Analyzation of Sorting Algorithms, Lab 5, Java II

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Abstract—Sorting has become a sub-field of computer science. The ability to take a set of data and arrange it into a specific sequence of ascending or descending values has created a multitude of algorithms and analyses. Our challenge was to create a class from which we could run at least 6 different sorting algorithms coded in Java and compare the performance of each on a given set of data.

 $\textbf{Keywords} \\ \textbf{--} Computer Society, IEEE tran, journal, \\ \textbf{--} ET_{E}X, Force Push, nanosecond unreliability.$

1 Introduction

OUR team created a class that contained 6 different sorting routines. We also created a separate class to handle the acquisition of the data to be sorted, the invoking of the Sort class to access the individual sort methods, and the recording of the time that it would take to perform each task. We also wanted to be able to provide different sizes of data sets to be able to compare sort times for varying values of n.

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1.1 Tasking the Sorts

Our task was not to code each sorting routine from scratch, although we did so with the Bubble sort and our own Force

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Push sort. Our task was, instead, to assemble the sorts from other documented sources such as the textbook (in the case of the Quick Sort) or the Internet as noted in each method. Once the sorting routines were gathered and assembled, a methodology had to be developed to provide fair and accurate testing and recording. Each sorting algorithm was structured to receive a reference to an array of type Point. The class Point was defined in a separate point.java class file and consisted of two double values representing the (x,y) Cartesian graphing pair. The array would be sorted in ascending order based on the x coordinate. No return type was required since the reference passed as an actual parameter was a reference to the array from the invoking class. Thus, the original array was the one being sorted. System time was recorded in nanoseconds before and after each sorting routine was called. The difference of the two times was logged as the time that it took that sort algorithm to perform a sort on the array with n elements.

1

1.2 Methodology

Our methodology for testing the different sorting algorithms was to create a sort

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testing class that would have an array of integers representing the different sizes of arrays (different size n arrays) to be sorted. This gave us great flexibility in deciding the different sizes of the data sets. Given a file with 100,000 test elements, this array of integers was used to create sub arrays from the full list.

Each sort was invoked given an array of sizes varying from 10 elements to 100,000 elements. Each sort was timed and that time was recorded. Because run time of a program can be influenced by other tasks the computer may be performing at the same time, we ran each test 10 times and took the average time for each test case.

We noticed that there were anomalies with the results when the test cases used

, regardless of the sort routine called. In fact, the first test of each routine resulted in times that were higher than the second run of the same routine even when the same data was being used. To counter this anomaly, we disregarded the results of the first few tests i.e.

$$n = 2, n = 3, n = 4, ... n = 9$$

1.3 Hypothesis

We were working on a hypothesis that an

$$O(n^2)$$

sort might have better real time performance on lower n elements that an

sort. Thus, we structured our test cases to cover more tests of

However, our tests did not show a significant benefit for using a non-recursive or

$$O(n^2)$$

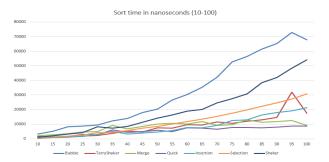
sort for lower values of n.

1.4 Graphs

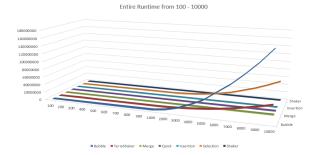
These were the averages of the test results for

$$n <= 100$$

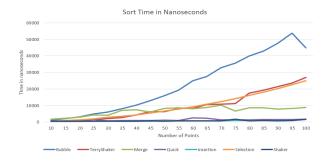
based on the random data from the original file:



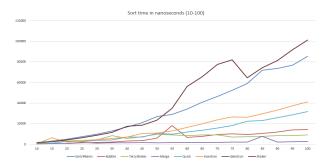
The overall graph of the 7 sort routines that we tested is as follows:



We also ran test cases for the same test sizes of n using data that was already sorted. This graph shows the results of this test for best-case scenario:



In addition, we ran the same tests for data that was in perfect reverse order. This graph shows the results of a worst-case scenario:



1.5 Run-time

For this sorting task analysis, we were required to include 5 well know sorts: Bubble, Merge, Insertion, Selection, and Quick. We were also required to select a sixth sorting algorithm of our choice. Our team chose the Cocktail shaker sort, which derives it's name from a visualization of the way the sort "bubbles" the largest item to the right, then the smallest item to the left, then the next largest item to its place on the right, etc... until all items are sorted. Our perception of the Cocktail Shaker sort was that it was an improvement over the Bubble sort because it correctly identified and placed the largest and smallest two items in their proper place on each pass. Further examination proved our initial perception to be incorrect. By definition, the Cocktail Shaker sort is indeed a Bubble sort that makes two passes: the first pass, from left to right, bubbles the largest item to the right until it is in it's proper place; the second pass, from right to left, bubbles the smallest item to the left. This alternation of directions can provide better performance in some cases. Our Bubble sort had a calculated run-time of

$$8n^2 + 5n + 2$$

and the Cocktail Shaker sort had a runtime of

$$8n^2 + 3n$$

While both the Bubble and Cocktail Shaker sorts have a Big-O of

$$n^2$$

the data shows that the Cocktail Shaker out-performs the Bubble sort. Given the misconception of the algorithm for the Cocktail Shaker sort, our team decided to try to implement what we thought was a commonsense improvement on the Bubble sort. We decided to call our sort the "Force Push" sort, given the way that each pass forces an item left and an item right. If a Jedi were standing in the middle of the area to be sorted, he would "Push" both items to either end and have them land in their appropriate location. So we created, from scratch, an algorithm to make a pass through the data and locate, but not swap, the largest and smallest items at the same time. At the end of the first iteration, we would swap the smallest element to the leftmost side, and swap the largest element to the rightmost side. The next iteration would then not need to include the first or last element when searching for the remaining minimum and maximum value and repeating the process. In our Force Push sort, the first iteration would take n passes, the second iteration would require

$$n-2$$

passes. Each iteration would decrease n by 2. Thus, the number of iterations would be

$$n+(n-2)+(n-4)+(n-6)...+(n-(n-2))$$

with only two element swaps per iteration. A Bubble sort would do n passes through the data on the first pass and n passes through the data on each successive iteration. Bubble sort can be improved to skip the area already sorted, thus decreasing the number of passes through the list of elements by 1 on each iteration.

This means that a better coded Bubble sort would take

$$n+(n-1)+(n-2)+(n-3)+...+(n-(n-1))$$

passes, with an average of n/2 swaps on the first iteration and

$$\frac{n-1}{2}$$

swaps on the second iteration, and so on. The Force Push sort is also in the Big-O of:

$$n^2$$

Our Force Push sort has a run-time formula of

$$8n^2 + 11n + 18$$

1.6 Conclusion

We found many different ideas and methods when analyzing different sort methods. Most of them are fascinating and many are ingenious. Data will always need to be sorted. We found that coding a simple sort routing can be quick and easy, as in the case of the Bubble sort. We did Bubble sort from memory and typed it straight into our class. We copied the Quick sort from the textbook and easily adapted it to our data type. Other sorts were copied straight from the Internet and adapted. But coding our Force Push sort required more testing and troubleshooting than anticipated. Our first attempt at coding the Force Push sort used recursion. This proved to be memory intensive and caused overflow errors. After the testing was done on the method's algorithm, we easily converted the implementation to an iterative approach. As can be seen from the data, we bested the Bubble sort, even though we were still in the family of

$$n^2$$

Java has a built-in sort, and we wanted to implement a TimSort from the Internet. But the ability to sort non-primitive data types using generics was beyond the scope of our knowledge. Implementing different sorts and having to adjust the coding to handle our Point data type for sorting whets the appetite to learn generics. The library of sorting methods that we have tested is already a valuable reference.

ACKNOWLEDGMENTS

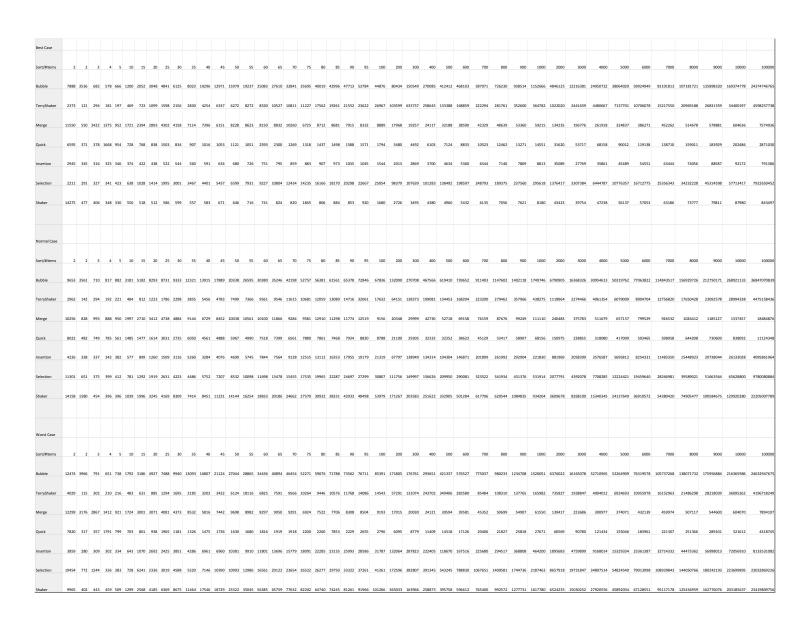
The authors of this document would like to thank Phd. C. Servin of EPCC and the LaTeX community for making documentation and presets readily available to us during this paper.

APPENDIX A CODE

At the end of this document we provide our Java source code for the review of those grading this paper and for those interested in running future tests with Superior hardware, give Moore's Law, despite it not being a law, holds true.

APPENDIX B DATA

On the next page we give the table of data for our project. The data was made by summing the time of iteration for each number of elements then dividing the sum by the number of the iterations thus getting a more precise measurement. This was necessary in light of the unreliability of java run-time at the scale of nanoseconds.



Source Code

by Jason Ivey & Terry Speicher using: LaTeX & Dr Java.

October 2016

1 Goal

To read the 100kpoint.txt file and sort the data within it providing reliable data on the run-time of each algorithm.

```
class Point {
        private double x;
private double y;
2
3
         public Point (double x, double y) {
4
5
           this.x = x;
6
           this.y = y;
8
         public double getX(){
9
          return this.x;
10
        public double getY(){
11
12
          return this.y;
13
14
    //Terry Speicher and Jason Ivey
4
     * Collection of sorting algorithms to be compared, along with a few utilitarian methods
       Qauthor Terry Speicher
5
       @author Jason Ivey
8
    public class SortAlgorithms {
9
10
      ^{/**}
* Constructor doesn't have to do anything
11
12
      public SortAlgorithms() {
13
14
      }
15
16
17
       * Not a complete comparison of all items to see if they are in order, but quickly checks
18
19
       * consecutive items to make sure that the smaller item is first.
       * @param data Point[] of some size to be processed

* @return boolean Return true if each consecutive item is less than the one after it.
20
21
22
      public boolean is Sorted (Point [] data) {
23
24
         boolean sortedYN = true;
25
        26
27
           if (data[i].getX() > data[i+1].getX())
    sortedYN = false;
28
29
           i++;
30
31
32
33
         return sortedYN;
34
         //end isSortedYN
35
36
37
       * Compare the first 100 items in each given Point array to see if they are equal. Used
38
39
       * as a testing routine to work out some bugs when learning how to pass the data by
40
       * reference.
       * @param data Point[] of some size to be processed * @param test Point[] of some size to be processed
41
42
```

```
* @return boolean Return true if the first 100 x coordinates are the same in the
 43
 44
         * two arravs
 45
 46
         public boolean looksEqual(Point [] data, Point [] test){
 47
           boolean lookSame = true;
for (int i = 0 ; i < 100 ; i++)
  if (data[i].getX() != test[i].getX())</pre>
 48
 49
 50
           lookSame = false;
return lookSame;
 51
 52
         } //end looksEqual
 53
 54
           ****** Begin BubbleSort **********
 55
 56
 57
         * BubbleSort was coded by the authors from scratch. This is the worst case sort because
 58
          st it does not even adjust the inner loop to start at the first nonsorted element – it
 59
 60
          * always starts at the beginning.
          * @param d Point[] of some size to be processed
 61
 62
         public void bubbleSort(Point [] d){
 63
           \begin{array}{lll} & \text{for}\,(\,\text{int}\ x = 0\,;\ x < \,d\,.\,\text{length}\,;\ x++)\{\\ & \text{for}\,(\,\text{int}\ y = 0\,;\ y < \,d\,.\,\text{length}\,-1;\ y++)\{\\ & \text{if}\,\,(\,d\,[y]\,.\,\text{getX}\,()\,>\,d\,[y+1]\,.\,\text{getX}\,()\,\,)\,\,\,\{ \end{array}
 64
 65
 66
                   Point temp = d[y];
 67
 68
                   d[y] = d[y+1];
                   d[y+1] = temp;
 69
 70
 71
             }
 72
          }
        }
 73
 74
           ******* End Bubble Sort ***********
      // ******* Begin QuickSort Methods ***********
 77
 78
 79
 80
         * Copied from the textbook. Improved on by adding a check in the partitioning
             routine to check "if (endOfLeftList != scan)" before swapping, because we found
 81
 82
            that the swap routine was sometimes being called to swap an element with itself if
          * the "endOfLeftList" equaled "scan".

* @param data Point[] of some size to be processed
 84
         public void quickSort (Point [] data) {
 86
 87
           doQuickSort (data, 0, data.length-1);
        } // end quickSort
 90
 91
 92
 93
         * sub method to start recursion process. Required to
         * pass in the two indexes for processing the sub array.
 94
 95
            Initially, we pass in the absolute beginning and
 96
            absolute ending elements.
         * @param array Point[] to be sorted
* @param start int that shows where the start of the subarray to be sorted is
 97
 98
          * ©param end int that shows where the end of the bubarray to be sorted is
 99
100
         private void doQuickSort(Point[] array, int start, int end){
101
           int pivotPoint;
102
           if (start < end) {
103
             pivotPoint = partition(array, start, end);
doQuickSort(array, start, pivotPoint -1);
doQuickSort(array, pivotPoint+1,end);
104
105
106
107
         } // end doQuickSort
108
109
110
         * Partition the given array from array[start] to array[end] and put * pivot element in middle. Then move all elements with values less * than the pivot point to the left of the pivot point and move
111
112
113
114
            all elements with values greater than the pivot point to the right
115
            of the pivot point
            @param array Point[] to be sorted
@param start int that shows where the start of the subarray to be partitioned is
116
117
118
          * @param end int that shows where the end of the bubarray to be partitioned is
119
         private int partition(Point[] array, int start, int end){
120
121
122
           double pivotValue;
123
           int endOfLeftList;
```

```
124
          int mid;
125
126
          mid = (start + end) / 2;
127
           //take the first element of the array and swap
128
           //it with the middle element. I don't know why, //except to keep with the idea of using the middle
129
130
           //element as the pivot
131
           swap(array, start, mid);
132
133
           pivotValue = array[start].getX();
134
           endOfLeftList = start;
for (int scan = start + 1; scan <= end; scan++){</pre>
135
136
             if (array[scan].getX() < pivotValue){
  endOfLeftList++;</pre>
137
138
                if (endOfLeftList != scan)
139
                  swap(array, endOfLeftList, scan);
140
141
          }//end of for
142
143
           swap(array, start, endOfLeftList);
144
145
           return endOfLeftList;
146
147
148
        }//end partition
149
150
         * Swap elements array[a] and array[b]
151
         * @param array Point[] with elements to be swapped
* @param a index of an element to be swapped
* @param b index of other element to be swapped
152
153
154
155
        private void swap(Point[] array, int a, int b){
156
157
158
           Point temp = array[a];
           array[a] = array[b];
array[b] = temp;
159
160
161
162
        }//end swap
163
      // ********** End QuickSort Methods
164
165
           ****** Begin MergeSort Method ************
167
168
169
                 With online help from:
                 Title: howotodoinjava.com
170
                 Author: Lokesh Gupta
171
                 Date: October 23, 2015
172
                 Availability: http://howtodoinjava.com/algorithm/merge-sort-java-example/
173
174
175
176
         * Merge Sort method to sort an array of data points
177
         * @param data Point[] of some size to be processed
178
179
180
        public void mergeSort(Point[] data){
181
182
           if (data.length <= 1){
183
            return;
184
185
           186
187
188
          \label{eq:System.arraycopy} System.arraycopy(data, 0, half1, 0, half1.length); \\ System.arraycopy(data, half1.length, half2, 0, half2.length); \\
189
190
191
           mergeSort(half1);
192
           mergeSort (half2);
193
194
           merge(half1, half2, data);
195
196
           return:
197
198
           * Main recursive method to merge sorting.

* divides Points array into smaller pieces.

* Makes use of merge(Point[], Point[], Point[]) to combine and sort data
199
200
201
202
203
204
```

```
* Merge elements of array[a] and array[b] into result
* @param half1 Point[] half to be combined with brother in order
* @param half2 Point[] brother of half1
* @param result Point[] to be returned when brothers combined and sorted.
205
206
207
208
209
210
        private static void merge(Point[] half1, Point[] half2, Point[] result){
211
          int x = 0;
           int y = 0;
212
          int merge = 0;
213
214
          while(x < half1.length && y < half2.length ){
  if(half1[x].getX() < half2[y].getX()){
    result[merge] = half1[x];</pre>
215
216
217
218
219
             else{
220
               result [merge] = half2[y];
221
222
223
224
             merge++;
225
           {\bf \acute{S}ystem.arraycopy}\,(\,half1\,,\,\,x\,,\,\,result\,\,,\,\,merge\,,\,\,half1\,.\,length\,\,-\,\,x\,)\,;
226
          System.arraycopy(half2, y, result, merge, half2.length - y);
227
228
             ******* End mergeSort Methods ***********
229
230
231
             ******* Begin insertionSort Methods ************
232
233
234
                With online help from:
235
                Title: http://www.java2novice.com/
236
                 Author: N/A
237
                Date: N/A
238
                 Availability: http://www.java2novice.com/java-interview-programs/insertion-sort/
239
240
241
242
        * Sort Point[] by inserting unsorted points into the correct positions in sorted Point[]
243
244
         * @param array Point[] array to be sorted
245
246
        public static void insertionSort(Point array[]) {
247
248
           int n = array.length; // Limiter initalized to length of Point[] array
249
           for (int j = 1; j < n; j++) {
             Point key = array[j]; // Key to comparison int i = j-1; // Interloop counter, one less than 'j' while ( (i > -1) && ( array[i].getX() > key.getX() ) ) {
    array [i+1] = array [i];
250
251
252
253
254
255
256
             array[i+1] = key; // Inserting key into sorted portion of array.
257
258
        }
259
260
            ******* End insertionSort Methods ***********
261
262
263
            ******* Begin selectionSort Methods ***********
264
265
266
                With online help from:
                Title: Sorting.java
267
                 Author: Lewis/Loftus
268
269
                Date: N/A
                 Availability: http://www.ics.uci.edu/~stasio/winter06/Lectures/Lec7code/ComparableExample
270
                /Sorting.java
271
272
273
         \ast Sorts the Point array of objects using the selection sort algorithm.
274
         * @param input Point[] array to be sorted
275
276
277
        public static void selectionSort (Point[] input)
278
279
280
           int min; // Min value
281
           Point temp; // Temp storage of Points
282
283
           for (int index = 0; index < input.length -1; index++)
284
```

```
min = index; // Initializing min
for (int scan = index+1; scan < input.length; scan++)
  if (input[scan].getX() < input[min].getX())</pre>
285
286
287
288
                 min = scan; // update min
289
290
              // Swap the values
temp = input[min];
291
              input [min] = input [index];
input [index] = temp;
292
293
294
         }
295
296
297
                                  Cocktail Shaker Sort
298
      //http://www.javacodex.com/Sorting/Cocktail-Sort
299
300
301
          * Cocktail Shaker Sort
302
          * @param array Point[] of some size to be processed
303
304
         public static void cocktailShakerSort( Point[] array ){
305
306
            boolean swapped;
307
            do {
308
              swapped = false;
              for (int i =0; i <= array.length - 2; i++) {
  if (array[ i ].getX() > array[ i + 1 ].getX()) {
309
310
                      test whether the two elements are in the wrong order
311
312
                    Point temp = array[i];
                    array[i] = array[i+1];
array[i+1]=temp;
313
314
315
                    swapped = true;
                 }
316
317
318
                   (!swapped) {
319
                  //we can exit the outer loop here if no swaps occurred.
320
                 break;
321
322
              swapped = false;
              for (int i= array.length - 2; i>=0;i--) {
    if (array[ i ].getX() > array[ i + 1 ].getX()) {
        Point temp = array[i];
    }
}
323
324
325
                    array[i] = array[i+1];
array[i+1]=temp;
326
327
328
                    swapped = true;
329
                 }
330
               //if no elements have been swapped, then the list is sorted
331
332
           } while (swapped);
333
334
335
          * Force Push Sort was coded by Terry Speicher to show how he *thought* the Cocktail
336
          * Shaker was supposed to work. It is kind of a "double ended" selection sort.

* @param p Point[] of some size to be processed
337
338
339
         public void forcePush (Point [] p) {
340
341
            Point temp = new Point(0.0, 0.0);
342
343
            \label{eq:point_swapper} \text{Point} \ \text{swapper} \ = \ \underset{}{\text{new}} \ \ \text{Point} \ (\, 0.0 \, , 0.0 \, ) \ ;
344
345
            for (int start = 0, end = p.length -1; start < end; start++, end--){ if (start != end) { //only 1 element. We are done.
346
347
                 if ( (end - start) == 1) {
348
                    //only two elements left
//compare and swap if necessary
349
350
                    if (p[start].getX() > p[end].getX()){
351
352
                      temp = p[start];
p[start] = p[end];
p[end] = temp;
353
354
355
                 } else {
356
                    //more than 2 elements - go through and find max and min and swap
357
                    double maxValue;
358
359
                    double minValue;
360
                    int minValueIndex;
361
                    int maxValueIndex;
362
                    minValue = maxValue = p[start].getX(); //set min and max to first element
363
364
                    minValueIndex = maxValueIndex = start;
365
```

```
for (int i = start + 1 ; i <= end ; i++){ //go through array to find max and mins if (p[i].getX() < minValue) { minValue = p[i].getX();
366
367
368
                      minValueIndex = i;
369
370
                      else
                      if (p[i].getX() > maxValue) {
maxValue = p[i].getX();
371
372
373
                      maxValueIndex = i;
374
                     //after this for statement, we know the locations of the max and min elements
375
376
377
                  swapper \,=\, p\,[\,start\,]\,;
                 p[start] = p[minValueIndex];
378
379
                 p[minValueIndex] = swapper;
380
                  //this handles the funny case where the max value was in the start position
381
                  ///but then we moved it when we swapped it with the element that actually should //end up in the start position. That was the swap above, which would leave our
382
383
                  //max element in the position with the index of 'minValueIndex'.
384
385
                  if (maxValueIndex == start)
386
                    maxValueIndex = minValueIndex;
387
388
                  swapper = p[end];
389
                 p[end] = p[maxValueIndex];
390
                  p[maxValueIndex] = swapper;
391
392
393
               }
394
395
396
          }//end main for loop
397
        } //end forcePush
398
399
400
         * Utilitarian method to print out the x value from an array of Points. This was used
401
         * to find the odd case in the Force Push sort where the largest element in the portion
402
         * of the array to be sorted was in the first slot and therefore got moved when the
403
         * smallest element was swapped into its place. The method was left in the class for
         * documentation purposes only.

* @param a Point[] of some size to be processed
404
405
406
        public void printArray(Point [] a) {
  for (int i = 0 ; i < a.length ; i++) {
    System.out.printf("%.2f",a[i].getX());</pre>
407
408
409
410
411
          System.out.println();
412
413
414
415
     }
     //Terry Speicher and Jason Ivey
 1
 2
 3
      * Read file into an array and present different sized sub arrays of those points
  4
      * to each different sort routines and record timed results

* @author Terry Speicher
  5
  6
      * @author Jason Ivey
  7
  8
 9
     import java.util.*;
 10
     import java.util.Arrays;
 11
     import java.io.*;
 12
 13
     public class SortAlgorithmsTester{
 14
        private static Point[] data = new Point[100000];
 15
 16
 17
 18
         * Main body of Sort Tester.
 19
         * @param args standard header String[]
 20
 21
        public static void main(String args[]) {
 22
 23
          /** Create int array representing the number of elements that will be taken from the
 24
            *
*/
[]
            * beginning of the array of random items
 25
 26
                  test Cases \ = \ \{2\,,2\,,3\,,4\,,5\,,10\,,15\,,20\,,25\,,30\,,35\,,40\,,45\,,50\,,55\,,60\,,65\,,70\,,75\,,80\,,85\,,90\,,95\,,100\,,
             200,300,400,500,600,700,800,900,1000,2000,3000,4000,5000,6000,7000,
 27
 28
             8000,9000,10000,data.length };
 29
 30
          // Set the number of times the data will be tested and averaged.
```

```
31
           int iterations = 10;
            //results table is 9 by however many testCases there are
 32
           long [][] resultsTable = new long[8][testCases.length];
 33
           boolean testing = false; //Turn on/off verbose intermediate findings
 34
 35
 36
           read(); //one time read of data[]
 37
           SortAlgorithms sort = new SortAlgorithms(); //create class of sort methods
 38
 39
            //Main counter for determined number of test cases
 40
 41
           for (int counter = 1; counter <= iterations; counter++){
 42
             //Visual output of loop # currently being processed
System.out.println("Iteration#" + counter + " of " + iterations);
 43
 44
 45
              //Print headings
 46
              if (testing) System.out.print("Sort/#Items,");
 47
              for (int i = 0; i < testCases.length; i++)
if (testing) System.out.print(testCases[i] + ",");</pre>
 48
 49
 50
              if (testing) System.out.println();
 51
              //Print each sort name and test results from each test of n elements
if (testing) System.out.print("BubbleSort,");
 52
 53
              for (int i = 0 ; i < testCases.length ; i++){
  long startTime = System.nanoTime();</pre>
 54
 55
                sort.bubbleSort(Arrays.copyOfRange(data,0,testCases[i]));
 56
 57
                long estimatedTime = System.nanoTime() - startTime;
 58
                if (testing) System.out.print(estimatedTime + ",
 59
                resultsTable[1][i] += estimatedTime;
 60
 61
              if (testing) System.out.println();
 62
 63
              //Keep printing name of sort and results
              if (testing) System.out.print("ForcePush,");
 64
              for (int i = 0 ; i < testCases.length ; i++){
  long startTime = System.nanoTime();</pre>
 65
 66
 67
                sort.forcePush(Arrays.copyOfRange(data,0,testCases[i]));
                long estimatedTime = System.nanoTime() - startTime;
if (testing) System.out.print(estimatedTime + ",");
 68
 69
 70
                resultsTable [2][i] +=estimatedTime;
 71
              if (testing) System.out.println();
 72
 73
              if (testing) System.out.print("MergeSort,");
 74
              for (int i = 0; i < testCases.length; i++){
 75
                long startTime = System.nanoTime();
                sort.mergeSort(Arrays.copyOfRange(data,0,testCases[i]));
 77
                long estimatedTime = System.nanoTime() - startTime;
 79
                if (testing) System.out.print(estimatedTime + ",");
                resultsTable[3][i] +=estimatedTime;
 80
 81
 82
              if (testing) System.out.println();
 83
              if (testing) System.out.print("QuickSort,");
 84
 85
              for (int i = 0; i < testCases.length; i++){
                long startTime = System.nanoTime();
 86
 87
                sort.quickSort(Arrays.copyOfRange(data,0,testCases[i]));
                long estimatedTime = System.nanoTime() - startTime;
 88
                if (testing) System.out.print(estimatedTime + ",");
 89
                resultsTable [4][i] +=estimatedTime;
 90
 91
 92
              if (testing) System.out.println();
 93
              if (testing) System.out.print("InsertionSort,");
for (int i = 0 ; i < testCases.length ; i++){
  long startTime = System.nanoTime();</pre>
 94
 95
 96
                sort.insertionSort\left(Arrays.copyOfRange\left(data,0,testCases\left[i\right]\right)\right);
 97
 98
                {\color{red} \textbf{long}} \hspace{0.2cm} \textbf{estimatedTime} \hspace{0.1cm} = \hspace{0.1cm} \textbf{System.nanoTime} \hspace{0.1cm} (\hspace{0.1cm}) \hspace{0.1cm} - \hspace{0.1cm} \textbf{startTime} \hspace{0.1cm} ;
                if (testing) System.out.print(estimatedTime + ",");
resultsTable[5][i] +=estimatedTime;
 99
100
101
102
              if (testing) System.out.println();
103
              if (testing) System.out.print("SelectionSort,");
104
              for (int i = 0; i < testCases.length; i++){
105
106
                long startTime = System.nanoTime();
107
                sort.selectionSort(Arrays.copyOfRange(data,0,testCases[i]));
108
                \begin{array}{lll} \textbf{long} & \textbf{estimatedTime} & = & \textbf{System.nanoTime()} & - & \textbf{startTime}; \end{array}
                if (testing) System.out.print(estimatedTime + ",
109
110
                resultsTable[6][i] +=estimatedTime;
111
```

```
112
             if (testing) System.out.println();
113
             if (testing) System.out.print("ShakerSort,");
for (int i = 0 ; i < testCases.length ; i++){
  long startTime = System.nanoTime();</pre>
114
115
116
               sort.cocktailShakerSort(Arrays.copyOfRange(data,0,testCases[i]));
117
               long estimatedTime = System.nanoTime() - startTime;
118
               if (testing) System.out.print(estimatedTime + ",");
119
120
               resultsTable[7][i] +=estimatedTime;
121
122
             if (testing) System.out.println();
123
124
          }
125
          //In verbose mode, print out time totals table
126
          if (testing) {
127
                                                                -Totals over " + iterations + " iterations
            System.out.println("-
128
                                              ----");
                                                         System.out.print("Sort/#Items,");
             for (int i = 0; i < testCases.length; i++)
System.out.print(testCases[i] + ",");
129
130
            System.out.println();
131
132
133
             for (int a = 1; a <= 7; a++){
134
               switch (a) {
135
                 case 1:
                    System.out.print("Bubble,");
136
137
                    break;
138
                 case 2:
139
                    System.out.print("ForcePush,");
140
                    break;
141
                 case 3:
142
                    System.out.print("Merge,");
143
                    break;
144
                 case 4:
145
                    System.out.print("Quick,");
146
                    break;
147
                 case 5:
148
                    System.out.print("Insertion,");
149
150
                 case 6:
                    System.out.print("Selection,");
151
152
                    break;
                 case 7:
153
154
                    System.out.print("Shaker,");
155
                    break;
156
                  //end switch case
157
               for (int j = 0; j < resultsTable[0].length; j++){
158
                 System.out.print(resultsTable[a][j] + ",");
159
160
161
               System.out.println();
162
163
         }
164
165
166
           //Same loop as above, but print out the averages instead of the total time
167
                                                          ----Averages over " + iterations + " iterations
          System.out.println("-
168
169
           \begin{array}{ll} System.out.print("Sort/\#Items");\\ for\ (int\ i=0\ ;\ i< testCases.length\ ;\ i++)\\ System.out.print(testCases[i]+""); \end{array} 
170
171
172
173
          System.out.println();
174
          for (int a = 1; a \le 7; a++){
175
176
            switch (a) {
177
               case 1:
                 System.out.print("Bubble,");
178
179
                 break;
180
               case 2:
                 System.out.print("ForcePush,");
181
182
                 break;
183
               case 3:
184
                 System.out.print("Merge,");
185
                 break;
186
               case 4:
187
                 System.out.print("Quick,");
                 break;
188
               case 5:
189
190
                 System.out.print("Insertion,");
```

```
191
                  break;
192
                case 6:
193
                   System.out.print("Selection,");
194
                   break:
195
                case 7:
196
                   System.out.print("Shaker,");
197
             break;
} //end switch case
198
199
             \begin{array}{lll} for \ (int \ j = 0 \ ; \ j < resultsTable [0]. \ length \ ; \ j++) \{ \\ System.out.print (resultsTable [a][j]/iterations+","); \end{array}
200
201
202
203
             System.out.println();
204
205
           }
206
207
        }
208
209
210
         * Read in the file with the pairs of point (x,y) coordinates and place in the data[]
211
212
213
         private static void read(){
214
215
           try {
  File file = new File("100000 Points.txt");
216
217
218
             Scanner scan = new Scanner (file);
219
220
             int x = 0;
              String line;
221
              String token [];
222
223
              while (scan.hasNext()) {
                line = scan.nextLine();
token = line.split("\t");
224
225
                data[x] = new Point(Double.parseDouble(token[0]), Double.parseDouble(token[1]));
^{226}
227
228
229
             //confirm number of elements read
System.out.println("We have " + x + " points");
230
231
232
             scan.close();
233
234
           catch (Exception e) {
             System.out.println("Error in reading file");
235
236
237
        }
238
239
240
     }
```