

Lecture 51(Heap)

By help of Recursion

↳ gives segmentation fault

Space Complexity : $O(N)$

Stack :

↳ solved and run properly

Space Complexity : $O(N)$

Segmentation fault : We want to access that memory

which does not exist

Recursion works in Stack Memory

Heap is space

Jyada hota hai.

Heap

Stack

Stack works in heap. In heap, only STL library Command works.

Heap → Max Heap

↳

Min Heap

Priority Queue.

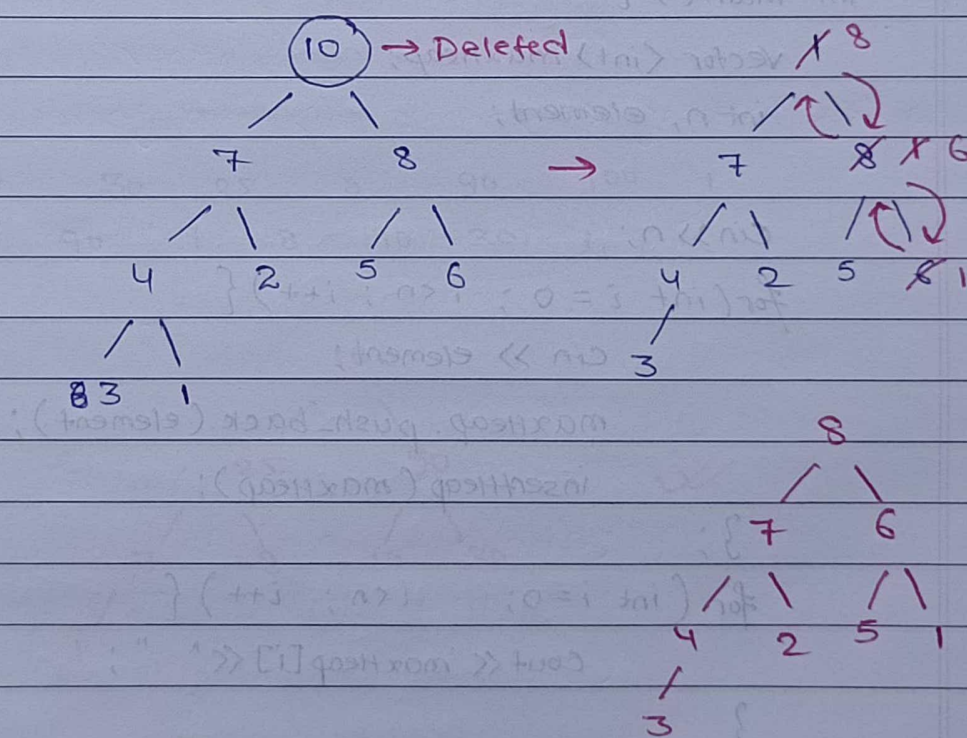
In Max Heap, $\text{parent Node} \geq \text{child Node}$

Max Heap & Min Heap can be represented in tree and Array both.

In Heap, creation always starts from top.

In Heap, Deletion always starts from top.

After Deleting the top Node, replace it from the last Node and then swap accordingly.



For creation, We have to use level order traversal in tree.

In Array, We easily access the nodes and know about the blank space.

$$\text{parent} = (i-1)/2$$

$$i = \text{child.idx} = \text{parent.idx} * 2 + 1$$

By tree, to find the last Node takes time.

But in array, we know about the last space.

Code :

```
#include <iostream>
```

```
#include <vector>
```

```
using namespace std;
```

```
int main() {
```

```
    vector<int> maxHeap;
```

```
    int n, element;
```

```
    cin >> n;
```

```
    for (int i = 0; i < n; i++) {
```

```
        cin >> element;
```

```
        maxHeap.push_back(element);
```

```
        insertHeap(maxHeap);
```

```
    };
```

```
    for (int i = 0; i < n; i++) {
```

```
        cout << maxHeap[i] << " ";
```

```
    }
```

```
    return 0;
```

```
}
```

```
void insertHeap(vector<int> &maxHeap) {
```

```
    int index = maxHeap.size() - 1;
```

```
    int parent;
```

```
    while (index > 0) {
```

```
        parent = (index - 1) / 2;
```

```
        // parent is small
```



```

if (maxHeap[parent] < maxHeap[index]) {
    swap(maxHeap[parent], maxHeap[index]);

```

```

    index = parent;

```

```

}

```

```

// parent is big

```

```

else

```

```

    break;

```

```

}

```

```

}

```

output :

8

10

7

30

25

8

90

100

1

100

25

90

7

8

10

30

1

100

25

90

7

8

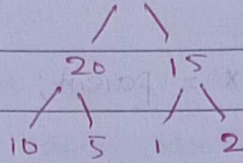
10

30

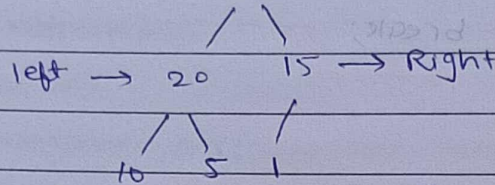
Time Complexity : $N \log N$

$N \rightarrow$ N element creation

$\log N \rightarrow$ Every Element Swap to the Max of tree height.

* DeletionDelete

2 → index



Heapify (maxheap, largest) {

largest = index;

left = 2 * index + 1;

right = 2 * index + 2;

if (maxHeap[left] > maxHeap[largest])

largest = left;

if (maxHeap[right] > maxHeap[largest])

largest = right;

if (largest != index) {

Swap (maxHeap[largest], maxHeap[index]);

Heapify (maxHeap, largest);

}

DeleteHeap(maxHeap); // good function call

```

void DeleteHeap (vector<int> &maxHeap) {
    // Replace first element with last element
    maxHeap[0] = maxHeap[maxHeap.size()-1];
    // Delete the last element
    maxHeap.pop_back();
    // Correct position
    Heapify(maxHeap, 0);
}

```

```

void Heapify (vector<int> &maxHeap, int index) {
    int largest = index;
    int left = 2 * index + 1;
    int right = 2 * index + 2;
    int size = maxHeap.size();
    // check left side
    if (left < size && maxHeap[left] > maxHeap[largest])
        largest = left;
    // check for right side
    if (right < size && maxHeap[right] > maxHeap[largest])
        largest = right;
    // swap
    if (largest != index) {
        swap(maxHeap[largest], maxHeap[index]);
        Heapify(maxHeap, largest);
    }
    return;
}

```


Time Complexity : $O(\log N)$

Space Complexity : $O(\log N)$

→ Max Heap :

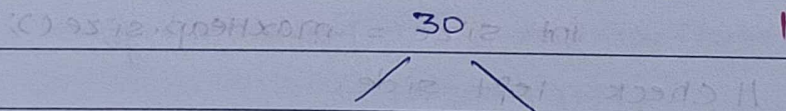
↳ Time Complexity : $N \log N$

Then we can't use sorting instead of using Max Heap.

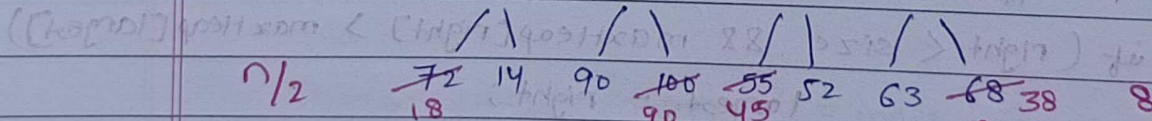
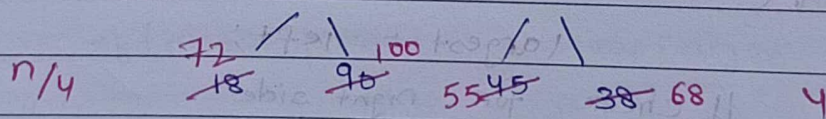
But, for Max Heap we can optimize $N \log N$

to N .

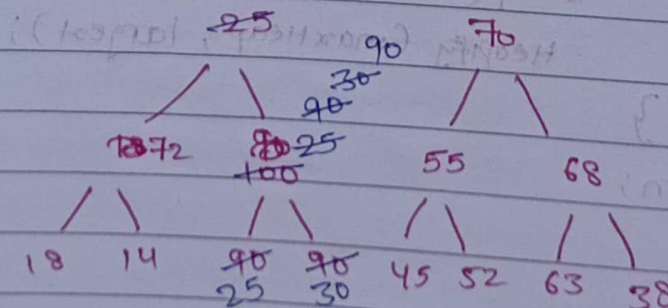
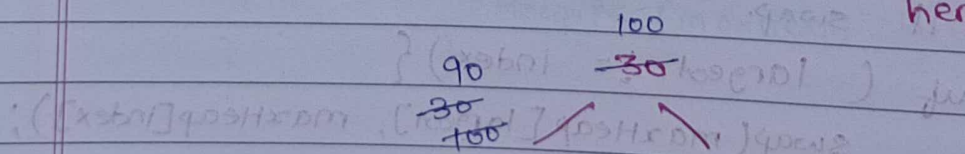
lets see how :-

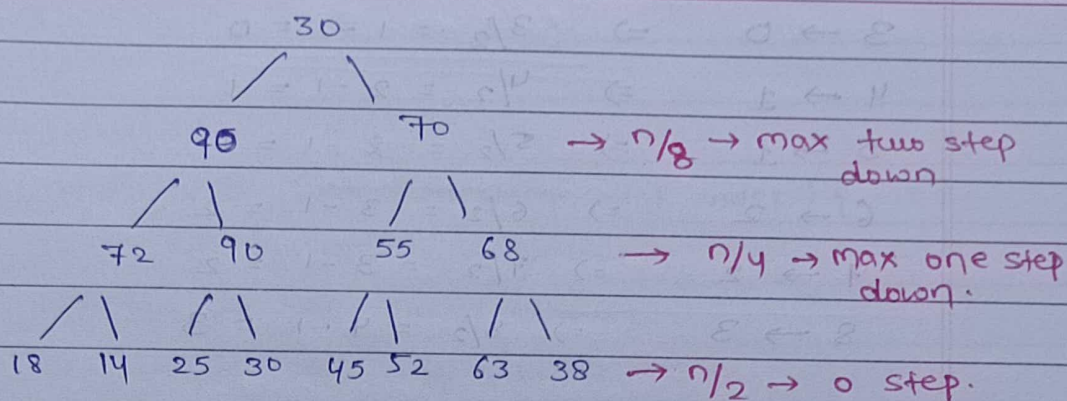


$n/8$ 25 70 4



Start from here





$$\left[\frac{n}{4} \times 1 + \frac{n}{8} \times 2 + \frac{n}{16} \times 3 + \frac{n}{32} \times 4 + \dots \right]$$

$$= n \left\{ \frac{1}{2^2} + \frac{2}{2^3} + \frac{3}{2^4} + \frac{4}{2^5} + \dots \right\}$$

\rightarrow sum nearly 1.

$$= O(n \times x)$$

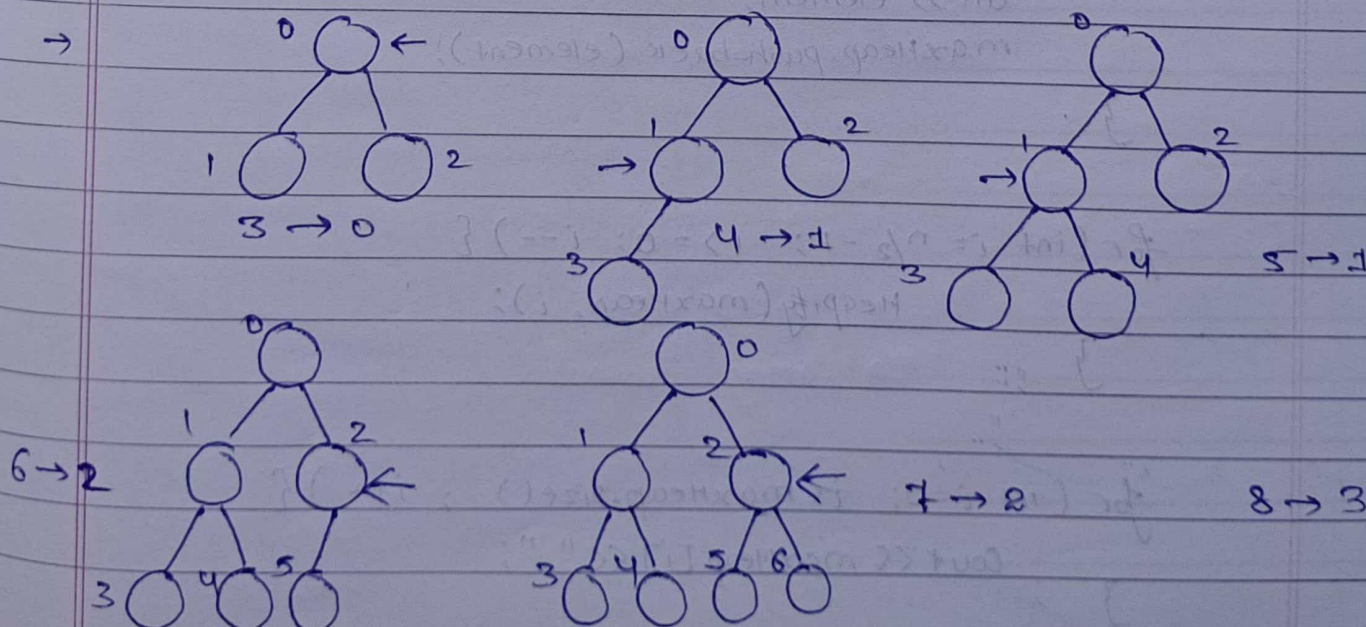
\rightarrow Constant

$$O(n)$$

\rightarrow Time Complexity

Space Complexity $O(\log n)$.

How to solve? How to find Starting Node.



$$3 \rightarrow 0 \Rightarrow 3/2 = 1 - 1 = 0$$

$$4 \rightarrow 1 \Rightarrow 4/2 = 2 - 1 = 1$$

$$5 \rightarrow 1 \Rightarrow 5/2 = 2 - 1 = 1$$

$$6 \rightarrow 2 \Rightarrow 6/2 = 3 - 1 = 2$$

$$7 \rightarrow 2 \Rightarrow 7/2 = 3 - 1 = 2$$

$$8 \rightarrow 3 \Rightarrow 8/2 = 4 - 1 = 3$$

→ `for (int i = size/2 - 1; i >= 0; i--) {
 Heapify(maxHeap, i);
 }`

* Code:

```
#include <iostream>
```

```
using namespace std;
```

```
int main() {
```

```
    vector<int> maxHeap;
```

```
    int n, element;
```

```
    cin >> n;
```

```
    for (int i = 0; i < n; i++) {
```

```
        cin >> element;
```

```
        maxHeap.push_back(element);
```

```
    }
```

```
    for (int i = n/2 - 1; i >= 0; i--) {
```

```
        Heapify(maxHeap, i);
```

```
    }
```

```
    for (int i = 0; i < maxHeap.size(); i++) {
```

```
        cout << maxHeap[i] << " ";
```

```
    }
```



```
return 0;
```

```
}
```

```
void Heapify (vector<int> &maxHeap, int index) {
```

```
    int largest = index;
```

```
    int left = 2 * index + 1;
```

```
    int right = 2 * index + 2;
```

```
    int size = maxHeap.size();
```

```
    if (left < size && maxHeap[left] > maxHeap[largest])
```

```
        largest = left;
```

```
    if (right < size && maxHeap[right] > maxHeap[largest])
```

```
        largest = right;
```

```
    if (largest != index) {
```

```
        Swap (maxHeap[largest], maxHeap[index]);
```

```
        Heapify (maxHeap, largest);
```

```
    }
```

```
    return;
```

```
}
```