Projeto 1 – SRC

Departamento de Eletrónica, Telecomunicações e Informática

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Índice

Exercício 9	3
Exercício 9.1	
Exercício 9.2	
Exercício 9.3	6
Resultados	θ
Exercício 10	8
Resultados	11
Script	
Conclusão	15

Exercício 9

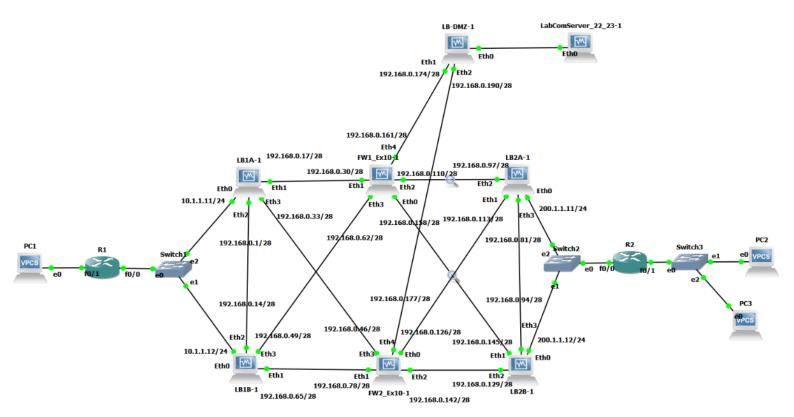


Figura 1-Rede

Neste exercício dividimos a rede 192.168.0.0/24 em várias subredes como mostra a figura 1.

Nas figuras 2 e 3 estão presentes as configurações para um Load Balancer e para uma firewall, uma vez que para os restantes Load Balancers e a outra firewall as configurações são semelhantes, pois apenas mudam endereços ip e algumas rotas estáticas. De salientar que na configuração dos Load Balancers, apesar de nós termos configurado rotas estáticas para as firewalls, não era necessário por causa do comando set load-balancing wan.

EXERCICIO 9: -----ROUTER INSIDE----conf term ip route 0.0.0.0 0.0.0.0 10.1.1.12 ip route 0.0.0.0 0.0.0.0 10.1.1.11 end write -----ROUTER OUTSIDE----conf term ip route 192.1.0.0 255.255.254.0 200.1.1.11 ip route 192.1.0.0 255.255.254.0 200.1.1.12 ip route 10.0.0.0 255.0.0.0 200.1.1.11 end write -----LB1B-----!Forçar o ping do inside a ir por baixo configure set protocols static route 10.2.2.0/24 next-hop 10.1.1.10 set protocols static route 0.0.0.0/0 next-hop 192.168.0.62 !N é preciso set protocols static route 0.0.0.0/0 next-hop 192.168.0.78 !N é preciso commit save set load-balancing wan interface-health eth1 nexthop 192.168.0.78 set load-balancing wan interface-health eth3 nexthop 192.168.0.62 set load-balancing wan rule 1 inbound-interface eth0 set load-balancing wan rule 1 interface eth1 weight 1 set load-balancing wan rule 1 interface eth3 weight 1 set load-balancing wan sticky-connections inbound set load-balancing wan disable-source-nat commit save set high-availability vrrp group FWCluster vrid 10 set high-availability vrrp group FWCluster interface eth2 set high-availability vrrp group FWCluster virtual-address 192.168.100.1/24 set high-availability vrrp sync-group FWCluster member FWCluster set high-availability vrrp group FWCluster rfc3768-compatibility commit Save set service conntrack-sync accept-protocol 'tcp,udp,icmp' set service conntrack-sync failover-mechanism vrrp sync-group FWCluster set service conntrack-sync interface eth2

Figura 2-Configurações routers e LB1B

set service conntrack-sync mcast-group 225.0.0.50 set service conntrack-sync disable-external-cache

commit save

```
-----FW1-----
configure
set protocols static route 0.0.0.0/0 next-hop 192.168.0.97
set protocols static route 10.2.2.0/24 next-hop 192.168.0.17
set protocols static route 10.2.2.0/24 next-hop 192.168.0.49 !
commit
save
set nat source rule 20 outbound-interface eth0
set nat source rule 20 source address 10.0.0.0/8 !
set nat source rule 20 translation address 192.1.0.1-192.1.0.10 !
set nat source rule 10 outbound-interface eth2
set nat source rule 10 source address 10.0.0.0/8 !192.168.0.0/8
set nat source rule 10 translation address 192.1.0.1-192.1.0.10
commit
save
set zone-policy zone INSIDE description "Inside (Internal Network)"
set zone-policy zone INSIDE interface eth1
set zone-policy zone INSIDE interface eth3
set zone-policy zone OUTSIDE description "Outside (Internet)"
set zone-policy zone OUTSIDE interface eth0
set zone-policy zone OUTSIDE interface eth2
set firewall name FROM-INSIDE-TO-OUTSIDE rule 10 action accept
set firewall name FROM-INSIDE-TO-OUTSIDE rule 10 protocol udp
set firewall name FROM-INSIDE-TO-OUTSIDE rule 10 destination port 5000-6000
set firewall name FROM-OUTSIDE-TO-INSIDE rule 10 action accept
set firewall name FROM-OUTSIDE-TO-INSIDE rule 10 state established enable
set zone-policy zone INSIDE from OUTSIDE firewall name FROM-OUTSIDE-TO-INSIDE
set zone-policy zone OUTSIDE from INSIDE firewall name FROM-INSIDE-TO-OUTSIDE
commit
save
```

Figura 3-Configurações FW1

Exercício 9.1

Através dos comandos set *load-balancing wan sticky-connections inbound* e *set load-balancing wan disable-source-nat* permite que não haja a necessidade de as Firewalls terem sincronização. Com o comando *set load-balancing wan sticky-connections* as interfaces vão memorizar de onde o tráfego foi recebido permitindo que o tráfego entre dispositivos seja sempre enviado pela mesma ligação e, com o comando *set load-balancing wan disable-source-nat*, o endereço ip de origem dos pacotes de saída será preservado quando forem enviados pela interface WAN de maneira a garantir que os pacotes sejam roteados corretamente de volta ao remetente. Por exemplo, se o request do ping do pc1 for por LB1A-FW1-LB2A até ao pc2, o mesmo, por causa dos comandos anteriores, vai saber por onde tem de enviar o reply para o pc1.

Exercício 9.2

O algoritmo IP Hash pode permitir a inexistência de sincronização em Load Balancers porque, neste algoritmo uma hash function é usada para mapear cada request para um servidor específico. A hash function assegura que o mesmo request é enviado para o mesmo servidor, assim, cada load balancer, através do endereço ip do cliente, pode calcular, independentemente, o servidor/firewall para qual o request dever ser enviado, eliminando a necessidade de sincronização.

Exercício 9.3

Um ataque DDOS é quando o atacante tenta sobrecarregar um servidor/sistema com um elevado número de pedidos/pacotes tornando esse mesmo sistema inválido ou extremamente lento devido à sobrecarga de recursos. Os dispositivos que tenham a sincronização ativa estão constantemente a trocar informação para se atualizarem. Se houver um ataque DDOS nesses dispositivos de sincronização, a rede pode tornar-se muito mais lenta, pelo número elevado de pacotes a serem encaminhados nas interfaces de sincronização, por exemplo, se houver um ataque DDOS no LB1A, como o LB1A e o LB1B estão sincronizados, esse ataque ia se propagar para o LB1B, ou seja, ia afetar ainda mais a rede levando a um efeito em cascada. Posto isto, a sincronização pode ser prejudicial durante um ataque de DDOS.

Resultados

Testámos os pings do inside para o outside e verificámos que o ping foi bemsucedido como se pode ver na figura 4.

Como se pode ver no print em baixo, com as capturas do wireshark (na figura 1 com a imagem da rede dá para ver onde estão a ser feitas estas capturas), o ping passou com sucesso nas firewalls e ainda conseguimos ver a translação dos ip´s na tabela NAT (o endereço ip de origem 10.2.2.100 transformou-se em 192.1.0.3).

```
PC1> ping 200.2.2.100 -P 17 -p 5001
84 bytes from 200.2.2.100 udp_seq=1 ttl=57 time=35.577 ms
84 bytes from 200.2.2.100 udp_seq=2 ttl=57 time=36.640 ms
84 bytes from 200.2.2.100 udp_seq=3 ttl=57 time=36.703 ms
84 bytes from 200.2.2.100 udp_seq=4 ttl=57 time=36.092 ms
84 bytes from 200.2.2.100 udp_seq=5 ttl=57 time=36.123 ms
```

Figura 4-Ping inside para Outside

Time	Source	Destination	Protocc L	engt Info			No.	Time	Source	Destination	Protocc L	.engt Info		
238 453.02688	3 192.1.0.3	200.2.2.100	UDP	98 5449 → 5001 Len=56				246 447.794534	200.2.2.100	192.1.0.3	UDP	98 5001 → 5449 Len=56		
239 454.064504	4 192.1.0.3	200.2.2.100	UDP	98 5449 → 5001 Len=56				247 448.832304	1 200.2.2.100	192.1.0.3	UDP	98 5001 → 5449 Len=56		
240 455.10305	5 192.1.0.3	200.2.2.100	UDP	98 5449 → 5001 Len=56				248 449.870724	200.2.2.100	192.1.0.3	UDP	98 5001 → 5449 Len=56		
241 455.324898	8 192.168.0.97	192.168.0.110	ICMP	74 Echo (ping) request	id=0x0754,	seq=256/1.		249 450.322248	192.168.0.145	192.168.0.158	ICMP	74 Echo (ping) request	id=0x0767, seq=	256/1.
242 455.325114	4 192.168.0.110	192.168.0.97	ICMP	74 Echo (ping) reply	id=0x0754,	seq=256/1.	ш	250 450.322493	192.168.0.158	192.168.0.145	ICMP	74 Echo (ping) reply	id=0x0767, seq=	256/1.
243 460.328172	2 192.168.0.97	192.168.0.110	ICMP	74 Echo (ping) request	id=0x0754,	seq=256/1.	ш	251 455.326059	192.168.0.145	192.168.0.158	ICMP	74 Echo (ping) request	id=0x0767, seq=	256/1.
244 460.328414	4 192.168.0.110	192.168.0.97	ICMP	74 Echo (ping) reply	id=0x0754,	seq=256/1.		252 455.326295	192.168.0.158	192.168.0.145	ICMP	74 Echo (ping) reply	id=0x0767, seq=	256/1
245 463.511466	6 192.1.0.3	200.2.2.100	UDP	98 8725 → 5001 Len=56				253 458.279144	200.2.2.100	192.1.0.3	UDP	98 5001 → 8725 Len=56		
246 464.548668	8 192.1.0.3	200.2.2.100	UDP	98 8725 → 5001 Len=56				254 459.316389	200.2.2.100	192.1.0.3	UDP	98 5001 → 8725 Len=56		
247 465.33170	7 192.168.0.97	192.168.0.110	ICMP	74 Echo (ping) request	id=0x0754,	seq=256/1.		255 460.329592	192.168.0.145	192.168.0.158	ICMP	74 Echo (ping) request	id=0x0767, seq=	256/1.
248 465.33191	3 192.168.0.110	192.168.0.97	ICMP	74 Echo (ping) reply	id=0x0754,	seq=256/1.		256 460.329813	192.168.0.158	192.168.0.145	ICMP	74 Echo (ping) reply	id=0x0767, seq=	256/1
249 465.585633	3 192.1.0.3	200.2.2.100	UDP	98 8725 → 5001 Len=56				257 460.353512	200.2.2.100	192.1.0.3	UDP	98 5001 → 8725 Len=56		
250 466.622510	6 192.1.0.3	200.2.2.100	UDP	98 8725 → 5001 Len=56				258 460.442544	PcsCompu_80:3d	PcsCompu_c0:84	ARP	60 Who has 192.168.0.14	5? Tell 192.168.	0.158
251 467.659902	2 192.1.0.3	200.2.2.100	UDP	98 8725 → 5001 Len=56				259 460.442833	PcsCompu_c0:84	PcsCompu_80:3d	ARP	60 192.168.0.145 is at	08:00:27:c0:84:a	4
252 470.335100	0 192.168.0.97	192.168.0.110	ICMP	74 Echo (ping) request	id=0x0754,	seq=256/1.		260 461.390145	200.2.2.100	192.1.0.3	UDP	98 5001 → 8725 Len=56		
253 470.335306	6 192.168.0.110	192.168.0.97	ICMP	74 Echo (ping) reply	id=0x0754,	seq=256/1.		261 462.427509	200.2.2.100	192.1.0.3	UDP	98 5001 → 8725 Len=56		
254 475.339043	1 192.168.0.97	192.168.0.110	ICMP	74 Echo (ping) request	id=0x0754,	seq=256/1.		262 465.333514	192.168.0.145	192.168.0.158	ICMP	74 Echo (ping) request	id=0x0767, seq=	256/1.
255 475.33926	5 192.168.0.110	192.168.0.97	ICMP	74 Echo (ping) reply	id=0x0754,	seq=256/1.		263 465.333732	192.168.0.158	192.168.0.145	ICMP	74 Echo (ping) reply	id=0x0767, seq=	256/1.
256 480.34276	7 192.168.0.97	192.168.0.110	ICMP	74 Echo (ping) request	id=0x0754,	seq=256/1.		264 470.336972	192.168.0.145	192.168.0.158	ICMP	74 Echo (ping) request	id=0x0767, seq=	256/1.
257 480.342994	4 192.168.0.110	192.168.0.97	ICMP	74 Echo (ping) reply	id=0x0754,	seq=256/1.		265 470.337196	192.168.0.158	192.168.0.145	ICMP	74 Echo (ping) reply	id=0x0767, seq=	256/1.
258 489 54372	4 PcsCompu 77:8d	PcsCompu 97.8f	ARP	60 Who has 192 168 0 97	? Tell 192	168 0 110		266 475 340523	192 168 0 145	192 168 0 158	TCMP	74 Echo (ning) request	id=0x0767 sea=	-256/1

Figura 5-Captura Wireshark

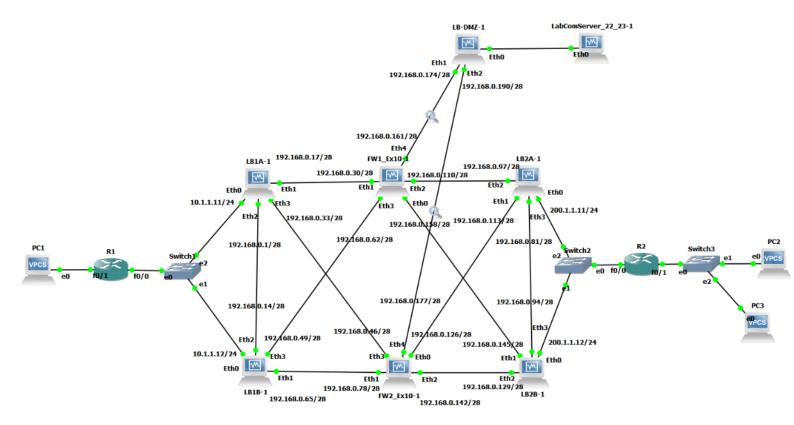


Figura 6-Rede2

Para este exercício, foram adicionadas novas configurações nas firewalls, uma nova VM, na zona DMZ, e foi adicionado também o PC3 que irá simular o endereço ip do atacante.

Em primeiro lugar, criámos uma nova zona(DMZ) e adicionámos em cada uma das firewalls uma rota estática para a rede da DMZ. Em seguida, definimos um conjunto de regras que permite os dispositivos da zona *INSIDE* comunicarem com a DMZ por TCP/UDP nos portos 22(SSH),443(HTTPS),21(FTP) e 53(DNS) e que permitem aos dispositivos da zona *OUTSIDE* comunicarem com a DMZ também por TCP/UDP, mas com acesso restrito, ou seja, neste caso, apenas o tráfego com destino nos portos de HTTPS (443) e DNS(53) é permitido.

Por fim, foram adicionadas configurações para dar drop/bloquear os pacotes com base no endereço de ip de origem(ip do atacante). De salientar que, de maneira a verificar sempre em primeiro lugar se o ip de origem é um ip identificado como sendo um atacante(dentro do range 200.2.2.20 – 200.2.2.30), definimos um menor sequence number para esta regra(rule 5).

Na figura abaixo estão presentes as configurações de uma das firewalls(FW1) e do Load Balancer que liga as firewalls à DMZ.

```
configure
set load-balancing wan interface-health eth1 nexthop 192.168.0.161
set load-balancing wan interface-health eth2 nexthop 192.168.0.177
set load-balancing wan rule 1 inbound-interface eth0
set load-balancing wan rule 1 interface eth1 weight 1
set load-balancing wan rule 1 interface eth2 weight 1
set load-balancing wan sticky-connections inbound
set load-balancing wan disable-source-nat
commit
save
```

Figura 7- Configurações Load Balancer

```
-----FW1-----
configure
set zone-policy zone DMZ description "DMZ Zone"
set zone-policy zone DMZ interface eth4
commit
save
set protocols static route 192.1.1.0/24 next-hop 192.168.0.174
save
set firewall name FROM-INSIDE-TO-DMZ rule 10 action accept
set firewall name FROM-INSIDE-TO-DMZ rule 10 protocol udp
set firewall name FROM-INSIDE-TO-DMZ rule 10 destination port 22,443,21,53
set firewall name FROM-OUTSIDE-TO-DMZ rule 10 action accept
set firewall name FROM-OUTSIDE-TO-DMZ rule 10 destination port https,domain
set firewall name FROM-OUTSIDE-TO-DMZ rule 10 protocol udp
set firewall name FROM-DMZ-TO-INSIDE rule 10 action accept
set firewall name FROM-DMZ-TO-INSIDE rule 10 state established enable
set firewall name FROM-DMZ-TO-OUTSIDE rule 10 action accept
set firewall name FROM-DMZ-TO-OUTSIDE rule 10 state established enable
set firewall name FROM-INSIDE-TO-DMZ rule 11 action accept
set firewall name FROM-INSIDE-TO-DMZ rule 11 protocol tcp
set firewall name FROM-INSIDE-TO-DMZ rule 11 destination port 22,443,21,53
set firewall name FROM-OUTSIDE-TO-DMZ rule 11 action accept
set firewall name FROM-OUTSIDE-TO-DMZ rule 11 destination port https,domain
set firewall name FROM-OUTSIDE-TO-DMZ rule 11 protocol tcp
set firewall name FROM-OUTSIDE-TO-DMZ rule 5 action drop
set firewall group address-group attackers address '200.2.2.20-200.2.2.30'
set firewall name FROM-OUTSIDE-TO-DMZ rule 5 source group address-group attackers
set zone-policy zone DMZ from OUTSIDE firewall name FROM-OUTSIDE-TO-DMZ
set zone-policy zone DMZ from INSIDE firewall name FROM-INSIDE-TO-DMZ
set zone-policy zone INSIDE from DMZ firewall name FROM-DMZ-TO-INSIDE
set zone-policy zone OUTSIDE from DMZ firewall name FROM-DMZ-TO-OUTSIDE
commit
save
```

Figura 8-Configurações FW1

Resultados

Como era suposto, o ping do atacante para a DMZ foi bloqueado pela firewall, porque o ip do atacante é 200.2.2.22 e como definimos como regra na firewall que os ip's de origem no range 200.2.2.20 – 200.2.2.30 iam ser bloqueados, então o ip do atacante é bloqueado. Nas capturas do wireshark(na figura 6, com a imagem da rede, dá para ver onde estão a ser feitas estas capturas), figura 9, dá para ver que nenhum pacote passou pela firewall.

```
PC3> ping 192.1.1.100 -P 6 -p 53
Connect 53@192.1.1.100 timeout
```

Figura 9-Ping Atacante

Time	Source	Destination	Protocc I	Lengi Info	No.	Time	Source	Destination	Protoct L	engt Info
8_ 3617.495	1 192.168.0.177	192.168.0.190	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1	1 2	22_ 3607.4882	192.168.0.161	192.168.0.174	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1.
8_ 3617.664	4 PcsCompu_50:38	PcsCompu_40:8f_	ARP	60 Who has 192.168.0.190? Tell 192.168.0.177	1 2	22 3612.4915	192.168.0.174	192.168.0.161	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1
8_ 3617.664	7 PcsCompu_40:8f	PcsCompu_50:38_	ARP	60 192.168.0.190 is at 08:00:27:40:8f:26	1 3	22 3612.4919	192.168.0.161	192.168.0.174	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1.
8_ 3622.498	5 192.168.0.190	192.168.0.177	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1		22 3612.5670	PcsCompu_60:51_	PcsCompu_74:aa	ARP	60 Who has 192.168.0.161? Tell 192.168.0.174
8_ 3622.498	7 192.168.0.177	192.168.0.190	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1		22_ 3612.5672	PcsCompu_74:aa_	PcsCompu_60:51	ARP	60 192.168.0.161 is at 08:00:27:74:aa:65
8 3627.501	9 192.168.0.190	192.168.0.177	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1	1 2	22 3617.4952	192.168.0.174	192.168.0.161	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1
8_ 3627.502	2 192.168.0.177	192.168.0.190	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1	1 3	22 3617.4954	192.168.0.161	192.168.0.174	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1
8 3632.505	5 192.168.0.190	192.168.0.177	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1	1 2	22 3617.6502	PcsCompu_74:aa_	PcsCompu_60:51	ARP	60 Who has 192.168.0.174? Tell 192.168.0.161
8 3632.505	3 192.168.0.177	192.168.0.190	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1		22 3617.6504	PcsCompu_60:51_	PcsCompu_74:aa	ARP	60 192.168.0.174 is at 08:00:27:60:51:cb
8_ 3632.674	L PcsCompu_40:8f	PcsCompu_50:38	ARP	60 Who has 192.168.0.177? Tell 192.168.0.190	1 2	22 3622.4986	192.168.0.174	192.168.0.161	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1
8 3632.674	3 PcsCompu_50:38	PcsCompu_40:8f	ARP	60 192.168.0.177 is at 08:00:27:50:38:55	1 3	22 3622.4989	192.168.0.161	192.168.0.174	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1
8_ 3637.509	L 192.168.0.190	192.168.0.177	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1	1 2	22 3627.5022	192.168.0.174	192.168.0.161	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1
8_3637.509	3 192.168.0.177	192.168.0.190	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1	1 3	22 3627.5025	192.168.0.161	192.168.0.174	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1
8_ 3642.512	L 192.168.0.190	192.168.0.177	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1	1 3	22 3632.5058	192.168.0.174	192.168.0.161	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1.
8_3642.512	3 192.168.0.177	192.168.0.190	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1	1 3	22 3632.5060	192.168.0.161	192.168.0.174	ICMP	74 Echo (ping) reply id=θxθ6b3, seq=256/1.
8_ 3642.752	5 PcsCompu_50:38	PcsCompu_40:8f_	ARP	60 Who has 192.168.0.190? Tell 192.168.0.177	1 3	22 3637.5088	192.168.0.174	192.168.0.161	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1
8_ 3642.752	9 PcsCompu_40:8f_	PcsCompu_50:38_	ARP	60 192.168.0.190 is at 08:00:27:40:8f:26	1 2	22_ 3637.5090	192.168.0.161	192.168.0.174	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1
8_ 3647.515	7 192.168.0.190	192.168.0.177	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1	1 3	22 3642.5124	192.168.0.174	192.168.0.161	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1
8_ 3647.516	0 192.168.0.177	192.168.0.190	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1	1 3	22 3642.5127	192.168.0.161	192.168.0.174	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1
8_ 3652.519	L., 192.168.0.190	192.168.0.177	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1	1 3	22 3647.5158	192.168.0.174	192.168.0.161	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1
8_ 3652.519	3., 192,168,0,177	192,168,0,190	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1	1 3	22 3647.5160	192.168.0.161	192.168.0.174	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1.

Figura 10-Wireshark Ping Atacante

Como é possível ver nas figuras abaixo, o ping do pc na zona *INSIDE* para a DMZ foi um sucesso. Caso o ping tivesse como destino uma outra porta que não as definidas nas regras, o ping seria bloqueado.

```
C1> ping 192.1.1.100 -P 6 -p 53
         53@192.1.1.100 seq=1 ttl=60 time=12.333
endData 53@192.1.1.100 seq=1 ttl=60 time=13.359
         53@192.1.1.100 seq=1 ttl=60 time=13.355
         53@192.1.1.100 seq=2 ttl=60 time=38.033
endData 53@192.1.1.100 seq=2 ttl
         53@192.1.1.100 seq=2 ttl
         53@192.1.1.100
endData 53@192.1.1.100
         53@192.1.1.100
         53@192.1.1.100
endData 53@192.1.1.100
         53@192.1.1.100
Connect
         53@192.1.1.100
endData 53@192.1.1.100 seq=5 ttl=60 time=18.536
         53@192.1.1.100 seq=5 ttl=60
```

Figura 11-Ping Inside

	Time	Source	Destination	Protocc	Lengt Info	No.	Time	Source	Destination	Protocc	Lengt Info
	21 4204.2019	10.2.2.100	192.1.1.100	TCP	74 17163 → 53 [SYN] Seq=0 Win=2920 Len=0 MSS	1	28 4202.6214.	. 10.2.2.100	192.1.1.100	TCP	66 [TCP Keep-Alive] 17163 → 53 [ACK] Seq=57
	21 4204.2026	192.1.1.100	10.2.2.100	TCP	74 53 → 17163 [SYN, ACK] Seq=0 Ack=1 Win=651	1 2	28 4202.6219.	. 192.1.1.100	10.2.2.100	TCP	66 [TCP Keep-Alive ACK] 53 → 17163 [ACK] Seq
	21 4204.2225	10.2.2.100	192.1.1.100	TCP	66 17163 → 53 [ACK] Seq=1 Ack=1 Win=5840 Len	2	28 4202.8279.	. 192.1.1.100	10.2.2.100	TCP	66 [TCP Retransmission] 53 → 17163 [FIN, ACK
	21 4204.2226	10.2.2.100	192.1.1.100	TCP	122 17163 → 53 [PSH, ACK] Seq=1 Ack=1 Win=584	1 2	28 4202.9120.	. 192.168.0.174	192.168.0.161	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1
	21 4204.2230	192.1.1.100	10.2.2.100	TCP	66 53 → 17163 [ACK] Seq=1 Ack=57 Win=65152 L	2	28 4202.9122.	. 192.168.0.161	192.168.0.174	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1
	21 4204.2430	10.2.2.100	192.1.1.100	TCP	66 17163 → 53 [FIN, PSH, ACK] Seq=57 Ack=1 W	2	28 4203.3044.	. 192.1.1.100	10.2.2.100	TCP	66 [TCP Retransmission] 53 → 17163 [FIN, ACK
	21 4204.2436	192.1.1.100	10.2.2.100	TCP	66 53 → 17163 [FIN, ACK] Seq=1 Ack=58 Win=65	1					74 [TCP Port numbers reused] 17163 → 53 [SYN
	21 4204.2636	10.2.2.100	192.1.1.100	TCP	66 [TCP Keep-Alive] 17163 → 53 [ACK] Seq=57	2	28 4203.6399.	. 192.1.1.100	10.2.2.100	TCP	74 53 → 17163 [SYN, ACK] Seq=0 Ack=1 Win=651
	21 4204.2641	192.1.1.100	10.2.2.100	TCP	66 [TCP Keep-Alive ACK] 53 → 17163 [ACK] Seq	2	28 4203.6598.	. 10.2.2.100	192.1.1.100	TCP	66 17163 → 53 [ACK] Seq=1 Ack=1 Win=5840 Len
	22 4204.4716		10.2.2.100		66 [TCP Retransmission] 53 → 17163 [FIN, ACK	1 2	28 4203 . 6803.	. 10.2.2.100	192.1.1.100	TCP	122 17163 → 53 [PSH, ACK] Seq=1 Ack=1 Win=584
	22 4204.9474				66 [TCP Retransmission] 53 → 17163 [FIN, ACK	1 2	28 4203.6809.	. 192.1.1.100	10.2.2.100	TCP	66 53 → 17163 [ACK] Seq=1 Ack=57 Win=65152 L
					74 [TCP Port numbers reused] 17163 → 53 [SYN	2	28 4203.7215.	. 10.2.2.100	192.1.1.100	TCP	66 17163 → 53 [FIN, PSH, ACK] Seq=57 Ack=1 W
-	22 4205.2615	192.1.1.100	10.2.2.100	TCP	66 53 → 17163 [ACK] Seq=1 Ack=2871193468 Win	2	28 4203 . 7222.	. 192.1.1.100	10.2.2.100	TCP	66 53 → 17163 [FIN, ACK] Seq=1 Ack=58 Win=65
	22 4205.2812	10.2.2.100	192.1.1.100	TCP	66 17163 → 53 [RST, ACK] Seq=2871193468 Ack=	1	28 4203.7421.	. 10.2.2.100	192.1.1.100	TCP	66 [TCP Keep-Alive] 17163 → 53 [ACK] Seq=57
-	22 4207.9153	192.168.0.190	192.168.0.177	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1	2	28 4203 . 7427.	. 192.1.1.100	10.2.2.100	TCP	66 [TCP Keep-Alive ACK] 53 → 17163 [ACK] Seq
	22 4207.9156	192.168.0.177	192.168.0.190	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1	1					66 [TCP Retransmission] 53 → 17163 [FIN, ACK
	22 4208.6220				74 [TCP Port numbers reused] 17163 → 53 [SYN	2					66 [TCP Retransmission] 53 → 17163 [FIN, ACK
	22 4208 . 6227				66 [TCP ACKed unseen segment] 53 → 17163 [AC	1 2					66 [TCP Retransmission] 53 → 17163 [FIN, ACK
	22 4208.6425	10.2.2.100	192.1.1.100	TCP	66 17163 → 53 [RST, ACK] Seq=735693199 Ack=1		28 4207.1762.	. 192.1.1.100	10.2.2.100	TCP	66 [TCP Retransmission] 53 → 17163 [FIN, ACK
-	22 4212.9185	192.168.0.190	192.168.0.177	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1	1 2	28 4207.9152.	. 192.168.0.174	192.168.0.161	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1
	22 4212 9187	192 168 0 177	192 168 0 190	TCMP	74 Echo (ning) renly id=0x06h3 seq=256/1	1 3	28 4207 9154	192 168 0 161	192 168 0 174	TCMP	74 Echo (ning) renly id=0x06h3 seq=256/1

Figura 12-Wireshark Ping Inside

Por fim, testámos o ping do *OUTSIDE* para a DMZ e verificamos que também teve resposta, tal como é possível ver nas figuras abaixo.

```
PC2> ping 192.1.1.100 -P 6 -p 53

Connect 53@192.1.1.100 seq=1 ttl=60 time=20.552 ms

SendData 53@192.1.1.100 seq=1 ttl=60 time=18.543 ms

Close 53@192.1.1.100 seq=1 ttl=60 time=18.507 ms

Connect 53@192.1.1.100 seq=2 ttl=60 time=38.047 ms

SendData 53@192.1.1.100 seq=2 ttl=60 time=18.518 ms

Close 53@192.1.1.100 seq=2 ttl=60 time=19.559 ms

Connect 53@192.1.1.100 seq=3 ttl=60 time=38.028 ms

SendData 53@192.1.1.100 seq=3 ttl=60 time=18.522 ms

Close 53@192.1.1.100 seq=3 ttl=60 time=19.553 ms

Connect 53@192.1.1.100 seq=4 ttl=60 time=38.017 ms

SendData 53@192.1.1.100 seq=4 ttl=60 time=18.495 ms

Close 53@192.1.1.100 seq=4 ttl=60 time=19.528 ms

Connect 53@192.1.1.100 seq=5 ttl=60 time=36.997 ms

SendData 53@192.1.1.100 seq=5 ttl=60 time=19.526 ms

Close 53@192.1.1.100 seq=5 ttl=60 time=19.526 ms

Close 53@192.1.1.100 seq=5 ttl=60 time=18.524 ms
```

Figura 13-Ping Outside

	Time	Source	Destination	Protocc	Lengt	Info					No
19	3717.5666	192.168.0.177	192.168.0.190	ICMP	74	Echo	(ping)	reply	id=0x06b3,	seq=256/1	
19	3722.5697	192.168.0.190	192.168.0.177	ICMP	74	Echo	(ping)	request	id=0x06b3,	seq=256/1	
19	3722.5699	192.168.0.177	192.168.0.190	ICMP	74	Echo	(ping)	reply	id=0x06b3,	seq=256/1	
19	3722.7860	PcsCompu_40:8f	PcsCompu_50:38	ARP	60	Who h	nas 192.	168.0.17	7? Tell 192.	168.0.190	
19	3722.7862	PcsCompu_50:38	PcsCompu_40:8f	ARP	60	192.1	168.0.1	77 is at (08:00:27:50:	38:55	
19	3727.5732	192.168.0.190	192.168.0.177	ICMP	74	Echo	(ping)	request	id=0x06b3,	seq=256/1	
19	3727.5735	192.168.0.177	192.168.0.190	ICMP	74	Echo	(ping)	reply	id=0x06b3,	seq=256/1	
19	3732.5769	192.168.0.190	192.168.0.177	ICMP	74	Echo	(ping)	request	id=0x06b3,	seq=256/1	
19	3732.5771	192.168.0.177	192.168.0.190	ICMP	74	Echo	(ping)	reply	id=0x06b3,	seq=256/1	
19	3737.5806	192.168.0.190	192.168.0.177	ICMP	74	Echo	(ping)	request	id=0x06b3,	seq=256/1	
19	3737.5809	192.168.0.177	192.168.0.190	ICMP	74	Echo	(ping)	reply	id=0x06b3,		
19	3742.5839	192.168.0.190	192.168.0.177	ICMP	74	Echo	(ping)	request	id=0x06b3,	seq=256/1	
19	3742.5841	192.168.0.177	192.168.0.190	ICMP	74	Echo	(ping)	reply	id=0x06b3,	seq=256/1	
19	3742.5926	PcsCompu_50:38	PcsCompu_40:8f	ARP	60	Who h	nas 192.	.168.0.190	0? Tell 192.	168.0.177	
19	3742.5929	PcsCompu_40:8f	PcsCompu_50:38	ARP	60	192.1	168.0.19	90 is at (08:00:27:40:	8f:26	
19	3747.5874	192.168.0.190	192.168.0.177	ICMP	74	Echo	(ping)	request	id=0x06b3,	seq=256/1	
19	3747.5877	192.168.0.177	192.168.0.190	ICMP	74	Echo	(ping)	reply	id=0x06b3,	seq=256/1	
19	3752.5910	192.168.0.190	192.168.0.177	ICMP	74	Echo	(ping)	request	id=0x06b3,	seq=256/1	
19	3752.5912	192.168.0.177	192.168.0.190	ICMP	74	Echo	(ping)	reply	id=0x06b3,	seq=256/1	
19	3757.5945	192.168.0.190	192.168.0.177	ICMP	74	Echo	(ping)	request	id=0x06b3,	seq=256/1	ı I
19	3757.5948	192.168.0.177	192.168.0.190	ICMP	74	Echo	(ping)	reply	id=0x06b3.	sea=256/1	1

	No.	Time	Source	Destination	Protocc	Lengt Info
		23 3746.0369	200.2.2.100	192.1.1.100	TCP	66 17163 → 53 [RST, ACK] Seq=974470464 Ack=1
		23 3746.0369	200.2.2.100	192.1.1.100	TCP	74 [TCP Port numbers reused] 17163 → 53 [SYN
		23 3746.0374	192.1.1.100	200.2.2.100	TCP	74 53 → 17163 [SYN, ACK] Seq=0 Ack=1 Win=651
		23 3746.0575	200.2.2.100	192.1.1.100	TCP	66 17163 → 53 [ACK] Seq=1 Ack=1 Win=5840 Len
		24 3746.0782	200.2.2.100	192.1.1.100	TCP	122 17163 → 53 [PSH, ACK] Seq=1 Ack=1 Win=584
		24 3746.0787	192.1.1.100	200.2.2.100	TCP	66 53 → 17163 [ACK] Seq=1 Ack=57 Win=65152 L
		24 3746.1193	200.2.2.100	192.1.1.100	TCP	66 17163 → 53 [FIN, PSH, ACK] Seq=57 Ack=1 W
		24 3746.1201	192.1.1.100	200.2.2.100	TCP	66 53 → 17163 [FIN, ACK] Seq=1 Ack=58 Win=65
		24 3746.1399	200.2.2.100	192.1.1.100	TCP	66 [TCP Keep-Alive] 17163 → 53 [ACK] Seq=57
		24 3746.1405	192.1.1.100	200.2.2.100	TCP	66 [TCP Keep-Alive ACK] 53 → 17163 [ACK] Seq
		24 3746.3483	192.1.1.100	200.2.2.100	TCP	66 [TCP Retransmission] 53 → 17163 [FIN, ACK
		24 3746.7109	PcsCompu_60:51	PcsCompu_74:aa	ARP	60 Who has 192.168.0.161? Tell 192.168.0.174
		24 3746.7111	PcsCompu_74:aa	PcsCompu_60:51	ARP	60 192.168.0.161 is at 08:00:27:74:aa:65
•		24 3746.8242	192.1.1.100	200.2.2.100	TCP	66 [TCP Retransmission] 53 → 17163 [FIN, ACK
		24 3747.5877	192.168.0.174	192.168.0.161	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1
		24 3747.5880	192.168.0.161	192.168.0.174	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1
		24 3747.7519	192.1.1.100	200.2.2.100	TCP	66 [TCP Retransmission] 53 → 17163 [FIN, ACK
					TCP	66 [TCP Retransmission] 53 → 17163 [FIN, ACK
		24 3752.5912	192.168.0.174	192.168.0.161	ICMP	74 Echo (ping) request id=0x06b3, seq=256/1
		24 3752.5915	192.168.0.161	192.168.0.174	ICMP	74 Echo (ping) reply id=0x06b3, seq=256/1
"		24 3753.2562	192.1.1.100	200.2.2.100	TCP	66 [TCP Retransmission] 53 → 17163 [FIN. ACK

Figura 14- Wireshark Ping Outside

Este script tem como objetivo criar automaticamente as regras de bloqueio da firewall após a identificação do endereço de ip dos atacantes.

Foi definido um range de endereços para simular a lista com ip´s dos atacantes. Em seguida, através do comando da linha do *tcpdump ...* capturamos os pacotes na interface enp0s3 filtrando os pacotes que têm o endereço ip no range definido pelas variáveis \$ip_start e \$ip_end\$ e parar após capturar 1000 pacotes. Ainda na mesma linha, o *while read attacker_ip* vai iterar sobre cada linha e vai atribuir o valor à variável *attacker_ip*.

Posto isto, se pacotes forem capturados nesta interface através do comando *tcpdump* significa que houve pacotes cujo ip origem estava na lista dos endereços ip´s dos atacantes e iniciamos uma conexão remota via ssh para as duas firewall para bloquear pacotes com base no endereço ip de origem(attacker_ip).

Infelizmente tentámos testar este script, mas não conseguimos chegar ao resultado pretendido, mas pelo menos fica uma ideia daquilo que pretendíamos fazer no script.

```
#!/bin/bash
# Set the firewall IP address
firewall1_ip_address="firewall1_ip_address"
firewall2_ip_address="firewall2_ip_address"
# Set the IP range to block
ip start="200.2.2.20"
ip_end="200.2.2.30"
echo "Searching for DDOS attacks"
# Start capturing network traffic with tcpdump and filter it based on source IP address
tcpdump -n -i \ enp0s3 \ src \ net \ p_start/\ p_end -c \ 1000 \ | \ grep -oE \ b([0-9]\{1,3\}.) \{3\}[0-9]\{1,3\}\ | \ sort -u \ | \ while \ read \ attacker_ip; \ attacker_ip
            # Connect to the firewall1 via SSH
            ssh firewall1_ip_address << EOF
                       configure
                  set firewall name BLOCK rule 10 action drop
                             set firewall name BLOCK rule 10 source address $attacker_ip
                       commit
                       save
                       exit
FOF
            # Connect to the firewall2 via SSH
           ssh firewall2_ip_address << EOF
                       configure
                 set firewall name BLOCK rule 10 action drop
                            set firewall name BLOCK rule 10 source address $attacker_ip
                       commit
                       save
                       exit
EOF
            # Print a message to confirm that the blocking rule has been added
            echo "Blocking rule added to the firewalls to block traffic from IP address $attacker_ip"
done
```

Figura 15- Script

Conclusão

Concluindo, achamos que conseguimos atingir todos os objetivos deste projeto com sucesso. Salientando que, apesar de não termos mostrado os resultados dos pings em todas as portas das regras, os mesmos tiveram, tal como os outros, sucesso.