

# Design and Analysis of Algorithms

## Tutorial 2



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# Question 1

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Give asymptotic upper bounds for  $T(n)$  by recursion tree approach. Make your bounds as tight as possible.

(a)

$$T(1) = 1$$

$$T(n) = T\left(\frac{n}{2}\right) + n \quad \text{if } n > 1$$

(b)

$$T(1) = T(2) = 1$$

$$T(n) = T(n - 2) + 1 \quad \text{if } n > 2$$

(c)

$$T(1) = 1$$

$$T(n) = T\left(\frac{n}{3}\right) + n \quad \text{if } n > 1$$

# Question 1

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(d)

$$T(1) = 1$$

$$T(n) = 4T\left(\frac{n}{2}\right) + n \quad \text{if } n > 1$$

(e)

$$T(1) = 1$$

$$T(n) = 3T\left(\frac{n}{2}\right) + n^2 \quad \text{if } n > 1$$

(f)

$$T(1) = 0, T(2) = 1$$

$$T(n) = T\left(\frac{n}{2}\right) + \log_2 n \quad \text{if } n > 2$$

# Question 2

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Prove the following problems by induction.

(a) Given

$$T(1) = 1$$

$$T(n) = T\left(\frac{n}{2}\right) + n \quad \text{if } n > 1$$

Prove  $T(n) \leq c \cdot n$  for some  $c$ .

(b) Given

$$T(1) = T(2) = 1$$

$$T(n) = T(n-2) + 1 \quad \text{if } n > 1$$

Prove  $T(n) \leq c \cdot n$  for some  $c$ .

(c) Given

$$T(1) = 1$$

$$T(n) = T\left(\frac{n}{3}\right) + n \quad \text{if } n > 1$$

Prove  $T(n) \leq c \cdot n$  for some  $c$ .

# Question 2

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(d) Given

$$T(1) = 1$$

$$T(n) = 4T\left(\frac{n}{2}\right) + n \quad \text{if } n > 1$$

Prove  $T(n) \leq c_1 \cdot n^2 - c_2 \cdot n$  for some  $c_1$  and  $c_2$ .

(e) Given

$$T(1) = 1$$

$$T(n) = 3T\left(\frac{n}{2}\right) + n^2 \quad \text{if } n > 1$$

Prove  $T(n) \leq c_1 \cdot n^2$  for some  $c$ .

(f) Given

$$T(1) = 0, T(2) = 1$$

$$T(n) = T\left(\frac{n}{2}\right) + \log_2 n \quad \text{if } n > 2$$

Prove  $T(n) \leq c \cdot \log^2 n$  for some  $c$ .

# Question 3

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Let  $A[1..n]$  be an array of positive integers.

Design a divide-and-conquer algorithm for computing the maximum value of  $A[j] - A[i]$  with  $j \geq i$ .

Analyze your algorithm running time.