A CPS Transformation for Gradual Programs with Evidence

CPSC 539B Final Project

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Source Language

Term Syntax



e ::=

x Variablesb Booleans

n Natural Numbers $\lambda x: T.e$ Functions

 $e_1 \ e_2$ Function Application

 $e_1 + e_2$ Addition

 $e_1\stackrel{?}{=}e_2$ Number Equality Test

if e_1 **then** e_2 **else** e_3 Conditionals

 $\langle e_1, e_2 \rangle$ Tuples

 $\pi_1 e$ Tuple First Projection

 $\pi_2 e$ Tuple Second Projection

 ε e Evidence Ascription

error Runtime Type Error

Type Syntax



$$T$$
 ::= Types \mid Nat \mid Bool \mid $T_1
ightarrow T_2 \mid$ \mid $T_1 imes T_2 \mid$ \mid ?

Type Rules



HASTYPEASCR $\Gamma \vdash e : T_2$ $\varepsilon \vdash T_1 \cong T_2$ $\Gamma \vdash \varepsilon e : T_1$

Consistent EV
$$T_3 \sqcap T_1 = T_3$$

$$T_3 \sqcap T_2 = T_3$$

$$\{T_3\} \vdash T_1 \cong T_2$$

Combining Evidence



$$T_1 \sqcap T_2 = T_3$$

(Precision Meet)

$$\frac{\text{MEETDYNL}}{7 \sqcap T = T}$$

$${\rm MeetRefl}$$

$$\overline{T \cap ? = T}$$

$$\overline{T \cap T = T}$$

MEETFUN

$$\frac{T_1 \sqcap T_1' = T_1''}{T_2 \sqcap T_2' = T_2''}$$

$$\frac{T_1 \to T_2 \sqcap T_1' \to T_2' = T_1'' \to T_2''}{T_1 \to T_2 \sqcap T_1' \to T_2''}$$

$$T_{1} \sqcap T'_{1} = T''_{1}$$

$$T_{2} \sqcap T'_{2} = T''_{2}$$

$$T_{1} \times T_{2} \sqcap T'_{1} \times T'_{2} = T''_{1} \times T''_{2}$$



RedAscr

$$\overline{\varepsilon_1(\varepsilon_2 e) \longrightarrow (\varepsilon_1 \sqcap \varepsilon_2) e}$$

RedAscrFail

$$\varepsilon_1 \sqcap \varepsilon_2$$
 undefined

$$\varepsilon_1 \left(\varepsilon_2 \, e \right) \longrightarrow \mathsf{error}$$

REDAPPEV

$$(\varepsilon_1 (\lambda x : T.e)) (\varepsilon_2 r) \longrightarrow (\operatorname{cod} \varepsilon_2) ([(\operatorname{dom} \varepsilon_1 \sqcap \varepsilon_2) r/x]e)$$

REDAPPEVFAIL

$$\operatorname{dom} \varepsilon_1 \sqcap \varepsilon_2 \operatorname{undefined}$$

$$(\varepsilon_1 (\lambda x : T.e))(\varepsilon_2 r) \longrightarrow error$$

Examples



```
{Nat}({Bool} true) + 0 typechecks!
{Bool} \vdash Bool \cong ? and {Nat} \vdash ? \cong Bool
But: fails at runtime!
```

The Target

Simplified λ^K



```
u, k ::=
                     X
                                             t ::=
                     n
                      b
                                                              V
                     \mathbf{fix} \times \mathbf{u}
                                                              let d in t
                     \lambda x_1 \dots x_i. t
                                                              u(arg)
                    \langle u_1, u_2 \rangle
                                                              if u then t_1 else t_2
                                                              halt [u]
d
                                                              error
                x := u
               x := \pi_1 u
                                             arg
                                                         u_1, \ldots, u_i
               x := \pi_2 u
              x := u_1 + u_2
              x := u_1 \stackrel{?}{=} u_2
```

The Translation

Translating Evidence



Integer constants BOOL, NAT, ARROW, PRODUCT, DYN

$$\llbracket \varepsilon \rrbracket = u$$

(CPS Representation of Runtime Evidence)

EVTRANSFORMBOOL

EVTRANSFORMNAT

EVTRANSFORMDYN

$$[\![\{\mathsf{Bool}\}]\!] = \langle \mathtt{BOOL}, \mathsf{0} \rangle \qquad [\![\{\mathsf{Nat}\}]\!] = \langle \mathtt{NAT}, \mathsf{0} \rangle \qquad [\![\{?\}]\!] = \langle \mathtt{DYN}, \mathsf{0} \rangle$$

$$\overline{[\![\{\mathsf{Nat}\}]\!] = \langle \mathtt{NAT}, \mathsf{0} \rangle}$$

$$\boxed{\llbracket\{?\}\rrbracket=\langle \mathtt{DYN},0\rangle}$$

EVTRANSFORMARR.

$$[\![\{T_1 \rightarrow T_2\}]\!] = \langle \mathtt{ARROW}, \langle [\![\{T_1\}]\!], [\![\{T_2\}]\!] \rangle \rangle$$

EVTRANSFORMPROD

$$[\![\{\mathit{T}_1 \rightarrow \mathit{T}_2\}]\!] = \langle \mathtt{PRODUCT}, \langle [\![\{\mathit{T}_1\}]\!], [\![\{\mathit{T}_2\}]\!] \rangle \rangle$$

Evidence Operations



 $MEET(u_1, u_2, k)$

Combines evidence representation u_1 and u_2 , gives result to k

Passes control error continuation if meet undefined

Similar for DOM, COD to decompose function types

Translating Values



$$\llbracket v \rrbracket = u$$

(CPS Translation of Closed Values)

VALTRANSFORMBOOL

$$\overline{\|b\|} = \langle \text{DYN}, b \rangle$$

$$[\![n]\!] = \langle \mathtt{DYN}, n \rangle$$

VALTRANSFORMFUN

$$[e]c = t$$

$$[\![\lambda x : T.e]\!] = \langle \mathtt{DYN}, \lambda x c.t \rangle$$

$$\llbracket \langle \mathbf{v}_1, \mathbf{v}_2 \rangle \rrbracket = \langle \mathtt{DYN}, \langle \llbracket \mathbf{v}_1 \rrbracket, \llbracket \mathbf{v}_2 \rrbracket \rangle \rangle$$

ValTransformEv

$$[r] = \langle \mathtt{DYN}, u \rangle$$

$$\llbracket \varepsilon r \rrbracket = \langle \llbracket \varepsilon \rrbracket, u \rangle$$

Translating Applications



TRANSFORMAPP

$$k_1 := (\lambda x_2. \operatorname{let} y_1 := \pi_1 x_1 \operatorname{in} \operatorname{let} z_1 := \pi_2 x_1 \operatorname{in} \operatorname{let} y_2 := \pi_1 x_2 \operatorname{in} \operatorname{let} z_2 := \pi_2 x_2 \operatorname{in} t_1 \\ t_1 := \operatorname{DOM}(y_1, \lambda y_1'. \operatorname{COD}(y_1, \lambda y_1''. \operatorname{MEET}(y_1', y_2, (\lambda y_3. z_1(\langle y_3, z_2 \rangle, (\lambda z_3. t_2)))))) \\ t_2 := \operatorname{let} z_3' := \pi_1 z_3 \operatorname{in} \operatorname{let} z_3'' := \pi_2 z_3 \operatorname{in} \operatorname{MEET}(y_1'', z_3', (\lambda z_4. k(\langle z_4, z_3'' \rangle))) \\ \mathbb{[e_2]} k_1 = t' \\ \mathbb{[e_1]} (\lambda x_1. t') = t \\ \mathbb{[e_1 e_2]} k = t$$

Correctness

Incorrectness