

Project:

Algorithmic complexity and graphs

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3. Complexities

Function 1

```
def function_1(n: int) -> None:
    temp_list = list()
    for i in range(n**2):
        temp = 0
        for j in range(i):
            temp += j
        temp_list.append(temp)
    sum(temp_list)
```

This algorithm uses a loop within a loop. For that reason, it might be tempting to say this algorithm has a runtime complexity of $O(n^2)$. However, the outer loop gets factored by 2, resulting in the same factorization of all operations within the loop itself.

The function starts off by initializing the object `temp_list` as a type list. As this operation is independent of the variable `n` and the list is created empty, we can consider it as $O(1)$. In the next line, we enter the outer loop, looping over the range of `n` factored by 2, resulting in $O(n^2)$ which is the base factorization of all following operations. Although the outer loop is factored by 2, the initializing of `temp` stays $O(1)$ but it will be called n^2 times at the end.

Following this is the initialization of the object `temp` in the inner for-loop, using the value of `i` as its base. The inner for-loop runs at most n^2 times, however, due to the factorization of n^2 by the outer loop, this for-loop runs with a complexity of $O(n^4)$ over all its lines.

`temp += j` is independent of the variable `n` thus being $O(1)$, but it will be called n^4 times at the end. This marks the end of the inner for-loop. The outer for-loop ends with `temp` being appended to `temp_list` which results in $O(1)$. The closing function of `function_1` is a sum of the content of `temp_list`. This operation is $O(n^2)$ as `temp_list` will be of size n^2 by the end of the for-loop.

The upper bound of the complexity of `function_1` is $O(n^4)$. A verification for this complexity can be seen in figure 1.

Written out complexity:

$$1 + n^2 * (1 + (n^2 * 1) + 1) + n^2$$

Shortened complexity:

$$1 + 3n^2 + n^4$$

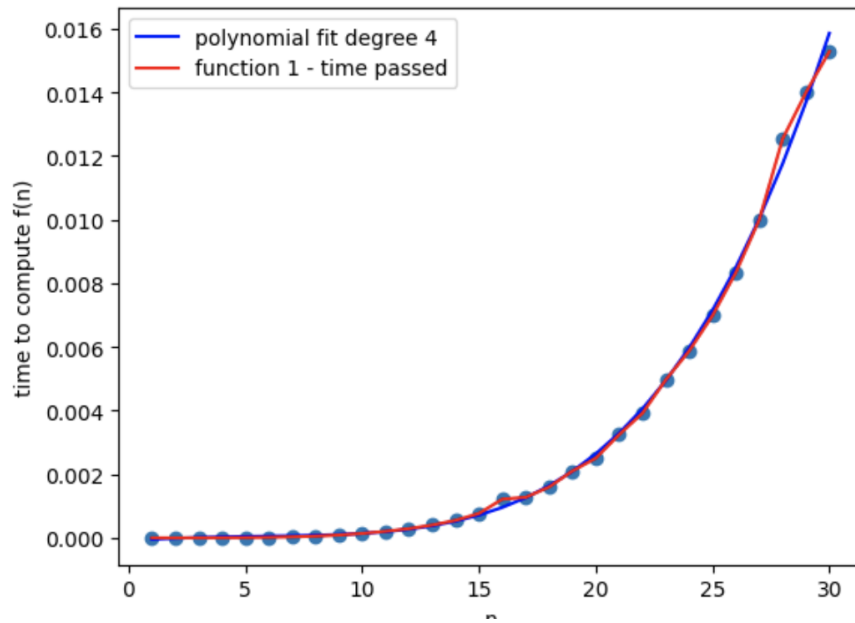


Figure 1. Polynomial fitting of n^4 to the measured time of function_1()

Function 2

```
def function_2(n: int) -> None:
    print(n)
    for i in range(n):
        temp_list = [j+i for j in range(n)]
        shuffle(temp_list)
        max(temp_list)
```

Even if the code seems a bit confusing at first, the nested loop should be a giveaway - there is a for-loop inside another for-loop. This algorithm has $O(n^2)$ runtime complexity, which means it has quadratic complexity. It starts off with an $O(1)$ print function and enters the outer for-loop right after, setting the baseline of $O(n)$ for the following lines. The following line is the most tricky one of the function, as it houses an inline for-loop, looping over the range of the variable n . The outer loop closes off by first shuffling the created list and retrieving the largest value. Both of these functions need to go over the list once thus having a complexity of $O(n)$ each.

The upper bound of the complexity of function_2 is $O(n^2)$. A verification for this complexity can be seen in figure 2.

Written out complexity:

$1 + n * (1 + n * (1) + n + n)$

Shortened complexity:

$1 + 3n^2$

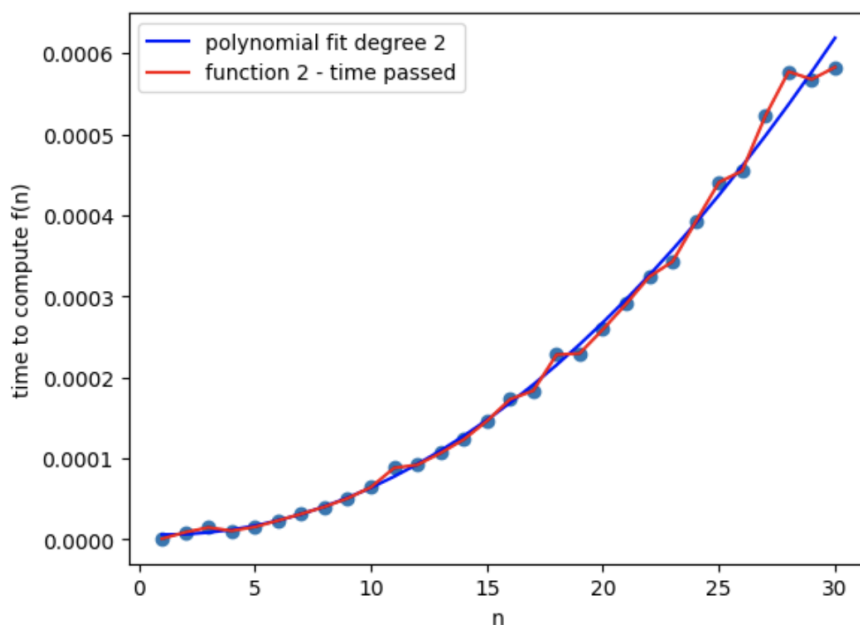


Figure 2. Polynomail fitting of n^2 to the measured time of function_2()