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*Kruskal's Algorithm-

Kruskal's Algorithm finds a safe edge to add to the growing forest by finding, of all edges that connect any two trees in the forest an edge (u,v) of highest weight.

Let C_1, C_2 denote the two trees that are connected by (u,v) . Since (u,v) must be high weight edge connecting C_1 to some other tree,

MST-KRUSKAL(G, w)

$A = \emptyset$

For each vertex $v \in G.V$

MAKE-SET(v)

sort the edges of $G.E$ into decreasing order by weight w

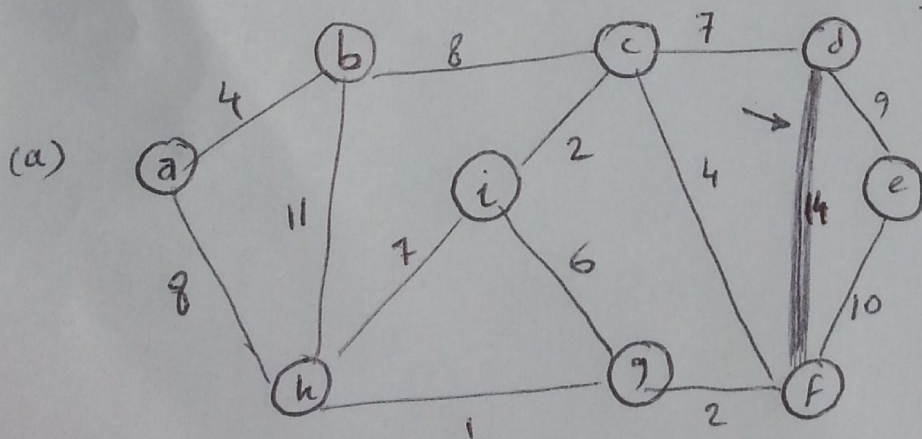
foreach edge $(u,v) \in G.E$, taken in decreasing order by weight

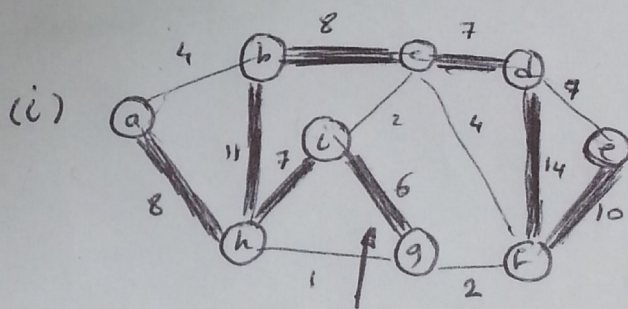
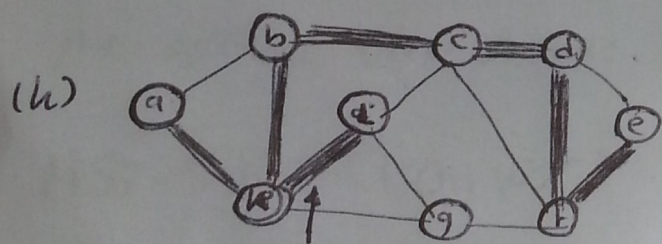
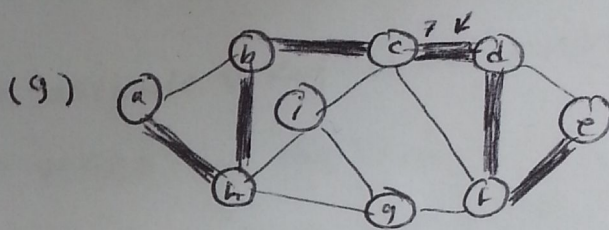
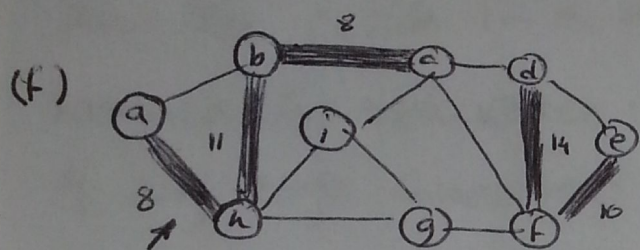
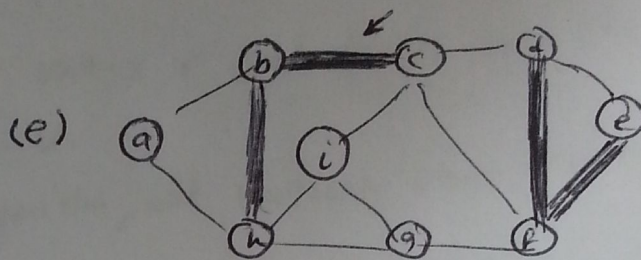
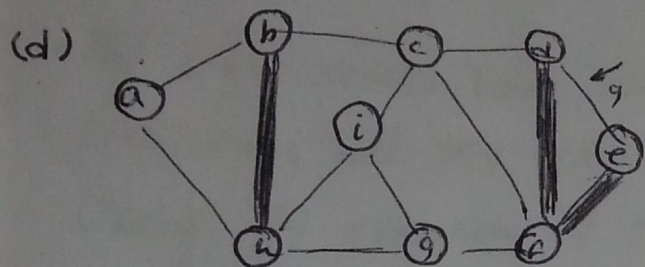
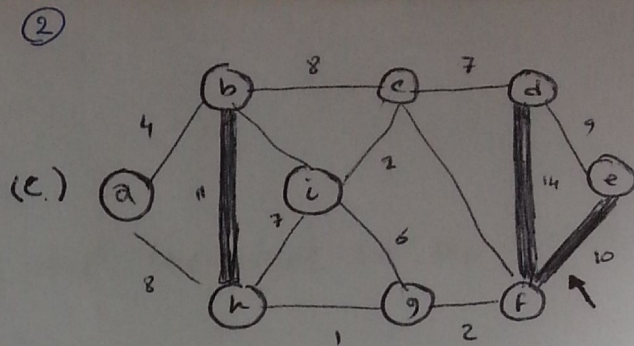
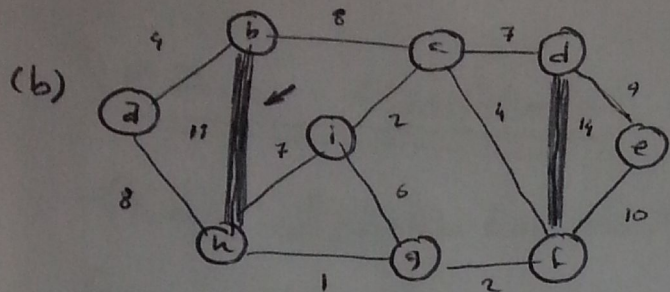
if FIND-SET(u) \neq FIND-SET(v)

$A = A \cup \{(u,v)\}$

UNION(u,v)

return A





* Prim's Algorithm :

- Prim's algorithm has the property that the edges in the set A always form a single tree.
- The tree starts from arbitrary vertex r and grows until the tree spans all the vertices in V .
- During the execution of the algorithm, all vertices that are not in the tree reside in a max-priority queue Q based on a key attribute. For each vertex v , the attribute $v.key$ is the maximum weight of any edge connecting v to a vertex in the tree.
- by convention $v.key = -\infty$ if there is no such edge.
- the attribute $v.\pi$ names the parent of v in the tree.

MST-PRIM(G, w, r)

for each $u \in G.V$

$u.key = -\infty$

$u.\pi = NIL$

$r.key = 0$

$Q = G.V$

while $Q \neq \emptyset$

$u = \text{Extract-Max}(Q)$

for each $v \in G.Adj[u]$

if $v \in Q$ and $w(u, v) > v.key$

$v.\pi = u$

$v.key = w(u, v)$

