

Objects and OOP

CS 350

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Concepts of Objects

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- To add a new variant, we needed to refactor *every single* definition that uses type-case to have a new case for the new variant

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 - Since you have to add the new method

OOP vs Algebraic Datatypes

	Add new variant	Add new operation
Datatypes	Needs global refactoring	Local additions only
Objects	Local additions only	Needs global refactoring

Curly-Obj-Immut

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- Like Tuples in Python

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 - Name comes from Smalltalk, history of OOP
- Self-reference: `this`
 - Inside an object's methods, there is a special variable called `this`
 - Refers to the object that the method is being called on, so we can get e.g. field values from methods

Example

Example

```
{letvar mkCircle
  {fun {r}
    {object {radius r}
      ;; All methods take one parameter, so we just ignore it
      {getArea {x} {* {get this radius}
                     {get this radius} 3.14}}}}}
{letvar unitCircle {mkCircle 1}
  {send unitCircle getArea 0}}}
```

- Produces 3.14

Implementing Objects

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General Strategy

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 - Just make new `Expr` and `Value` variants with data for object-fields and methods
 - Store fields as values
 - Store methods as closures
- Tricky bit: making sure there's a value for `this` in scope for method calls
 - Making sure it's the *right* value

New Expression Variants

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(define-type Expr
  ...
  (Object [fields : (Listof (Symbol * Expr))]
          [methods : (Listof (Symbol * (Symbol * Expr)))])
  (Get [obj-expr : Expr]
        [field-name : Symbol])
  (Send [obj-expr : Expr]
         [method-name : Symbol]
         [arg-expr : Expr]))
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- Nothing fancy, just the literal tree representation of the previous syntax

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 - Send has expression for the object whose method we're calling, the name of the method, and an expression for the argument

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- Value version of object
 - Fields: just like in Expr, except each name has a value, not an expression
 - Methods: list of name-value pairs, where each value is assumed to be a closure

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- Assignment 6 will be integrating Objects and Stores
 - Building a small language like Python or JavaScript

Interpreting Object Creation

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```
(define (interp [env : Env]
               [e : Expr] : Value
  (type-case Expr e
    [(Object fields methods)
     ;; Each named field expression gets turned into a name-value pair
     (ObjV (map (lambda ([pr : (Symbol * Expr)]) (pair (fst pr) (interp (snd pr) env)))
               ;; Each method gets turned into a closure
               (lambda ([pr : (Symbol * (Symbol * Expr))])
                 (pair (fst pr) (ClosureV (fst (snd pr))
                                           (lambda () (interp (snd (snd pr)) env))))))
           ]))
```

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```
(define (find [l : (Listof (Symbol * 'a))] [name : Symbol]) : 'a
  (type-case (Listof (Symbol * 'a)) l
    [empty
     (error 'find (string-append "not found: " (symbol->string name)))]
    [(cons p rst-l)
     (if (symbol=? (fst p) name)
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 - Works for fields and methods

Interpreting Field Lookup

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```
[(getE obj-expr field-name)
 (type-case Value (interp env obj-expr)
  [(ObjV fields methods)
   (find fields field-name)]
  [else (error 'interp "not an object")])])]
```

Interpreting Method Calls

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```
(Send obj-expr method-name arg-expr)
  (let [(obj-val (interp env obj-expr))
        (arg-val (interp env arg-expr))]
    (type-case Value obj-val
      [(ObjV fields methods)
       (let ([param-body (find methods method-name)])
         (interp (extendEnv (bind 'this obj-val)
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 - The entire object's value bound to 'this

Passing the self reference

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- Like `self` in Python, `this` in Java/C++/JS

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{letvar empty {object {{length 0}}
                      {{head {x} errorNoEmptyListHead}
                      {tail {x} this}}} ;; Note the self reference
{letvar cons {fun {h t}
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                       {{head {x} h}
                       {tail {x} t}}}}}
}}
```

- Instead of having a type with two variants, each list carries its own information
 - Its length (e.g. whether empty or not)
 - Its head (errors if empty)
 - Its tail (empty if empty)
- Can do `{get someList length}` or `{send someList tail 0}` on empty or cons, and will work in either case

Example: OOP Lists

```
{letvar empty {object {{length 0}}
                      {{head {x} errorNoEmptyListHead}
                      {{tail {x} this}}}} ;; Note the self reference
{letvar cons {fun {h t}
               {object {length {+ 1 {get t length}}}
                       {{head {x} h}
                       {{tail {x} t}}}}}
}}
```

- Instead of having a type with two variants, each list carries its own information
 - Its length (e.g. whether empty or not)
 - Its head (errors if empty)
 - Its tail (empty if empty)
- Can do {get someList length} or {send someList tail 0} on empty or cons, and will work in either case
 - Return of methods carried around with the list

Example: Functions as Objects

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```
{object {} {call {x} body}}
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- Object with a single “call” method

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```

- Can even do recursion

```
;; factorial with objects  
{define fact  
  {object {}  
    {call {x}  
      {if0 x  
        1  
        {* x {send this call {- x 1}}}}}}}}
```