



# Lab 2 Help

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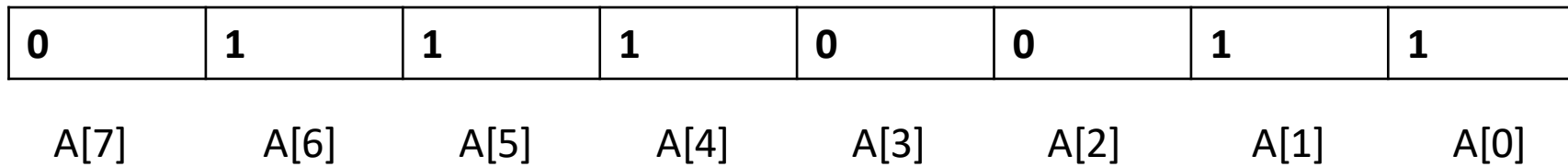
Oklahoma State University

Electrical and Computer Engineering Department

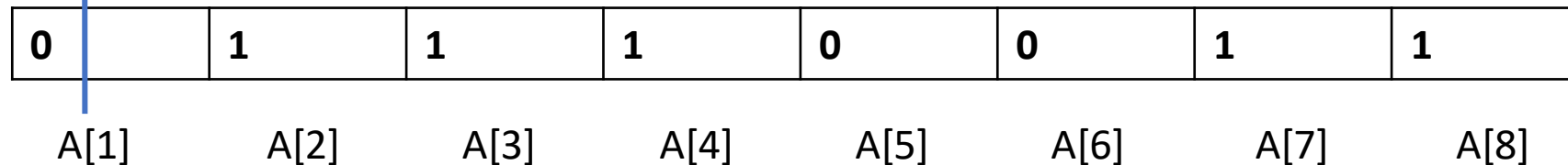
Stillwater, OK 74078 USA

# Bit Numbering Convention

Our convention for an arbitrary 8 bit field:  
Logic [7:0] A = 8'b01110011;



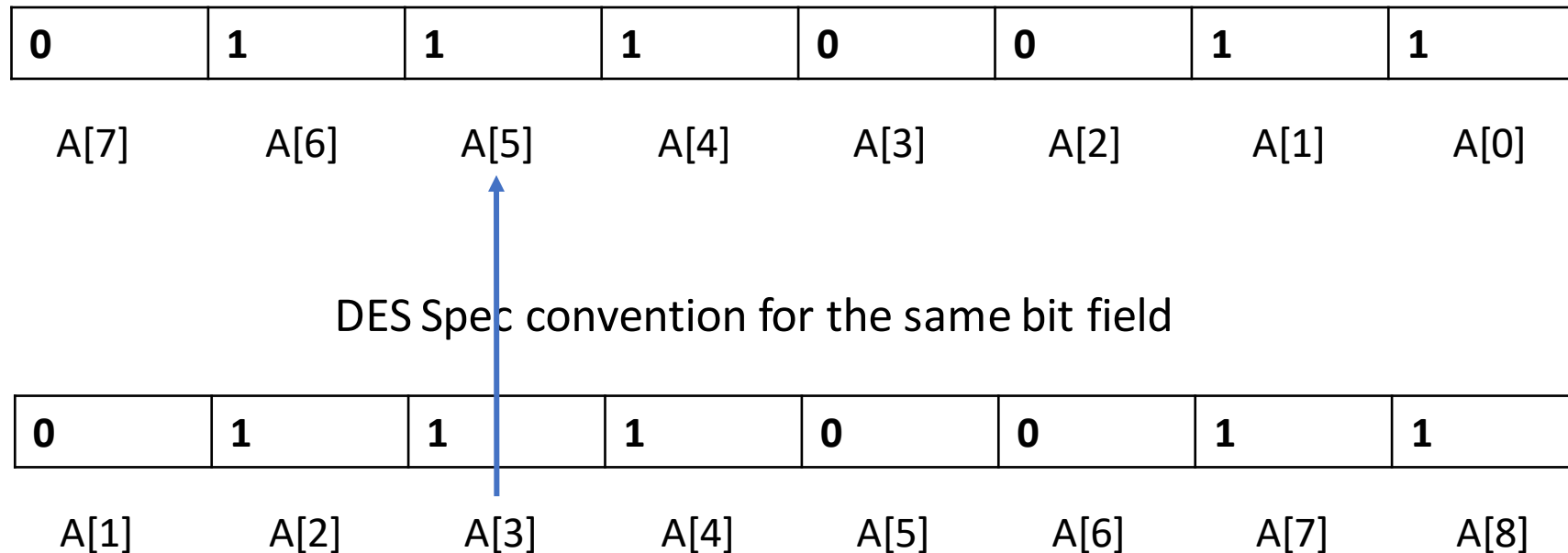
DES Spec convention for the same bit field



# Bit Numbering Convention

Reconciling the different number conventions:

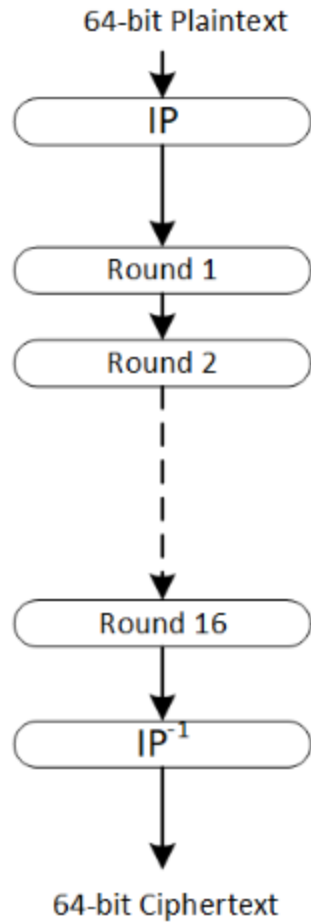
- Bit # in our numbering system = (Max number of the DES numbering) – (DES value we want to convert)



What is A[3] from DES numbering to our numbering scheme?

- Max number = 8, DES Value = 3
- Bit number in our numbering scheme =  $8 - 3 = 5$

# Bit Numbering Convention IP Block



|    |    |    |    |    |    |    |   |
|----|----|----|----|----|----|----|---|
| 58 | 50 | 42 | 34 | 26 | 18 | 10 | 2 |
| 60 | 52 | 44 | 36 | 28 | 20 | 12 | 4 |
| 62 | 54 | 46 | 38 | 30 | 22 | 14 | 6 |
| 64 | 56 | 48 | 40 | 32 | 24 | 16 | 8 |
| 57 | 49 | 41 | 33 | 25 | 17 | 9  | 1 |
| 59 | 51 | 43 | 35 | 27 | 19 | 11 | 3 |
| 61 | 53 | 45 | 37 | 29 | 21 | 13 | 5 |
| 63 | 55 | 47 | 39 | 31 | 23 | 15 | 7 |

Table 5: Initial Permutation (IP) Function

```
assign out_block[63] = inp_block[64-58];
assign out_block[62] = inp_block[64-50];
assign out_block[61] = inp_block[64-42];
assign out_block[60] = inp_block[64-34];
assign out_block[59] = inp_block[64-26];
assign out_block[58] = inp_block[64-18];
assign out_block[57] = inp_block[64-10];
assign out_block[56] = inp_block[64-2];
assign out_block[55] = inp_block[64-60];
assign out_block[54] = inp_block[64-52];
assign out_block[53] = inp_block[64-44];
assign out_block[52] = inp_block[64-36];
assign out_block[51] = inp_block[64-28];
assign out_block[50] = inp_block[64-20];
assign out_block[49] = inp_block[64-12];
assign out_block[48] = inp_block[64-4];
assign out_block[47] = inp_block[64-62];
assign out_block[46] = inp_block[64-54];
assign out_block[45] = inp_block[64-46];
assign out_block[44] = inp_block[64-38];
assign out_block[43] = inp_block[64-30];
assign out_block[42] = inp_block[64-22];
assign out_block[41] = inp_block[64-14];
assign out_block[40] = inp_block[64-6];
assign out_block[39] = inp_block[64-64];
assign out_block[38] = inp_block[64-56];
assign out_block[37] = inp_block[64-48];
assign out_block[36] = inp_block[64-40];
assign out_block[35] = inp_block[64-32];
assign out_block[34] = inp_block[64-24];
assign out_block[33] = inp_block[64-16];
assign out_block[32] = inp_block[64-8];
assign out_block[31] = inp_block[64-57];
assign out_block[30] = inp_block[64-49];
assign out_block[29] = inp_block[64-41];
assign out_block[28] = inp_block[64-33];
assign out_block[27] = inp_block[64-25];
assign out_block[26] = inp_block[64-17];
```

# Where Should I Start?

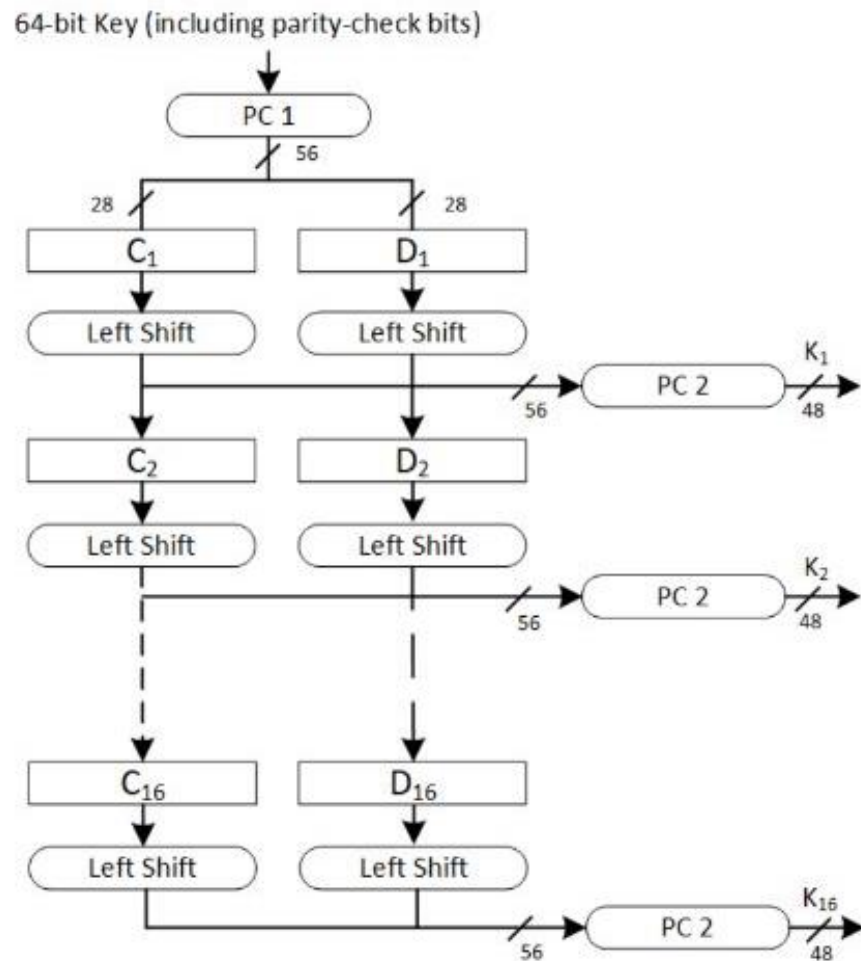
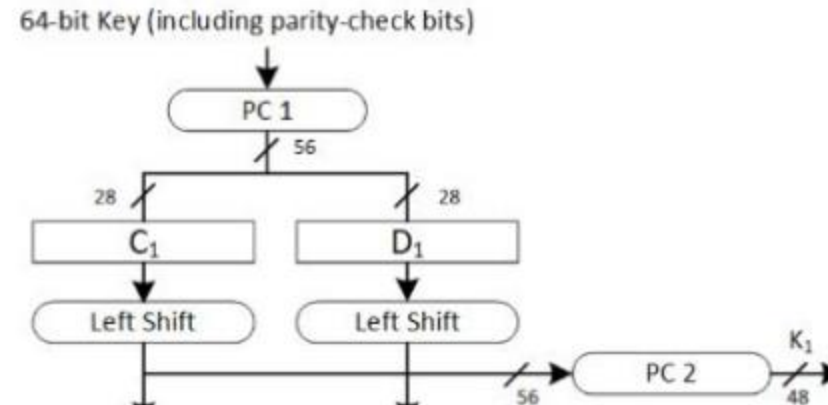


Figure 3: DES SubKey Permutation Diagram



Break anything you do in engineering into bite sized pieces.  
This way we can just repeat what we have done.

First step: Implement PC 1

PC 1 shuffles bits almost exactly like the IP block, except it  
throws away parity bits and splits the result in half

# Where Should I Start?

| Left  |    |    |    |    |    |    |
|-------|----|----|----|----|----|----|
| 57    | 49 | 41 | 33 | 25 | 17 | 9  |
| 1     | 58 | 50 | 42 | 34 | 26 | 18 |
| 10    | 2  | 59 | 51 | 43 | 35 | 27 |
| 19    | 11 | 3  | 60 | 52 | 44 | 36 |
| Right |    |    |    |    |    |    |
| 63    | 55 | 47 | 39 | 31 | 23 | 15 |
| 7     | 62 | 54 | 46 | 38 | 30 | 22 |
| 14    | 6  | 61 | 53 | 45 | 37 | 29 |
| 21    | 13 | 5  | 28 | 20 | 12 | 4  |

```
module PC1 (key, left_block, right_block);  
  
  input logic [63:0] key;  
  output logic [27:0] left_block;  
  output logic [27:0] right_block;  
  
  assign left_block[27] = key[64-57];  
  assign left_block[26] = key[64-49];  
  
  //fill in the rest  
  
  assign right_block[27] = key[64-63];  
  assign right_block[26] = key[64-55];
```

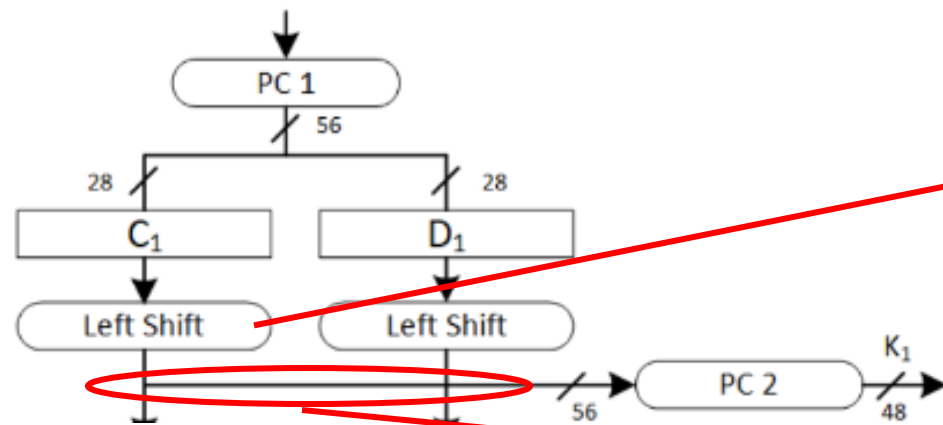
Table 3: Permutation Choice 1 (PC-1) Function<sup>2</sup>

|    |    |    |    |    |    |    |    |
|----|----|----|----|----|----|----|----|
| 14 | 17 | 11 | 24 | 1  | 5  | 3  | 28 |
| 15 | 6  | 21 | 10 | 23 | 19 | 12 | 4  |
| 26 | 8  | 16 | 7  | 27 | 20 | 13 | 2  |
| 41 | 52 | 31 | 37 | 47 | 55 | 30 | 40 |
| 51 | 45 | 33 | 48 | 44 | 49 | 39 | 56 |
| 34 | 53 | 46 | 42 | 50 | 36 | 29 | 32 |

Table 4: Permutation Choice 2 (PC-2) Bit Function

# Left Shifting and Concatenation

64-bit Key (including parity-check bits)

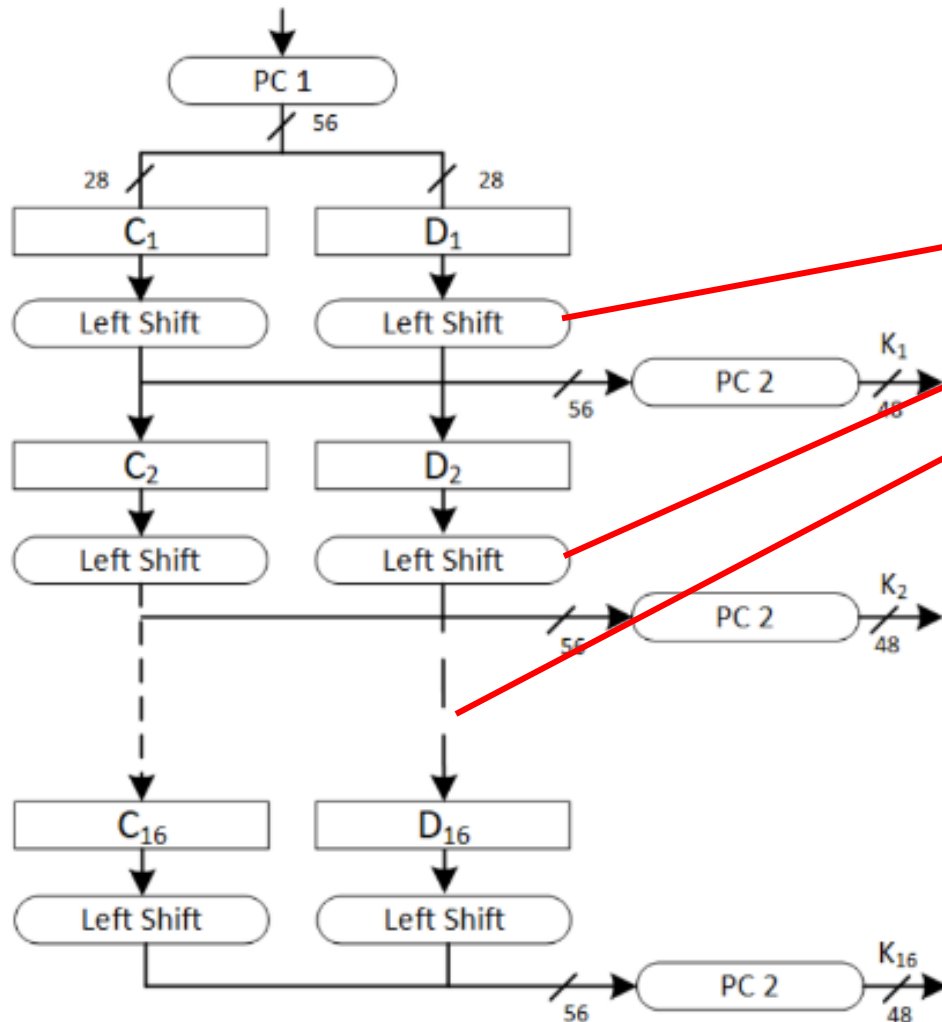


ShiftedC1 = {C1[26:0],C1[27]} //1 bit left circular shift

logic [55:0] Concatenated;  
Concatenated = {C1,D1}

# Extend to All Rounds

64-bit Key (including parity-check bits)



| Iteration # | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 |
|-------------|---|---|---|---|---|---|---|---|---|----|----|----|----|----|----|----|
| Left Shifts | 1 | 1 | 2 | 2 | 2 | 2 | 2 | 2 | 1 | 2  | 2  | 2  | 2  | 2  | 2  | 1  |

Table 2: Left Rotation within SubKey Generation

Shifted = {C<sub>3</sub>[25:0], C<sub>3</sub>[27:26]} // 2 bit left circular shift

Note: Remember that the input to the next round is the shifted values from the previous round NOT the original output from PC1

Figure 3: DES SubKey Permutation Diagram



# How do we check our results?

```
PS C:\Users\amerigo\Documents\GitHub\dldLab2_Spring23\Lab2\java> javac -d . -classpath . DES.java
PS C:\Users\amerigo\Documents\GitHub\dldLab2_Spring23\Lab2\java> java DES
Original plain Text: 2579DB866C0F528C
Key: 433E4529462A4A62

Encryption:

After initial permutation: 5646B9278C13B66C
After splitting: L0=5646B927 R0=8C13B66C

Round 1 8C13B66C F3EFC169 208066A253BA
Round 2 F3EFC169 25DAF255 C0B6508F6DC2
Round 3 25DAF255 1890CFBF 44D6422CC355
Round 4 1890CFBF AFB98FA0 62D142D3C4C6
Round 5 AFB98FA0 8F76DBD7 28C143CC8789
Round 6 8F76DBD7 C176D0E5 21411B9A764D
Round 7 C176D0E5 C7401A8C 2501917AD3A0
Round 8 C7401A8C B748825A 170891906D2B
Round 9 B748825A 61239171 084949255DD5
Round 10 61239171 FE28B577 01690D8B80F3
Round 11 FE28B577 CDB650DE 012D81C7CF05
Round 12 CDB650DE 8B8270E5 512CA11A07DC
Round 13 8B8270E5 DDBEE19 D1A480D9D185
Round 14 DDBEE19 5F82D63F 5086864266A9
Round 15 5F82D63F B35B4964 709006FA390D
Round 16 B35B4964 850AC7BE C03E202F8437

Cipher Text: FCB54739A1832EC5
```

Original text (unencrypted)

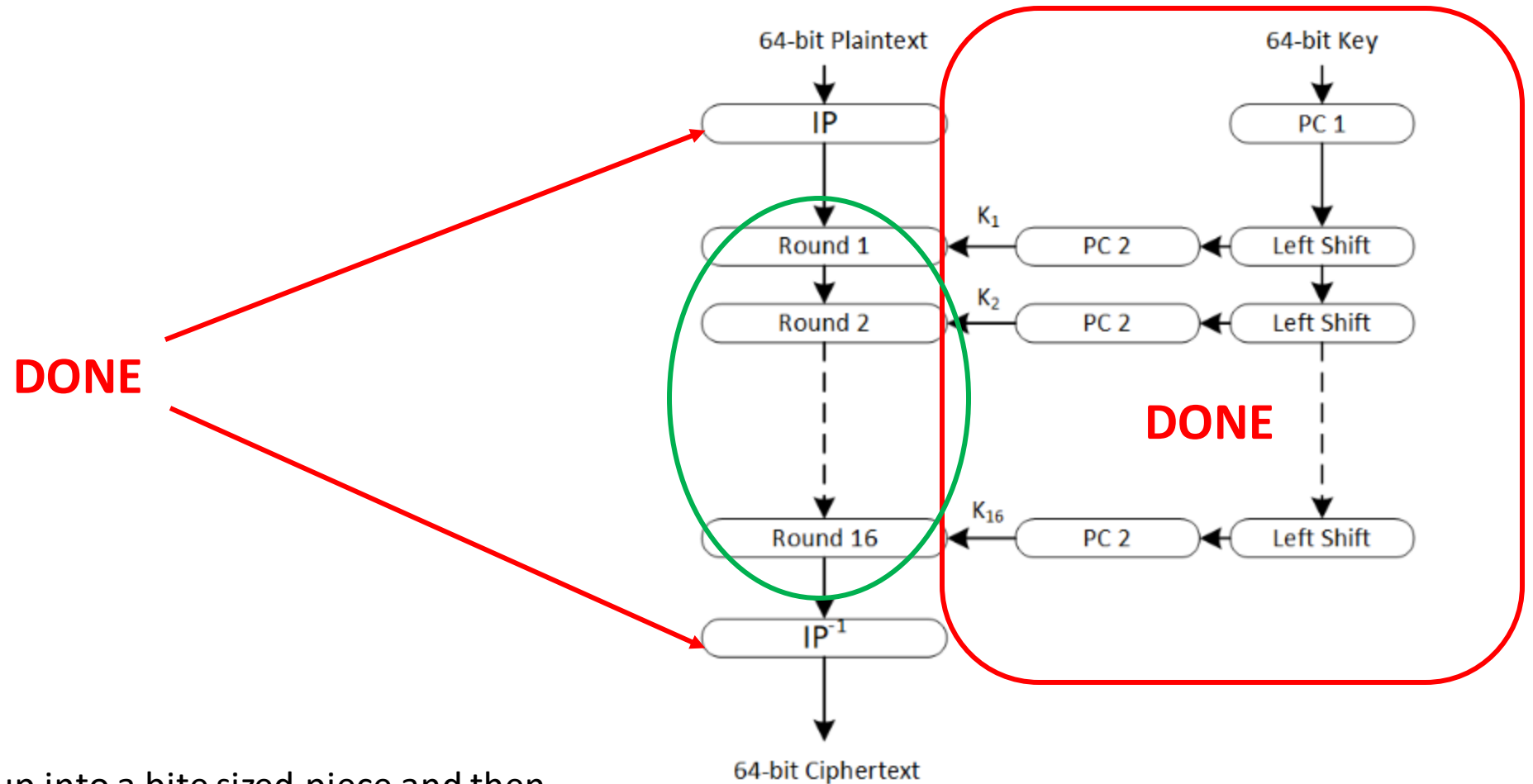
64 bit Key (in Hex) to send to  
GenerateKeys()

Output of IP Block

C1, D1, and SubKey1

Encrypted Output

# Encryption Using the SubKeys



Rounds: Break up into a bite sized piece and then instantiate it 16 times

Figure 2: DES Block Diagram

# Round

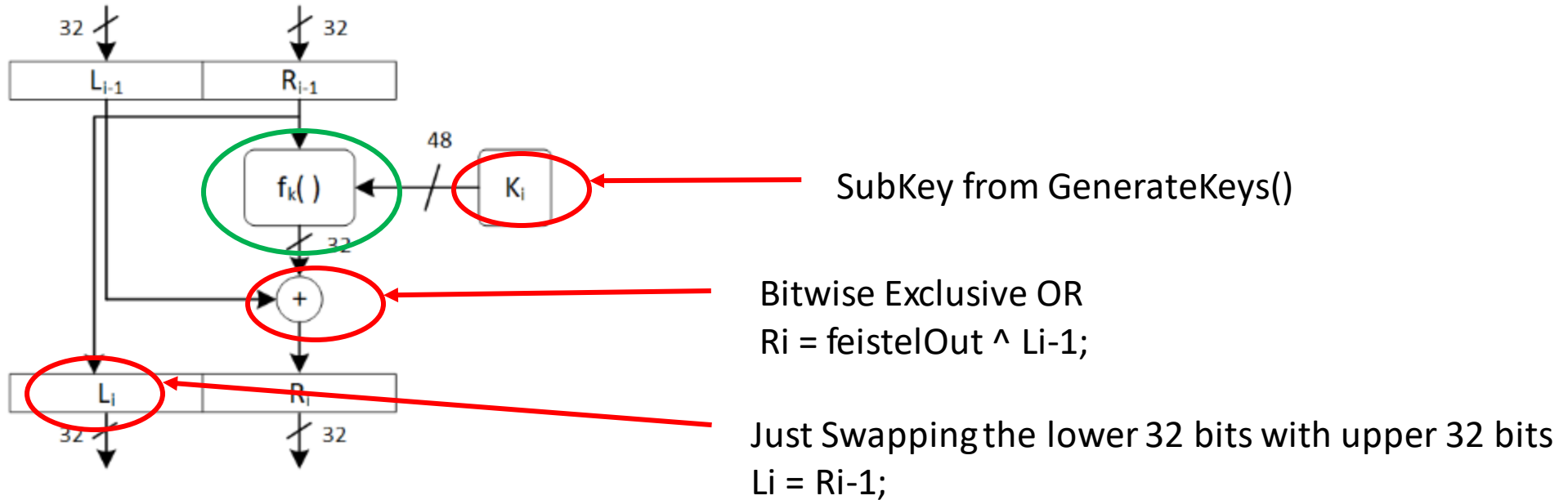
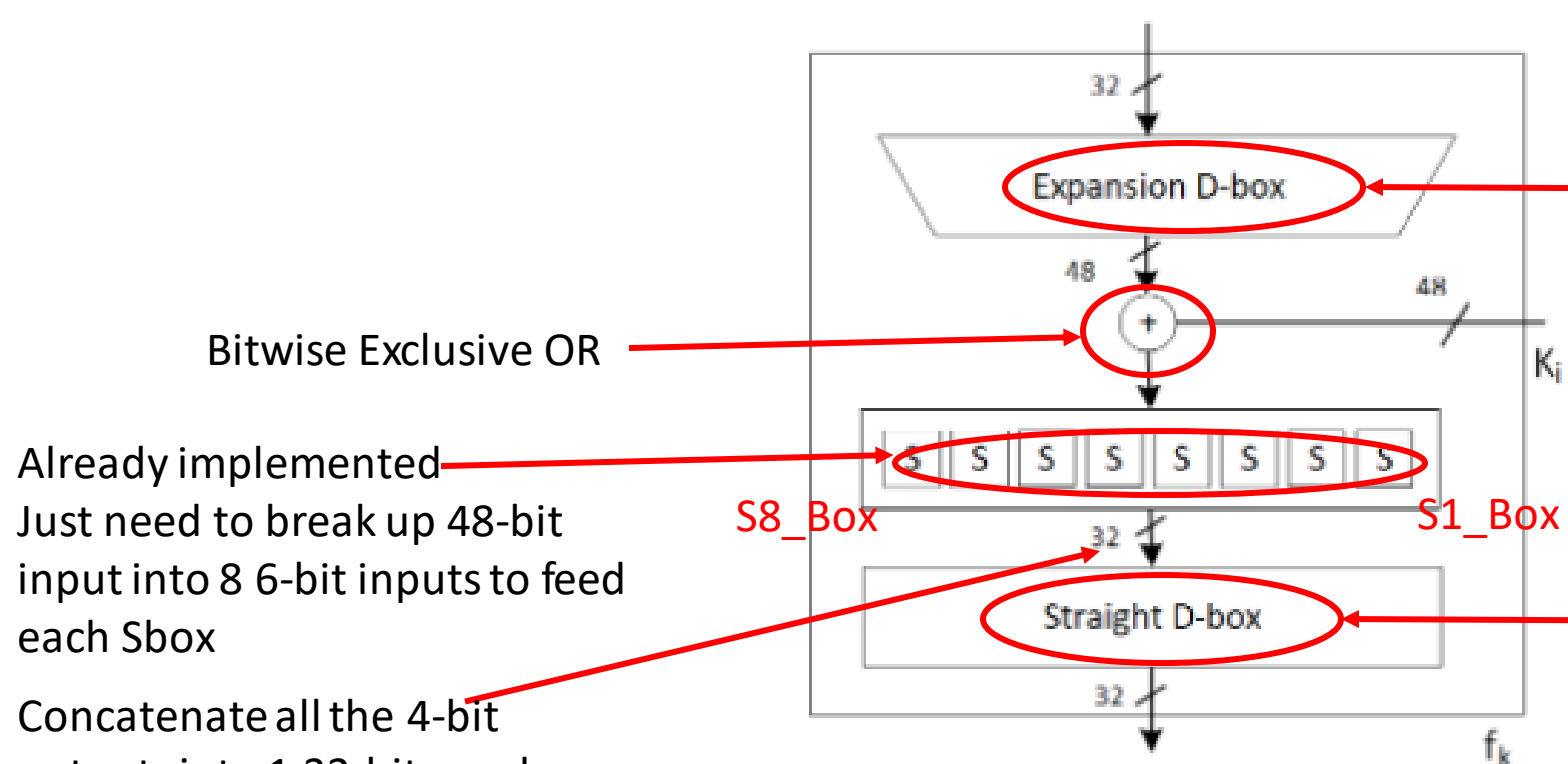


Figure 5: DES Round Block Diagram

# Feistel Block



Already implemented  
Just need to break up 48-bit  
input into 8 6-bit inputs to feed  
each Sbox

Concatenate all the 4-bit  
outputs into 1 32-bit number:  
StraightIn = {S8\_out, S7\_out,  
..., S1\_out};

|    |    |    |    |    |    |
|----|----|----|----|----|----|
| 32 | 1  | 2  | 3  | 4  | 5  |
| 4  | 5  | 6  | 7  | 8  | 9  |
| 8  | 9  | 10 | 11 | 12 | 13 |
| 12 | 13 | 14 | 15 | 16 | 17 |
| 16 | 17 | 18 | 19 | 20 | 21 |
| 20 | 21 | 22 | 23 | 24 | 25 |
| 24 | 25 | 26 | 27 | 28 | 29 |
| 28 | 29 | 30 | 31 | 32 | 1  |

Table 7: Expansion Function or Permutation (EP) Function

|    |    |    |    |    |    |    |    |
|----|----|----|----|----|----|----|----|
| 16 | 7  | 20 | 21 | 29 | 12 | 28 | 17 |
| 1  | 15 | 23 | 26 | 5  | 18 | 31 | 10 |
| 2  | 8  | 24 | 14 | 32 | 27 | 3  | 9  |
| 19 | 13 | 30 | 6  | 22 | 11 | 4  | 25 |

Table 8: Straight Diffusion-box (D-box) Function

Figure 6: Feistel ( $f_K$ ) Block Diagram

# Nearing the end

- Repeat for all 16 rounds
- Output of 16th round should be the input of FP()
- Output of FP will be the encrypted data
- Create an input to your system which dictates if you are encrypting or decrypting. Modify your system so that you can encrypt or decrypt based on this signal.
  - Hint1: the difference between encrypting and decrypting is just reversing the order of the subkeys.
  - Hint2: the ternary operator can operate on multiple bitfields at once if you use concatenation:
    - `assign {resultA, resultB} = decide ? {option1A, option1B} : {option2A, option2B};`

# Extra Credit – DES in CBC Mode

There several modes in DES encryption, implement a switch input that selects between ECB (Electronic Code Book - Default) and CBC (Cipher Block Chaining) modes.

- CBC encryption process also requires a 64-bit Initialization Vector (IV), otherwise the process is the same

The definitive step in CBC is to XOR the plaintext block with an IV – a unique, fixed-length conversion function – to create a random, or pseudorandom, output. That is, it is simply that the plaintext is XOR'ed with the Initialization Vector (IV) right in the beginning of the process.

$$Plaintext_{new} = IV \oplus Plaintext_{input}$$

- The key to this is that the IV must be processed in the beginning during encryption and towards the end in decryption.

