1. Aggregate Analysis

we have a sequence of n-operations and, the cost function is defined as-

s defined as-
$$i, \text{ for } i = 2^{k} (k = 0, 1, 2, \dots)$$

$$C_{i} = \begin{cases} 1, \text{ otherwise} \end{cases}$$

Now

Aggregate cost of n-operations

= cost of operations for i=2k + cost of operations for i +2k

$$= \frac{\text{Cost of operations } \pi}{2^{j}} + \left[n - \left[\log n\right] + 1\right] \times 1$$

$$= \frac{2^{\lfloor \log_2 n \rfloor} \cdot 2^{-1}}{2^{-1}} + n - \lfloor \log_2 n \rfloor - 1$$

$$= 2n-1 + n - \lfloor \log_2 n \rfloor - 1$$

Thus, amortized cost per operation = # operations

.. By aggregate conclysis, the amortized cost for operation is O(2).

1. Amortized Analysis (Accounting Method)

the actual cost of ith operation is $G = \begin{cases} i, & \text{for } i = 2^{K} \text{ any } K = 0,1; \dots \\ 1, & \text{otherwise}. \end{cases}$

We define amortized cort as follows: $\hat{C}_i = \begin{cases} 2, & \text{for } i = 2^k \text{ any } k = 0,12, \dots \\ 3, & \text{otherwise.} \end{cases}$ Wi

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- for an operation, i), with $i=2^k$, it uses all previously accumulated credits plus one unit of its own amoratized cost to pay its true cost. It then assigns remaining one unit as credit.
 - for an operation, i, with i = 2k, it uses one unit to pay its actual cost and assign remaining two as credit.

This covers the actual cost of the operation Now, we need to show that there will be enough credit to pay the lost for 25 operations - for 1st operation (i=1) , it is clear.

- New Suppose 170 and after 25-th operation, there is one unit Beth operations 2^{j-1} and 2^j, there are 2^{j-1}-1 operations & of credit. all one not-power-of-2. So, total credit $= 1 + 2 \cdot (2^{j-1})$ $= 2^{j} - 1, \text{ before } 2^{j+1} \text{ operation}$
 - : This accumulated credit alongwith the one of its own is enough to cover its true cost. Therefore, the trad amortized upt per operation is an upper bound for the total actual cost.