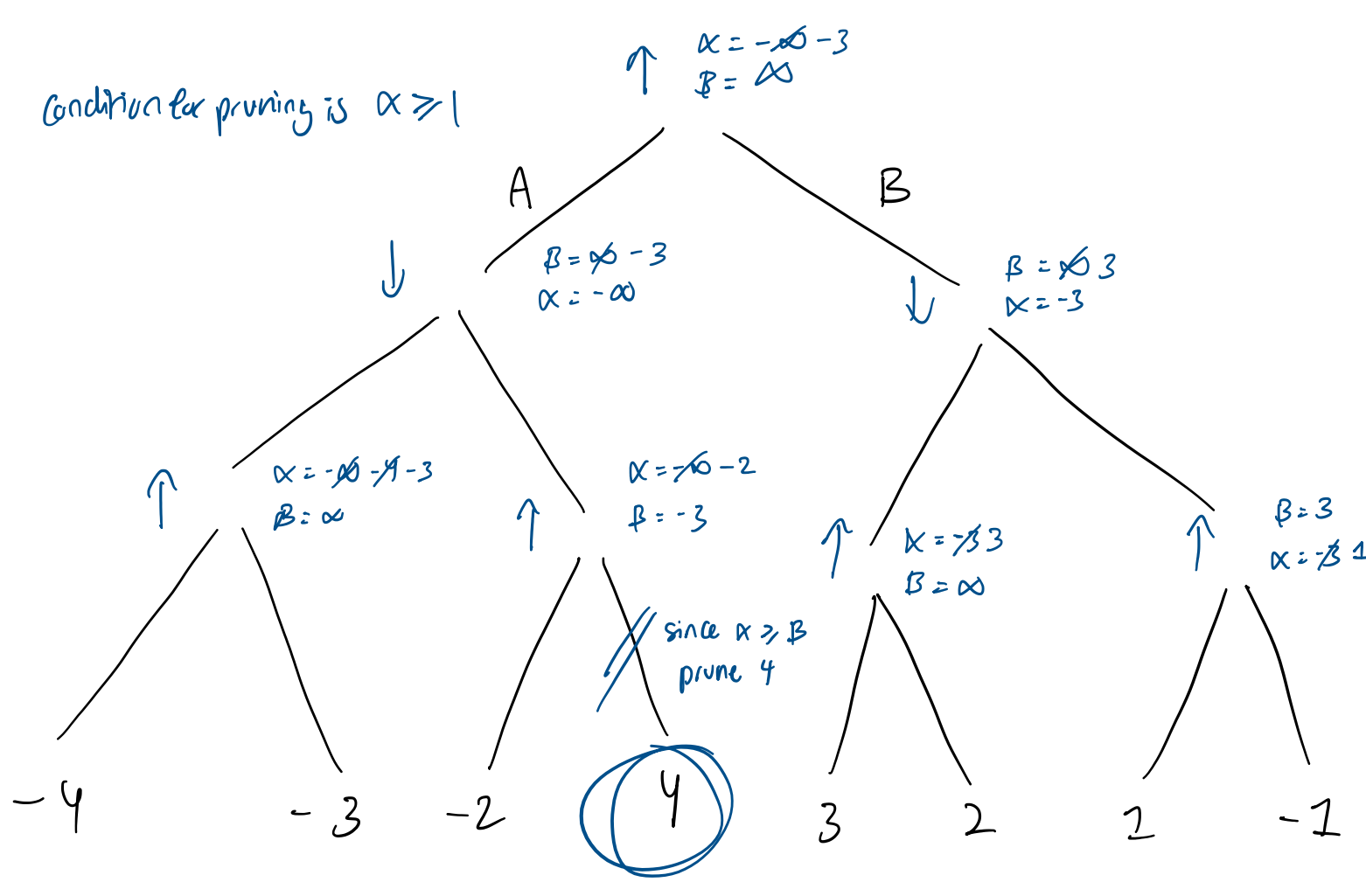


1c) option B is optimal for player 1 as  $1 > -3$  and root node is a max node so player will choose higher of two values therefore option B is optimal for player 1

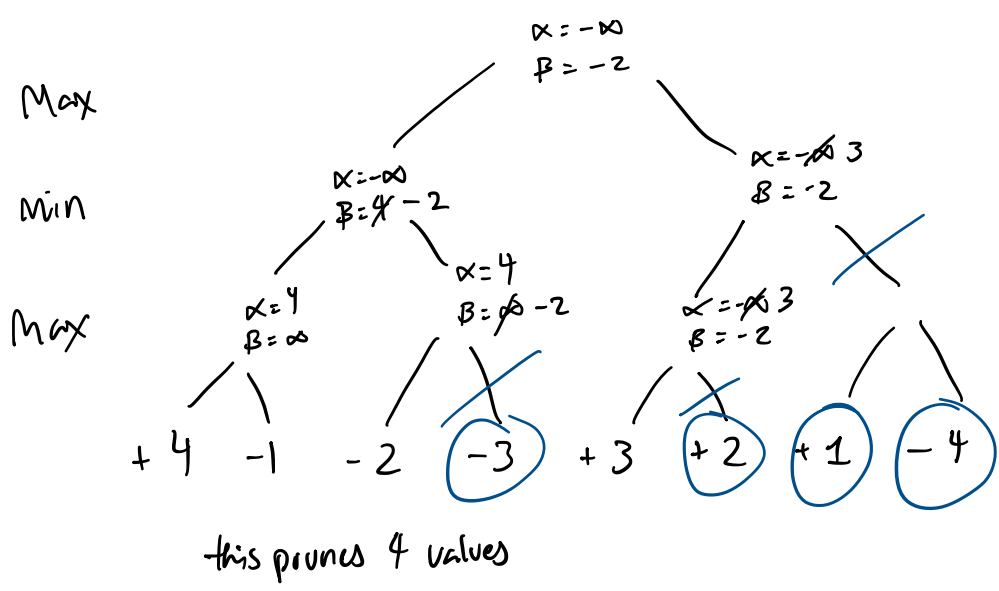
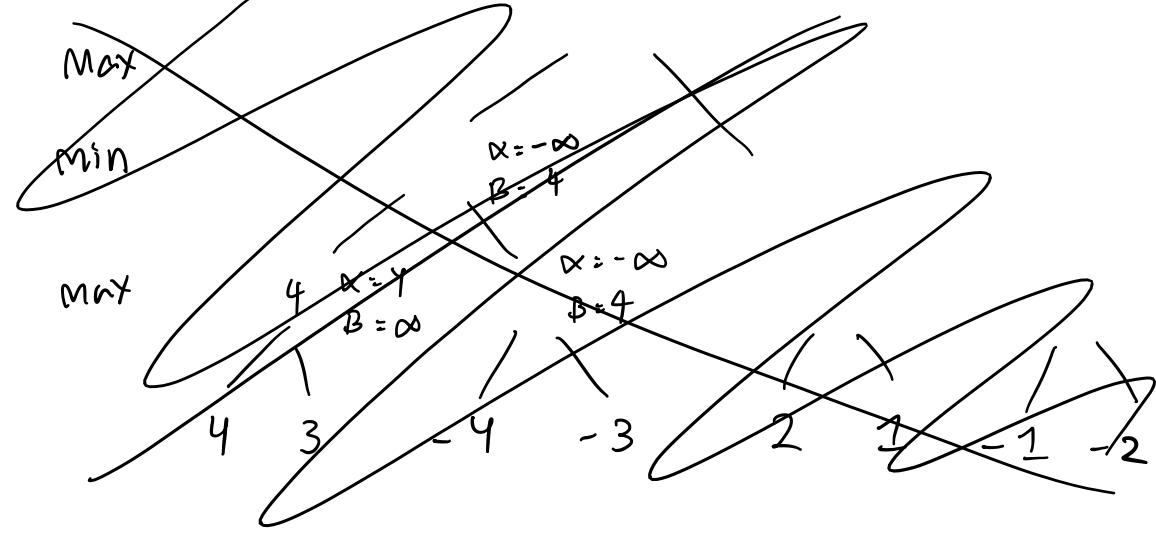
1d) expected outcome is minimum value at root node which is 1

1e) For  $\alpha = -\infty$  and  $\beta = \infty$

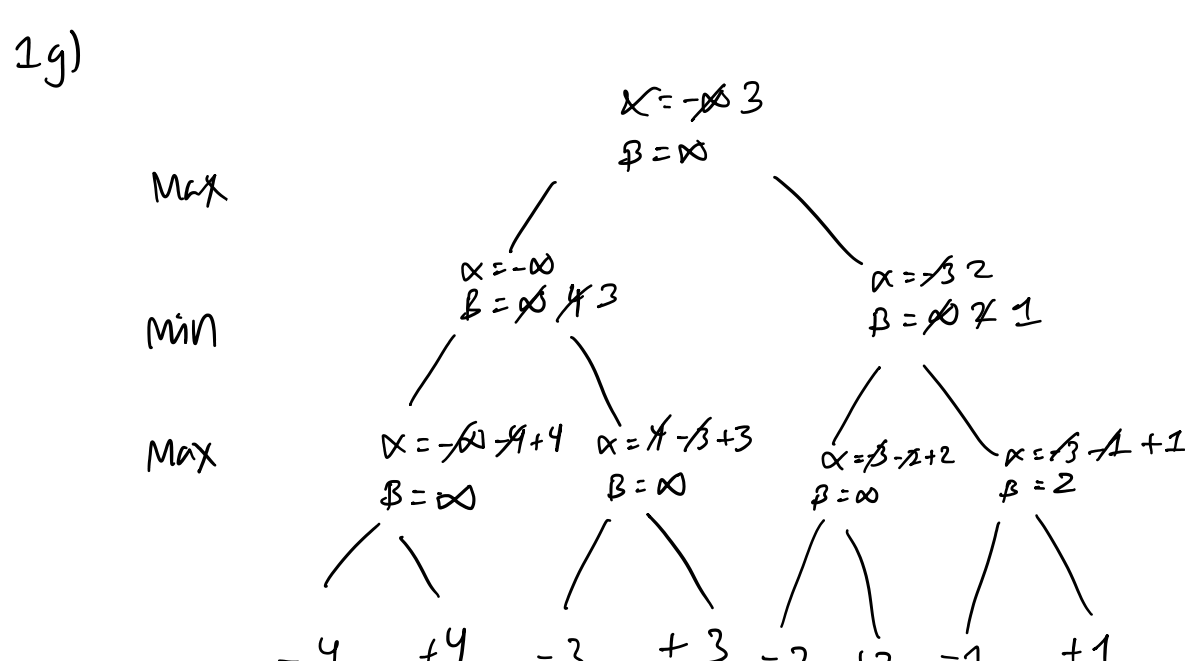


1f) For max nodes put best values early & for min nodes put worst values early

For max pruning try and have max values on left & min on right best case is branch A right & B right can be pruned



this prunes 4 values



2) a)

Variables

P1 → person in position 1

P2 → person in position 2

P3 → person in position 3

Domain

each var can take any 3 bikers {A, B, C}

but no biker can be repeated

constraints

1) each person needs appear at least once

and  $P1 \neq P2 \neq P3$

2) Adjacent degrees must differ

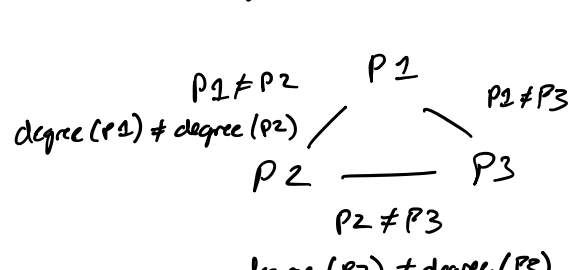
$\text{degree}(P1) \neq \text{degree}(P2)$

$\text{degree}(P2) \neq \text{degree}(P3)$

3)  $\text{Carol} \neq P3$

b) Nodes = variables

Edges = constraints



c) Try  $P1 = A$  ✓

→ Try  $P2 = A$  × same as P1

Try  $P2 = B$  → degree constraint met ✓

→ Arrows indicate backtracking

→ Try  $P3 = A$  × same as P1

→ Try  $P3 = B$  × same as P2

→ Try  $P3 = C$  × Carol is last violates constraint

→ Try  $P2 = C$  × violates degree constraint

→ Try  $P1 = B$  ✓

Try  $P2 = A$  ✓

→ Try  $P3 = C$  × Carol last violates constraint

→ Try  $P2 = B$  × ~~violates~~  $P1 \neq P2$

Try  $P2 = C$  ✓

→ Try  $P3 = A$  ×  $\text{degree}(P2) \neq \text{degree}(P3)$

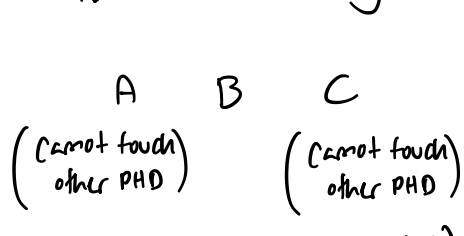
→ Try  $P1 = C$  ✓

→ Try  $P2 = A$  ×  $\text{degree}(P1) \neq \text{degree}(P2)$

Try  $P2 = B$  ✓

Try  $P3 = A$  ✓

d) MRPV = min remaining value



P1 P2 P3

A cannot go P2

B can go anywhere

C cannot go P2 or P3

So start Assigning one with most constraints

Try  $P1 = C$

remaining domains

$P2 \in \{A, B\}$

$P3 \in \{A, B\}$

Try  $P2 = A$  (A has less options left) × degree constraint

✓ Try  $P2 = B$  ✓

Try  $P3 = A$