



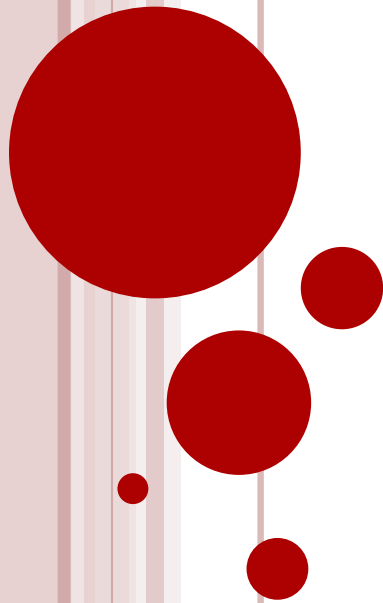
CORSO DI LAUREA IN
INGEGNERIA INFORMATICA



SOCIAL NETWORKS ANALYSIS

A.A. 2022/23

GRAPHS AND NETWORKS

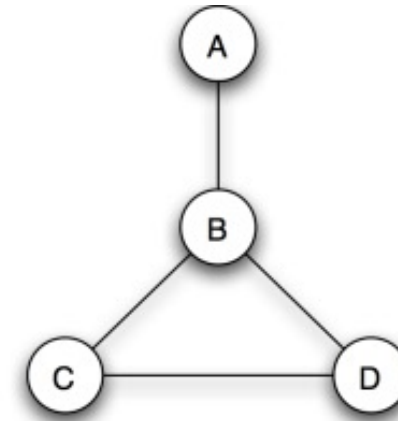


GRAPH THEORY

- Graph Theory is a consolidated branch of Mathematics that allows to describe sets of elements together with binary relations between these elements
 - Started in the XVIII century by Euler
 - ❖ Solution to the Königsberg's bridges problem
- Graph theory provides a unifying language to describe the structure of all kinds of networks
- Nowadays, the possibility of gather data on large scale and work with massive data sets pave the way to new approaches and new problems on graphs
 - Recognize statistical properties characterizing the structure of a network and provide methods to measure them
 - Create network models and describe network formation process
 - Predict the behaviour of a network based on their models and structural properties

GRAPH'S DEFINITION

- A *graph* consists of
 - A set of nodes (vertices)
 - A set of edges (links)
 - ❖ An edge connects two nodes

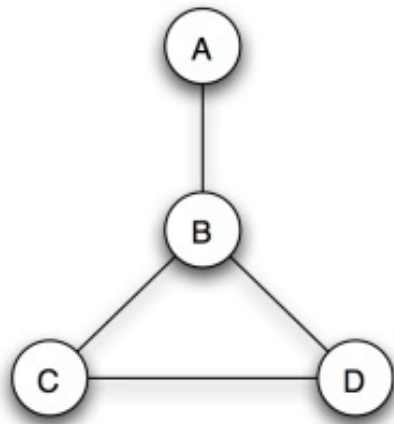


- Two nodes are adjacent (**neighbors**) if they are connected by an edge
 - C and D are adjacent through the edge (C, D)
 - B is a neighbor of both C and D

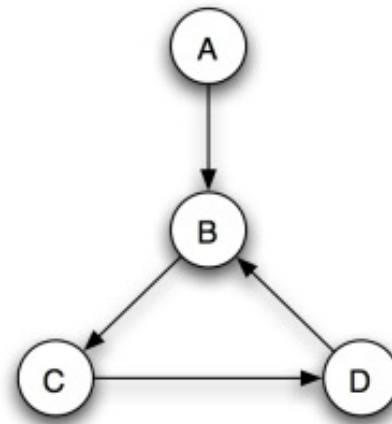
DIRECTED AND UNDIRECTED GRAPHS

- Edges can be oriented

- Directed edge: the relation holds only between the head and the tail of the edge
- Undirected edge: the relation holds in both the directions



Undirected graphs



Directed graphs

- Directed and undirected edges differ substantially

- Different models of network formation and maintenance
- Different algorithms

GRAPHS AS MODELS OF NETWORKS

- Directed and undirected graphs describe different kinds of networks
- Directed graphs
 - The relation between the nodes originates from a unilateral decision
 - ❖ Eg: link to a web page, followers, citation of an article
- Undirected graphs
 - The relation between the nodes comes from a decision of both the elements
 - ❖ Eg: friendship, alliance, acquaintance, connection, ecc.

WEIGHTED OR SIGNED GRAPHS

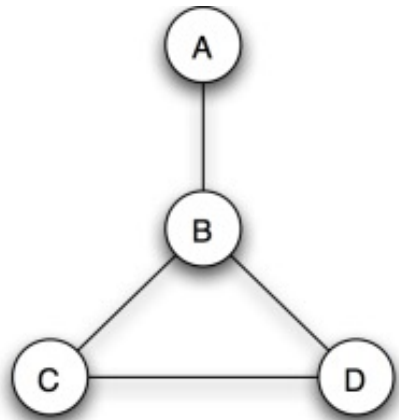
- We can put additional informations to edges
 - sign (friends/enemies)
 - weight (strength of the social connection)
 - distance (connection length)
 - delay (transmission time)
 - reliability (transmission error rate)
 - cost (cost of the link usage)
- In a *weighted* graph each edge has a number associated that is its weight
- In a *signed* graph each edge has a positive/negative sign

GRAPH REPRESENTATION

- Graph Theory is a mathematical theory interesting by itself
 - Studies characteristics and properties of graphs
- A graph is a pair of sets
 - $G = (V, E)$
 - ❖ V = vertex set
 - ❖ E = edge set
- More used representations
 - Adjacency matrix
 - ❖ matrix $n \times n$ ($n = |V|$)
 - ❖ Element $(i, j) = 1$ if there is an edge between i and j
 - $w_{i,j}$ is the weight of edge (i, j) if the graph is weighted
 - Lists of vertices V and edges E
 - ❖ For each vertex v we have the list of vertices adjacent to v
 - ❖ For directed graphs we have separate lists of incoming and outgoing edges
 - Graphic representation

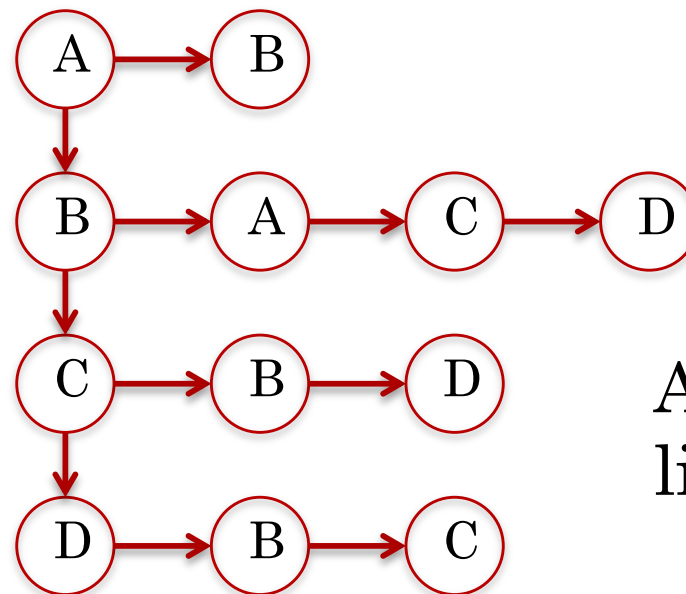
AN EXAMPLE OF GRAPH REPRESENTATION

7



	A	B	C	D
A	0	1	0	0
B	1	0	1	1
C	0	1	0	1
D	0	1	1	0

Adjacency
matrix

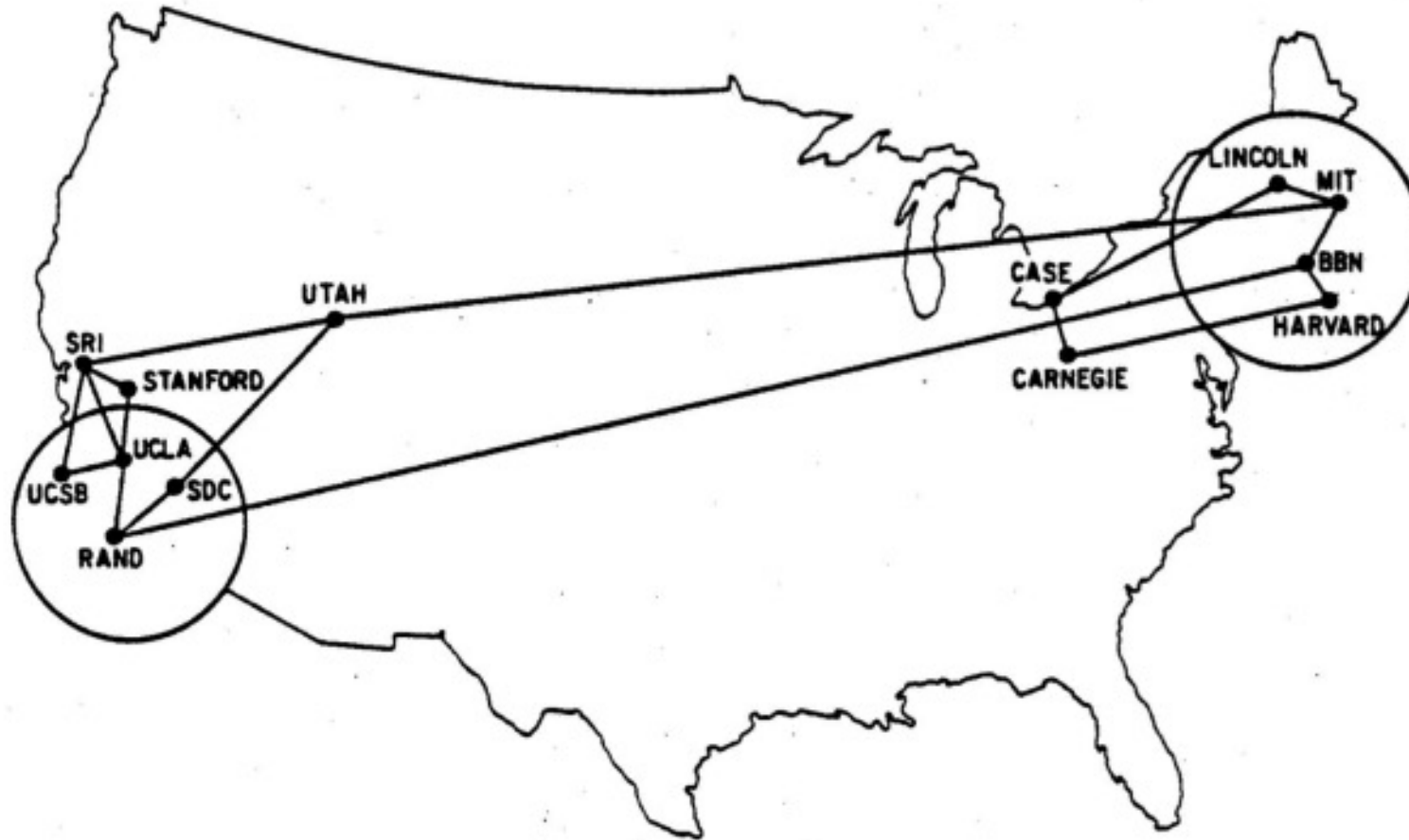


Adjacency
lists

VERTICES AND EDGES

- When studying networks vertices and edges can represent entities of the real world
 - Some network abstraction are commonly used
- Some examples
 - Communication networks
 - ❖ Network devices, communication links
 - Social networks
 - ❖ people, friendships/social connections
 - ❖ companies, commercial relations
 - Information networks
 - ❖ Web sites, hyperlinks
 - Biological networks
 - ❖ Species, predator-prey links
 - ❖ molecules, chemical bonds

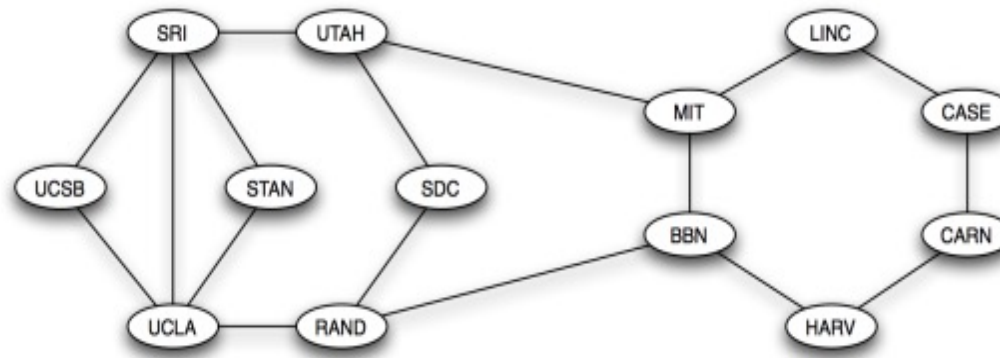
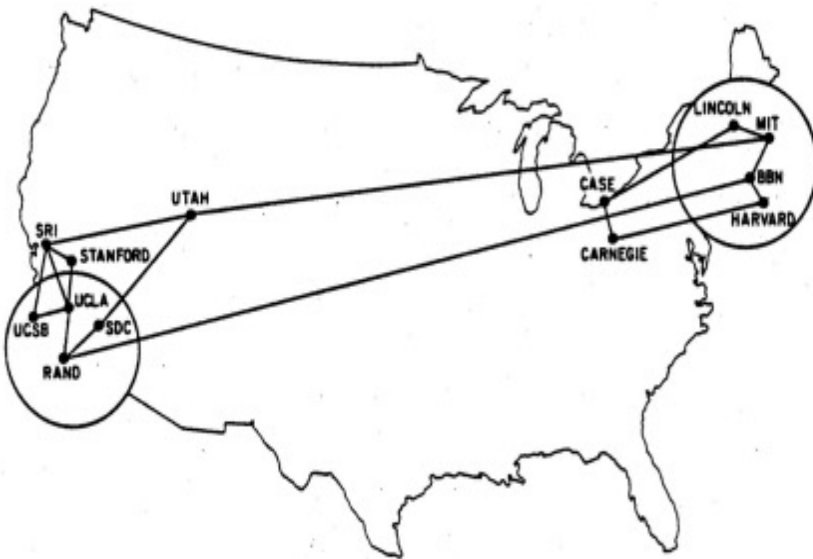
ARPANET: FIRST VERSION OF INTERNET



- Created in 1970 with 13 nodes

THE GRAPH OF ARPANET

- We are interested only in connectivity
 - Distances can be represented as edge weights



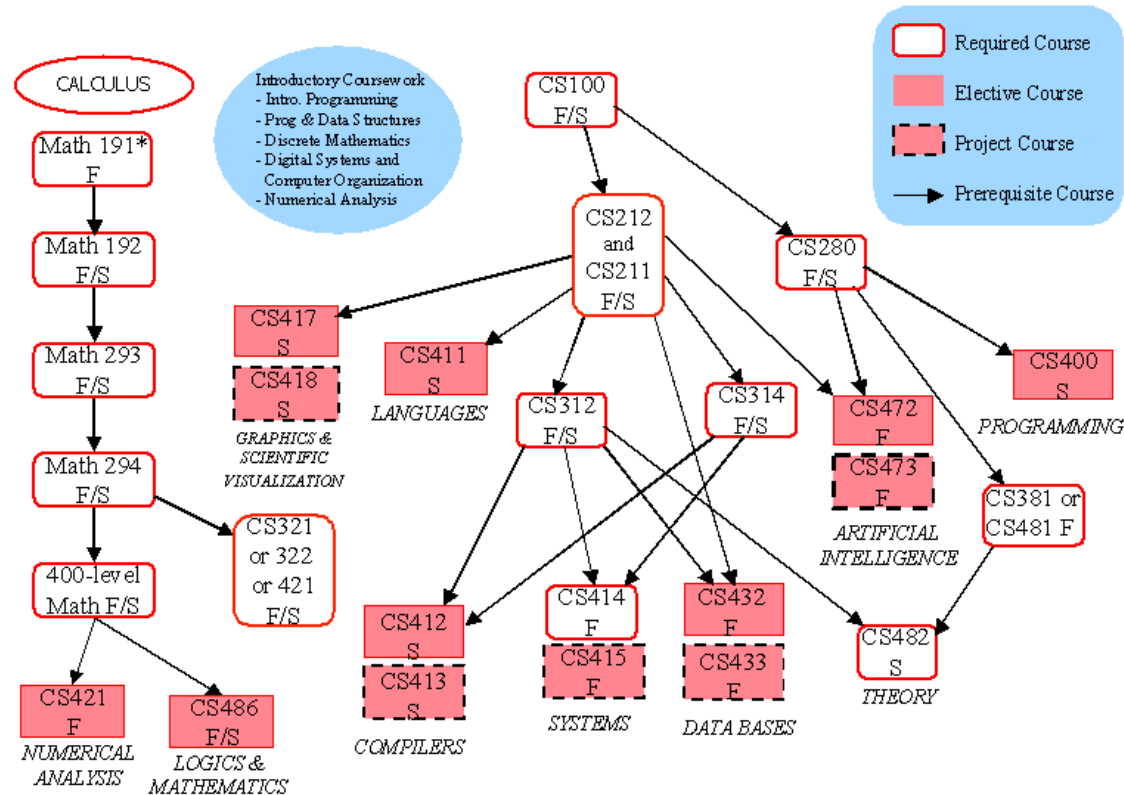
TRANSPORTATION NETWORKS



- Most of the graph terminology derives from metaphors based on transportation
 - “shortest path”, “diameter”, “flow”
- Questions we are interested in
 - The network structure can support the required performances? How much robust is it? Is it exposed to cascading failures?

DEPENDENCY GRAPHS

Undergraduate Computer Science Courses for Majors



* Calc. Sequence In Arts Is 111; 112; 221; 222

- Nodes are tasks and directed edges are dependencies
- To design complex software systems or industrial processes we need to carefully analyze the dependency graph to define a good scheduling policy and avoid deadlocks

STRUCTURAL NETWORKS



- The internal frameworks of mechanical structures such as buildings, vehicles, or human bodies are based on such networks
- rigidity theory studies the stability of such structures from a graph-based perspective

MOST RELEVANT CONCEPTS ON GRAPHS

- “Graph theory is a terminological jungle in which every newcomer may plant a tree“

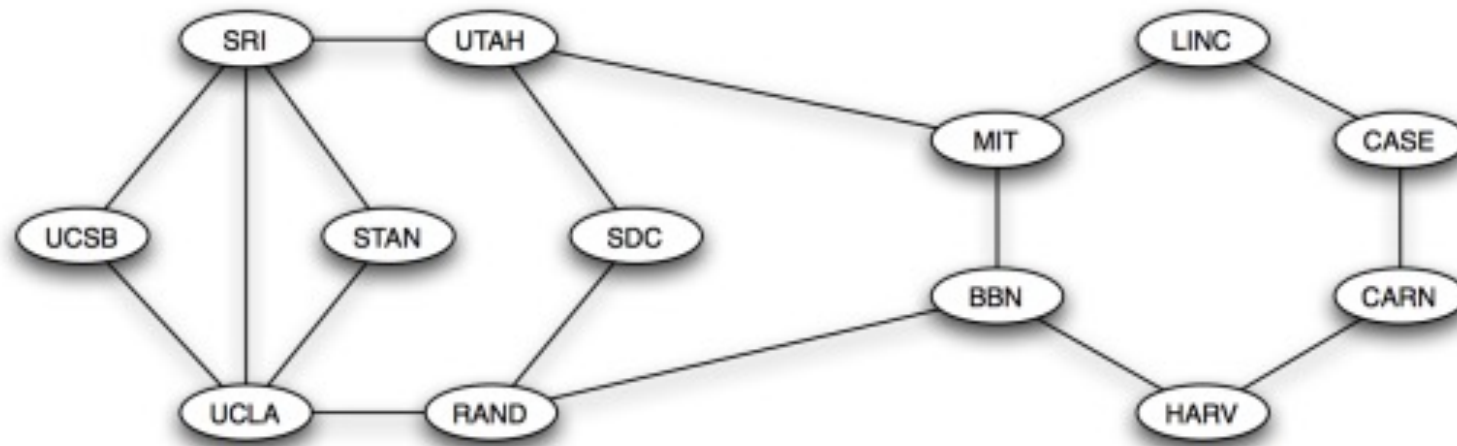
(Social scientist John Barnes)

- We will introduce only concepts that are relevant for the scope of this course
 - paths
 - cycles
 - connectivity
 - components (giant component)
 - distance

PATHS

- A characteristic of the social networks is that nodes can influence each other indirectly
 - Influence can travel along the connections of the network
- Several things can travel along edges of a graph
 - Vehicles
 - information
 - influence or popularity
 - diseases
- A *path* is a sequence of nodes such that every two consecutive nodes in the path are connected by an edge
 - You can also see the path as a sequence of edges, where two consecutive edges share an endpoint
 - If there exists a path between u and v then they are in an undirected relation
- Multiplying the adjacency matrix for itself k times we can find how many paths of length k there are between each pair of nodes
 - $M^k(u, v) = \# \text{ paths of length } k \text{ between } u \text{ and } v$

P



- MIT – BBN – RAND – UCLA is a path
- UCSB – UCLA – RAND – MIT is not a path
- A path can go through the same node several times
 - SRI – STAN – UCLA – SRI – UTAH – MIT
- A *simple path* never goes through the same node twice
- A *shortest path* is a path that goes through the minimum number of edges

CYCLES

- A *cycle* is a simple path that starts and ends in the same node
 - LINC – CASE – CARN – HARV – BBN – MIT – LINC is a cycle
 - A cycle has at least three edges
- In communication networks and transportation networks each node lies on one or many cycles
 - Redundancy introduced to increase the robustness of the network
 - The network is guaranteed to work even in presence of a limited number of faults
- In a social network cycles are very common but accidental and we don't care of them
 - Es: the bestfriend of the cousin of my wife is the sister of my officemate

DIRECTED PATHS AND CYCLES

- Directed paths and cycles can be defined similarly to the undirected case
 - We have only to take care of the direction of the edges in the path
- Sometimes we consider undirected cycles even in directed graphs
 - We simply ignore the edges' directions
 - Useful if we are interested in the existence of a relation, independently from who activated it

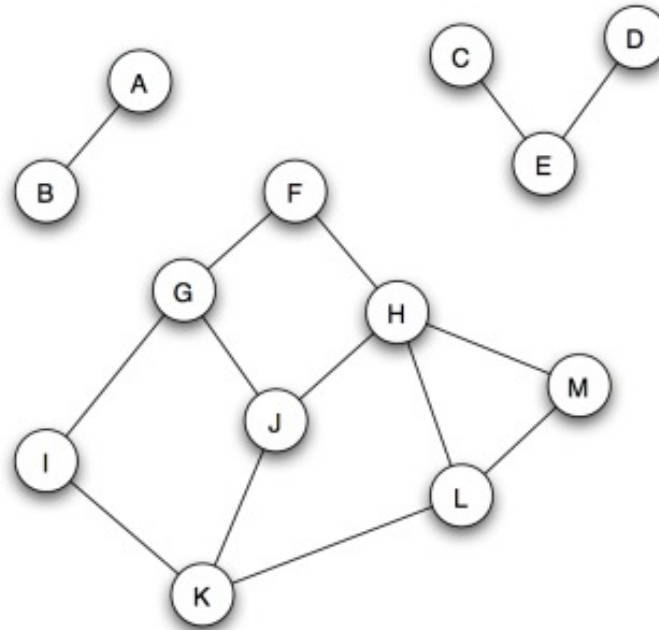
CONNECTIVITY

- Two nodes u e v are connected if there exists a path from u to v
- A graph is *connected* if there exists a path between each pair of nodes in the graph
 - The graph of ARPANET is connected
 - Communication networks are usually connected
- In most of the cases graphs are disconnected
 - Social networks
 - Collaboration networks
 - *etc.*
- A directed graph is *strongly connected* if there exists a directed path between each pair of nodes in the graph

COMPONENTS

- If a graph is not connected it can be partitioned in subgraphs that are connected
- A *connected component* C in a undirected graph is a subset of nodes such that
 - There exists a path between each pair of nodes in C
 - For each node u not in C , there exists at least one vertex v in C such that there exists no path between u and v
- A *strongly connected component* S in a directed graph is a subset of nodes such that
 - There exists a directed path between each pair of nodes in S
 - For each node u not in S , there exists at least one vertex v in S such that there exists no directed path between u and v
- An edge belongs to a component if both its endpoints belong to the component
- The edge is a *bridge* if its removal makes the component disconnected

COMPONENTS



- This graph has three connected components
 - $\{A,B\}$, $\{C,D,E\}$, $\{F,G,...,M\}$
- $\{H, L, M\}$ is not a component
- (D, E) is a bridge

COMPONENT ANALYSIS

- Analyzing the components of a graph we can gain useful informations on the global structure of the network
 - Which edges tie different components?
 - How information spreads in the network?
 - Which role has each node?



Graph of the scientific collaborations in a research center

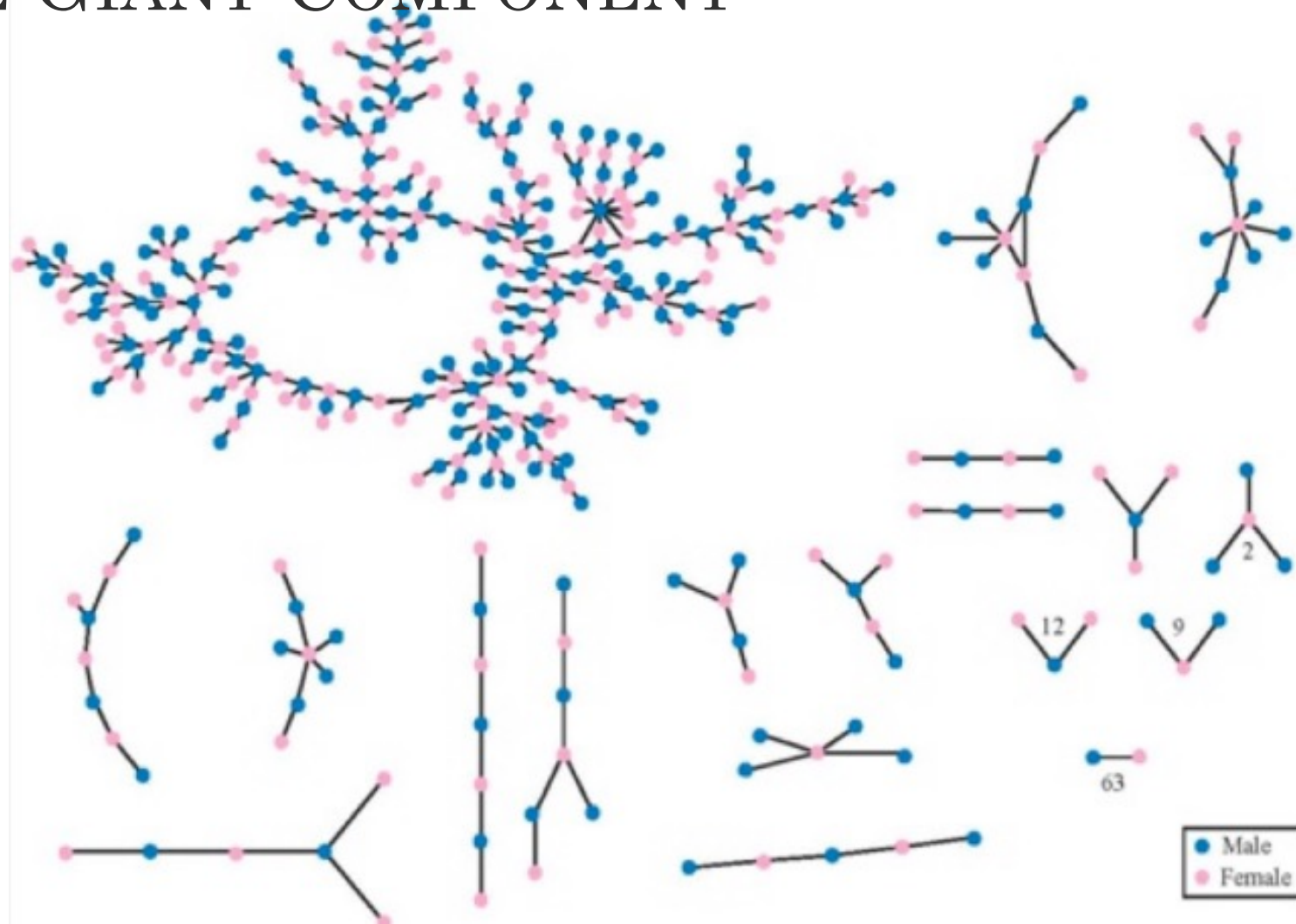
COMPONENT ANALYSIS

- In many cases it's crucial to recognize components in a network
- A component is a region of the network densely connected
 - We would like to recognize the most densely connected components and their borders
- Another important type of analysis
 - We consider only edges with weight greater than a given threshold
 - Gradually increasing the threshold the graph will break into several components
 - The remaining components include nodes that are strongly tied

THE GIANT COMPONENT

- Several graphs are not connected but they include a very large component
 - E.g. The graph of the love relations in a high school, the graph of the Web
- A *giant component* is a connected component containing a large fraction of the nodes
 - Usually the giant component is unique
- When two giant components merge it can give rise to dramatic events
 - The graph of the human relations between populations before the America discovery had two giant components
 - Their fusion brought to the extermination of one of the two components

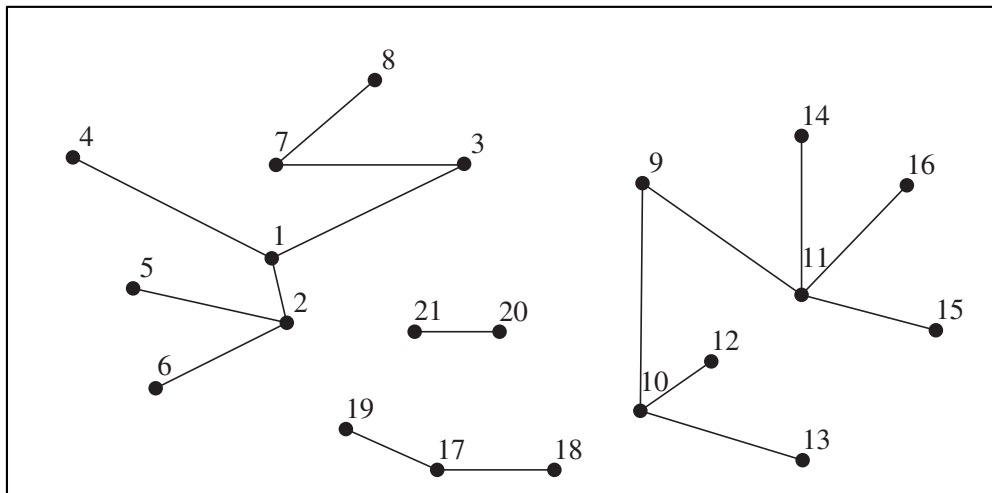
THE GIANT COMPONENT



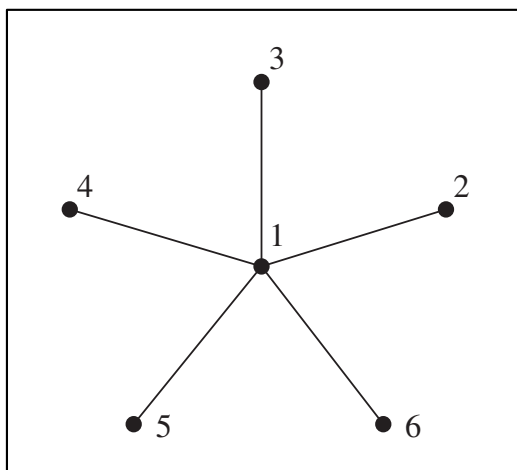
- The existence of a giant component in this network implies a higher risk of diffusion of sexual diseases

PARTICULAR CLASSES OF GRAPHS

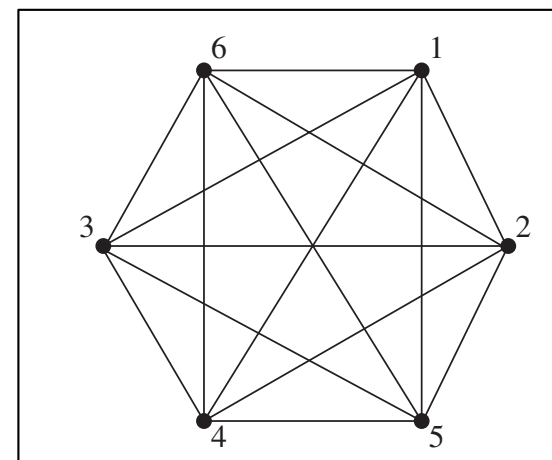
- There are specific topologies that occur very often and there were very well studied



Trees and Forests



Star



Clique

NEIGHBORHOODS

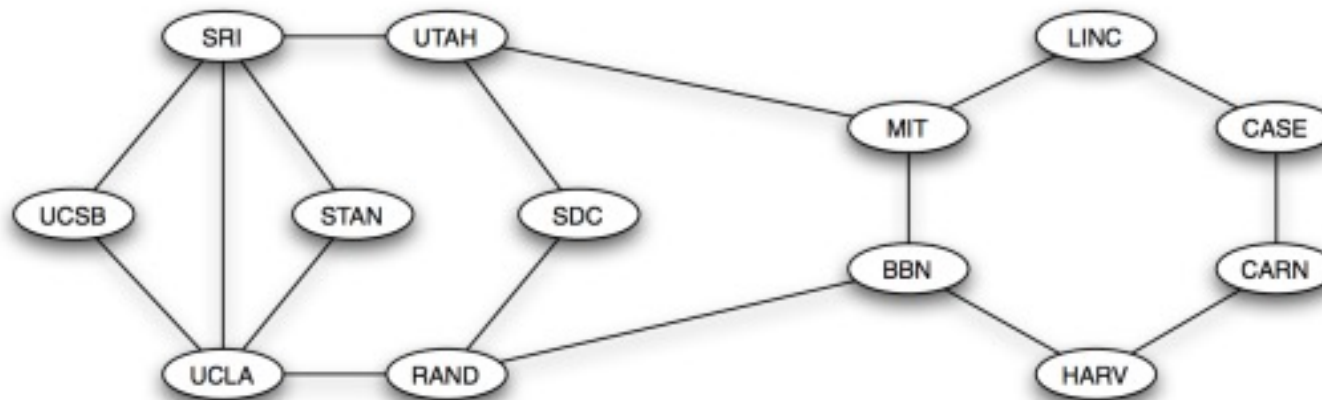
- The **neighborhood** of a node is the set of nodes that are adjacent to it
- The neighborhood of a set of nodes S is the set of nodes that do not belong to S but they are adjacent to some nodes in S

DEGREES

- The *degree* of a node is the number of its neighbors
 - Number of edges adjacent to the node
 - It's equal to the size of the neighborhood
- In a directed graph we distinguish between in-degree and out-degree
 - In-degree: number of incoming edges
 - Out-degree: number of outgoing edges

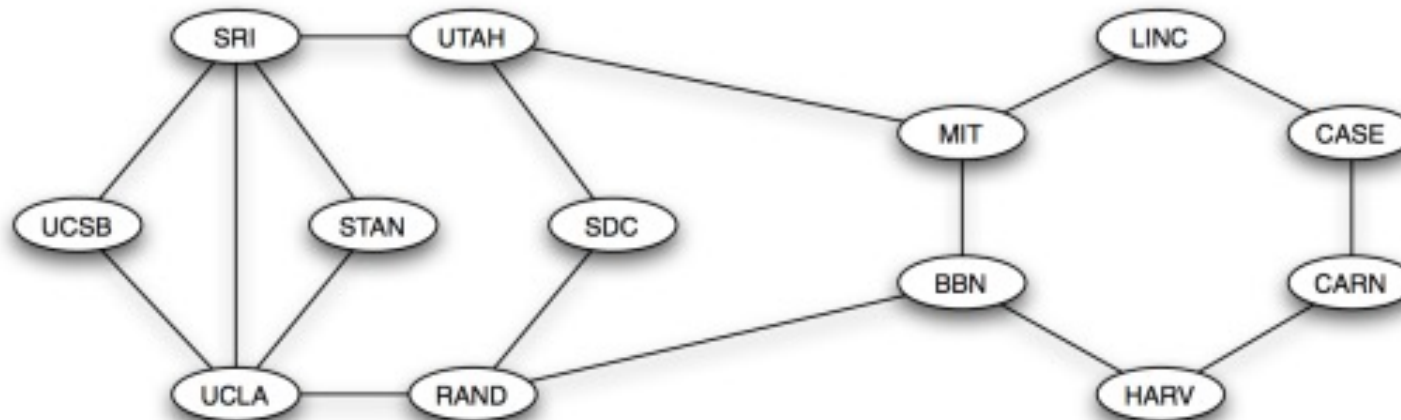
DISTANCES

- The *distance* between a pair of nodes is the length of the shortest path connecting the nodes
 - We assume each edge has weight 1
- The *diameter* of a graph is the largest distance between a pair of nodes in the graph
 - Which is the distance between MIT and SDC?
 - Which is the diameter of the network?



HOW TO COMPUTE MINIMAL DISTANCES

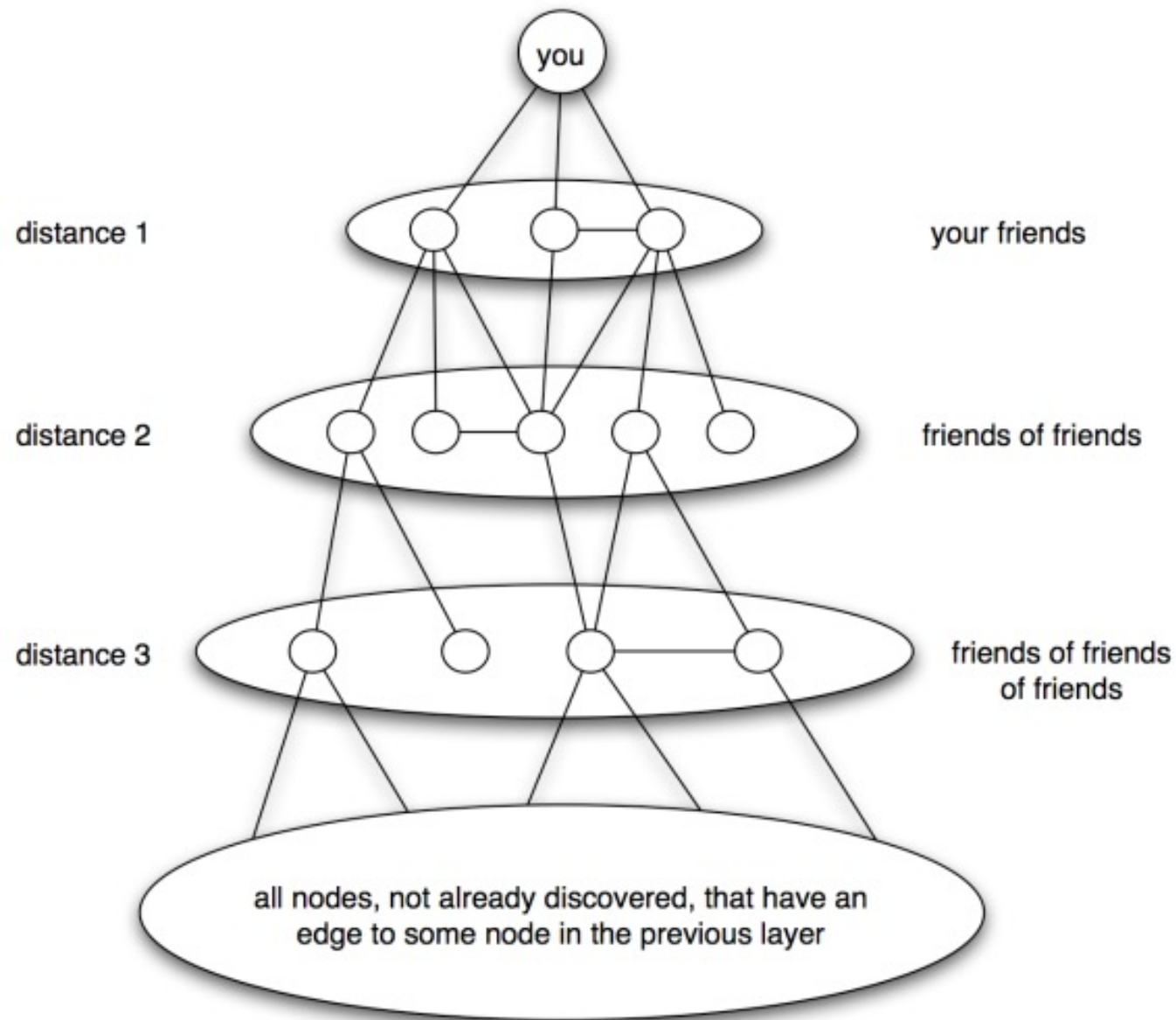
- Given a graph, how can we compute the minimal distances from a node to all the others?
 - We need an efficient algorithm
- How can we approach the problem?
 - BFS



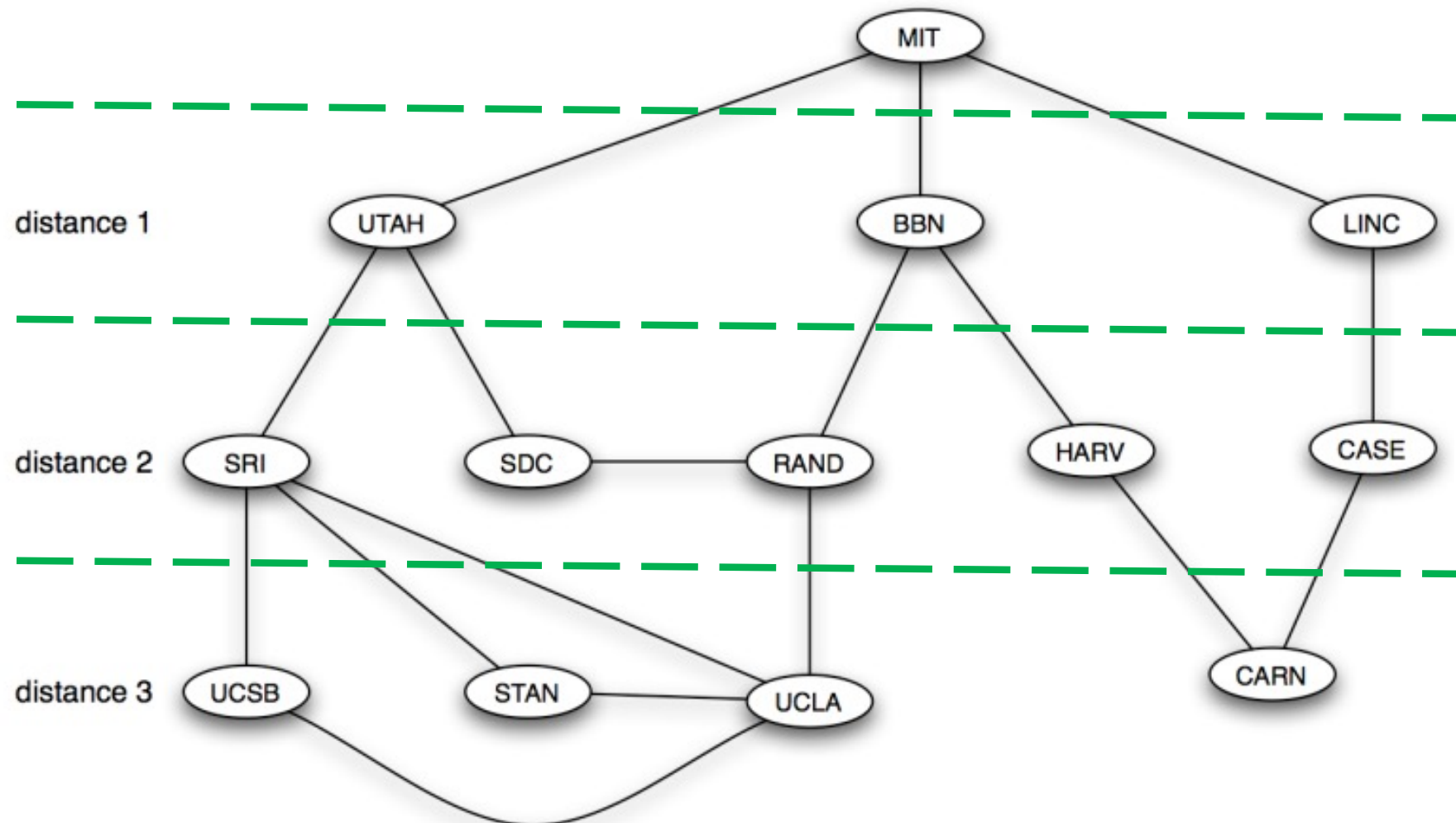
BREADTH-FIRST SEARCH (BFS)

- Starting from the source node (*root*)
 - Find all nodes adjacent to root
 - ❖ These nodes are at “distance 1”
 - Find all the nodes that are adjacent to nodes at distance 1 and not yet visited
 - ❖ These nodes are at “distance 2”
 - ...
 - Find all the nodes that are adjacent to nodes at distance j and not yet visited
 - ❖ These nodes are at “distance $j+1$ ”
 - Stops when there are no other adjacent vertices not visited

BFS TREE



BFS ON THE GRAPH OF ARPANET

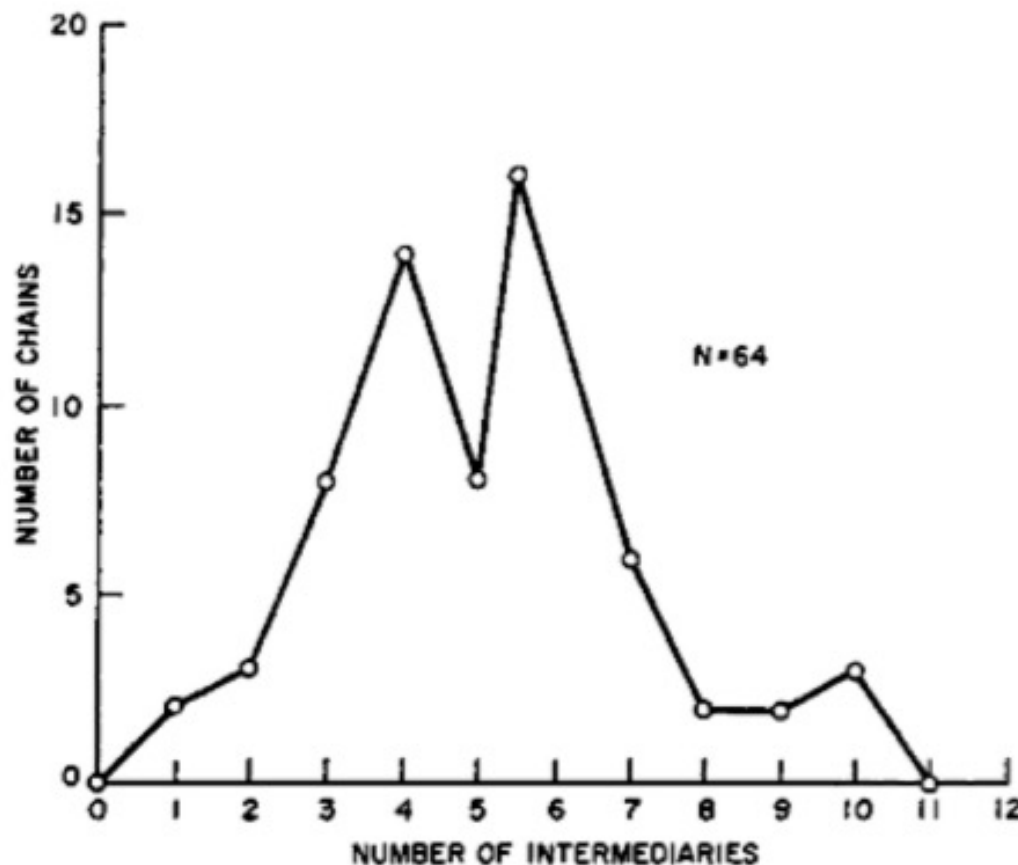


SMALLWORLD PHENOMENON

- Hypothesis: in large scale networks most of the nodes belong to a giant component and they are connected through very short paths
 - Informations/infections spread very fast
- First experiment realized by Stanley Milgram in 1960s (research budget \$680)
 - 296 people, randomly chosen in the USA, was asked to deliver a letter to a given recipient
 - ❖ They received a profile of the recipient (address, work, education, place of origin, interests, ecc.)
 - ❖ They could send the letter to one of their friends or acquaintances
 - The experiment measured the average number of hops for each letter

SIX DEGREES OF SEPARATION

- In the Milgram's experiment only 64 letters were delivered to the recipient
 - Average number of hops < 6

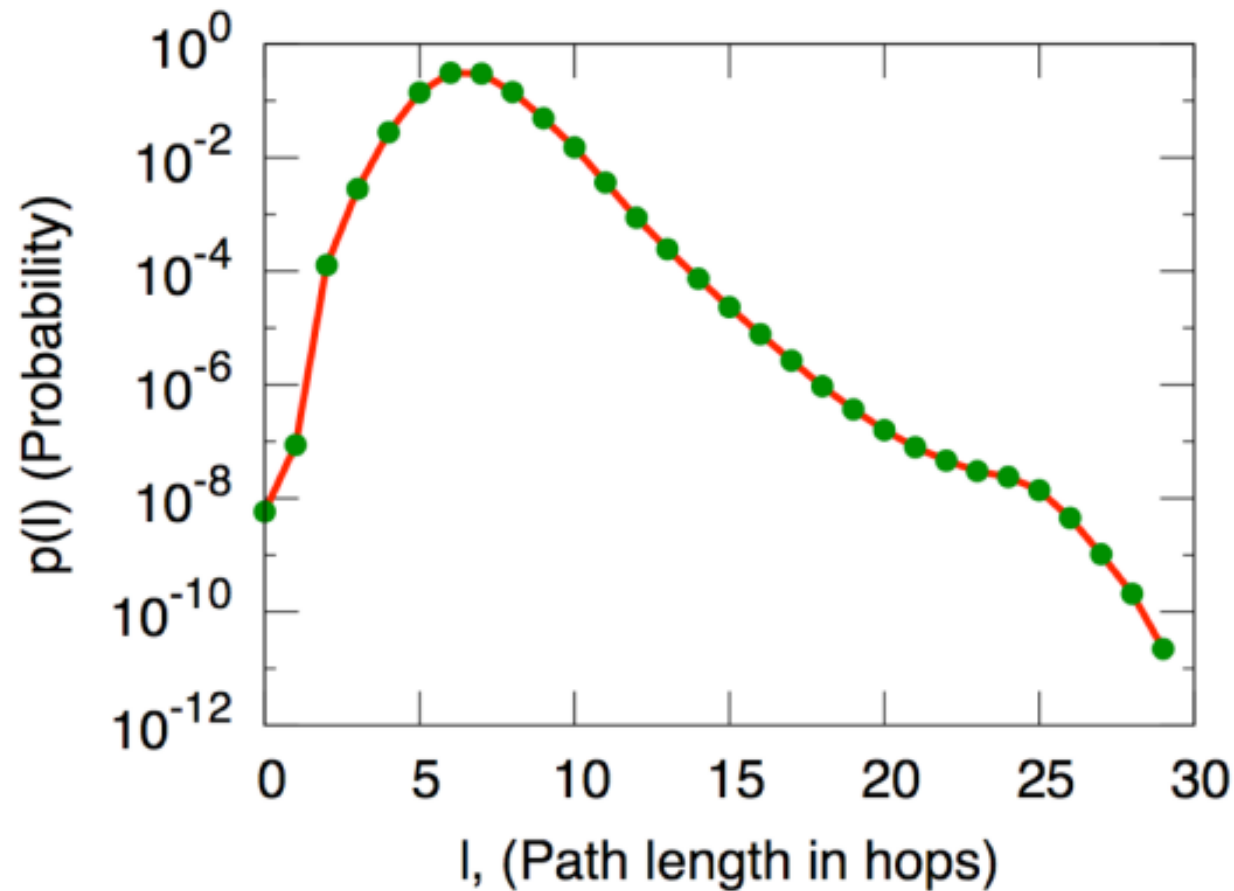


The experiment has been largely criticized in the following years but ...

SIX DEGREES OF SEPARATION

- In recent years several experiments confirmed Milgram's results
- In 2008 Leskovec and Horvitz realized a new version of the Milgram's experiment
 - They used data related to Messenger's connections of 240 milion users in a period of 30 days
 - Their graph has a giant component with average distance equal to 6.6
- In each run they selected a random sample of 1000 users and computed minimal distances with BFS

LESKOVEC AND HORVITZ RESULTS



expected average distance = 6.6, median = 7

SIX DEGREES OF CERTAINTY

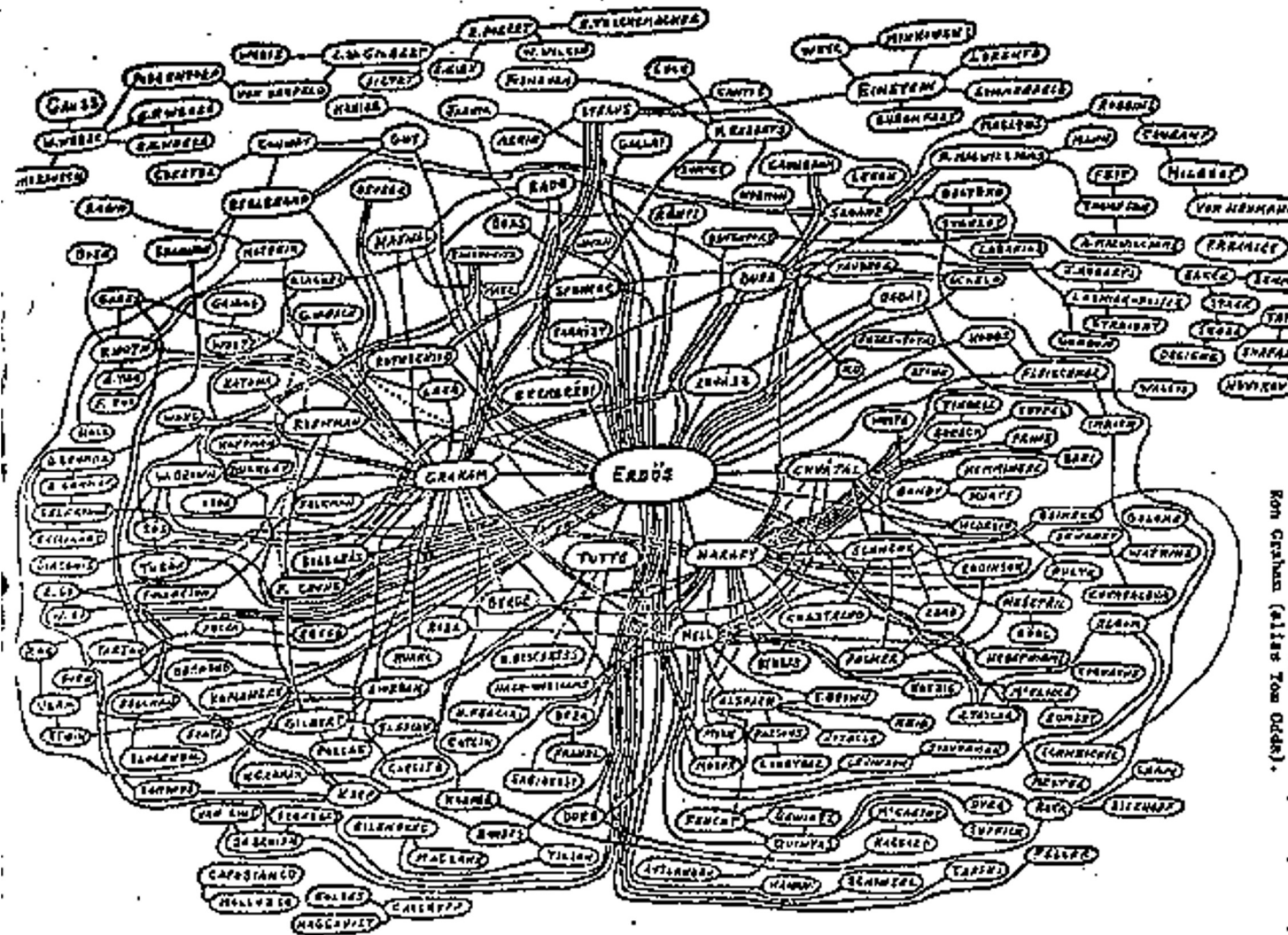


Figure 1
To appear in *Topics in Graph Theory* (P. Harary, ed.) New York Academy of Sciences (1979).

- Graph of the scientific collaborations rooted in Paul Erdős
 - Most of the mathematicians (and computer scientists) have Erdős number < 5

HOW TO DESCRIBE A GRAPH WITHOUT REPRESENTING EXPLICITLY

39

- A large scale network can have millions of nodes
 - It cannot be represented explicitly
 - ❖ We need a set of quantitative parameters that can describe the graph
 - ❖ We use these parameters to compare graphs without representing them
- Some of the most used parameters
 - diameter, average path length
 - Clustering indices
 - Centrality measures
 - Node degree distribution

DEGREE DISTRIBUTION

- The distribution of the node degrees is fundamental characteristic of the network
- $\Pr(d)$ = fraction of nodes with degree d
 - Probability that a node randomly chosen has degree d
- k -regular graphs have a degenerate degree distribution
- Random graphs have a Poisson degree distribution
- In several networks the degree distribution is a power law
 - $\Pr(d) = cd^{-\gamma}$

DIAMETER AND AVERAGE PATH LENGTH

- The diameter is an upper bound to the length of the shortest path between each pair of nodes in a connected graph
- The average path length is the average of the shortest path lengths between all pairs of nodes
- Comparing the diameter with the average path length we can obtain useful informations
 - If they are not comparable then there are very few pairs of nodes that are very far apart

CLUSTERING INDICES

- An interesting information about a social network is how much connected and close it is
 - How many of my friends are friends each other?
- These characteristics can be measured through the **clustering indices**
- Two alternative definitions
 - **Overall clustering**: fraction of node pairs that are adjacent and they have a common neighbor
 - **Individual clustering (of node u)**: fraction of pairs of u 's neighbors that are adjacent
 - ❖ Average Clustering is the average of the individual clusterings of all nodes in the graph

CENTRALITY MEASURES

- They measure the relevance (centrality) of a node in the network related to a given process
 - We can use them to compare nodes
- There are several centrality measures that can model different processes
 - Degree centrality
 - Closeness centrality
 - Betweenness centrality
 - Katz-prestige centrality
 - Eigenvector centrality

NETWORK DATA-SETS

- On the web there are several datasets of large scale networks
 - Collaboration graphs
 - ❖ Wikipedia, World of Warcraft, Citation graphs
- Who-talks-to-Whom Graphs
 - Microsoft IM, Cell phone graphs
- Information networks
 - Snapshots of the Web, social networks, blogging sites
- Technological networks
 - Power grids, communication links, Internet
- Networks in the Natural world
 - Food webs, neural interconnections, cell metabolism
- SNAP is a general purpose network analysis and graph mining library led by Jure Leskovec at Stanford University
 - <http://snap.stanford.edu/data>
 - There is a repository with lots of data on large scale networks