

Solving $A\bar{x} = \bar{b}$: Iterative methods

◦ Recap : We have looked at direct methods

◦ Gaussian elimination

◦ LU decomposition $\rightarrow A^{-1} \rightarrow \bar{x} = A^{-1}\bar{b}$

In principle, exact methods.

◦ Alternative approach : Iterative methods.

◦ Iterate closer and closer to true solution (but not exact)

◦ Often faster, or with smaller memory footprint, than direct methods when matrices are large.

◦ Iterative scheme : Need to think about req. for the iterative procedure to converge.

◦ Methods :

- Jacobi's method (not to be confused with Jacobi's rotation method for $A\bar{x} = \lambda\bar{x}$)
- Gauss-Seidel
- Over-relaxation methods

- Checking convergence (when to stop iterations)

- Typically, monitor the relative change in vector norm

$$\epsilon = \left| \frac{|\bar{x}^{(n+1)}|_l - |\bar{x}^{(n)}|_l}{|\bar{x}^{(n)}|_l} \right|$$

- Intuitively: stop when $\bar{x}^{(n+1)}$ is almost identical to $\bar{x}^{(n)}$,
i.e. when it doesn't change anymore.

- Can use different norms:

$$|\bar{x}|_l \equiv \left[\sum_{i=1}^N |x_i|^l \right]^{\frac{1}{l}}$$

- $l=1$: $|\bar{x}|_1 = |x_1| + |x_2| + \dots + |x_N|$ (sum of abs. vals)

- $l=2$: $|\bar{x}|_2 = \sqrt{x_1^2 + x_2^2 + \dots + x_N^2}$ (Eucl. length)

- $l=\infty$: $|\bar{x}|_\infty = \max_i(x_i)$ (Max element of \bar{x})

$$\left[\begin{array}{l} \text{Intuition:} \\ 0.9^{1000} + 1.2^{1000} + 1.19^{1000} + 0.7^{1000} \approx 1.2^{1000} \\ \text{Formal:} \\ \lim_{l \rightarrow \infty} [x_1^l + x_2^l + \dots + x_N^l]^{\frac{1}{l}} = [\max(x_i)^l]^{\frac{1}{l}} \\ = \max(x_i) \end{array} \right]$$

- Alternative way to monitor convergence :

Monitor residual \bar{r} :

$$\bar{r} = A\bar{x}^{(n+1)} - \bar{b}$$

Note :

$$\begin{aligned}\bar{r} &= A\bar{x}^{(n+1)} - \bar{b} \\ &= A\bar{x}^{(n+1)} - A\bar{x} \\ &= A(\bar{x}^{(n+1)} - \bar{x}) \\ &= A\bar{e}^{(n+1)} \\ &\quad \uparrow \text{error vector}\end{aligned}$$

- So would check $\frac{|\bar{r}|_1}{|\bar{b}|_1}$ at every iteration.

- Pro : This directly monitors how far we are from a vector that satisfies $A\bar{x} = \bar{b}$, not just relative change from $\bar{x}^{(n)}$ to $\bar{x}^{(n+1)}$

- Con : More expensive computationally , since it requires another matrix-vector multiplication.

The Jacobi method (for iterative solution of $A\bar{x} = \bar{b}$)

Summary:

$$\left[\begin{array}{l} \circ \text{ Decompose } A: A = L + D + U \\ \circ \text{ Iterate to find } \bar{x}: \bar{x}^{(m+1)} = -D^{-1}[L+U]\bar{x}^{(m)} + D^{-1}\bar{b} \\ \circ \text{ Start from some guess } \bar{x}^{(0)} \end{array} \right]$$

- Always converges if A is diagonally dominant, i.e.

$$|a_{ii}| > \sum_{j \neq i} |a_{ij}| \quad \text{for all rows } i$$

or A is positive definite, i.e. all eigenvalues > 0 .

(Can converge in other cases too.)

- Rearrange A as sum $L + D + U$:

$$\begin{bmatrix} \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot \end{bmatrix} = \begin{bmatrix} \cdot & 0 & & \\ \cdot & \cdot & & \\ \cdot & \cdot & \cdot & \\ \cdot & \cdot & \cdot & \cdot \end{bmatrix} + \begin{bmatrix} \cdot & & 0 & \\ & \cdot & & 0 \\ 0 & & \cdot & \\ & 0 & & \cdot \end{bmatrix} + \begin{bmatrix} \cdot & \cdot & \cdot & \\ & 0 & \cdot & \cdot \\ & & 0 & \cdot \\ & & & 0 \end{bmatrix}$$
$$A = L + D + U$$

- Note: Unrelated to LU decomposition, which was a product
 $A = LU$

- $A\bar{x} = (L + D + U)\bar{x} = \bar{b}$

$$\Rightarrow D\bar{x} = -(L + U)\bar{x} + \bar{b}$$

$$\Rightarrow \bar{x} = -D^{-1}(L + U)\bar{x} + D^{-1}\bar{b}$$

- Treat this as iterative equation, with $\bar{x}^{(n+1)}$ on the left-hand side, and $\bar{x}^{(n)}$ on the right-hand side.

- $T = D^{-1}[L + U]$ is called the "update matrix".

- Trivial to get D^{-1} since D is diagonal:

$$D^{-1} = \text{diag}\left(\frac{1}{d_{11}}, \frac{1}{d_{22}}, \dots, \frac{1}{d_{NN}}\right) = \begin{bmatrix} \frac{1}{d_{11}} & & & \\ & \frac{1}{d_{22}} & & \\ & & \ddots & \\ & & & \frac{1}{d_{NN}} \end{bmatrix}$$

- 4x4 example: $\bar{x} = \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix}$

- $x_1^{(n+1)} = \left[b_1 - a_{12}x_2^{(n)} - a_{13}x_3^{(n)} - a_{14}x_4^{(n)} \right] / a_{11}$

- $x_2^{(n+1)} = \left[b_2 - a_{21}x_1^{(n)} - a_{23}x_3^{(n)} - a_{24}x_4^{(n)} \right] / a_{22}$

- $x_3^{(n+1)} = \left[b_3 - a_{31}x_1^{(n)} - a_{32}x_2^{(n)} - a_{34}x_4^{(n)} \right] / a_{33}$

- $x_4^{(n+1)} = \left[b_4 - a_{41}x_1^{(n)} - a_{42}x_2^{(n)} - a_{43}x_3^{(n)} \right] / a_{44}$

- In general: $x_i^{(n+1)} = \left[b_i - \sum_{\substack{j=1 \\ j \neq i}}^N a_{ij}x_j^{(n)} \right] / a_{ii}$

Gauss-Seidel method

- Consider 4×4 example for the Jacobi method, but now use the $(m+1)$ versions of already computed vector elements:

$$x_1^{(m+1)} = \left[b_1 - a_{12}x_2^{(m)} - a_{13}x_3^{(m)} - a_{14}x_4^{(m)} \right] / a_{11}$$

$$x_2^{(m+1)} = \left[b_2 - a_{21}\underline{x_1^{(m+1)}} - a_{23}x_3^{(m)} - a_{24}x_4^{(m)} \right] / a_{22}$$

$$x_3^{(m+1)} = \left[b_3 - a_{31}\underline{x_1^{(m+1)}} - a_{32}\underline{x_2^{(m+1)}} - a_{34}x_4^{(m)} \right] / a_{33}$$

$$x_4^{(m+1)} = \left[b_4 - a_{41}\underline{x_1^{(m+1)}} - a_{42}\underline{x_2^{(m+1)}} - a_{43}\underline{x_3^{(m+1)}} \right] / a_{44}$$

- We're doing forward subst.
- On matrix form we're rewriting $A\bar{x} = (L+D+U)\bar{x} = \bar{b}$ as

$$D\bar{x} = -L\bar{x} - U\bar{x} + \bar{b}$$

and turning it into an iterative equation as

$$D\bar{x}^{(m+1)} = -L\bar{x}^{(m+1)} - U\bar{x}^{(m)} + \bar{b}$$

- Component form:

$$\left[\Leftrightarrow \bar{x}^{(m+1)} = -(L+D)^{-1} U\bar{x}^{(m)} + (L+D)^{-1} \bar{b} \right]$$

but this just looks confusing...

$$x_i^{(m+1)} = \left[-\sum_{j=1}^{i-1} a_{ij}x_j^{(m+1)} - \sum_{j=i+1}^N a_{ij}x_j^{(m)} + b_i \right] / a_{ii}$$

• Must be computed in order
• Typically better conv. than Jacobi's method

Successive over-relaxation (SOR)

- Modified version of Gauss-Seidel, with better convergence
- But: has a free parameter ω that must be chosen/tuned for the specific problem.

- Schematically:

$$[\text{new component}] = [\text{old component}] + \overset{\omega}{\downarrow} [\text{weight}] \left[\text{change new-old from Gauss-Seidel algo} \right]$$

- For $\omega=1$ we get back Gauss-Seidel
- Can be proven: For $\omega \in (1, 2]$ we have better convergence than G-S, but optimal choice is problem specific.
- For $\omega > 2$, SOR fails. ("Wobbles" out of control?)
- Component formulation:

$$x_i^{(m+1)} = x_i^{(m)} + \frac{\omega}{a_{ii}} \left[-\sum_{j=1}^{i-1} a_{ij} x_j^{(m+1)} - \sum_{j=i+1}^N a_{ij} x_j^{(m)} + b_i - a_{ii} x_i^{(m)} \right]$$

$$i=1, 2, \dots, N$$

- Matrix formulation (confusing-looking)

(evaluated in this order)

$$\bar{x}^{(m+1)} = (\omega L + D)^{-1} \cdot \left[-[\omega U + (\omega - 1)D] \bar{x}^{(m)} + \omega \bar{b} \right]$$

• Why are iterative methods often useful?

- Each iteration only requires matrix-vector multiplication, which has complexity

$$\mathcal{O}(N^1) < \mathcal{O}(N^k) < \mathcal{O}(N^2)$$

↑
If A is diagonal

↑
if A is dense matrix

- If the method requires M iterations for convergence, we have a total cost of

$$\mathcal{O}(N^k M) \text{ operations}$$

which is typically much less than the typical $\mathcal{O}(N^3)$

cost of direct methods (assuming number of iterations M is smaller than matrix size N)

• Second advantage: (in practice) can often get more accurate results than direct (exact) methods, because iterative methods are less susceptible to round-off errors.