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Groups: Busic Definitions
Def: Let G be a set. A binary operation on Gr is a function:
*: G x G - D G
 For psychological reasons, notationally we write *(a,b) as a *b, for all a,b & G.

The pair (6,*) is called a binary structure.
Def: A group is a set Go, together with a binary operation * on G, such that the following hold:
           1. (Associativity): (a*b) *c = a*(b*c) taibic & G.

2. (Existence of identity): Je&G s.t. 9*e=e*a=a tae G.

3. (Existence of inverses): For all a & G., Jb&G s.t. a*b=b*a=e.
Del: A binary operation & on a set G is called commutative, if Hoube G, axb = b*a. A group (G,*) is called commutative, if Hoube G, axb = b*a.
                                                                                 (or Abelian)
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Examples: Binary operation,
+: ZxZ ->Z
+(ab)=a+b

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Notes (Romarks: 1. In Botje's notes there is a more complete treatment of binary structures. Due to time we are forusing/skipping struct to graps, 2. A Warning: In Boltje's notes (Zn, tn) is used to denote the go we are colling (Z/nZ1+). (Zn, +n) is slightly different in that Boltje defines Zn as the set of standard representatives of the equivalence classes ZnZ, and defines at the = remainder of orth when divided by n This way he avoids having to talk whent equivalence classes. 3. Notation/Example: Let X be a set. F(X, X) = functions from X to X Sym(x) = bijective functions from X to X

If $X = \{1,2,3,...,n\}$, Sym(X) is denoted by Sym(n) or Symn and called the symmetric group. Composition of functions, o, is the binary operation we will usually consider on the sets $F(X_1X_1)$ and $Sym(X_1)$.

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an juvija. Yes	a. Va= 1 Wat pred matrix has	(a,, u1, -, an) + (-a,,-u2,-n-an) = Yes	$a \cdot \alpha = 1$ $a \cdot \alpha = 1$ $a \cdot \alpha = 1$	the all a except ato	70 at at a	1585]
40}	Yes No inv.	No, no in.	yes Yes		Jes winv.	Resurp? Why Not?

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(X,X)7 蒙蒙 Not every function No has an invesse Yes how invesses Yes

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Group- new methomatical object, set with extra structure. We define functions between

Def: Let (G, *) and (H, 0) be two groups. A homomorphism, f, from G to H, is a function f: G - H, such that for all xiy E 6, $f(x*A) = f(x) \cdot f(A)$

A bijective homomorphism is called an isomorphism. [Att) (6,*) and (H.o) are isomorphic if there exists an isomorphism between them. Their is denoted (6,*) \cong (H.o) or $G\cong H$ if the binary operations are clear.

Remarks: 1. For all yps (6,*), the identity function ide is a yp hom. from 2. Two gps being isom, means intuitively that they are the "yame" gp but just viewed from different perspectives.

Examples: M. The Exponential f: (R+) - (Rro..), offined by $f(x) = e^x$ is a gP hom become $f(x) \in \mathbb{R}$.

 $f(x+y) = e^{x+y} = e^{x} \cdot e^{y} = f(x) \cdot f(y)$

is an iromorphism becomes f hus an inverse given by $q(x) = \log x$

2. The function TB: (R,+) -> (R,+) given by T(xy)=x is a gp Now because for all (xy), (Z,w) = R, T((x,y)+(Z,w))=T((x+Z1,y+w)) = T(x1y) + T(Z,w) 3. If $A \in M_{n\times m}(\mathbb{R})$, $\Phi_A: (\mathbb{R}^m, +) \longrightarrow (\mathbb{R}^n, +)$, $\Phi_A(v) = Av$ is $\mathcal{A}P$ how. Prop: Let (6, *), (H.o), and (M, rd) be three groups. Let f: G-DH and g: H-DM be homomorphisms. Then the composition, gof: G-DM is a housemerphism. Modified xiy E.G. Then need to show gof (x*y) = (gof (x) 1] (gof (y). Have gof(x*y) = g(f(x*y)) $= g(f(x) \circ f(y)) \text{ because } f \mapsto g \text{ him.}$ $= g(f(x)) \prod g(f(y)) \text{ because } g \mapsto a \text{ hom.}$ $= (a px.) \prod (-a) r$ = (g.fx) [] (g.f)(y) [] Prop. Let (6, *) he a gp. Then the identity elevent of Go is unique. proof: Let e, e'é 6 be two identity els ments of G. The 4 e=e*e' and e'=e*e' by the identity property. e = e'.]

Prop: Let (G,*) he a yp. For all a EG, the inverse of a is any unique.

proof. Let b, c be town inverses of a, so

a*b=b*a=e

and

a*c=c*a=e

Then we need to show b=c. We have:

b=e*b by id. Prop.

=(c*a*b) by associativity

=c*(a*b) by associativity

=c*e by about

=c w id. prop.

Remark: By propositions, giver a gr (6.x) we can unambiguously write down the instity element. For an abstract gr (6.x), e is resumbly used to denote the identity element. For grr you already know, the used to denote may already have a symbol. We may also given as 6, identity element may already have a symbol. We may also given as 6, unambiguously write a for the inverse of a in Gr.

Prop (cancellation law): Let $a_1b_1e \in G_1$, whose $(G_1/*)$ is a gp. If a*c = a*b (resp c*a = b*a), then c = b.

Proof: If a*c = a*b, then $c = (\bar{a}'*a)*c = \bar{a}'*(a*c) = \bar{a}'*(a*b) = (\bar{a}'*a)*b = b$ Similarly for other one.

Prop: Let (G, *) be a group. Then for all a, b & G, $(a*b)^{-1} = b^{-1} * a^{-1}$ proof: We have $(a + b) * (b' * \hat{a}') = a * (b*b') * \hat{a}'$ = a * e * a 1 (6' + a') + (a+b) = 6' + (a' + a) * 6 Thought (axb) = 5 x a'. Propilet (6,+) he " gp. Ther e'=e. proof; e*e = e, 50 e=e. Remork/Notation: Writing or group additively, means the group operation is written with a +. In this case, O usually denotes the identity element and -a people of the invers of a. Otherwise, the group may be written with 1,0, **, etc., and I the identity is lore, along with Inverse painty of a being denoted by a!.

Therefore the following notation

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Additive group (6,+), a & G, n & Z $Nd = \begin{cases} a_1 a_1 - a_1 & \text{n-times} & \text{if} & \text{n>0} \\ 0 & \text{if} & \text{n=0} \\ -a_1 - a_1 - a_1 - a_1 & \text{n+times} & \text{if} & \text{n} < 0 \end{cases}$ Mult. group (G1.), af G. $n \in \mathbb{Z}$ $\alpha^{n} = \begin{cases}
\alpha \cdot \alpha \cdot \dots \cdot \alpha & n + t \text{ ine}, & f \neq n > c
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\alpha \cdot \alpha \cdot \dots \cdot \alpha &$ With this notation, by definition, we have amon = and and (an) = amon factor, min & I. Prop: Lt (6.+), (HID) he two JPS, and let f: G - Abe a hom. (a) Let egas EG, exett dente the identity elements. Then. +(ea) = CH (b) For all af 6, f(ai) = f(ai). proof: (a) Let h & H. Then hafles) = hafles of fles) By carrellation law got $h = hof(e_6)$ Similarly h = f(eg) th. Though fleg1 = ex. $f(\tilde{a}') \prod f(\tilde{a}) = f(\tilde{a}' \star \tilde{a}) = f(e_{\tilde{a}}) = e_{H}$ f(a) 0 f(a") = f(a + a") = f(ea) = eH

Therefore $f(\tilde{a}') = f(\tilde{a})' = f(\tilde{a})'$

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