CSC 212: Data Structures and Abstractions Basic Sorting Algorithms

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Announcements

- Programming #2
 - ✓ posted
 - ✓ Gradescope active
 - focus on classes and dynamic arrays (std::vector)
- · If programming is still a significant issue ...
 - consider taking this class next semester and focus on addressing the issue

Looking for internships/jobs?

https://careers.google.com/how-we-hire/interview/

https://www.amazon.jobs/en/landing_pages/softwaredevelopment-topics

https://www.facebook.com/careers/life/preparing-for-your-software-engineering-interview-at-facebook/

Worst-case, Average-case, Best-case

Warming up: Analyze this code

```
unsigned int argmin(const std::vector<int> &values) {
    unsigned int length = values.size();
    assert(length > 0);
    unsigned int idx = 0;
    int current = values[0];
    for (unsigned int i = 1; i < length; i ++) {</pre>
        if (values[i] < current) {</pre>
            current = values[i];
             idx = i;
    return idx;
```

```
T(n) = ? based on number of comparisons
```

Warming up: Analyze this code

```
bool argk(const std::vector<int> &values, int k, unsigned int &idx) {
   unsigned int length = values.size();
   for (unsigned int i = 0 ; i < length ; i ++) {
      if (values[i] == k) {
        idx = i;
        return true;
      }
   }
   return false;
}</pre>
```

```
T(n) = ? based on number of comparisons
```

Different types of analysis



Worst-case: maximum time of algorithm on any input

Average-case: expected time of algorithm over all inputs

Best-case: minimum time of algorithm on some (optimal) input

Different types of analysis

- While **asymptotic analysis** describes T(n) as n approaches infinity ...
 - ✓ asymptotic notation: big O, big Omega, big Theta
- Case analysis looks into the different input types
 - √ best-case, worst-case, average-case

Both analysis types are orthogonal to each other

Worst-case, Average-case, Best-case

• Ex: factorial of a number (iterative algorithm)

• Ex: sequential search (return first occurrence)

• Ex: sequential search (return last occurrence)

Basic Sorting Algorithms

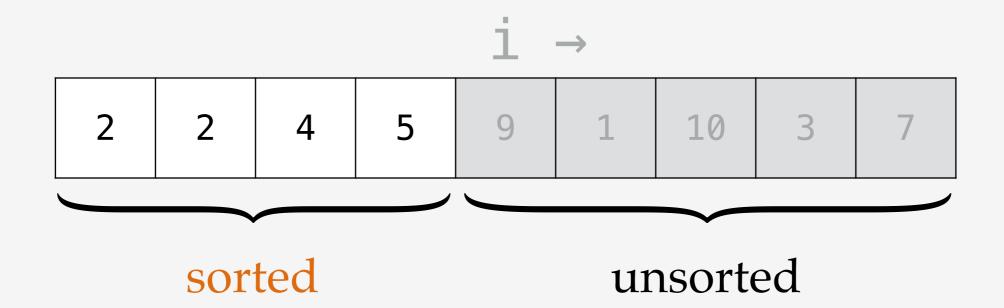
Sorting

- Given **n** elements that can be compared according to a **total order** relation
 - we want to rearrange them in non-increasing/non-decreasing order
 - ✓ for example (non-decreasing):
 - input: sequence of items $A = [k_0, k_1, ..., k_{n-1}]$
 - **output:** permutation of A $B \mid B[0] \leq B[1] \leq ...B[n-1]$

Central problem in computer science

Insertion Sort

- Array is divided into sorted and unsorted parts
 - √ algorithm scans array from left to right
- Invariants
 - elements in sorted are in ascending order
 - delements in **unsorted** have not been seen



Insertion Sort Demo

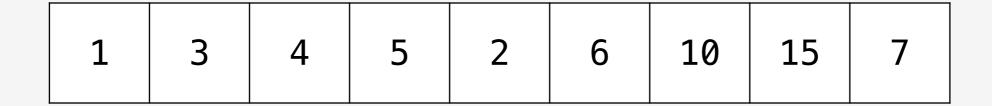
```
void insertionsort(int *A, unsigned int n) {
    int temp;
    unsigned int i, j;
    // grows the left part (sorted)
    for (i = 0 ; i < n ; i ++) {
        // inserts A[j] in sorted part
        for (j = i ; j > 0 ; j --) {
            if (A[j] < A[j-1]) {
                temp = A[j];
                A[j] = A[j-1];
                A[j-1] = temp;
            else
                break:
```

Analysis — Insertion Sort (comparisons)

- Running time depends on the input
- Worst-case?
 - √ input reverse sorted
- Best-case?
 - √ input already sorted
- Average-case?
 - \checkmark expect every element to move O(n/2) times

Partially sorted arrays

An **inversion** is a pair of keys that are out of order



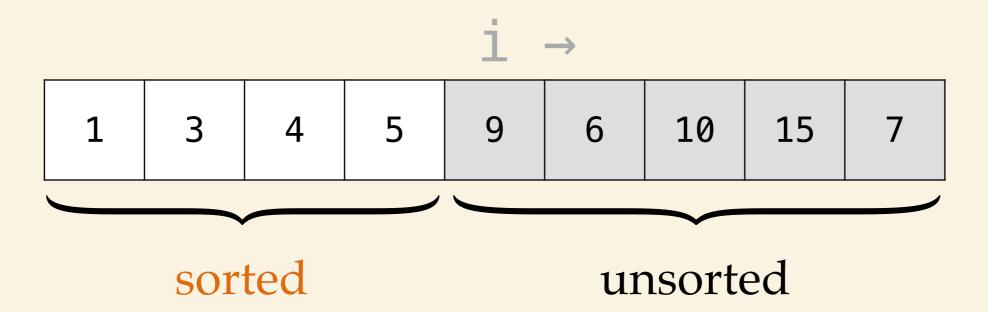
"array is partially sorted if the number of pairs that are out-of-order is O(n)"

For partially-sorted arrays, insertion sort runs in linear time.

$$\Theta(n)$$

Selection Sort

- Array is divided into sorted and unsorted parts
 - √ algorithm scans array from left to right
- Invariants
 - delements in sorted are fixed and in ascending order
 - no element in unsorted is smaller than any element in sorted



Selection Sort Demo

```
void selectionsort(int *A, unsigned int n) {
    int temp;
    unsigned int i, j, min;
    // grows the left part (sorted)
    for (i = 0 ; i < n ; i ++) {
        min = i;
        // find min in unsorted part
        for (j = i+1; j < n; j ++) {
            if (A[j] < A[min]) {</pre>
                min = j;
        // swap A[i] and A[min]
        temp = A[i];
        A[i] = A[min];
        A[min] = temp;
```

Summary

	Best-Case	Average-Case	Worst-Case
Selection Sort	θ(n²)	θ(n²)	θ(n²)
Insertion Sort	θ(n)	θ(n²)	θ(n²)