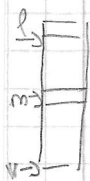


Ex

int besch (int a, int b, int c)



int l, v, m;

$$0 = 0$$
$$U' = N - 1/k$$

while $1 \leq u \leq 1$

$$m = (v + w) / 2$$
$$\text{nil} \{ a \in W \mid a = u = w \};$$

eise 3

if $(a \cdot u) > r$

then $v = m - 1$

```
else C = m + 1
```

```
if (a == 0) return 0
```

else return -1

$$K_{\text{residue}} = v - l$$

unverändert: β_0 ist nicht
 im 2022y statt β_0 ~~ist~~ ev
 und für β_0 $\beta_0 = 0$

Pos-condição: retorna 1 se não existir
se existe retorna 2 posição

Meior custo: estar no meio $T(N) \Rightarrow \infty$

Pior caso: No. currentes: $T(N) = 2 + 1 + \sum_{i=1}^{\log_2 N} \binom{5}{i} + 2$

$$= 3 + \log_2(N^4) +$$

$$O(\log(N))$$

8)

int repetidos (int v[], int N)

3

```
int rep = 0;
```

ent i, j.

```
int i, j;  
for(i=0; i<N-1; i++) { rep, i++ }
```

$$\text{for } i = 1, \dots, N \text{ do}$$

if $(v \in U \Rightarrow v \in JJ) \text{ ref} = 1;$

return rep;

4. No melhor caso, $1^\circ = 2^\circ \Rightarrow T(N) = 12$
 \Rightarrow inicializa 2 \times 350 + configura 1 = 1200

For case: No. of repetitions: $2 + \sum_{i=0}^{n-2} (2 + (\sum_{j=i+1}^{N-1} 3) + 1 + 1) \cdot (i+1)$
 $= 4 + \sum_{i=0}^{n-2} 4 + 3(N-i-1) = 4 + 4 \sum_{i=0}^{n-2} (N-i-1) - 3 \sum_{i=0}^{n-2} i$