

## Exercise 1

### Question 1

Flow tcp1 is competing with flow tcp4 on link n1-n2 and link n2-n4. Flow tcp2 has less RTT and it gets higher share of the bandwidth on link n2-n4, which is recorded at n5.

### Question 2

Flow tcp1 is using Slow Start mechanism.

### Question 3

Between 0.5 – 2.0 sec, tcp1 is the only flow active, but it is still in the SS phase and must compete with the other flows starting at 2.0 sec before it can discover the maximum bandwidth.

## Exercise 2

### Question 1

- 2000-bytes and 3500-bytes payload would cause fragmentation, because the MTU is 1500 bytes
- 192.168.1.103 has fragmented the original datagram
- 2 fragments have been created when data size is specified as 2000.

| No. | Time      | Source                | Destination     | Protocol | Length | Info   |
|-----|-----------|-----------------------|-----------------|----------|--------|--|
| 1   | 0.000000  | 192.168.1.1           | 255.255.255.255 | UDP      | 215    | 36861 → 7437 Len=173   |
| 2   | 1.433512  | fe80::ec49:66ff:fe... | ff02::1         | ICMPv6   | 78     | Router Advertisement from e8:de:27:4d:a1:40                                |
| 3   | 3.071798  | 192.168.1.1           | 255.255.255.255 | UDP      | 215    | 36861 → 7437 Len=173   |
| 4   | 3.262721  | 192.168.1.103         | 8.8.8.8         | ICMP     | 98     | Echo (ping) request id=0xd805, seq=0/0, ttl=64 (reply in 5)                |
| 5   | 3.287081  | 8.8.8.8               | 192.168.1.103   | ICMP     | 98     | Echo (ping) reply id=0xd805, seq=0/0, ttl=122 (request in 4)               |
| 6   | 3.993354  | 192.168.1.106         | 224.0.0.251     | MDNS     | 108    | Standard query 0x0000 PTR _homekit._tcp.local, "QM" question OPT           |
| 7   | 3.993713  | fe80::18fe:846:f78... | ff02::fb        | MDNS     | 128    | Standard query 0x0000 PTR _homekit._tcp.local, "QM" question OPT           |
| 8   | 4.264498  | 192.168.1.103         | 8.8.8.8         | ICMP     | 98     | Echo (ping) request id=0xd805, seq=1/256, ttl=64 (reply in 9)              |
| 9   | 4.286191  | 8.8.8.8               | 192.168.1.103   | ICMP     | 98     | Echo (ping) reply id=0xd805, seq=1/256, ttl=122 (request in 8)             |
| 10  | 5.269254  | 192.168.1.103         | 8.8.8.8         | ICMP     | 98     | Echo (ping) request id=0xd805, seq=2/512, ttl=64 (reply in 11)             |
| 11  | 5.291311  | 8.8.8.8               | 192.168.1.103   | ICMP     | 98     | Echo (ping) reply id=0xd805, seq=2/512, ttl=122 (request in 10)            |
| 12  | 6.045195  | 192.168.1.1           | 255.255.255.255 | UDP      | 215    | 36861 → 7437 Len=173   |
| 13  | 9.113418  | 192.168.1.1           | 255.255.255.255 | UDP      | 215    | 36861 → 7437 Len=173   |
| 14  | 9.169162  | 192.168.1.103         | 255.255.255.255 | DB-LS... | 268    | Dropbox LAN sync Discovery Protocol, JavaScript Object Notation            |
| 15  | 9.169892  | 192.168.1.103         | 192.168.1.255   | DB-LS... | 268    | Dropbox LAN sync Discovery Protocol, JavaScript Object Notation            |
| 16  | 10.558043 | 192.168.1.103         | 8.8.8.8         | IPv4     | 1514   | Fragmented IP protocol (proto=ICMP 1, off=0, ID=a13d) (Reassembled in #17) |
| 17  | 10.558045 | 192.168.1.103         | 8.8.8.8         | ICMP     | 562    | Echo (ping) request id=0xd905, seq=0/0, ttl=64 (reply in 19)               |
| 18  | 10.610386 | 8.8.8.8               | 192.168.1.103   | IPv4     | 1482   | Fragmented IP protocol (proto=ICMP 1, off=0, ID=dfd0) (Reassembled in #19) |
| 19  | 10.612610 | 8.8.8.8               | 192.168.1.103   | ICMP     | 594    | Echo (ping) reply id=0xd905, seq=0/0, ttl=122 (request in 17)              |
| 20  | 10.649226 | fe80::ec49:66ff:fe... | ff02::1         | ICMPv6   | 78     | Router Advertisement from e8:de:27:4d:a1:40                                |
| 21  | 11.563299 | 192.168.1.103         | 8.8.8.8         | IPv4     | 1514   | Fragmented IP protocol (proto=ICMP 1, off=0, ID=aaf2) (Reassembled in #22) |
| 22  | 11.563302 | 192.168.1.103         | 8.8.8.8         | ICMP     | 562    | Echo (ping) request id=0xd905, seq=1/256, ttl=64 (reply in 24)             |
| 23  | 11.600293 | 8.8.8.8               | 192.168.1.103   | IPv4     | 1412   | Fragmented IP protocol (proto=ICMP 1, off=0, ID=a2e9) (Reassembled in #24) |
| 24  | 11.609956 | 8.8.8.8               | 192.168.1.103   | ICMP     | 594    | Echo (ping) reply id=0xd905, seq=1/256, ttl=122 (request in 22)            |
| 25  | 12.082915 | 192.168.1.1           | 255.255.255.255 | UDP      | 215    | 36861 → 7437 Len=173   |
| 26  | 12.568393 | 192.168.1.103         | 8.8.8.8         | IPv4     | 1514   | Fragmented IP protocol (proto=ICMP 1, off=0, ID=4a07) (Reassembled in #27) |
| 27  | 12.568394 | 192.168.1.103         | 8.8.8.8         | ICMP     | 562    | Echo (ping) request id=0xd905, seq=2/512, ttl=64 (reply in 29)             |

### Question 2

Yes, since the MTU for the last link is 1500 bytes, we cannot be sure who fragmented the original 3500 bytes reply.

### Question 3

- Packet 39: ID=0x7a7b, length=1500, flag=0x01(MF bit on), offset=0

- Packet 40: ID=0x7a7b, length=1500, flag=0x01(MF bit on), offset=1480
- Packet 41: ID=0x7a7b, length=568, flag=0, offset=2960

#### Question 4

No fragmentation of fragments has occurred for the reply from 8.8.8.8 for data size of 3500. We cannot be sure the fragmentation from 192.168.1.103 to 8.8.8.8 since the reassembly is only done at the destination.

#### Question 5

The fragments would be incomplete, and the receiver would discard.

### Exercise 3

#### Question 1

- Node 0 communicates with Node 5 follow route 0 - 1 - 4 - 5
- Node 2 communicates with Node 5 follow route 2 - 3 - 5

It does NOT change over time.

#### Question 2

- At time  $t=1.0$ , link 1-4 goes down, node 0 – node 5 not changed, node 0 cannot reach node 5
- At time  $t=1.2$ , link 1-4 goes up, node 0 can reach node 5, node 1 can reach node 4

Traffic between node 2 and node 5 does not affected.

#### Question 3

The DV routing protocol discovers a different route (0-1-2-3-5) when link 1-4 goes down. When 1-4 goes up, the routing protocol is 0-1-4-5.

#### Question 4

The cost of link 1-4 increase to 3, the cost of 0-1-4-5 is 4, and the cost of 0-1-2-3-5 is 4. Therefore, the flow between node 0 and node 5 now uses route 0- 1-2-3-5, because it has the lowest cost.

#### Question 5

For node 0 and node 5, the total cost does not change. But from node 2 to node 5, there are two routes of equal total cost (2-3-5 and 2-1-4-5). Since the network is using multipath routing, node 2 will split traffic equally on both paths.