Quantum Computation Project

The essence of problems where quantum computers fail to outperform classical computers

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Introduction

- Quantum Computers cannot significantly outperform classical computers while trying to solve **NP-complete** problems.
- Narrowed the problem statement by taking one such problem: 3-colouring graph problem
- Real-life applications of graph colouring
- Analysed the problem via two approaches to see if there was a notable improvement in time complexity



Classical Approach: Backtracking

- For various NP-complete problems, a brute force approach is observed to be the fastest classical algorithm. For some, a refined brute-force techniques like backtracking can be adopted.
- Upper-bound time remains the same but the average time taken reduces
- \bullet Time complexity: $O(3^n)$, where n is the number of vertices

Method:

- Assign colours to each vertex one by one
- Before each assignment, verify if any of the adjacent vertices have the same colour
- If there are no violations proceed else backtrack a step and reassign colours



Grover's Search

Speeds up an unstructured search problem quadratically by using Grover's Amplitude Amplification Trick. Widely used as a sub-routine in many other algorithms to improve computation time.

- ullet Given a list of N items, it achieves the search problem in roughly \sqrt{N} steps, which is significantly faster than the classical approach, which is N steps.
- It takes advantage of the superposition principle of qubits
- It applies the Oracle function and the Diffusion Operator multiple times to the superposed state to amplify the amplitude of the qubit state we are searching for.



Case Study and Comparative Analysis

The 3 colouring problem can be solved via a quantum method by converting it to an equivalent problem that yield the same result.

- Compute the chromatic number and utilise Grover's search
- NK method: generation of superpositions of all possibilities and filtering
- Reducing the problem to a SAT problem

Chromatic number computation

Chromatic number is the minimum number of colours required to colour the graph in such a way that it doesn't violate the adjacent colour node rule.

- A quantum dynamic programming approach using Byskov's algorithm that generates all the Maximal Independent Sets. As Byskov's algorithm is a branching method, we can use Grover's search on it.
- This algorithm computes the chromatic number with a time complexity of O(1.9140ⁿ).
- This is still in exponential time!

NK method

This is a generalised approach to colour an n-vertex graph with k colours.

- Assign K qubits to each node and generate superpositions of all possible colourings
- Filter these possibilities by applying the outer product, horizontal concatenation and element wise product
- Eliminate possibilities that violate the rule
- A row entry should not contain 2 as it implies that two nodes have the same colour
- The time complexity will be $O(N^2)$ where N is the number of qubits.

If we expressed the time complexity in terms of k and n then it would be $O(k^{2n})$, which is still exponential time!

Reduction to SAT problem

- A Reduction-Based Problem Mapping for Quantum Computing relies on the property that NP-Complete problems can be reduced to each other in polynomial time.
- This can be implemented by choosing a particular problem S and transforming the general input in polynomial time and the computed output is converted back to the desired output for the original input.
- The unrestricted Boolean-SAT is chosen as the general problem which can be solved with the help of Grover's Quantum Search Algorithm.



Figure: End-to-End Reduction Framework

Our Approach

- We have built a quantum-based 3-colouring problem solver using Qiskit.
- The solver is capable of deciding the 3-colourability of a graph while generating a possible solution.
- The solver is based on the Reduction to SAT Problem approach. This
 approach was the most alluring since it involved solving the SAT problem.
- There are a total of three major modules: *Graph to SAT converter, Quantum Circuit Generator*, and *SAT solution to Coloured graph converter*.
- The Quantum Circuit Generator is the most important component of the solver. It has two sub-modules: the Phase Oracle Generator and the Diffusion Operator Generator.



Quantum Circuit Generator

- **Step 1:** Qubits are used to represent variables from the SAT problem. Initially, in the circuit Hadamard gates are used to put all the qubits in uniform superposition.
- Step 2 (Phase Oracle): This oracle flips the phase of states based on the boolean function represented by the SAT problem. The oracle is made with a combination of X gates and Toffoli gates.
- Step 3 (Diffusion Operator): The Diffusion operator is build to scale with the number of the variables in the SAT problem. It is composed of X gates, Hadamard gates and Toffoli gates.
- Step 4: The Phase Oracle + Diffusion Operator unit is repeated multiple times (upto $2^{N/2}$ where N is number of variables in the SAT problem).
- **Step 5:** Run the quantum circuit on a quantum computer or a simulator (Eg: AER Matrix Product State Simulator)



Quantum Circuit Example

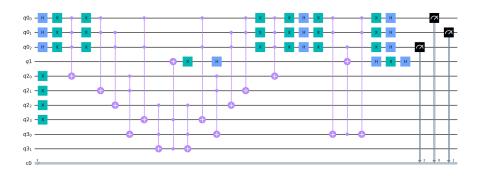
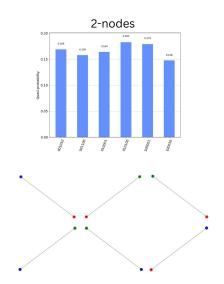
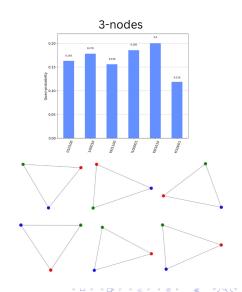


Figure: Quantum circuit generated by our solver for a trivial single node graph



Results with a 2-node and 3-node Graph

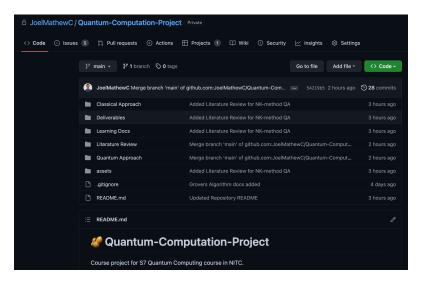




Our Contributions

- Reviewed the most recent literature revolving the topic
- Documented the entire learning process as we realised one of the main challenges in this field is the lack of systematic documentation and comprehensive code structuring
- Implemented all of the components from scratch and without the usage of pre-defined libraries to avoid black-box learning.
- This implementation is built to be scalable and it can be improved upon to generate more efficient solvers.
- Comprehensive analysis to prove the given problem statement

Repository



Conclusion

- By taking an instance of an NP-Complete problem, we have seen that quantum computers do not significantly outperform classical computers.
- Any computational problem that can be solved by a classical computer can
 also be solved by a quantum computer. Conversely, any problem that can be
 solved by a quantum computer can also be solved by a classical computer
 with the principle that sufficient time is given.
- For all applications, a classic computer will not only continue to **co-exist** with quantum computers but also dominate in most of the arenas.
- Though Quantum Computers will not be able to solve all our problems as we expect it to, it is **not** a reason for research to **not** be done in this domain.



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Thank you