ECE408 Final Project

CNN convolution layer gpu accelerator

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ECE408 FINAL PROJECT

Introduction

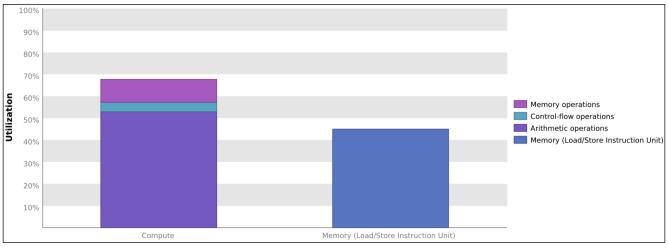
This project was to implement and optimize a fast convolution forward pass kernel of input size 70x70. A baseline implementation of such was done using sequential code on the CPU and another for code based on the GPU. The GPU was what was the baseline for our optimizations. Each optimization was to decrease the runtime while keeping the correctness of the baseline data the same. The optimization that yielded the best results was using a parallel shared matrix multiplication and a regular shared matrix multiplication for different layer sizes.

4.1 Implicit Unrolling & Data parallelism

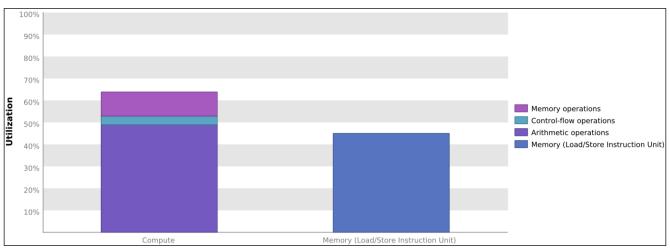
Description

Performance of matrix multiplication with kernel unrolling is limited by global memory in two ways. First, the unrolling factor is approximately K*K. In the 10000 dataset, the unrolled matrix X cannot fit into global memory. Therefore, a method of mini-batch is used in pervious optimization. We reduce the amount of parallelism in the B (batch) dimension so that unrolled matrix X can fit into global memory. Second, kernel unrolling method requires too much global memory access as discussed in earlier checkpoints.

The first part of this optimization is to avoid explicit kernel unrolling to solve the limitation in global memory size and accessing. Unrolling can be performed when loading the tiles in the matrix multiplication kernel. This reduces the runtime from ~400ms to ~80ms.



Regular Matrix Multiplication

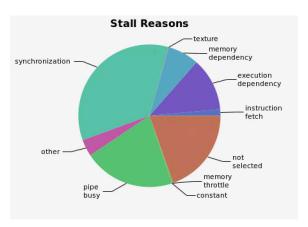


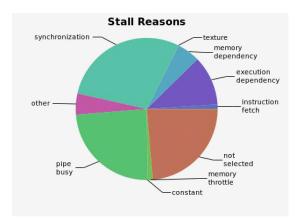
Parallelized Batch Matrix Multiplication

The second part of this optimization is input parallelism in B (batch) dimension. Note that fact that we can using one weight for all pictures. In the tiling loading step, we try to avoid some unnecessary data loading. To do this, instead of letting each block dealing with one input matrix, a third block dimension is added so that we can assign multiple input matrices to one block and use the same weight matrices in the shared memory. To match with this, we added a new dimension for matrix X's tile, indexed by threadIdx.z. While the shared memory for weight stays the same. This is to some extent similar to tiling in the sense that we partition the input data and reuse the same data in shared memory as much as possible.

Data parallelism only improve the runtime of layer 1 by ~5ms together with for loop unrolling and restrict tuning. It also makes layer 2 slower. According to our understanding, the main reason for the latency is because of an increased number of for loop iteration in on each thread. Since each CUDA block can only have a maximum of 1,024 threads, we have no

choice but to decrease tile size in order to introduce a third dimension to blockDim. This works good for input matrices with small dimensions but has significant drawbacks if the input matrices size is large.





Regular Matrix Multiplication

Parallelized Batch Matrix Multiplication

As you can see from the figures above, Regular Matrix Multiplication has more synchronization overhead and memory dependency. This agrees with out reasoning about the potential improvement that Parallelized Batch Matrix Multiplication can bring to the kernel — we load more often from global memory in Regular Matrix Multiplication, and each time we load, we need to synchronize the thread.

This parallel approach has a little portion of control divergence:

▼ Line / File	new-forward.cuh - /mxnet/src/operator/custom
130	Divergence = 12.7% [277000 divergent executions out of 2184000 total executions]

```
subTileB[b][ty][tx] = (Col < numBColumns && temp_row < numBRows
&& X_b < numBatch) ? x4d(X_b, X_c, X_h + X_p, X_w + X_q) : 0;</pre>
```

This is the line of code to load input matrices into shared memory. We cannot avoid this divergence because total batch number might not be multiple of Batch partition size.

Output

Loading fashion-mnist data... done

Loading model... done

New Inference **Op Time:** 0.025720 **Op Time:** 0.083995

Correctness: 0.7653 Model: ece408

As we can see from the output. Layer 1's runtime is less than milestone 4 while layer 2's runtime is longer in contrast.

all kernels that collectively consume more than 90% of the program time:

[CUDA memcpy HtoD]

void mshadow::cuda::MapPlanLargeKernel<mshadow::sv::saveto, int=8, int=1024, mshadow::expr::Plan<mshadow::gpu, int=4, float>, float>, mshadow::expr::Plan<mshadow::expr::BinaryMapExp<mshadow::op::mul, mshadow::expr::ScalarExp<float>, mshadow::Tensor<mshadow::gpu, int=4, float>, float, int=1>, float>>(mshadow::gpu, unsigned int, mshadow::Shape<int=2>, int=4, int)

volta_sgemm_128x128_tn

all CUDA API calls that collectively consume more than 90% of the program time:

cudaStreamCreateWithFlags

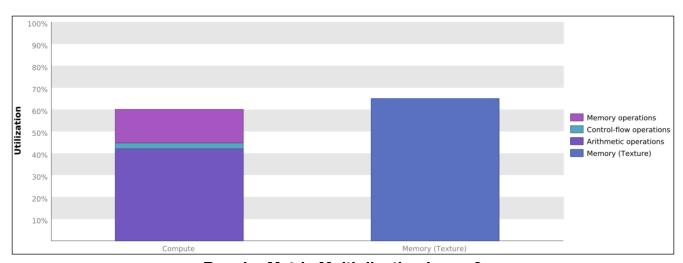
cudaMemGetInfo

cudaFree

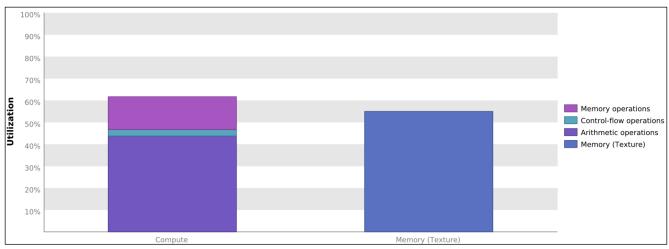
4.2 Tuning with restrict & loop unrolling

Description

Because all the dimensions are fixed in our case, we can further improve our instructions. Instead of doing control flow in run time, we can calculate the number of iterations for all the for loops. To do this we use the # pragma unroll directives in front deterministic for loops to tell the complier to generate instructions by repeating commands without explicit for loop. Besides, we also include <code>__restrict__</code> keyword to hint the compiler that we only access the non-overlapping arrays via their pointers. Lastly, <code>const</code> keyword in used for all read only variables to ensure safe access.



Regular Matrix Multiplication Layer 2



For loop unrolling & restrict Matrix Multiplication

As shown the figures above, before doing these subtle tuning. The matrix multiplication kernel is bounded by both instruction and memory latency. After this optimization Compute utilization increase a little bit, and the kernel is bounded by only memory bandwidth now.

In terms of runtime, adding unroll and restrict along only reduce ~1ms. But after including multiple kernels for different layers. It seems that parallelized matrix multiplication benefits more. A ~5ms speed up is achieve with layer 1 parallelized matrix multiplication. A possible reason for this is in order to add parallel in B (batch) dimension, we reduced the TILE_WIDTH in tiled matrix multiplication. Therefore, introduced more iterations in for loops.

Output The output of 4.2 is combined with 4.3

4.3 Multiple kernel implementations for different layer sizes

Description

The optimization Involved using separate kernel Implementations for different layers. The first layer was executed using a parallel matrix multiplication kernel using a (16,16,4) block dimension. Layer 2 was executed using a regular matrix multiplication kernel using a (32,32,1) block dimension. These block dimensions were chosen such that the thread counts are divisible by a warp size (32) to reduce control divergence.

As mentioned in 4.1. Data parallelism will slow down the overall runtime. Therefore, we adopt multiple kernel so that we only apply parallelized batch to kernel 1. This combined optimization brings around 50 ms decrease to the total run time.

Using data parallelism means we need to sacrifice the boost of tiled matrix multiplication since the product of tile size and number of batches handled in parallel is at most 1,024, the max number of threads we can have in a block. We have tried different combinations, (32, 32, 1), (16, 16, 4), (8, 8, 16), etc., and (16, 16, 4) gives us the greatest improvement to the overall performance.

Output

Loading fashion-mnist data... done

Loading model... done

New Inference **Op Time:** 0.025109 **Op Time:** 0.053508

Correctness: 0.7653 Model: ece408

GPU Activities

all kernels that collectively consume more than 90% of the program time:

[CUDA memcpy HtoD]

void mshadow::cuda::MapPlanLargeKernel<mshadow::sv::saveto, int=8, int=1024, mshadow::expr::Plan<mshadow::gpu, int=4, float>, float>, mshadow::expr::Plan<mshadow::expr::BinaryMapExp<mshadow::op::mul, mshadow::expr::ScalarExp<float>, mshadow::float>, float>, float>, float>, float>, float>, int=1>, float>>(mshadow::gpu, unsigned int, mshadow::Shape<int=2>, int=4, int)

volta_sgemm_128x128_tn

all CUDA API calls that collectively consume more than 90% of the program time:

cudaStreamCreateWithFlags

cudaMemGetInfo

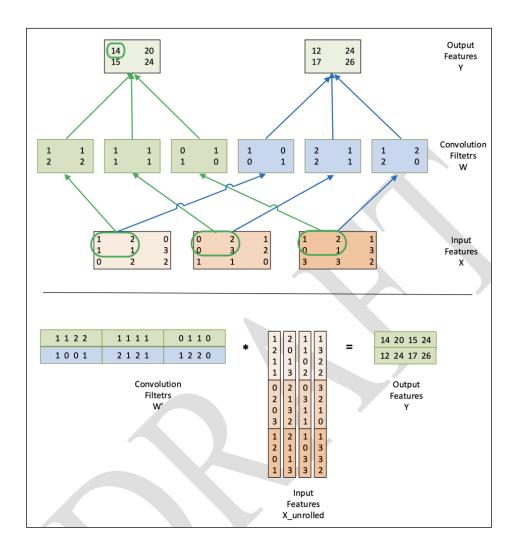
cudaFree

3.1 Unroll and Matrix Multiplication

Description

The optimization unrolls both input data matrices and weight matrices so that we can perform a simple shared memory matrix multiplication on the data to generate the output. The unrolling of weight matrices needs no additional execution since the row-major layout, we only need to change the indexing when we access it as unrolled matrices. For the input data, we run an additional kernel to perform the unroll. The unrolled matrix exceeds the assigned CUDA memory size when we perform op2 for 10000 dataset. So, we partition the input data into small sections and deal with each section sequentially. This slows down the performance of this optimization for the last dataset.

This optimization is based on Chapter 16:



Unrolling and matrix multiplication for convolution

We utilize this optimization because we can transfer a time-complex matrix convolution operation into a faster matrix multiplication operation, but the runtime is about the same for the datasets. According to our understanding, the reason is mainly that the memory utilization grows a lot for this optimization compared to the basic approach. There are much more global and shared memory reads and writes, so even our code becomes more parallel, the overall performance is not better.

We use two kernels for this optimization, one matrix unroll and one matrix multiplication. Matrix multiplication kernel takes more time to finish than the other.

Matrix Multiplication:

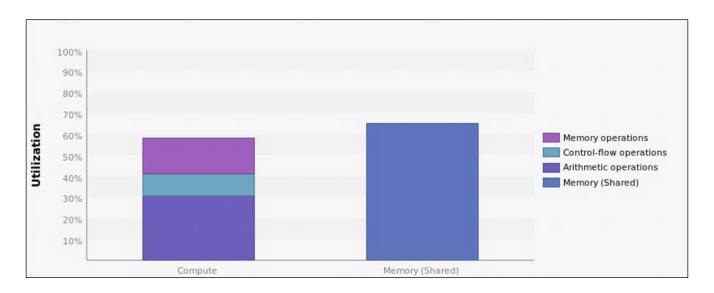


Fig 3.1.1

This kernel is Bounded by Memory Bandwidth.

As we can see from **figure 3.1.1**, which is kernel performance chart for matrix multiplication, the memory utilization is higher than compute utilization. This is obvious with our implementation since the matrix multiplication kernel accesses shared memory a lot.

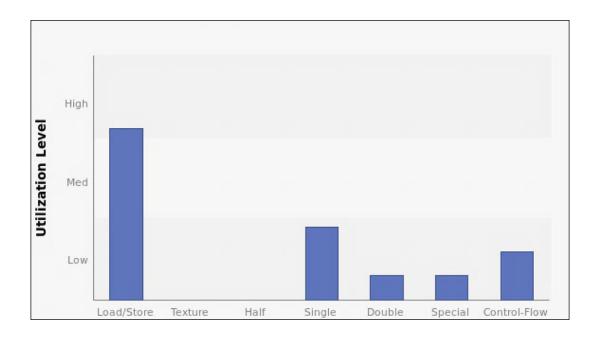


Fig 3.1.2

Figure 3.1.2 is function unit utilization chart, this also indicates that the matrix multiplication kernel mainly spends time on memory load and store.

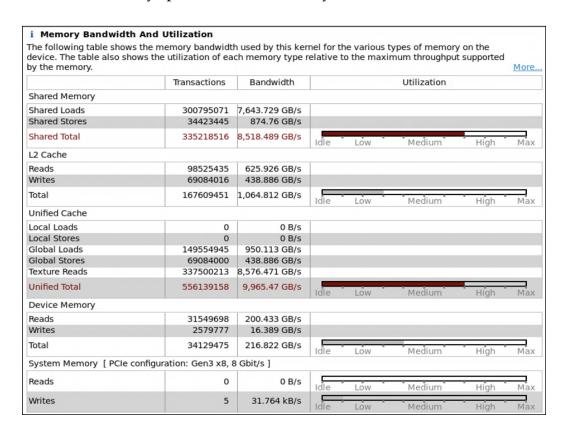


Fig 3.1.3

Figure 3.1.3 is the memory usage of Matrix multiplication kernel. Both shared memory and unified cache memory is accessed frequently. We are utilizing the memory bandwidth pretty well.

Matrix Unroll:

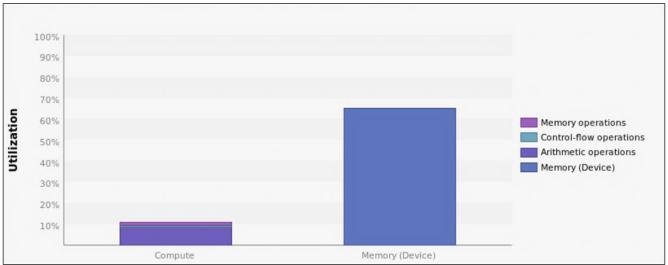


Fig 3.1.4

As we conclude from **figure 3.1.4**, this kernel's utilization distribution is very unbalanced. This is consistent with our design since we unroll the matrix, which mainly moves data from one place in memory to another. For further improvement to our optimization, we can abandon this kernel and instead access data from original matrix by changing the indexing in matrix multiplication kernel.

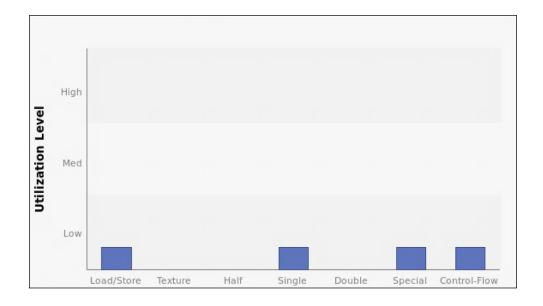


Fig 3.1.5

Figure 3.1.5 shows that utilization level is low for all performances in the kernel. The same reason applies here.

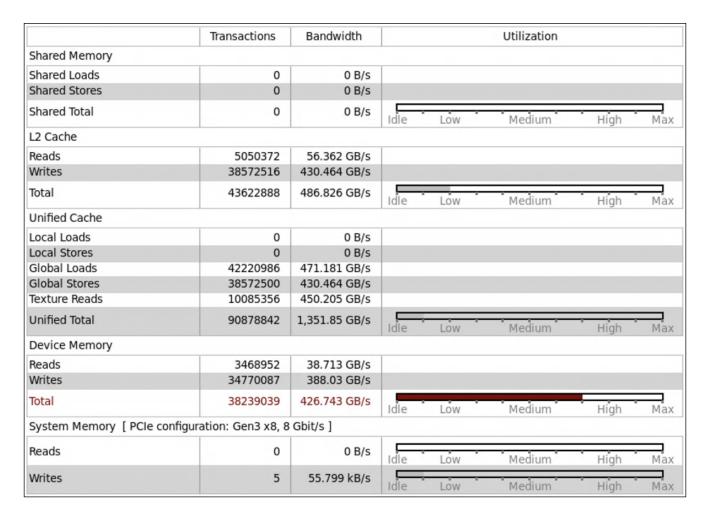


Fig 3.1.5

Memory access is exactly as expected: we perform a lot of global memory access in this kernel.

GPU Activities:

mxnet::op::matrixMultiplyShared(float*, float*, float*, int, int, int, int, int, int)

Time(%) Time Calls Avg Min Max 34.31% 67.957ms 5 13.591ms 3.5026ms 17.328ms

mxnet::op::unroll_Kernel(int, int, int, int, int, float*, float*)

Time(%) Time Calls Avg Min Max 22.23% 39.680ms 5 7.9360ms 2.1032ms 10.816ms

[CUDA memcpy HOTD]

Time(%) Time Calls Avg Min Max 19.72% 35.198ms 20 1.7599ms 1.0880us 32.926ms

API Calls:

cudaStreamCreateWithFlags

Time(%)	Time	Calls	Avg	Min	Max
41.65%	3.37645s	22	153.48ms	15.188us	1.68743s
cudaMem(GetInfo				
Time(%)	Time	Calls	Avg	Min	Max
30.50%	2.47277s	24	103.03ms	72.464us	2.46793s

cudaFree

Time(%) Time Calls Avg Min Max 20.58% 1.73656s 20 86.828ms 1.2690us 454.89ms

cudaMalloc

Time(%) Time Calls Avg Min Max 4.40% 356.67ms 71 5.0236ms 7.7960us 121.33ms

3.2 Shared Memory Convolution

Description

Kernel performance in checkpoint 3 is bounded by memory. Since all data read and write are global, there is a high demand on memory bandwidth. This optimization use shared memory to reduce memory dependency on convolution kernel execution. Similar to lab: 3D convolution, we preload the data and weight matrices on shared memory.

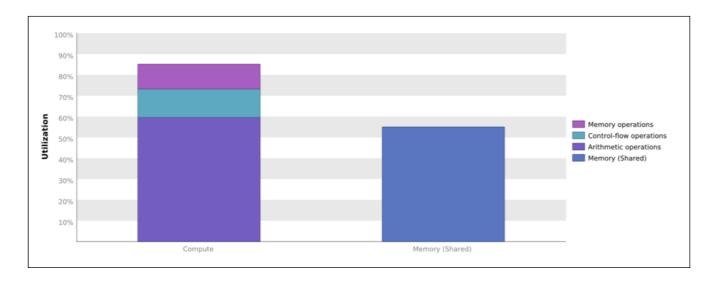


Fig 3.2.1

As we can see from **figure 3.2.1**, which is kernel performance chart for shared memory convolution, the compute utilization is higher than memory utilization. It kernel is bounded by computation. As further proved in **figure 3.2.3** resource analysis. TITIAN V supports 65536 registers for each block and 32 blocks in a SM. The kernel used 10240 registers. Therefore, only 6 block is simultaneously executing on a SM.

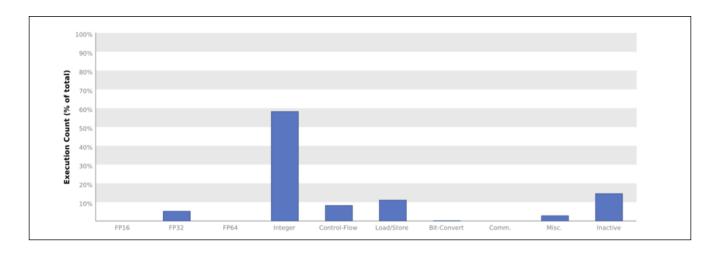


Fig 3.2.2

Figure 3.2.2 is function unit utilization chart, this also indicates that the shared memory kernel mainly spends time on memory load and store.

Line 43 Divergence = 37.5% [3456000 divergent executions out of 9216000 total executions]

if(t < C * W_u nroll && b < B) // when loading shared memory

Line 47 Divergence = 50% [9216000 divergent executions out of 18432000 total executions]

```
for(p = 0; p < K; p++){
   for(q = 0; q < K; q++){ // for loop for convolution</pre>
```

One major disadvantage of share memory method is the control divergence when transferring data from global memory to shared memory.

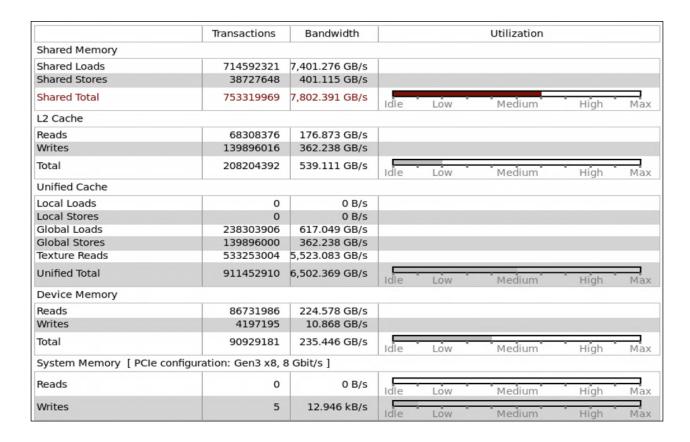


Fig 3.2.3

Figure 3.2.3 is the memory usage of Shared Memory Convolution kernel. Share Memory is properly utilized.

GPU Activities:

mxnet::op::forward_kernel_shared(float*, float const *, float const *, int, int, int, int, int)						
Time(%)	Time	Calls	Avg	Min	Max	
70.57%	170.98ms	2	85.489ms	39.489ms	131.49ms	
[CUDA mer	ncpy HtoD]					
Time(%)	Time	Calls	Avg	Min	Max	
14.03%	33.995ms	20	1.6998ms	1.0880us	31.782ms	

void mshadow::cuda::MapPlanLargeKernel<mshadow::sv::saveto, int=8, int=1024, mshadow::expr::Plan<mshadow::Tensor<mshadow::gpu, int=4, float>, float>, mshadow::expr::BinaryMapExp<mshadow::op::mul, mshadow::expr::ScalarExp<float>, mshadow::Tensor<mshadow::gpu, int=4, float>, float>, int=1>, float>>(mshadow::gpu, unsigned int, mshadow::Shape<int=2>, int=4, int)

Time(%) Time Calls Avg Min Max 7.03% 17.040ms 2 8.5200ms 3.0484ms 13.992ms

API Calls:

32.04%

cudaStreamCreateWithFlags

2.53134s

22

Time(%) 40.84%	Time 3.22616s	Calls 22	Avg 146.64ms	Min 14.968us	Max 1.71205s
cudaMemG	etInfo				
Time(%)	Time	Calls	Avg	Min	Max

115.06ms

71.350us

2.52673s

3.3 Constant Cache Optimization

Description

The optimization was to use constant cache memory to load the weighted matrix, then perform the convolution by reading through the constant Cache. The reason for this was because the cache reduces memory traffic by being read-only thus making memory access much quicker. By loading our weighted matrix into a constant cache, it is expected that most of the reads will be accessed from the Cache during convolution execution thus improving performance.

The optimization resulted in a similar performance as it were to the baseline kernel (milestone 3.1). The correctness was the same; however, the computational time was also the same. A major factor to this result is the control divergence seen in line 54 of the kernel. There was 50% control divergence, meaning half the threads in the same warp were executing the same branching behavior. Due to high divergence, the compute utilization of the kernel was low to the point where most of the GPU is idle, as seen in the far-right bar in **figure 3.3.1**. We can see that the code reaches 20% execution count for Inactive instructions. Meaning that these threads did not execute any code due to control divergence. As we can see from **figure 3.3.2**, the kernel is also bounded by instruction and memory latency. This shows that the kernel is not fully utilizing the TITAN V to achieve the optimal throughput.

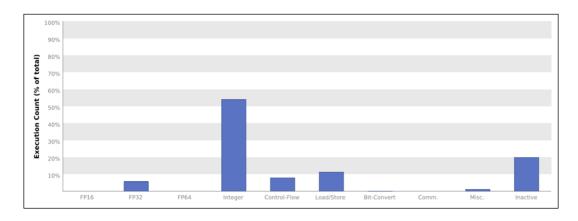


Fig 3.3.1

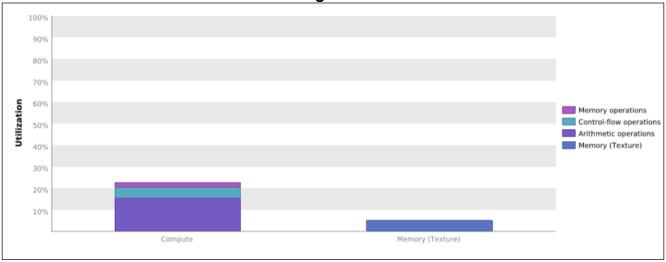


Fig 3.3.2

The reason for trying this optimization was because the constant cache is faster than global memory. By placing the weighted matrix into constant memory, it can be accessed quickly. This is ideal since the weighted matrix is used to calculate every value of the output, so by making these reads quicker, it should cause the overall kernel to be quicker. We believe that the idea is correct; however, the implementation will need to work on optimizing the number of blocks being executed by the kernel to reduce the memory and instruction latency. Along with that, the 50% control divergence causes drastic slow down during execution. We believe by optimizing the control flow branches and finding the optimal grid and block sizes, this optimization can be quicker.

GPU Activities:

mxnet::op::fo	orward_kernel	l(float *, float o	const *, float c	onst *, int, int,	int, int, int, int);		
Time(%)	Time	Calls	Avg	Min	Max		
81.75%	41.815ms	2	20.957ms	4.065ms	37.849ms		
[CUDA mem	[CUDA memcpy HOTD]						
Time(%)	Time	Calls	Avg	Min	Max		
11.41%	5.8487ms	20	293.44us	1.088us	3.557ms		

Void mshadow::cuda::Ma	oPlanLargeKernel<1	mshadow::sv::saveto.	int=8, int=1024,)
		,	-, - , ,

Time(%)	Time	Calls	Avg	Min	Max
2.55%	1.3098ms	2	654.92us	282.43us	1.0274ms

Volta_sgemm_32x32_slicedlx4_tn

Time(%)	Time	Calls	Avg	Min	Max
1.69%	864.54us	2	432.27us	10.047us	854.49us

API Calls:

cudaStreamCreateWithFlags

Time(%)	Time	Calls	Avg	Min	Max
43.77%	3.20445s	22	145.66ms	14.999us	1.65972s
cudaMemGe	etInfo				
Time(%)	Time	Calls	Avg	Min	Max
31.57%	2.31112s	22	105.05ms	93.281us	2.30600s
cudaFree					
Time(%)	Time	Calls	Avg	Min	Max
21.96%	1.60793s	18	89.329ms	1.2550us	431.47ms
cudaMemcp	3				
Time(%)	Time	Calls	Avg	Min	Max
0.01%	497.0us	12	41.466us	6.405us	89.563us
1.34	m 0 1 1				
cudaMemcp					
Time(%)	Time	Calls	Avg	Min	Max
0.00%	81.937us	2	40.968us	38.458us	43.479us

2 GPU implementation

all kernels that collectively consume more than 90% of the program time:

```
mxnet::op::forward_kernel(float*, float const *, float const *, int, int, int, int, int, int, volta_scudnn_128x64_relu_interior_nn_v1 volta_gcgemm_64x32_nt
```

[CUDA memcpy HtoD]

```
void mshadow::cuda::MapPlanLargeKernel<mshadow::sv::saveto, int=8, int=1024,
mshadow::expr::Plan<mshadow::Tensor<mshadow::gpu, int=4, float>, float>,
mshadow::expr::Plan<mshadow::expr::BinaryMapExp<mshadow::op::mul,
mshadow::expr::ScalarExp<float>, mshadow::Tensor<mshadow::gpu, int=4, float>, float, int=1>,
float>>(mshadow::gpu, unsigned int, mshadow::Shape<int=2>, int=4, int)
```

volta sgemm 128x128 tn

void op_generic_tensor_kernel<int=2, float, float, float, int=256, cudnnGenericOp_t=7, cudnnNanPropagation_t=0, cudnnDimOrder_t=0, int=1>(cudnnTensorStruct, float*, cudnnTensorStruct, float const *, cudnnTensorStruct, float const *, float, float,

all CUDA API calls that collectively consume more than 90% of the program time:

cudaStreamCreateWithFlags

cudaMemGetInfo

cudaFree

difference between kernels and API calls:

Kernels are C functions defined by the user to execute N times in parallel by N CUDA threads. Therefore, a kernel launch is to execute the user defined C function. API functions are the functions provided by CUDA to execute some operations on CUDA GPU. API function calls are when users adopt the provided CUDA functions.

output of rai running MXNet on the CPU:

```
* Running /usr/bin/time python m1.1.py
Loading fashion-mnist data... done
Loading model... done
New Inference
EvalMetric: {'accuracy': 0.8154}
```

program runtime:

17.09user 4.83system 0:09.02elapsed 243%CPU (0avgtext+0avgdata 6044912maxresident)k

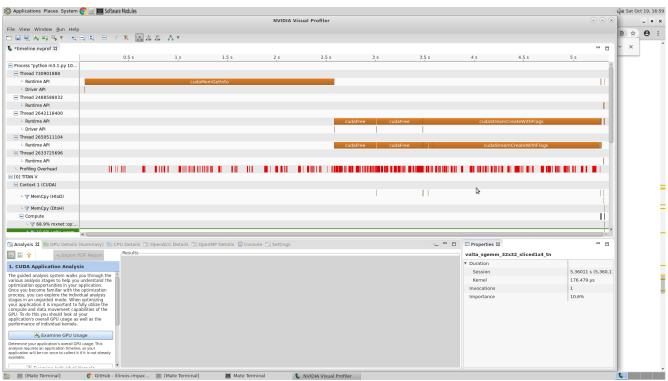
output of rai running MXNet on the GPU:

```
* Running /usr/bin/time python m1.2.py
Loading fashion-mnist data... done
Loading model... done
New Inference
EvalMetric: {'accuracy': 0.8154}
```

program runtime:

5.05user 3.23system 0:04.76elapsed 174%CPU (0avgtext+0avgdata 2963832maxresident)k

NVVP results



1 CPU implementation

Whole program execution time:

105.64user 9.82system 1:34.79elapsed 121%CPU (0avgtext+0avgdata 6044512maxresident)k

Op times:

Op Time: 13.013842 Op Time: 77.417661