

# CSC 212: Data Structures and Abstractions

## Merge Sort

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# Announcements

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- Assignment 3 due October 16th
- MEC Project due
  - October 25th

# Divide and Conquer

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- **Divide** the problem into smaller subproblems
- **Conquer** recursively
  - ✓ ... each subproblem
- **Combine** Solutions

# Example

10	2	3	7	4	13	11	9
----	---	---	---	---	----	----	---

- ✓ sorting with insertion sort is  $(n^2)$
- ✓ we can divide the array into two halves and sort them separately

2	3	7	10
---	---	---	----

4	9	11	13
---	---	----	----

- ✓ each subproblem could be sorted in  $\approx n^2/4$
- ✓ sorting both halves will require  $\approx 2(n^2/4)$  🤔
- ✓ we need an additional operation to combine both solutions

Time “reduced” from  $\approx n^2$  to  $\approx n^2/2 + n$

# Merge Sort

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- **Divide** the array into **two** halves
  - ✓ just need to calculate the mid point
- Conquer **Recursively** each half
  - ✓ call Merge Sort on each half (i.e. solve 2 smaller problems)
- **Merge** Solutions
  - ✓ after both calls are finished, proceed to **merge** the solutions

# Divide and Conquer

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0	1	2	3	4	5	6	7
4	2	1	5	6	3	7	8



# Merge Sort: pseudocode

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```
if (hi <= lo) return;
```

```
int mid = lo + (hi - lo) / 2;
```

```
mergesort(A, lo, mid);  
mergesort(A, mid+1, hi);
```

```
merge(A, lo, mid, hi);
```

# Merge Sort

---

```
1 void r_mergesort(int *A, int *aux, int lo, int hi) {
2     // base case (single element or empty list)
3     if (lo == hi) return;
4     if (div == lo) return;
5     int mid = lo + (hi - lo) / 2;
6     // recursively sort halves
7     r_mergesort(A, aux, lo, mid);
8     r_mergesort(A, aux, mid+1, hi);
9     // merge results
10    merge(A, aux, lo, mid, hi);
11 }
```

```
1 void mergesort(int *A, int n) {
2     int *aux = new int[n];
3     r_mergesort(A, aux, 0, n-1);
4     delete [] aux;
5 }
```



# Merging two sorted arrays

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1	2	5	10	15
---	---	---	----	----

*i*

3	6	9	16	20
---	---	---	----	----

*j*

1	2	3	5	6	9	10	15	16	20
---	---	---	---	---	---	----	----	----	----

A secondary array is necessary to guarantee a **lineartime** operation

# Merge

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```
void merge(int *A, int *aux, int lo, int mid, int hi) {  
    // copy array  
    std::memcpy(aux+lo, A+lo, (hi-lo+1) * sizeof(int));  
    // merge  
    int i = lo, j = mid + 1;  
    for (int k=lo; k<=hi;k++) { if (i  
        > mid) A[k] =aux[j++]; else  
        if (j>hi)A[k]=aux[i++];  
        else if (aux[j] < aux[i]) A[k] = aux[j++];  
        else A[k] = aux[i++];  
    }  
}
```

# Analysis (recurrence)

---

```
if (hi <= lo) return;
```

```
int mid = lo + (hi - lo) / 2;
```

```
mergesort(A, lo, mid);
```

```
mergesort(A, mid+1, hi);
```

```
merge(A, lo, mid, hi);
```

Worst case?

Average case?

Best case?

# Recursion Tree (trace)

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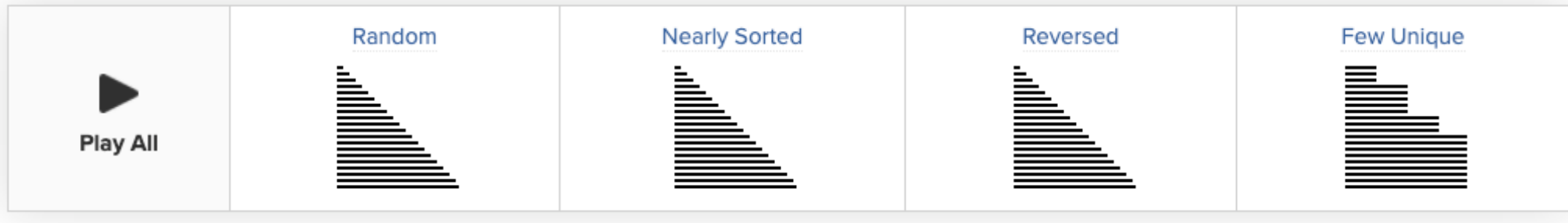
```
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    int *aux = new int[n];  
    r_mergesort(A, aux, 0, n-1);  
    delete [] aux;  
}
```

```
void r_mergesort(int *A, int *aux, int lo, int hi) {  
    if (hi <= lo) return;  
    int mid = lo + (hi - lo) / 2;  
    r_mergesort(A, aux, lo, mid);  
    r_mergesort(A, aux, mid+1, hi);  
    merge(A, aux, lo, mid, hi);  
}
```

# Animation

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<https://www.toptal.com/developers/sorting-algorithms/merge-sort>



# Comments on Merge Sort

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- Major disadvantage

- ✓ it is not **in-place**
- ✓ in-place algorithm exists but it is complex and inefficient

- Improvements

- ✓ use insertion sort for small arrays
  - avoid overhead on small instances (~10 elements)
- ✓ stop if already sorted
  - avoids unnecessary merge
  - works well with partially sorted arrays

# In-place Sorting

# Example

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- Think about reversing an array or string
  - ✓ **solution 1:** use an additional array of equal size
    - what is the required extra memory?
  - ✓ **solution 2:** exchange first and last and work recursively on the inner part
    - can do it iteratively as well
    - what is the required extra memory?



# In-place sorting

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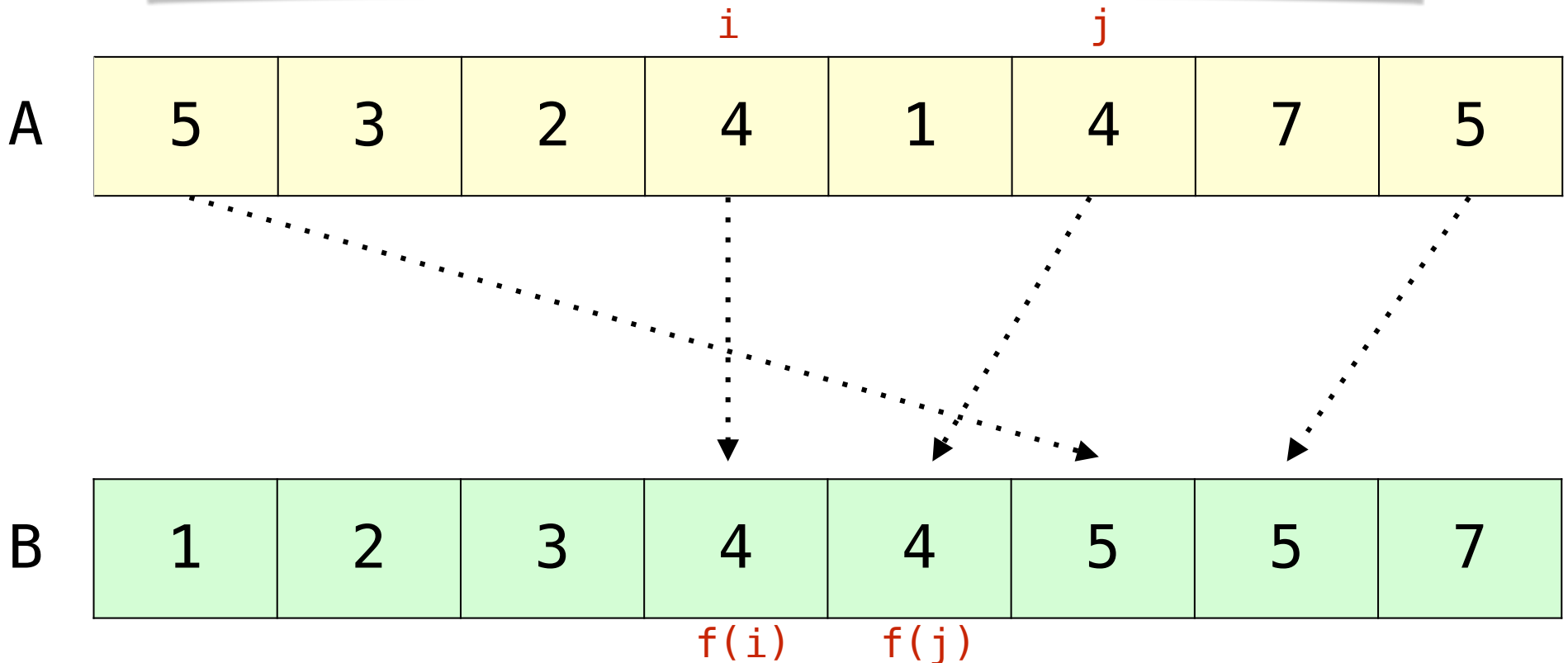
- A sorting algorithm is **in-place** if it uses  **$O(\log n)$**  extra memory
- Are selection and insertion sorts **in-place**?

# Stable Sorting

# Stability

- A sorting algorithm is **stable** if it preserves the order of **equal** elements

Consider sorting (in ascending order) a list **A** into a sorted list **B**. Let  $f(i)$  be the index of element  $A[i]$  in **B**. The sorting algorithm is stable if for any pair  $(i,j)$  such that  $A[i] = A[j]$  and  $i < j$ , then  $f(i) < f(j)$ .



DL 2273	Detroit	5:30 am	Departed
WN 6240	Chicago - MDW	5:55 am	Departed
AA 489	Philadelphia	6:00 am	Departed
DL 1263	Atlanta	6:00 am	Departed
UA 6208	Washington - IAD	6:00 am	Departed
WN 1138	Baltimore	6:05 am	Departed
AA 5202	Washington - DCA	6:14 am	Departed
B6 475	Orlando	6:15 am	Departed
UA 4894	New York/Newark	6:15 am	Departed
AA 1703	Charlotte	6:17 am	Departed
WN 28	Orlando	6:55 am	Departed
AA 3410	Chicago - ORD	7:02 am	Departed
WN 6235	Tampa	7:05 am	Departed
UA 3615	Chicago - ORD	7:30 am	Departed
AA 1735	Philadelphia	8:02 am	Departed
AA 632	Charlotte	8:07 am	At 9:45 am
WN 6247	Fort Lauderdale	8:30 am	Departed
WN 2640	Washington - DCA	8:45 am	Departed
WN 3420	Chicago - MDW	8:45 am	Departed
AA 4280	Washington - DCA	8:49 am	At 10:20 am
WN 846	Baltimore	9:20 am	Departed
DL 305	Detroit	10:40 am	On time
AA 774	Philadelphia	10:51 am	On time
AA 1981	Charlotte	11:01 am	On time
WN 3020	Baltimore	11:20 am	On time
AA 5524	Washington - DCA	11:46 am	At 2:35 pm
AC 7379	Toronto	11:50 am	On time
AA 5550	Charlotte	11:54 am	On time
DL 5090	Detroit	12:32 pm	On time
WN 6296	Baltimore	12:35 pm	On time
DL 2225	Atlanta	12:48 pm	On time
AA 4424	Washington - DCA	1:38 pm	On time

sort  
then  
sort  
again

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DL 2225	Atlanta	12:48 pm	On time
WN 1138	Baltimore	6:05 am	Departed
WN 846	Baltimore	9:20 am	Departed
WN 3020	Baltimore	11:20 am	On time
WN 6296	Baltimore	12:35 pm	On time
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AA 3410	Chicago - ORD	7:02 am	Departed
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# Stability

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› Is **selection sort** stable?



✓ long distance swaps

✓ try: 5 1 2 4 4 3 2 1

› Is **insertion sort** stable?



✓ equal items never pass each other (depends on correct implementation)

# Sorting Algorithms

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	Best-Case	Average-Case	Worst-Case	Stable?	In-place?
Selection Sort	$\theta(n^2)$	$\theta(n^2)$	$\theta(n^2)$	No	Yes
Insertion Sort	$\theta(n)$	$\theta(n^2)$	$\theta(n^2)$	Yes	Yes
Merge Sort	$\theta(n \log n)$	$\theta(n \log n)$	$\theta(n \log n)$	Yes	No