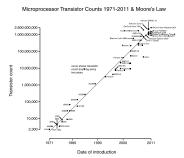
COMP20200 Unix Programming Lecture 11

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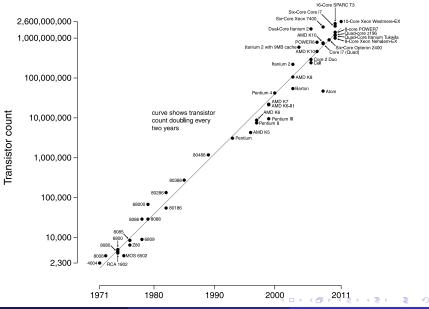
Moors law

- Gordon E. Moore in 1965 predicted "the number of transistors on integrated circuits doubles approximately every two years."
- Adjusted to doubling performance every 18 months.



• Need parallelism to use all these transistors: multi-core, many-core.

Microprocessor Transistor Counts 1971-2011 & Moore's Law



Parallelism

- Blocking IO
 - writing to disk or network
 - Could wait, asynchronous IO or spawn a separate thread
- Server
 - receiving multiple clients requests (eg. web server)
- User Interface
 - Respond to user input while program engine busy
- Performance on machines
 - Desktop, mobile, supercomputer clusters.
 - SMPs (shared memory multiprocessor), GPU-accelerators, clusters
 - Program must be parallel for performance.
 - SIMD (single instruction, multiple data)

Processes and Threads

- fork() returns two completely independent copies of the original process.
 - Each process has own address space, own copies of variables independent of same variables in other process.
 - provides memory protection and therefore stability
 - problem when multiple processes working on same task/problem
- Light-weight process (LWP): POSIX threads or pthreads
 - set of C programming language types, functions and constants
 - Around 100 Pthreads procedures, all prefixed "pthread_"

POSIX threads

```
#include <pthread.h>
#include <stdio.h>
#include <stdlib.h>
#include <assert.h>
#define NUM_THREADS
void *TaskCode(void *argument) {
   int tid;
   tid = *((int *) argument);
   printf("Hello World! It's me, thread %d!\n", tid);
   /* optionally: insert more useful stuff here */
   return NULL:
```

```
int main(void) {
   pthread_t threads[NUM_THREADS];
   int thread_args[NUM_THREADS];
   int rc, i;
   /* create all threads */
   for (i=0; i< NUM\_THREADS; ++i) {
      thread_args[i] = i;
      printf("In main: creating thread %d\n", i);
      rc = pthread_create(&threads[i], NULL, TaskCode,
                         (void *)&thread_args[i]);
      assert(0 = rc);
   /* wait for all threads to complete */
   for (i=0; i<NUM\_THREADS; ++i) {
      rc = pthread_join(threads[i], NULL);
      assert(0 = rc);
   exit (EXIT_SUCCESS);
```

Pthreads

When compiling with gcc, must add -pthread flag.

```
\$ gcc -Wall -o pthread_example pthread_example.c -pthread \$
```

Output:

```
In main: creating thread 0
In main: creating thread 1
In main: creating thread 2
Hello World! It's me, thread 1!
Hello World! It's me, thread 0!
In main: creating thread 3
In main: creating thread 4
Hello World! It's me, thread 2!
Hello World! It's me, thread 4!
Hello World! It's me, thread 3!
```

Race conditions

```
THREAD 1 THREAD 2

a = data; b = data;

a++; b--;

data = a: data = b:
```

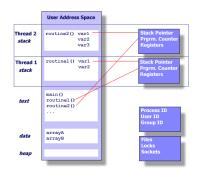
- If above executed serially, no problem, expected result.
- If in parallel it is completely non-deterministic. data may be +1, 0, -1.
- Solution: block other threads when writing to data
- Pthreads use a data type called mutex to achieve this.

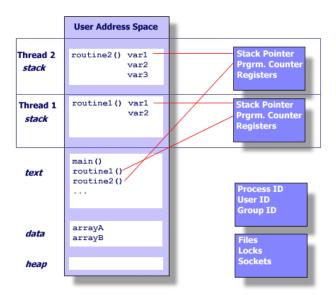
mutex example

```
THREAD 1
                                  THREAD 2
pthread_mutex_lock (&mut);
                                  pthread_mutex_lock (&mut);
                                  /* blocked */
a = data;
                                  /* blocked */
a++:
                                  /* blocked */
data = a;
                                  /* blocked */
pthread_mutex_unlock (&mut);
                                  b = data:
                                  b--;
                                  data = b:
                                  pthread_mutex_unlock (&mut)
```

POSIX threads

- Independent stream of instructions
- High performance on SMP
- Duplicate only the bare essential resources
- Accomplished by the thread maintaining its own:
 - Stack pointer
 - Registers
 - Scheduling properties (such as policy or priority)
 - Set of pending and blocked signals
 - Thread specific data.





In UNIX environment, a thread:

- Exists within a process and uses the process resources
- Has its own independent flow of control as long as its parent process exists
- Duplicates only the essential resources it needs to be independently schedulable
- May share the process resources with other threads that act equally independently (and dependently)
- Dies if the parent process dies
- Is "lightweight" because most of the overhead has already been accomplished through the creation of its process.

Threads share the process resources

Because threads within the same process share resources:

- Changes made by one thread to shared resources will be seen by all threads.
 - eg. closing a file
- Two pointers having the same value point to the same data.
- Reading and writing to the same memory locations is possible
 - but requires explicit synchronization by the programmer.