

Estufa en Piloto



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1 General

1.1 Plantilla

```

1  /*  AUTHOR: Estufa en Piloto  */
2  #include <bits/stdc++.h>
3  using namespace std;
4
5  #ifdef LOCAL
6  #define DBG(x) cerr << #x << " = " << (x) << endl
7  #define RAYA cerr << "===== " << endl
8  #else
9  #define DBG(x)
10 #define RAYA
11 #endif
12
13 typedef long long ll;
14 typedef vector<ll> vi; typedef pair<ll,ll> ii;
15 typedef vector<ii> vii; typedef vector<bool> vb;
16 #define FIN ios::sync_with_stdio(0);cin.tie(0);cout.tie(0)
17 #define forr(i, a, b) for(ll i = (a); i < (b); i++)
18 #define forn(i, n) forr(i, 0, n)
19 #define SZ(x) int((x).size())
20 #define pb push_back
21 #define mp make_pair
22 #define all(c) (c).begin(),(c).end()
23 #define esta(x,c) ((c).find(x) != (c).end())
24 const int INF = 1<<30; // const ll INF = (1LL<<60);
25 const int MOD = 1e9+7; // const int MOD = 998244353;
26 const int MAXN = 2e5+5;
27
28 int main(){
29     FIN;
30
31
32
33     return 0;
34 }

```

1.2 Compilar

```

1  # Llenar lo siguiente en el build de geany
2

```

```

3 # Compile
4 g++ -std=c++17 -DLOCAL -g -O2 -Wconversion -Wshadow -Wall -Wextra -
    D_GLIBCXX_DEBUG -c %f
5
6 # Build
7 g++ -std=c++17 -DLOCAL -g -O2 -Wconversion -Wshadow -Wall -Wextra -
    D_GLIBCXX_DEBUG -o %e %f

```

1.3 Informacion

```

1
2 DATA STRUCTURE - C++
3
4 ----- VECTOR -----
5
6
7 vector<ll> v; //DECLARACION DEL VECTOR, TAMANO 0.
8
9 vector<ll> v(n); //DECLARACION DEL VECTOR DE TAMANO N
10
11 v.push_back(n); //PUSHEA N AL FINAL DEL VECTOR, O(1)
12
13 sort(v.begin(),v.end()); //ORDENA EL VECTOR, O(n.logn)
14
15 v.pop_back(); //ELIMINA EL ULTIMO VALOR DEL VECTOR, O(1);
16
17 v.size(); //RETORNA EL TAMANO DEL VECTOR
18
19 v.clear(); //BORRA TODO EL VECTOR
20
21 v.resize(n); //ESTABLECEMOS EL TAMANO DEL VECTOR A TAMANO N
22
23 vector<pair<ll,ll>> v; //VECTOR DE PAIR
24
25 v.push_back(make_pair(a,b)); //PUSHEA EL PAR (A,B).
26
27 v[i].first; //RETORNA EL PRIMER VALOR DEL PAIR
28
29 v[i].second; //RETORNA EL SEGUNDO VALOR DEL PAIR
30
31 for(auto u : v) //RECORRE TODOS LOS VALORES DEL VECTOR
32
33 BUENO PARA STACKEAR COSAS, PARA USARLO COMO LISTA DE ADYACENCIAS

```

```

34 EN LOS PROBLEMAS DE GRAFOS, PARA USARLOS COMO ARREGLO.
35 SI SE ORDENA UN VECTOR DE PAIR, LOS ORDENA DE MENOR A MAYOR SEGUN
36 EL PRIMER VALOR, Y LUEGO DE MENOR A MAYOR SEGUN EL SEGUNDO.
37
38 -----
39
40 ----- SET -----
41
42
43 set<int> s; //DECLARACION DEL SET, EL SET GUARDA LOS ELEMENTOS UNA SOLA
44 //VEZ, Y MIENTRAS LOS VA GUARDANDO, LOS VA ORDENANDO.
45
46 multiset<int> s; //DECLARACION DEL MULTISSET. EL MULTISSET PUEDE GUARDAR
47 //VARIAS COPIAS DEL MISMO NUMERO, Y MIENTRAS LOS VA
48 //GUARDANDO, LOS VA ORDENANDO.
49
50 s.insert(n); //INSERTA EL VALOR N EN EL SET, SI YA ESTABA INSERTADO, NO
51 //NO SE MODIFICA NADA. O(logn)
52
53 s.erase(n); //ELIMINA EL VALOR N DEL SET, O(log n)
54
55 s.size(); //TAMANO DEL SET
56
57 s.clear(); //BORRA TODOS LOS ELEMENTOS DEL SET O(s.size())
58
59 auto it=s.begin(); //DECLARACION DEL ITERATOR IT PARA EL SET
60
61 it=s.find(n); //HALLA LA POSICION DEL VALOR N DENTRO DEL SET, SI NO
62 //CREO QUE RETORNA s.end(); (FINAL DEL SET)
63
64 it.lower_bound(n); //HERRAMIENTA FUERTISIMA DEL SET. CON BUSQUENA
65 //BINARIA OBTIENE EL PRIMER VALOR DEL SET QUE ES MAYOR O IGUAL
66 //A N, O(logn)
67
68 it=upper_bound(n); //HERRAMIENTA TAMBIEN MUY FUERTE, PERO NO LA SUELO
69 //USAR MUCHO, SUELE PODER SER REEMPLAZADA POR EL LOWER_BOUND
70 //EN O(logn) DEVUELVE EL PRIMER VALOR MAYOR Estricto QUE
71 //N
72
73 for(auto it : s) //RECORRE TODOS LOS VALORES DEL SET
74
75 // Custom comparators
76 struct Edge {

```

```

77     int a, b, w;
78 };
79 struct cmp { // Funcion comparadora para el set
80     bool operator()(const Edge &x, const Edge &y) const { return x.w < y.w
81         ; }
82 };
83 int main() {
84     int M = 4;
85     set<Edge, cmp> v;
86     ...
87
88     // FACIL DE USAR. PUEDE REEMPLAZAR EL PRIORITY QUEUE PERO ES UN POQUITO
89     // MAS
90     // LENTO. BUENO PARA PROBLEMA QUE CONVIENE IR ORDENANDO EN TIEMPO REAL,
91     // EN VEZ DE LEER TODO Y LUEGO ORDENARLO. SE USA MUCHO EN ALGORITMOS
92     // PARA
93     // GRAFOS COMO DIJKSTRA Y PRIM.
94
95     -----
96     ----- QUEUE -----
97
98     queue<int> q; //DECLARACION DEL QUEUE
99
100    q.push(n); //PUSHEA AL FINAL DEL QUEUE EL NUMERO N
101
102    q.front(); //OBTIENE EL NUMERO QUE SE ENCUENTRA PRIMERO EN LA QUEUE
103
104    q.pop() //BORRA EL NUMERO QUE SE ENCUENTRA PRIMERO EN EL QUEUE
105
106    -----
107
108    ----- PRIORITY QUEUE -----
109
110    priority_queue<int> pq; //DECLARACION DEL PRIORITY_QUEUE. A MEDIDA QUE
111        //VAMOS INSERTANDO VALORES, LOS VA ORDENANDO EN
112        //ORDEN CRECIENTE.
113
114    priority_queue<int, vector<int>, greater<int>> pq; // PRIORITY_QUEUE DE
115        MINIMO

```

```

116    pq.push(n); //INSERTA EL VALOR N EN EL LUGAR QUE CORRESPONDE. O(log n)
117
118    pq.top(); //OBTIENE EL VALOR QUE SE ENCUENTRA PRIMERO EN LA PQ.
119
120    pq.pop(); //ELIMINA EL ELEMENTO QUE SE ENCUENTRA PRIMERO EN LA PQ.

```

1.4 Medir Tiempo

```

1 unsigned t0, t1;
2 t0=clock();
3
4 // Aca Va el codigo
5
6 t1 = clock();
7 double time = (double(t1-t0)/CLOCKS_PER_SEC);
8 cout << "Execution_Time:_" << time << endl;

```

1.5 Plantilla Old

```

1 #include <bits/stdc++.h>
2 using namespace std;
3 //freopen("input.txt", "r", stdin);
4 //freopen("output.txt", "w", stdout);
5
6 // neal Debugger
7 template<typename A, typename B> ostream& operator<<(ostream &os, const
8     pair<A, B> &p) { return os << '(' << p.first << ",_" << p.second <<
9     ')'; }
10
11 template<typename T_container, typename T = typename enable_if<!is_same<
12     T_container, string>::value, typename T_container::value_type>::type
13     > ostream& operator<<(ostream &os, const T_container &v) { os << '{'
14     ; string sep; for (const T &x : v) os << sep << x, sep = ",_";
15     return os << '}'; }
16
17 void dbg_out() { cerr << endl; }
18 template<typename Head, typename... Tail> void dbg_out(Head H, Tail... T
19     ) { cerr << '_' << H; dbg_out(T...); }
20
21 #ifdef LOCAL
22 #define dbg(...) cerr << "(" << #__VA_ARGS__ << "):", dbg_out(
23     __VA_ARGS__)
24 #else
25 #define dbg(...)
26 #endif

```

```

18
19 typedef long long ll;
20 #define FIN ios::sync_with_stdio(0);cin.tie(0);cout.tie(0)
21 #define forr(i, a, b) for(ll i = (a); i < (ll) (b); i++)
22 #define forn(i, n) forr(i, 0, n)
23 #define pb push_back
24 #define mp make_pair
25 #define all(c) (c).begin(),(c).end()
26 #define DBG(x) cerr << #x << "□=□" << (x) << endl
27 #define DBGV(v,n) forn(i,n) cout << v[i] << "□"; cout << endl
28 #define esta(x,c) ((c).find(x) != (c).end())
29 #define RAYA cerr << "===== " << endl
30 const ll MOD = (ll)(1e9+7); // 998244353
31 const ll INF = (ll)(1<<30); // (1LL<<60)
32 const int MAXN = (int)(2e5+5);
33
34
35 int main(){
36     FIN;
37
38
39     return 0;
40 }

```

2 Data Structures

2.1 Compresion De Coordenadas

```

1 // Compresion de coordenadas
2 // Complejidad: O(n log n) para construir, O(log n) para obtener
3 struct compresion {
4     vector<int> todos;
5     compresion(vector<int> v) {
6         todos = v; sort(all(todos));
7         todos.erase(unique(all(todos)),todos.end());
8     }
9     int obtener(int x) { // Esto se podria hacer tambien con un
10         unordered_map en O(1)
11         return (int)(lower_bound(all(todos),x)-todos.begin());
12     };
13 }

```

2.2 Dsu

```

1 // DSU struct con path compression y union por size
2 // Complejidad: O(ack(n)) por operacion, donde ack(n) es la funcion
3 // inversa de Ackermann, casi O(1)
4 struct DSU {
5     vi link, sz;
6
7     DSU(int tam) {
8         link.resize(tam+5), sz.resize(tam+5);
9         forn(i, tam+5) link[i] = i, sz[i] = 1;
10    }
11
12    ll find(ll x){ return link[x] = (link[x] == x ? x : find(link[x])); }
13
14    bool same(ll a, ll b) { return find(a) == find(b); }
15
16    void join(ll a, ll b) {
17        a = find(a), b = find(b);
18        if(a == b) return;
19        if(sz[a] < sz[b]) swap(a,b);
20        sz[a] += sz[b];
21        link[b] = a;
22    }
23 };

```

2.3 Fenwick

```

1 // Description: Fenwick Tree (BIT) for range queries and point updates
2 // Time: O(log n) for both queries and updates
3 // Usage: BIT bit(v); bit.query(l,r); bit.update(pos,val);
4 struct BIT {
5     vector<ll> prefix, a;
6     BIT(vector<ll> &v) {
7         int n = v.size(); prefix.resize(n+1); a = v;
8         vector<ll> aux(n+1,0);
9         forn(i,n) aux[i+1] = aux[i] + v[i];
10        forr(i,1,n+1) prefix[i] = aux[i] - aux[i - (i&(-i))];
11    }
12    ll query(int l, int r) { //[a,b] 0-indexed
13        ll ans = 0; r++;
14        while(r) ans += prefix[r], r -= r&(-r);
15        while(l) ans -= prefix[l], l -= l&(-l);
16        return ans;
17    }
18    void update(int pos, ll val) {
19        int i = pos + 1; ll upd = val - a[pos];
20        while(i < prefix.size()) prefix[i] += upd, i += i&(-i);
21        a[pos] = val;
22    }
23 };

```

2.4 Implicit Treap

```

1 //~ Puede realizar todas las siguientes operaciones en O(log N)
2
3 //~ Dividir el arreglo en dos
4 //~ Mergear dos arreglos en uno
5 //~ Insertar elemento en cualquier posicion
6 //~ Si queremos intertar 'x' en la posicion i-esima, spliteamos en
7 //~ dos arreglos L = [0,i-1] : R = [i,n]; y mergeamos T = (L,x) y
8 //~ T = (T,R)
9 //~ Eliminar elemento en cualquier posicion
10 //~ Si queremos eliminar el i-esimo elemento, spliteamos en tres
11 //~ arreglos L = [0,i-1], M = [i,i], R = [i+1,n] y mergeamos T = (L,R)
12 //~ Aplicar cualquier funcion de Segment Tree o Lazy
13 //~ Para realizar la query [l,r], tenemos que primero splitear y
14 //~ obtener el intervalo [l,r]. El valor de su raiz sera la respuesta.
15 //~ Luego mergeamos para volver a la normalidad.

```

```

16
17 typedef struct item *pitem;
18
19 struct item {
20     ll key, prior, cont, mini, sum;
21     bool rev; // (parameters for lazy prop)
22     pitem l, r;
23     item(ll key):
24         key(key),prior(rand()),cont(1),l(NULL),r(NULL),rev(0) {}
25 };
26
27 void push(pitem it) { //Lazy for reverse array
28     if(it && it->rev) {
29         swap(it->l,it->r);
30         if(it->l) it->l->rev ^= true;
31         if(it->r) it->r->rev ^= true;
32         it->rev = false;
33     }
34 }
35
36 int cont(pitem it) {return it ? it->cont : 0;}
37 ll mini(pitem it) {return it ? it->mini : 1<<30;}
38 ll sum(pitem it) {return it ? it->sum : 0LL;}
39
40 void upd_cont(pitem it){
41     if(it) {
42         it->cont = cont(it->l) + cont(it->r) + 1;
43         it->mini = min(it->key,min(mini(it->l),mini(it->r)));
44         it->sum = it->key + sum(it->l) + sum(it->r);
45     }
46 }
47
48 void merge(pitem &t, pitem l, pitem r) { //Merge l and r, new root t
49     push(l); push(r);
50     if(!l || !r) t = l ? l : r;
51     else if(l->prior > r->prior) merge(l->r,l->r,r), t=l;
52     else merge(r->l,l,r->l), t = r;
53     upd_cont(t);
54 }
55
56 void split(pitem t, pitem& l, pitem& r, int sz){ //sz:desired size of l
57     if(!t) { l = r = 0; return;}
58     push(t);

```

```

59     if(sz <= cont(t->l)) split(t->l,l,t->l,sz), r = t;
60     else split(t->r,t->r,r,sz-1-cont(t->l)), l = t;
61     upd_cont(t);
62 }
63
64 int find_min(pitem t) { //Devuelve la posicion del minimo
65     push(t);
66     if(t->mini == t->key) return cont(t->l);
67     else if(t->l && t->l->mini == t->mini) return find_min(t->l);
68     else return cont(t->l)+1+find_min(t->r);
69 }
70
71 void set_treap(pitem &root, int p, ll val) {
72     pitem aux1 = NULL, aux2 = NULL;
73     split(root,root,aux2,p+1); split(root,root,aux1,p);
74     aux1->key = val;
75     merge(root,root,aux1); merge(root,root,aux2);
76 }
77
78 ll get_sum(pitem &root, int l, int r) {
79     pitem aux1 = NULL, aux2 = NULL;
80     split(root,root,aux2,r+1); split(root,root,aux1,l);
81     ll ans = aux1->mini; // aux1->mini for minimum
82     merge(root,root,aux1); merge(root,root,aux2);
83     return ans;
84 }
85
86 void DFS(pitem t) { //Good for debug
87     DBG(t->key);
88     if(t->l != NULL) DBG(t->l->key); if(t->r != NULL) DBG(t->r->key);
89     RAYA;
90     if(t->l != NULL) DFS(t->l); if(t->r != NULL) DFS(t->r);
91 }

```

2.5 Indexed Set

```

1 #include <bits/stdc++.h>
2 #include <ext/pb_ds/assoc_container.hpp>
3 typedef long long ll;
4
5 using namespace __gnu_pbds;
6 using namespace std;
7

```

```

8 typedef tree<ll,null_type,less<ll>,rb_tree_tag,
9     tree_order_statistics_node_update> indexed_set;
10
11 int main() {
12     indexed_set s;
13     s.insert(2); s.insert(3); s.insert(5);
14     auto x=s.find_by_order(2); // retorna el valor en la posicion 2 (5)
15     // (se puede desreferenciar con *)
16     int pos=int(s.order_of_key(3)); // retorna la posicion que se
17     // encuentra el 3 (1)
18     return 0;
19 }

```

2.6 Indexed Set Casero

```

1 // Indexed Set (Treap) - O(log n) for both queries and updates
2 typedef struct item *pitem;
3
4 struct item {
5     ll key, prior, cont, mini, sum;
6     bool rev; // (parameters for lazy prop)
7     pitem l, r;
8     item(ll key):
9         key(key),prior(rand()),cont(1),l(NULL),r(NULL),rev(0) {}
10 };
11
12 void push(pitem it) { //Lazy for reverse array
13     if(it) {
14         if(it->rev) {
15             swap(it->l,it->r);
16             if(it->l)it->l->rev ^= true;
17             if(it->r)it->r->rev ^= true;
18             it->rev = false;
19         }
20     }
21 }
22
23 int cont(pitem it) {return it ? it->cont : 0;}
24 ll mini(pitem it) {return it ? it->mini : 1<<30;}
25 ll sum(pitem it) {return it ? it->sum : 0LL;}
26
27 void upd_cont(pitem it){
28     if(it) {

```



```

29     it->cont = cont(it->l) + cont(it->r) + 1;
30     it->mini = min(it->key,min(mini(it->l),mini(it->r)));
31     it->sum = it->key + sum(it->l) + sum(it->r);
32 }
33 }
34
35 void merge(pitem &t, pitem l, pitem r) { //Merge l and r, new root t
36     push(l); push(r);
37     if(!l || !r) t = l ? l : r;
38     else if(l->prior > r->prior) merge(l->r,l->r,r), t=l;
39     else merge(r->l,l,r->l), t = r;
40     upd_cont(t);
41 }
42
43 void split(pitem t, pitem& l, pitem& r, int sz){ //sz:desired size of l
44     if(!t) { l = r = 0; return;}
45     push(t);
46     if(sz <= cont(t->l)) split(t->l,l,t->l,sz), r = t;
47     else split(t->r,t->r,r,sz-1-cont(t->l)), l = t;
48     upd_cont(t);
49 }
50
51 int treap_lower_bound(pitem &t, ll val) { // retorna cantidad de
52     elementos menores estrictos
53     if(t == NULL) return 0;
54     if(t->key >= val) return treap_lower_bound(t->l,val);
55     return cont(t->l) + 1 + treap_lower_bound(t->r,val);
56 }
57
58 int find_by_order(pitem &t, int k) { // retorna el k-esimo elemento
59     if(t == NULL) return -1;
60     if(cont(t->l) + 1 == k) return t->key;
61     if(cont(t->l) + 1 < k) return find_by_order(t->r,k-cont(t->l)-1);
62     else return find_by_order(t->l,k);
63 }
64
65 void insert(pitem &t, ll val) {
66     int tam = treap_lower_bound(t,val);
67     pitem r = NULL;
68     split(t,t,r,tam);
69     merge(t,t,new item(val)); merge(t,t,r);
70 }

```

```

71 void erase(pitem &t, ll val) {
72     int tam = treap_lower_bound(t,val);
73     pitem r = NULL, aux = NULL;
74     split(t,t,r,tam+1); split(t,t,aux,tam);
75     merge(t,t,r);
76 }
77
78 void DFS(pitem t) { //Good for debug
79     if(t == NULL) return;
80     DFS(t->l); cerr << t->key << "\n"; DFS(t->r);
81 }

```

2.7 Linked List

```

1 // Linked list based on arrays
2 struct linked_list {
3     ll n;
4     vi l, r;
5     linked_list(int n): n(_n) {
6         l.resize(n); r.resize(n);
7         forn(i, n){
8             l[i] = (i-1+n)%n;
9             r[i] = (i+1)%n;
10        }
11    }
12    void erase(int pos){
13        r[l[pos]] = r[pos];
14        l[r[pos]] = l[pos];
15    }
16 };

```

2.8 Mo Dquery

```

1 // Mo's Algorithm for range queries
2 // Time: O(n*sqrt(n))
3 // Usage: mo mo; mo.read_queries(m); mo.get_mo(v);
4
5 struct query{
6     int l, r, ind;
7     bool operator <(query q) const {
8         return r < q.r;
9     }
10 };
11

```

```

12 struct mo {
13     const int block = 448; // sqrt(MAXN)
14     vector<vector<query>> q;
15     vector<ll> ans, cont;
16
17     mo(){
18         ans.resize(MAXN), cont.resize(MAXN);
19         q.resize(MAXN/block+5);
20     }
21
22     void read_queries(int m) { // m --> number of queries
23         ans.resize(m,0);
24         forn(i, m) {
25             int l, r; cin >> l >> r;
26             l--; r--; // For 0-indexed
27             q[l / block].pb({l,r,(int)i});
28         }
29     }
30
31     void add(ll &sum, ll val) {
32         if(cont[val] == 0) sum++;
33         cont[val]++;
34     }
35
36     void erase(ll &sum, ll val) {
37         if(cont[val] == 1) sum--;
38         cont[val]--;
39     }
40
41     void get_mo(vector<ll> &v) {
42         ll sum = 0;
43         forn(i,q.size()) {
44             sort(all(q[i]));
45             int l, r; l = r = block * i;
46             for(query u : q[i]) { // Solve for [l,r]
47                 while(r <= u.r) add(sum,v[r]), r++;
48                 while(l <= u.l) erase(sum,v[l]), l++;
49                 while(l > u.l) l--, add(sum,v[l]);
50                 ans[u.ind] = sum;
51             }
52             while(l < r) erase(sum,v[l]), l++;
53         }
54     }

```

```

55 };

```

2.9 Segtree 2D

```

1 struct st2d{ // 0 indexado, [l, r]
2     // PARA HACER BUILD HAY QUE HACER LOS N*N updates
3     int n;
4     vector<vi> st;
5     ll NEUT=0; // NEUTRO DE LA OPERACION
6     ll op(ll a, ll b){return a+b;} // OPERACION
7
8     st2d(int _n): n(_n) { st.resize(2*n+5, vi(2*n+5, 0)); }
9
10    void upd(int x, int y, ll v){
11        st[x+n][y+n]=v;
12        for(int j=y+n;j>1;j>=1)st[x+n][j>=1]=op(st[x+n][j],st[x+n][j^1]);
13        for(int i=x+n;i>1;i>=1)for(int j=y+n;j>=1) st[i>=1][j]=op(st[i][j],st[i^1][j]);
14    }
15
16    ll query(int x0, int x1, int y0, int y1){
17        ll r=NEUT; x1++, y1++;
18        for(int i0=x0+n,i1=x1+n;i0<i1;i0>=1,i1>=1){
19            int t[4],q=0;
20            if(i0&1)t[q++]=i0++;
21            if(i1&1)t[q++]--i1;
22            forn(k,q) for(int j0=y0+n,j1=y1+n;j0<j1;j0>=1,j1>=1){
23                if(j0&1)r=op(r,st[t[k]][j0++]);
24                if(j1&1)r=op(r,st[t[k]][--j1]);
25            }
26        }
27        return r;
28    }
29 };

```

2.10 Segtree Point Query

```

1 typedef ll tipo;
2 struct segtree {
3     vector<tipo> t; int tam;
4     tipo NEUT = 0;
5     tipo op(tipo a, tipo b){ return a + b; }
6
7     void build(vector<tipo> v, int n) { // build the tree

```

```

8      // root en 1, ojas en el intervalo [tam, 2*tam-1]
9      tam = sizeof(int) * 8 - __builtin_clz(n);
10     tam = 1 << tam; //Primera potencia de 2 mayor a n
11     t.resize(2 * tam, 0); for(i, n) t[tam+i] = v[i];
12     for(int i = tam - 1; i > 0; i--)
13         t[i] = op(t[i << 1], t[i << 1 | 1]); //Poner la operacion del
            seg
14 }
15
16 void modify(int l, int r, tipo value) {
17     for (l += tam, r += tam; l <= r; l >>= 1, r >>= 1) {
18         if (l&1) t[l] = op(t[l], value); //Poner la operacion del seg
19         if (!(r&1)) t[r] = op(t[r], value); //Poner la operacion del seg
20         l++, r--;
21     }
22 }
23
24 tipo query(int p) {
25     tipo res = NEUT;
26     for (p += tam; p > 0; p >>= 1) res = op(res, t[p]); //Operacion del
            seg
27     return res;
28 }
29 };

```

```

16     //Tener cuidado aca, podria ser que en cosas no conmutativas, para p
        impar: p > p^1.
17 }
18
19 tipo query(int l, int r) { // op on interval [l, r]
20     tipo res = NEUT;
21     for (l += tam, r += tam; l <= r; l >>= 1, r >>= 1) {
22         if (l&1) res = op(res, t[l++]);
23         if (!(r&1)) res = op(res, t[r--]);
24     }
25     return res;
26 }
27 };

```

2.11 Segtree Range Query

```

1 typedef long long tipo;
2 struct segtree {
3     vector<tipo> t; int tam;
4     tipo NEUT = 0; // Neutral element of operation
5     tipo op(tipo a, tipo b){ return a+b; } // Operation to make
6
7     void build(vector<tipo>v, int n) { // build the tree
8         // root en 1, ojas en el intervalo [tam, 2*tam-1]
9         tam = sizeof(int) * 8 - __builtin_clz(n); tam = 1<<tam;
10        t.resize(2*tam,NEUT); for(i, n) t[tam+i] = v[i];
11        for(int i = tam - 1; i > 0; i--) t[i] = op(t[i<<1], t[i<<1|1]);
12    }
13
14    void update(int p, tipo value) { // set value at position p
15        for (t[p += tam] = value; p > 1; p >>= 1) t[p>>1] = op(t[p], t[p^1])
            ;

```

2.12 Find Kth Min

```

1 struct Vertex {
2     Vertex *l, *r;
3     ll sum;
4
5     Vertex(ll val) : l(nullptr), r(nullptr), sum(val) {}
6     Vertex(Vertex *l, Vertex *r) : l(l), r(r), sum(0) {
7         if (l) sum += l->sum;
8         if (r) sum += r->sum;
9     }
10 };
11
12 Vertex* build(vector<int> &a, ll tl, ll tr) {
13     if (tl == tr) return new Vertex(a[tl]);
14     ll tm = (tl + tr) / 2;
15     return new Vertex(build(a, tl, tm), build(a, tm+1, tr));
16 }
17
18 ll get_sum(Vertex* v, ll tl, ll tr, ll l, ll r) {
19     if (l > r) return 0;
20     if (l == tl && tr == r) return v->sum;
21     ll tm = (tl + tr) / 2;
22     return get_sum(v->l, tl, tm, l, min(r, tm))
23         + get_sum(v->r, tm+1, tr, max(l, tm+1), r);
24 }
25
26 Vertex* update(Vertex* v, ll tl, ll tr, ll pos, ll new_val) {
27     if (tl == tr) return new Vertex(new_val);
28     ll tm = (tl + tr) / 2;
29     if (pos <= tm) return new Vertex(update(v->l, tl, tm, pos, new_val),
30         v->r);
31     else return new Vertex(v->l, update(v->r, tm+1, tr, pos, new_val));
32 }
33
34 void compresion(vector<int> ValoresQueAparecen, vector<int> &todos) {
35     for(int x : ValoresQueAparecen) todos.pb(x);
36     sort(todos.begin(), todos.end());
37     todos.erase(unique(todos.begin(), todos.end()), todos.end());
38 }
39 // Retorna el valor de f(x)
40

```

```

41 int obtener(vector<int> &todos, int x) {
42     return (int)(lower_bound(todos.begin(), todos.end(), x) - todos.begin());
43 }
44
45 struct kth_min {
46     int n;
47     vector<int> v, todos, v_comp, real_num;
48     vector<Vertex*> estado;
49     void read(int z) {
50         n = z;
51         v.resize(n); v_comp.resize(n); real_num.resize(n);
52         for(int i, n) cin >> v[i];
53     }
54     void prepare() {
55         compresion(v, todos);
56         for(int i, n) v_comp[i] = obtener(todos, v[i]);
57         for(int i, n) real_num[v_comp[i]] = v[i];
58         vector<int> cont(n, 0);
59         estado.pb(build(cont, 0, n-1));
60         for(int u : v_comp) {
61             cont[u]++;
62             estado.pb(update(estado.back(), 0, n-1, u, cont[u]));
63         }
64     }
65     int find_kth(int l, int r, int k) {
66         //tl and tr --> 0 indexed, k --> 1 indexed
67         Vertex* posl = estado[l];
68         Vertex* posr = estado[r+1];
69
70         int tl = 0, tr = n-1;
71         while(tl != tr) {
72             int tm = (tr + tl) / 2;
73             if (posr->l->sum - posl->l->sum >= k) {
74                 posr = posr->l;
75                 posl = posl->l;
76                 tr = tm;
77             }
78             else {
79                 k -= posr->l->sum - posl->l->sum;
80                 posr = posr->r;
81                 posl = posl->r;
82                 tl = tm+1;
83             }
84         }
85     }
86 }

```

```

84     }
85     return real_num[t1];
86 }
87 };

```

2.13 Lazy Arithmetic Sum

```

1  typedef long long tipo;
2
3  struct node {
4      tipo ans, l, r, a=0, d=0;
5      bool upd;
6      node() {ans = a = d = 0;}
7      node(tipo val, int pos) {ans = val; l = r = pos;}
8      void set_lazy(tipo _a, tipo _d, tipo start) {
9          upd=true, a += (l-start)*_d + _a; d += _d;
10     }
11 };
12
13 struct lazy {
14
15     #define l(x) int(x<<1)
16     #define r(x) int(x<<1|1)
17
18     vector <node> t; int tam;
19
20     node op(node a, node b) { //Operacion de query
21         node aux; aux.ans = a.ans + b.ans;
22         aux.l = a.l; aux.r = b.r;
23         return aux;
24     }
25
26     void push(tipo p) {
27         if(t[p].upd == true) {
28             tipo d = t[p].r - t[p].l + 1;
29             t[p].ans += d * t[p].a;
30             t[p].ans += t[p].d * (d-1)*(d)/2;
31             if(t[p].l < t[p].r) {
32                 t[l(p)].set_lazy(t[p].a,t[p].d,t[p].l);
33                 t[r(p)].set_lazy(t[p].a,t[p].d,t[p].l);
34             }
35             t[p].a = 0; t[p].d = 0; t[p].upd = false;
36         }

```

```

37     }
38
39     node query(tipo l, tipo r, int p = 1) {
40         push(p);
41         if(l > t[p].r || r < t[p].l) return node();
42         if(l <= t[p].l && t[p].r <= r) return t[p];
43         return op(query(l,r,l(p)),query(l,r,r(p)));
44     }
45
46     void update(tipo l, tipo r, tipo a, tipo d, int p = 1) {
47         push(p); node &cur = t[p];
48         if(l > cur.r || r < cur.l) return;
49         if(l <= cur.l && cur.r <= r) {
50             cur.set_lazy(a,d,l); push(p); return;
51         }
52         update(l, r, a, d, l(p)); update(l, r, a, d, r(p));
53         cur = op(t[l(p)], t[r(p)]);
54     }
55
56     void build(vector<tipo> v, int n) { // iterative build
57         tam = sizeof(int) * 8 - __builtin_clz(n); tam = 1<<tam;
58         t.resize(2*tam); v.resize(tam);
59         forn(i,tam) t[tam+i] = node(v[i],int(i));
60         for(int i = tam - 1; i > 0; i--) t[i] = op(t[l(i)],t[r(i)]);
61     }
62 };

```

2.14 Lazy Extreme

```

1  typedef long long tipo;
2  const int NEUT = 0; // REMINDER !!!
3
4  struct node {
5      tipo ans, l, r, lazy = 0;
6      bool upd = false;
7      node() {ans = lazy = 0; upd = false; l = r = -1;} // REMINDER !!! SET
8      NEUT
9      node(tipo val, int pos) : ans(val), l(pos), r(pos) {} // Set node
10     void set_lazy(tipo x) {lazy += x; upd = true;}
11 };
12
13 struct segtree_lazy {
14     #define l(x) int(x<<1)

```

```

14 #define r(x) int(x<<1|1)
15
16 vector <node> t; int tam;
17
18 node op(node a, node b) {
19     node aux; aux.ans = a.ans + b.ans; //Operacion de query
20     aux.l = a.l; aux.r = b.r;
21     return aux;
22 }
23
24 void push(int p) {
25     node &cur = t[p];
26     if(cur.upd == true) {
27         cur.ans += cur.lazy * (cur.r-cur.l+1); //Operacion update
28         if(cur.l < cur.r) {
29             t[l(p)].lazy += cur.lazy; t[l(p)].upd = true;
30             t[r(p)].lazy += cur.lazy; t[r(p)].upd = true;
31         }
32         cur.lazy = 0; cur.upd = false; //Poner el neutro del update
33     }
34 }
35
36 node query(int l, int r, int p = 1) {
37     push(p); node &cur = t[p];
38     if(l > cur.r || r < cur.l) return node(); // Return NEUT
39     if(l <= cur.l && cur.r <= r) return cur;
40     return op(query(l,r,l(p)),query(l,r,r(p)));
41 }
42
43 void update(int l, int r, tipo val, int p = 1) { // root at p = 1
44     push(p); node &cur = t[p];
45     if(l > cur.r || r < cur.l) return;
46     if(l <= cur.l && cur.r <= r) {
47         cur.set_lazy(val); push(p); return;
48     }
49     update(l, r, val, l(p)); update(l, r, val, r(p));
50     cur = op(t[l(p)], t[r(p)]);
51 }
52
53
54 void build(vector <tipo> v, int n) { // iterative build
55     tam = sizeof(int) * 8 - __builtin_clz(n); tam = 1<<tam;
56     t.resize(2*tam); v.resize(tam);

```

```

57     for(i,tam) t[tam+i] = node(v[i],i);
58     for(int i = tam - 1; i > 0; i--) t[i] = op(t[l(i)],t[r(i)]);
59 }
60 };

```

2.15 Merge Sort Tree

```

1 //~ Puede resolver Kth Min en O(Q * log^3 N)
2 //~ La funcion query(l,r,k) retorna la cantidad de numeros en el
3 //~ intervalo que son menores o iguales a k en [L,R] en O(log^2 N)
4
5 typedef long long tipo;
6
7 struct segtree {
8     struct node {
9         tipo ans = 0, l, r, inv = 0; // Poner el neutro del update
10        vector <tipo> cur, left, right;
11        tipo nomatch = 0; // No match en el intervalo de query
12        node base(node aux) { aux.ans = aux.inv = 0, aux.cur.clear(); return
13            aux;} //Poner el neutro de la query
14        void set_node(tipo x, tipo pos) {ans = 0, cur.pb(x); l = r = pos;}
15        void combine(node a, node b) {
16            left = a.cur, right = b.cur;
17            int p1 = 0, p2 = 0; inv = ans = 0;
18            while(p1 < a.cur.size() && p2 < b.cur.size()) {
19                if(a.cur[p1] <= b.cur[p2]) cur.pb(a.cur[p1++]);
20                else cur.pb(b.cur[p2++]), inv += a.cur.size() - p1;
21            }
22            while(p1 < a.cur.size()) cur.pb(a.cur[p1++]);
23            while(p2 < b.cur.size()) cur.pb(b.cur[p2++]);
24            inv += a.inv + b.inv;
25            l = min(a.l,b.l); r = max(a.r,b.r);
26        }
27    };
28    vector <node> t;
29
30 int BS(int k, vector <tipo> &aux) {
31     int a = -1, b = aux.size();
32     while(b-a > 1) {
33         int med = (a+b)/2;
34         if(aux[med] > k) b = med;
35         else a = med;
36     }

```

```

36     return b; //return aux.size()-b for K-Query
37 }
38
39 ll ask(int p, tipo l, tipo r, tipo k) {
40     if(l > t[p].r || r < t[p].l) return 0LL;
41     if(l <= t[p].l && t[p].r <= r) {
42         return BS(k,t[p].cur);
43     }
44     return ask(2*p+1,l,r,k)+ask(2*p+2,l,r,k);
45 }
46
47 void update(int p, tipo pos, tipo val) {
48     if(t[p].r < pos || t[p].l > pos) return;
49     if(t[p].l == t[p].r) { t[p].set_node(val,pos); return; }
50     update(2*p+1, pos, val); update(2*p+2, pos, val);
51     t[p].combine(t[2*p+1], t[2*p+2]);
52 }
53
54 void build(tipo a, tipo b, int p, vector<tipo> &v) {
55     if(a==b) {t[p].set_node(v[a],a); return;}
56     tipo med=(a+b)/2;
57     build(a, med, 2*p+1, v); build(med+1, b, 2*p+2, v);
58     t[p].combine(t[2*p+1], t[2*p+2]);
59 }
60 ll query(tipo l, tipo r, tipo k) {return ask(0,l,r,k);}
61 void modificar(tipo pos, tipo val) {update(0,pos,val);}
62 void construir(vector<tipo> &v, int n) { t.resize(4*n); build(0,n
    -1,0,v); }
63
64 tipo find_kth(int l, int r, tipo k) {
65     tipo minA = (tipo)(-1e9-7), maxA = (tipo)(1e9+7);
66     while(maxA - minA > 1) {
67         tipo med = (maxA + minA)/2;
68         int ans = query(l,r,med);
69         if(ans >= k) maxA = med;
70         else minA = med;
71     }
72     return maxA;
73 }
74 };

```

2.16 Persistent

```

1 struct Vertex {
2     Vertex *l, *r;
3     ll sum;
4
5     Vertex(ll val) : l(nullptr), r(nullptr), sum(val) {}
6     Vertex(Vertex *l, Vertex *r) : l(l), r(r), sum(0) {
7         if (l) sum += l->sum;
8         if (r) sum += r->sum;
9     }
10 };
11
12 Vertex* build(ll a[], ll tl, ll tr) {
13     if (tl == tr)
14         return new Vertex(a[tl]);
15     ll tm = (tl + tr) / 2;
16     return new Vertex(build(a, tl, tm), build(a, tm+1, tr));
17 }
18
19 ll get_sum(Vertex* v, ll tl, ll tr, ll l, ll r) {
20     if (l > r)
21         return 0;
22     if (l == tl && tr == r)
23         return v->sum;
24     ll tm = (tl + tr) / 2;
25     return get_sum(v->l, tl, tm, l, min(r, tm))
26         + get_sum(v->r, tm+1, tr, max(l, tm+1), r);
27 }
28
29 Vertex* update(Vertex* v, ll tl, ll tr, ll pos, ll new_val) {
30     if (tl == tr)
31         return new Vertex(new_val);
32     ll tm = (tl + tr) / 2;
33     if (pos <= tm)
34         return new Vertex(update(v->l, tl, tm, pos, new_val), v->r);
35     else
36         return new Vertex(v->l, update(v->r, tm+1, tr, pos, new_val));
37 }
38
39 vector<Vertex*> estado;

```

2.17 Segtree

```

1 typedef long long tipo;

```

```

2  const tipo NEUT = 0; // REMINDER !!!
3
4  struct node {
5      int l, r; tipo ans;
6      node() {ans = NEUT, l = r = -1;} // REMINDER !!! Set NEUT
7      node(tipo val, int pos): l(pos), r(pos),ans(val) {} // Set node
8      void update(tipo val) {ans = val;} // Define update function
9  };
10
11 struct segtree { // Segtree for Sum Range Query
12     #define l(x) int(x<<1)
13     #define r(x) int(x<<1|1)
14
15     vector<node> t; int tam;
16
17     node op(node a, node b) { //Operacion de query
18         node aux; aux.ans = a.ans + b.ans;
19         aux.l = a.l; aux.r = b.r;
20         return aux;
21     }
22
23     node query(int l, int r, int p = 1) {
24         node &cur = t[p];
25         if(l > cur.r || r < cur.l) return node(); // Return NEUT
26         if(l <= cur.l && cur.r <= r) return cur;
27         return op(query(l,r,l(p)),query(l,r,r(p)));
28     }
29
30     void update(int pos, tipo val, int p = 1) { // root at p = 1
31         node &cur = t[p];
32         if(cur.r < pos || cur.l > pos) return;
33         if(cur.l == cur.r) { cur.update(val); return; }
34         update(pos, val, l(p)); update(pos, val, r(p));
35         cur = op(t[l(p)], t[r(p)]);
36     }
37
38     void build(vector<tipo> v, int n) { // iterative build
39         tam = sizeof(int) * 8 - __builtin_clz(n); tam = 1<<tam;
40         t.resize(2*tam); v.resize(tam);
41         forn(i,tam) t[tam+i] = node(v[i],int(i));
42         for(int i = tam - 1; i > 0; i--) t[i] = op(t[l(i)],t[r(i)]);
43     }
44 };

```

```

45
46 //~ Max Sum Query
47 struct node {
48     tipo ans, pref, suff, sum, l, r;
49     node() {ans = pref = suff = sum = NEUT, l = r = -1;}
50     node(tipo val, int pos) : ans(val), pref(val), suff(val), sum(val), l
51         (pos), r(pos) {}
52     void update(tipo val) {ans = pref = suff = sum = val;}
53 };
54
55 node op(node a, node b) {
56     node aux;
57     aux.pref = max(a.pref, a.sum + b.pref);
58     aux.suff = max(b.suff, b.sum + a.suff);
59     aux.sum = a.sum + b.sum;
60     aux.ans = max(max(a.ans,b.ans),a.suff + b.pref);
61     aux.l = a.l; aux.r = b.r;
62     return aux;
63 }

```


2.18 Sqrt Decomposition

```

1 // This code solves the same problem as lazy arithmetic sum
2
3 struct block {
4     ll l, r, a, b, ans;
5     block(ll _l, ll _r) {
6         l = _l; r = _r; ans = a = b = 0;
7     }
8     block() {}
9 };
10
11 const int BSIZE = 350; // Size of each block
12
13 struct segment_block {
14     int n;
15     vi v;
16     vector<block> t;
17
18     void refresh_block(block &u) { //clean block and refresh vector
19         block new_b = block(u.l,u.r);
20         forr(i,u.l,u.r+1) {
21             v[i] += u.a + u.b * (i - u.l);
22             new_b.ans += v[i];
23         }
24         u = new_b;
25     }
26
27     ll query_block(block &u, ll l, ll r) { // The answer of one blk
28         if(l <= u.l && u.r <= r) { // full intersection
29             ll last = u.a + u.b * (u.r - u.l);
30             return u.ans + (u.a + last) * (u.r-u.l+1) / 2;
31         }
32         // partial intersection
33         refresh_block(u);
34         ll ans = 0;
35         forr(i,max(l,u.l),min(u.r+1,r+1)) {
36             ans += v[i];
37         }
38         return ans;
39     }
40
41     void update_block(block &u, ll l, ll r, ll a, ll b) { // upd one blk

```

```

42         if(l <= u.l && u.r <= r) { // full intersection
43             u.a += (u.l - l) * b + a;
44             u.b += b;
45             return;
46         }
47         // partial intersection
48         refresh_block(u);
49         forr(i,max(l,u.l),min(u.r+1,r+1)) {
50             v[i] += (i - l) * b + a;
51             u.ans += (i - l) * b + a;
52         }
53     }
54
55     ll query(ll l, ll r) {
56         ll ans = 0;
57         for(block &u : t) {
58             if(r < u.l || l > u.r) continue;
59             ans += query_block(u,l,r);
60         }
61         return ans;
62     }
63
64     void update(ll l, ll r, ll a, ll b) {
65         for(block &u : t) {
66             if(r < u.l || l > u.r) continue;
67             update_block(u,l,r,a,b);
68         }
69     }
70
71     void build(vi _v, int _n) {
72         v = _v; n = _n;
73         int tam = 0;
74         while(tam <= n-1) {
75             t.pb(block(tam,min(n-1,tam+BSIZE-1)));
76             tam += BSIZE;
77             refresh_block(t.back());
78         }
79     }
80
81     void debug() { // good for check if updates are correct
82         for(block u : t) {
83             refresh_block(u);
84             forr(i,u.l,u.r+1) cerr << v[i] << " ";

```

```

85     }
86     cerr << "\n"; RAYA;
87 }
88 };

```

2.19 Sum Min K Multiset

```

1  typedef int tipo;
2
3  struct MultisetWithSum {
4      multiset<tipo> ms;
5      tipo sum = 0;
6      void insert(tipo x) {ms.insert(x); sum += x;}
7      void erase(tipo x) {ms.erase(ms.find(x)); sum -= x;}
8      int size() const { return int(ms.size()); }
9      tipo highest() const { return *ms.rbegin(); }
10     tipo lowest() const { return *ms.begin(); }
11 };
12
13 struct SumMinKMultiset {
14     MultisetWithSum low, high;
15     void insert(tipo x) {low.insert(x);}
16     void erase(tipo x) { (low.ms.find(x) != low.ms.end() ? low : high).
17         erase(x); }
18     tipo sumMinK(int k) { adjust(k); return low.sum; }
19     tipo kth_element(int k) { adjust(k); return low.highest(); }
20     void adjust(int k) { // Low tenga k elementos
21         assert(low.size() + high.size() >= k);
22         while (low.size() < k) {
23             tipo toMove = high.lowest();
24             low.insert(toMove);
25             high.erase(toMove);
26         }
27         while (high.size() > 0 && high.lowest() < low.highest()) {
28             tipo toMove = high.lowest();
29             high.erase(toMove); low.insert(toMove);
30             toMove = low.highest();
31             low.erase(toMove); high.insert(toMove);
32         }
33         while (low.size() > k) {
34             tipo toMove = low.highest();
35             low.erase(toMove);
36             high.insert(toMove);

```

```

36     }
37 }
38 };

```

2.20 Treap

```

1  // Treap, es levemente mas poderoso que un set, pero mas lento
2  // Complejidad: O(log n) por operacion
3  typedef struct item *pitem;
4  struct item {
5      int pr,key,cnt;
6      pitem l,r;
7      item(int key):key(key),pr(rand()),cnt(1),l(0),r(0) {}
8  };
9  int cnt(pitem t){return t?t->cnt:0;}
10 void upd_cnt(pitem t){if(t)t->cnt=cnt(t->l)+cnt(t->r)+1;}
11 void split(pitem t, int key, pitem& l, pitem& r){ // l: < key, r: >= key
12     if(!t)l=r=0;
13     else if(key<=t->key)split(t->l,key,l,t->l),r=t;
14     else split(t->r,key,t->r,r),l=t;
15     upd_cnt(t);
16 }
17 void insert(pitem& t, pitem it){
18     if(!t)t=it;
19     else if(it->pr>t->pr)split(t,it->key,it->l,it->r),t=it;
20     else insert(it->key<t->key?t->l:t->r,it);
21     upd_cnt(t);
22 }
23 void merge(pitem& t, pitem l, pitem r){
24     if(!l||!r)t=l?l:r;
25     else if(l->pr>r->pr)merge(l->r,l->r,r),t=l;
26     else merge(r->l,l,r->l),t=r;
27     upd_cnt(t);
28 }
29 void erase(pitem& t, int key){
30     if(t->key==key)merge(t,t->l,t->r);
31     else erase(key<t->key?t->l:t->r,key);
32     upd_cnt(t);
33 }
34 void unite(pitem &t, pitem l, pitem r){
35     if(!l||!r){t=l?l:r;return;}
36     if(l->pr<r->pr)swap(l,r);
37     pitem p1,p2;split(r,l->key,p1,p2);

```

```

38     unite(l->l,l->l,p1);unite(l->r,l->r,p2);
39     t=l;upd_cnt(t);
40 }
41 pitem kth(pitem t, int k){
42     if(!t)return 0;
43     if(k==cnt(t->l))return t;
44     return k<cnt(t->l)?kth(t->l,k):kth(t->r,k-cnt(t->l)-1);
45 }
46 pair<int,int> lb(pitem t, int key){ // position and value of lower_bound
47     if(!t)return {0,1<<30}; // (special value)
48     if(key>t->key){
49         auto w=lb(t->r,key);w.first+=cnt(t->l)+1;return w;
50     }
51     auto w=lb(t->l,key);
52     if(w.first==cnt(t->l))w.second=t->key;
53     return w;
54 }

```

3 Dp

3.1 Cht Dynamic

```

1 // Convex Hull trick dinamico.
2 // O(log n) para agregar e consultar.
3 // Para minimo, cambiar el signo de m y b.
4 typedef ll tc;
5 const tc is_query=-(1LL<<62); // special value for query
6 struct Line {
7     tc m,b;
8     mutable multiset<Line>::iterator it,end;
9     const Line* succ(multiset<Line>::iterator it) const {
10         return (++it==end? NULL : &*it);}
11     bool operator<(const Line& rhs) const {
12         if(rhs.b!=is_query)return m<rhs.m;
13         const Line *s=succ(it);
14         if(!s)return 0;
15         return b-s->b<(s->m-m)*rhs.m;
16     }
17 };
18 struct HullDynamic : public multiset<Line> { // for maximum
19     bool bad(iterator y){
20         iterator z=next(y);
21         if(y==begin()){
22             if(z==end())return false;
23             return y->m==z->m&&y->b<=z->b;
24         }
25         iterator x=prev(y);
26         if(z==end())return y->m==x->m&&y->b<=x->b;
27         return (x->b-y->b)*(z->m-y->m)>=(y->b-z->b)*(y->m-x->m);
28     }
29     iterator next(iterator y){return ++y;}
30     iterator prev(iterator y){return --y;}
31     void add(tc m, tc b){
32         // m *= -1; b *= -1; --> For mininum
33         iterator y=insert((Line){m,b});
34         y->it=y;y->end=end();
35         if(bad(y)){erase(y);return;}
36         while(next(y)!=end()&&bad(next(y)))erase(next(y));
37         while(y!=begin()&&bad(prev(y)))erase(prev(y));
38     }
39     tc eval(tc x){

```

```

40     Line l=*lower_bound((Line){x,is_query});
41     return l.m*x+l.b; // -1*(l.m*x + l.b) for minimum
42 }
43 };

```

3.2 Cht Li Chao Tree

```

1  //~ Li Chao Tree
2  //~ Permite resolver CHT dinamico
3  //~ Desventaja: si evaluamos  $x > 10^6$ , no aguanta el segment tree
4
5  const int N = int(1e6 + 1);
6  const ll INF = ll(1e18+10);
7
8  struct punto { ll x, y; };
9
10 vector <punto> tree(4*N, {0, INF});
11
12 ll f(punto line, ll x) { return line.x * x + line.y; }
13
14
15 void insert(punto line, ll lo = 1, ll hi = N, int i = 1){
16     ll m = (lo + hi) / 2;
17     bool left = f(line, lo) < f(tree[i], lo);
18     bool mid = f(line, m) < f(tree[i], m);
19
20     if(mid) swap(tree[i], line);
21
22     if(hi - lo == 1) return;
23     else if(left != mid) insert(line, lo, m, 2*i);
24     else insert(line, m, hi, 2*i+1);
25 }
26
27 ll query(ll x, ll lo = 1, ll hi = N, int i = 1){
28     int m = (lo+hi)/2;
29     ll curr = f(tree[i], x);
30     if(hi-lo==1) return curr;
31     if(x<m) return min(curr, query(x, lo, m, 2*i));
32     else return min(curr, query(x, m, hi, 2*i+1));
33 }

```

3.3 Chull Trick

```

1  // Convex Hull Trick

```

```

2  // 0(log n) para agregar linea y consultar
3  // Las linease agregan en orden creciente de pendiente
4  typedef long long tipo;
5  typedef __int128 ull;
6
7  struct punto {
8      tipo x, y;
9      punto operator -(const punto &p) const {return {x-p.x,y-p.y};}
10     punto operator +(const punto &p) const {return {x+p.x,y+p.y};}
11     tipo operator *(const punto &p) const {return x*p.x + y*p.y;}
12     ull operator ^(const punto &p) const { // Producto Cruz
13         return (ull)x * p.y - (ull)y * p.x; } // (ull) --> __int128
14 };
15
16 struct chull { // Agregar siempre pendientes en orden creciente
17     ll op = 1; // 1 para minimo, -1 para maximo
18     vector <punto> hull;
19
20     chull(bool maxi) { // true para maximo, false para minimo
21         if(maxi) op *= -1;
22         // add_line(0,0); // Push base case {0,0} if necessary
23     }
24     ll get(tipo x) {
25         punto query = {op * x, op * 1LL};
26         int a=0, b=SZ(hull);
27         while(b-a > 1) {
28             int med = (a+b)/2;
29             if(query * hull[med-1] >= query * hull[med]) a = med;
30             else b = med;
31         }
32         return op * (query * hull[a]);
33     }
34
35     bool check(punto aux, int last) {
36         return op * ((hull[last]-hull[last-1])^(aux-hull[last])) <= 0;
37     }
38
39     void add_line(ll x, ll y) {
40         if(SZ(hull)) assert(x >= hull.back().x); // Chequeo de
41                                     pendientes crecientes
42
43         punto aux = {x,y}; int last = SZ(hull)-1;
44         while(last > 0 && check(aux,last)) {

```

```

44         hull.pop_back(); last--;
45     }
46     hull.pb(aux);
47 }
48 };

```

3.4 Count Cycles

```

1 // Count the number of cycles in a graph
2 // Time: O(n*2^n)
3 // Usage: count_cycles(n); // n --> number of nodes
4 const int MAXN = 20;
5 int g[MAXN][MAXN];
6
7 bool bit(int mask, int i) {
8     if((mask & 1<<i) != 0) return true;
9     return false;
10 }
11
12 int first(int mask, int n) {
13     for(int i = 0; i < n; i++) {
14         if(bit(mask,i)) {return i;}
15     }
16     return -1;
17 }
18
19 int count(int mask, int n) {
20     int ans = 0;
21     for(int i = 0; i < n; i++) {
22         if(bit(mask,i)) ans++;
23     }
24     return ans;
25 }
26
27 ll count_cycles(int n) {
28     ll dp[1<<MAXN][MAXN];
29     memset(dp,0,sizeof(dp));
30     ll ans = 0;
31     for(int mask = 1; mask < (1<<n); mask++) {
32         if(count(mask,n) == 1) {
33             dp[mask][first(mask,n)] = 1;
34             continue;
35         }

```

```

36     int f = first(mask,n);
37     for(int last = 0; last < n; last++) {
38         if(!bit(mask,last) || last == f) continue;
39         for(int next = 0; next < n; next++) {
40             if(g[last][next]) {
41                 dp[mask][last] += dp[mask ^ (1<<last)][next];
42             }
43         }
44         if(count(mask,n) >= 3 && g[first(mask,n)][last]) ans += dp[
45             mask][last];
46     }
47     return ans / 2;
48 }

```

3.5 Divide And Conquer Dp Opt

```

1 // Divide and conquer optimization for DP
2 // Complexity: O(nklogn)
3 // Usage: solve_dac(n,k); // n --> number of elements, k --> number of
4 // groups
5 ll costo[MAXN][MAXN];
6
7 vector <ll> last(MAXN), dp(MAXN);
8
9 void calc_costo(int n, vector <ll> &v) {
10     // Dependiendo el problema, aca dentro realizamos la funcion para
11     // precalcular el costo de cada intervalo y guardarlo en 'costo'
12 }
13
14 void compute(int l, int r, int optl, int optr) {
15     if(l > r) return;
16     int med = (l+r)/2;
17     pair<ll,ll> best = {INF,-1};
18     for(int p = optl; p <= min(med,optr); p++) {
19         best = min(best,{last[p] + costo[p+1][med],p});
20     }
21     dp[med] = best.first;
22     int opt = best.second;
23     compute(l,med-1,optl,opt);
24     compute(med+1,r,opt,optr);
25 }

```

```

26
27 ll solve_dac(int n, int k) { // divide and conquer optimization
28     for(int i = 0; i < n; i++) last[i] = costo[0][i];
29     for(int i = 2; i <= k; i++) {
30         fill(all(dp), INF);
31         compute(0, n-1, 0, n-1);
32         last = dp;
33     }
34     return dp[n-1];
35 }

```

3.6 Elevator Problem

```

1  /*
2  * Elevator Problem
3  *
4  * Dado los pesos maximos de las personas y el peso maximo del
5  * ascensor, cual es el minimo numeros de viajes necesarios para subir
6  * a todas las personas:  $O(N * 2^N)$ 
7  *
8  */
9
10 ll elevator_problem(int n, ll x, vector<ll> &peso) {
11     vector<pair<ll, ll>> best(1<<n); best[0] = {1, 0};
12     forr(s, 1, (1<<n)) {
13         best[s] = {n+1, 0};
14         forn(p, n) {
15             if(s & (1<<p)) {
16                 pair<ll, ll> option = best[s^(1<<p)];
17                 if(option.second + peso[p] <= x) {
18                     option.second += peso[p];
19                 }
20                 else{
21                     option.first++; option.second = peso[p];
22                 }
23                 best[s] = min(best[s], option);
24             }
25         }
26     }
27     return best[(1<<n)-1].first;
28 }

```

3.7 Knapsack

```

1  //~ Problema de la mochila -  $O(n*w)$ 
2
3  struct item { int w, v; };
4
5  ll knapsack(vector<item> &v, int n, int w) {
6      vector<ll> dp(w+1, 0), aux(w+1, 0);
7      forn(i, n) {
8          forn(j, w+1) {
9              aux[j] = dp[j];
10             if(j-v[i].w < 0) continue;
11             aux[j] = max(aux[j], dp[j-v[i].w] + v[i].v);
12         }
13         forn(j, w+1) dp[j] = aux[j];
14     }
15     ll ans = 0;
16     forn(i, w+1) ans = max(ans, dp[i]);
17     return ans;
18 }

```

3.8 Knapsack Big W

```

1  // Description: Given a set of items, each with a weight and a value,
2  // determine the number of each item to include in a collection so that
3  // the total weight is less than or equal to a given limit and the
4  // total value is as large as possible.
5  // Time:  $O(n * MAXV)$ 
6  // Usage: knapsack(v, n, w); // v --> vector of items, n --> number of
7  // items, w --> maximum weight
8
9  struct item { int w, v; };
10
11 ll knapsack(vector<item> &v, int n, int w) {
12     int MAXV = 100005;
13     vector<ll> dp(MAXV+1, INF), aux(MAXV+1, INF);
14     vector<bool> visto(MAXV, false);
15     visto[0] = true; dp[0] = 0;
16     forn(i, n) {
17         forn(j, MAXV) {
18             aux[j] = dp[j];
19             if(j-v[i].v < 0) continue;
20             if(visto[j-v[i].v] == true) {
21                 aux[j] = min(aux[j], dp[j-v[i].v] + v[i].w);
22                 visto[j] = true;
23             }
24         }
25     }
26 }

```

```

19     }
20     }
21     forn(j,MAXV) dp[j]=aux[j];
22 }
23 ll ans = 0;
24 forn(i,MAXV+1) if(dp[i] <= w && visto[i]==true) ans=i;
25 return ans;
26 }

```

3.9 Knuth Division Dp Opt

```

1 const ll INF = (ll)(1e18+10);
2 const int MAXN = 5005;
3 vector <ll> prefix;
4
5 // Sea Pos(i,j) la posicion del corte que optimiza la dp, intervalo [i,j]
6 // Se puede aplicar optimizacion de Knuth si se cumple que:
7 // Pos(i,j-1) <= pos(i,j) <= pos(i+1,j)
8 // Objetivo: llevar un array a elementos individuales mediante cortes.
9
10 ll suma(int i, int j) {
11     return prefix[j+1] - prefix[i]; // Suma de los elementos del inter
12 }
13
14 ll solve_Knuth(int n) {
15     vector <ll> pos(n+1);
16     vector <vector <ll> > dp(MAXN,vector<ll>(MAXN,INF));
17     for(int i = 0; i < n; i++) dp[i][i] = 0, pos[i] = i;
18
19     for(int len = 1; len < n; len++) {
20         vector <ll> aux(n+1,0);
21         for(int i = 0; i < n - len; i++) {
22             for(int k = pos[i]; k <= pos[i+1]; k++) {
23                 ll cost = suma(i,i+len) + dp[i][k] + dp[k+1][i+len];
24                 if(cost < dp[i][i+len]) {
25                     dp[i][i+len] = cost; aux[i] = k;
26                 }
27             }
28         }
29         pos = aux;
30     }
31     return dp[0][n-1];

```

```

32 }

```

3.10 Lcs

```

1 // Longest Common Subsequence
2 // O(n*m) time, O(n*m) space
3 // Usage: LCS(s1,s2)
4
5 string LCS(string s1, string s2) { // O(n*m)
6     const int n = SZ(s1), SZ(s2);
7     int dp[n+1][m+1]; pair<int,int> last[n+1][m+1];
8     forn(i,n+1) {
9         forn(j,m+1) {
10             if(i==0 || j==0) dp[i][j] = 0;
11             else if(s1[i-1] == s2[j-1])
12                 dp[i][j] = dp[i-1][j-1]+1, last[i][j] = {i-1,j-1};
13             else if(dp[i-1][j] > dp[i][j-1])
14                 dp[i][j] = dp[i-1][j], last[i][j] = {i-1,j};
15             else dp[i][j] = dp[i][j-1], last[i][j] = {i,j-1};
16         }
17     }
18     string ans; pair<int,int> cur = {n,m};
19     while(cur.first != 0 && cur.second != 0) {
20         int x = cur.first, y = cur.second;
21         if(x-last[x][y].first==1 && y-last[x][y].second==1) ans += s1[x-1];
22         cur = last[x][y];
23     }
24     reverse(all(ans));
25     return ans;
26 }

```

3.11 Lis

```

1 // Longest Increasing Subsequence
2 // O(n*log(n)) time, O(n) space
3 // Usage: lis(a,strict)
4
5 // strict = 1: estrictamente creciente
6 ll lis(vi &a, int strict = 0){
7     vi temp; temp.pb(a[0]);
8     forr(i, 1, SZ(a)){
9         ll x = a[i];
10         if(x >= temp.back()+strict) temp.pb(x);

```

```

11     else {
12         auto it = upper_bound(all(temp), x-strict);
13         *it = x;
14     }
15 }
16 return SZ(temp);
17 }

```

3.12 Merge Sort

```

1 // Merge sort con inversiones
2 // Complejidad: O(n log n)
3 // Usage: merge_sort(0,n-1,v)
4
5 ll inv = 0; // Inversiones
6
7 vector<int> merge_sort(int li, int ri, vector<int> &v) {
8     if(li == ri) return {v[li]};
9     vector<int> ans;
10    int med = (li+ri)/2;
11    vector<int> l = merge_sort(li,med,v);
12    vector<int> r = merge_sort(med+1,ri,v);
13    int a = 0, b = 0;
14    while(a < l.size() && b < r.size()) {
15        if(l[a] <= r[b]) ans.pb(l[a++]);
16        else { inv += (ll)(l.size() - a); ans.pb(r[b++]); }
17    }
18    while(a < l.size()) ans.pb(l[a]), a++;
19    while(b < r.size()) ans.pb(r[b]), b++;
20    return ans;
21 }

```

4 Geometry

4.1 Angular Sort

```

1 // Angular sort
2 struct frac {
3     ll num, den;
4     frac() {}
5     frac(ll x, ll y) {
6         ll m = __gcd(abs(x),abs(y));
7         num = x / m; den = y / m;
8     }
9     int cuad(ll n, ll d) {
10        if(n >= 0 && d >= 0) return 1;
11        if(n >= 0 && d < 0) return 2;
12        if(n < 0 && d < 0) return 3;
13        if(n < 0 && d >= 0) return 4;
14    }
15    bool operator <(frac &p) {
16        if(cuad(num,den) != cuad(p.num,p.den)) {
17            return cuad(num,den) < cuad(p.num,p.den); }
18        return num * p.den < p.num * den;
19    }
20    bool operator ==(frac &p) {
21        return (num * p.den == den * p.num) && (cuad(num,den) == cuad(p.
22            num,p.den));
23    }
24 };
25 struct punto {
26     ll x, y;
27     frac pend;
28     ll val;
29     punto() {}
30     punto(ld a, ld b, ll z) {
31         x = a, y = b, val = z, pend = frac(b,a);}
32     bool operator <(punto p) {
33         return pend < p.pend;
34     }
35 };

```

4.2 Calipers


```

1 // Devuelve la distancia entre los puntos mas lejanos en O(N)
2 // Cuando se usa con nuestra convex_hull, hacer reverse porque vienen CW
3 ld callipers(vector<punto> p, int n){
4     ld r=0; // prereq: Convex, ccw, NO COLLINEAR POINTS
5     for(int i=0,j=n-2?0:1;i<j;++i){
6         for(;;j=(j+1)%n){
7             r=max(r,(p[i]-p[j]).mod());
8             if((p[(i+1)%n]-p[i])^(p[(j+1)%n]-p[j])<=EPS)break;
9         }
10    }
11    return r;
12 }

```

4.3 Closest Points

```

1 typedef long double tipo;
2
3 struct punto {
4     tipo x, y; int ind;
5     punto operator -(punto p) const {return {x-p.x,y-p.y};}
6     bool operator <(punto p) const {return x != p.x ? x < p.x : y < p.y;
7     };}
8     tipo mod() {return sqrtl(x*x + y*y);}
9     tipo mod2() {return x*x + y*y;}
10 };
11 tuple<tipo,int,int> closest(vector<punto> &p) { //closest and indices
12     int n = p.size();
13     set<punto> s;
14     tipo best = (tipo)(1e18); int ansi, ansj;
15     int j = 0;
16     forn(i,n) {
17         tipo d = ceil(sqrt(best));
18         while(p[i].x - p[j].x >= best) s.erase({p[j].y, p[j].x, j}), j
19             ++;
20         auto it1 = s.lower_bound({p[i].y - d, p[i].x});
21         auto it2 = s.upper_bound({p[i].y + d, p[i].x});
22         for(auto it = it1; it != it2; ++it) {
23             tipo dx = p[i].x - it->x;
24             tipo dy = p[i].y - it->y;
25             if(dx * dx + dy * dy < best) {
26                 best = dx * dx + dy * dy;
27                 ansi = i, ansj = it->ind;
28             }
29         }
30     }
31     return {best, ansi, ansj};
32 }

```

```

27     }
28     }
29     s.insert({p[i].y, p[i].x, (int)i});
30 }
31 return {sqrtl(best),ansi,ansj};
32 }

```

4.4 Closest Points Monogon

```

1 const ll INF = (ll)(1e15+10); // (1e18+10)
2
3 template<typename T>
4 using minpq = priority_queue<T, vector<T>, greater<T>>;
5
6 struct punto {
7     ll x, y;
8     punto (ll x = 0, ll y = 0) : x(x), y(y) {}
9     punto operator -(const punto &p) const {
10         return punto(x - p.x, y - p.y); }
11     ll len2() const { return x * x + y * y; }
12     bool const operator <(const punto &p) const {
13         return (x != p.x) ? x < p.x : y < p.y;
14     }
15 };
16
17 istream& operator>>(istream &is, punto &p) {
18     return is >> p.x >> p.y;
19 }
20
21 ll closest(vector<punto> &p, int l, int r) { //Init l = 0 and r = p.size
22     ()-1
23     if(l == r) return INF;
24     int m = (l + r) / 2;
25     ll mid = p[m].x;
26     ll d = min(closest(p, l, m), closest(p, m + 1, r));
27     vector<punto> A, B, C;
28     forr(i, l, m + 1) {
29         ll r = mid - p[i].x;
30         if(r * r <= d) A.push_back(p[i]);
31     }
32     forr(i, m + 1, r + 1) {
33         ll r = mid - p[i].x;
34         if(r * r <= d) B.push_back(p[i]);
35     }
36 }

```

```

34     }
35     auto cmpy = [&](punto a, punto b) { return a.y < b.y; };
36     merge(all(A), all(B), back_inserter(C), cmpy);
37     forr(i, 0, C.size()) {
38         int j = i + 1;
39         while(j < C.size() && (C[i] - C[j]).y * (C[i] - C[j]).y <= d) {
40             d = min(d, (C[i] - C[j]).len2());
41             j++;
42         }
43     }
44     inplace_merge(p.begin() + 1, p.begin() + m + 1, p.begin() + r + 1,
45                 cmpy);
46     return d;
47 }
48 void sorting(vector <punto> &p) { sort(all(p)); } // Don't Forget

```

4.5 Complex

```

1  /*
2  * Numeros complejos
3  *
4  * Tener en cuentas que:
5  *
6  * - No se pueden leer complejos directamente con cin >>;
7  * - Una vez definidas las macros, no podemos definir variables con el
8  *   mismo nombre
9  *
10 */
11
12 typedef complex<double> complejo;
13 #define x real()
14 #define y imag()
15
16 complejo a, b, c;
17 abs(a); // Devuelve sqrt(a.x^2 + a.y^2)
18 arg(a); // Devuelve el angulo entre ( -pi ; pi ]
19 conj(a); // Devuelve a.x - i*a.y
20 a*b; // Devuelve (a.x*b.x-a.y*b.y) + i*(a.x*b.y + a.y*b.x)
21 polar(1.0,pi); // Devuelve el complejo de modulo 1 y angulo pi
22 (conj(a)*b).y // Producto cruz -> a.x*b.y - a.y*b.x
23 (conj(b-a)*(c-a)).y/abs(b-c) // Devuelve distancia punto a -> recta bc

```

4.6 Convex Hull

```

1  /*
2  Convex Hull:
3  Complejidad O(N).
4  La hull no tiene puntos colineales
5  Los devuelve en orden clockwise -> Para rotating calipers hacer reverse
6  */
7
8  const ld EPS = 1e-10;
9  void convex_hull(vector<punto> &a) {
10     if(SZ(a) == 1) return;
11     sort(all(a));
12     punto p1 = a[0];
13     vector<punto> up, down;
14     up.pb(p1); down.pb(p1);
15     forr(i,1,SZ(a)) {
16         int n = SZ(up), m = SZ(down);
17         while(n > 1 && ((up[n-1]-up[n-2])^(a[i]-up[n-1])) >= -EPS) {
18             up.pop_back(); n--;
19         } up.pb(a[i]);
20         while(m > 1 && ((down[m-1]-down[m-2])^(a[i]-down[m-1])) <= EPS) {
21             down.pop_back(); m--;
22         } down.pb(a[i]);
23     } // Cambiar EPS a 0 para mejor precision en enteros.
24     a.clear();
25     for(punto u : up) a.pb(u);
26     for(int i = SZ(down)-2; i > 0; i--) a.pb(down[i]);
27 }

```

4.7 Formulas

```

1  // Shoelace formula
2  tipo area(vector <punto> &v) {
3      tipo ans = 0.0; int n = v.size();
4      forn(i,n) ans += v[i] ^ v[(i+1)%n];
5      return fabs(ans/2.0);
6  }
7
8  tipo dist_point_line(punto &p, recta &r) {
9      punto p1 = r.p, p2 = r.p + r.v;
10     return fabs((p1-p)^(p2-p))/r.v.mod();

```

```

11 }
12
13 punto project(punto a, punto b) { //Proyeccion de b sobre a
14     return ((a*b)/a.mod2()) * a;
15 }
16
17 tipo find_alpha(recta r, punto p) {
18     return r.v.x != 0 ? (p.x-r.p.x)/r.v.x : (p.y-r.p.y)/r.v.y;
19 }

```

4.8 Intersection All

```

1 struct recta { // Puede usarse para segmentos ([p,p+v] o alpha = [0,1])
2     punto v, p; // v -> director, p -> punto por donde pasa
3     recta(punto p1, punto p2) { v = (p2-p1); p = p1;}
4     recta() {}
5     recta(tipo A, tipo B, tipo C) { // Transform Ax + By + C = 0
6         v = {-B,A}; A != 0 ? p = {-C / A,0} : p = {0,-C / B};
7     }
8     bool is_in(punto q){return fabs((q.x-p.x)*v.y - (q.y-p.y)*v.x) < EPS
9         ;}
10    punto eval(double x) {return x * v + p;}
11 };
12
13 bool inter_recta(recta &r1, recta &r2, punto &ans) {
14     // Retorna false si son paralelas, sino guarda el punto en ans
15     if(fabs(r1.v ^ r2.v) < EPS) return false;
16     tipo alpha = tipo((r2.p - r1.p)^r2.v) / tipo(r1.v^r2.v);
17     ans = r1.p + alpha * r1.v;
18     return true;
19 }
20
21 bool inter_seg(recta &r1, recta &r2, punto &ans) {
22     if(r1.p == r2.p || r1.p == r2.p+r2.v) {ans = r1.p; return true;}
23     if(r1.p+r1.v == r2.p || r1.p+r1.v == r2.p+r2.v) {
24         ans = r1.p+r1.v; return true; } //Casos que coincidan extremos
25     if(fabs(r1.v ^ r2.v) < EPS) return false; // son paralelos
26     tipo alpha = tipo((r2.p - r1.p)^r2.v) / tipo(r1.v^r2.v);
27     tipo beta = tipo((r1.p - r2.p)^r1.v) / tipo(r2.v^r1.v);
28     if(alpha < -EPS || beta < -EPS) return false;
29     if(alpha > 1.0+EPS || beta > 1.0+EPS) return false;
30     ans = r1.p + alpha * r1.v; return true;
31 }

```

```

31
32 struct circ { punto c; tipo r; };
33
34 tipo dist_point_line(punto &p, recta &r) {
35     punto p1 = r.p, p2 = r.p + r.v;
36     return fabs((p1-p)^2)/r.v.mod();
37 }
38
39 punto project(punto a, punto b) { //Proyeccion de b sobre a
40     return ((a*b)/a.mod2()) * a;
41 }
42
43 vector <punto> inter_circ_line(recta r, circ c) {
44     vector <punto> ans; tipo dist = dist_point_line(c.c,r);
45     if(dist > c.r+EPS) return ans;
46     (c.c-r.p) * r.v != 0 ? r.p = r.p : r.p = r.p + r.v;
47     punto aux = c.c - r.p, dir = project(r.v,aux);
48     if(fabs(dist-c.r) <= EPS) {ans.pb(r.p + dir); return ans;}
49     tipo factor = sqrt(c.r*c.r - dist*dist)/dir.mod();
50     ans.pb(r.p + dir + factor * dir); ans.pb(r.p + dir - factor * dir);
51     return ans;
52 }
53
54 tipo intersection_area(circ a, circ b) {
55     punto aux = (b.c - a.c); tipo dist=aux.mod(), dist2=aux.mod2();
56     if(a.r + b.r - dist < -EPS) return 0;
57     if(fabs(a.r - b.r) - dist > -EPS) {
58         return min(a.r, b.r) * min(a.r, b.r) * 2*acosl(0); }
59     tipo alpha = acosl((dist2 + a.r*a.r - b.r*b.r) / (2 * dist * a.r));
60     tipo beta = acosl((dist2 + b.r*b.r - a.r*a.r) / (2 * dist * b.r));
61     tipo ans1 = (alpha - sinl(alpha+alpha)*0.5) * a.r * a.r;
62     tipo ans2 = (beta - sinl(beta+beta)*0.5) * b.r * b.r;
63     return ans1 + ans2;
64 }
65
66 vector <punto> inter_circ_circ(circ a, circ b) {
67     vector <punto> ans;
68     if(a.c==b.c) {
69         return abs(a.r-b.r) <= EPS ? vector<punto>{a.c,a.c,a.c} : ans; }
70     b.c = b.c - a.c; punto aux = a.c; a.c = a.c - a.c;
71     recta r(-2*b.c.x , -2*b.c.y , a.r*a.r - b.r*b.r + b.c.x*b.c.x + b.c.
72         y*b.c.y);
73     ans = inter_circ_line(r,a);

```

```

73     forn(i,ans.size()) ans[i] = ans[i] + aux; return ans;
74 }

```

4.9 Kd Tree

```

1 struct pt { // for 3D add z coordinate
2     ll x,y;
3     pt(ll x, ll y):x(x),y(y){}
4     pt(){}
5     ll norm2(){return *this**this;}
6     pt operator-(pt p){return pt(x-p.x,y-p.y);}
7     ll operator*(pt p){return x*p.x+y*p.y;}
8     // 2D from now on
9     bool operator<(pt p) const { // for convex hull
10         return x<p.x-EPS || (abs(x-p.x)<=EPS && y<p.y-EPS); }
11 };
12
13 ll manhattan(pt a, pt b){ return abs(a.x-b.x)+abs(a.y-b.y); }
14
15 // given a set of points, answer queries of nearest point in O(log(n))
16 bool onx(pt a, pt b){return a.x<b.x;}
17 bool ony(pt a, pt b){return a.y<b.y;}
18 struct Node {
19     pt pp;
20     ll x0=INF, x1=-INF, y0=INF, y1=-INF;
21     Node *first=0, *second=0;
22     ll distance(pt p){
23         ll x=min(max(x0,p.x),x1);
24         ll y=min(max(y0,p.y),y1);
25         return manhattan(pt(x,y), p);
26     }
27     Node(vector<pt>&& vp):pp(vp[0]){
28         for(pt p:vp){
29             x0=min(x0,p.x); x1=max(x1,p.x);
30             y0=min(y0,p.y); y1=max(y1,p.y);
31         }
32         if(SZ(vp)>1){
33             sort(all(vp),x1-x0>=y1-y0?onx:ony);
34             int m=SZ(vp)/2;
35             first=new Node({vp.begin(),vp.begin()+m});
36             second=new Node({vp.begin()+m,vp.end()});
37         }
38     }

```

```

39 };
40 struct KdTree {
41     Node* root;
42     KdTree(const vector<pt>& vp):root(new Node({all(vp)})) {}
43     pair<ll,pt> search(pt p, Node *node){
44         if(!node->first){
45             //avoid query point as answer
46             //if(p==node->pp) {INF,pt()};
47             return {manhattan(p, node->pp),node->pp};
48         }
49         Node *f=node->first, *s=node->second;
50         ll bf=f->distance(p), bs=s->distance(p);
51         if(bf>bs)swap(bf,bs),swap(f,s);
52         auto best=search(p,f);
53         if(bs<best.first) best=min(best,search(p,s));
54         return best;
55     }
56     // Return nearest point and its distance to p
57     pair<ll,pt> nearest(pt p){return search(p,root);}
58 };

```

4.10 Point In Poly

```

1 bool point_in_poly(vector <punto> &v, punto p) { // O(n) for convex
2     unsigned i, j, mi, mj, c = 0;
3     for(i = 0, j = v.size()-1; i < v.size(); j = i++) {
4         if((v[i].y <= p.y && p.y < v[j].y) || (v[j].y <= p.y && p.y < v[
5             i].y)) {
6             mi = i; mj = j; if(v[mi].y > v[mj].y) swap(mi,mj);
7             if(((p-v[mi]) ^ (v[mj]-v[mi])) < 0 ) c ^= 1;
8         }
9     }
10    return c;
11 }
12 bool point_in_poly_convex(vector <punto> &v, punto p) { // O(log n)
13     bool ans = true;
14     if( ((v[1]-v[0]) ^ (v[2]-v[1])) >= 0 ) reverse(all(v));
15     int a = 1, b = v.size()-1;
16     while(b-a > 1) {
17         int med = (a+b)/2;
18         if( ((v[med]-v[0]) ^ (p-v[med])) <= 0 ) a = med;
19         else b = med;

```

```

20     }
21     if( ((v[a]-v[0]) ^ (p - v[a])) > 0) ans = false;
22     if( ((v[b]-v[a]) ^ (p - v[b])) > 0) ans = false;
23     if( ((v[0]-v[b]) ^ (p - v[0])) > 0) ans = false;
24     return ans;
25 }

```

4.11 Template Punto

```

1  typedef long double tipo; //Cambiar a long long para operar en enteros
2  double EPS = (double)(1e-10);
3
4  struct punto { // Puede usarse para vectores
5      tipo x, y;
6      punto const operator -(const punto &p) const {return {x-p.x,y-p.y};}
7      punto const operator +(const punto &p) const {return {x+p.x,y+p.y};}
8      tipo operator *(const punto &p) const {return x*p.x + y*p.y;}
9      tipo operator ^(const punto &p) const {return x*p.y - y*p.x;}
10     bool operator == (const punto &p) const {
11         return (abs(x-p.x) < EPS && abs(y-p.y) < EPS); // Para double
12     }
13     bool operator <(punto p) const {return x != p.x ? x < p.x : y < p.y
14         ;}
15     tipo arg() {return atan2(y,x);}
16     tipo mod() {return sqrtl(x*x + y*y);}
17     tipo mod2() {return x*x + y*y;}
18 };
19
20 punto operator*(tipo k, const punto &p) {return {k*p.x, k*p.y};}
21
22 ostream &operator << (ostream &os, const punto &p) { //Para imprimir
23     return os << "(" << p.x << "," << p.y << ")";
24 }
25
26 istream &operator >> (istream &is, punto &p) { //Para leer
27     return is >> p.x >> p.y;
28 }
29
30 struct frac { // Por si es necesario trabajar con enteros
31     tipo n, d;
32     ll mcd(ll a, ll b) {
33         a = abs(a); b = abs(b);
34         while(a > 0 && b > 0) {

```

```

34         if(a >= b) a %= b;
35         else b %= a;
36     }
37     return a == 0 ? b : a;
38 }
39 frac(ll x, ll y) {
40     ll g = mcd(x,y);
41     n = x/g; d = y/g;
42     if(d < 0) d *= -1, n *= -1;
43 }
44 frac() {}
45 bool operator ==(frac &F) const {return n*F.d == d*F.n;}
46 bool operator < (frac &F) const {return n*F.d < d*F.n;}
47 };

```

5 Grafos

5.1 2Sat

```

1 // 2sat (2-satisfiability) - O(n + m) (Korasaju)
2
3 struct two_sat { // 2*x representa a x, y 2*x+1 a ~x
4     int tot;
5     vector<vector<int>> g, g_trans;
6     vb used, assignment;
7     vector<int> order, comp;
8
9     two_sat(int _tot): tot(_tot){
10         g.resize(tot); g_trans.resize(tot);
11     } // tot = total de nodos (normales + negados), en general n = tot/2
12
13     void dfs1(int v) {
14         used[v] = true;
15         for(auto u: g[v]) if(!used[u]) dfs1(u);
16         order.pb(v);
17     }
18
19     void dfs2(int v, int cl) { // Korasaju para encontrar las SCC
20         comp[v] = cl;
21         for(auto u: g_trans[v]) if(comp[u] == -1) dfs2(u, cl);
22     }
23
24     bool solve() {
25         order.clear(); used.assign(tot, false);
26         comp.assign(tot, -1);
27         forn(i, tot) if(!used[i]) dfs1(i);
28
29         int comp_act = 0;
30         forn(i, tot){
31             auto v = order[tot-i-1];
32             if(comp[v] == -1) dfs2(v, comp_act++);
33         }
34
35         assignment.assign(tot/2, false);
36         forn(i, tot/2){
37             if(comp[2*i] == comp[2*i + 1]) return false;
38             assignment[i] = comp[2*i] > comp[2*i+1]; // asignacion
39                                     greedy de variables

```

```

39     }
40     return true;
41 }
42
43 void add_edge(int from, int to){ // implicancia comun from->to
44     g[from].pb(to);
45     g_trans[to].pb(from); // agregar en el TRANSPUESTO
46 }
47
48 void add_or(int v1, int v2) { // agrega (v1 or v2)
49     add_edge(v1 ^ 1, v2); // ~v1 -> v2
50     add_edge(v2 ^ 1, v1); // ~v2 -> v1
51 }
52 // setear variable x en true/false: add_or(x, x)/add_or(~x, ~x)
53 };

```

5.2 Bellman Ford

```

1 // Description: Bellman-Ford algorithm for finding the shortest path
2 // from a source to all other nodes in a graph. It can also detect
3 // negative cycles.
4 // Time: O(VE)
5
6 typedef long long tipo;
7 const int MAXN = 3000;
8 tipo INF = (tipo)(1e18+7);
9
10 struct arista {
11     int x, y; tipo w; // Edge from x to y, w = weight
12 };
13
14 struct nodo {
15     int p; tipo d; //f -> parent, d -> distance
16 };
17
18 vector<nodo> ans(MAXN);
19 vector<int> ciclo;
20
21 bool bFord(vector<arista> &lista, int n, int start) {
22     int m = lista.size();
23     forn(i,n) ans[i].p = i, ans[i].d = INF;
24     ans[start].d = 0; int x;
25     forn(i,n) {

```

```

24     x = -1;
25     for(arista u : lista) {
26         if(ans[u.y].d > ans[u.x].d + u.w) {
27             ans[u.y].d = ans[u.x].d + u.w;
28             ans[u.y].p = u.x;
29             x = u.y;
30         }
31     }
32 }
33 if(x == -1) return false;
34 else {
35     for(i,n) x = ans[x].p;
36     for(int v = x ;; v = ans[v].p) {
37         ciclo.push_back(v);
38         if(v == x && ciclo.size() > 1) break;
39     }
40     reverse(all(ciclo));
41     return true;
42 }
43 }

```

5.3 Bfs

```

1  const int MAXN = 200005;
2  typedef long long tipo;
3  tipo INF = (tipo)(1e18+7);
4
5  struct nodo {
6      tipo d; bool visto; //d -> distance, visto -> seen
7  };
8
9  vector<nodo> BFS(int start, int n, vector<vector<int>> &g) {
10     vector<nodo> ans(n); queue<int> q;
11     for(i, n) ans[i] = {INF, false};
12     ans[start] = {0, true}; q.push(start);
13     while(!q.empty()) {
14         int v = q.front(); q.pop();
15         for(int u : g[v]) {
16             if(ans[u].visto) continue;
17             ans[u] = {ans[v].d+1, true}; q.push(u);
18         }
19     }
20     return ans;

```

```

21 }

```

5.4 Biconnected

```

1  struct edge{ ll u, v, id; };
2
3  struct graph{
4      int n;
5      vector<vii> adj; // (to, id)
6      vector<edge> edges;
7      graph(int _n) : n(_n), adj(_n) {}
8
9      void add_edge(ll u, ll v){
10         ll id = SZ(edges);
11         adj[u].pb({v, id}); adj[v].pb({u, id});
12         edges.pb({u, v, id});
13     }
14
15     int add_node(){ adj.pb({}); return n++; }
16     vii& operator[] (ll u) { return adj[u]; }
17 };
18
19 struct BCC{
20     int n; graph adj;
21     vector<vi> comps;
22     vi num, low, art, stk, bridge;
23     BCC(graph &_adj) : n(_adj.n), adj(_adj){
24         num.resize(n), low.resize(n), art.resize(n), bridge.resize(SZ(
25             adj.edges));
26         for (ll u = 0, t; u < n; ++u) if (!num[u]) dfs(u, -1, t = 0);
27     }
28
29     void dfs(ll v, ll p, ll &t){
30         num[v] = low[v] = ++t;
31         stk.pb(v);
32
33         for(auto [u, id] : adj[v]) if (u != p){
34             if (!num[u]){
35                 dfs(u, v, t);
36                 low[v] = min(low[v], low[u]);
37
38                 if(low[u] > num[v]) bridge[id] = true;
39                 if(low[u] >= num[v]){

```

```

39         art[v] = (num[v] > 1 || num[u] > 2);
40
41         comps.pb({v});
42         while (comps.back().back() != u)
43             comps.back().pb(stk.back()), stk.pop_back();
44     }
45 }
46 else low[v] = min(low[v], num[u]);
47 }
48 }
49
50 // build the block cut tree
51 pair<vi, graph> build_tree(){
52     graph tree(0); vi id(n);
53
54     forn(v, n) if (art[v]) id[v] = tree.add_node();
55
56     for (auto &comp : comps){
57         ll node = tree.add_node();
58         for(ll v: comp){
59             if (!art[v]) id[v] = node;
60             else tree.add_edge(node, id[v]);
61         }
62     }
63     return {id, tree};
64 }
65 };

```

5.5 Componentes Fuertemente Conexas

```

1  /*
2  Algoritmo para hallar las componentes fuertemente conexas.
3  Una componente es fuertemente conexa cuando para todo nodo
4  perteneciente a la componente, se puede llegar a cualquier
5  otro nodo tambien perteneciente a la componente.
6  */
7  struct SCC {
8      int n, scc;
9      vector<vi> g, gr, ans;
10     vb visto;
11     vi order, comp_act, component;
12
13     SCC(vector<vi> &_g): n(SZ(_g)) {

```

```

14     g = _g;
15     gr.resize(n), visto.resize(n), component.resize(n);
16     forn(v, n) for(auto u: g[v]) gr[u].pb(v); // me lo creo aca el
17         grafo traspuesto
18
19     find_scc();
20
21     void DFS1 (int v) {
22         visto[v] = true;
23         for (int u : g[v]) if(!visto[u]) DFS1(u);
24         order.pb(v);
25     }
26
27     void DFS2 (int v) {
28         visto[v] = true;
29         comp_act.pb(v);
30         for (int u : gr[v]) if(!visto[u]) DFS2(u);
31     }
32
33     void find_scc() {
34         fill(all(visto), false);
35         forn(i, n) if(!visto[i]) DFS1(i);
36         fill(all(visto), false);
37         forn(i, n) {
38             int v=order[n-i-1];
39             if(!visto[v]) {
40                 DFS2(v);
41                 ans.pb(comp_act);
42                 comp_act.clear();
43             }
44         }
45         scc = SZ(ans); // cantidad de scc's
46         forn(i, scc) for(auto v: ans[i]) component[v] = i;
47     }
48 };

```

5.6 Dfs

```

1  const int MAXN = 200005;
2  vector< vector<int> > g; // graph represented as an adjacency list
3  vector <bool> visto(MAXN, false);
4

```



```

5 void dfs(int v) {
6     visto[v] = true;
7     for (int u : g[v]) if (!visto[u]) dfs(u);
8 }

```

5.7 Dijkstra

```

1 typedef ll tipo;
2 const int MAXN = 200005;
3
4 struct arista {
5     int x; tipo w; // x -> next node, w = weight
6 };
7
8 struct nodo {
9     tipo d, v, a; // d -> distance, v -> actual node, a = previous node
10    bool operator<(const nodo& x) const {return d > x.d;}
11 };
12
13 vector<nodo> Dijkstra(int start, int n, vector<vector<arista>> &g) {
14     vector<nodo> ans(n);
15     vector<bool> visto(n, false);
16     priority_queue<nodo> p; p.push({0,start,-1});
17     while(!p.empty()) {
18         nodo it=p.top(); p.pop();
19         if(visto[it.v]) continue;
20         else {
21             ans[it.v] = it; visto[it.v] = true;
22             for(arista u : g[it.v]) {
23                 if(!visto[u.x]) p.push({it.d + u.w, u.x, it.v});
24             }
25         }
26     }
27     return ans;
28 }

```

5.8 Dominator Tree

```

1 // Dominator Tree
2 // Description: Given a directed graph, find the dominator tree of the
   graph
3 // Time: O(nlogn)
4 // Usage: Dominator_Tree dt(n,g,root); // n --> number of nodes, g -->
   graph, root --> root of the tree

```

```

5 // If u is an ancestor of v in the dominator tree, then u dominates v
6
7
8 #define rz(n) resize(n)
9
10 struct Dominator_Tree {
11     vector<vector<int>> g, tree, rg, bucket;
12     vector<int> sdom, par, dom, dsu, label;
13     vector<int> arr, rev;
14     int n, T = 0, root;
15
16     int Find(int u,int x=0) {
17         if(u==dsu[u])return x?-1:u;
18         int v = Find(dsu[u],x+1);
19         if(v<0)return u;
20         if(sdom[label[dsu[u]]]<sdom[label[u]]) {
21             label[u] = label[dsu[u]];
22         }
23         dsu[u] = v;
24         return x ? v : label[u];
25     }
26
27     void Union(int u,int v) { //Add an edge u-->v
28         dsu[v] = u;
29     }
30
31     void dfs0(int u) {
32         T++; arr[u] = T; rev[T] = u;
33         label[T] = T; sdom[T] = T; dsu[T] = T;
34         for(int i = 0; i < SZ(g[u]); i++) {
35             int w = g[u][i];
36             if(!arr[w]) {
37                 dfs0(w);
38                 par[arr[w]]=arr[u];
39             }
40             rg[arr[w]].pb(arr[u]);
41         }
42     }
43
44     Dominator_Tree(int _n, vector<vector<int>> _g, int _root) {
45         n = _n; root = _root;
46         g = _g; tree.rz(n+5); rg.rz(n+5); bucket.rz(n+5); sdom.rz(n+5);
47         par.rz(n+5); dom.rz(n+5); dsu.rz(n+5); label.rz(n+5);

```

```

48     arr.rz(n+5); rev.rz(n+5);
49     dfs0(root);
50
51
52     for(int i = n; i >= 1; i--) {
53         for(int j = 0; j < SZ(rg[i]); j++) {
54             sdom[i] = min(sdom[i], sdom[Find(rg[i][j])]);
55         }
56         if(i>1) bucket[sdom[i]].pb(i);
57         for(int j = 0; j < SZ(bucket[i]); j++) {
58             int w = bucket[i][j], v = Find(w);
59             if(sdom[v] == sdom[w]) dom[w] = sdom[w];
60             else dom[w] = v;
61         }
62         if(i>1) Union(par[i], i);
63     }
64
65     forr(i, 2, n+1){
66         if(dom[i] != sdom[i]) dom[i] = dom[dom[i]];
67         tree[rev[i]].pb(rev[dom[i]]);
68         tree[rev[dom[i]]].pb(rev[i]);
69     }
70
71 }
72 };

```

5.9 Dynamic Connectivity

```

1 // Imprime la cantidad de componentes conexas
2 // Resuelve el problema offline, para agregar una query es necesario
3 // agregar DynCon.query() en el instante que se quiera conocer la
4 // respuesta.
5
6
7 struct UnionFind {
8     int n, comp;
9     vector<int> uf, si, c;
10    UnionFind(int n=0):n(n), comp(n), uf(n), si(n, 1){
11        forr(i, 0, n) uf[i]=i;
12    int find(int x){return x==uf[x]?x:find(uf[x]);}
13    bool join(int x, int y){
14        if((x=find(x))==(y=find(y)))return false;
15        if(si[x]<si[y])swap(x, y);

```

```

16        si[x]+=si[y]; uf[y]=x; comp--; c.pb(y);
17        return true;
18    }
19    int snap(){return c.size();}
20    void rollback(int snap){
21        while(c.size()>snap){
22            int x=c.back(); c.pop_back();
23            si[uf[x]]-=si[x]; uf[x]=x; comp++;
24        }
25    }
26 };
27 enum {ADD, DEL, QUERY};
28 struct Query {int type, x, y;};
29 struct DynCon {
30     vector<Query> q;
31     UnionFind dsu;
32     vector<int> mt;
33     map<pair<int, int>, int> last;
34     DynCon(int n):dsu(n){}
35     void add(int x, int y){
36         if(x>y)swap(x, y);
37         q.pb((Query){ADD, x, y}); mt.pb(-1); last[{x, y}]=q.size()-1;
38     }
39     void remove(int x, int y){
40         if(x>y)swap(x, y);
41         q.pb((Query){DEL, x, y});
42         int pr=last[{x, y}]; mt[pr]=q.size()-1; mt.pb(pr);
43     }
44     void query(){q.pb((Query){QUERY, -1, -1}); mt.pb(-1);}
45     void process(){ // answers all queries in order
46         if(!q.size())return;
47         forr(i, 0, q.size())if(q[i].type==ADD&&mt[i]<0)mt[i]=q.size();
48         go(0, q.size());
49     }
50     void go(int s, int e){
51         if(s+1==e){
52             if(q[s].type==QUERY) // answer query using DSU
53                 cout << dsu.comp << "\n";
54             return;
55         }
56         int k=dsu.snap(), m=(s+e)/2;
57         for(int i=e-1; i>=m; --i)if(mt[i]>=0&&mt[i]<s)dsu.join(q[i].x, q[i].y);

```

```

58     go(s,m);dsu.rollback(k);
59     for(int i=m-1;i>=s;--i)if(mt[i]>=e)dsu.join(q[i].x,q[i].y);
60     go(m,e);dsu.rollback(k);
61 }
62 };

```

5.10 Eulerian Cycle

```

1  // Todos tienen que cumplir: SZ(path) == tot_edges+1
2  //
3  // Existencia directed:
4  // Path: Que el start tenga out_deg=in_deg+1
5  //       Que el final tenga in_deg=out_deg+1
6  //       Que el resto tenga in_deg=out_deg
7  //
8  // Cycle: Todos los nodos in_deg=out_deg
9  //
10 // Existencia undirected:
11 // Path: Que el start y final tengan deg impar, el resto todo par
12 // Cycle: Todos los nodos in_deg=out_deg
13
14 // Directed version (uncomment commented code for undirected)
15 struct edge {
16     int y;
17     // list<edge>::iterator rev;
18     edge(int _y):y(_y){}
19 };
20 list<edge> g[MAXN];
21 void add_edge(int a, int b){
22     g[a].push_front(edge(b));//auto ia=g[a].begin();
23     // g[b].push_front(edge(a));auto ib=g[b].begin();
24     // ia->rev=ib;ib->rev=ia;
25 }
26 vector<int> p;
27 void go(int x){
28     while(SZ(g[x])){
29         int y=g[x].front().y;
30         //g[y].erase(g[x].front().rev);
31         g[x].pop_front();
32         go(y);
33     }
34     p.pb(x);
35 }

```

```

36 vector<int> get_path(int x){ // get a path that begins in x
37 // check that a path exists from x before calling to get_path!
38     p.clear();go(x);reverse(all(p));
39     return p;
40 }

```

5.11 Blossom

```

1  /* Encuentra un matching en un grafo no bipartito
2  * en O(nm) mediante el algoritmo de Edmond's blossom
3  * Los matchings estan en (i, mate[i]), con i < mate[i]
4  */
5  vi Blossom(vector<vi> &g) {
6      int n = SZ(g), timer = -1;
7      vi mate(n, -1), label(n), parent(n),
8          orig(n), aux(n, -1), q;
9      auto lca = [&](int x, int y) {
10         for (timer++; ; swap(x, y)) {
11             if (x == -1) continue;
12             if (aux[x] == timer) return x;
13             aux[x] = timer;
14             x = (mate[x] == -1 ? -1 : orig[parent[mate[x]]]);
15         }
16     };
17     auto blossom = [&](int v, int w, int a) {
18         while (orig[v] != a) {
19             parent[v] = w; w = mate[v];
20             if (label[w] == 1) label[w] = 0, q.pb(w);
21             orig[v] = orig[w] = a; v = parent[w];
22         }
23     };
24     auto augment = [&](int v) {
25         while (v != -1) {
26             int pv = parent[v], nv = mate[pv];
27             mate[v] = pv; mate[pv] = v; v = nv;
28         }
29     };
30     auto bfs = [&](int root) {
31         fill(all(label), -1);
32         iota(all(orig), 0);
33         q.clear();
34         label[root] = 0; q.pb(root);
35         forn(i, SZ(q)){

```

```

36     int v = q[i];
37     for (auto x : g[v]) {
38         if (label[x] == -1) {
39             label[x] = 1; parent[x] = v;
40             if (mate[x] == -1)
41                 return augment(x), 1;
42             label[mate[x]] = 0; q.pb(mate[x]);
43         } else if (label[x] == 0 && orig[v] != orig[x]) {
44             int a = lca(orig[v], orig[x]);
45             blossom(x, v, a); blossom(v, x, a);
46         }
47     }
48 }
49 return 0;
50 };
51 // Time halves if you start with (any) maximal matching.
52 forn(i, n) if(mate[i] == -1) bfs(i);
53 return mate;
54 }

```

5.12 Dinic

```

1 // Dinic: Max Flow en  $O(V^2 E)$ . Para el grafo bipartito con source
2 // y sink dummy, funciona en  $O(\sqrt{V} E)$ . Equivalente a Hopcroft-Karp.
3
4 // Matching: aristas saturadas (que no incluyan source/sink)
5 // Min cut: nodos con  $\text{dist} \geq 0$  vs nodos con  $\text{dist} < 0$ 
6 // MVC: Nodos izquierda con  $\text{dist} < 0$  + nodos derecha con  $\text{dist} > 0$ 
7 // Maximum Independent Set: complemento de MVC (N-MVC)
8 struct Dinic{
9     int nodes,src,dst;
10    vector<int> dist,q,work;
11    struct edge { int to,rev; ll f,cap; };
12    vector<vector<edge>> g;
13    Dinic(int x):nodes(x),dist(x),q(x),work(x),g(x){}
14    void add_edge(int s, int t, ll cap){
15        g[s].pb({t,SZ(g[t]),0,cap});
16        g[t].pb({s,SZ(g[s])-1,0,0});
17    }
18    bool dinic_bfs(){
19        fill(all(dist),-1);dist[src]=0;
20        int qt=0;q[qt++]=src;
21        forn(qh, qt){

```

```

22            int u=q[qh];
23            forn(i,SZ(g[u])){
24                edge &e=g[u][i];int v=g[u][i].to;
25                if(dist[v]<0&&e.f<e.cap)dist[v]=dist[u]+1,q[qt++]=v;
26            }
27        }
28        return dist[dst]>=0;
29    }
30    ll dinic_dfs(int u, ll f){
31        if(u==dst)return f;
32        for(int &i=work[u];i<SZ(g[u]);i++){
33            edge &e=g[u][i];
34            if(e.cap<=e.f)continue;
35            int v=e.to;
36            if(dist[v]==dist[u]+1){
37                ll df=dinic_dfs(v,min(f,e.cap-e.f));
38                if(df>0){e.f+=df;g[v][e.rev].f-=df;return df;}
39            }
40        }
41        return 0;
42    }
43    ll max_flow(int _src, int _dst){
44        src=_src;dst=_dst;
45        ll result=0;
46        while(dinic_bfs()){
47            fill(all(work),0);
48            while(ll delta=dinic_dfs(src,INF))result+=delta;
49        }
50        return result;
51    }
52 };

```

5.13 Hopcroft Karp

```

1 /**
2  * Fast bipartite matching algorithm. Graph $g$ should be a list
3  * of neighbors of the left partition, and $btoa$ should be a vector
4  * full of
5  * -1's of the same size as the right partition. Returns the size of
6  * the matching. $btoa[i]$ will be the match for vertex $i$ on the right
7  * side,
8  * or $-1$ if it's not matched.
9  * Usage: vi btoa(m, -1); hopcroftKarp(g, btoa);

```

```

8  * Time:  $O(\sqrt{V}E)$ 
9  */
10 bool dfs(int a, int L, vector<vi> &g, vi &btoa, vi &A, vi &B) {
11     if(A[a] != L) return 0;
12     A[a] = -1;
13     for(auto b: g[a]) if(B[b] == L + 1) {
14         B[b] = 0;
15         if (btoa[b] == -1 || dfs(btoa[b], L+1, g, btoa, A, B))
16             return btoa[b] = a, 1;
17     }
18     return 0;
19 }
20
21 int hopcroftKarp(vector<vi> &g, vi &btoa) { // bipartite matching rapido
22     int res = 0;
23     vi A(SZ(g)), B(SZ(btoa)), cur, next;
24     for(;;) {
25         fill(all(A), 0); fill(all(B), 0);
26         /// Find the starting nodes for BFS (i.e. layer 0).
27         cur.clear();
28         for(auto a : btoa) if(a != -1) A[a] = -1;
29         forn(a, SZ(g)) if(A[a] == 0) cur.pb(a);
30         /// Find all layers using bfs.
31         for(int lay = 1;; lay++) {
32             bool islast = 0;
33             next.clear();
34             for(auto a: cur) for(auto b: g[a]) {
35                 if (btoa[b] == -1) {
36                     B[b] = lay; islast = 1;
37                 } else if (btoa[b] != a && !B[b]) {
38                     B[b] = lay; next.pb(btoa[b]);
39                 }
40             }
41             if(islast) break;
42             if(next.empty()) return res;
43             for(auto a : next) A[a] = lay;
44             cur.swap(next);
45         }
46         /// Use DFS to scan for augmenting paths.
47         forn(a, SZ(g)) res += dfs(a, 0, g, btoa, A, B);
48     }
49 }

```

5.14 Hungarian

```

1  /* Algoritmo  $O(n^3)$  para assignment problem
2  * Inspirando el el notebook vasito, quien se inspiro
3  * en http://e-maxx.ru/algorithm/assignment\_hungary
4  */
5  typedef long double td; typedef vector<td> vd;
6  const td INF=1e100; //for maximum set INF to 0, and negate costs
7  bool zero(td x){return fabs(x)<1e-9;} //change to x==0, for ints/ll
8  struct Hungarian{
9      int n; vector<vd> cs; vi L, R;
10     Hungarian(int N, int M):n(max(N,M)),cs(n,vd(n)),L(n),R(n){
11         forn(x,N)forn(y,M)cs[x][y]=INF;
12     }
13     void set(int x,int y,td c){cs[x][y]=c;}
14     td assign() {
15         int mat = 0; vd ds(n), u(n), v(n); vi dad(n), sn(n);
16         forn(i,n)u[i]=*min_element(all(cs[i]));
17         forn(j,n){v[j]=cs[0][j]-u[0];forr(i,1,n)v[j]=min(v[j],cs[i][j]-u[i]);}
18         L=R=vi(n, -1);
19         forn(i,n)forn(j,n)
20             if(R[j]==-1&&zero(cs[i][j]-u[i]-v[j])){L[i]=j;R[j]=i;mat++;break;}
21         for(;mat<n;mat++){
22             int s=0, j=0, i;
23             while(L[s] != -1)s++;
24             fill(all(dad),-1);fill(all(sn),0);
25             forn(k,n)ds[k]=cs[s][k]-u[s]-v[k];
26             for(;;){
27                 j = -1;
28                 forn(k,n)if(!sn[k]&&(j==-1||ds[k]<ds[j]))j=k;
29                 sn[j] = 1; i = R[j];
30                 if(i == -1) break;
31                 forn(k,n)if(!sn[k]){
32                     auto new_ds=ds[j]+cs[i][k]-u[i]-v[k];
33                     if(ds[k] > new_ds){ds[k]=new_ds;dad[k]=j;}
34                 }
35             }
36             forn(k,n)if(k!=j&&sn[k]){auto w=ds[k]-ds[j];v[k]+=w,u[R[k]]-=w;}
37             u[s] += ds[j];
38             while(dad[j]>=0){int d = dad[j];R[j]=R[d];L[R[j]]=j;j=d;}

```

```

39     R[j]=s;L[s]=j;
40 }
41     td value=0;for(n,i,n)value+=cs[i][L[i]];
42     return value;
43 }
44 };

```

5.15 Maximum Bipartite Matching

```

1  /*
2  Algoritmo de Kuhn para bipartite matching, se le pasa el grafo
3  y encuentra el matching maximo en O(VE).
4  La pareja de cada nodo se guarda en match[i], sino -1
5  */
6
7  struct bipartite_matching {
8      int n;
9      vector<vi> g; // 0-indexed
10     vb vis; vi match;
11
12     bipartite_matching(int _n, vector<vi> _g): n(_n), g(_g) {
13         vis.resize(n), match.resize(n);
14     }
15
16     bool dfs(int node){
17         if(vis[node])return 0;
18         vis[node] = 1;
19         for(auto nx : g[node]){
20             if(match[nx]==-1 || dfs(match[nx])){
21                 match[node] = nx; match[nx] = node;
22                 return 1;
23             }
24         }
25         return 0;
26     }
27
28     int solve() { // toma los nodos de 0 a n-1
29         fill(all(match),-1);
30         while(true) {
31             fill(all(vis),false);
32             bool cont = 0;
33             for(n,i,n) if(match[i] == -1) cont |= dfs(i);
34             if(cont==0) break;

```

```

35     }
36     int matching = 0;
37     for(n,i,n) if(match[i] != -1) matching++;
38     return matching; // hay que dividir por 2 si quiero la cantidad de
39                         parejas
40 }
41 };

```

5.16 Max Flow

```

1  /*
2  Algoritmo de Edmonds-Karp, halla el max flow en O(VE^2).
3  Para eso, va eligiendo caminos de aumento con la menor cantidad de
4  aristas
5  en el grafo residual.
6
7  El min_cut esta formado por las aristas que unen un nodo alcanzable por
8  s en el grafo residual final, con un nodo inalcanzable.
9  */
10
11 typedef long long tipo; // el tipo en que se mide el flow
12
13 struct max_flow { // Edmonds-Karp, O(VE^2)
14     int n;
15     vector<vector<int>> g;
16     vector<vector<tipo>> cap;
17     // si n es grande (>5000), hay que usar un map de capacidades
18
19     max_flow(int n): n(n){
20         g.resize(n), cap.resize(n, vi(n));
21     }
22     void add_edge(int x, int y, tipo z){
23         g[x].pb(y), g[y].pb(x);
24         cap[x][y] += z;
25     }
26     tipo bfs(int s, int t, vector<int> &parent){
27         fill(all(parent), -1);
28         queue<pair<int, tipo>> q; q.push({s, INF});
29         parent[s] = s;
30
31         while(!q.empty()){
32             int cur = q.front().first;
33             tipo flow = q.front().second; q.pop();

```

```

33     for(int next: g[cur]){
34         if(parent[next] == -1 && cap[cur][next]) {
35             parent[next] = cur;
36             tipo new_flow = min(flow, cap[cur][next]);
37             if(next == t) return new_flow;
38             q.push({next,new_flow});
39         }
40     }
41 }
42 return 0; // no encuentre aug paths
43 }
44
45 tipo get_max_flow(int s, int t){
46     tipo flow = 0, new_flow;
47     vector<int> parent(n);
48
49     while((new_flow = bfs(s, t, parent))){
50         flow += new_flow;
51         int cur = t;
52         while(cur != s){
53             int prev = parent[cur];
54             cap[prev][cur] -= new_flow;
55             cap[cur][prev] += new_flow;
56             cur = prev;
57         }
58     }
59     return flow;
60 }
61
62 void remove_dups(vector<int> &a){
63     sort(all(a)); a.erase(unique(all(a)), a.end());
64 }
65
66 vector<pair<int, int>> get_min_cut(int s, int t) {
67     // cambiar ans a set<pair<int, int>>
68     // si no se quiere usar remove dups
69     vector<pair<int, int>> ans;
70     for(auto &x: g) remove_dups(x);
71
72     get_max_flow(s, t); // si ya se corrio afuera, comentar
73
74     vector<int> parent(n);
75     bfs(s, t, parent); // hago bfs en el grafo residual final

```

```

76     forn(v, n){
77         for(auto u: g[v]){
78             if(parent[v] != -1 && parent[u] == -1){
79                 ans.pb({v, u});
80             }
81         }
82     }
83     return ans;
84 }
85 };
86
87 // encontrar paths disjuntos (si tienen cap 1)
88 // ini_cap es lo mismo que cap antes de correr el flow.
89 vector<int> path;
90 bool find_paths(int v, max_flow &mf){
91     path.pb(v);
92     if(v == mf.n-1) return true;
93     for(int u : mf.g[v]) {
94         if(mf.ini_cap[v][u] && !mf.cap[v][u]){
95             mf.cap[v][u] = 1; // marco la arista como usada (
96                 dessaturandola)
97             if(find_paths(u, mf)) return true;
98             mf.cap[v][u] = 0;
99         }
100     }
101     path.pop_back();
102     return false;
103 }

```

5.17 Max Flow Min Cost

```

1  typedef long long tipo;
2  typedef double tipo_cost;
3  const int MAXN = 505;
4  tipo INF = (tipo)(1e9+7);
5  tipo_cost MAX_COST = (double)(1e16);
6
7  struct arista {
8      int v; // Next node
9      tipo_cost c; // Cost
10 };
11
12 vector < vector<arista> > g(MAXN);

```

```

13 map < pair<int,int>, tipo > cap;
14 map < pair<int,int>, tipo_cost > cost;
15
16 void add_edge(int x, int y, tipo z, tipo_cost c) {
17     g[x].pb({y,c});
18     g[y].pb({x,-1.0*c});
19     cap[{x,y}] += z;
20     cost[{x,y}] += c;
21     cost[{y,x}] -= c;
22 }
23
24 void Bellman_Ford(int s, int t, vector <int> &parent, vector <tipo_cost>
    &d) {
25     fill(all(parent),-1);
26     fill(all(d),MAX_COST);
27     vector <bool> inq(MAXN,false);
28     queue <int> q; q.push(s); parent[s] = s; inq[s] = true; d[s] = 0;
29
30     while(q.size() > 0) {
31         int u = q.front(); q.pop(); inq[u] = false;
32         for(arista next : g[u]) {
33             if(cap[{u,next.v}]>0 && d[next.v] > d[u] + next.c) {
34                 d[next.v] = d[u] + next.c;
35                 parent[next.v] = u;
36                 if(!inq[next.v]) {
37                     inq[next.v] = true;
38                     q.push(next.v);
39                 }
40             }
41         }
42     }
43 }
44
45
46 tipo_cost max_flow_min_cost(int s, int t) {
47     tipo flow = 0;
48     tipo_cost min_cost = 0;
49     vector <int> parent(MAXN);
50     vector <tipo_cost> d(MAXN);
51
52     while(true) {
53         Bellman_Ford(s,t,parent,d);
54         if(d[t] == MAX_COST) break;

```

```

55     tipo new_flow = INF;
56     int cur = t;
57     while(cur != s) {
58         new_flow = min(new_flow,cap[{parent[cur],cur}]);
59         cur = parent[cur];
60     }
61
62     flow += new_flow;
63     min_cost += d[t] * (tipo_cost)(new_flow);
64     cur = t;
65     while(cur != s) {
66         int prev = parent[cur];
67         cap[{prev,cur}] -= new_flow;
68         cap[{cur,prev}] += new_flow;
69         cur = prev;
70     }
71 }
72
73 return min_cost;
74 }

```

5.18 Max Flow Min Cost Multiedge

```

1 typedef long long tipo;
2 typedef long long tipo_cost;
3 const int MAXN = 505;
4 tipo INF = (tipo)(1e9+7);
5 tipo_cost MAX_COST = (tipo_cost)(1e16);
6
7 struct arista {
8     int v;
9     tipo cap;
10    tipo_cost c; // Cost
11    int id;
12 };
13
14 struct prev_node {
15     int prev, id;
16     void setting() {prev = -1, id = -1;}
17 };
18
19 vector < vector<arista> > g(MAXN);
20 vector <tipo> cap;

```



```

21
22 void add_edge(int x, int y, tipo z, tipo_cost c) {
23     int id = cap.size();
24     g[x].pb({y,z,c,id});
25     g[y].pb({x,0,-1*c,id+1});
26     cap.pb(z); cap.pb(0);
27 }
28
29 void Bellman_Ford(int s, int t, vector <prev_node> &parent, vector <
    tipo_cost> &d) {
30     for(prev_node u : parent) u.setting();
31     fill(all(d),MAX_COST);
32     vector <bool> inq(MAXN,false);
33     queue <int> q; q.push(s); parent[s] = {s,-1}; inq[s] = true; d[s] = 0;
34
35     while(q.size() > 0) {
36         int u = q.front(); q.pop(); inq[u] = false;
37         for(arista next : g[u]) {
38             if(cap[next.id]>0 && d[next.v] > d[u] + next.c) {
39                 d[next.v] = d[u] + next.c;
40                 parent[next.v] = {u,next.id};
41                 if(!inq[next.v]) {
42                     inq[next.v] = true;
43                     q.push(next.v);
44                 }
45             }
46         }
47     }
48 }
49
50
51 tipo_cost max_flow_min_cost(int s, int t, int total) {
52     tipo flow = 0;
53     tipo_cost min_cost = 0;
54     vector <prev_node> parent(MAXN);
55     vector <tipo_cost> d(MAXN);
56
57     while(true) {
58         Bellman_Ford(s,t,parent,d);
59         if(d[t] == MAX_COST) break;
60
61         tipo new_flow = INF;
62         int cur = t;

```

```

63         while(cur != s) {
64             new_flow = min(new_flow,cap[parent[cur].id]);
65             cur = parent[cur].prev;
66         }
67
68         flow += new_flow;
69         min_cost += d[t] * (tipo_cost)(new_flow);
70         cur = t;
71         while(cur != s) {
72             int prev = parent[cur].prev;
73             int id = parent[cur].id;
74             cap[id] -= new_flow;
75             cap[(id^1)] += new_flow;
76             cur = prev;
77         }
78     }
79     if(flow < total) return -1;
80     return min_cost;
81 }

```

5.19 Min Cost Flow

```

1  /*
2  Max flow min cost, trabaja con Dijkstra y funciones potenciales
3  en  $O(n^3 m)$ .
4  Tolera multiedge, pero no ciclos negativos de costo (esto es complicado
   en gral.)
5
6  Si es necesario el min cost flow (para un flow K fijo),
7  agregar una sink mas t', unir t->t' con una arista de cap K y costo 0
8  y correr flow entre s y t'.
9  */
10 typedef ll tf; // tipo que se usa para el flow
11 typedef ll tc; // tipo que se usa para el cost
12 const tf INFFLOW=1e9;
13 const tc INFCOST=1e9;
14 struct MCF {
15     int n;
16     vector<tc> prio, pot; vector<tf> curflow; vector<int> prevedge,
        prevnode;
17     priority_queue<pair<tc, int>, vector<pair<tc, int>>, greater<pair<tc,
        int>>> q;
18     struct edge{int to, rev; tf f, cap; tc cost;};

```

```

19 vector<vector<edge>> g;
20 MCF(int _n:n(_n),prio(_n),pot(_n),curflow(_n),prevedge(_n),prevnode(
    _n),g(_n)){
21 void add_edge(int s, int t, tf cap, tc cost) {
22     g[s].pb({t,SZ(g[t]),0,cap,cost});
23     g[t].pb({s,SZ(g[s])-1,0,0,-cost});
24 }
25 pair<tf,tc> get_flow(int s, int t) {
26     tf flow=0; tc flowcost=0;
27     while(1){
28         q.push({0, s});
29         fill(all(prio),INFCOST);
30         prio[s]=0; curflow[s]=INFFLOW;
31         while(!q.empty()) {
32             auto cur=q.top(); q.pop();
33             tc d=cur.first; int u=cur.second;
34             if(d!=prio[u]) continue;
35             forn(i, SZ(g[u])){
36                 edge &e=g[u][i];
37                 int v=e.to;
38                 if(e.cap<=e.f) continue;
39                 tc nprio=prio[u]+e.cost+pot[u]-pot[v];
40                 if(prio[v]>nprio) {
41                     prio[v]=nprio;
42                     q.push({nprio, v});
43                     prevnode[v]=u; prevedge[v]=i;
44                     curflow[v]=min(curflow[u], e.cap-e.f);
45                 }
46             }
47         }
48         if(prio[t]==INFCOST) break;
49         forn(i, n) pot[i]+=prio[i];
50         tf df=min(curflow[t], INFFLOW-flow);
51         flow+=df;
52         for(int v=t; v!=s; v=prevnode[v]) {
53             edge &e=g[prevnode[v]][prevedge[v]];
54             e.f+=df; g[v][e.rev].f-=df;
55             flowcost+=df*e.cost;
56         }
57     }
58     return {flow,flowcost};
59 }
60 };

```

5.20 Floyd Warshall

```

1 //~ Minino camino entre todos los nodos
2 //~ O(n^3) - Tambien detecta ciclos negativos
3
4 void floyd(vector<vector<int>>&matriz, ll n) {
5     forn(k,n) forn(i,n) forn(j,n) {
6         matriz[i][j]=min(matriz[i][j],matriz[i][k]+matriz[k][j]);
7     }
8 }

```

5.21 Kruskal

```

1 struct arista {
2     int u, v;
3     ll w;
4     bool operator<(const arista& other) const { return w < other.w; }
5 };
6 struct DSU {
7     vi link, sz;
8
9     DSU(int tam) {
10         link.resize(tam+5), sz.resize(tam+5);
11         forn(i, tam+5) link[i] = i, sz[i] = 1;
12     }
13
14     ll find(ll x){ return link[x] = (link[x] == x ? x : find(link[x])); }
15
16     bool same(ll a, ll b) { return find(a) == find(b); }
17
18     void join(ll a, ll b) {
19         a = find(a), b = find(b);
20         if(a == b) return;
21         if(sz[a] < sz[b]) swap(a,b);
22         sz[a] += sz[b];
23         link[b] = a;
24     }
25 };
26
27 vector<arista> kruskal(int n, vector<arista> &g) {
28     sort(all(g)); DSU dsu(n);
29     vector<arista> mst;
30     for(arista e : g) {

```

```

30     if(!dsu.same(e.v, e.u)){
31         dsu.join(e.v, e.u); mst.pb(e);
32     }
33 }
34 return mst;
35 }

```

5.22 Prim

```

1 //Prim Algorithm
2
3 //Dado un grado no direccionado, halla el arbol tal que la suma de
4 //los pesos de las aristas sea minimo
5
6 //Metodologia: insertamos un nodo cualquiera, agregamos en un priority
7 //queue las aristas, ordenadas desde el menor peso hasta el mayor.
8 //insertamos la menor arista y agregamos el otro extremo como un nodo
9 //nuevo. Iteramos estos pasos hasta obtener las n-1 aristas.
10
11 //El algoritmo retorna la lista con las arista que generan el minimum
12 //Spanning Tree
13
14 vector <pair<ll,ll>> v[500005];
15 vector <pair<ll,ll>> a;
16 vector <bool> visto(500005,false);
17 map <pair<ll,ll>,bool> mapa;
18
19 struct arista
20 {
21     ll n, m; // n=nodo padre, m=nodo nuevo
22     ll peso;
23 };
24
25 struct compare
26 {
27     bool operator()(arista a, arista b)
28     {
29         if(a.peso!=b.peso) return a.peso<b.peso;
30         else return true;
31     }
32 };
33
34 vector <pair<ll,ll>> Prim(ll n)

```

```

35 {
36     vector <pair<ll,ll>> lista;
37     visto[n]=true;
38     set <arista,compare> s;
39     forn(i,v[n].size())
40     {
41         arista aux;
42         aux.n=n; aux.m=v[n][i].F; aux.peso=v[n][i].S;
43         s.insert(aux);
44     }
45     while(s.size()!=0)
46     {
47         auto it=s.begin();
48         if(visto[it->m]==true) {s.erase(it); continue;}
49         visto[it->m]=true;
50         ll ini=it->n, fin=it->m, p=it->peso;
51         lista.pb(mp(ini,fin));
52         forn(i,v[fin].size())
53         {
54             arista aux;
55             aux.n=fin; aux.m=v[fin][i].F; aux.peso=v[fin][i].S;
56             if(visto[aux.m]==false) s.insert(aux);
57         }
58     }
59     return lista;
60 }

```

5.23 Toposort

```

1 // Topological sort
2 // Complexity: O(n+m)
3 vi tsort(vector<vi> &g, int n){ // lexicographically smallest
4     // topological sort
5     vi r; priority_queue<ll> q;
6     vi d(2*n);
7     forn(i, n) for(ll j: g[i]) d[j]++;
8     while(!q.empty()){
9         ll x=-q.top(); q.pop(); r.pb(x);
10        for(ll j: g[x]){
11            d[j]--;
12            if(!d[j]) q.push(-j);
13        }

```

```

14     }
15     return r; // if not DAG it will have less than n elements
16 }

```

6 Math

6.1 Bell

```

1 // Los numeros de Bell son utiles para calcular de cuantas maneras
2 // podemos subdivir el conjunto utilizando todos los elementos
3
4 // O(n^2)
5 // Numeros de Bell
6 // Bell numbers = {1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, ...}
7 // bell[i][j] = bell[i-1][j-1] + bell[i][j-1]
8 // bell[i][0] = bell[i-1][i-1]
9
10 vector<int> bellNumber(int n) {
11     vector<vector<int>> bell(n+1,vector<int>(n+1));
12     vector<int> ans(n+1,0); // Numeros de bell desde 0 hasta n
13     bell[0][0] = 1;
14     forr(i, 1, n+1){
15         bell[i][0] = bell[i-1][i-1];
16         forr(j, 1, i+1) bell[i][j] = bell[i-1][j-1] + bell[i][j-1];
17     }
18     forn(i, n+1) ans[i] = bell[i][0];
19     return ans;
20 }

```

6.2 Berlekamp Massey

```

1 // Berlekamp-Massey algorithm
2 // Description: Finds the shortest linear recurrence relation that
3 // generates a given sequence.
4 // Time: O(N^2)
5 // Output: A vector of integers, where the first element is the constant
6 // term and the last element is the coefficient of x^(n-1).
7
8 vi BM(vi x){
9     vi ls,cur; ll lf,ld;
10    forn(i,SZ(x)){
11        ll t=0;
12        forn(j,SZ(cur)) t=(t+x[i-j-1]*(ll)cur[j])%MOD;
13        if((t-x[i])%MOD==0) continue;
14        if(!SZ(cur)){
15            cur.resize(i+1); lf=i; ld=(t-x[i])%MOD;
16            continue;
17        }
18    }
19 }

```

```

15     }
16     ll k=-(x[i]-t)*be(1d,MOD-2,MOD)%MOD;
17     vi c(i-lf-1); c.pb(k);
18     forr(j,0,SZ(ls))c.pb(-ls[j]*k%MOD);
19     if(SZ(c)<SZ(cur))c.resize(SZ(cur));
20     forr(j,0,SZ(cur))c[j]=(c[j]+cur[j])%MOD;
21     if(i-lf+SZ(ls)>=SZ(cur)) ls=cur,lf=i,ld=(t-x[i])%MOD;
22     cur=c;
23 }
24 forn(i,SZ(cur)) cur[i]=(cur[i]%MOD+MOD)%MOD;
25 return cur;
26 }

```

6.3 Catalan

```

1 // Catalan dp
2
3 vector<ll> catalan(int n) { //O(N^2)
4     vector<ll> c(n+1,0); c[0] = c[1] = 1;
5     forr(i,2,n+1) forn(j,i) c[i] += c[j] * c[i-j-1];
6     return c;
7 }
8
9 // Catalan formula
10 // binom(2n,n) / (n+1) --- (2n!) / [(n+1)! * n!]
11
12 ll catalan(ll n) {
13     return (((fact[2*n] * ifact[n+1])%MOD)*ifact[n])%MOD;
14 }

```

6.4 Coeficientes Binomiales

```

1 //~ Solucion DP
2
3 vector<vector<ll>> binomial(int n) { // O (N^2)
4     vector<vector<ll>> binom(n+1,vector<ll>(n+1));
5     binom[0][0] = binom[1][0] = binom[1][1] = 1;
6     forr(i,2,n+1){ forn(j,i+1) {
7         if(j==0 || j==i) binom[i][j]=1;
8         else binom[i][j] = binom[i-1][j-1] + binom[i-1][j];
9     }
10 }
11 return binom;
12 }

```

```

13
14 //~ =====
15
16 //~ Solucion O(N + Log N)
17
18 vector<ll> fact, ifact;
19
20 ll be(ll b, ll e, ll m) {
21     ll r=1; b%=m;
22     while(e){if(e&1LL)r=r*b%m;b=b*m/m;e/=2;}
23     return r;
24 }
25 ll inv_mod(ll x, ll m){ return be(x,m-2,m); }
26
27 void init_fact(int maxn){
28     fact.resize(maxn), ifact.resize(maxn);
29     fact[0] = 1; forr(i, 1, maxn) fact[i] = (fact[i-1]*i)%MOD;
30     ifact[maxn-1] = inv_mod(fact[maxn-1], MOD);
31     for(int i = maxn-2; i >= 0; i--) ifact[i] = (ifact[i+1]*(i+1))%MOD;
32 }
33
34 ll binom(ll n, ll k){
35     if(n < 0 or k < 0 or n < k) return 0;
36     ll ans = fact[n];
37     ans *= ifact[k]; ans %= MOD;
38     ans *= ifact[n-k]; ans %= MOD;
39     return ans;
40 }

```

6.5 Fft

```

1 typedef long long tipo;
2
3 struct CD{
4     double r,i;
5     CD(double r=0,double i=0):r(r),i(i){}
6     double real()const{return r;}
7     void operator /=(const tipo c){r/=c; i/=c;}
8 };
9 CD operator*(const CD& a, const CD& b){return CD(a.r*b.r-a.i*b.i,a.r*b.i
10 +a.i*b.r);}
11 CD operator+(const CD& a, const CD& b){return CD(a.r+b.r,a.i+b.i);}
12 CD operator-(const CD& a, const CD& b){return CD(a.r-b.r,a.i-b.i);}

```

```

12 const double pi=acos(-1);
13 const tipo MAXN=1<<21;
14 CD cp1[MAXN+9],cp2[MAXN+9];
15 tipo R[MAXN+9];
16
17 void dft(CD* a, tipo n, bool inv){
18     forn(i, n) if(R[i]<i) swap(a[R[i]],a[i]);
19     for(int m=2;m<=n;m*=2){
20         double z=2*pi/m*(inv?-1:1);
21         CD wi=CD(cos(z),sin(z));
22         for(int j=0;j<n;j+=m){
23             CD w(1);
24             for(int k=j,k2=j+m/2;k2<j+m;k++,k2++){
25                 CD u=a[k]; CD v=a[k2]*w;a[k]=u+v;a[k2]=u-v;w=w*wi;
26             }
27         }
28     }
29     if(inv) forn(i,n) a[i]/=n;
30 }
31 vector<tipo> multiply(vector<tipo> &p1, vector<tipo> &p2){
32     tipo n=SZ(p1)+SZ(p2)+1;
33     tipo m=1,cnt=0;
34     while(m<=n) m+=m,cnt++;
35     forn(i, m){R[i]=0;forr(j,0,cnt)R[i]=(R[i]<<1)|((i>>j)&1);}
36     forn(i, m) cp1[i]=cp2[i]=0;
37     forn(i, SZ(p1)) cp1[i]=p1[i];
38     forn(i, SZ(p2)) cp2[i]=p2[i];
39     dft(cp1,m,0); dft(cp2,m,0);
40     forn(i, m) cp1[i]=cp1[i]*cp2[i];
41     dft(cp1,m,1);
42     vector<tipo> res;
43     n -= 2;
44     forn(i,n) res.pb((ll)floor(cp1[i].real()+.5));
45     return res;
46 }

```

6.6 Fft Mod

```

1 /**
2  Implementacion iterativa de FFT and FFTmod. O(N log N)
3
4  La precision de la FFT depende de la respuesta.
5  x <= 5e14 (double), x <= 1e18(long double)

```

```

6 donde x = max(ans[i]) para la FFT, y x = N*mod for FFTmod
7 **/
8
9 typedef double ld;
10 typedef complex<ld> cd;
11
12 void fft(vector<cd> &a){
13     int n=SZ(a), L=31-__builtin_clz(n);
14     vi rev(n);
15     forn(i, n) rev[i]=(rev[i/2]+((i&1)<<L))/2;
16     forn(i, n) if(i<rev[i]) swap(a[i],a[rev[i]]);
17     static vector<cd> root(2,1);
18     for(static int k=2;k<n;k*=2){
19         root.resize(n);
20         cd z=polar(1.0,acos(-1)/k);
21         forr(i, k, 2*k) root[i]=root[i/2]*((i&1)?z:1);
22     }
23     for(int k=1;k<n;k*=2){
24         for(int i=0;i<n;i+=2*k){
25             forn(j, k){
26                 cd z=root[j+k]*a[i+j+k];
27                 a[i+j+k]=a[i+j]-z;
28                 a[i+j]+=z;
29             }
30         }
31     }
32 }
33
34 vi mul(vi &a, vi &b){
35     int sa=SZ(a), sb=SZ(b);
36     if(sa==0||sb==0) return {};
37     int n=1<<(32-__builtin_clz(sa+sb-2));
38     vector<cd> f(n),h(n);
39     forn(i, n) f[i]=cd(ld(i<sa?a[i]:0),ld(i<sb?b[i]:0));
40     fft(f);
41     forn(i, n){
42         ll j=(-i)&(n-1);
43         h[i]=(f[j]*f[j]-conj(f[i]*f[i]))*cd(0,-0.25/n);
44     }
45     fft(h);
46     vi c(sa+sb-1);
47     forn(i, sa+sb-1) c[i]=ll(round(h[i].real()));
48     return c;

```

```

49 }
50
51 vi mul_mod(vi &a, vi &b, ll mod){
52     int sa=SZ(a), sb=SZ(b);
53     if(sa==0||sb==0) return {};
54     vi a1(sa),a2(sa),b1(sb),b2(sb);
55     forn(i, sa) a1[i]=(a[i]&((1<<15)-1)),a2[i]=(a[i]>>15);
56     forn(i, sb) b1[i]=(b[i]&((1<<15)-1)),b2[i]=(b[i]>>15);
57     vi c1=mul(a1,b1),c2=mul(a1,b2),c3=mul(a2,b1),c4=mul(a2,b2);
58     vi c(sa+sb-1);
59     forn(i, sa+sb-1) c[i]=(c1[i]+((1<<15)*(c2[i]%mod+c3[i]%mod))
60         +((1<<30)*(c4[i]%mod))%mod);
61     return c;
62 }

```

6.7 Fht

```

1 ll c1[MAXN+9],c2[MAXN+9]; // MAXN must be power of 2 !!
2 void fht(ll* p, int n, bool inv){
3     for(int l=1;2*l<=n;l*=2)for(int i=0;i<n;i+=2*l)forn(j,l){
4         ll u=p[i+j],v=p[i+l+j];
5         if(!inv)p[i+j]=u+v,p[i+l+j]=u-v; // XOR
6         else p[i+j]=(u+v)/2,p[i+l+j]=(u-v)/2;
7         //if(!inv)p[i+j]=v,p[i+l+j]=u+v; // AND
8         //else p[i+j]=-u+v,p[i+l+j]=u;
9         //if(!inv)p[i+j]=u+v,p[i+l+j]=u; // OR
10        //else p[i+j]=v,p[i+l+j]=u-v;
11    }
12 }
13 // like polynomial multiplication, but XORing exponents
14 // instead of adding them (also ANDing, ORing)
15 vi multiply(vi &p1, vi &p2){
16     int n=1<<(32-__builtin_clz(max(SZ(p1),SZ(p2))-1));
17     forn(i,n)c1[i]=0,c2[i]=0;
18     forn(i,SZ(p1))c1[i]=p1[i];
19     forn(i,SZ(p2))c2[i]=p2[i];
20     fht(c1,n,false);fht(c2,n,false);
21     forn(i,n)c1[i]*=c2[i];
22     fht(c1,n,true);
23     return vi(c1,c1+n);
24 }

```

6.8 Fracciones

```

1 struct frac {
2     ll n, d;
3     frac(ll x, ll y) {
4         ll g = __gcd(x,y);
5         n = x/g; d = y/g;
6         if(d < 0) n *= -1LL, d *= -1LL;
7     }
8     bool const operator <(const frac &a) const {
9         return n * a.d < d * a.n; }
10    bool const operator ==(const frac &a) const {
11        return n * a.d == d * a.n; }
12    frac const operator +(const frac &a) const {
13        return frac(n*a.d + a.n*d, d*a.d); }
14    frac const operator -(const frac &a) const {
15        return frac(n*a.d - a.n*d, d*a.d); }
16    frac const operator *(const frac &a) const {
17        return frac(n * a.n, d * a.d); }
18    frac const operator /(const frac &a) const {
19        return frac(n * a.d, d * a.n); }
20 };

```

6.9 Gauss Sistema Ecuaciones

```

1 //~ const int INF = 2; // it doesn't actually have to be infinity or a
2   big number
3
4 /*
5 Sistema de ecuaciones lineales modulares. Resuelve Ax=B.
6 En las filas estan las ecuaciones y en las columnas las variables.
7 En ans queda la solucion en caso de que sea unica. La matriz A es de n x
8   m.
9 Esta matriz tiene n ecuaciones y m-1 variables, la ultima columna es la
10  matrix b.
11 Complejidad O(min(n, m)nm). Si n == m, es O(n^3).
12 */
13 ll mod(ll x) {x%=MOD; if(x<0) x+=MOD; return x;}
14 ll add(ll a, ll b){return mod(a+b);}
15 ll sub(ll a, ll b){return mod(a-b);}
16 ll mul(ll a, ll b){return mod(mod(a)*mod(b));}
17 ll be(ll a, ll p) {
18     ll ans=1; for(; p; p/=2, a = mul(a,a)) if(p&1) ans=mul(ans,a);
19     return ans;
20 }

```

```

18 }
19 ll divi(ll a, ll b){return mul(a,be(b,MOD-2));}
20
21 int gauss(vector<vi> &a, vi &ans) {
22     int n = SZ(a), m = SZ(a[0])-1;
23
24     vector<int> where(m, -1);
25     for(int col=0, row=0; col<m && row<n; ++col) {
26         int sel = row;
27         forr(i, row, n){ // if(abs(a[i][col]) > abs(a[sel][col]))
28             if(a[i][col] > a[sel][col]) sel = i;
29         }
30         if(a[sel][col] == 0) continue; // if(abs(a[sel][col]) < EPS)
31         forr(i, col, m+1) swap(a[sel][i], a[row][i]);
32         where[col] = row;
33
34         forn(i, n){
35             if(i != row) {
36                 ll c = divi(a[i][col], a[row][col]); // double c = a[i][
37                     col / a[row][col];
38                 forr(j, col, m+1) a[i][j] = sub(a[i][j], mul(a[row][j],
39                     c));
40             }
41         }
42         ++row;
43     }
44
45     ans.assign(m, 0);
46     forn(i, m) if(where[i] != -1) ans[i] = divi(a[where[i]][m], a[where[
47         i]][i]);
48     forn(i, n){
49         ll sum = 0; // double sum = 0;
50         forn(j, m) sum = add(sum, mul(ans[j], a[i][j]));
51         if(sum - a[i][m] != 0) return 0; // if(abs(sum-a[i][m]) > EPS)
52         // No hay soluciones
53     }
54     forn(i, m) if(where[i] == -1) return INF; // Infinitas soluciones
55     return 1; // Exactamente una solucion
56 }
57
58 // Encuentra las bases que generan la solucion
59 vector<vi> find_bases(vector<vi> &A, vi &ans) {
60     int n = SZ(A), m = SZ(ans);

```

```

58 vi var_ind;
59 forn(i,SZ(ans)) if(ans[i] == 0) var_ind.pb(i);
60
61 vector <vi> bases(SZ(var_ind),vi(m,0));
62 int r = 0;
63 for(auto u : var_ind) {
64     bases[r][u] = 1;
65     forn(i,n) {
66         forn(j,m) {
67             if(A[i][j] != 0) {
68                 bases[r][j] = divi(-A[i][u], A[i][j]);
69                 break;
70             }
71         }
72     }
73     r++;
74 }
75 return bases;
76 }

```

6.10 Inverse Matrix

```

1 // modular functions in gauss
2
3 pair<ll,vector<vi>> inverse_matrix(vector<vector<ll> > &x) {
4     // returns det and the inverse of the matrix,
5     // if det == 0, no inverse
6     int n = SZ(x); vector <vi> I(n,vi(n,0));
7     ll det = 1;
8     forn(i,n) I[i][i] = 1;
9     forn(i,n) {
10         int pos = -1;
11         forr(fila,i,n) if(x[fil][i] > 0) pos = fila; // Not null pivot
12         if(pos == -1) {return {0,I};} // If pos = -1, singular matrix
13         if(pos != i) {
14             swap(x[i],x[pos]), swap(I[i],I[pos]);
15             det = mul(det,-1);
16         }
17         ll pivot = x[i][i];
18         det = mul(det,pivot);
19         forn(j,n) {
20             x[i][j] = divi(x[i][j],pivot);
21             I[i][j] = divi(I[i][j],pivot);

```



```

22     }
23     forn(filas,n) {
24         if(filas == i) continue;
25         pivot = x[filas][i];
26         forn(j,n) {
27             x[filas][j] = sub(x[filas][j],mul(pivot,x[i][j]));
28             I[filas][j] = sub(I[filas][j],mul(pivot,I[i][j]));
29         }
30     }
31 }
32 return {det,I};
33 }

```

6.11 Iterate Submask

```

1 // 3 ^ n iteration
2 forn(mask, (1<<n)){
3     int sm = 0;
4     do {
5         // do something with submask
6     } while (sm=(sm-mask)&mask);
7 }

```

6.12 Lagrange

```

1 ll lagrange(vii a, ll x){
2     /* Recibe una muestra a de n puntos (xi,yi)
3     y evalua el polinomio de Lagrange de
4     en el punto x desconocido */
5
6     ll ans = 0;
7     int n = SZ(a);
8     forn(i, n){
9         ll term = a[i].second;
10        forn(j, n){
11            if(j != i)
12                term = term*(x-a[j].first)/(a[i].first-a[j].first);
13        }
14        ans += term;
15    }
16    return ans;
17 }

```

6.13 Linear Rec

```

1 const ll LOG = 60;
2 // Linear Recurrence
3 // Description: Calculates the kth term of a linear recurrence relation
4 // init 0(n^2log) query(n^2 logk)
5 // input: terms: first n term; trans: transition function (calcular con
6 //         BM); MOD; LOG=mxlog of k
7 // output calc(k): kth term mod MOD
8 // example: {1,1} {2,1} an=2*a_(n-1)+a_(n-2); calc(3)=3 calc(10007)
9 //           =71480733
10 struct LinearRec {
11     ll n; vi terms, trans; vector<vi> bin;
12     vi add(vi &a, vi &b){
13         vi res(n*2+1);
14         forn(i, n+1) forn(j, n+1) res[i+j]=(res[i+j]*1LL+(ll)a[i]*b[j])%
15             MOD;
16         for(ll i=2*n; i>n; --i){
17             forn(j, n) res[i-1-j]=(res[i-1-j]*1LL+(ll)res[i]*trans[j])%
18                 MOD;
19             res[i]=0;
20         }
21         res.erase(res.begin()+n+1,res.end());
22         return res;
23     }
24     LinearRec(vi &terms, vi &trans):terms(terms),trans(trans){
25         n=SZ(trans);vi a(n+1);a[1]=1;
26         bin.pb(a);
27         forr(i,1,LOG)bin.pb(add(bin[i-1],bin[i-1])); // Precompute
28         powers of a
29     }
30     ll calc(ll k){
31         vi a(n+1);a[0]=1;
32         forn(i,LOG) if((k>>i)&1)a=add(a,bin[i]);
33         ll ret=0;
34         forn(i,n) ret=(ret+a[i+1]*terms[i])%MOD;
35         return ret;
36     }
37 };

```

6.14 Matrix Fast Pow

```

1 typedef vector<vector<ll>> Matrix;
2
3 Matrix ones(int n) {

```

```

4   Matrix r(n,vector<ll>(n));
5   forn(i,n)r[i][i]=1;
6   return r;
7 }
8 Matrix operator*(Matrix &a, Matrix &b) {
9     int n=a.size(),m=b[0].size(),z=a[0].size();
10    Matrix r(n,vector<ll>(m));
11    forn(i,n)forn(j,m)forn(k,z)
12        r[i][j]+=a[i][k]*b[k][j],r[i][j]%=MOD;
13    return r;
14 }
15 Matrix be(Matrix b, ll e) {
16    Matrix r=ones(b.size());
17    while(e){if(e&1LL)r=r*b;b=b*b;e/=2;}
18    return r;
19 }

```

6.15 Matrix Reduce

```

1 double reduce(vector<vector<double> > &x) { // returns determinant
2     int n=x.size(),m=x[0].size();
3     int i=0,j=0;double r=1.;
4     while(i<n&&j<m){
5         int l=i;
6         forr(k,i+1,n)if(abs(x[k][j])>abs(x[l][j]))l=k;
7         if(abs(x[l][j])<EPS){j++;r=0.;continue;}
8         if(l!=i){r=-r;swap(x[i],x[l]);}
9         r*=x[i][j];
10        for(int k=m-1;k>j;k--)x[i][k]/=x[i][j];
11        forr(k,0,n){
12            if(k==i)continue;
13            for(int l=m-1;l>j;l--)x[k][l]-=x[k][j]*x[i][l];
14        }
15        i++;j++;
16    }
17    return r;
18 }

```

6.16 Ntt

```

1 struct ntt {
2     const int mod = 998244353;
3     const int root = 15311432;
4     const int root_1 = 469870224;

```

```

5     const int root_pw = 1 << 23;
6
7     void fft(vector<int> &a, bool invert) {
8         int n = SZ(a);
9
10        for(int i = 1, j = 0; i < n; i++) {
11            int bit = n >> 1;
12            for (; j&bit; bit >>= 1) j ^= bit;
13            j ^= bit;
14
15            if (i < j) swap(a[i], a[j]);
16        }
17
18        for(int len = 2; len <= n; len <= 1) {
19            int wlen = invert ? root_1 : root;
20            for (int i = len; i < root_pw; i <= 1)
21                wlen = (int)(1LL * wlen * wlen % mod);
22
23            for(int i = 0; i < n; i += len) {
24                int w = 1;
25                for (int j = 0; j < len / 2; j++) {
26                    int u = a[i+j], v = (int)(1LL * a[i+j+len/2] * w %
27                        mod);
28                    a[i+j] = u + v < mod ? u + v : u + v - mod;
29                    a[i+j+len/2] = u - v >= 0 ? u - v : u - v + mod;
30                    w = (int)(1LL * w * wlen % mod);
31                }
32            }
33
34            if (invert) {
35                int n_1 = (int) inv_mod(n, mod);
36                for (int &x : a) x = (int)(1LL * x * n_1 % mod);
37            }
38        }
39
40        vector<int> multiply(vector<int> const &a, vector<int> const &b){
41            vector<int> fa(all(a)), fb(all(b));
42            int n = 1;
43            while(n < SZ(a)+SZ(b)) n <= 1;
44            fa.resize(n), fb.resize(n);
45            fft(fa, false); fft(fb, false); // transformo
46            forn(i, n) fa[i] = (int) (1LL*fa[i]*fb[i]%mod);

```

```

47     fft(fa, true); // anti-transformo
48
49     vector<int> result(n);
50     forn(i, n) result[i] = fa[i];
51     return result;
52 }
53 };

```

6.17 Polynomial

```

1  typedef int tp; // type of polynomial
2  template<class T=tp>
3  struct poly { // poly<> : 1 variable, poly<poly<>>: 2 variables, etc.
4      vector<T> c;
5      T& operator[](int k){return c[k];}
6      poly(vector<T>& coefs):c(coefs){}
7      poly(initializer_list<T> coefs):c(coefs){}
8      poly(int k):c(k){}
9      poly(){}
10     poly operator+(poly<T> o){
11         int m=c.size(),n=o.c.size();
12         poly res(max(m,n));
13         forn(i,m)res[i]=res[i]+c[i];
14         forn(i,n)res[i]=res[i]+o.c[i];
15         return res;
16     }
17     poly operator*(tp k){
18         poly res(c.size());
19         forn(i,c.size())res[i]=c[i]*k;
20         return res;
21     }
22     poly operator*(poly o){
23         int m=c.size(),n=o.c.size();
24         poly res(m+n-1);
25         forn(i,m)for(n,j,n)res[i+j]=res[i+j]+c[i]*o.c[j];
26         return res;
27     }
28     poly operator-(poly<T> o){return *this+(o*-1);}
29     T operator()(tp v){
30         T sum(0);
31         for(int i=c.size()-1;i>=0;--i)sum=sum*v+c[i];
32         return sum;
33     }

```

```

34 };
35
36 // example: p(x,y)=2*x^2+3*x*y-y+4
37 // poly<poly<>> p={{4,-1},{0,3},{2}}
38 // printf("%d\n",p(2)(3)) // 27 (p(2,3))
39 set<tp> roots(poly<> p){ // only for integer polynomials
40     set<tp> r;
41     while(!p.c.empty()&&!p.c.back())p.c.pop_back();
42     if(!p(0))r.insert(0);
43     if(p.c.empty())return r;
44     tp a0=0,an=abs(p[p.c.size()-1]);
45     for(int k=0;!a0;a0=abs(p[k++]));
46     vector<tp> ps,qs;
47     forr(i,1,sqrt(a0)+1)if(a0%i==0)ps.pb(i),ps.pb(a0/i);
48     forr(i,1,sqrt(an)+1)if(an%i==0)qs.pb(i),qs.pb(an/i);
49     for(auto pt:ps)for(auto qt:qs)if(pt%qt==0){
50         tp x=pt/qt;
51         if(!p(x))r.insert(x);
52         if(!p(-x))r.insert(-x);
53     }
54     return r;
55 }
56 pair<poly<>,tp> ruffini(poly<> p, tp r){ // returns pair (result,rem)
57     int n=p.c.size()-1;
58     vector<tp> b(n);
59     b[n-1]=p[n];
60     for(int k=n-2;k>=0;--k)b[k]=p[k+1]+r*b[k+1];
61     return mp(poly<>(b),p[0]+r*b[0]);
62 }
63 // only for double polynomials
64 pair<poly<>,poly<>> polydiv(poly<> p, poly<> q){ // returns pair (
65     result,rem)
66     int n=p.c.size()-q.c.size()+1;
67     vector<tp> b(n);
68     for(int k=n-1;k>=0;--k){
69         b[k]=p.c.back()/q.c.back();
70         forn(i,q.c.size())p[i+k]-=b[k]*q[i];
71         p.c.pop_back();
72     }
73     while(!p.c.empty()&&abs(p.c.back())<EPS)p.c.pop_back();
74     return mp(poly<>(b),p);
75 }
// only for double polynomials

```

```

76 poly<> interpolate(vector<tp> x, vector<tp> y){ //TODO TEST
77     poly<> q={1},S={0};
78     for(tp a:x)q=poly<>({-a,1})*q;
79     forn(i,x.size()){
80         poly<> Li=ruffini(q,x[i]).fst;
81         Li=Li*(1.0/Li(x[i])); // change for int polynomials
82         S=S+Li*y[i];
83     }
84     return S;
85 }

```

6.18 Print Int128

```

1 string toString(__int128 num) {
2     string str;
3     do {
4         int digit = num % 10;
5         str = to_string(digit) + str;
6         num = (num - digit) / 10;
7     } while (num != 0);
8     return str;
9 }
10
11 ostream& operator<<(std::ostream& os, __int128 t) {
12     string str = toString(t);
13     return os << str;
14 }

```

6.19 Random

```

1 mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
2 int my_random(int l, int r){ // Random number in [l, r]
3     return uniform_int_distribution<int>(l, r)(rng);
4 }

```

6.20 Shortest Path Matrix Exp

```

1 #define INF 2e18
2
3 typedef vector<vector<ll> > Matrix;
4
5 ll funcion(ll x, ll y){
6     if(x!=INF && y!=INF) return (x+y);
7     else return INF;

```

```

8 }
9
10 Matrix operator*(Matrix &a, Matrix &b) {
11     int n=a.size(),m=b[0].size(),z=a[0].size();
12     Matrix r(n,vector<ll>(m,INF));
13     forn(i,n) forn(j,m) forn(k,z)
14         r[i][j]=min(r[i][j],funcion(a[i][k],b[k][j]));
15     return r;
16 }
17
18 Matrix be(Matrix b, ll e) {
19     Matrix r(b.size(),vector<ll>(b[0].size())); e--;
20     forn(i,b.size()){ forn(j,b[0].size()) r[i][j]=b[i][j]; }
21     while(e){if(e&1LL)r=r*b;b=b*b;e/=2;}
22     return r;
23 }

```

6.21 Simplex

```

1 struct Simplex {
2     vector<int> X,Y;
3     vector<vector<double>> A;
4     vector<double> b,c;
5     double z;
6     int n,m;
7     void add_equation(vector<double> v, double val){ A.pb(v); b.pb(val);
8     }
9     void set_objective(vector<double> _c){ c = _c; }
10    void pivot(int x,int y){
11        swap(X[y],Y[x]);
12        b[x]/=A[x][y];
13        forr(i,0,m)if(i!=y)A[x][i]/=A[x][y];
14        A[x][y]=1/A[x][y];
15        forr(i,0,n)if(i!=x&&abs(A[i][y])>EPS){
16            b[i]-=A[i][y]*b[x];
17            forr(j,0,m)if(j!=y)A[i][j]-=A[i][y]*A[x][j];
18            A[i][y]=-A[i][y]*A[x][y];
19        }
20        z+=c[y]*b[x];
21        forr(i,0,m)if(i!=y)c[i]-=c[y]*A[x][i];
22        c[y]=-c[y]*A[x][y];
23    }
24    pair<double,vector<double>> simplex() {

```

```

24 // maximiza c^T x, dado Ax<=b, x>=0
25 // retorna par (maximum value, solution vector)
26 n=b.size();m=c.size();z=0.;
27 X=vector<int>(m);Y=vector<int>(n);
28 forn(i,m)X[i]=i;
29 forn(i,n)Y[i]=i+m;
30 while(1){
31     int x=-1,y=-1;
32     double mn=-EPS;
33     forn(i,n)if(b[i]<mn)mn=b[i],x=i;
34     if(x<0) break;
35     forr(i,m)if(A[x][i]<-EPS){y=i;break;}
36     assert(y>0 && "No_solution_found"); // no solution to Ax<=b
37     pivot(x,y);
38 }
39 while(1){
40     double mx=EPS;
41     int x=-1,y=-1;
42     forn(i,m)if(c[i]>mx)mx=c[i],y=i;
43     if(y<0)break;
44     double mn=1e200;
45     forn(i,n)if(A[i][y]>EPS&&b[i]/A[i][y]<mn)mn=b[i]/A[i][y],x=i;
46     ;
47     assert(x>0 && "c^T x is unbounded"); // c^T x is unbounded
48     pivot(x,y);
49 }
50 vector<double> r(m);
51 forn(i,n)if(Y[i]<m)r[Y[i]]=b[i];
52 return mp(z,r);
53 };

```

6.22 Simpson

```

1 double integrate(double f(double), double a, double b, int n=10000){
2     double r=0,h=(b-a)/n,fa=f(a),fb;
3     forr(i,0,n) {fb=f(a+h*(i+1));r+=fa+4*f(a+h*(i+0.5))+fb;fa=fb;}
4     return r*h/6.;
5 }

```

7 Number Theory

7.1 Binexp Invmod

```

1 ll be(ll x, ll y, ll m) {
2     if (y == 0) return 1;
3     ll p = be(x, y/2, m) % m;
4     p = (p * p) % m;
5     return (y%2 == 0)? p : (x * p) % m;
6 }
7
8 ll be_it(ll b, ll e, ll m) {
9     ll r=1; b%=m;
10    while(e){if(e&1LL)r=r*b%m;b=b*b%m;e/=2;}
11    return r;
12 }
13
14 ll inv_mod(ll x, ll m) {return be(x,m-2,m);}

```

7.2 Criba Phi Euler

```

1 bool is_composite[1000];
2 int phi[1000];
3 vector<long long> prime;
4
5 void sieve (int n) {
6     fill(is_composite, is_composite + n, false);
7     phi[1] = 1;
8     for (int i = 2; i < n; ++i) {
9         if (!is_composite[i]) {
10             prime.push_back (i);
11             phi[i] = i - 1; //i is prime
12         }
13         for (int j = 0; j < prime.size () && i * prime[j] < n; ++j) {
14             is_composite[i * prime[j]] = true;
15             if (i % prime[j] == 0) {
16                 phi[i * prime[j]] = phi[i] * prime[j]; //prime[j]
17                                     divides i
18                 break;
19             } else {
20                 phi[i * prime[j]] = phi[i] * phi[prime[j]]; //prime[j]
21                                     does not divide i
22             }
23         }
24     }
25 }

```

```

21     }
22 }
23 }

```

7.3 Criba Primos

```

1 // Criba lineal, obtiene los primos menores al parametro
2 vi min_prime; // min_prime[i] contiene el menor primo que divide a i,
  // util para factorizar en log(i)
3
4 vi criba(ll n) {
5     vb prime(n+1,true);
6     min_prime.resize(n+1,INF);
7     vi primos;
8     for(ll p=2; p*p<=n; p++){
9         if(!prime[p]) continue;
10        for(ll i=p*p; i<=n; i += p) {
11            prime[i] = false;
12            min_prime[i] = min(min_prime[i],p);
13        }
14    }
15    forr(i, 2, n+1){
16        if(prime[i]) primos.pb(i), min_prime[i] = i;
17    }
18    return primos; // lista de primos hasta n
19 }

```

7.4 Diofantica

```

1 // Algoritmo de euclides extendido para encontrar gcd de a y b.
2 // En x e y se guarda una solucion particular de la ecuacion ax+by=gcd(a
  // ,b)
3 ll gcd(ll a, ll b, ll &x, ll &y) {
4     if (b == 0) {
5         x = 1; y = 0;
6         return a;
7     }
8     ll x1, y1;
9     ll d = gcd(b, a % b, x1, y1);
10    x = y1;
11    y = x1 - y1 * (a / b);
12    return d;
13 }
14

```

```

15 // Soluciones a la ecuacion ax+by=c, retorna true si hay solucion
16 // x0 e y0 son soluciones particulares de la ecuacion.
17 // g es el gcd(a,b), hay solucion si y solo si divide a c.
18 bool find_any_solution(ll a, ll b, ll c, ll &x0, ll &y0, ll &g) {
19     g = gcd(abs(a), abs(b), x0, y0);
20     if (c % g) return false;
21
22     x0 *= c / g;
23     y0 *= c / g;
24     if (a < 0) x0 = -x0;
25     if (b < 0) y0 = -y0;
26     return true;
27 }
28 // Todas las soluciones de ax+by=c son de la forma: (x, y) = (x0+k*b/g,
  // y0-k*a/g), para todo k

```

7.5 Discrete Log

```

1 /* Algoritmo de baby-step giant-step para
2  * hallar x tal que a^x = b (mod m)
3  * en O(sqrt(m))
4  */
5 ll discrete_log(ll a, ll b, ll m) {
6     a %= m, b %= m;
7     ll n = sqrt(m) + 1;
8
9     ll an = 1;
10    for (ll i = 0; i < n; ++i)
11        an = (an * 11l * a) % m;
12
13    unordered_map<ll, ll> vals;
14    for (ll q = 0, cur = b; q <= n; ++q) {
15        vals[cur] = q;
16        cur = (cur * 11l * a) % m;
17    }
18
19    for (ll p = 1, cur = 1; p <= n; ++p) {
20        cur = (cur * 11l * an) % m;
21        if (vals.count(cur)) {
22            ll ans = n * p - vals[cur];
23            return ans;
24        }
25    }
26 }

```

```

26     return -1;
27 }

```

7.6 Discrete Root

```

1 // Finds the primitive root modulo p
2 ll generator(ll p) {
3     vi fact;
4     ll phi = p-1, n = phi;
5     for (ll i = 2; i * i <= n; ++i) {
6         if (n % i == 0) {
7             fact.push_back(i);
8             while (n % i == 0)
9                 n /= i;
10        }
11    }
12    if (n > 1)
13        fact.push_back(n);
14
15    forr(res, 2, p+1){
16        bool ok = true;
17        for (ll factor : fact) {
18            if (be(res, phi / factor, p) == 1) {
19                ok = false;
20                break;
21            }
22        }
23        if (ok) return res;
24    }
25    return -1;
26 }
27
28 // This program finds all numbers x such that x^k = a (mod n)
29 vi discrete_root(ll n, ll k, ll a){
30     if (a == 0) return {0};
31
32     ll g = generator(n);
33
34     // Baby-step giant-step discrete logarithm algorithm
35     ll sq = (ll) sqrt (n + .0) + 1;
36     vector<pair<ll, ll>> dec(sq);
37     forr(i, 1, sq+1)
38         dec[i-1] = {be(g, i * sq * k % (n - 1), n), i};

```

```

39     sort(all(dec));
40     ll any_ans = -1;
41     forn(i, sq){
42         ll my = be(g, i * k % (n - 1), n) * a % n;
43         auto it = lower_bound(all(dec), mp(my, 0));
44         if (it != dec.end() && it->first == my) {
45             any_ans = it->second * sq - i;
46             break;
47         }
48     }
49     if (any_ans == -1) return {};
50
51     // Print all possible answers
52     ll delta = (n-1) / __gcd(k, n-1);
53     vi ans;
54     for (ll cur = any_ans % delta; cur < n-1; cur += delta)
55         ans.push_back(be(g, cur, n));
56     sort(ans.begin(), ans.end());
57     return ans;
58 }

```

7.7 Divisors

```

1 vi find_divisors(ll n, vi &primos) {
2     vector <pair<ll,ll>> factor;
3     for(ll prime : primos) {
4         int cont = 0;
5         while(n % prime == 0) {
6             cont++;
7             n /= prime;
8         }
9         if(cont > 0) factor.pb({prime,cont});
10    }
11
12    if(n > 1) factor.pb({n,1});
13
14    vi divisores = {1};
15    for(auto [p,exp] : factor) {
16        int tam = SZ(divisores);
17        forn(i,exp) {
18            forn(j,tam) {
19                int pos = SZ(divisores)-tam;
20                divisores.pb(divisores[pos] * p);

```

```

21     }
22   }
23 }
24
25 sort(all(divisores));
26 return divisores;
27 }

```

7.8 Fast Factorization Mrabin Prho

```

1 typedef long double ld;
2 mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().count());
3 ll myrand(ll a, ll b) { return uniform_int_distribution<ll>(a, b)(rng);
4   }
5 struct Factorization {
6   inline ll mul(ll a, ll b, ll c){
7     ll s = a * b - c * ll((ld) a / c * b + 0.5);
8     return s < 0 ? s + c : s;
9   }
10
11   ll be(ll a, ll k, ll mod) {
12     ll res = 1;
13     for(; k >= 1, a = mul(a, a, mod)) if (k & 1) res = mul(res,
14       a, mod);
15     return res;
16   }
17   bool miller(ll n) { // chequea si un numero es primo
18     auto test = [&](ll n, ll a) {
19       if (n == a) return true;
20       if (n % 2 == 0) return false;
21
22       ll d = (n - 1) >> __builtin_ctzll(n - 1);
23       ll r = be(a, d, n);
24       while (d < n - 1 && r != 1 && r != n - 1) {
25         d <<= 1;
26         r = mul(r, r, n);
27       }
28       return r == n - 1 || d & 1;
29     };
30
31     if(n == 2) return 1;

```

```

32   for(auto p: vi{2, 3, 5, 7, 11, 13}) if(!test(n, p)) return 0;
33   return 1;
34 }
35
36 ll pollard(ll n) { // devuelve un factor no trivial de n
37   auto f = [&](ll x) { return ll((__int128(x) * x + 1) % n); };
38
39   ll x = 0, y = 0, t = 30, prd = 2;
40   while (t++ % 40 || __gcd(prd, n) == 1) { // speedup: no tomar
41     gcd en cada iteracion
42     if (x == y) x = myrand(2, n - 1), y = f(x);
43     ll tmp = mul(prd, abs(x - y), n);
44     if (tmp) prd = tmp;
45     x = f(x), y = f(f(y));
46   }
47   return __gcd(prd, n);
48 }
49
50 vi factorize(ll n) {
51   vi res;
52   auto dfs = [&](auto &dfs, ll x) {
53     if (x == 1) return;
54     if (miller(x)) res.pb(x);
55     else {
56       ll d = pollard(x);
57       dfs(dfs, d); dfs(dfs, x / d);
58     }
59   };
60   dfs(dfs, n); sort(all(res));
61   return res;
62 };

```

7.9 Modint

```

1 // Template de benq para enteros que tienen que ser tomados MOD
2 // se define un tipo mi como "typedef modular<ll> mi"
3 template<class T> struct modular {
4   T val;
5   explicit operator T() const { return val; }
6   modular() { val = 0; }
7   template<class U> modular(const U& v) {
8     val = (-MOD <= v && v <= MOD) ? v : v % MOD;

```



```

9         if (val < 0) val += MOD;
10    }
11    friend ostream& operator<<(ostream& os, const modular& a) { return
        os << a.val; }
12    friend bool operator==(const modular& a, const modular& b) { return
        a.val == b.val; }
13    friend bool operator!=(const modular& a, const modular& b) { return
        !(a == b); }

14
15    modular operator-() const { return modular(-val); }
16    modular& operator+=(const modular& m) { if ((val += m.val) >= MOD)
        val -= MOD; return *this; }
17    modular& operator-=(const modular& m) { if ((val -= m.val) < 0) val
        += MOD; return *this; }
18    modular& operator*=(const modular& m) { val = (ll)val*m.val%MOD;
        return *this; }
19    friend modular be(modular a, ll p) {
20        modular ans = 1; for (; p; p /= 2, a *= a) if (p&1) ans *= a;
21        return ans;
22    }
23    friend modular inv(const modular& a) { return be(a,MOD-2); }
24    // inv is equivalent to return be(b,b.mod-2) if prime
25    modular& operator/=(const modular& m) { return (*this) *= inv(m); }
26
27    friend modular operator+(modular a, const modular& b) { return a +=
        b; }
28    friend modular operator-(modular a, const modular& b) { return a -=
        b; }
29    friend modular operator*(modular a, const modular& b) { return a *=
        b; }
30
31    friend modular operator/(modular a, const modular& b) { return a /=
        b; }
32 };
33
34 typedef modular<ll> mi;
35 typedef vector<mi> vmi;

```

7.10 Mod Ops

```

1 ll mod(ll x) {x%=MOD; if(x<0) x+=MOD; return x;}
2 ll add(ll a, ll b){return mod(a+b);}
3 ll sub(ll a, ll b){return mod(a-b);}

```

```

4 ll mul(ll a, ll b){return mod(mod(a)*mod(b));}
5 ll be(ll a, ll p) {
6     ll ans=1; for(; p; p/=2, a = mul(a,a)) if(p&1) ans=mul(ans,a);
7     return ans;
8 }
9 ll divi(ll a, ll b){return mul(a,be(b,MOD-2));}

```

7.11 Moebius

```

1 vector<int> prime;
2 bool is_composite[MAXN];
3 int moebius[MAXN];
4
5 void sieve_moebius(int n) {
6     fill(is_composite, is_composite + n, false);
7     moebius[1] = 1;
8     forr(i, 2, n){
9         if(!is_composite[i]){
10             prime.pb(i); moebius[i] = - 1; // i is prime
11         }
12         for(int j = 0; j < SZ(prime) && i * prime[j] < n; ++j) {
13             is_composite[i * prime[j]] = true;
14             if (i % prime[j] == 0) {
15                 moebius[i * prime[j]] = 0; // prime[j] divides i
16                 break;
17             } else {
18                 moebius[i * prime[j]] = moebius[i] * moebius[prime[j]];
19                 // prime[j] does not divide i
20             }
21         }
22     }
23 }

```

7.12 Teo Chino Resto

```

1 // Teorema chino del resto
2 // Resuelve el sistema de ecuaciones modulares
3 // ai*xi = bi (mi)
4 ii extendedEuclid(ll a, ll b){ //a * x + b * y = gcd(a,b)
5     ll x,y;
6     if (b==0) return {1, 0};
7     auto p = extendedEuclid(b, a%b);
8     x = p.second, y = p.first-(a/b)*x;
9     if(a*x+b*y == -__gcd(a,b)) x=-x, y=-y;

```

```

10     return {x, y};
11 }
12 pair<ii, ii> dioph(ll a, ll b, ll r) {
13     // a*x+b*y=r donde r es multiplo de gcd(a,b);
14     ll d=__gcd(a,b);
15     a/=d; b/=d; r/=d;
16     auto p = extendedEuclid(a,b);
17     p.first*=r; p.second*=r;
18     assert(a*p.first+b*p.second==r);
19     return {p, {-b,a}}; // solutions: p+t*ans.second
20 }
21 ll inv(ll a, ll mod) { // inverso de a modulo mod
22     assert(__gcd(a,mod)==1);
23     ii sol = extendedEuclid(a,mod);
24     return ((sol.first%mod)+mod)%mod;
25 }
26
27 #define mod(a,m) (((a)%m+m)%m)
28 ii sol(tuple<ll,ll,ll> c){ //requires inv, diophantine
29     auto [a, x1, m] = c; ll d = __gcd(a,m);
30     if(d==1) return {mod(x1*inv(a,m),m), m};
31
32     return x1%d ? mp(-1LL,-1LL) : sol({a/d, x1/d, m/d});
33 }
34
35 // cond[i] = {ai,bi,mi} ai*xi=bi (mi); assumes lcm fits in ll
36 // Mucho cuidado con el overflow, usar __int128 si lcm no entra en ll
37 pair<ll,ll> crt(vector<tuple<ll,ll,ll>> cond) { // returns: (sol, lcm)
38     ll x1=0,m1=1,x2,m2;
39     for(auto t:cond){
40         tie(x2,m2) = sol(t);
41         if((x1-x2)%__gcd(m1,m2)) return {-1,-1};
42         if(m1==m2) continue;
43         ll k=dioph(m2,-m1,x1-x2).first.second, l=m1*(m2/__gcd(m1,m2));
44         x1 = mod((__int128)m1*k+x1,l); m1=l;
45     }
46     return sol({1,x1,m1});
47 }

```

8 Strings

8.1 Aho Corasick

```

1 struct vertex {
2     map<char,int> next, go;
3     int p,link,nextleaf;
4     char pch;
5     vector<int> leaf;
6     vertex(int p=-1, char pch=-1,int nextleaf=-1):
7         p(p),pch(pch),link(-1),nextleaf(nextleaf){}
8 };
9
10 vector<vertex> t;
11 vector <vector <int> > g(MAXN); // Suffix-links tree
12
13 void aho_init(){ //do not forget!!
14     t.clear(); t.pb(vertex());
15 }
16
17 void add_string(string s, int id) {
18     int v=0;
19     for(char c:s) {
20         if(!t[v].next.count(c)){
21             t[v].next[c]=t.size();
22             t.pb(vertex(v,c));
23         }
24         v=t[v].next[c];
25     }
26     t[v].leaf.pb(id);
27 }
28
29 int go(int v, char c);
30
31 int get_link(int v) {
32     if(t[v].link < 0) {
33         if(v == 0 || t[v].p == 0) t[v].link = 0;
34         else t[v].link = go(get_link(t[v].p),t[v].pch);
35         g[t[v].link].pb(v);
36     }
37     return t[v].link;
38 }
39

```

```

40 int go(int v, char c){
41     if(!t[v].go.count(c)) {
42         if(t[v].next.count(c)) t[v].go[c]=t[v].next[c];
43         else t[v].go[c] = v == 0 ? 0 : go(get_link(v),c);
44     }
45     return t[v].go[c];
46 }
47
48 int init_next_leaf(int v) {
49     if(v == 0) t[v].nextleaf=0;
50     if(t[get_link(v)].leaf.size()) return t[v].nextleaf=get_link(v);
51
52     else return t[v].nextleaf = t[get_link(v)].nextleaf != -1 ?
53         t[get_link(v)].nextleaf :
54         init_next_leaf(get_link(v));
55 }
56
57 void construct_links() { forn(i,t.size()) get_link(i); }

```

8.2 Dynamic Hash

```

1 ll be(ll x, ll y, ll m) {
2     if (y == 0) return 1;
3     ll p = be(x, y/2, m) % m;
4     p = (p * p) % m;
5     return (y%2 == 0)? p : (x * p) % m;
6 }
7
8 ll inv_mod(ll x, ll m) {return be(x,m-2,m);}
9
10 struct BIT {
11     vector<ll> prefix, a; ll M;
12     BIT() {}
13     BIT(vector<ll> &v, ll MOD) {
14         int n = v.size(); prefix.resize(n+1); a = v; M = MOD;
15         vector<ll> aux(n+1,0);
16         forn(i,n) aux[i+1] = (aux[i] + v[i]) % M;
17         forr(i,1,n+1) prefix[i] = (aux[i] + M - aux[i - (i&(-i))]) % M;
18     }
19     ll query(int l, int r) { //[a,b] 0-indexed
20         ll ans = 0; r++;
21         while(r) ans += prefix[r], r -= r&(-r), ans %= M;
22         while(l) ans += M - prefix[l], l -= l&(-l), ans %= M;

```

```

23     return ans;
24 }
25 void update(int pos, ll val) {
26     int i = pos + 1; ll upd = (val + M - a[pos]) % M;
27     while(i < prefix.size()) prefix[i] += upd, prefix[i] %= M, i += i&(-i);
28     a[pos] = val;
29 }
30 };
31
32 struct Hashing {
33     int n;
34     const ll MOD[2] = {999727999, 1070777777};
35     vector<ll> prefix[2], rev[2], pot[2], inv_pot[2];
36     ll P = 31, invP[2] = {inv_mod(P,MOD[0]), inv_mod(P,MOD[1])};
37     BIT s[2], rs[2];
38     Hashing() {}
39     Hashing(string &pal) {
40         n = pal.size();
41         forn(k,2) {
42             prefix[k].resize(n,0); rev[k].resize(n,0);
43             pot[k].resize(n,1); inv_pot[k].resize(n,1);
44             forn(i,n) {
45                 if(i) pot[k][i] = (pot[k][i-1] * P) % MOD[k];
46                 if(i) inv_pot[k][i] = (inv_pot[k][i-1] * invP[k]) % MOD[k];
47                 prefix[k][i] = (ll)(pal[i]-'a') * pot[k][i] % MOD[k];
48                 rev[k][i] = (ll)(pal[n-i-1]-'a') * pot[k][i] % MOD[k];
49             }
50             s[k] = BIT(prefix[k],MOD[k]);
51             rs[k] = BIT(rev[k],MOD[k]);
52         }
53     }
54     pair<ll,ll> get(int l, int r) { //[l,r] 0-indexed
55         ll x = (s[0].query(l,r) * inv_pot[0][l]) % MOD[0];
56         ll y = (s[1].query(l,r) * inv_pot[1][l]) % MOD[1];
57         return {x,y};
58     }
59     pair<ll,ll> getr(int l, int r) { //[l,r] 0-indexed
60         l = n-1-l, r = n-1-r; swap(l,r);
61         ll x = (rs[0].query(l,r) * inv_pot[0][l]) % MOD[0];
62         ll y = (rs[1].query(l,r) * inv_pot[1][l]) % MOD[1];
63         return {x,y};
64     }

```

```

65 void update(int pos, char a) {
66     ll val = (ll)(a-'a');
67     forn(k,2) {
68         s[k].update(pos,(val * pot[k][pos]) % MOD[k]);
69         assert(n-1-pos >= 0);
70         rs[k].update(n-1-pos,(val * pot[k][n-1-pos]) % MOD[k]);
71     }
72 }
73 };

```

8.3 Find Kth Substr Repetitions

```

1  typedef string str;
2  #define RB(x) (x<n?r[x]:0)
3
4  void csort(vector<int>& sa, vector<int>& r, int k){
5      int n=sa.size();
6      vector<int> f(max(255,n),0),t(n);
7      forr(i,0,n)f[RB(i+k)]++;
8      int sum=0;
9      forr(i,0,max(255,n))f[i]=(sum+=f[i])-f[i];
10     forr(i,0,n)t[f[RB(sa[i]+k)]++]=sa[i];
11     sa=t;
12 }
13
14 vector<int> suffix_array(string& s){ // O(n logn)
15     //~ s += '$';
16     int n=s.size(),rank;
17     vector<int> sa(n),r(n),t(n);
18     forr(i,0,n)sa[i]=i,r[i]=s[i];
19     for(int k=1;k<n;k*=2){
20         csort(sa,r,k);csort(sa,r,0);
21         t[sa[0]]=rank=0;
22         forr(i,1,n){
23             if(r[sa[i]]!=r[sa[i-1]]||RB(sa[i]+k)!=RB(sa[i-1]+k))rank++;
24             t[sa[i]]=rank;
25         }
26         r=t;
27         if(r[sa[n-1]]==n-1)break;
28     }
29     return sa;
30 }
31

```

```

32 vector<ll> prefix;
33
34 ll suma(int l, int r) {
35     return prefix[r+1] - prefix[l];
36 }
37
38 int BS(int l, int r, char &z, vector<int> &sa, str &pal, ll pos, ll k)
39 {
40     ll n = pal.size();
41     ll a = l-1, b = r;
42     while(b-a > 1) {
43         ll med = (a+b) / 2;
44         if((med-l+1)*(n-pos) - suma(l,med) >= k) b = med;
45         else a = med;
46     }
47     z = pal[sa[b]+pos];
48     return b;
49 }
50
51 int find_left(int l, int r, char z, vector<int> &sa, str &pal, int pos)
52 {
53     ll a = l-1, b = r;
54     ll n = pal.size();
55     while(b-a > 1) {
56         ll med = (a+b)/2;
57         if(sa[med]+pos >= n || pal[sa[med]+pos] != z) a = med;
58         else b = med;
59     }
60     return b;
61 }
62
63 int find_right(int l, int r, char z, vector<int> &sa, str &pal, int pos)
64 {
65     ll a = l, b = r+1;
66     ll n = pal.size();
67     while(b-a > 1) {
68         ll med = (a+b)/2;
69         if(sa[med]+pos >= n || pal[sa[med]+pos] != z) b = med;
70         else a = med;
71     }
72     return a;
73 }

```

```

72
73 void solve(str &pal, str &ans, int l, int r, int pos, ll k, vector <int>
    &sa) {
74     if(k <= 0) return;
75     vector <int> cont(30,0);
76     ll acum = 0; ll n = pal.size();
77     char z = '#';
78     int cr = BS(l,r,z,sa,pal,pos,k);
79     ans += z;
80
81     int new_start = find_left(l,cr,z,sa,pal,pos);
82     int new_end = find_right(cr,r,z,sa,pal,pos);
83     ll resta = (new_start - l) * (n - pos) - suma(l,new_start-1);
84     if(sa[new_start] + pos == n) new_start++;
85     solve(pal,ans,new_start,new_end,pos+1,k-resta-(new_end-new_start+1),
        sa);
86 }
87
88 string get_kth(string pal, ll k) {
89     vector <int> sa = suffix_array(pal);
90     prefix.resize(sa.size()+1,0);
91     for(int i = 0; i < sa.size(); i++) {
92         prefix[i+1] = prefix[i] + sa[i];
93     }
94     string ans;
95     solve(pal,ans,0,pal.size()-1,0,k,sa);
96
97     return ans;
98 }

```

8.4 Hashing

```

1 struct Hash {
2     const ll P=1777771;
3     const ll MOD[2] = {999727999, 1070777777};
4     const ll PI[2] = {325255434, 10018302};
5
6     vector<ll> h[2],pi[2];
7     Hash(string& s){
8         forn(k,2)h[k].resize(s.size()+1),pi[k].resize(s.size()+1);
9         forn(k,2){
10             h[k][0]=0;pi[k][0]=1;
11             ll p=1;

```

```

12         forr(i,1,s.size()+1){
13             h[k][i]=(h[k][i-1]+p*s[i-1])%MOD[k];
14             pi[k][i]=(1LL*pi[k][i-1]*PI[k])%MOD[k];
15             p=(p*P)%MOD[k];
16         }
17     }
18 }
19 ll get(ll s, ll e){ // [s, e] (s y e van de 0 a n-1)
20     e++;
21     ll h0=(h[0][e]-h[0][s]+MOD[0]);
22     h0=(1LL*h0*pi[0][s])%MOD[0];
23     ll h1=(h[1][e]-h[1][s]+MOD[1]);
24     h1=(1LL*h1*pi[1][s])%MOD[1];
25     return (h0<<32)|h1;
26 }
27 };

```

8.5 Manacher

```

1 vi d1; //d1[i] = max odd palindrome centered on i
2 vi d2; //d2[i] = max even palindrome centered on i
3 //s aabbaacaabbaa
4 //d1 1111117111111
5 //d2 0103010010301
6 void manacher(string& s){
7     ll l=0,r=-1,n=SZ(s);
8     d1.resize(n), d2.resize(n);
9     forn(i, n){
10         ll k = (i>r ? 1 : min(d1[l+r-i],r-i));
11         while(i+k<n && i-k>=0 && s[i+k]==s[i-k]) k++;
12         d1[i] = k--;
13         if(i+k>r) l=i-k,r=i+k;
14     }
15     l=0; r=-1;
16     forn(i, n){
17         ll k = (i>r ? 0 : min(d2[l+r-i+1],r-i+1)); k++;
18         while(i+k<=n && i-k>=0 && s[i+k-1]==s[i-k]) k++;
19         d2[i] = --k;
20         if(i+k-1>r) l=i-k, r=i+k-1;
21     }
22 }

```

8.6 Palindromic Tree

```

1 struct palindromic_tree{
2     static const int SIGMA=26;
3     struct Node{
4         int len, link, to[SIGMA];
5         ll cnt;
6         Node(int len, int link=0, ll cnt=1):len(len),link(link),cnt(cnt)
7         {
8             memset(to,0,sizeof(to));
9         }
10    };
11    vector<Node> ns;
12    int last;
13    palindromic_tree():last(0){ns.pb(Node(-1));ns.pb(Node(0));}
14    void add(int i, string &s){
15        int p=last, c=s[i]-'a';
16        while(s[i-ns[p].len-1]!=s[i])p=ns[p].link;
17        if(ns[p].to[c]){
18            last=ns[p].to[c];
19            ns[last].cnt++;
20        }else{
21            int q=ns[p].link;
22            while(s[i-ns[q].len-1]!=s[i])q=ns[q].link;
23            q=max(1,ns[q].to[c]);
24            last=ns[p].to[c]=SZ(ns);
25            ns.pb(Node(ns[p].len+2,q,1));
26        }
27    };

```

8.7 Prefix Function

```

1 vector<int> prefix_function(string s) {
2     int n = (int)s.length();
3     vector<int> pi(n);
4     forr(i,1,n) {
5         int j = pi[i-1];
6         while(j > 0 && s[i] != s[j]) j = pi[j-1];
7         if(s[i] == s[j]) j++;
8         pi[i] = j;
9     }
10    return pi;
11 }

```

8.8 Rabin Karp

```

1 /*
2     Dadas dos strings s y t, podemos hallar cuantas veces aparece el
3     string s en el string t
4     en O(|s|+|t|), mediante el uso de hashing.
5 */
6 vector<ll> rabin_karp(string s, string t) {
7     ll p=31, m=1e9+7, n=2147483647;
8     ll S=s.size(), T=t.size();
9     vector<ll> p_pow;
10    p_pow.pb(1);
11    forn(i,max(S,T)) p_pow.pb((p_pow.back()*p)%m);
12    vector<ll> h; h.pb(0);
13    forn(i,T) h.pb((h.back()+(t[i]-'a'+1)*p_pow[i])%m);
14    ll h_s=0;
15    forn(i,S) h_s=(h_s+(s[i]-'a'+1)*p_pow[i])%m;
16    vector<ll> o;
17    forn(i,T-S+1) {
18        ll cur_h=(h[i+S]+m-h[i])%m;
19        if(cur_h==(h_s*p_pow[i])%m) o.pb(i);
20    }
21    return o;
22 }

```

8.9 Suffix Array

```

1 #define RB(x) (x<n?r[x]:0)
2 void csort(vector<int>& sa, vector<int>& r, int k){
3     int n=sa.size();
4     vector<int> f(max(255,n),0),t(n);
5     forr(i,0,n)f[RB(i+k)]++;
6     int sum=0;
7     forr(i,0,max(255,n))f[i]=(sum+=f[i])-f[i];
8     forr(i,0,n)t[f[RB(sa[i]+k)]++]=sa[i];
9     sa=t;
10 }
11
12 vector<int> suffix_array(string& s){ // O(n logn)
13     s += '$';
14     int n=s.size(),rank;
15     vector<int> sa(n),r(n),t(n);

```

```

16     forr(i,0,n)sa[i]=i,r[i]=s[i];
17     for(int k=1;k<n;k*=2){
18         csort(sa,r,k);csort(sa,r,0);
19         t[sa[0]]=rank=0;
20         forr(i,1,n){
21             if(r[sa[i]]!=r[sa[i-1]]||RB(sa[i]+k)!=RB(sa[i-1]+k))rank++;
22             t[sa[i]]=rank;
23         }
24         r=t;
25         if(r[sa[n-1]]==n-1)break;
26     }
27     return sa;
28 }
29
30 vector<int> lcp_construction(string const& s, vector<int> const& p) {
31     int n = s.size();
32     vector<int> rank(n, 0);
33     for (int i = 0; i < n; i++) rank[p[i]] = i;
34     int k = 0;
35     vector<int> lcp(n-1, 0);
36     for (int i = 0; i < n; i++) {
37         if(rank[i] == n - 1) {
38             k = 0;
39             continue;
40         }
41         int j = p[rank[i] + 1];
42         while (i + k < n && j + k < n && s[i+k] == s[j+k]) k++;
43         lcp[rank[i]] = k;
44         if(k) k--;
45     }
46     return lcp;
47 }
48
49 bool substr_search(string &text, string &pal, vector<int> &sa) {
50     int n = text.size(), a = -1, b = n-1;
51     while(b-a > 1) {
52         int med = (a+b)/2, pos = sa[med]; bool d = false;
53         int tam = min((int)pal.size(),int(n-sa[med]));
54         string check = text.substr(sa[med],tam);
55         if(check < pal) a=med;
56         else b = med;
57     }
58     int tam = min(int(pal.size()),int(n-sa[b]));

```

```

59     string fin = text.substr(sa[b],tam);
60     return fin == pal ? true : false;
61 }
62
63 ll count_substring(string &pal) {
64     ll n = pal.size();
65     vector<int> sa = suffix_array(pal);
66     vector<int> lcp = lcp_construction(pal,sa);
67     ll ans = n * (n+1) / 2;
68     for(int u : lcp) ans -= (ll)u;
69     return ans;
70 }
71
72 #define SZ(a) ((int)a.size())
73 string lcsbstr(string s, string t) {
74     string r = s + '#' + t + '$';
75     int best = 0; int pos = 0;
76     vector<int> sa = suffix_array(r);
77     vector<int> lcp = lcp_construction(r, sa);
78     forr(i, 1, SZ(sa)) {
79         if(isInRange(sa[i] - 1, 0, SZ(s)) && isInRange(sa[i], SZ(s) + 1,
80             SZ(r) - 1)) {
81             if(lcp[i-1] > best) best = lcp[i-1], pos = sa[i];
82         }
83         if(isInRange(sa[i], 0, SZ(s)) && isInRange(sa[i] - 1, SZ(s) + 1,
84             SZ(r) - 1)) {
85             if(lcp[i-1] > best) best = lcp[i-1], pos = sa[i];
86         }
87     }
88     return r.substr(pos,best);
89     // To find de LCSbstr between n string, we can concatenate the n
90     // strings and apply min_segtree in lcp ans intervals with at least
91     // one suffix of each string.

```

8.10 Suffix Automaton

```

1 struct state {
2     int len, link;
3     map<char,int> next;
4 };
5
6 const int MAXN = 300010;

```

```

7 state st[MAXN*2];
8 int sz, last;
9
10 void sa_init(){
11     forn(i,2*MAXN) {
12         st[i].next.clear();
13         st[i].link = 0;
14         st[i].len = 0;
15     }
16     last= st[0].len=0; sz=1;
17     st[0].link=-1;
18 }
19
20 void sa_extend(char c) {
21     int cur = sz++;
22     st[cur].len = st[last].len + 1;
23     int p;
24     for(p = last; p != -1 && !st[p].next.count(c); p = st[p].link) {
25         st[p].next[c] = cur;
26     }
27     if (p == -1) st[cur].link = 0;
28     else {
29         int q = st[p].next[c];
30         if(st[p].len + 1 == st[q].len) st[cur].link = q;
31         else {
32             int clone = sz++;
33             st[clone].len = st[p].len + 1;
34             st[clone].next = st[q].next;
35             st[clone].link = st[q].link;
36             while (p != -1 && st[p].next[c] == q) {
37                 st[p].next[c] = clone;
38                 p = st[p].link;
39             }
40             st[q].link = st[cur].link = clone;
41         }
42     }
43     last = cur;
44 }
45
46 /* ..... */
47
48 vector <bool> visto(2*MAXN,false);
49

```

```

50 void count_substr(int cur) {
51     // Find #distinct substrings from vertex cur
52     visto[cur] = true;
53     st[cur].total = 1;
54     for(pair<char,int> u : st[cur].next) {
55         if(!visto[u.second]) count_substr(u.second);
56         st[cur].total += st[u.second].total;
57     }
58 }
59
60 /* ..... */
61
62 string ans;
63
64 void find_kth(int v, ll k) { // Find kth substring (all different)
65     if(k <= 0) return;
66     ll acum = 0;
67     for(pair<char,int> u : st[v].next) {
68         if(acum + st[u.second].total >= k) {
69             ans += u.first;
70             find_kth(u.second, k-acum-1);
71             return;
72         }
73         acum += st[u.second].total;
74     }
75 }

```

8.11 Trie

```

1 const int K = 26;
2
3 struct Vertex {
4     int next[K];
5     int leaf;
6
7     Vertex() {
8         fill(begin(next), end(next), -1);
9         leaf = 0;
10    }
11 };
12
13 vector<Vertex> trie(1);
14

```



```

15 void add_string(string const& s) {
16     int v = 0;
17     for (char ch : s) {
18         int c = ch - 'a';
19         if (trie[v].next[c] == -1) {
20             trie[v].next[c] = trie.size();
21             trie.emplace_back();
22         }
23         v = trie[v].next[c];
24     }
25     trie[v].leaf++;
26 }
27
28 bool find(string &pal) {
29     int cur = 0;
30     for(char u : pal) {
31         cur = trie[cur].next[u-'a'];
32         if(cur == -1) return false;
33     }
34     return true;
35 }

```

8.12 Z Function

```

1 vector<int> z_function(string s) {
2     int n = (int) s.size();
3     vector<int> z(n);
4     for (int i = 1, l = 0, r = 0; i < n; ++i) {
5         if (i <= r)
6             z[i] = min (r - i + 1, z[i - l]);
7         while (i + z[i] < n && s[z[i]] == s[i + z[i]])
8             z[i]++;
9         if (i + z[i] - 1 > r)
10            l = i, r = i + z[i] - 1;
11     }
12     return z;
13 }

```

9 Testing

9.1 A

```

1 // incorrect solution for finding second smallest element
2 #include <bits/stdc++.h>
3 using namespace std;
4 // #warning TODO: remember about bigger N
5 const int MAX_N = 1e6 + 5;
6 int a[MAX_N];
7 set<int> asd[MAX_N], asdfasd[3*MAX_N];
8 int main() {
9     int n;
10    scanf("%d", &n);
11    for(int i = 1; i <= n; ++i) {
12        scanf("%d", &a[i]);
13    }
14    sort(a + 1, a + n);
15    printf("%d\n", a[2]);
16 }

```

9.2 Brute

```

1 // slow solution for finding second smallest element
2 #include <bits/stdc++.h>
3 using namespace std;
4 int main() {
5     int n;
6     cin >> n;
7     vector<int> a(n);
8     for(int& x : a) {
9         cin >> x;
10    }
11    for(int x : a) {
12        int count_smaller = 0;
13        for(int y : a) {
14            if(y < x) {
15                ++count_smaller;
16            }
17        }
18        if(count_smaller == 1) {
19            cout << x;
20            return 0;
21        }
22    }
23 }

```

```

21     }
22 }
23 assert(false);
24 }

```

9.3 Gen

```

1 // generating a random sequence of distinct elements
2 #include <bits/stdc++.h>
3 using namespace std;
4
5 int rand(int a, int b) {
6     return a + rand() % (b - a + 1);
7 }
8
9 int main(int argc, char* argv[]) {
10     srand(atoi(argv[1])); // atoi(s) converts an array of chars to int
11     int n = rand(2, 10);
12     printf("%d\n", n);
13     set<int> used;
14     for(int i = 0; i < n; ++i) {
15         int x;
16         do {
17             x = rand(1, 10);
18         } while(used.count(x));
19         printf("%d_", x);
20         used.insert(x);
21     }
22     puts("");
23 }

```

9.4 Gen Tree

```

1 // generating a tree in a simple way
2 #include <bits/stdc++.h>
3 using namespace std;
4
5 int rand(int a, int b) {
6     return a + rand() % (b - a + 1);
7 }
8
9 int main(int argc, char* argv[]) {
10     srand(atoi(argv[1]));
11     int n = rand(2, 20);

```

```

12     printf("%d\n", n);
13     for(int i = 2; i <= n; ++i) {
14         printf("%d_%d\n", rand(1, i - 1), i);
15     }
16 }

```

9.5 Gen Tree2

```

1 // generating a tree in a not-so-stupid way
2 #include <bits/stdc++.h>
3 using namespace std;
4
5 int rand(int a, int b) {
6     return a + rand() % (b - a + 1);
7 }
8
9 int main(int argc, char* argv[]) {
10     srand(atoi(argv[1]));
11     int n = rand(2, 20);
12     printf("%d\n", n);
13     vector<pair<int,int>> edges;
14     for(int i = 2; i <= n; ++i) {
15         edges.emplace_back(rand(1, i - 1), i);
16     }
17
18     vector<int> perm(n + 1); // re-naming vertices
19     for(int i = 1; i <= n; ++i) {
20         perm[i] = i;
21     }
22     random_shuffle(perm.begin() + 1, perm.end());
23
24     random_shuffle(edges.begin(), edges.end()); // random order of edges
25
26     for(pair<int, int> edge : edges) {
27         int a = edge.first, b = edge.second;
28         if(rand() % 2) {
29             swap(a, b); // random order of two vertices
30         }
31         printf("%d_%d\n", perm[a], perm[b]);
32     }
33 }

```

9.6 In

```

1 4
2 10 30 20 40

```

9.7 Readme.Md

```

1 ##### How to use the stress tester (Thanks Errichto)
2 - 'a.cpp' - incorrect solution
3 - 'brute.cpp' - correct slow solution
4 - 'in' - sample input
5 - 'gen.cpp' - test generator
6 - 's.sh' - script that tests a.cpp and brute.cpp against each other
7 - 'gen_tree.cpp' - stupid tree generator
8 - 'gen_tree2.cpp' - better tree generator
9
10 My compilation flags:
11
12 1. fast running time
13     'g++ -O2 -std=c++17 -Wno-unused-result -Wshadow -Wall -o a a.cpp'
14
15 2. check for mistakes
16     'g++ -std=c++17 -Wshadow -Wall -o a a.cpp
17     -fsanitize=address -fsanitize=undefined -D_GLIBCXX_DEBUG -g'

```

9.8 Stress

```

1 for((i = 1; ; ++i)); do
2     echo $i
3     ./gen $i > int
4     # ./a < int > out1
5     # ./brute < int > out2
6     # diff -w out1 out2 || break
7     diff -w -y <(. /a < int) <(. /brute < int) || break
8 done

```

10 Tree

10.1 Binary Lifting

```

1 struct binary_lifting {
2     vector<vi> jump;
3     binary_lifting(vi par){
4         int n = SZ(par);
5         jump.resize(LOG, vi(n));
6         jump[0] = par;
7
8         forr(j, 1, LOG){
9             forn(i, n) jump[j][i] = jump[j-1][jump[j-1][i]];
10        }
11    }
12
13    int lift(int v, int k){
14        for(int i = LOG-1; i >= 0; i--)
15            if(k & (1<<i)) v = jump[i][v];
16        return v;
17    }
18 };

```

10.2 Centroid

```

1 struct centroid {
2     vector<vector<int>> g; int n;
3     vector<vector<int>> c_tree;
4     bool tk[MAXN];
5     int fat[MAXN]; // father in centroid decomposition
6     int szt[MAXN]; // size of subtree
7     int centro = -1;
8     int calcsz(int x, int f){
9         szt[x]=1;
10        for(auto y:g[x])if(y!=f&&!tk[y])szt[x]+=calcsz(y,x);
11        return szt[x];
12    }
13    void cdfs(int x=0, int f=-1, int sz=-1){ // 0(nlogn)
14        if(sz<0)sz=calcsz(x,-1);
15        for(auto y:g[x])if(!tk[y]&&szt[y]*2>=sz){
16            szt[x]=0;cdfs(y,f,sz);return;
17        }
18        tk[x]=true;fat[x]=f;

```

```

19     for(auto y:g[x])if(!tk[y])cdfs(y,x);
20 }
21 centroid(vector <vector<int>> gg, int nn) {
22     g = gg; n = nn; memset(tk,false,sizeof(tk));cdfs();
23     c_tree.resize(n);
24     forn(i,n) {
25         if(fat[i] == -1) centro = i;
26         else c_tree[fat[i]].pb(i);
27     }
28 }
29 };

```

10.3 Centroid Cses

```

1 vector<int> g[MAXN];int n;
2 bool tk[MAXN];
3 int fat[MAXN]; // father in centroid decomposition
4 int szt[MAXN]; // size of subtree
5 int calcsz(int x, int f){
6     szt[x]=1;
7     for(auto y:g[x])if(y!=f&&!tk[y])szt[x]+=calcsz(y,x);
8     return szt[x];
9 }
10 void cdfs(int x=0, int f=-1, int sz=-1){ // 0(nlogn)
11     if(sz<0)sz=calcsz(x,-1);
12     for(auto y:g[x])if(!tk[y]&&szt[y]*2>=sz){
13         szt[x]=0;cdfs(y,f,sz);return;
14     }
15     tk[x]=true;fat[x]=f;
16     for(auto y:g[x])if(!tk[y])cdfs(y,x);
17 }
18 void centroid(){memset(tk,false,sizeof(tk));cdfs();}

```

10.4 Diameter

```

1 vector < vector <int> > g(MAXN);
2 vector <bool> visto(MAXN,false);
3
4 void DFS(int v, int &new_start, int &ans, int l) {
5     visto[v] = true;
6     if(l > ans) ans = l, new_start = v;
7     for(int u : g[v]) if(!visto[u]) DFS(u,new_start,ans,l+1);
8 }
9

```

```

10 int find_diameter() {
11     int ans = 0, new_start;
12     DFS(0,new_start,ans,0); fill(all(visto),false);
13     DFS(new_start,new_start,ans,0);
14     return ans;
15 }

```

10.5 Dsu On Tree

```

1 #define rs(n) resize(n)
2
3 struct DSU_on_Tree { // count #vertex with color c in subtree v
4     vector<map<ll, ll>> h; // histogram [color,count]
5     vi color, link, ans;
6     vii sum;
7     vector<vi> g;
8
9     DSU_on_Tree(int n, vector<vi> _g, vi _c) : g(_g), color(_c) {
10         ans.rs(n+5); link.rs(n+5); h.rs(n+5); sum.rs(n+5);
11         forn(i,n+5) link[i] = i, sum[i] = {0,0};
12         dfs(0);
13     }
14
15     int find(int x) {
16         return link[x] = (link[x] == x ? x : find(link[x]));
17     }
18
19     void add(int v, int c, int cnt) { // add [color,cnt] to root v
20         ll &cur = h[v][c];
21         cur += cnt;
22         if(sum[v].first == cur) sum[v].second += c;
23         else if(sum[v].first < cur) sum[v] = {cur,c};
24     }
25
26     void merge(int i, int j){
27         if(SZ(h[i]) < SZ(h[j])) swap(i,j);
28         link[j] = i;
29         for(auto x : h[j]){
30             int c, cnt; tie(c,cnt) = x;
31             add(i,c,cnt);
32         }
33         h[j].clear();
34     }

```

```

35
36 void dfs(ll v, ll p = -1){
37     add(v,color[v],1);
38     for(auto u : g[v]) {
39         if(u == p) continue;
40         dfs(u, v);
41         merge(find(v), find(u));
42     }
43     ans[v] = sum[find(v)].second; // here solve for vertex v
44 }
45 };

```

10.6 Dynamic Connectivity

```

1 struct UnionFind {
2     int n,comp;
3     vi uf,si,c;
4     UnionFind(int n=0):n(n),comp(n),uf(n),si(n,1){
5         forn(i,n) uf[i]=(int)i;
6     }
7     int find(int x){return x==uf[x]?x:find(uf[x]);}
8     bool join(int x, int y){
9         if((x=find(x))==find(y))return false;
10        if(si[x]<si[y])swap(x,y);
11        si[x]+=si[y];uf[y]=x;comp--;c.pb(y);
12        return true;
13    }
14    int snap(){return SZ(c);}
15    void rollback(int snap){
16        while(SZ(c)>snap){
17            int x=c.back();c.pop_back();
18            si[uf[x]]-=si[x];uf[x]=x;comp++;
19        }
20    }
21 };
22 enum {ADD,DEL,QUERY};
23 struct Query {int type,x,y};
24 struct DynCon {
25     vector<Query> q;
26     UnionFind dsu;
27     vector<int> mt;
28     map<pair<int,int>,int> last;
29     DynCon(int n):dsu(n){}

```

```

30 void add(int x, int y){
31     if(x>y)swap(x,y);
32     q.pb({ADD,x,y});mt.pb(-1);last[{x,y}]=SZ(q)-1;
33 }
34 void remove(int x, int y){
35     if(x>y)swap(x,y);
36     q.pb({DEL,x,y});
37     int pr=last[{x,y}];mt[pr]=SZ(q)-1;mt.pb(pr);
38 }
39 void query(int x, int y){q.pb({QUERY,x,y});mt.pb(-1);} // modificar
    que query se tiene que usar
40 void process(){ // answers all queries in order
41     if(!SZ(q)) return;
42     forn(i,SZ(q))if(q[i].type==ADD&&mt[i]<0)mt[i]=SZ(q);
43     go(0,SZ(q));
44 }
45 void go(int s, int e){
46     if(s+1==e){
47         if(q[s].type==QUERY) // answer query using DSU q[s]
48             assert(false); // poner alguna operacion entre q[s].x y
                q[s].y
49         return;
50     }
51     int k=dsu.snap(),m=(s+e)/2;
52     for(int i=e-1;i>m;--i)if(mt[i]>=0&&mt[i]<s)dsu.join(q[i].x,q[i]
        ].y);
53     go(s,m);dsu.rollback(k);
54     for(int i=m-1;i>s;--i)if(mt[i]>=e)dsu.join(q[i].x,q[i].y);
55     go(m,e);dsu.rollback(k);
56 }
57 };

```

10.7 Find Centroid

```

1 vector < vector <int> > g(MAXN);
2 vector <bool> is_centroid(MAXN,true);
3 vector <int> sz(MAXN,0);
4
5 void DFS(int v, int prev, int n) {
6     sz[v] = 1;
7     for(int u : g[v]) {
8         if(u == prev) continue;
9         DFS(u,v,n);

```

```

10     sz[v] += sz[u];
11     if(sz[u] > n/2) is_centroid[v] = false;
12 }
13 if(n - sz[v] > n/2) is_centroid[v] = false;
14 }
15
16 vector<int> find_centroid(int n) {
17     // Centroid: A node so that each subtree has at most floor(n/2)
18     // nodes
19     vector<int> ans;
20     DFS(1,-1,n);
21     forr(i,1,n+1) if(is_centroid[i]) ans.pb(i);
22     return ans;
23 }

```

10.8 Hld

```

1 struct hld {
2     #define rs(x) resize(x)
3     int n;
4     vector<vector<int>> g;
5     vector<tipo> values, sz, heavy, parent, depth, top, r, pos;
6     segtree st; // iterative is faster
7
8     hld(int _n, vector<vector<int>> _g, vector<tipo> _v) {
9         n = _n; g = _g; values = _v;
10        sz.rs(n); parent.rs(n); depth.rs(n); top.rs(n); pos.rs(n);
11        heavy.rs(n);
12        forn(i,n) heavy[i] = i;
13        dfs(0);
14        get_order(0,0);
15        st.build(r,r.size());
16    }
17
18    ll dfs(int v, int p = -1, int d = 0) {
19        depth[v] = d; parent[v] = p;
20        for(int &u : g[v]) {
21            if(u == p) continue;
22            ll tam = dfs(u,v,d+1);
23            // move heavy node to the beginning
24            if(sz[u] > sz[heavy[v]]) heavy[v] = u, swap(u,g[v][0]);
25            sz[v] += tam;
26        }

```

```

27        sz[v] += 1;
28        return sz[v];
29    }
30
31    void get_order(int v, int t, int p = -1) { //first child is heavy
32        pos[v] = r.size(); top[v] = t; r.pb(values[v]);
33        for(int u : g[v]) {
34            if(u == p) continue;
35            if(u == heavy[v]) get_order(u,t,v);
36            else get_order(u,u,v);
37        }
38    }
39
40    void op(ll &res, int a, int b) { // Set operation
41        res = max(res, st.query(pos[a],pos[b]));
42        // ----- for segtree recursive use -----
43        // res = max(res, st.query(pos[a],pos[b]).ans);
44    }
45
46    tipo query_hld(int a, int b) {
47        tipo res = NEUT;
48        for(; top[a] != top[b]; b = parent[top[b]]) {
49            if(depth[top[a]] > depth[top[b]]) swap(a, b);
50            op(res,top[b],b);
51        }
52        if(depth[a] > depth[b]) swap(a, b);
53        op(res,a,b);
54        return res;
55    }
56
57    void update_hld(int p, tipo val) { st.update(pos[p],val); }
58 }

```

10.9 Isomorphism Centroid

```

1 struct centroid {
2     vector<vector<int>> g;
3     vector<bool> is_centroid;
4     vector<int> sz;
5
6     centroid(vector<vector<int>> &graph, int tam) {
7         g = graph;
8         is_centroid.resize(tam,true);

```

```

9         sz.resize(tam,0);
10     }
11
12     void DFS(int v, int prev, int n) {
13         sz[v] = 1;
14         for(int u : g[v]) {
15             if(u == prev) continue;
16             DFS(u,v,n);
17             sz[v] += sz[u];
18             if(sz[u] > n/2) is_centroid[v] = false;
19         }
20         if(n - sz[v] > n/2) is_centroid[v] = false;
21     }
22
23     vector<int> find_centroid(int n) {
24         // Centroid: A node so that each subtree has at most floor(n/2)
25         // nodes
26         vector<int> ans;
27         DFS(0,-1,n);
28         forn(i,n) if(is_centroid[i]) ans.pb(i);
29         return ans;
30     }
31
32     struct iso {
33         vector<vector<int>> t1, t2;
34         map<vector<int>,int> mapa;
35         int idx = 0;
36
37         iso(vector<vector<int>> &tree1, vector<vector<int>> &tree2) {
38             t1 = tree1; t2 = tree2;
39         }
40
41         int dfs(int cur, int p, vector<vector<int>> &g) {
42             vector<int> v;
43             for(int u : g[cur]) {
44                 if(u != p) v.push_back(dfs(u,cur,g));
45             }
46             sort(all(v));
47             if(!mapa.count(v)) mapa[v] = idx++;
48             return mapa[v];
49         }
50     }

```

```

51     bool check_iso(vector<int> &centroid1, vector<int> &centroid2) {
52         int s1 = dfs(centroid1[0],-1,t1);
53         for(int u : centroid2) {
54             int s2 = dfs(u,-1,t2);
55             if(s1 == s2) return true;
56         }
57         return false;
58     }
59 };

```

10.10 Lca

```

1  /*
2  Tiempo de build: O(nlogn)
3  Memoria: O(nlogn)
4  Tiempo por query: O(logn)
5  */
6  const int MAXN = 200005, LOG = 20;
7
8  struct LCA {
9      int n, root;
10     vector<vector<int>> g;
11     int jmp[MAXN][LOG], depth[MAXN]; // jmp[i][j] tiene el 2^j-esimo
12     // ancestro de i
13     void lca_dfs(int x) {
14         for(int u : g[x]) {
15             if(u == jmp[x][0]) continue;
16             jmp[u][0] = x; depth[u] = depth[x]+1;
17             lca_dfs(u);
18         }
19     }
20
21     LCA(int tam, vector<vector<int>> &tree, int r): n(tam),root(r),g(
22     tree) {
23         depth[root] = 0;
24         memset(jmp,-1,sizeof(jmp)); jmp[root][0] = root;
25         lca_dfs(root);
26         forr(k, 1, LOG){ forn(i, n){
27             if(jmp[i][k-1]<0) jmp[i][k] = -1;
28             else jmp[i][k] = jmp[jmp[i][k-1]][k-1];
29         }
30     }

```

```

30
31 int lca(int x, int y){
32     if(depth[x] < depth[y]) swap(x,y);
33     for(int i = LOG-1; i >= 0; i--) {
34         if(depth[x]-(1<<i) >= depth[y]) x = jmp[x][i];
35     }
36     if(x == y) return x;
37     for(int i = LOG-1; i >= 0; i--) {
38         if(jmp[x][i] != jmp[y][i]) x = jmp[x][i], y = jmp[y][i];
39     }
40     return jmp[x][0];
41 }
42
43 int dist(int x, int y) {
44     return depth[x] + depth[y] - 2 * depth[lca(x,y)];
45 }
46 };

```

10.11 Link Cut Tree

```

1 const int MAXN = 5e5+5;
2 int ls[MAXN],rs[MAXN],fa[MAXN],siz[MAXN],st[MAXN],ch[MAXN][2],siz2[MAXN]
   ];
3 bool rev[MAXN];
4 #define ls(x) ch[x][0]
5 #define rs(x) ch[x][1]
6 inline bool notrt(int x){return ls(fa[x])==x||rs(fa[x])==x;}
7 inline void pushdown(int x){
8     if(rev[x]){
9         if(ls(x)) ls(ls(x))^=rs(ls(x))^=ls(ls(x))^=rs(ls(x)),rev[ls(x)]
           ]^=1;
10        if(rs(x)) ls(rs(x))^=rs(rs(x))^=ls(rs(x))^=rs(rs(x)),rev[rs(x)]
           ]^=1;
11        rev[x]=0;
12    }
13 }
14 inline void rotate(int x){
15     int y=fa[x],z=fa[y];bool l=rs(y)==x,r=!l;
16     if(notrt(y)) ch[z][rs(z)==y]=x;if(ch[x][r]) fa[ch[x][r]]=y;
17     fa[y]=x,fa[x]=z,ch[y][l]=ch[x][r],ch[x][r]=y;
18     siz[y]=siz[ls(y)]+siz[rs(y)]+1+siz2[y],siz[x]=siz[ls(x)]+siz[rs(x)]
       ]+1+siz2[x];
19 }

```

```

20 inline void splay(int x){
21     int y=x,z=1;st[1]=y;
22     while(notrt(y)) st[++z]=y=fa[y];
23     while(z) pushdown(st[z--]);
24     while(notrt(x)){
25         y=fa[x],z=fa[y];
26         if(notrt(y)) if(ls(z)==y^ls(y)==x) rotate(x);else rotate(y);
           rotate(x);
27     }
28 }
29 inline void access(int x){int y=0;while(x) splay(x),siz2[x]+=siz[rs(x)]-
   siz[y],rs(x)=y,siz[x]=siz[ls(x)]+siz[rs(x)]+1+siz2[x],x=fa[y=x];}
30 inline void make(int x){access(x),splay(x),ls(x)^=rs(x)^=ls(x)^=rs(x),
   rev[x]^=1;}
31 inline void split(int x,int y){make(x),access(y),splay(y);}
32 inline void link(int x,int y){make(x),make(y),fa[x]=y,siz2[y]+=siz[x];}
33 inline void cut(int x,int y){split(x,y),fa[x]=ls(y)=0,siz[y]=siz[ls(y)]+
   siz[rs(y)]+1+siz2[y];}

```

10.12 Sm To Large

```

1 // Small to large technique.
2 struct query { // Queries to answer, v: vertex
3     ll v, h, idx;
4 };
5
6 vi g[MAXN];
7 vector<query> q[MAXN]; // Queries to answer
8 vi ans(MAXN, -1); // Answer to each query
9
10 unordered_map<ll, ll> cnt[MAXN] // The structures to store the
   information and merge
11
12 // Merge operation, merge small to large and return the large one
13 int merge(int v, int u){
14     if(SZ(cnt[v]) < SZ(cnt[u])) swap(u, v); // now v is the large one
15
16     // Merge cnt[u] into cnt[v]
17     for(auto [x, y]: cnt[u]){
18         // Do something with x, y and cnt[v]
19     }
20     cnt[u].clear(); // Clear the small one to mantain memory in O(n)
21

```



```
22     return v; // return the large node
23 }
24
25 // Process the queries of v, v_repr is the representative of v (large
    node after merging v and its children)
26 void process_queries(int v, int v_repr){
27     for(auto &[_v, k, i]: q[v]) ans[i] = cnt_geq[v_repr][k];
28 }
29
30 string s;
31
32 int dfs(int v, int p){
33     int v_repr = v // Initialize the representative of v
34
35     // Initialize the storage structs of only v
36
37     for(auto u: g[v]){
38         if(u == p) continue;
39         int u_repr = dfs(u, v); // Get the representative of u
40         v_repr = merge(v_repr, u_repr); // Merge u_repr into v_repr
41     }
42
43     process_queries(v, v_repr); // Offline process the queries of v
44
45     return v_repr;
46 }
47
48
49 void solve(){
50     int n, m; cin >> n >> m;
51     color.resize(n); forn(i, n) cin >> color[i];
52
53     forn(i, n-1){
54         int u, v; cin >> u >> v; u--, v--;
55         g[u].pb(v); g[v].pb(u);
56     }
57
58     forn(i, m){
59         int v, k; cin >> v >> k; v--;
60         q[v].pb({v, k, i}); // Add the query to the list of queries of v
61     }
62     dfs(0, -1);
63 }
```

```
64     forn(i, m) cout << ans[i] << "\n";
65 }
66
67 int main(){
68     FIN;
69
70     solve();
71
72     return 0;
73 }
```