# **Practice 3**

# **Finite fields**

 $GF(2^n)$  and AES key schedule (128 / 192 bits)

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## **Description**

In this session we will work with finite fields. Please do the following programming exercises on your own. Use only one programming language (C/C++, Java, Python). It is recommended that you try to solve the exercises on your own. You can discuss the solution to the exercises with your colleagues but you should not copy source code. If copying is detected, that may immediately lead to a grade less than 6.

### **Programming Exercises**

- 1. Design a function that receives  $3 \le n \le 8$  for  $GF(2^n)$  and as outputs the multiplication table. Consider the following requirements.
  - a. The user can choose to represent each element in  $GF(2^n)$  as a polynomial or as an hexadecimal number.
  - b. The output must be stored in a file.
- 2. Implement the key schedule for AES, considering a key size of 128 bits. Your program must receive the key in hexadecimal and must store the 10 subkeys (derived from the key) in a file.
- 3. Repeat the previous point, but now for a key size of 192 bits.

#### **Products**

- → The most important parts of the source code.
- → Include screen capture of your programs showing how they work.
- → Please write a small user manual to know how to run your programs.

# Table generation

# Polynomial multiplied with x^1

In table.hpp

```
// Multiplies a polynomial 'a' with x^1 in its binary
// representation modulated 'm', under the field 'n'
inline
short multByX(short a, short m, int n){

    m &= ~((-1) << n);

    return (a & (1 << (n - 1))) ? ((a << 1) ^ m) : (a << 1);
}</pre>
```

## Polynomial multiplied with x^times

In table.hpp

```
// Multiplies a polinomial 'a' with x^times in its binary
// representation modulated 'm', under the field 'n'
inline
short multByManyX(short a, short m, int n, int times){

   if (times == 0) return a;
   if (times == 1) return multByX(a, m, n);

   for (int i = 0; i < times; ++i)
        a = multByX(a, m, n);

   return a;
}</pre>
```

## Multiply two polynomials module m

In table.hpp

#### Table creation

In table.hpp > table > writeTable()

# **AES Key Schedule**

# The last word rotation (first step)

In aes.hpp

```
// Rotate the given word
inline
void shiftColumn(unsigned int & column){
   auto aux = column >> 24;
   column = column << 8 | aux;
}</pre>
```

## The sbox substitution (second step)

In aes.hpp

```
// Substitute the values of the column with the sbox's values
inline
void sBoxSubs(unsigned int & column){
   unsigned char * aux = (unsigned char *)(&column);

   for (int i = 0; i < 4; ++i)
      aux[i] = sBox[aux[i] >> 4][aux[i] & 0xF];
}
```

# XOR with the round coefficient (third step)

In aes.hpp

```
// XOR with the round coeffient
inline
void roundCoeffiecient(unsigned int & column, int i){
    column ^= rc[i] << 24;
}</pre>
```

# XOR with all the columns (last step)

In aes.hpp

```
// XOR with the rest of the columns
inline
void finalBuild(std::vector<unsigned int> & key, unsigned int column){
   for (int i = 0; i < key.size(); ++i){
        key[i] ^= column;
        column = key[i];
   }
}</pre>
```

## **AES key schedule round**

In aes.hpp

```
// AES key schedule round
inline
void round(std::vector<unsigned int> & key, int i){
   unsigned int column = key[key.size() - 1];
   shiftColumn(column);
   sBoxSubs(column);
   roundCoeffiecient(column, i);
   finalBuild(key, column);
}
```

## **AES key schedule**

In aes.hpp

```
inline
auto keySchedule(std::string keyString){
    if (keyString.size() != 32 && keyString.size() != 48)
        return std::vector<unsigned int>(0);
    auto key = getKeyFromStr(keyString);
    std::vector<unsigned int> result(key.size() == 4 ? 44 : 52);
    for (int i = 0; i < result.size(); i += key.size()){</pre>
         for (int j = 0; j < \text{key.size}() \&\&\& i + j < \text{result.size}(); ++j)
             result[i + j] = key[j];
        if (i + key.size() >= result.size())
             break;
        round(key, i / key.size());
    return result;
```

## **AES Key Schedule pseudocode**

```
KeyExpansion(byte key[4*Nk], word w[Nb*(Nr+1)], Nk)
begin
  word temp
  i = 0
  while (i < Nk)
      w[i] = word(key[4*i], key[4*i+1], key[4*i+2], key[4*i+3])
      i = i+1
   end while
   i = Nk
  while (i < Nb * (Nr+1)]
      temp = w[i-1]
      if (i \mod Nk = \theta)
         temp = SubWord(RotWord(temp)) xor Rcon[i/Nk]
      else if (Nk > 6 \text{ and i mod } Nk = 4)
         temp = SubWord(temp)
      end if
      w[i] = w[i-Nk] xor temp
      i = i + 1
   end while
end
Note that Nk=4, 6, and 8 do not all have to be implemented;
they are all included in the conditional statement above for
conciseness.
               Specific implementation requirements for the
Cipher Key are presented in Sec. 6.1.
```

Figure 11. Pseudo Code for Key Expansion.<sup>2</sup>

# Instructions

# Compilation

Just use the command make:

make

or use the compiler:

```
g++ ./src/p3 -o ./build/p3
```

And change to the build folder, where are some input examples:

cd build

### Usage

The program must be used as follows:

```
./p3 < in.txt
```

The in.txt must be have the following structure:

An integer O,  $0 \le O \le 1$ , that selects between the table creation (0) or the aes schedule (1).

If the option 0 (table creation) is selected, the next line must be a integer n,  $3 \le n \le 8$ , the order of the Galois Field, and the next line a string: bin or hex, the way of output the numbers. The result will be in table.txt

If the option 1 (aes schedule) is selected, the nex line must be a string with a valid hexadecimal representation of a 128 bits key or 192 bits key, without spaces and with 0 values as 00 (in order to have 32 or 48 characters respectively).

## **Examples**

#### Example 1

in.txt

```
0
3
bin
```

table.txt

#### Example 2

in.txt

```
0
4
hex
```

#### table.txt

```
m(x): 13
123456789abcdef
2 4 6 8 a c e 3 1 7 5 b 9 f d
3 6 5 c f a 9 b 8 d e 7 4 1 2
4 8 c 3 7 b f 6 2 e a 5 1 d 9
5 a f 7 2 d 8 e b 4 1 9 c 3 6
6 c a b d 7 1 5 3 9 f e 8 2 4
7 e 9 f 8 1 6 d a 3 4 2 5 c b
8 3 b 6 e 5 d c 4 f 7 a 2 9 1
9 1 8 2 b 3 a 4 d 5 c 6 f 7 e
a 7 d e 4 9 3 f 5 8 2 1 b 6 c
b 5 e a 1 f 4 7 c 2 9 d 6 8 3
c b 7 5 9 e 2 a 6 1 d f 3 4 8
d 9 4 1 c 8 5 2 f b 6 3 e a 7
ef1d32c97684ab5
fd2964b1ec3875a
```

#### Example 3

in.txt

```
1
a0000000000000000000000000000000000
```

#### key.txt

```
Round: 0
A0 00 00 00
00 00 00
```

```
00 00 00 B0
00 00 00 00
Round: 1
C2 C2 C2 C2
E7 E7 E7 E7
63 63 63 D3
63 63 63 63
Round: 2
54 96 54 96
81 66 81 66
98 FB 98 4B
46 25 46 25
Round: 3
63 F5 A1 37
32 54 D5 B3
A7 5C C4 8F
D6 F3 B5 90
Round: 4
06 F3 52 65
41 15 C0 73
C7 9B 5F D0
4C BF 0A 9A
Round: 5
99 6A 38 5D
31 24 E4 97
7F E4 BB 6B
01 BE B4 2E
Round: 6
31 5B 63 3E
4E 6A 8E 19
4E AA 11 7A
4D F3 47 69
Round: 7
A5 FE 9D A3
94 FE 70 69
B7 1D 0C 76
FF 0C 4B 22
Round: 8
DC 22 BF 1C
AC 52 22 4B
24 39 35 43
F5 F9 B2 90
Round: 9
74 56 E9 F5
B6 E4 C6 8D
44 7D 48 0B
69 90 22 B2
Round: 10
1F 49 A0 55
9D 79 BF 32
73 0E 46 4D
8F 1F 3D 8F
```

#### **Example 4**

in.txt

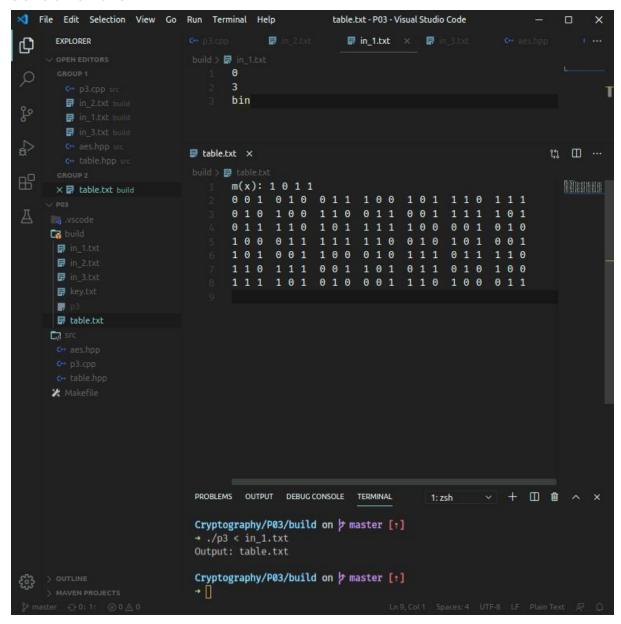
```
1
8e73b0f7da0e6452c810f32b809079e562f8ead2522c6b7b
```

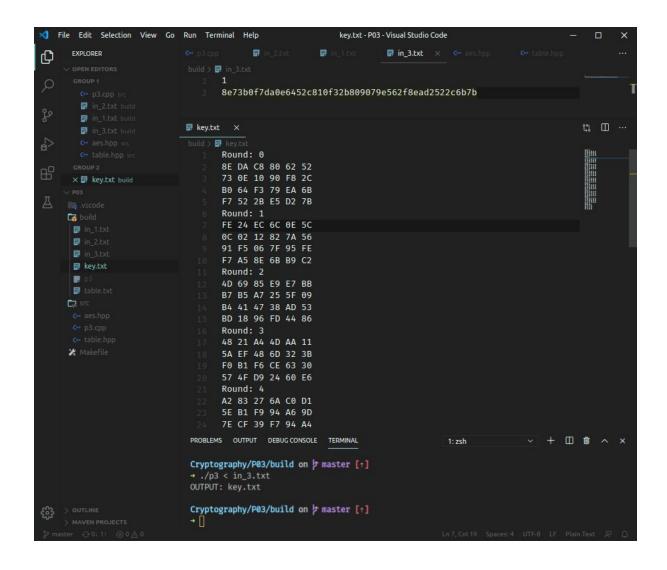
#### key.txt

```
Round: 0
8E DA C8 80 62 52
73 0E 10 90 F8 2C
B0 64 F3 79 EA 6B
F7 52 2B E5 D2 7B
Round: 1
FE 24 EC 6C 0E 5C
0C 02 12 82 7A 56
91 F5 06 7F 95 FE
F7 A5 8E 6B B9 C2
Round: 2
4D 69 85 E9 E7 BB
B7 B5 A7 25 5F 09
B4 41 47 38 AD 53
BD 18 96 FD 44 86
Round: 3
48 21 A4 4D AA 11
5A EF 48 6D 32 3B
F0 B1 F6 CE 63 30
57 4F D9 24 60 E6
Round: 4
A2 83 27 6A C0 D1
5E B1 F9 94 A6 9D
7E CF 39 F7 94 A4
D5 9A 43 67 07 E1
Round: 5
EC 6F 48 22 E2 33
17 A6 5F CB 6D F0
86 49 70 87 13 B7
EB 71 32 55 52 B3
Round: 6
40 2F 67 45 A7 94
BE 18 47 8C E1 11
EB A2 D2 55 46 F1
28 59 6B 3E 6C DF
Round: 7
82 AD CA 8F 28 BC
1F 07 40 CC 2D 3C
75 D7 05 50 16 E7
0A 53 38 06 6A B5
Round: 8
E9 44 8E 01
```

```
8B 8C CC 00
A0 77 72 22
6F 3C 04 02
```

#### **Screenshots**





The complete source code can be found at my personal Github : github.com/JoelHernandez343/Cryptography

And inside the .zip file in the drive folder.