# Secure Compilation Lecture 2

Renate Robin Eilers Cristina Matache Baber Rehman

June 24, 2019

This is the second talk presented by Amal Ahmed in OPLSS 2019, University of Oregon, USA.

## 1 Introduction

### 1.1 Source Language

#### 1.1.1 Types

We just have integers and functions in source language.

$$\sigma ::= \text{int} \mid \sigma_1 \to \sigma_2$$

#### 1.1.2 Terms

#### 1.2 Target Language

**Theorem 1.1** (Type Preservation). If  $\Gamma \vdash e_S : \sigma \text{ and } \Gamma \vdash e_S : \alpha \leadsto e_T \text{ then } \Gamma_S^+ \vdash e_T : \sigma^+$ 

For correctness, we want to show  $e_S \approx e_T$ . This is not contexual equivalence because source language and target language are two different languages. There are many ways to prove compiler correction. We want to say that when:

$$e_S \approx e_T$$
 then  $\sigma \approx \sigma^+$ 

$$\begin{split} & \text{V} \llbracket \ \sigma \ \rrbracket = \{ \ (V_S \ , \ V_T) \ j \ . \vdash V_S : \sigma \wedge . \ ; \ . \vdash V_T : \sigma^+ \ ..... \ \} \\ & \text{V} \llbracket \ \text{ints} \ \rrbracket = \{ \ (n_S \ , n_T) \ \} \\ & \text{V} \llbracket \ \sigma_1 \rightarrow \sigma_2 \ \rrbracket = \{ \ (\lambda \mathbf{x} : \sigma_1 \ . \ e_S \ \} \end{split}$$