


SLAM

Introduction & Data Structures

EDAA - G06

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Problem Definition



What's already done

- Filter noise/undesirable effects in data:
 - Multiple reflections;
 - Echos;
 - Self reflections;
 - Multipath errors.
- Represent map using Octomaps;
- Implement dynamic probabilistic mapping algorithms based on sonar data;

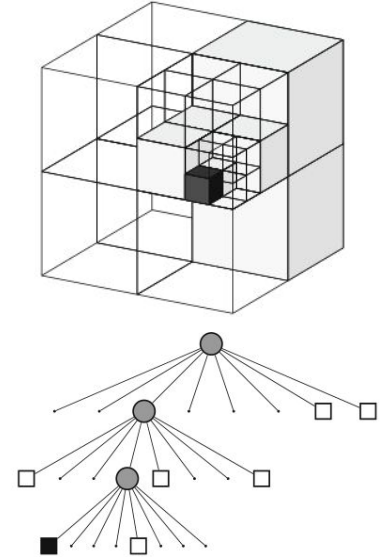


Fig 1. Example of octree storing occupied (black) and free (white) cells (Hornung et al.)

What's next in plan

- Expand the implementation to support 3D mapping;
- Support incremental map growth;
- Improve performance;
- Implement new data structures;
 - Hash Sets/Hash Maps
- Implement new algorithms;
 - Set collision resolution techniques;
 - Explore ray casting solutions for 3D;
 - Explore more edge detection algorithms;

Data structure

Sets using Hash Maps

- The operation for set merge uses about 60% of ray cast's execution time;
- Simple to implement;
- Low space usage;
- Can be tinkered with to better fit our problem definition.

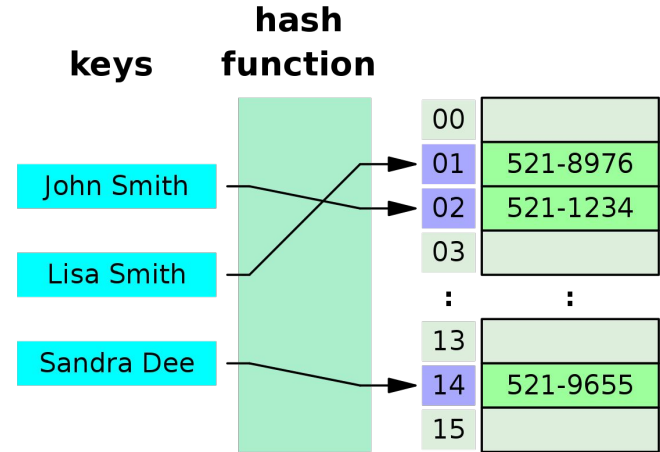


Fig 2. Hash Map Representation

Hash Sets Details

- The index to store the value corresponds to its key and is calculated using the following formula: $\text{hash} \% \text{size}$.
- Each Hash Map only needs to maintain a list of its stored elements and the number of elements it's currently storing.
- Stores each element once and only once.
- The elements to be stored need a good hash function;
- Set collision resolutions technics change the way elements with the same hash are stored.

| | |
|-------|---------|
| Key 1 | Value 1 |
| Key 2 | Value 2 |
| Key 3 | Value 3 |

Fig 3. Hash Map Reference

Collision Resolution - Linear

- If two elements have the same hash, the next index is tested (essentially **hash = hash + 1**);
- Very simple to implement and verify;
- Not the most efficient algorithm.
- Poor hash functions affect this algorithm the most.
- **Deleted node is not removed, but marked as "removed".**

Insert (55)

$55 \% 7 = 6$

| | |
|---|----|
| 0 | 47 |
| 1 | 55 |
| 2 | 93 |
| 3 | 10 |
| 4 | |
| 5 | 40 |
| 6 | 76 |

Fig 4. [Linear Probing](#)

Collision Resolution - Quadratic

- If two elements have the same hash, we try **index = hash + nCollision²**.
- Also easy to implement, but can be a bit harder to verify;
- Can create **fewer** collisions than linear probing.
- **Deleted node is not removed, but marked as "removed"**.

| | |
|---|-----|
| 0 | 700 |
| 1 | 50 |
| 2 | 85 |
| 3 | |
| 4 | |
| 5 | 92 |
| 6 | 76 |

Insert 92:

Collision occurs at 1.

Collision occurs at $1 + 1*1$ position

Insert at $1 + 2*2$ position.

Fig 5. Quadratic Probing

Collision Resolution - Double Hashing

- If two elements have the same hash, we try **hash + nCollision*hash**.
- Hard to verify;
- Creates **fewer** collisions in hashes with the same index. The next index to test in these cases will be different, contrary to the previous algorithms.
- **Deleted node is not removed, but marked as "removed"**.
- Hash must **never** equal 0.

Lets say, $\text{Hash1}(\text{key}) = \text{key} \% 13$

$\text{Hash2}(\text{key}) = 7 - (\text{key} \% 7)$

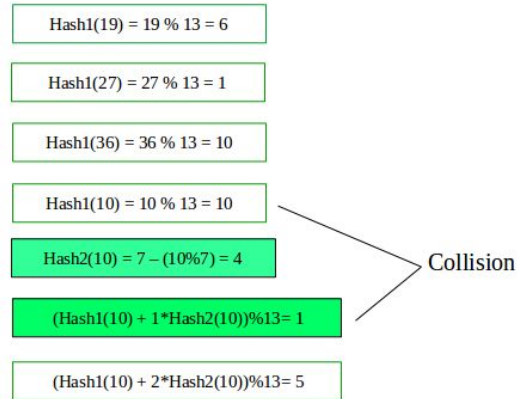


Fig 6. Quadratic Probing

Collision Resolution - Separate Chaining

- If two elements have the same hash, we store both in a **bucket**.
- **Harder** to implement but **easy** to verify;
- Collisions are handled **easily** and thus, the insert operation works the **fastest** in this algorithm.
- **Deleted node can be deleted.**
- Uses **more** space as there are additional data structures.
- **Open hashing**, which means caching is harder, thus leading to more cache misses.

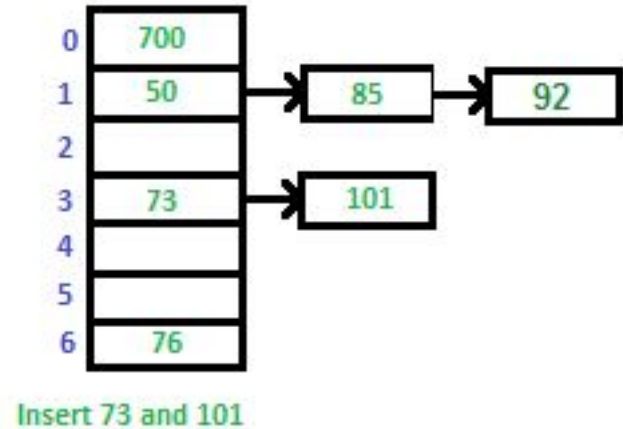


Fig7. [Separate Chaining](#)

HashMap Resizing

- When a certain **threshold** of occupied indexes is reached, the map needs to be **resized**.
- This leads to **fewer** hash collisions since there are more possible indexes to store the values.

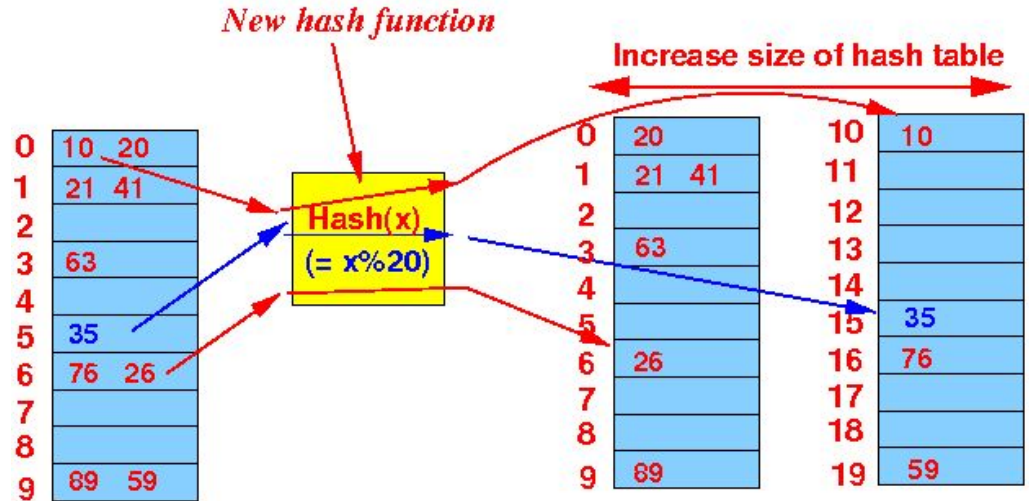


Fig 8. Resizing HashMap

HashMap Resizing

- After **resize**, the previously stored values need to be relocated.
- This can lead to a **high overhead** since we need to calculate the hash of the values all over again.
- To mitigate this problem, **the hash of each element is stored** alongside it, so it can be reutilized when calculating the new index position

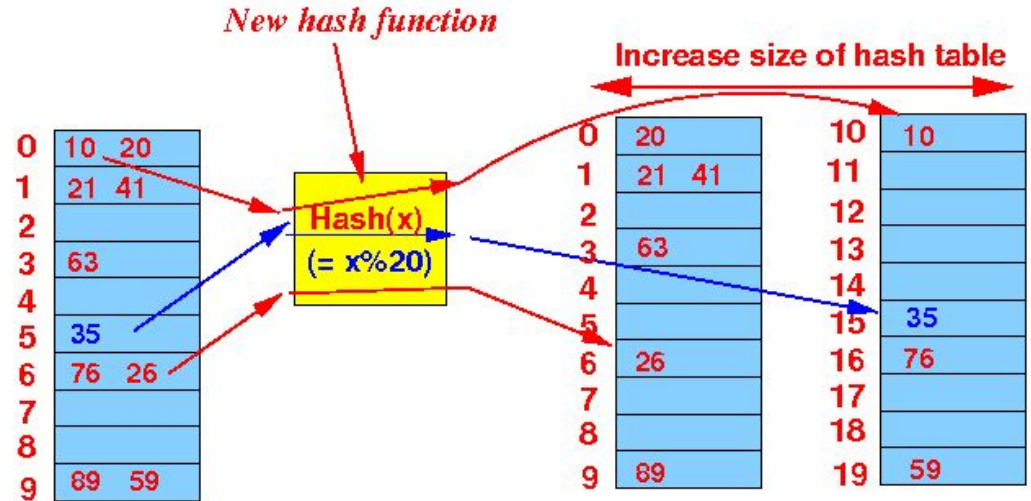


Fig 9. [Resizing HashMap](#)

HashMaps in Our Project

- Our **Hashmaps** will be used instead of STD hashmaps:
 - The STD hashmaps were made to be as generic as possible, and this can lead to some overhead in our project (the merge of sets take about 60% of raycast's execution time).
- The initial size can be **calculated** to reduce the number of resizings done:
 - By using the number of raycasts we want to perform, we can estimate the number of elements each hashmap will have.
- It is **uncertain** which collision resolution technique will prove to be the best
 - They will be measured and compared.

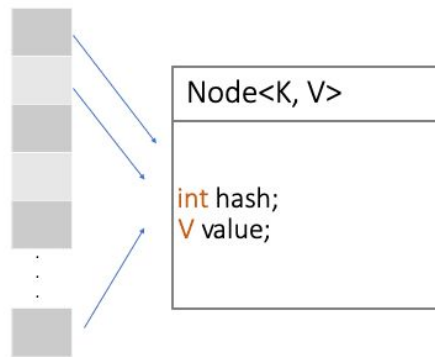


Fig 10. Example of a table entry in our implementation



Questions?

