

## ECE 2372 - Homework 4

### << — VC BOUND — >>

Please upload your solutions to Canvas by February 25, 2024

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#### Problem 1:

The VC dimension, rather implicitly, depends both on the input space  $\mathcal{X}$  and  $\mathcal{H}$ — since  $\mathcal{H}$  is defined as the mapping of  $\mathcal{X}$  to a binary set of  $\{-1, +1\}$ .

For a fixed  $\mathcal{H}$ , consider two input spaces  $\mathcal{X}_1 \subseteq \mathcal{X}_2$ . Let  $d_{VC}(\mathcal{X}_1)$  and  $d_{VC}(\mathcal{X}_2)$  denote the VC dimension of  $\mathcal{H}$  with respect to these input spaces. Show that  $d_{VC}(\mathcal{X}_1) \leq d_{VC}(\mathcal{X}_2)$ .

**Sol.**

For  $x_1$ , the VC dimension is  $d_{VC}(x_1)$ , which means in  $x_1$  there are  $d_{VC}(x_1)$  points can be shattered.

$x_1 \subseteq x_2$ ,  $x_2$  has all the points of  $x_1$ . Therefore,  $x_2$  has at least  $d_{VC}(x_1)$  points can be shattered. Which means

$$d_{VC}(x_1) \leq d_{VC}(x_2)$$

#### Problem 2:

Suppose  $\mathcal{H}$  is obtained by evaluating some polynomial of degree  $d$  and comparing the result to a threshold, i.e.,

$$\mathcal{H} = \left\{ h : h(x) = \text{sign} \left( \sum_{i=0}^d c_i x_i \right) \right\}$$

for some  $c_0, \dots, c_d \in \mathbb{R}$ , and  $\mathcal{X} \in \mathbb{R}$  as well. Prove that  $d_{VC}$  of this  $\mathcal{H}$  is exactly  $d + 1$  by showing that

(a) There are  $d + 1$  points which are shattered by  $\mathcal{H}$ .

(b) There are no  $d + 2$  points which are shattered by  $\mathcal{H}$ . [Hint: Try relating this to a linear classifier in  $d$  dimensions and use the result from first problem]

a)

Assume that

$$\begin{aligned} c &= (c_0, c_1, \dots, c_d)^T \\ x &= (x_0, x_1, \dots, x_d)^T \end{aligned}$$

And we rewrite  $H$ :

$$H = \{h: h(x) = \text{sign}(c^T x)\}$$

Suppose there are  $d + 1$  points in  $x$ ,

$$x = \begin{pmatrix} - & (x^{(1)})^T & - \\ - & (x^{(2)})^T & - \\ \dots & \dots & \dots \\ - & (x^{(d+1)})^T & - \end{pmatrix} y = \begin{pmatrix} y_1 \\ y_2 \\ \dots \\ y_{d+1} \end{pmatrix}$$

Let  $x$  be:

$$x = \begin{pmatrix} 1 & 0 & 0 & \dots & 0 \\ 1 & 1 & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots \\ 1 & 0 & 0 & \dots & 1 \end{pmatrix}$$

$x$  is non-singular. Therefore for  $\forall y$ , we can always find

$$c = x^{-1}y$$

that satisfied:

$$\text{sign}(xc) = y$$

And we proved that there are  $d + 1$  points shattered by  $H$ .

b)

$$x_2 = x_1 \cup \{\text{new points}\}$$

From problem 1,

$$\begin{aligned} x_1 &\subseteq x_2 \\ d_{VC}(x_1) &\leq d_{VC}(x_2) \end{aligned}$$

Suppose

$$x = \begin{pmatrix} - & (x^{(1)})^T & - \\ - & (x^{(2)})^T & - \\ \dots & \dots & \dots \\ - & (x^{(d+1)})^T & - \\ - & (x^{(d+2)})^T & - \end{pmatrix}$$

Since  $x^{(j)} \in \mathbb{R}^{d+1}$ , for  $d + 2$  points, there are at least one point  $x^{(j)}$  can be represented as

$$x^{(j)} = \sum_{i \neq j} \alpha_i x^{(i)}$$

And

$$y^{(j)} = \text{sign}(c^T x^{(j)}) = \sum_{i \neq j} \alpha_i \text{sign}(c^T x^{(i)})$$

For a given  $x^{(i)}, x^{(j)}$  (or  $y^{(j)}$ ) is determined by  $x^{(i)}$  and not free.

**Problem 3:** Suppose the VC dimension of our  $\mathcal{H}$  is  $d_{VC} = 3$  (e.g., linear classifiers in  $\mathbb{R}^2$ ). And suppose that we have an algorithm for selecting some  $h^* \in \mathcal{H}$  based on a training sample size  $n$ .

- Using the VC generalization bound, give an upper bound on  $R(h^*)$  that holds with probability at least 0.95 when  $n = 100$ . Repeat this for  $n = 1,000$  and  $10,000$ .
- How large does  $n$  need to be to obtain a bound of the form:

$$R(h^*) \leq \hat{R}(h^*) + 0.01$$

that holds with probability at least 0.95? How does this compare to the “rule of thumb” that we stated in the class?

a)

$$d_{VC} = 3, \delta = 1 - 0.95 = 0.05$$

$$R(h^*) \leq \hat{R}(h^*) + \sqrt{\frac{8d_{VC}}{n} \log \frac{8n}{\delta}}$$

$n = 100$ :

$$Upper\ Bound = \sqrt{\frac{8 * 3}{100} \log \frac{8 * 100}{0.05}} = 1.524$$

$n = 1000$ :

$$Upper\ Bound = \sqrt{\frac{8 * 3}{1000} \log \frac{8 * 1000}{0.05}} = 0.536$$

$n = 10000$ :

$$Upper\ Bound = \sqrt{\frac{8 * 3}{10000} \log \frac{8 * 10000}{0.05}} = 0.185$$

b)

$$\sqrt{\frac{8 * 3}{n} \log \frac{8 * n}{0.05}} = 0.01$$

If  $n = 10^6$ , Bound  $\approx 0.021$ , if  $n = 10^7$ , Bound  $\approx 0.007$ , therefore,

$$n \in (10^6, 10^7)$$

Compare to the 'rule of thumb', we see that for every  $n$ ,

$$n \geq 10d_{VC} = 30$$