CIVIL-557

Decision aid methodologies in transportation

Lab 5:

Heuristics:

Column Generation

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■ Lab I & 2 : Gurobi, VRP

Lab 3: Labeling Algorithm, SPP

Lab 4: Heuristics

Lab 5: Coding a heuristic for the VRP that uses the

labeling algorithm to solve an SPP to iteratively

add columns to a master problem that we solve

with Gurobi 2220 0000





 Labeling algorithm to solve shortest path problem (SPP) given rewards for all customers as input

Today:

SPP is the **subproblem**when solving the vehicle routing problem
(VRP) using **branch and price**

or:

when solving the VRP using a **column** generation heuristic





- Branch & Price:
 - I. Take **VRP MILP** (set partitioning formulation)

Minimize
$$\sum_{r \in \Omega} c_r \lambda_r$$
 $s.t. \sum_{r \in \Omega} a_{ir} \lambda_r = 1 \qquad \forall i \in N,$ $\sum_{r \in \Omega} \lambda_r = |K|,$ $\lambda_r \in \{0,1\} \qquad \forall r \in \Omega.$





Branch & Price:

2. Start with some **dummy variable** routes:

Example:

 $\lambda_i = ext{route from depot to customer } i ext{ and back}$ $cost_i = 2 * ext{distance(depot, } i)$ $a_{ij} = 1 ext{ if } i = j, 0 ext{ else}$ $\Omega^0 = \{\lambda_1, ..., \lambda_n\} ext{ set of all dummy routes}$





- Branch & Price:
 - 3. Solve the **relaxed** restricted **master problem**:

$$egin{aligned} extit{Minimize} & \sum_{r \in \Omega} c_r \lambda_r \ & s.t. \sum_{r \in \Omega} a_{ir} \lambda_r = 1 \ & orall i \in extit{N}, \qquad (arphi_i) \ & \sum_{r \in \Omega} \lambda_r = |\mathcal{K}|, \ & \lambda_r \in [0,1] \ & orall r \in \Omega^0 \end{aligned}$$





Branch & Price:

4. Read values of the **dual variables** (reduced cost of visiting a customer) corresponding to constraints (φ_i) :

$$reward_i = \varphi_i$$

- 5. Solve the SPP with these rewards
 - Compute shortest path = route *r*
 - Compute cost c_r = total length of route r
 - Compute incidence vector a_{ir} of route r
 - Compute the reduced cost of the solution: total cost minus the collected rewards

$$c_r - \sum_i a_{ir} \varphi_i$$
 if < 0, adding route r to MP decreases total cost



- Branch & Price:
 - 6. Add the route to the master problem
 - 7. Resolve the master problem, read dual variables, solve subproblem with new rewards, etc.
 - 8. Repeat until reduced costs of SPP are zero (no more possible improvements)
 - 9. Resolve the master problem as **MILP**
 - Solution = column generation heuristic
 - ⇒ not guaranteed to be optimal





Branch & Price:

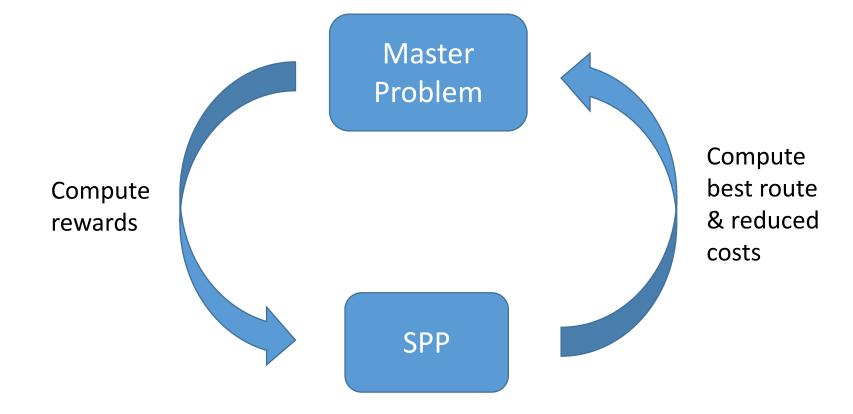
In **Branch & Price** step 9 is instead:

- 9. Resolve the master problem as a **relaxation**. We might get some routes with fractional values.
 - \Rightarrow **Branch** on fractional variables
 - ⇒ Repeat process for both child nodes
 - ⇒ Continue until no more fractional values in solution
 - ⇒ Guaranteed to be **optimal**





Column generation:







- In this exercise class we will do only the column generation heuristic
- Start from code for solving the SPP with loads, vehicle capacity and time window constraints from week 3
- Extend it to fit in a column generation procedure





Minimize
$$\sum_{r \in \Omega} c_r \lambda_r$$
 $s.t. \sum_{r \in \Omega} a_{ir} \lambda_r = 1 \qquad orall i \in \mathcal{N},$ $\sum_{r \in \Omega} \lambda_r = |\mathcal{K}|,$ $\lambda_r \in [0,1] \qquad orall r \in \Omega^0$





$$Minimize \sum_{r \in \Omega} c_r \lambda_r$$

$$s.t. \sum_{r \in \Omega} a_{ir} \lambda_r \trianglerighteq 1 \qquad \forall i \in N,$$

$$\sum_{r \in \Omega} \lambda_r \lozenge \bigwedge \bigwedge \bigwedge \bigwedge Not \ necessary$$

$$\lambda_r \in [0,1] \qquad \forall r \in \Omega^0$$





```
def initialize_master(n):
   # Create a new model
   m = gp.Model("Master Problem")
   # Create variables for dummy routes (create a dicitionary of variables)
   # Set objective
   # Add constraint that every costumer needs to be visited by at least one
   # vehicle (store constraints in a dictionary as you would for variables):
   # Optimize model
   m.setParam("OutputFlag", 0)
   m.optimize()
   # read dual variables of costumer constraints (create a dictionary called
   # "reward" with customers as keys)
   # set the reward for the depot to be 0
    return m, reward, visit, lambda vars
```





```
# creating gurobi variables
m.addVar(lb, ub, vtype=GRB.CONTINUOUS / GRB.BINARY, name)
# setting objective
m.setObjective(objective expression, GRB.MINIMIZE / GRB.MAXIMIZE)
# adding constraints
m.addConstr(variable >= value)
# sums
gp.quicksum(variable[i] for i in range(number))
# solution value of variables
variable.x
# dual values of constraints
constraint.Pi
```





Exercise 2: Finish the column generation implementation

```
# initialize master problem
master, first_reward, visit, lambda_vars = initialize_master(n)

iteration = 0
reward = first_reward
red_cost = -100000
tours = dict()
```





```
# while there exist positive reduced costs, generate new columns
while red_cost < -1e-13:</pre>
    iteration += 1
    if iteration % 10 == 0:
        print(f"iteration = {iteration}, reduced costs = {red_cost}")
    l, red_cost = labeling_algorithm(reward)
    # generate tour
    tour = []
    while l.parent:
        tour.append(l.customer)
        l = l.parent
    tour.append(Start_id)
    tour.reverse()
    tour.append(Start_id)
    tours[iteration] = tour
    tour_length = total_distance(tour, distance)
    # create new column a[:, r] for the new found route
    # add the route as new variable to the master problem
    # update customer constraints:
    # resolve master problem
    master.optimize()
    if iteration % 10 == 0:
        print("obj function = ", master.ObjVal)
    # read dual variables of costumer constraints (rewards)
```





```
# count how many times an element appears in a list
List.count(element)

# adding a new variable to a gurobi model (including objective coefficient)
model.addVar(lb, ub, obj=coeff, vtype, name)

# update the coefficient of a variable in a constraint
model.chgCoeff(constraint, variable, coefficient)
```





Exercise 3: Dealing with cycles

```
# Feasibilty check
def feasible(l, i):
    """Returns if extending from label l to node i is feasible"""
    #Is it feasible to go to node i from label l?
    if l.customer == i:
        return False
    if distance[(l.customer,i)] +l.time > late[i]:
        return False
    if l.load + demand[i] > Q:
        return False
    return True
```



