

Lecture - 2.6

Height of some students are given:

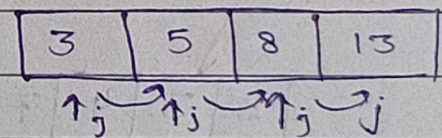
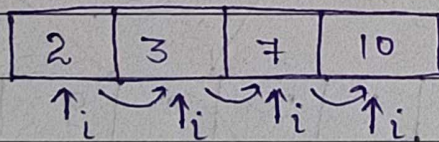
4 3 7 1 2 8 9

We have to arrange the student in ascending order by sorting.

Previously, we learn Bubble Sort, Insertion Sort, Selection Sort.

Now, we can sort this by a new sorting method called Merge Sort.

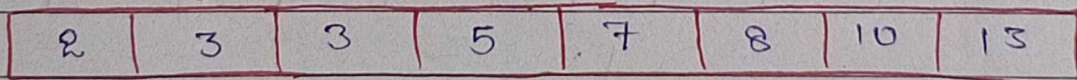
* Merge Sort :



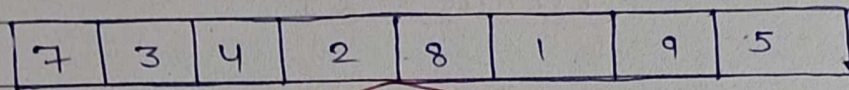
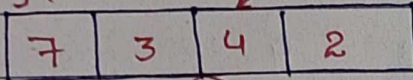
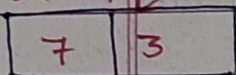
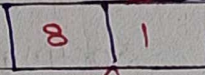
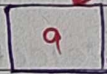
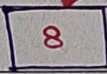
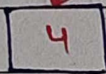
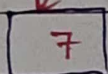
$$2 < 3, \quad 3 < 5$$

$$3 < 7$$

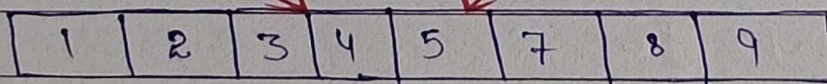
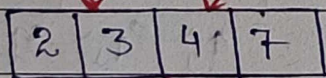
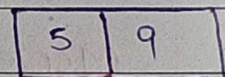
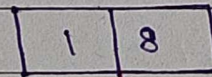
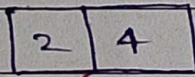
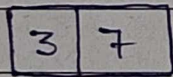
Depth first search



- Divide the array in two parts of equal size.
- Divide it until at the last the size of array becomes 1.
- Now, start Merging the single size array in sorted order and then merge all upper step.

 $f(0, n/2)$  $f(n/2+1, n)$  $f(0, n/4)$  $f(n/4+1, n/2)$  $f(n/2+1, 3n/4)$  $f(3n/4+1, n)$ 

Now, it can reverse the process with Comparing and sorting.



↳ This is sorted
by Merge sort.


```

Merge-sort (intarr, intlow, inthigh) {

```

```

    if (high == low)

```

```

        return ;

```

```

    mid = (low + high) / 2 ;

```

```

    Merge-sort (arr, low, mid);

```

```

    Merge-sort (arr, mid+1, end);

```

```

    Merge-array (arr, low, mid, high);

```

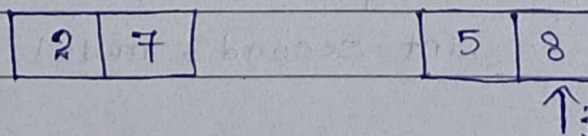
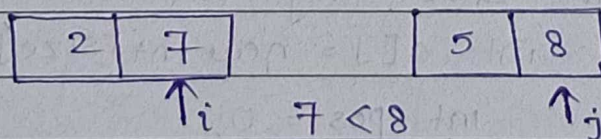
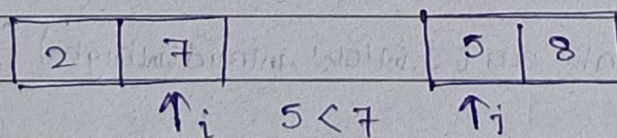
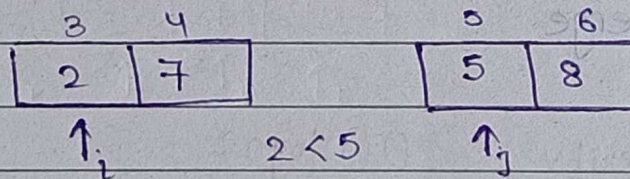
```

}

```

Merge-sort function can divide the array until the size of array becomes 1.

Merge-array can merge the after dividing with size = 1 and merge the element in sorted order.



2 5 7 8

↳ Final answer after sorting

Size :

$$n \rightarrow n/2^0$$

↓

$$n/2 \rightarrow n/2^1$$

↓

$$n/4 \rightarrow n/2^2$$

↓

↓

↓

$$n/2^k = 1$$

$$\frac{n}{2^k} = 1$$

$$n = 2^k$$

$$k = \log_2 n$$

→ Merge sort can follow divide and conquer method.

n swap in each height.

$$\text{height} = \log n.$$

$$\text{Total Time} = n \log n.$$

for Merge

$$\Rightarrow \text{Time Complexity} : O(n \log n)$$

Merge_array (int arr[], int low, int mid, int high) {

int size = high - low + 1;

int a[] = new int[size];

int pos = 0;

int first = low;

int second = mid + 1;

while (first <= mid && second <= high) {

Teacher's Signature


```
if (arr[first] <= arr[second]) {
```

```
    a[pos] = arr[first];
```

```
    pos++;
```

```
    first++;
```

```
}
```

```
else {
```

```
    a[pos] = arr[second];
```

```
    pos++;
```

```
    second++;
```

```
}
```

```
}
```

```
while (first <= mid) {
```

```
    a[pos] = arr[first];
```

```
    pos++;
```

```
    first++;
```

```
}
```

```
while (second <= high) {
```

```
    a[pos] = arr[second];
```

```
    pos++;
```

```
    second++;
```

```
}
```

```
for (int i = 0; i < arr.length; i++) {
```

```
    arr[i] = a[i];
```

```
}
```

```
pos = 0, mid = low;
```

```
while (mid <= high)
```

```
    arr[mid++] = a[pos++];
```

```
}
```


Space Complexity:

At every step we make array
 $\rightarrow O(n)$

Stack size: $\log n$

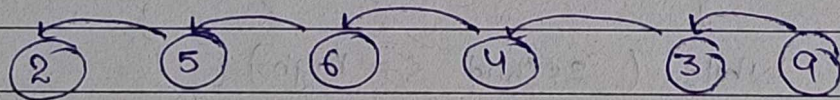
$(n + \log n)$

* Quick Sort

0	1	2	3	4	5
2	5	6	4	3	9

Choose one element as pivot element

Put the pivot element at correct place.



2 3 4 5 6 9

\rightarrow Now the array is sorted.

Pivot Element: Choose any element as pivot element. But we can organise in correct way.

So, we choose last element as pivot element.

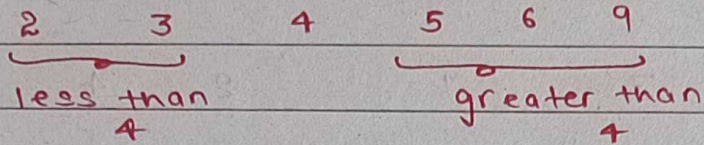
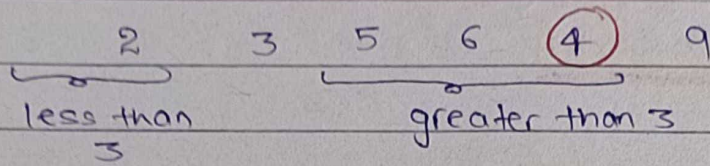
Example-1

2 5 6 4 3 9

2 5 6 4 3 9

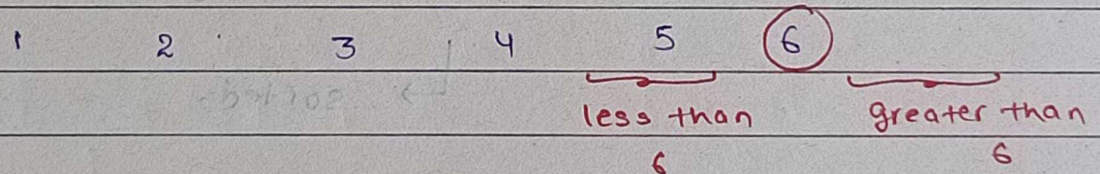
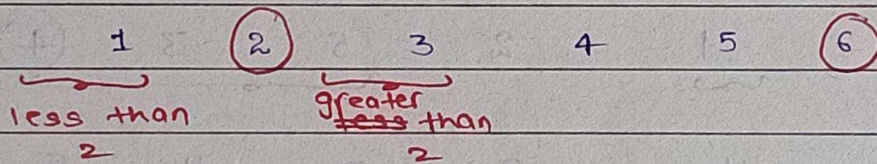
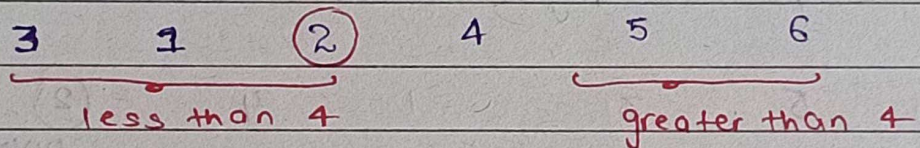
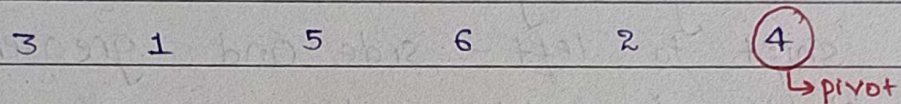
↑
pivot

Teacher's Signature.....



Now, Array is Sorted.

Example - 2



Now, Array is Sorted.