# ECE482 — Introduction to Operating Systems

Homework 5

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Non-programming exercises:

- Write in a neat and legible handwriting
- Clearly explain the reasoning process
- Write in a complete style (subject, verb and object) Progamming exercises:
- Write a single README file per homework
- Push to git and create a release with tag h5

ECE4821: submit together with ECE4820

# **ECE4820 Exercises**

## **Ex. 1** — Simple questions

- 1. A system has two processes and three identical resources. Each process needs a maximum of two resources. Can a deadlock occur? Explain.
- 2. A computer has six tape drives, with n processes competing for them. Each process may need two drives. For which values of n is the system deadlock free?
- 3. A real-time system has four periodic events with periods of 50, 100, 200, and 250 msec each. Suppose the four events require 35, 20, 10, and x msec of CPU time, respectively. What is the largest value x for which the system is schedulable?
- 4. Round-robin schedulers normally maintain a list of all runnable processes, with each process occurring exactly once in the list. What would happen if a process occurred more than once in the list? Would there be any reason for allowing this?
- 5. Can a measure of whether a process is likely to be CPU bound or I/O bound be detected by analyzing the source code. How to determine it at runtime?

#### Ex. 2 — Deadlocks

Assuming three resources consider the following snapshot of a system.

Process	Allocated	Maximum	Available
$P_1$	010	753	332
$P_2$	200	322	
$P_3$	302	902	
$P_4$	211	222	
P <sub>5</sub>	002	433	

- 1. Determine the content of the Request matrix.
- 2. Is the system in a safe state?
- 3. Can all the processes be completed without the system being in an unsafe state at any stage?

### **Ex. 3** — The reader-writer problem

In the *reader-writer problem*, some data could be accessed for reading but also sometimes for writing. When processes want to read the data they get a *read lock* and a *write lock* for writing. Multiple processes

could get a read lock at the same time while a write lock should prevent anybody else from reading or writing the data until the write lock is released.

To solve the problem we decide to use a global variable count together with two semaphores: count\_lock for locking the count variable, and db\_lock for locking the database. To get a write lock we can proceed as follows:

- 1. Explain how to get a read lock, and write the corresponding pseudocode.
- 2. Describe what is happening if many readers request a lock.

To overcome the previous problem we will block any new reader when a writer becomes available.

- 3. Explain how to implement this idea using another semaphore called read\_lock.
- 4. Is this solution giving any unfair priority to the writer or the reader? Can the problem be considered as solved?

# **ECE4821** Exercise

**Ex. 4** — *Programming* 

Implement the Banker's algorithm.