

**HW 10 Due 3 nov 2017**

1. Prove that the class of Turing-acceptable languages is closed under union, intersection, and reversal. For each property, give a detailed sketch of the proof, by saying how you would build a Turing machine that accepts the resulting language, given the Turing machine(s) that accept the original language(s). 50
2. Prove or disprove that the set of Turing-acceptable languages is closed under concatenation. 15
3. Consider a new type of *deterministic* machine, having one read-only input tape and two stacks. The tape is read-only, it cannot be written, but the head can move left, right, or do nothing. Each stack operates, independently of the other, as in a deterministic pushdown automaton:

$$M = (K, \Sigma, \Gamma_1, \Gamma_2, z_1, z_2, \delta, s)$$

where  $K$  is a finite set of states,  $\Sigma$  is a finite input alphabet,  $\Gamma_1$  and  $\Gamma_2$  are two finite stack alphabets ( $\Gamma_1$  for the first stack,  $\Gamma_2$  for the second stack),  $z_1 \in \Gamma_1$  and  $z_2 \in \Gamma_2$  are the initial symbols for the two stacks,  $s \in K$  is the initial state.  $h$  is a special halting state not in  $K$ , just like in a Turing machine.

- (a) Give an appropriate definition for the transition function  $\delta$ , for a configuration of this machine, for the “yields in one step” operator, and for the language accepted by this machine.
- (b) These machines can accept the same languages as a class of automata you already know: deterministic pushdown automata, pushdown automata, or Turing machines? Prove your answer formally. 50