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1<sup>st</sup> Given Name Surname dept. name of organization (of Aff.) name of organization (of Aff.) City, Country email address 2<sup>nd</sup> Given Name Surname dept. name of organization (of Aff.) name of organization (of Aff.) City, Country email address 3<sup>rd</sup> Given Name Surname dept. name of organization (of Aff.) name of organization (of Aff.) City, Country email address

Abstract—Machine to Machine (M2M) communication has gained much interest in the recent past and the issues related to energy efficiency are central to all wireless networks, including wireless sensor, for an energy economy, we need to limit emissions so some nodes will make requests to their neighbors. This article will show how to make sure neighbors are woken up during queries in a hybrid network architecture or how to make sure that the nodes that will receive the information will be woken up?

Index Terms—M2M, hybrid network, hybrid network, wifi, adhoc

#### I. INTRODUCTION

### II. OPERATION

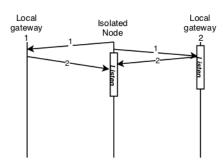


Figure 1. Discover Phase

Phase 1: Discovery Isolated nodes broadcast messages to discover(message 1) a Local Gateway. The end-nodes send messages every (X + random Integer) seconds. If a Gateway receive a discovery message, the Local Gateway will accept (message 2) and communicate to the end-nodes the communication slots.

### Phase 2: Registration

The end-node will confirm the pairing to the Local gateway (message 3). The end-node is registered to the Local gateway. There might be several Local gateway inside the system. However, there must never be multiple registrations on the same one. After the local gateway receive the message 3, gateway send to the isolates node the message 4 to give his acknowledgment. Message 5 is the first message from the isolated node, which contains datas to put an end the pairing phase. When all this stuff is done, the local gateway is able

Isolated Local gateway
Node 2

4

5

6

7

Figure 2. Registering/pairing phase

to send the message 6: a request of data, then the node will answer his data.

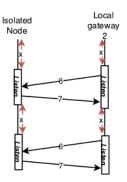


Figure 3. harvest phase

Phase 3: Collection The Local gateway will ask data to the end-nodes. If so, the isolated node sends the data to the Local gateway. This one will then forward it to another Gateway.

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- 1) Accuracy: Each message contains information on the next slot, the time of listening on this one, as well as the id of the senders and receivers. If a message is lost or collapsed, it will be re-sent managed by timers.
- 2) Particular case: When two gateways are in range from the isolated node and there both answer to him with the message 2. The isolated node pick one randomly. To manage multiple isolated node,

## III. ALGORITHM AND SIMULATION

## A. Algorithm

1) Messages format: The set of system messages are of the form:

There are several messages in the system. It is considered that each message corresponds to a function carrying the name of the message, initializing the type and the source of the message, and taking in parameter the destination and the data.

- 2) Message from IN to GW:
- The discover message
  - message set up for the discovery of a GW by an IN
  - destination = udef: broadcast mode
  - data = udef : nothing
- Messages pair
  - Pairing message from an IN to a GW.
  - data = udef : nothing
- Messages data\_response
  - Answer from the data request of a GW.
- 3) les messages  $GW \to IN$ : For all the messages coming from the GW the data part is structured as follows:
  - answer\_frequency: frequency on which the IN must respond
  - next\_slot : delay by the next listening window
  - next\_duration : fixed time of the next listening window
  - next\_frequency : Frequency of the next listening window
  - data: data space specific to the exchange

### The message is like that:

The different messages GW rightarrow IN are thus:

- the messages candidate
  - texttt GW reply message after receiving discover from IN
  - texttt data = udef: no info
- messages data\_request
  - data request message
  - data = udef if only one data available or data = requested\_data in the case of multiple data

- 4) Liste des fonctions utilises:
- a) Fonction d'mission: void send(frequency ,
  message)
- b) Fonction de rception: la fonction listen coute sur la frquence frequency un temps dfinit par time. Le prototype de cette fonction est :

les valeurs des paramtres de cette fonction sont :

- frequency: frquence d'coute
- source : id de l'metteur du message
  - source = udef : coute de tous les nœuds sur la frquence dfinie
- message\_type : type de message attendu
  - message\_type = udef : coute de tous les types
    de messages
- time\_listen : dure de la fentre de rception
  - time\_listen = udef : fentre infinie

## Valeurs de retour :

- message message reu
  - passage du message dans sa totalit
  - message == udef : pas de rception respectant les contraintes
- time temps restant bas sur time\_listen
  - time == udef : dans le cas de time\_listen = udef

```
5) Algorithm 1 - 1:
```

## Algorithm 1 Initialization of communication variables of IN

```
1: procedure init_var(msg)
        gw \leftarrow msg.source
2:
        next\_time \leftarrow msg.next\_slot
3:
        timer \leftarrow msg.next\_duration
4:
5: end procedure
6:
7: procedure flush\_var()
        gw \leftarrow udef
8:
        msg \leftarrow udef
9:
10:
        next\_time \leftarrow udef
        timer \leftarrow timer\_disco
11:
12: end procedure
```

## Algorithm 2 Algorithm IN 1-1

```
1: while (true) do
2:
       flush\_var()
       while (msg = udef) do
3:
          send(freq\_listen, discover(udef, udef))
4:
          (msg,t) = listen(udef, candidate, timer +
5:
   rnd())
      end while
6:
       initVar(msg)
7:
       send(freq\_send, pair(gw, udef))
8:
9:
       while (gw! = udef) do
10:
          sleep(next\ time)
11:
          (msg,t) = listen(gw, data\_request, timer)
12:
          if msq! = udef then
13:
              initVar(msg)
14:
              send(fdate\_response(gw, local\_data)
15:
          else flush var()
16:
          end if
17:
       end while
18:
19: end while
```

## Algorithm 3 Initialization of communication variables of GW

```
1: procedure init_var()
        freg\ send \rightarrow chose()
2:
        timer \leftarrow chose()
3:
4:
        freq\_listen \leftarrow chose()
        freq\_next \leftarrow chose()
5:
6: end procedure
7:
8: procedure flush\_var()
        timer \leftarrow timer\_disco
9:
10:
        freq\_listen \leftarrow freq\_disco
        freq\_send \leftarrow freq\_disco
11:
12:
        in \leftarrow udef
13: end procedure
```

## Algorithm 4 Algorithm gw 1-1

```
1: LoRaWAN_join()
 2: flush\_var()
 3: while (true) do
 4:
       if (in == udef) then
5:
           (msg,t) = listen(,udef,discover,timer)
       end if
 6:
 7:
       if (msg! = udef) then
          in \leftarrow msq.source
 8:
9:
           init\_var()
10:
           send(candidate(in, slot, duration, udef))
           (msg,t) = listen(,in,pair,timer)
11:
12:
          if (msg == udef) then
              flush\_var()
13:
          end if
14:
       end if
15:
16:
       if (in! = udef) then
           init\_var()
17:
           send(data\_request(in, slot, durationt, udef))
18:
           (msg,t) = listen(in, data\_response, timer)
19:
          if (msq! = udef) then
20:
              send\_data(id + ":" + local\_data + ";" +
21:
   in + ":" + msg.data)
22:
          else
              send\_data(id + " : " + local\_data + ";" +
23:
   in + ":" + udef)
24:
              flush\_var()
           end if
25:
       end if
26:
27: end while
```

- 6) Algorithm of GW:
- B. Simulation

### IV. CONCLUSION

[6], [5], [2], [3], [4], [1]

#### ACKNOWLEDGMENT

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