COMPSYS 304 Assignment 3

Task 1 Cache measurement

Name of the processor: Intel(R) Core(TM) i5-7300HQ CPU @ 2.50GHz

LEVEL 1 Cache size:

Instruction cache size: 32768 – 32 KB Data cache size: 32768 – 32 KB

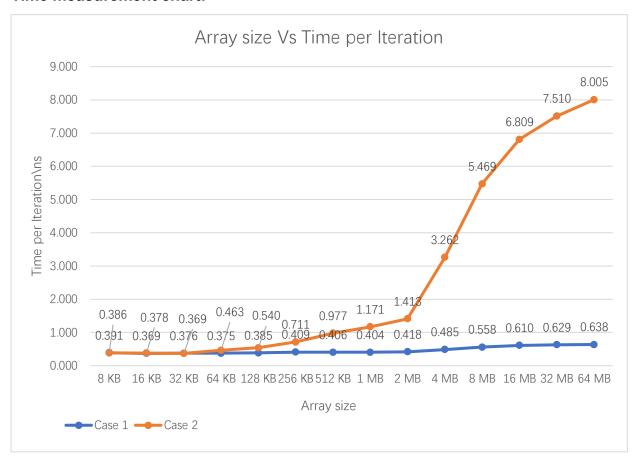
LEVEL 2 Cache size: 262144 – 256 KB

LEVEL 3 Cache size: 6291456 – 6 MB

Time measurement table:

N	size of a	time per iteration /ns	time per iteration /ns
		Case 1	Case 2
2048	8 KB	0.391	0.386
4096	16 KB	0.369	0.378
8192	32 KB	0.376	0.369
16384	64 KB	0.375	0.463
32768	128 KB	0.385	0.540
65536	256 KB	0.409	0.711
131072	512 KB	0.406	0.977
262144	1 MB	0.404	1.171
524288	2 MB	0.418	1.413
1048576	4 MB	0.485	3.262
2097152	8 MB	0.558	5.469
4194304	16 MB	0.610	6.809
8388608	32 MB	0.629	7.510
16777216	64 MB	0.638	8.005

Time measurement chart:



Performance analysis:

For case 1, the access pattern is going through the array linearly which provides a good spatial locality. All the subsequent elements of the array will be loaded into the cache so the next memory access can use the data in the cache block directly. Therefore, the time per iteration has a slow uniform growth after the working set is larger than the cache size.

For case 2, the access pattern is going through the array randomly which will greatly increase the miss rate as the next memory access will most likely not be in the cache when the working set is larger than the cache size. Therefore, the time per iteration has a rapid and uneven growth after the working set is larger than the cache size.

The time per iteration of both case 1 and case 2 are similar when the array size is below 32 KB. Because all the data of the array can be loaded into L1 cache, so both case 1 and case 2 have similar miss rate. However, after the array size is larger than L1 cache size (32KB) but smaller than L3 cache size (6 MB), the time per iteration of case 2 has a relative uniform growth as the average memory access time of L2/3 is higher than L1. Finally, when the working set is greater than L3, the time taken is significantly increased because a large number of main memory accesses are required.

Task 2 Matrix product

Matrix 1 Time taken: 2.03 secs

In this implementation, the access pattern of array A is row-major which provides a good spatial locality as the subsequent elements will be loaded into the same cache line. However, the access pattern of array B is column major. The size of array is 1000×1000 which is 8000 bytes per row. Most of the sequential elements will not load into the same cache line. Therefore, the performance is reduced by not using all words of each cache line.

Matrix 2 Time taken: 1.29 secs

Matrix 3 Time taken: 0.81 secs

Based on the original approach, instead of doing multiplication on the N x N matrices, we use blocking algorithm that has a submatrix of size $\mathbf{k} \times \mathbf{k}$ ($\mathbf{k} = 16$). This can be done by adding two outer for-loop with an increment of \mathbf{k} . Therefore, the $\mathbf{k} \times \mathbf{k}$ submatrix of array \mathbf{b} and a row of length \mathbf{k} of array \mathbf{c} can fit in the cache. Both array \mathbf{b} and array \mathbf{c} can reuse \mathbf{k} times each time the data are brought in.

Appendix:

```
Q
      jinkai@jinkai-OMEN-by-HP-Laptop-15-ce0xx: ~/Desktop/com...
jinkai@jinkai-OMEN-by-HP-Laptop-15-ce0xx:~/Desktop/compsys304 A3/task 2$ make
make: Nothing to be done for 'default'.
jinkai@jinkai-OMEN-by-HP-Laptop-15-ce0xx:~/Desktop/compsys304_A3/task_2$ ./matr
ix1
time:
        2.03 secs
The first element is: 0.000000 and last element is: 499000500.000000
jinkai@jinkai-OMEN-by-HP-Laptop-15-ce0xx:~/Desktop/compsys304_A3/task_2$ ./matr
ix2
time:
       1.29 secs
The first element is: 0.000000 and last element is: 499000500.000000
jinkai@jinkai-OMEN-by-HP-Laptop-15-ce0xx:~/Desktop/compsys304_A3/task_2$ ./matr
ix3
time:
       0.81 secs
The first element is: 0.000000 and last element is: 499000500.000000
jinkai@jinkai-OMEN-by-HP-Laptop-15-ce0xx:~/Desktop/compsys304_A3/task_2$
```

```
jinkai@jinkai-OMEN-by-HP-Laptop-15-ce0xx: /proc
                                                                   Q
GNU_LIBC_VERSION
                                        glibc 2.31
GNU_LIBPTHREAD_VERSION
                                        NPTL 2.31
POSIX2_SYMLINKS
                                        1
LEVEL1_ICACHE_SIZE
                                        32768
LEVEL1_ICACHE_ASSOC
                                        R
LEVEL1_ICACHE_LINESIZE
                                        64
LEVEL1_DCACHE_SIZE
                                        32768
LEVEL1_DCACHE_ASSOC
                                        8
LEVEL1_DCACHE_LINESIZE
                                        64
LEVEL2_CACHE_SIZE
LEVEL2_CACHE_ASSOC
                                        262144
LEVEL2_CACHE_LINESIZE
LEVEL3_CACHE_SIZE
LEVEL3_CACHE_ASSOC
                                        64
                                        6291456
                                        12
LEVEL3 CACHE LINESIZE
                                       64
LEVEL4 CACHE SIZE
                                        0
LEVEL4 CACHE ASSOC
                                        Θ
LEVEL4 CACHE LINESIZE
                                        0
IPV6
                                        200809
RAW SOCKETS
                                        200809
POSIX IPV6
                                        200809
_POSIX_RAW_SOCKETS
                                        200809
jinkai@jinkai-OMEN-by-HP-Laptop-15-ce0xx:/proc$ ^C
jinkai@jinkai-OMEN-by-HP-Laptop-15-ce0xx:/proc$
```