408	0.1364	0.1362
6	0.7327	0.2352	0.2292	0.2296	0.2281

In SHKV-Auditor, empirical  $\epsilon_{lb}$  decreases with iterations as the mean converges, with the first iteration showing the highest leakage.





# Summary

- We introduce a KV-Auditor framework to estimate the  $\epsilon_{lb}$  of LDP protocols for key-value data.
- Based on this framework, we propose HKV-Auditor and VKV-Auditor for non-interactive protocols, SKV-Auditor for interactive protocols.
- The upper bound for GRR is tighter than that for UE, indicating greater room for improvement in UE.



# Thank you!

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[arxiv.org/abs/2508.11495](https://arxiv.org/abs/2508.11495)

[github.com/JingnanXu97/KV-Auditor](https://github.com/JingnanXu97/KV-Auditor)

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