

```
for (int i = Count_zero + Count_one; i < n; i++) {
    arr[i] = 2;
}
```

(Lecture - 10)

(Sorting : Basic)

* Selection Sort :

arr : { 8 , 4 , 13 , 7 , 5 }

Output : 4 5 7 8 13

8 | 4 13 7 5
 ↓ min
 Choose minimum from this side

Swap this With 1st element

4 | 8 13 7 5
 ↑ minimum

choose min

Swap With 2nd element

4 5 | 13 7 8
 ↑ min

Choose Min

Swap With 3rd element

4 5 7 | 13 8

Teacher's Signature ↑ min

chose minimum
swap with 4th element.

4 5 7 8 13

↳ final result.

code?

```
for (int i = 0; i < n; i++) {
    int index = i;
    for (int j = i; j < n-1; j++) {
        if (arr[index] > arr[j]) {
            index = j;
        }
    }
    int temp = arr[index];
    arr[index] = arr[i];
    arr[i] = temp;
}
```

*

Bubble sort

arr[]: { 7 2 4 3 6 5 }

In this we compare two value and swap it
In every pass one element goes at its right place.
The last element goes at right place

We have to do (n-1) pass to complete this.

pass 1:

(7) (2) 4 3 6 5

$7 > 2$ (swap)

2 (7) (4) 3 6 5

$7 > 4$ (swap)

2 4 (7) (3) 6 5

$7 > 3$ (swap)

2 4 3 (7) (6) 5

$7 > 6$ (swap)

2 4 3 6 (7) (5)

$7 > 5$ (swap)

2 4 3 6 5 | 7

→ got its right place.

pass 2:

(2) (4) 3 6 5 | 7

$2 < 4$ (No swap)

2 (4) (3) 6 5 | 7

$4 > 3$ (swap)

2 3 (4) (6) 5 | 7

$4 < 6$ (No swap)

2 3 4 (6) (5) | 7

$6 > 5$ (swap)

2 3 4 5 | 6 7

pass 3:

(2) (3) 4 5 | 6 7

$2 < 3$ (No swap)

2 (3) (4) 5 | 6 7

$3 < 4$ (No swap)

2 3 (4) (5) | 6 7

$4 < 5$ (No swap)

pass 4 :

(2) (3) 4 | 5 6 7

$2 < 3$ (No swap)

2 (3) (4) | 5 6 7

$3 < 4$ (No swap)

pass 5 :

(2) (3) | 4 5 6 7

$2 < 3$ (No swap)

2 3 4 5 6 7

↳ sorted

code :

```
for (int i = 0; i < n - 1; i++) {
    for (int j = 0; j < (n - 1) - i; j++) {
        if (arr[j] > arr[j + 1]) {
            int temp = arr[j];
            arr[j] = arr[j + 1];
            arr[j + 1] = temp;
        }
    }
}
```


* Insertion Sort :

arr[] : 4 6 3 11 7 2

sorted

(1-0)

4 | 6 3 11 7 2
 ↑ place at correct place in left

4 6 | 3 11 7 2 (2-0)

3 4 6 | 11 7 2 (3-0)

3 4 6 11 | 7 2 (4-0)

3 4 6 7 11 | 2 (5-0)

2 3 4 6 7 11

Code : → sorted.

```

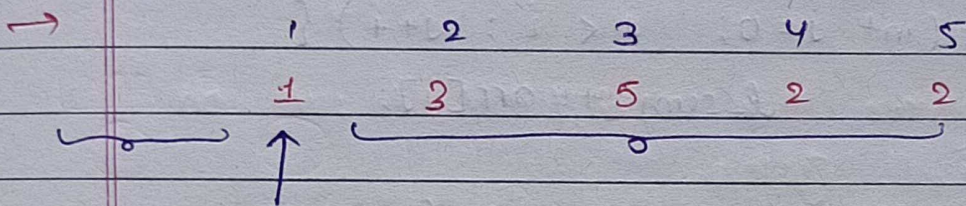
for (int i = 0 ; i < n-1 ; i++) {
    for (int j = i ; j >= 0 ; j--) {
        if (arr[j] > arr[j+1]) {
            int temp = arr[j];
            arr[j] = arr[j+1];
            arr[j+1] = temp;
        }
        else
            break;
    }
}

```


* Problem : Equilibrium point

The left side subarray and right side subarray of an element is equal, then the element is known as Equilibrium point:

Example:

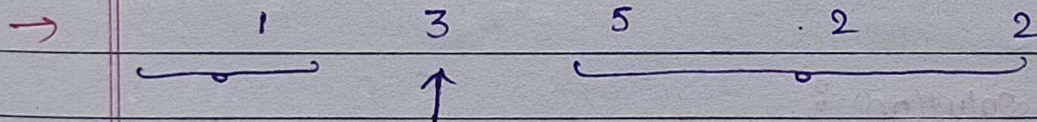


$$l_sum = 0$$

$$right = 3 + 5 + 2 + 2 = 12$$

$$l_sum \neq right$$

1 is not Equilibrium point

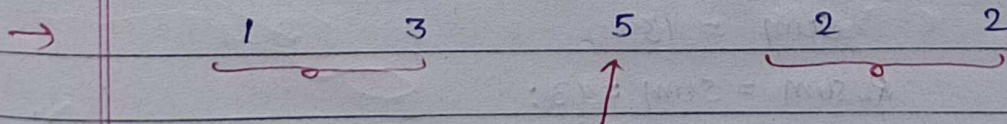


$$l_sum = 1$$

$$right = 5 + 2 + 2 = 9$$

$$l_sum \neq right$$

3 is not Equilibrium point



$$l_sum = 1 + 3 = 4$$

$$right = 2 + 2 = 4$$

$$l_sum = right$$

5 is the Equilibrium point.

Code: (Brute force) T.C: $O(N^2)$

```

for (int i = 0; i < n; i++) {
    int r-sum = 0;
    int l-sum = 0;
    for (int j = i + 1; j < n; j++) {
        r-sum += arr[j];
    }
    for (int j = 0; j < i; j++) {
        l-sum += arr[j];
    }
    if (l-sum == r-sum) {
        return i + 1;
    }
}
return -1;

```

Optimal solution :

1 3 5 2 2

calculate sum

sum = 13

R-sum = sum = 13;

→

↑

R-sum = 13 - 1 = 12
L-sum = 0

> Not Equal.

1 3 5 2 2
 ↑

$$R\text{-sum} = 12 - 3 = 9$$

$$L\text{-sum} = 1$$

> Not Equal

1 3 5 2 2
 ↑

$$R\text{-sum} = 9 - 5 = 4$$

$$L\text{-sum} = 4$$

> Equal

Equilibrium point.

Time complexity : $O(N)$

Code :

```

if (n == 1)
    return 1;

long totalsum = 0;
int res = -1;
long sum = 0;
for (int i = 0; i < n; i++)
    totalsum += arr[i];

for (int i = 1; i < n; i++) {
    sum += arr[i-1];

    if (totalsum - sum - arr[i] == sum) {
        res = i;
        break;
    }
}

return res == -1 ? -1 : res + 1;

```