

CSC 212: Data Structures and Abstractions

Merge Sort

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Announcements

- Assignment 3 due March 11th
- Midterm Exam
 - ✓ March 10th
 - ✓ Look at syllabus for a topic list
- Midterm Exam Review
 - ✓ Live Coding Session March 9th
 - ✓ Possibly 8th as well, keep an eye on EdStem

Divide and Conquer

- **Divide** the problem into smaller subproblems
- **Conquer** recursively
 - ✓ ... each subproblem
- **Combine** Solutions

Example

10	2	3	7	4	13	11	9
----	---	---	---	---	----	----	---

- ✓ sorting with insertion sort is $\Theta(n^2)$
- ✓ we can divide the array into two halves and sort them separately

2	3	7	10
---	---	---	----

4	9	11	13
---	---	----	----

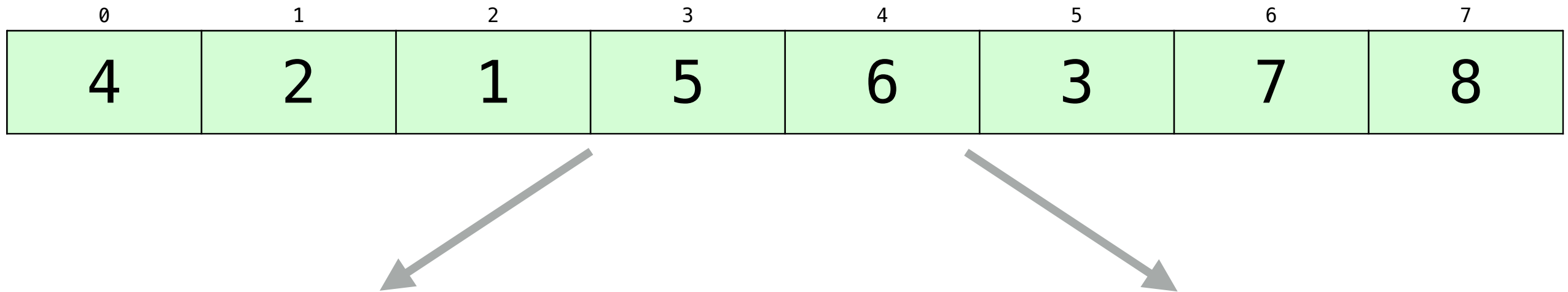
- ✓ each subproblem could be sorted in $\approx n^2/4$
- ✓ sorting both halves will require $\approx 2(n^2/4)$ 🤔
- ✓ we need an additional operation to combine both solutions

Time “reduced” from $\approx n^2$ to $\approx n^2/2 + n$

Merge Sort

- **Divide** the array into **two** halves
 - ✓ just need to calculate the mid point
- Conquer **Recursively** each half
 - ✓ call Merge Sort on each half (i.e. solve 2 smaller problems)
- **Merge** Solutions
 - ✓ after both calls are finished, proceed to **merge** the solutions

Divide and Conquer



Merge Sort: pseudocode

```
if (hi <= lo) return;
```

```
int mid = lo + (hi - lo) / 2;
```

```
mergesort(A, lo, mid);  
mergesort(A, mid+1, hi);
```

```
merge(A, lo, mid, hi);
```

Merge Sort

```
1 void r_mergesort(int *A, int *aux, int lo, int hi) {
2     // base case (single element or empty list)
3     if (hi <= lo) return;
4     // divide
5     int mid = lo + (hi - lo) / 2;
6     // recursively sort halves
7     r_mergesort(A, aux, lo, mid);
8     r_mergesort(A, aux, mid+1, hi);
9     // merge results
10    merge(A, aux, lo, mid, hi);
11 }
```

```
1 void mergesort(int *A, int n) {
2     int *aux = new int[n];
3     r_mergesort(A, aux, 0, n-1);
4     delete [] aux;
5 }
```


Merging two sorted arrays

1	2	5	10	15
---	---	---	----	----

i

3	6	9	16	20
---	---	---	----	----

j

1	2	3	5	6	9	10	15	16	20
---	---	---	---	---	---	----	----	----	----

A secondary array is necessary to guarantee a **linear time** operation

Merge

```
void merge(int *A, int *aux, int lo, int mid, int hi) {  
    // copy array  
    std::memcpy(aux+lo, A+lo, (hi-lo+1) * sizeof(int));  
    // merge  
    int i = lo, j = mid + 1;  
    for (int k = lo ; k <= hi ; k++) {  
        if (i > mid) A[k] = aux[j++];  
        else if (j > hi) A[k] = aux[i++];  
        else if (aux[j] < aux[i]) A[k] = aux[j++];  
        else A[k] = aux[i++];  
    }  
}
```

Analysis (recurrence)

```
if (hi <= lo) return;
```

```
int mid = lo + (hi - lo) / 2;
```

```
mergesort(A, lo, mid);
```

```
mergesort(A, mid+1, hi);
```

```
merge(A, lo, mid, hi);
```

Worst case?

Average case?

Best case?

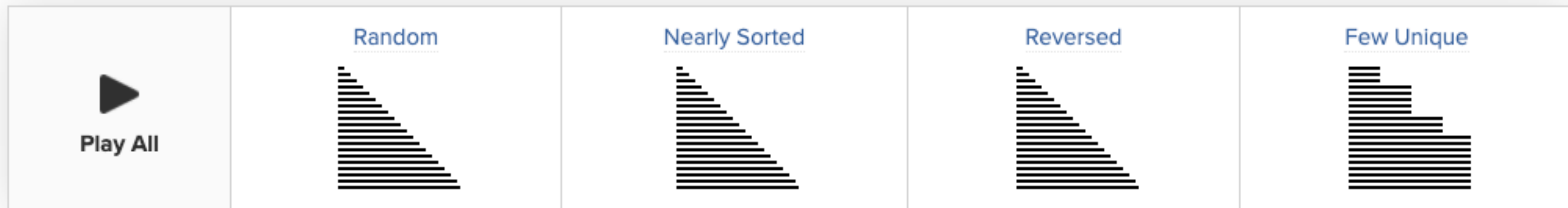
Recursion Tree (trace)

```
void mergesort(int *A, int n) {  
    int *aux = new int[n];  
    r_mergesort(A, aux, 0, n-1);  
    delete [] aux;  
}
```

```
void r_mergesort(int *A, int *aux, int lo, int hi) {  
    if (hi <= lo) return;  
    int mid = lo + (hi - lo) / 2;  
    r_mergesort(A, aux, lo, mid);  
    r_mergesort(A, aux, mid+1, hi);  
    merge(A, aux, lo, mid, hi);  
}
```

Animation

<https://www.toptal.com/developers/sorting-algorithms/merge-sort>



Comments on Merge Sort

- Major disadvantage

- ✓ it is not **in-place**
- ✓ in-place algorithm exists but it is complex and inefficient

- Improvements

- ✓ use insertion sort for small arrays
 - avoid overhead on small instances (~10 elements)
- ✓ stop if already sorted
 - avoids unnecessary merge
 - works well with partially sorted arrays

In-place Sorting

Example

- Think about reversing an array or string
 - ✓ **solution 1:** use an additional array of equal size
 - what is the required extra memory?
 - ✓ **solution 2:** exchange first and last and work recursively on the inner part
 - can do it iteratively as well
 - what is the required extra memory?

In-place sorting

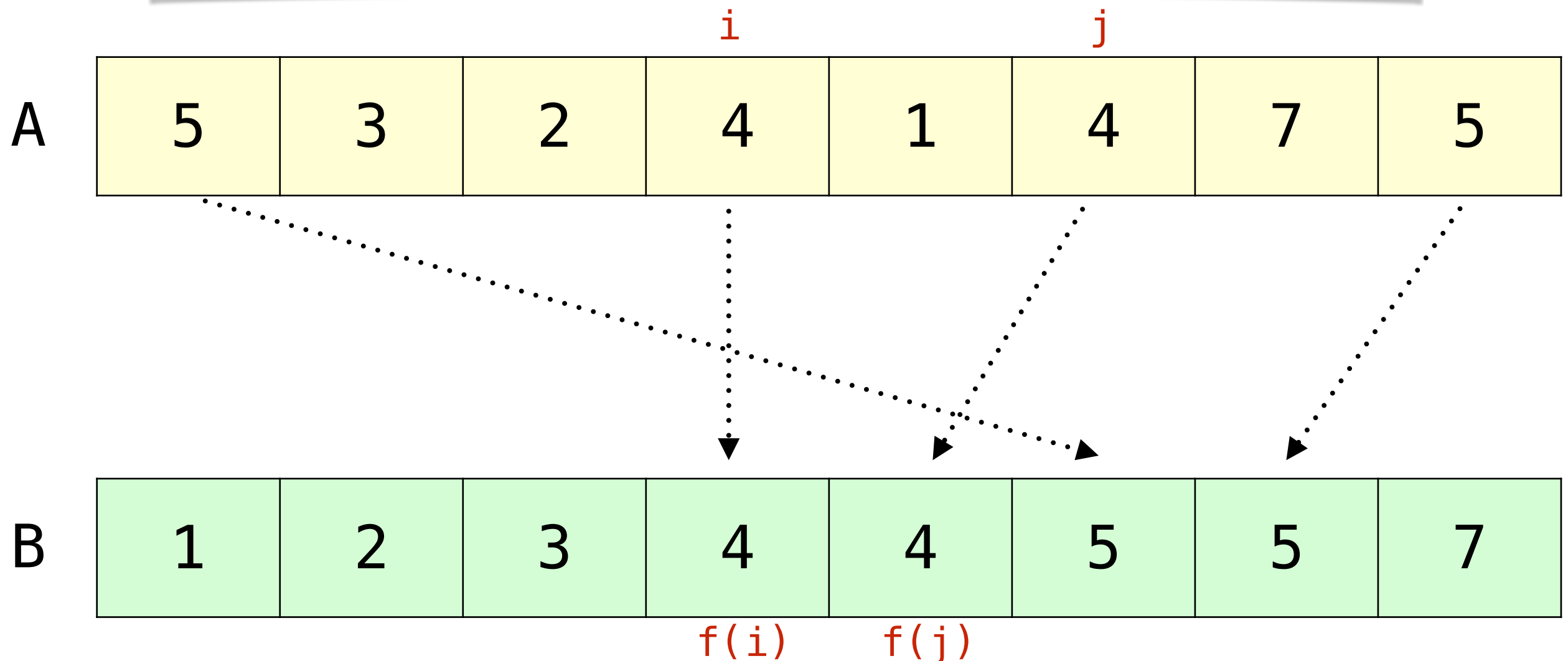
- A sorting algorithm is **in-place** if it uses **$O(\log n)$** extra memory
- Are selection and insertion sorts **in-place**?

Stable Sorting

Stability

- A sorting algorithm is **stable** if it preserves the order of **equal** elements

Consider sorting (in ascending order) a list **A** into a sorted list **B**. Let $f(i)$ be the index of element $A[i]$ in **B**. The sorting algorithm is stable if for any pair (i,j) such that $A[i] = A[j]$ and $i < j$, then $f(i) < f(j)$.



DL 2273	Detroit	5:30 am	Departed
WN 6240	Chicago - MDW	5:55 am	Departed
AA 489	Philadelphia	6:00 am	Departed
DL 1263	Atlanta	6:00 am	Departed
UA 6208	Washington - IAD	6:00 am	Departed
WN 1138	Baltimore	6:05 am	Departed
AA 5202	Washington - DCA	6:14 am	Departed
B6 475	Orlando	6:15 am	Departed
UA 4894	New York/Newark	6:15 am	Departed
AA 1703	Charlotte	6:17 am	Departed
WN 28	Orlando	6:55 am	Departed
AA 3410	Chicago - ORD	7:02 am	Departed
WN 6235	Tampa	7:05 am	Departed
UA 3615	Chicago - ORD	7:30 am	Departed
AA 1735	Philadelphia	8:02 am	Departed
AA 632	Charlotte	8:07 am	At 9:45 am
WN 6247	Fort Lauderdale	8:30 am	Departed
WN 2640	Washington - DCA	8:45 am	Departed
WN 3420	Chicago - MDW	8:45 am	Departed
AA 4280	Washington - DCA	8:49 am	At 10:20 am
WN 846	Baltimore	9:20 am	Departed
DL 305	Detroit	10:40 am	On time
AA 774	Philadelphia	10:51 am	On time
AA 1981	Charlotte	11:01 am	On time
WN 3020	Baltimore	11:20 am	On time
AA 5524	Washington - DCA	11:46 am	At 2:35 pm
AC 7379	Toronto	11:50 am	On time
AA 5550	Charlotte	11:54 am	On time
DL 5090	Detroit	12:32 pm	On time
WN 6296	Baltimore	12:35 pm	On time
DL 2225	Atlanta	12:48 pm	On time
AA 4424	Washington - DCA	1:38 pm	On time

sort
then
sort
again

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Stability

▸ Is **selection sort** stable? 

- ✓ long distance swaps

- ✓ try: 5 1 2 4 4 3 2 1

▸ Is **insertion sort** stable? 

- ✓ equal items never pass each other (depends on correct implementation)

Sorting Algorithms

	Best-Case	Average-Case	Worst-Case	Stable?	In-place?
Selection Sort	$\theta(n^2)$	$\theta(n^2)$	$\theta(n^2)$	No	Yes
Insertion Sort	$\theta(n)$	$\theta(n^2)$	$\theta(n^2)$	Yes	Yes
Merge Sort	$\theta(n \log n)$	$\theta(n \log n)$	$\theta(n \log n)$	Yes	No