

12- Convergence of simulated annealing: how to formulate it, how to compute it.

→ Does simulated annealing find the global optimum and under which conditions?

Simulated annealing can be analysed mathematically, it converges in probability.

It gives a solution arbitrary close to the global optimum with a probability arbitrary close to 1

→ The conditions are:

- Movements are invertible
- Any point in the search space can be reached from any other point in a finite number of movements
- The initial temperature must be large enough (do exploitation before exploration)
- The temperature should not decrease too fast, at iteration t :

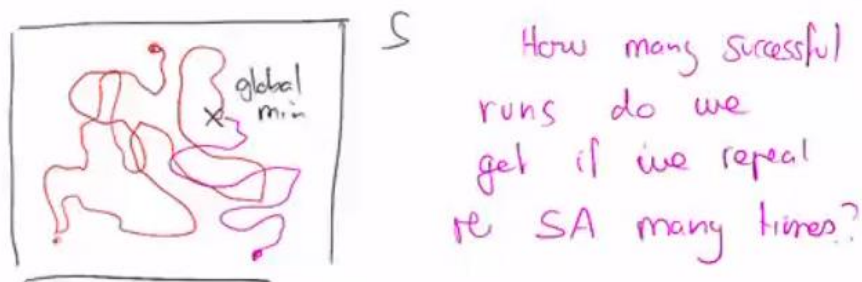
$$T(t) \sim \frac{C}{\log t} \quad \text{for } t \gg 1$$

C is an unknown constant that reflects the variation of E in S .

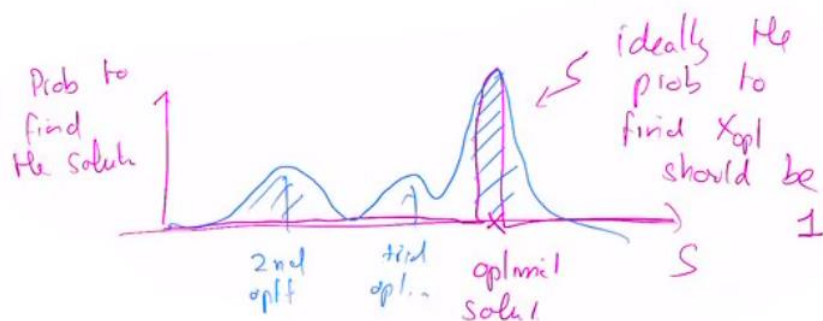
!!! If temperature is decreased too fast, we risk not converging (quenching) !!!

→ In practice, $\log t$ is too slow. So, temperature will decrease faster

How to compute it:



If we get 100% of success there is convergence, in practice it may end up sub optimal



We want to compute the probability $P(t, i)$ that the SA is at point $i \in S$ at iteration t

$$P(t+1, j) = \sum_i P(t, i) W_{ij}(t)$$



W_{ij} is the transition probability from i to j . Since we have

100 possible values of i and j , W_{ij} is a

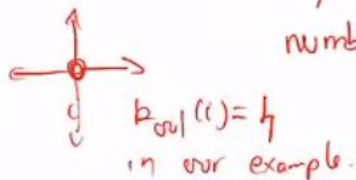
100x100 matrix.

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For SA, we have:

$$W_{ij} = \begin{cases} 0 & \text{if } j \text{ is not a neighbor of } i \\ \frac{1}{k_{out}(i)} \times P_{\text{metropolis}}(E_i, E_j, T(t)) & \end{cases}$$

\uparrow number of neighbors of i $\uparrow \min(e^{-(E_j - E_i)/T(t)})$



We also add the following

$$W_{ii} = 1 - \sum_{j \text{ neighbor of } i} W_{ij}$$