

Generative Recursion and Tail Recursion

CS 350

Dr. Joseph Eremondi

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Broad Goals

- Objectives

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```
int x = startVal;  
for (int i = 0; i < n; i++)  
{  
    x = f(x);  
}
```

Example: Reversing a List The Naive Way

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- $O(n^2)$: Each append has to walk through the whole list

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- When reach the end, have a reversed list

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- Helper function

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 - Pass the updated value as the new accumulator