

Lecture - 36

(Stack - Medium)

Problem 8.1 Get Minimum element from stack

You have N element

↪ find the min-m element while implementing a stack.

Required time Complexity : $O(1)$

11	13	7	1	11
11	13	7	1	11
7	13	7	1	11
13	7	1	1	11
4	13	7	1	11
2	4	13	7	1
			Ans = 1	

We check one by one all the element

and find min-m

but time Complexity is $O(N)$

and required is $O(1)$

This approach fails.

We can solve this by Another Approach.

Approach 2 :

2	4	13	7	1	11
---	---	----	---	---	----

Make 2 stack

Push(11)	11			1	
Push(1)	1			1	
push(7)	7			2	
push(13)	13			2	
push(4)	4			2	
push(2)	2	→ Min-m		2	

↳ This stack is
Strong for Min-m

We will push all the
element in stack in
Constant time.

We will return the top element from
2nd stack which is minimum.

But here we can take extra space for
second stack.

Space Complexity : O(N)
Not Required →

We want time complexity and space complexity
both as O(1).

Approach 3:

$$\frac{\text{freq} * \text{size}}{\text{size}} + \text{original_value} = \text{value}$$

frequency = divide

original value = Mod

0	1	2	3	4
2	1 6	12 2 7	13 3 8	3

↓
Go to
2nd
index

1+5 2+5

↓

$$7 \% 5 - 2$$

↓

Original
value

$$+ 5$$

Explanation

Limit of Array element \rightarrow (0 to 4)

We take a higher value

Let we take 5.

0	1	2	3	4
2	1	2 7	3	3

↑
Go to index 2

$$\text{Element} = 2 \% 5 = 2$$

Go to index 2

Add 5 to that index

$$2 + 5 = 7$$

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0	1	2	3	4
2	* 6	* 7	3	3

↑

Element = $1 \frac{1}{5}$ ($\leftarrow 1$ Go to index 1)

Add 5

$$1 + 5 = 6$$

0	1	2	3	4
2	6	* 12	3	3

↑

original Element = $7 \frac{1}{5} = 2$ ($\leftarrow 2$ Go to index 2)

Add 5

$$7 + 5 = 12$$

0	1	2	3	4
2	6	12	* 8	3

↑

original Element = $3 \frac{3}{5}$ ($\leftarrow 3$ Go to index 3)

Add 5

$$8 + 5 = 13$$

0	1	2	3	4
2	6	12	* 13	3

↑

Original Element = $3 \frac{3}{5} = 3$

Go to index 3

Add 5

$$8 + 5 = 13$$

For Finding Frequency, divide all element by 5

Frequency =

0	1	2	2	0
---	---	---	---	---

$$\begin{array}{lll} 2/5 = 0 & 6/5 = 1 & 12/5 = 2 \\ 13/5 = 2 & 3/5 = 0 \end{array}$$

* Same Concept we use in stack.

Range of Number is 1 to 100.

Take a number bigger than Range.

$$\text{Let } n = 101$$

originalvalue + minElement * 101
 ↓
 for this = mod for this = divide.

$$1 \leq \text{num} \leq 100$$

Why we take 101, why not 102, ...;

↳ because when we mod we get

0 to 100 which is limit of num.

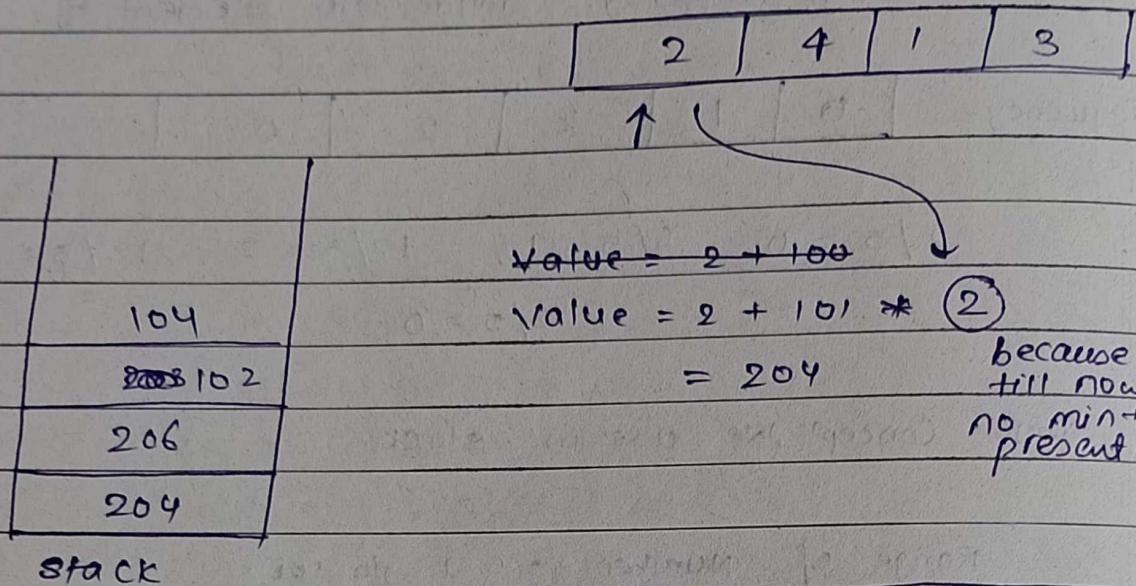
$$[\text{Original value} + \text{minElement} * 101] = \text{value}$$

↓

put in stack

$$\text{Value \% 100} = \text{original_value}$$

$$\text{Value / 101} = \text{min_element}.$$



$$\text{Value} = 2 + 104$$

$$\text{Value} = 2 + 101 * \textcircled{2}$$

$$= 204$$

because
till now
no min-m
present.

2	4	1	3
---	---	---	---

$$\min_m \text{ value} = 204$$

$$(2, 4) \rightarrow \min = 2$$

$$101$$

$$\text{Val} = 4 + 101 * 2$$

$$= \textcircled{2}$$

$$= 206$$

2	4	1	3
---	---	---	---

min-m element

$$= \frac{206}{101}$$

$$(2, 1) \rightarrow \min = 1$$

$$\text{Val} = 1 + 101 * \textcircled{2}$$

$$= \textcircled{2}$$

$$= \cancel{205} 102$$

2	4	1	3
---	---	---	---

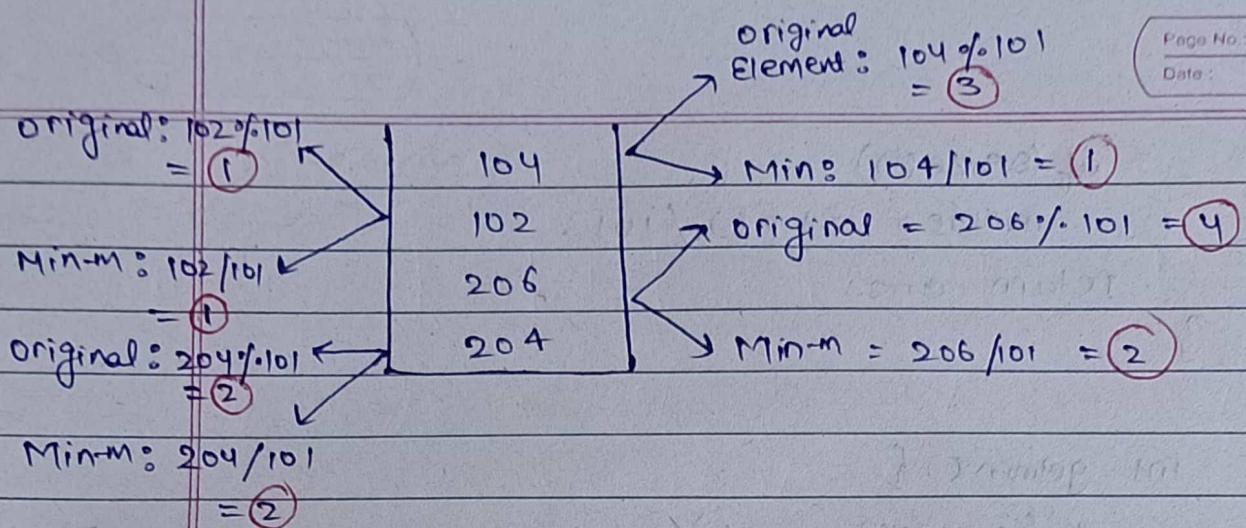
$$\min = \frac{102}{101}$$

$$(1, 3) \rightarrow \min = 1$$

$$= 1$$

$$\text{Val} = 3 + 101 * 1$$

$$= 104$$



By this we can get both by one stack.

→ original element

→ Min-M Number.

* Implementation of push:

```
void push (int x) {
    if (s.size() == 0) {
        minEle = x;
        s.push (x + (101 * x));
    } else {
        minEle = min (x, s.top() / 101);
        s.push (x + (101 * minEle));
    }
}
```

* Implementation of pop:

```
int pop() {
    if (s.size() == 0)
        return -1;
    int ans = s.top();
    ans = ans % 101;
    s.pop();
}
```

```

if (s.size())
    minEle = s.top() / 10;
return ans;
}

int getmin() {
    if (s.size() == 0)
        return -1;
    else
        return s.top() / 10;
}

```

Problem 2: String Manipulation

Sequence of n words.

If two same word come together
then we destroy both.

After all destruction happen,
how many numbers of words left.

$V[] = \{ ab, aa, aa, bcd, ab \}$

J_i destroyed

Output = 3

*

$v = [tom, jerry, jerry, tom]$

\downarrow
destroy

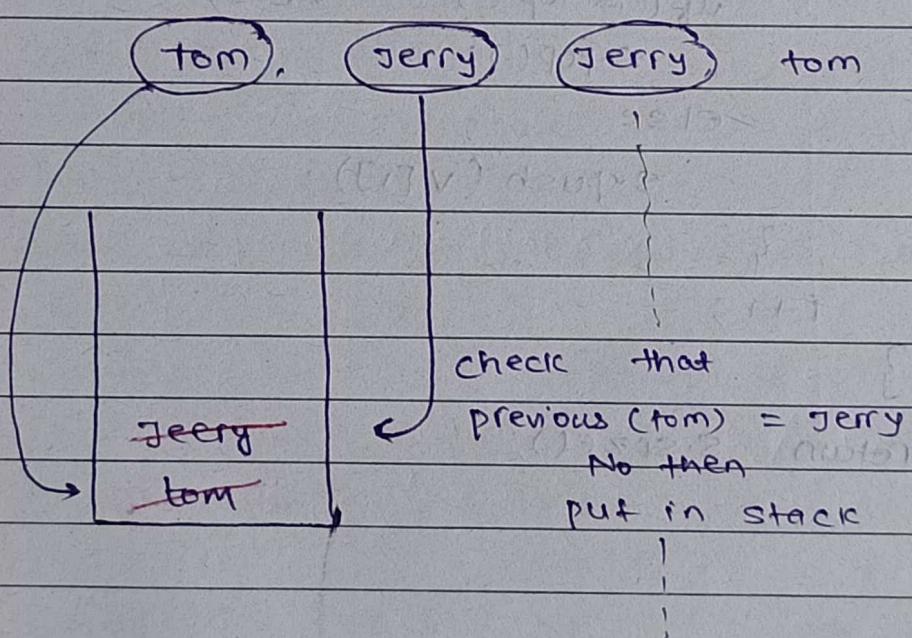
$v = [tom, tom]$

\downarrow
destroy

$v = []$

output = 0

* How to solve?



Again jerry

Check stock.top() == jerry
yes

then delete jerry
pop (top)

Again tom, check

stack.top() == tom

yes
pop (top).

then, Count the number of element present
in stack and return.

Code :

```
Stack <string> s;
int i = 0;
while (i < v.size()) {
    if (s.size() == 0)
        s.push(v[i]);
    else {
        if (s.top() == v[i])
            s.pop();
        else
            s.push(v[i]);
    }
    i++;
}
return s.size();
```

Problem : 3 Make the array beautiful.

Given a Array \rightarrow with negative & non-negative integers.

Condition for beautiful ?

two adjacent integer,

$arr[i]$ and $arr[i+1]$ are either negative or non-negative.

You have to do some operation until your array becomes beautiful :

If two adjacent integer are different i.e. one of them is negative and other is non-negative, remove them.

Empty array is also a beautiful array.

Example

Input : 4 destroyed

Output : 4

(2 -2)

Output : 4 1

Example :

Input :

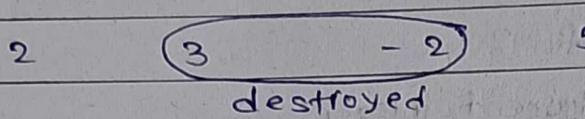
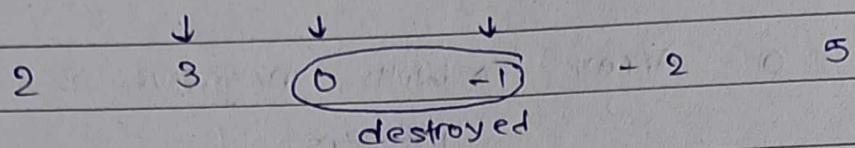
(2 -2)

destroyed

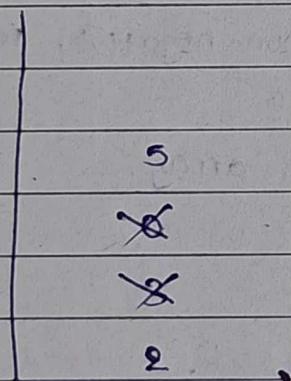
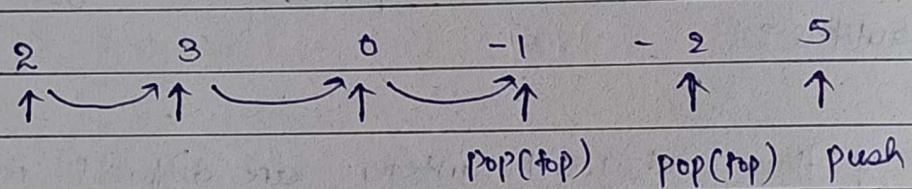
(1 -1)

destroyed.

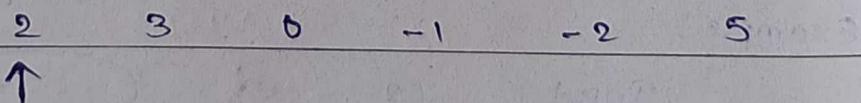
By stack :



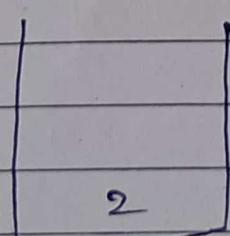
$[2, 5] \rightarrow \text{answer.}$



detailed explanation :



2 : (stack empty)
Push (2)



2 3 0 -1 -2 5 8.3.03.



+ve

3	
2	

top = 2 (+ve)

push(3)

3 (()) (empty) 3 } 5

(empty) deque

2 3 0 -1 -2 5



0	
3	
2	

0 +ve

top = 3 (+ve)

push

2 3 0 -1 -2 5

(empty) about ↑

X	
3	
2	

-1 -ve

top = 0 +ve

pop(top)

2 3 0 -1 -2 5

(empty) ↑

X	
3	
2	

-2 (-ve)

top = 3 +ve

pop(top)

2 3 0 -1 -2 5



5	
2	

5 (+ve)

top = 2 +ve

push

Now, return the element which
are present in stack.

Code :

```

Vector <int> makeBeautiful (vector <int> arr) {
    stack <int> s ;
    for (int i=0; i<arr.size(); i++) {
        if (s.empty()) {
            s.push(arr[i]);
        }
        else {
            if ((s.top() >= 0 && arr[i] < 0) || 
                (s.top() < 0 && arr[i] >= 0)) {
                s.pop();
            }
            else {
                s.push(arr[i]);
            }
        }
    }
}

```

```

vector <int> beautifulArray;
while (! s.empty()) {
    beautifulArray.push_back(s.top());
    s.pop();
}

```

```
reverse (beautifulArray.begin(), beautifulArray.end());
```

```
return beautifulArray;
```

Problem - 4

* Next Greater Element :

Given an Array :

return an Array with next greater element.

Ex :-

$$\text{arr}[] = [1, 3, 2, 4]$$

$$\text{ans} = [$$

~~1, 3, 2, 4~~

for ①

for ① → Next greater = 3

$$\text{ans} = \{ 3, \text{ } \} ,$$

for ② → Next greater = 4

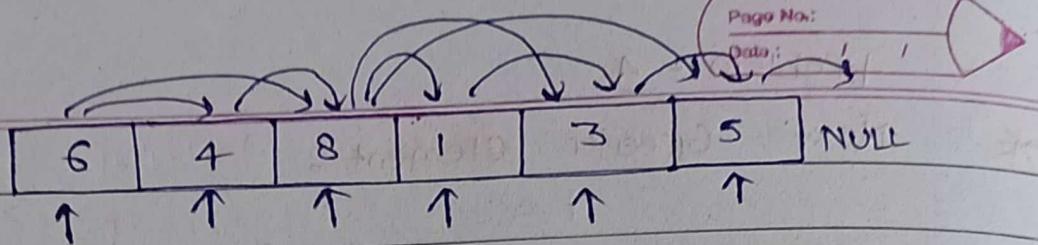
$$\text{ans} = \{ 3, 4, \text{ } \}$$

for ③ → Next greater = 4

$$\text{ans} = \{ 3, 4, 4, \text{ } \}$$

for ④ → Next greater = Not present (-1)

$$\text{ans} = \{ 3, 4, 4, -1 \} .$$



Output : | 8 | 8 | -1 | 3 | 5 | -1 |

* |

6	4	8	1	3	5
8	8	-1	3	5	-1

5	6 → put 1 st element
*	4 → 4 < 6 put in stack
*	8 → 8 > 4 , 8 > 6
8	1 → 1 < 8 put in stack
*	3 → 3 > 1
*	5 → 5 > 3

* |

6	4	8	3	1	2	5
↑	↑	↑	↑	↑	↑	↑
Next greater 8	8	-1	5	2	5	-1

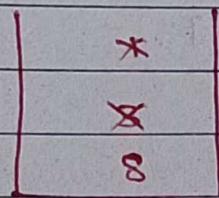
6 → put in stack

*	
*	

4 → 4 < 6 (then put in stack)

8 → 8 > 4 (then
for 4 '8' is next greater
pop(4))

Teacher's Signature.....



For 6 '8' is next greater

`pop (top);`

$3 \rightarrow 3 < 8 \text{ push}$

$1 \rightarrow 1 < 3 \text{ push}$

$2 \rightarrow 2 > 1 \text{ for } 1$

2 is Next greater

$5 \rightarrow 5 > 2$

then for 2

(5 is greater.)

for root (-1) is next greater

Code:

```
stack<int> s;
int i = 0;
while (i < n) {
    if (s.size() == 0) // stack is empty
        s.push(i);
    else {
        if (arr[s.top()] >= arr[i])
            s.push(i);
        else {
```

```

while (s.size() && arr[s.top()] < arr[i]) {
    ans[s.top()] = arr[i];
    s.pop();
    s.push(i);
}
i++;
}

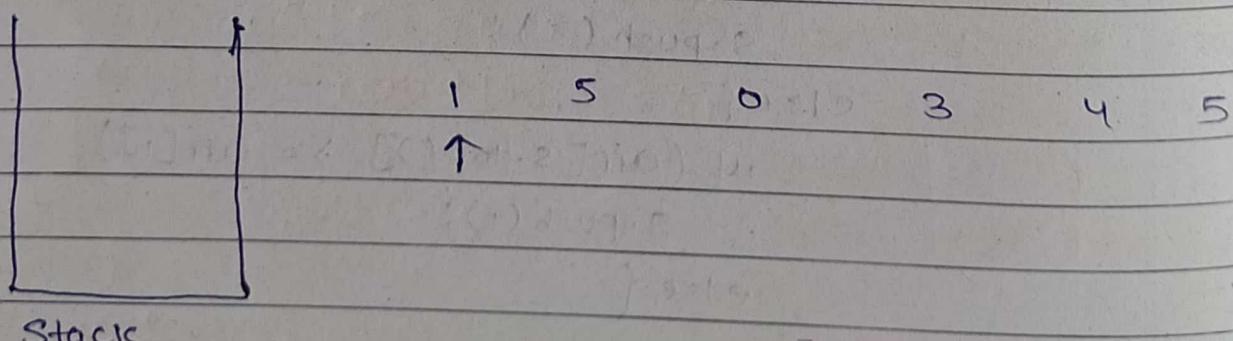
while (s.size()) {
    ans[s.top()] = -1;
    s.pop();
}
return ans;
}
}

```

* Problem - 5 : Smallest number on left :

$$a = [1, 5, 0, 3, 4, 5]$$

Output : -1 1 -1 0 3 4



Stack

Teacher's Signature.....

① → Stack is empty
then store -1

5
1

$$\text{Output} = [-1, \textcircled{1}]$$

1 5 0 3 4 5
↑

⑤ → Stack is not empty.
if smaller than arr[i] found
then return.
 1 is smaller. and then

$$\text{Output} = [-1, 1]$$

1 5 0 3 4 5

⑥ → Stack is not empty

Check, stack if smaller than arr[i] then return.

otherwise, return -1

$$5 > 0$$

$$1 > 0$$

return -1

$$\text{Output} = [-1, 1, -1]$$

③ stack is not empty

check smaller in stack

found (0) return 0.

$$\text{Output} = [-1, 1, -1, \textcircled{0}]$$

3
0
5
1

4 ④ stack is not empty
 3 check smaller than 4
 0 return 3
 5
 1

Output = [-1, 1, -1, 0, 3,

5 ⑤ stack is not empty
 check smaller than 5
 return 4

Output = [-1, 1, -1, 0, 3, 4]

Loop Ends

Return output.

Code :

```
Vector<int> leftsmaller (int n, int a[]) {
    Vector<int> result (n, -1);
    Stack<int> s;
    for (int i=0; i<n; i++) {
        while (!s.empty() && s.top() >= a[i]) {
            s.pop();
        }
        if (!s.empty()) {
            result[i] = s.top();
        }
        s.push(a[i]);
    }
    return result;
}
```