

README.pdf
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CSE241 lab 2

Part 1:

Initially, I ensured that $m = \text{maxSize}$ is prime.

For calculating $h2$, I used `int h2 = (int) Math.ceil((m-1)*(k*A2-Math.floor(k*A2)));`
Multiplying $(k*A2 - \text{Math.floor}(k*A2))$ by $(m-1)$ gives a double between 0 and $m-1$, not inclusive.
The outer `Math.ceil` operation then gives an integer between 1 and $m-1$, inclusive.

These two aspects (multiplying by $(m-1)$ instead of m , and using `Math.ceil` instead of `Math.floor`) were necessary to avoid having $h2$ equal 0 or m , which would prevent my methods from checking the entire table.

Later when I implemented the dynamically growing table that grew by powers of two this boundary was no longer needed, and I just needed to ensure that $h2$ is odd. I added this check that $h2$ is odd. However I left the $h2$ formula untouched because it worked like that and I didn't want to risk breaking anything.

Part 2:

To also store the `toHashCode` of a `Record`, I added an instance variable and a method to modify it in the `Record` class.

Having the same hash values does not necessarily mean that the strings are the same, as there may be rare cases when two strings generate the same hash value. Thus, it is still necessary to check that the strings are the same due to these rare cases.

To implement the doubling procedure, I added the instance variable `n` to the `StringTable` class, to keep track of how full the `StringTable` was. `n` was incremented when a record was inserted into a null slot. When a new record was to be inserted, if $(n * 4 > \text{the current table size})$, I called `doubleSize()`. This created a new table that was double the original size, re-set `n` to 0, and re-hashed all the current records into the new table.