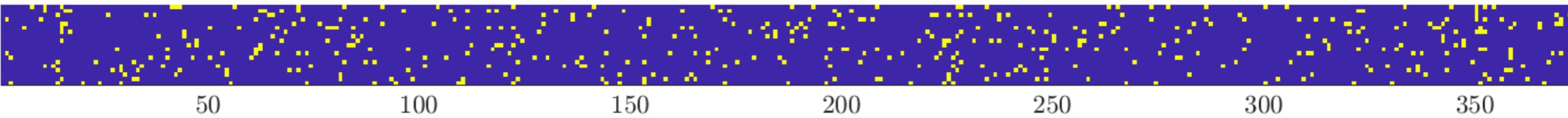


# Transfer entropy for network reconstruction in a simple dynamical model

Roy Goodman  
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# How I Spent My Sabbatical



NYU

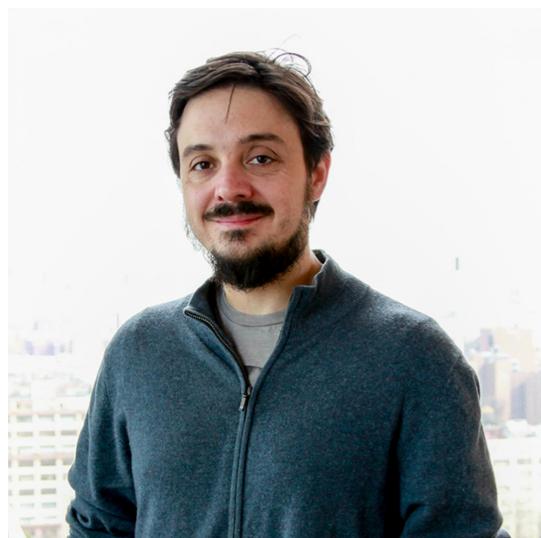
TANDON SCHOOL  
OF ENGINEERING



My location



My commute



My host: Mau Porfiri



<http://www.aimspress.com/journal/mine>

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*Research article*

**Topological features determining the error in the inference of networks using transfer entropy**

Roy H. Goodman<sup>1,\*</sup> and Maurizio Porfiri<sup>2,\*</sup>

# The Porfiri Lab

A big group working in a lot of areas, both theoretical and through laboratory experiments:

- Fluid mechanics: fluid structure interactions during water impact
- Artificial Muscles and Soft Robotics
- Telerehabilitation
- Network-based modeling of infectious diseases
- Fish schooling
- Using robotics and zebrafish to study substance-abuse disorders
- Information-theoretic analysis of social science datasets
- more...

A unifying theme in this work is using methods from information theory for modeling and analysis

# What is this talk about?

A lot of words to define here:

Transfer entropy for network reconstruction  
in a simple dynamical model

- Network: a *graph* composed of vertices and edges, the subject at the heart of *graph theory*
- Transfer entropy: a quantity describing the transfer of information from one evolving variable to another, from information theory
- The dynamical model: to be described, a simple probabilistic dynamics.
- Reconstruction: figure out properties of the graph based on the dynamics

18 months ago I knew  $\epsilon$  about any of this

# Graph Theory

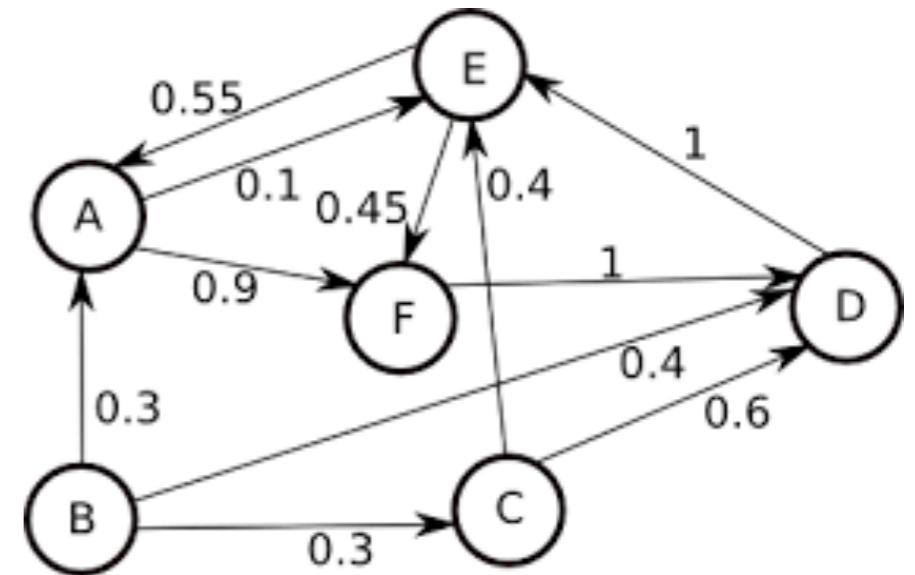
Graph theory is central to the mathematics of Computer Science, describing the connections between interacting agents.

A graph is defined by a set  $\mathcal{V}$  of vertices, connected by a set  $\mathcal{E}$  of edges. The graph at right is both *directed* and *weighted*.

The weight matrix has entries

$W_{ij}$  defined as follows

- If an edge exists from node  $j$  to node  $i$  the entry is the associated weight
- If no edge exists, the entry zero.



$$W = \begin{bmatrix} 0 & 0.3 & 0 & 0 & 0.55 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0.3 & 0 & 0 & 0 & 0 \\ 0 & 0.4 & 0.6 & 0 & 0 & 1 \\ 0.1 & 0 & 0.4 & 1 & 0 & 0 \\ 0.9 & 0 & 0 & 0 & 0.45 & 0 \end{bmatrix}$$

An important notion for us will be the *weighted incoming degree*  $\delta_i = \sum_j W_{ij}$ .

In this example  $\delta_E = 1 + 0.4 + 0.1 = 1.5$

# Information Theory

- Initiated by Claude Shannon's 1948 "A Mathematical Theory of Communication"
- Originally used to study the transmission of signals down noisy channels and to develop optimal strategies for encoding information
- Recently become popular tool for analyzing dynamical systems

Fundamental quantity:

Consider a discrete random variable  $X$  drawn from a sample space  $\mathcal{X}$

The *information* associated with the event  $X=x$  measures how "surprising" it is that  $X=x$

$$I(x) = -\log(\Pr(X = x))$$

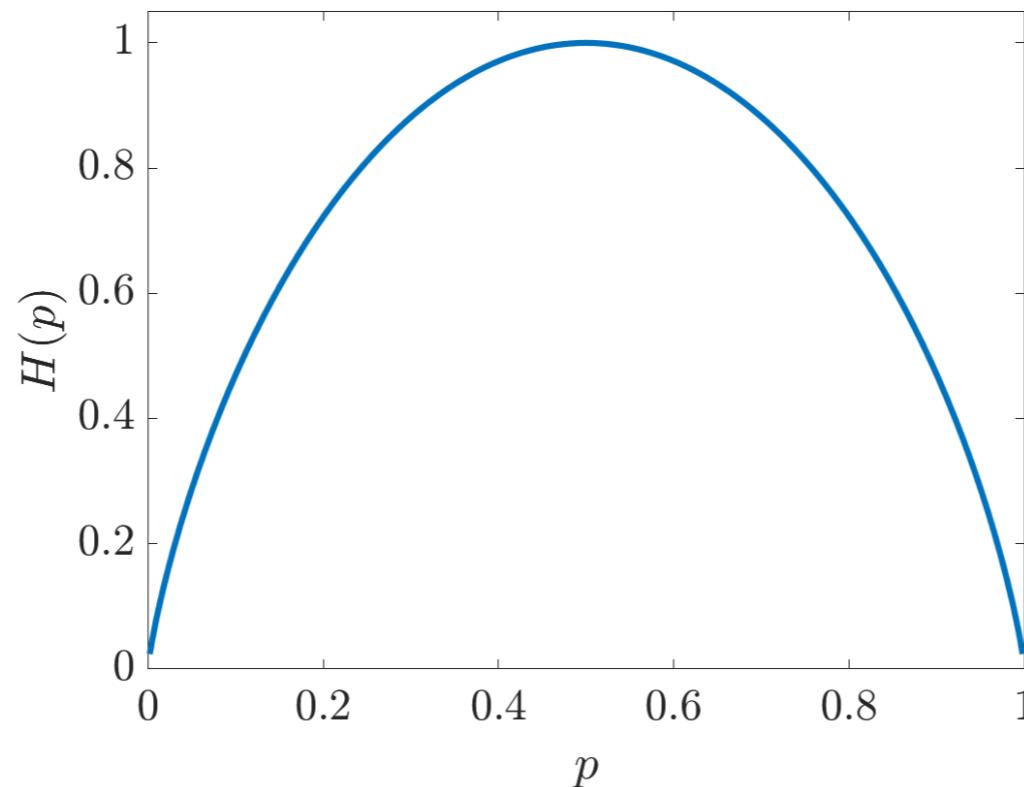
The *Shannon entropy* of the random variable is the expectation value of the information

$$H(X) = E[I(X)] = -\sum_{x \in \mathcal{X}} \Pr(X = x) \log \Pr(X = x)$$

# Basic example: biased coin toss

Consider a biased coin that gives heads with probability  $p$  and tails with probability  $1-p$

$$H(p) = -p \log p - (1 - p) \log (1 - p)$$



- When  $p \approx 0$  or  $p \approx 1$ , entropy small since surprising outcomes rarely occur
- When  $p \approx 0.5$ , entropy large since both outcomes equally likely

# Transfer entropy Schreiber 2000

Consider two random variables  $X$  and  $Y$ .

Define the *joint entropy* and the *conditional entropy*:

$$H(X, Y) = - \sum_{x \in \mathcal{X}, y \in \mathcal{Y}} \Pr(X = x, Y = y) \log \Pr(X = x, Y = y)$$

$$H(X|Y) = - \sum_{x \in \mathcal{X}, y \in \mathcal{Y}} \Pr(X = x, Y = y) \log \Pr(X = x|Y = y)$$

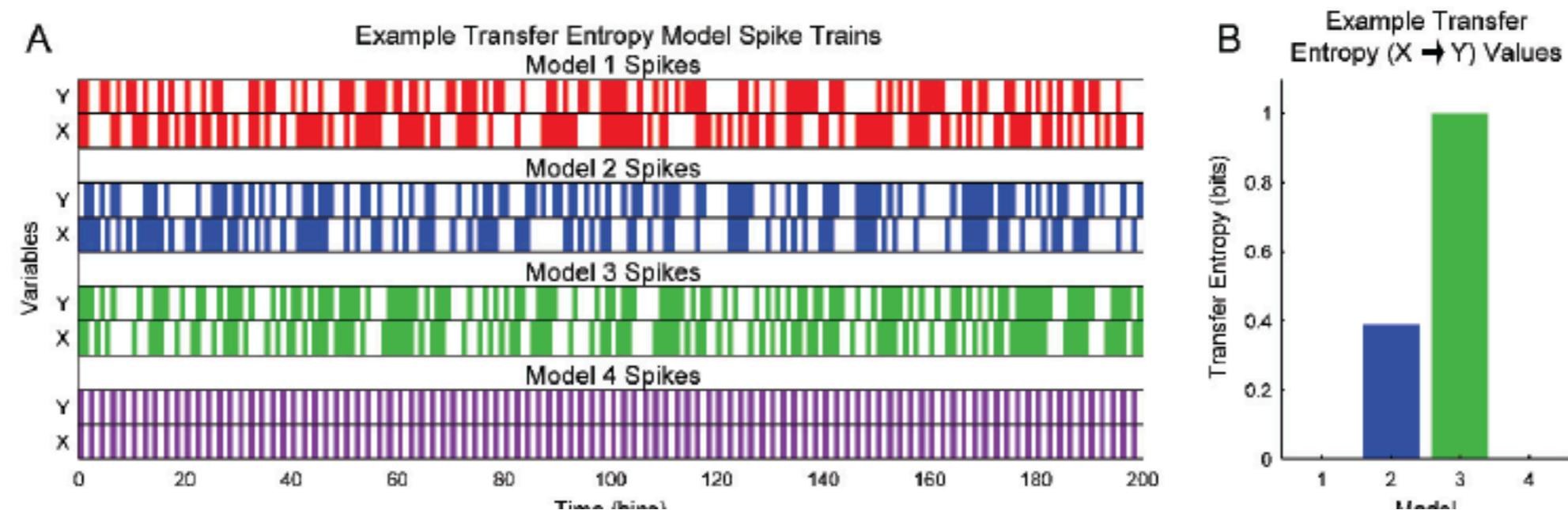
Entropy can be defined analogously for stationary stochastic processes  
*Transfer entropy* from  $Y$  to  $X$  is the difference between the entropy of  $X(t+1)$  conditioned on  $X(t)$  and that conditioned on both  $X(t)$  and  $Y(t)$

$$\begin{aligned} TE^{Y \rightarrow X} &= H(X(t+1)|X(t)) - H(X(t+1)|X(t), Y(t)) \\ &= \sum_{\substack{x_+ \in \mathcal{X} \\ x \in \mathcal{X} \\ y \in \mathcal{Y}}} \left\{ \Pr[X(t+1) = x_+, X(t) = x, Y(t) = y] \times \log \frac{\Pr[X(t+1) = x_+ | X(t) = x, Y(t) = y]}{\Pr[X(t+1) = x_+ | X(t) = x]} \right\} \end{aligned}$$

Transfer entropy measures the *reduction in the uncertainty* of predicting  $X(t + 1)$  from both  $X(t)$  and  $Y(t)$  relative to predicting it from  $X(t)$  alone.

# Contrived transfer entropy example

Source: A Tutorial for Information Theory in Neuroscience  
Nicholas M. Timme and Christopher Lapish 2018

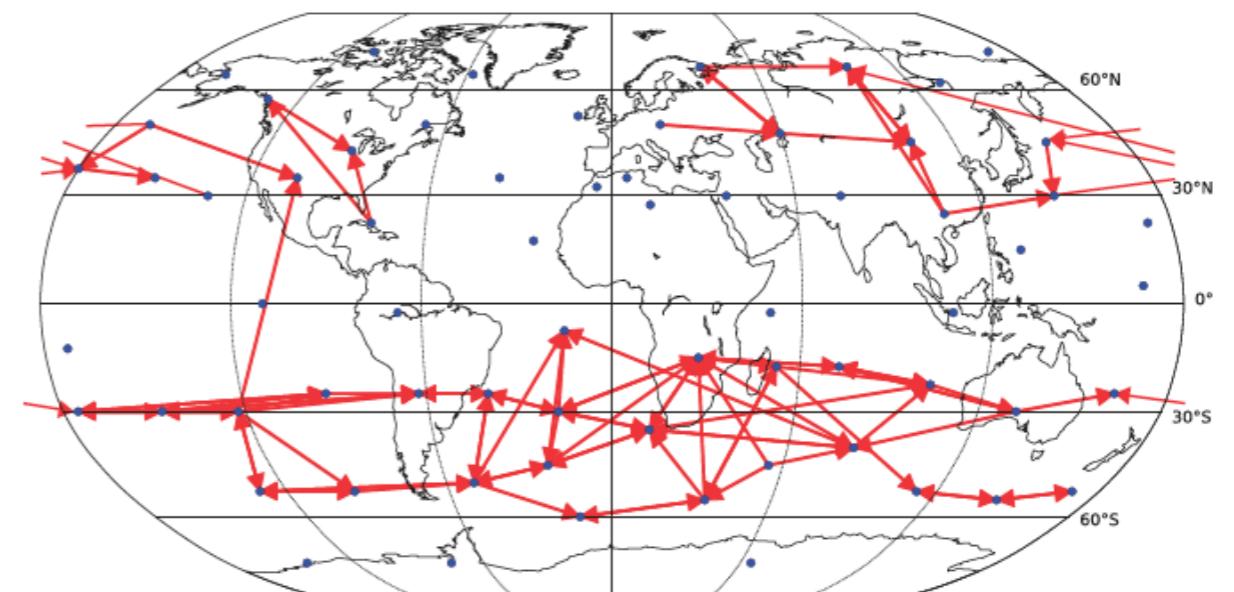
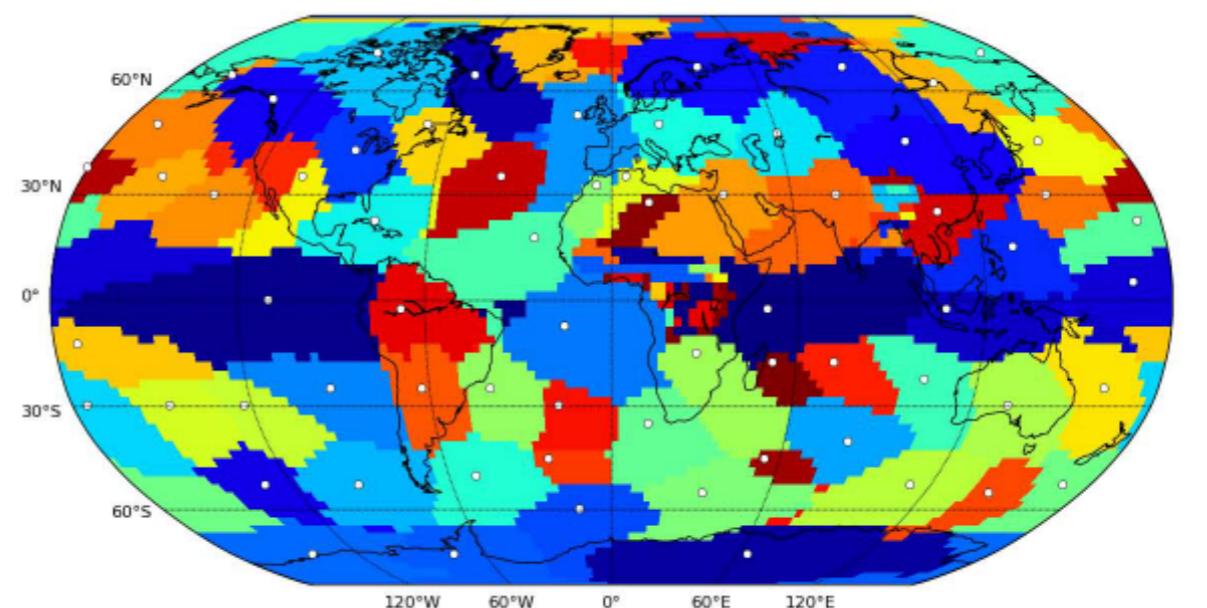


- Series 1:  $X$  &  $Y$  uncorrelated:  $TE \approx 0$
- Series 2 & 3:  $X$  tends to fire before  $Y$ :  $TE$  large
- Series 4:  $X(t+1)$  determined entirely from  $X(t)$ , no improvement from knowing  $Y(t)$ :  $TE = 0$

# Transfer entropy used to infer climate network

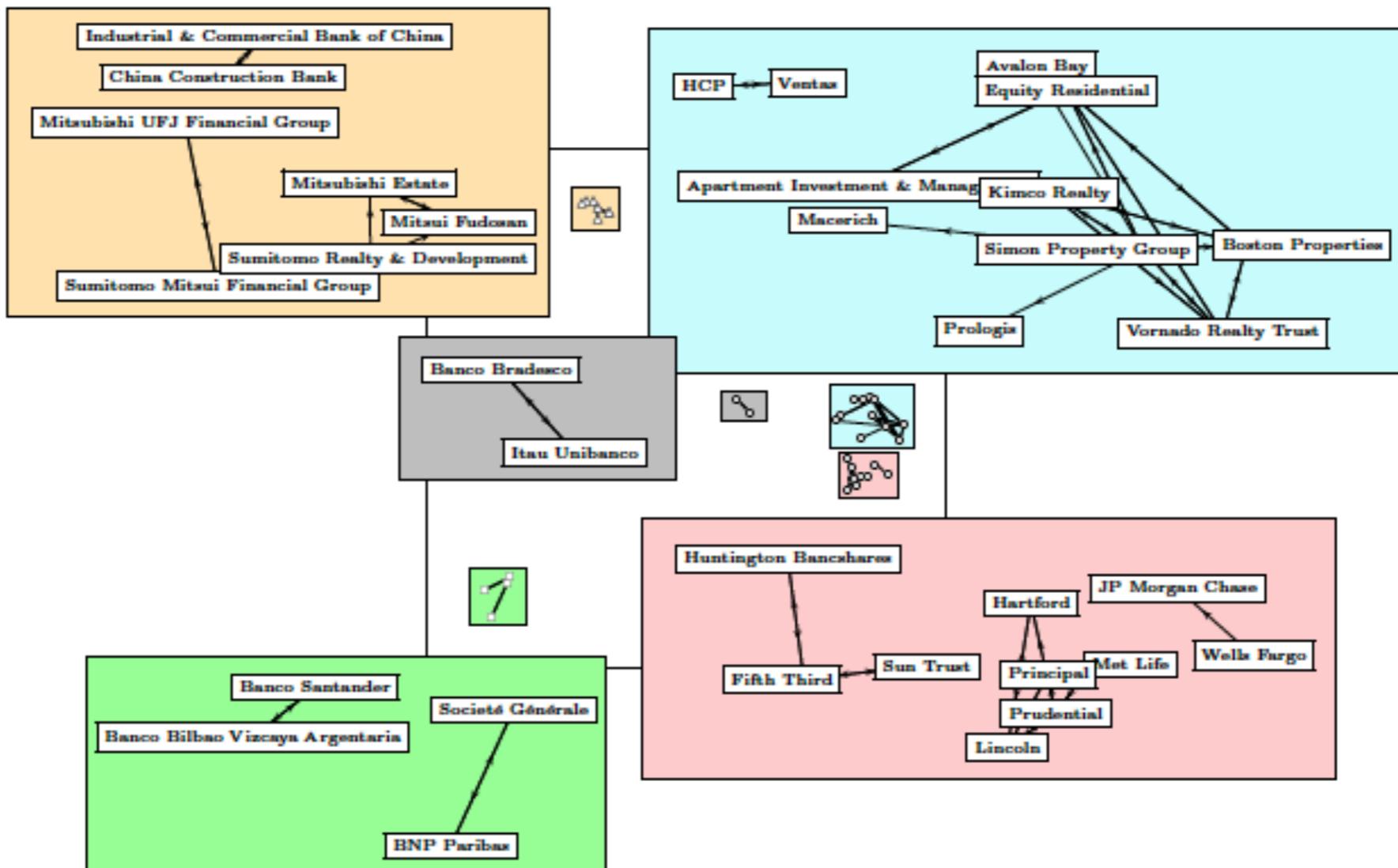
Hlinka et al. 2013

- Divide the earth into N patches using principal component analysis
- Look at time series of surface-area temperature deviations
- Compute transfer entropies, estimating PDFs using specially-tuned kernel density estimators
- Threshold to find links with non-negligible transfer entropy



# Transfer Entropy used to analyze connections between corporations

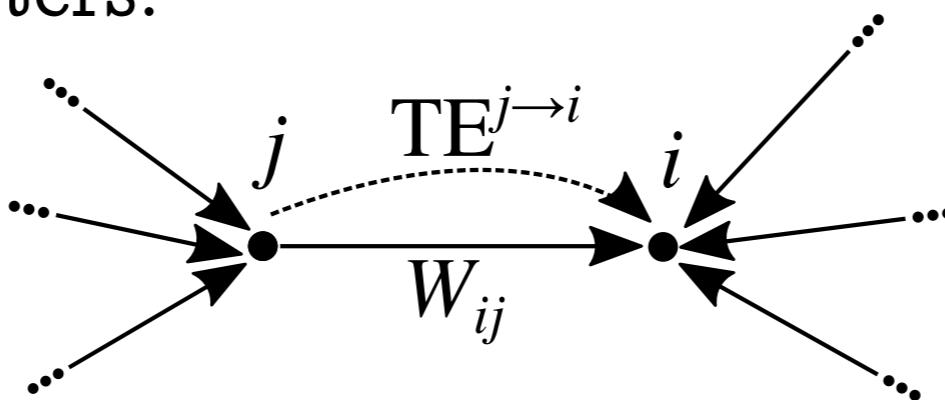
Sandoval 2014



Analysis based on time series of stock prices. Transfer entropy predicts that the price of a stock has an influence on other stocks in its sector and geographic area

# Advantages and disadvantages of Transfer Entropy

- Model free
- Nonlinear
- Relatively simple to compute
- Requires the estimation of underlying probability distribution, e.g. by binning or kernel density estimation
- Lots of hidden parameters to fiddle with that might effect the computation
- Inherently dyadic: quantifies the interaction between two agents while ignoring the effects of others. Unclear how much this matters.



# Random Boolean Network (RBN) Model

Porfiri & Ruiz Marín 2018

A system of nodes  $X^i(t)$ ,  $i=1\dots n$ , each of which can take values in  $\{0,1\}$

At time step  $t+1$ , the state of node  $i$  depends on the state of the system at time step  $t$ , according to

$$\Pr [X^i(t+1) = 1 | X^1(t) = x^1, \dots, X^N(t) = x^N] = \epsilon \left[ 1 + \sum_{j=1}^N W_{ij} x^j \right]$$

Conceived as a model “Policy Diffusion”—How the passage of laws in one jurisdiction influences the passage of laws in other jurisdictions.

The terms  $W_{ij} \geq 0$  represent the “network of influence.”



Sample time series

# Porfiri & Ruiz Marín's result (2018)

Setting  $\epsilon \ll 1$  allows the use of perturbation methods to approximately calculate Transfer Entropy in terms of the weights

Inverting gives an approximate formula for  $W_{ij}$  in terms of transfer entropy

$$TE^{j \rightarrow i} = \epsilon^2 G^{(2)}(W_{ij}) + \mathcal{O}(\epsilon^3)$$

where

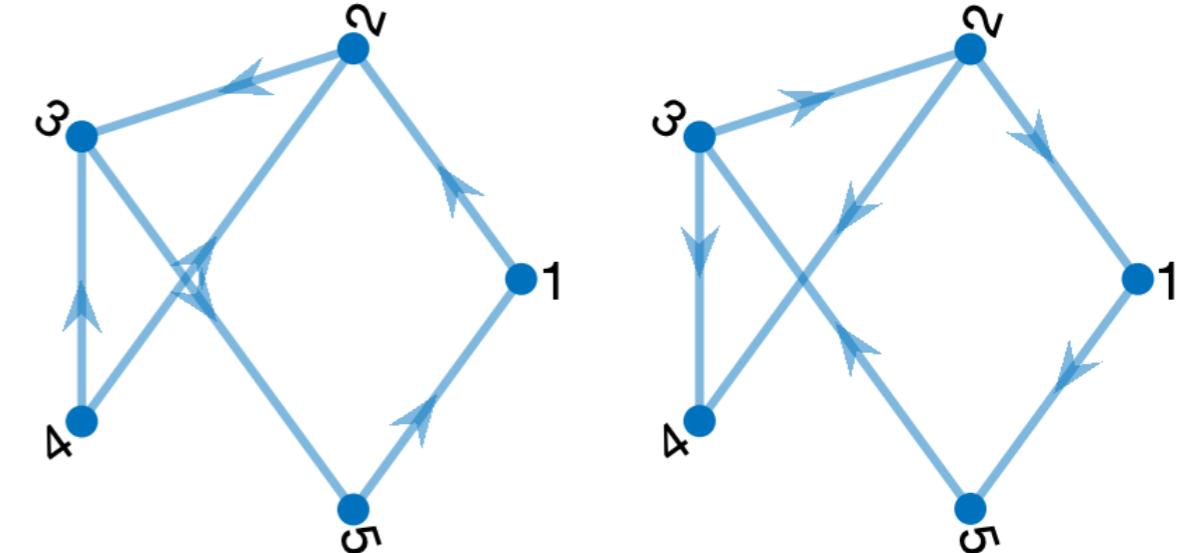
$$G^{(2)}(x) = -x + (1+x) \log(1+x) \approx \frac{1}{2}x^2 \text{ for } x \ll 1$$

**Takeaway:** To leading order transfer entropy from  $j$  to  $i$  depends on the strength of the weight from  $j$  to  $i$

**Obvious next question:** By calculating next term in the expansion, can we quantify the effect the global topology of the network has on the computed transfer entropy?

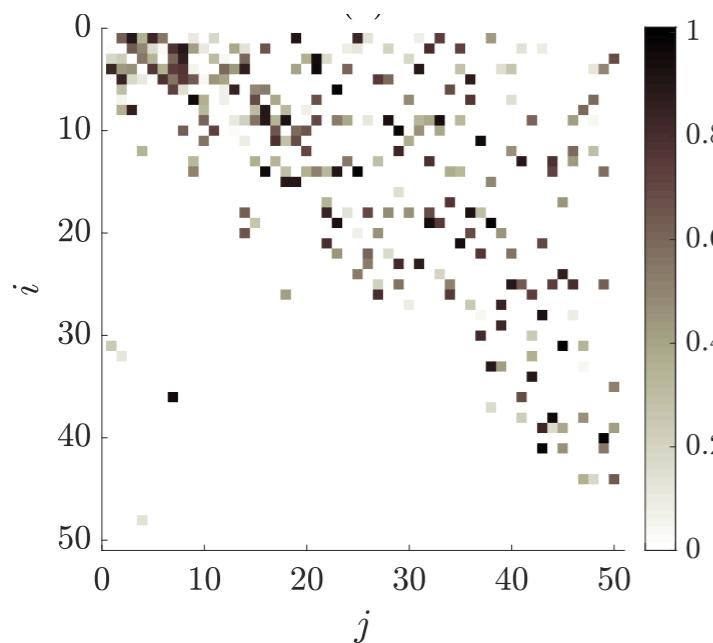
# An example to show that next-order terms matter

The transpose  $\Gamma^\top$  of a directed graph  $\Gamma$  has the same vertices, oppositely-directed edges, and a weight matrix  $W^\top$



Two networks  $\Gamma$  and  $\Gamma^\top$  that behave differently

A 50-vertex, 282-edge directed Barabási-Albert network with weight matrix  $W$ , random weights

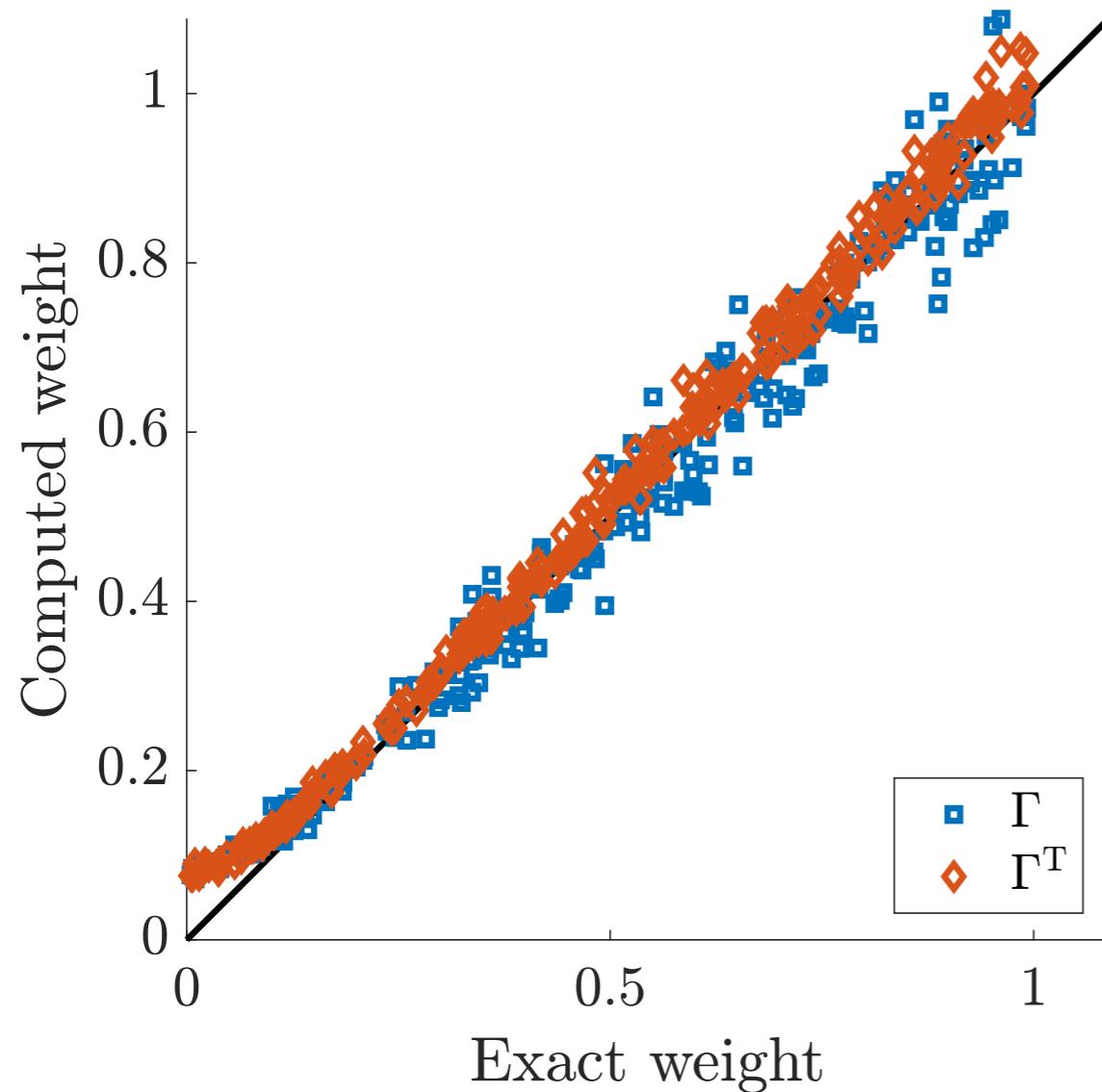


Procedure:

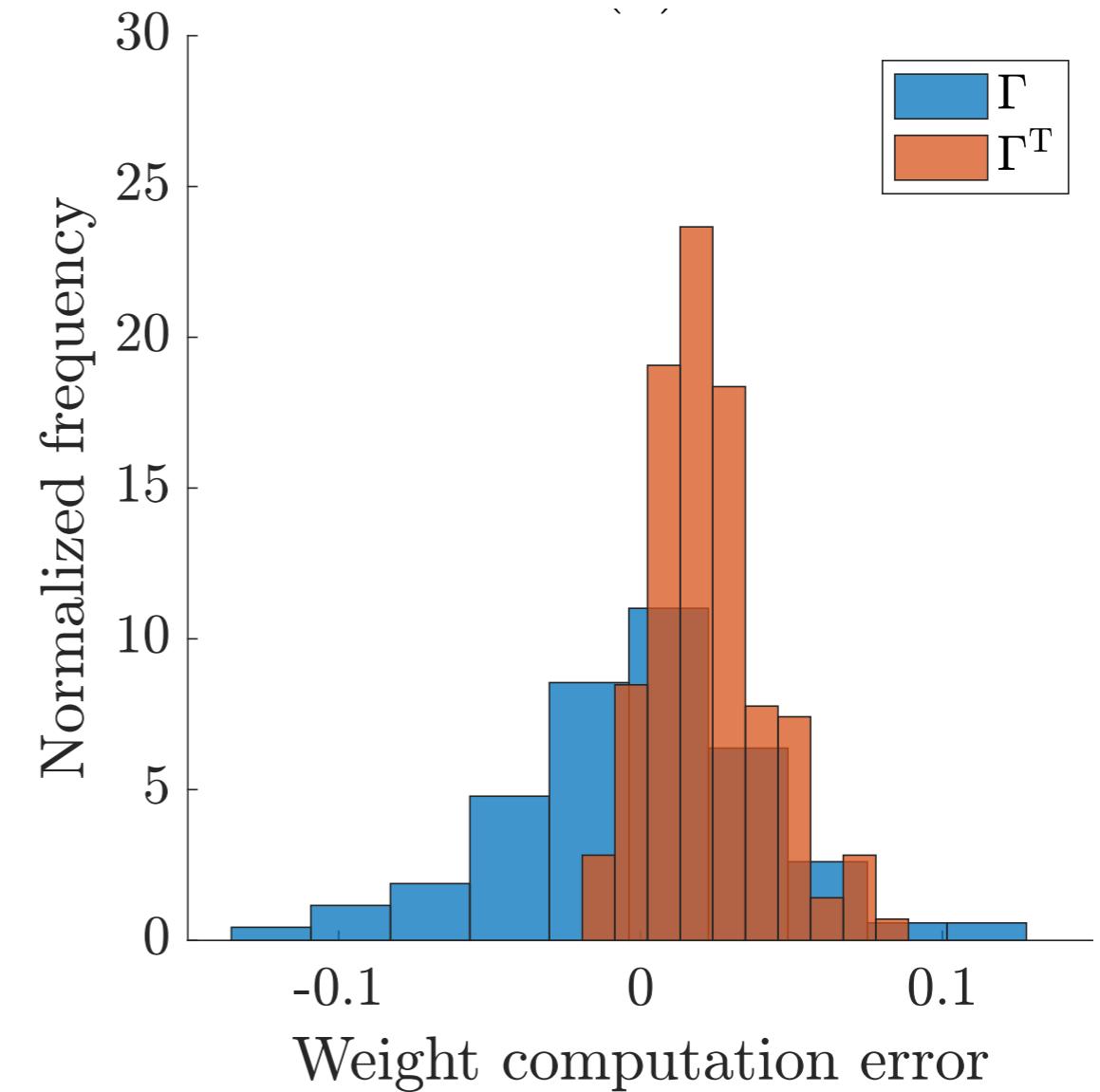
- Compute dynamics for  $10^5$  steps
- Repeat 100 times
- Compute TE from time series
- Estimate  $W_{ij}$  from formula

# Result: Distribution of error much wider for $\Gamma$ than $\Gamma^T$

Nonzero computed  
weights vs exact weights



Error in computed  
weights



# Strategy for analyzing the dynamics

- Recast the system as a Markov chain, dependent on a small parameter  $\epsilon$
- Calculate the stationary vector of the Markov chain via perturbation theory in  $\epsilon$
- Use the stationary probability vector and the transition law to derive a formula for transfer entropy in terms of the weights  $W_{ij}$
- Invert to get approximation for the weights in terms of the pairwise transfer entropy
- Apply to numerically-generated time series of the model

# To analyze: recast as a Markov Chain

## So, a quick review

Consider a discrete-time finite-state Markov chain  $Z(t)$ , i.e.

- $Z(t), t \in \mathbb{N}$ , takes values  $z_1, \dots, z_M$  in a space  $\mathcal{Z}$  of cardinality  $M$
- $v(t) \in \mathbb{R}_+^M$  is the probability vector  $v_i(t) = \Pr[Z(t) = z_i]$
- Transition matrix  $P$  with entries

$$P_{ij} = \Pr[Z(t+1) = z_j | Z(t) = z_i]$$

- Then  $v$  evolves according to

$$v(t+1) = v(t)P$$

- The long term behavior is  $v(t) \xrightarrow{t \rightarrow \infty} \pi$ , where the stationary vector  $\pi$  satisfies

$$\pi = \pi P$$

under mild assumptions on  $P$

# Recasting as a Markov chain: Main Idea

The states are binary vectors

$$\left\{ Z_1 = \begin{pmatrix} 0 \\ 0 \\ 0 \\ \vdots \\ 0 \end{pmatrix}, Z_2 = \begin{pmatrix} 1 \\ 0 \\ 0 \\ \vdots \\ 0 \end{pmatrix}, Z_3 = \begin{pmatrix} 0 \\ 1 \\ 0 \\ \vdots \\ 0 \end{pmatrix}, Z_4 = \begin{pmatrix} 1 \\ 1 \\ 0 \\ \vdots \\ 0 \end{pmatrix}, \dots, Z_{2^N} = \begin{pmatrix} 1 \\ 1 \\ 1 \\ \vdots \\ 1 \end{pmatrix} \right\}$$

The state vector is a  $2^N$ -dimensional probability vector  $v$  with components

$$v_i(t) = \Pr[X(t) = Z_i]$$

The transition law is determined from the dynamics

# Recasting as a Markov Chain: Ugly Details

Our RBN Model:

$$\Pr [X^i(t+1) = 1 | X^1(t) = x^1, \dots, X^N(t) = x^N] = \epsilon \left[ 1 + \sum_{j=1}^N w_{ij} x^j \right]$$

Letting  $z = [x^1, \dots, x^N]$  this becomes

$$\Pr [X^i(t+1) = 1 | X(t) = z] = \epsilon [1 + e_i^\top W z]$$

Allowing  $x^i_+ \in \{0, 1\}$

$$\begin{aligned} \Pr [X^i(t+1) = \underbrace{x^i_+}_{\in \{0, 1\}} | Z(t) = z] &= \underbrace{(1 - x^i_+)}_{=} + \epsilon \underbrace{(2x^i_+ - 1)}_{= \begin{cases} 1, & x^i_+ = 0 \\ 0, & x^i_+ = 1 \end{cases}} [1 + e_i^\top W z] \\ &= \begin{cases} -1, & x^i_+ = 0 \\ 1, & x^i_+ = 1 \end{cases} \end{aligned}$$

Taking a product of such terms yields the Markov transition matrix

$$\begin{aligned} P_{ij}(t) = \Pr [Z(t+1) = z_j | Z(t) = z_i] &= \prod_{k=1}^N \left\{ (1 - e_k^\top z_j) + \epsilon (2e_k^\top z_j - 1) [1 + e_k^\top W z_i] \right\} \\ &= P^{(0)} + \epsilon P^{(1)} + \epsilon^2 P^{(2)} + \mathcal{O}(\epsilon^3) \end{aligned}$$

# Setting up perturbation calculation

Expand  $\pi = \pi^{(0)} + \epsilon \pi^{(1)} + \epsilon^2 \pi^{(2)} + \dots$

$$(\pi^{(0)} + \epsilon \pi^{(1)} + \epsilon^2 \pi^{(2)}) (\mathbf{P}^{(0)} + \epsilon \mathbf{P}^{(1)} + \epsilon^2 \mathbf{P}^{(2)}) = \pi^{(0)} + \epsilon \pi^{(1)} + \epsilon^2 \pi^{(2)} + \dots$$

Separating by orders yields

$$\mathcal{O}(1) : \quad \pi^{(0)} (\mathbf{I} - \mathbf{P}^{(0)}) = 0$$

$$\sum_{j=1}^N \pi_j^{(0)} = 1,$$

$$\mathcal{O}(\epsilon) : \quad \pi^{(1)} (\mathbf{I} - \mathbf{P}^{(0)}) = \pi^{(0)} \mathbf{P}^{(1)},$$

$$\sum_{j=1}^N \pi_j^{(1)} = 0$$

$$\mathcal{O}(\epsilon^2) : \quad \pi^{(2)} (\mathbf{I} - \mathbf{P}^{(0)}) = \pi^{(0)} \mathbf{P}^{(2)} + \pi^{(1)} \mathbf{P}^{(1)}$$

$$\sum_{j=1}^N \pi_j^{(2)} = 0$$

Actually solving this was really hard...

# Solving the perturbation series I

The matrices  $P^{(j)}$  are  $2^N \times 2^N$  so we can only write them down explicitly for small  $N$ , need to work “in the abstract,” hybrid pencil&paper/Mathematica workflow

$$P_{ij}^{(0)} = \prod_{k=1}^N [(1 - e_k^\top z_j)] = \begin{cases} 1, & \|z_j\| = 0, \\ 0, & \|z_j\| > 0. \end{cases} \quad \text{where } \|Z\| = \sum_{k=1}^N x_k$$

$$P_{ij}^{(1)} = \sum_{r=1}^N \left\{ (2e_r^\top z_j - 1) [1 + e_r^\top W z_i] \prod_{\substack{k=1 \\ k \neq r}}^N (1 - e_k^\top z_j) \right\} = \begin{cases} -[N + 1_N^\top W z_i], & \|z_j\| = 0 \\ [1 + z_j^\top W z_i], & \|z_j\| = 1 \\ 0, & \|z_j\| > 1 \end{cases}$$

$$P_{ij}^{(2)} = \sum_{\substack{r,s=1 \\ r>s}}^N \left\{ (2e_r^\top z_j - 1) [1 + e_r^\top W z_i] (2e_s^\top z_j - 1) [1 + e_s^\top W z_i] \prod_{\substack{k=1 \\ k \neq r,s}}^N (1 - e_k^\top z_j) \right\}$$

$$= \begin{cases} \sum_{\substack{r,s=1 \\ r>s}}^N \{ [1 + e_r^\top W z_i] [1 + e_s^\top W z_i] \}, & \|z_j\| = 0, \\ -[1 + z_j^\top W z_i] [N - 1 + (1_N^\top - z_j^\top) W z_i], & \|z_j\| = 1, \\ \{ 1 + z_j^\top W z_i + [e_{\mathcal{I}_1(z_j)}^\top W z_i] [e_{\mathcal{I}_2(z_j)}^\top W z_i] \}, & \|z_j\| = 2, \\ 0, & \|z_j\| > 2, \end{cases}$$

# Solving the Perturbation Series II

Solve for the stationary vector order by order

$$\pi_i^{(0)} = \begin{cases} 1, & \|z_i\| = 0, \\ 0, & \|z_i\| > 0, \end{cases}$$

$$\pi_i^{(1)} = \begin{cases} -N, & \|z_i\| = 0, \\ 1, & \|z_i\| = 1, \\ 0, & \|z_i\| > 1, \end{cases}$$

$$\pi_i^{(2)} = \begin{cases} \binom{N}{2} - 1_N^\top W 1_N, & \|z_i\| = 0, \\ -(N-1) + z_i^\top W 1_N, & \|z_i\| = 1, \\ 1, & \|z_i\| = 2, \\ 0, & \|z_i\| > 2, \end{cases}$$

# Calculating Transfer Entropy

from node 2 to node 1

$$TE^{2 \rightarrow 1} = \sum_{x_+^1, x^1, x^2} \Pr[X_{(t+1)}^1 = x_+^1, X_t^1 = x^1, X_t^2 = x^2] \log \frac{\Pr[X_{(t+1)}^1 = x_+^1 | X_t^1 = x^1, X_t^2 = x^2]}{\Pr[X_{(t+1)}^1 = x_+^1 | X_t^1 = x^1]}$$

- Use the transition rule and the asymptotic expansions to approximate the various probabilities and conditional probabilities
- Use some tricks to avoid dividing by small numbers
- Get terms like

$$\Pr[X_{t+1}^1 = x_+^1 | X_t^1 = x^1, X^2(t) = x^2] = \\ (1 - x_+^1) + (2x_+^1 - 1) \left( \epsilon [1 + W_{11}x^1 + W_{12}x^2] + \epsilon^2 \sum_{j=3}^N W_{1j} \right) + \mathcal{O}(\epsilon^3)$$

# The next-order correction

Finally, we arrive at

$$TE^{j \rightarrow i} = \epsilon^2 G^{(2)}(W_{ij}) + \epsilon^3 G_{ij}^{(3)}(W) + \mathcal{O}(\epsilon^4)$$

where

$$G^{(2)}(x) = -x + (1+x) \log(1+x)$$

$$G_{ij}^{(3)}(W) = W_{ij} (W_{ij} - d_i - d_j) + \log(1+W_{ij}) (d_i - W_{ij} + (1+W_{ij})d_j)$$

and  $d_j$  is the weighted in-degree of node  $j$

$$d_j = \sum_{k=1}^N W_{jk}.$$

# Solving for the weights: one more perturbation expansion

Letting

$$T_{ij} \equiv \frac{TE^{j \rightarrow i}}{\epsilon^2} = G^{(2)}(W_{ij}) + \epsilon G_{ij}^{(3)}(W) + \mathcal{O}(\epsilon^2)$$

and

$$W = W^{(0)} + \epsilon W^{(1)} + \mathcal{O}(\epsilon^2)$$

we arrive at desired formula

$$W_{ij}^{(0)} = [G^{(2)}]^{-1}(T_{ij})$$

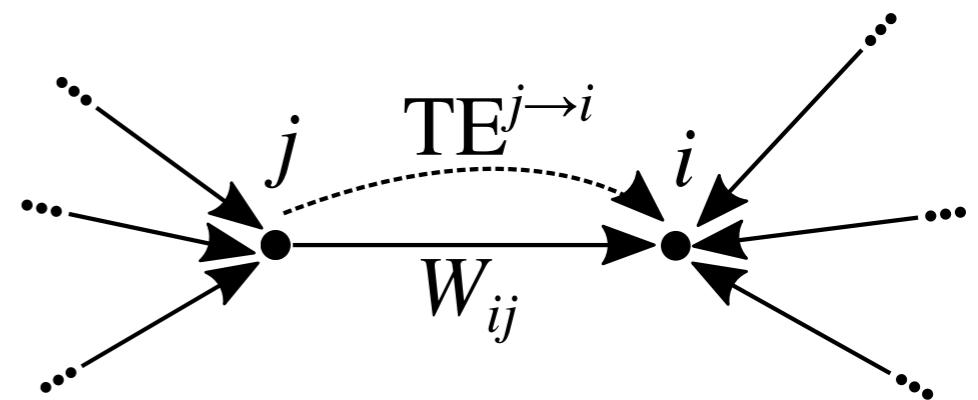
$$W_{ij}^{(1)} = -\frac{G_{ij}^{(3)}(W^{(0)})}{\frac{d}{dw} G^{(2)}(w)|_{w=W_{ij}^{(0)}}}$$

# Interpreting the correction term

If  $W_{ij} \ll 1$  then the correction to the computed transfer entropy is

$$G_{ij}^{(3)}(W) \sim \frac{W_{ij}^2}{2} (d_j - d_i + W_{ij}) + \mathcal{O}(\|W\|^3)$$

The difference in weighted in-degree of the two nodes

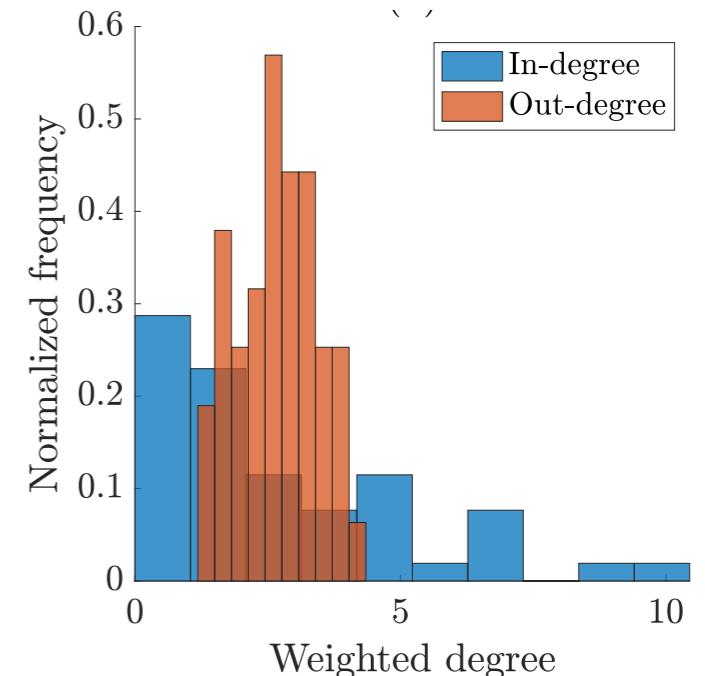


This yields a correction to the computed weight

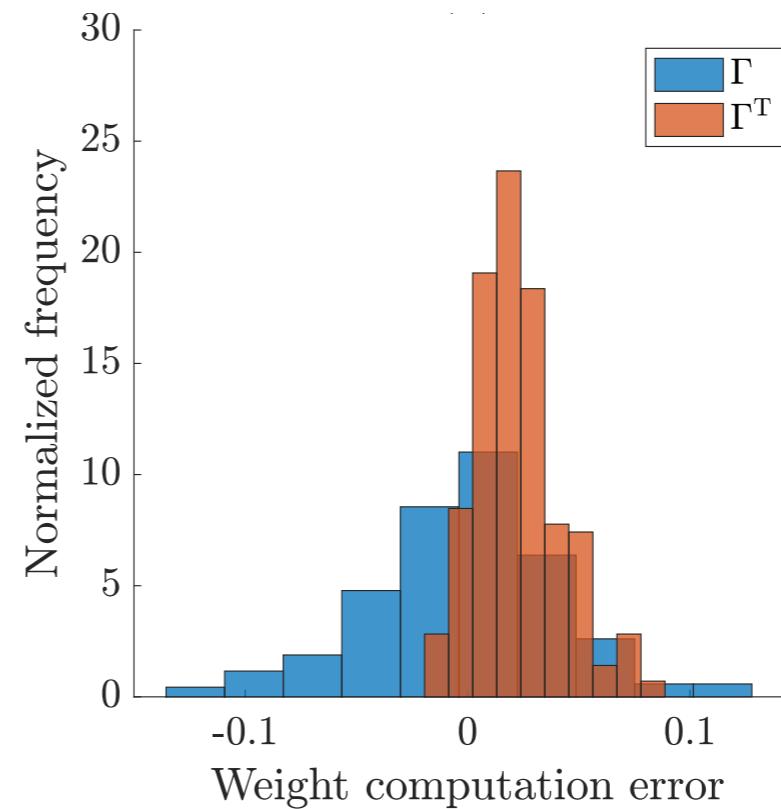
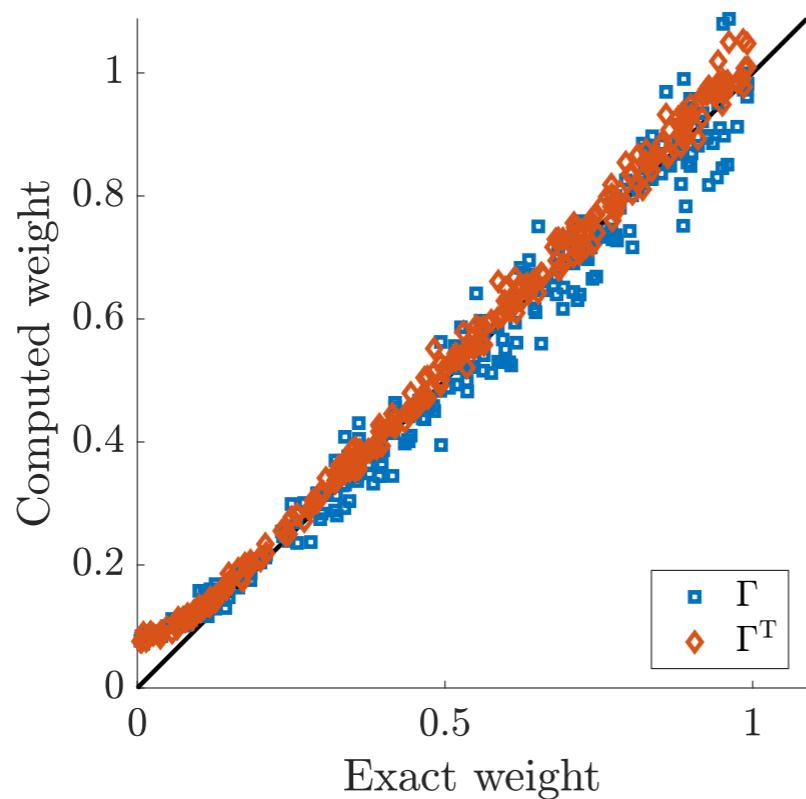
$$W_{ij}^{(1)} \sim \frac{W_{ij}^{(0)}}{2} \left( d_j^{(0)} - d_i^{(0)} + W_{ij}^{(0)} \right) + \mathcal{O} \left( \|w^{(0)}\|^2 \right)$$

# Return to numerical example

The network  $\Gamma$  was constructed so that its in-degrees vary more widely than its out-degrees

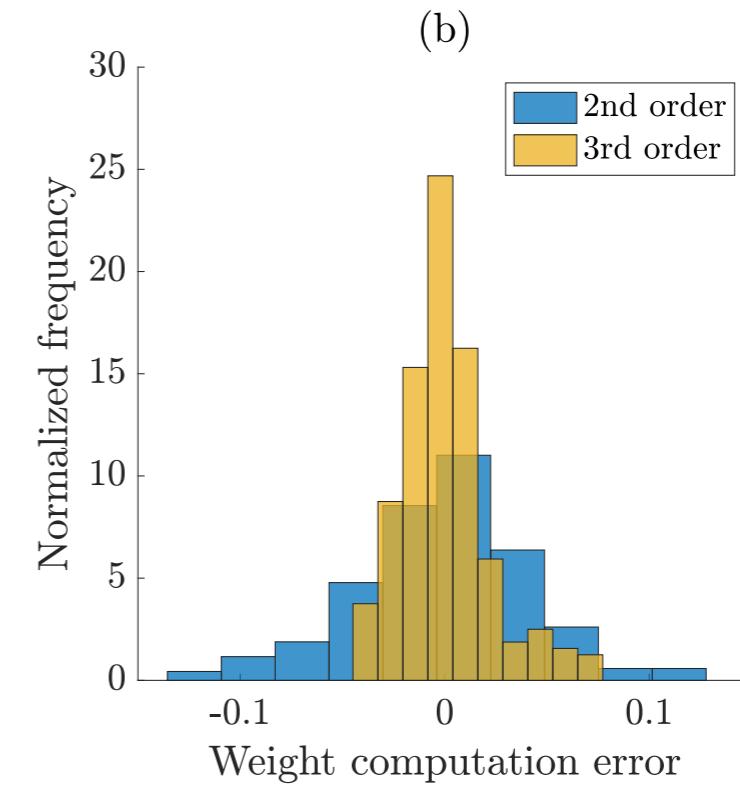
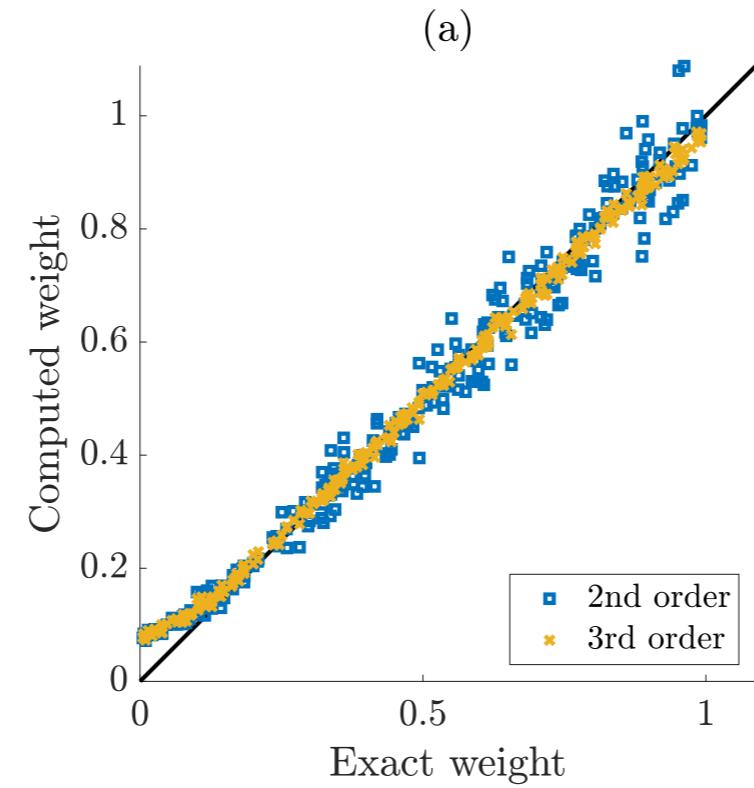


This leads to a larger variance in the computed weights for  $\Gamma$  than for  $\Gamma^\top$

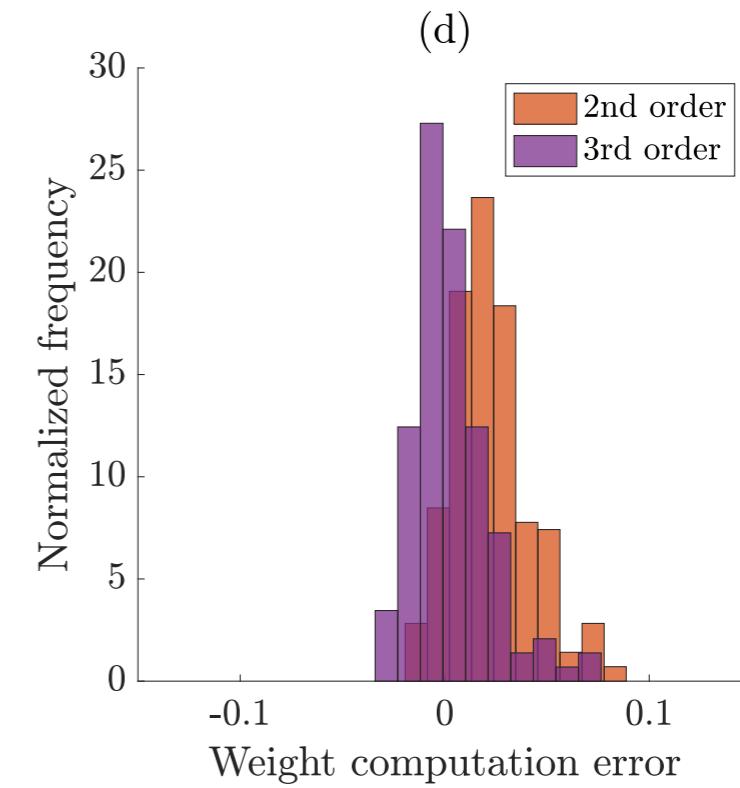
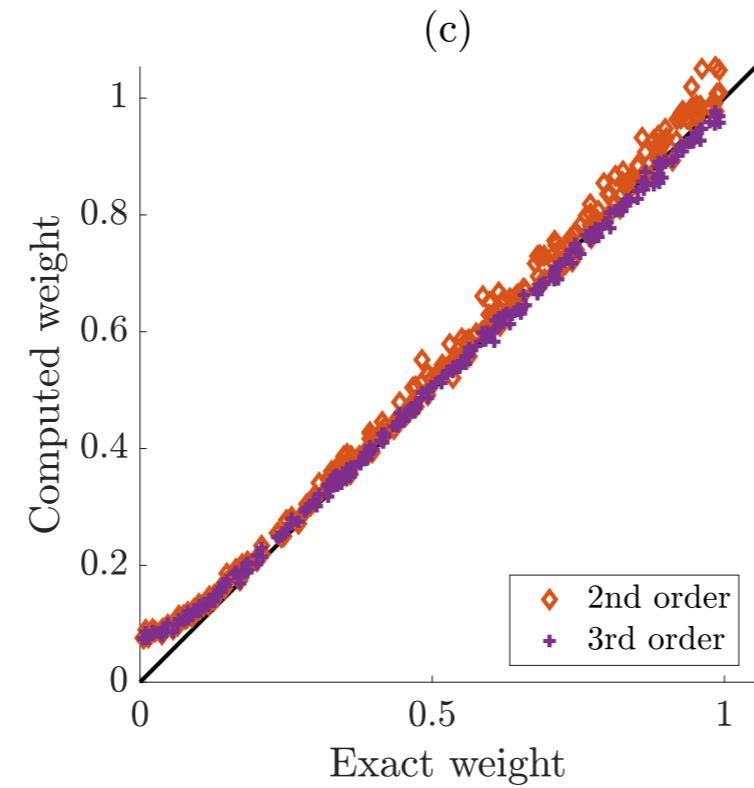


# Putting in the correction

$\Gamma$



$\Gamma^\top$



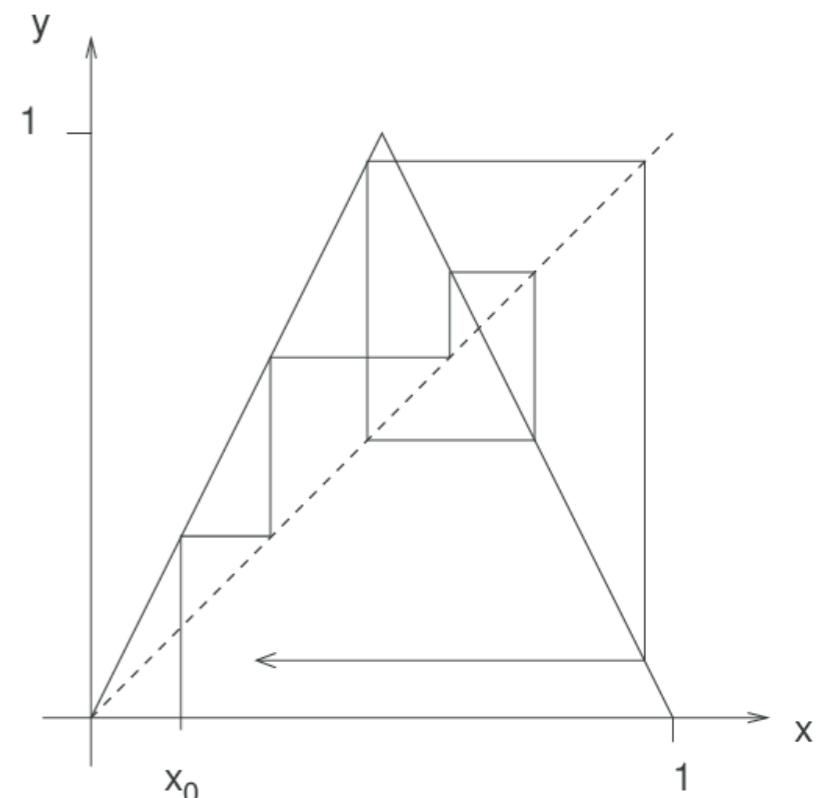
# Question: Is this a general phenomenon?

Does this long and painful calculation actually tell us anything about connection between the network structure and the accuracy of the computation?

Let's look at another model.

The *tent map* is a simple chaotic system

$$x_{t+1} = F(x_t) \equiv \begin{cases} 2x_t, & 0 \leq x_t < \frac{1}{2} \\ 2(1 - x_t) & \frac{1}{2} \leq x_t \leq 1 \end{cases}$$



We consider a system of coupled tent maps with noise

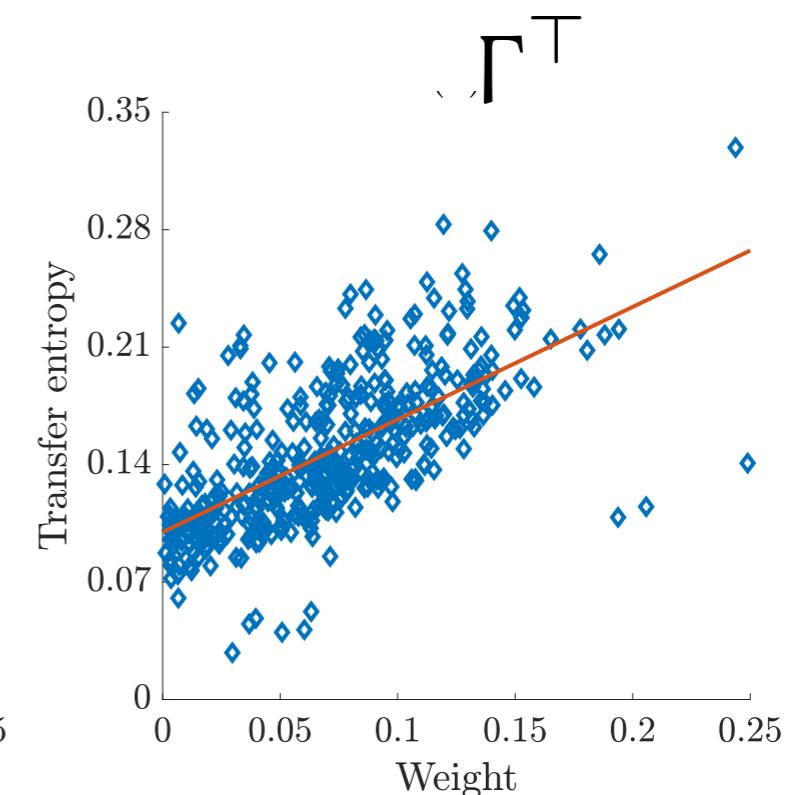
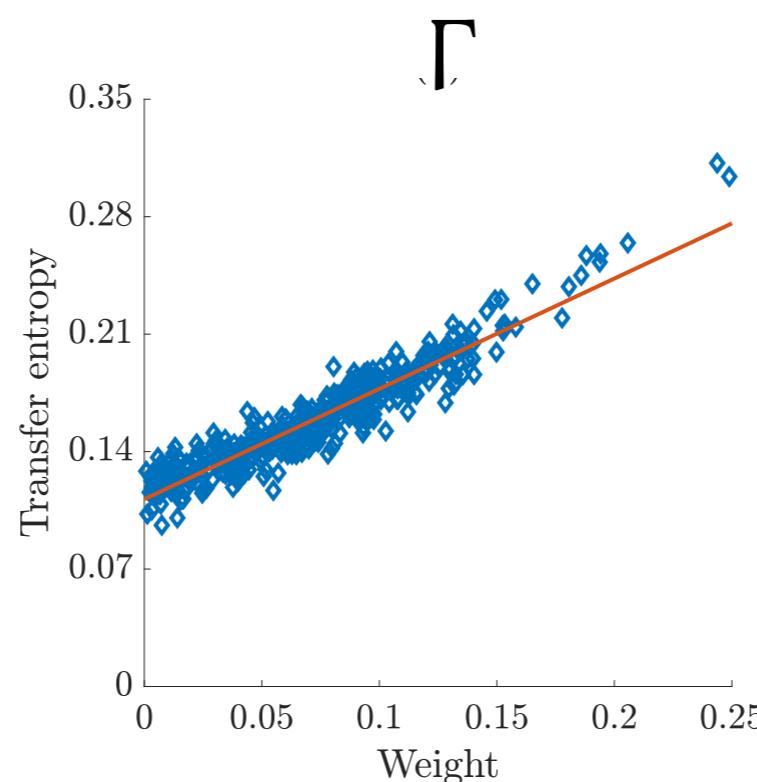
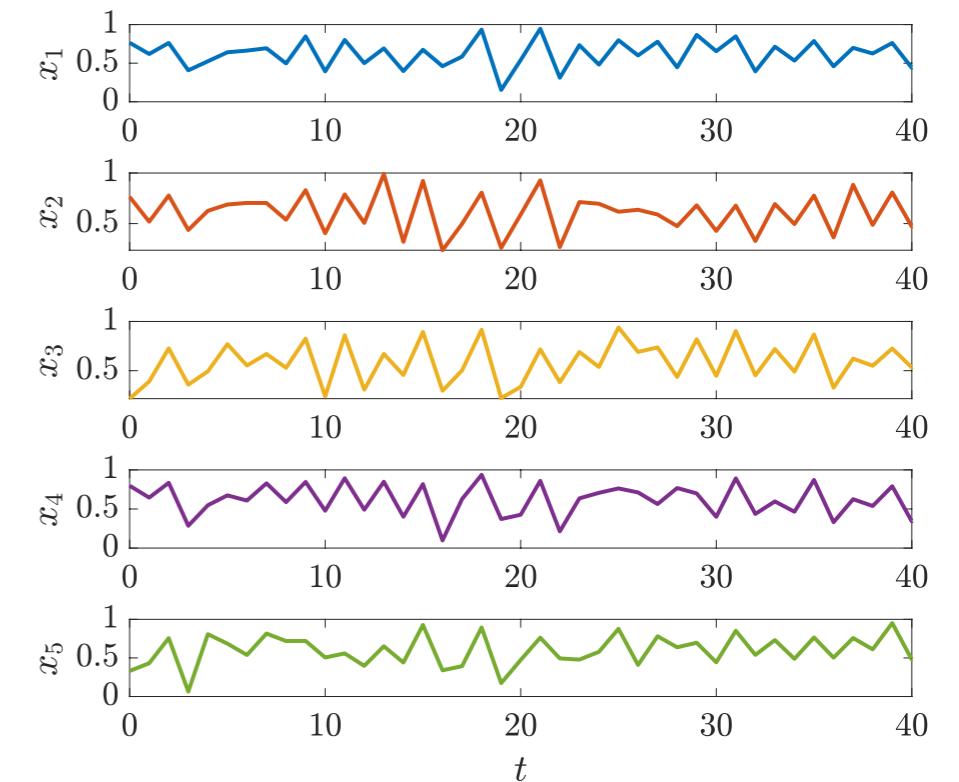
$$x_{t+1}^i = F(x_t^i) + \epsilon \sum_{j=1}^N W_{ij} (F(x_t^j) - F(x_t^i)) + \sigma n^i(t)$$

# The coupled tent map example

Simulate for two 30-node networks  $\Gamma$  and  $\Gamma^\top$  and observe oscillators going in and out of synchronization

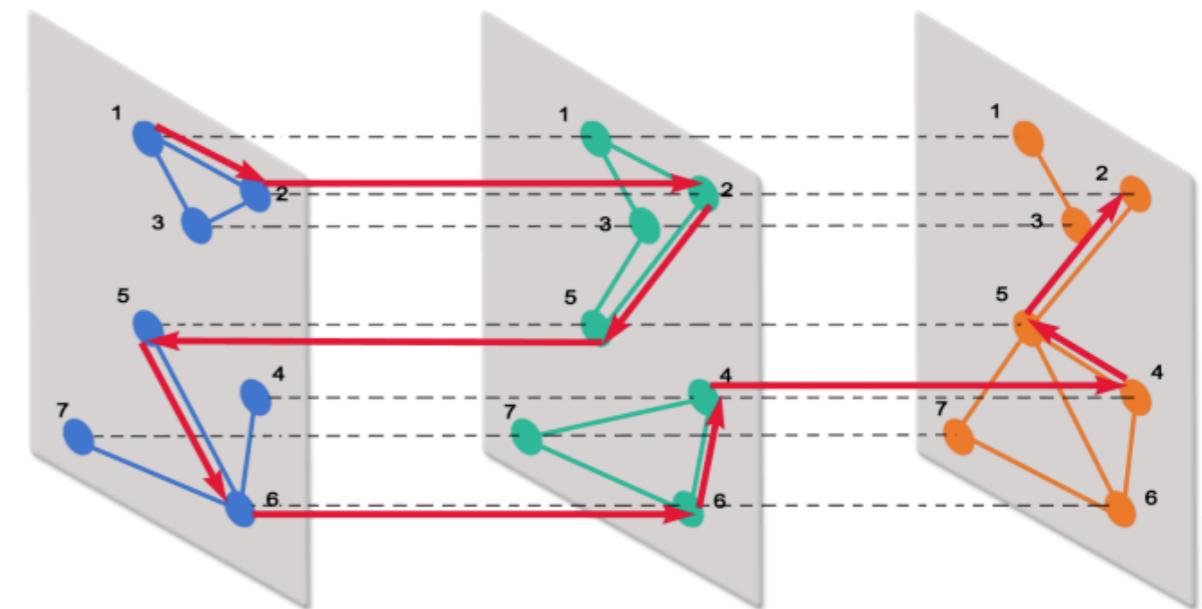
The networks constructed such that their distribution of weighted in-degrees very different

Variance in computed transfer entropy increases with variance in weighted in-degree



# Extending the model: a multi-layered graph

There are many generalizations of the concept of a graph, e.g. multigraph in which there exist different types of edges. organized in layers



Maurizio's interest: Fish communicate over multiple channels

Stimulus fish



Light, eyes, instantaneous

Responding fish



Fluid flow, lateral line, delayed

# Multilayered delay RBN model

$$\Pr(x_i(t) = 1) = \epsilon \left( 1 + \sum_{m=1}^M \sum_{j=1}^N W_{ij}^{(m)} x_j(t-m) \right)$$

Rewrite as a suspended system with coordinate

$$y(t) = \begin{pmatrix} x(t) \\ x(t-1) \\ \vdots \\ x(t+1-M) \end{pmatrix} = \begin{pmatrix} y_{(1)}(t) \\ y_{(2)}(t) \\ \vdots \\ y_{(M)}(t) \end{pmatrix} \in \{0, 1\}^{MN}$$

The first  $N$  coordinates are updated probabilistically

$$\Pr(y_i(t) = 1) = \epsilon \left( 1 + \sum_{j=1}^{MN} W_{ij} y_j(t-1) \right), \text{ where } W = [W^{(1)} \ W^{(2)} \ \dots \ W^{(M)}]$$

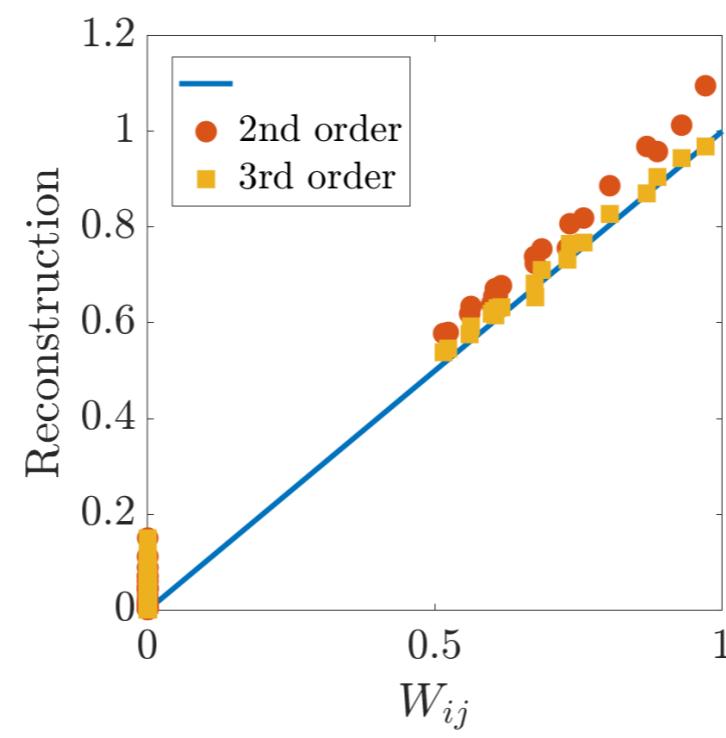
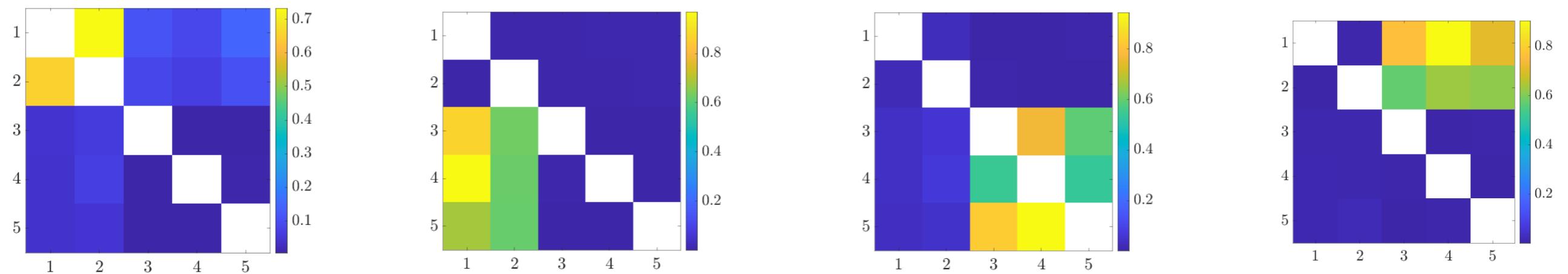
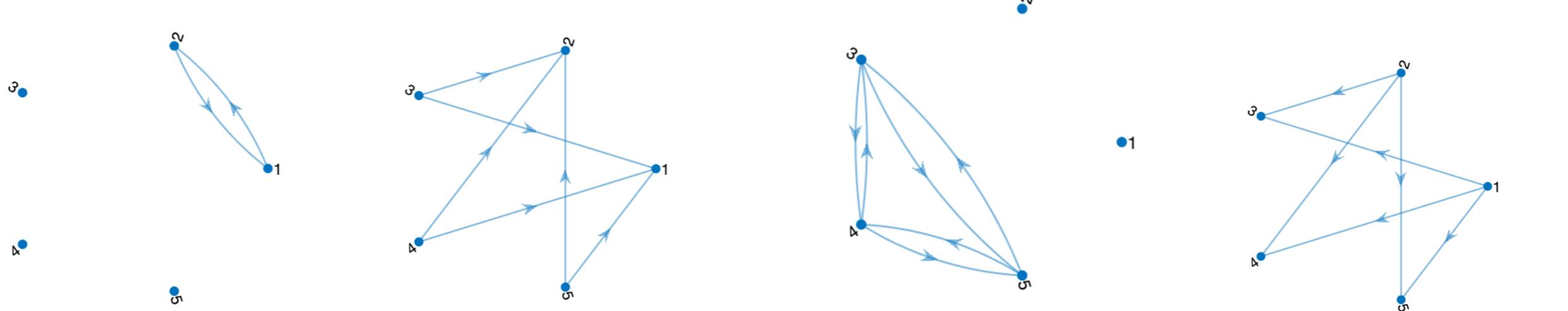
and the remaining  $(N-1)M$  deterministically

$$y_{(m)}(t) = y_{(m-1)}(t-1), \text{ for } m = 2, \dots, M$$

# The result of the calculation

- Repeat our procedure
  - Reformulate as a Markov chain on a  $2^{MN}$  dimensional state space
  - Calculate stationary vector
  - Use the transition law and the stationary vector to compute transfer entropy in terms of weights  $W_{ij}$
- This is the worst calculation I had to do in my entire life
- The result is entirely analogous to what I obtained in the first problem

# One last numerical experiment





why,  
thank  
you.