

Chapter 2

Introduction to Number Theory

Divisibility

- We say that a nonzero b divides a if a = mb for some m, where a, b, and m are integers
- b divides a if there is no remainder on division
- The notation b | a is commonly used to mean b divides a
- If b | a we say that b is a divisor of a

The positive divisors of 24 are 1, 2, 3, 4, 6, 8, 12, and 24 13 | 182; - 5 | 30; 17 | 289; - 3 | 33; 17 | 0

Properties of Divisibility

- If $a \mid 1$, then $a = \pm 1$
- If $a \mid b$ and $b \mid a$, then $a = \pm b$
- Any $b \neq 0$ divides 0
- If a | b and b | c, then a | c

If b | g and b | h, then b | (mg + nh) for arbitrary integers m and n

Properties of Divisibility

- To see this last point, note that:
 - If $b \mid g$, then g is of the form $g = b * g_1$ for some integer g_1
 - If $b \mid h$, then h is of the form $h = b * h_1$ for some integer h_1
- So:
 - $mg + nh = mbg_1 + nbh_1 = b * (mg_1 + nh_1)$ and therefore b divides mg + nh

```
b = 7; g = 14; h = 63; m = 3; n = 2

7 \mid 14 and 7 \mid 63.

To show 7 (3 * 14 + 2 * 63),

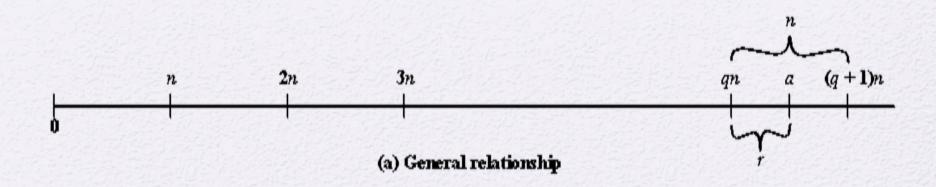
we have (3 * 14 + 2 * 63) = 7(3 * 2 + 2 * 9),

and it is obvious that 7 \mid (7(3 * 2 + 2 * 9)).
```

Division Algorithm

Given any positive integer n and any integer a,
if we divide a by n we get an integer quotient q
and an integer remainder r that obey the
following relationship:

$$a = qn + r$$
 $0 \le r < n; q = \lfloor a/n \rfloor$



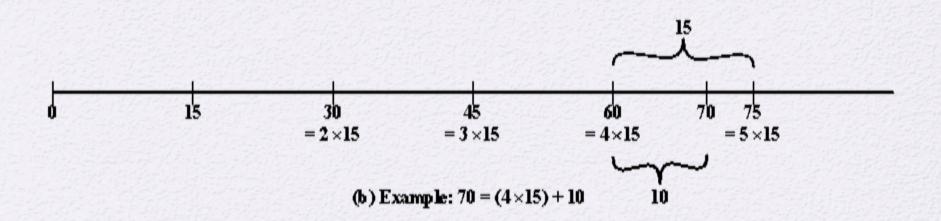


Figure 2.1 The Relationship a = qn + r; $0 \le r \le n$

Euclidean Algorithm



- One of the basic techniques of number theory
- Procedure for determining the greatest common divisor of two positive integers
- Two integers are relatively prime if their only common positive integer factor is 1

Greatest Common Divisor (GCD)

- The greatest common divisor of a and b is the largest integer that divides both a and b
- We can use the notation gcd(a,b) to mean the greatest common divisor of a and b
- We also define gcd(o,o) = o
- Positive integer c is said to be the gcd of a and b if:
 - c is a divisor of a and b
 - Any divisor of a and b is a divisor of c
- An equivalent definition is:

gcd(a,b) = max[k, such that k | a and k | b]

GCD

- Because we require that the greatest common divisor be positive, gcd(a,b) = gcd(a,-b) = gcd(-a,b) = gcd(-a,-b)
- In general, gcd(a,b) = gcd(|a|, |b|)

$$gcd(60, 24) = gcd(60, -24) = 12$$

- Also, because all nonzero integers divide o, we have gcd(a,o) = | a |
- We stated that two integers a and b are relatively prime if their only common positive integer factor is 1; this is equivalent to saying that a and b are relatively prime if gcd(a,b) = 1

8 and 15 are relatively prime because the positive divisors of 8 are 1, 2, 4, and 8, and the positive divisors of 15 are 1, 3, 5, and 15. So 1 is the only integer on both lists.

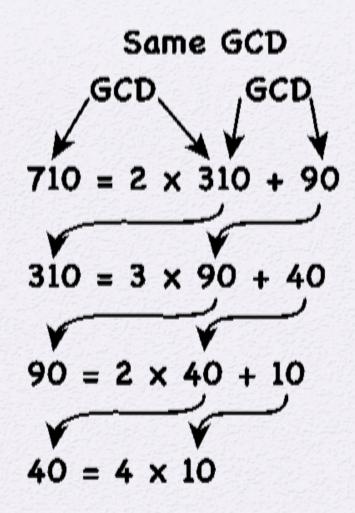


Figure 2.3 Euclidean Algorithm Example: gcd(710, 310)

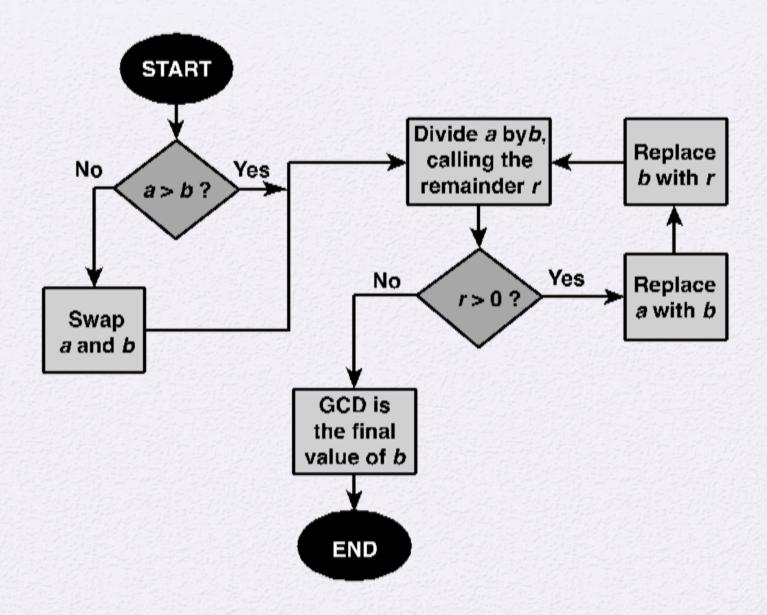


Figure 2.2 Euclidean Algorithm

Table 2.1

Euclidean Algorithm Example

| D | ividend | Divisor | Quotient | Remainder |
|------------------|-----------|-------------------|--------------|-------------------|
| a = 1 | 160718174 | b = 316258250 | $q_1 = 3$ | $r_1 = 211943424$ |
| b = | 316258250 | $r_1 = 211943424$ | $q_2 = 1$ | $r_2 = 104314826$ |
| r ₁ = | 211943424 | $r_2 = 104314826$ | $q_3 = 2$ | $r_3 = 3313772$ |
| r ₂ = | 104314826 | $r_3 = 3313772$ | $q_4 = 31$ | $r_4 = 1587894$ |
| r ₃ = | 3313772 | $r_4 = 1587894$ | $q_5 = 2$ | $r_5 = 137984$ |
| $r_4 =$ | 1587894 | $r_5 = 137984$ | $q_6 = 11$ | $r_6 = 70070$ |
| r ₅ = | 137984 | $r_6 = 70070$ | $q_7 = 1$ | $r_7 = 67914$ |
| r ₆ = | 70070 | $r_7 = 67914$ | $q_8 = 1$ | $r_8 = 2156$ |
| r ₇ = | 67914 | $r_8 = 2156$ | $q_9 = 31$ | $r_9 = 1078$ |
| r ₈ = | 2156 | $r_9 = 1078$ | $q_{10} = 2$ | $r_{10} = 0$ |

(This table can be found on page 30 in the textbook)

Modular Arithmetic

- The modulus
 - If a is an integer and n is a positive integer, we define a mod n to be the remainder when a is divided by n; the integer n is called the modulus
 - Thus, for any integer a:

$$a = qn + r$$
 $0 \le r < n; q = [a/n]$
 $a = [a/n] * n + (a mod n)$

 $11 \mod 7 = 4$; $-11 \mod 7 = 3$

Modular Arithmetic

- Congruent modulo n
 - Two integers a and b are said to be congruent modulo n if (a mod n) = (b mod n)
 - This is written as $a \equiv b \pmod{n}$
 - Note that if $a \equiv o \pmod{n}$, then $n \mid a$

```
73 \equiv 4 \pmod{23}; 21 \equiv -9 \pmod{10}
```

Properties of Congruences

Congruences have the following properties:

```
1. a \equiv b \pmod{n} if n \mid (a - b)

2. a \equiv b \pmod{n} implies b \equiv a \pmod{n}

3. a \equiv b \pmod{n} and b \equiv c \pmod{n} imply a \equiv c \pmod{n}
```

- To demonstrate the first point, if n(a b), then (a b) = kn for
 - So we can write a = b + kn
 - Therefore, $(a \mod n) = (remainder when b + kn is divided by n) = (remainder when b is divided by n) = (b mod n)$

```
23 \equiv 8 \pmod{5} because 23 - 8 = 15 = 5 * 3
- 11 \equiv 5 \pmod{8} because - 11 - 5 = -16 = 8 * (-2)
81 \equiv 0 \pmod{27} because 81 - 0 = 81 = 27 * 3
```

some k

Modular Arithmetic

Modular arithmetic exhibits the following properties:

1.
$$[(a \mod n) + (b \mod n)] \mod n = (a + b) \mod n$$

2.
$$[(a \mod n) - (b \mod n)] \mod n = (a - b) \mod n$$

3.
$$[(a \mod n) * (b \mod n)] \mod n = (a * b) \mod n$$

- We demonstrate the first property:
 - Define $(a \mod n) = r_a$ and $(b \mod n) = r_b$. Then we can write $a = r_a + jn$ for some integer j and $b = r_b + kn$ for some integer k
 - Then:

Remaining Properties:

Examples of the three remaining properties:

```
11 mod 8 = 3; 15 mod 8 = 7

[(11 mod 8) + (15 mod 8)] mod 8 = 10 mod 8 = 2

(11 + 15) mod 8 = 26 mod 8 = 2

[(11 mod 8) - (15 mod 8)] mod 8 = -4 mod 8 = 4

(11 - 15) mod 8 = -4 mod 8 = 4

[(11 mod 8) * (15 mod 8)] mod 8 = 21 mod 8 = 5

(11 * 15) mod 8 = 165 mod 8 = 5
```

Modular Arithmetic Inverses

- Modular Additive Inverse
 - If (a + b)mod n = o, then a and b are modular additive inverses of each other mod n.

- Modular Multiplicative Inverses
 - If (a * b) mod n = 1, then a and b are modular multiplicative inverses of each other mod n.

Modular Arithmetic Inverses

- Modular additive inverses can be used to carry out subtraction operations mod n
 - $(a b) \mod n = (a + (-b)) \mod n$, where -b is the modular additive inverse of b.
- Modular multiplicative inverses can be used to carry out division operations mod n
 - (a/b)mod $n = (a * (b^{-1}))$ mod n, where b^{-1} is the modular multiplicative inverse of b.

Table 2.2(a)

Arithmetic Modulo 8

| + | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|---|---|---|---|---|---|---|---|---|
| 0 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| 1 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 0 |
| 2 | 2 | 3 | 4 | 5 | 6 | 7 | 0 | 1 |
| 3 | 3 | 4 | 5 | 6 | 7 | 0 | 1 | 2 |
| 4 | 4 | 5 | 6 | 7 | 0 | 1 | 2 | 3 |
| 5 | 5 | 6 | 7 | 0 | 1 | 2 | 3 | 4 |
| 6 | 6 | 7 | 0 | 1 | 2 | 3 | 4 | 5 |
| 7 | 7 | 0 | 1 | 2 | 3 | 4 | 5 | 6 |

Table 2.2(b)

Multiplication Modulo 8

| × | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|---|---|---|---|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| 2 | 0 | 2 | 4 | 6 | 0 | 2 | 4 | 6 |
| 3 | 0 | 3 | 6 | 1 | 4 | 7 | 2 | 5 |
| 4 | 0 | 4 | 0 | 4 | 0 | 4 | 0 | 4 |
| 5 | 0 | 5 | 2 | 7 | 4 | 1 | 6 | 3 |
| 6 | 0 | 6 | 4 | 2 | 0 | 6 | 4 | 2 |
| 7 | 0 | 7 | 6 | 5 | 4 | 3 | 2 | 1 |

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(This table can be found on page 33 in the textbook)

Table 2.2(c)

Additive and Multiplicative Inverse Modulo 8

| W | -w | w^{-1} |
|---|----|----------|
| 0 | 0 | |
| 1 | 7 | 1 |
| 2 | 6 | |
| 3 | 5 | 3 |
| 4 | 4 | <u>=</u> |
| 5 | 3 | 5 |
| 6 | 2 | |
| 7 | 1 | 7 |

if
$$(a + b) \equiv (a + c) \pmod{n}$$
 then $b \equiv c \pmod{n}$

if $(a \times b) \equiv (a \times c) \pmod{n}$ then $b \equiv c \pmod{n}$

Properties of Modular Arithmetic for Integers in Z_n

| Property | Expression |
|-----------------------|--|
| Commutative Laws | $(w+x) \bmod n = (x+w) \bmod n$ |
| | $(w \times x) \bmod n = (x \times w) \bmod n$ |
| Associative Laws | $[(w+x)+y] \bmod n = [w+(x+y)] \bmod n$ |
| ASSOCIATIVE Laws | $[(w \times x) \times y] \bmod n = [w \times (x \times y)] \bmod n$ |
| Distributive Law | $[w \times (x + y)] \mod n = [(w \times x) + (w \times y)] \mod n$ |
| Identities | $(0+w) \bmod n = w \bmod n$ |
| Identities | $(1 \times w) \bmod n = w \bmod n$ |
| Additive Inverse (-w) | For each $w \in \mathbb{Z}_n$, there exists a z such that $w + z \equiv 0 \mod n$ |

Extended Euclidean Algorithm

- The extended Euclidean algorithm can be used to calculate the modular multiplicative inverse
 - Assume a > b
 - 2. If gcd(a,b) = 1 i.e. they are mutually prime
 - Then the modular multiplicative inverse of b mod a = y; y derived from solving for ax + by = d using the extended Euclidean algorithm
- From the previous example the modular multiplicative inverse of 550 mod 1759 is 355 since gcd(1759,550) = 1

Extended Euclidean Algorithm Example

Used to find x and y such that ax + by = d = gcd(a,b)

| | Extended Euclidean Algorithm | | | | | | | | | | |
|--|-------------------------------|---|--|--|--|--|--|--|--|--|--|
| Calculate | Which satisfies | Calculate | Which satisfies | | | | | | | | |
| $r_{-1} = a$ | | $x_{-1} = 1; y_{-1} = 0$ | $a = ax_{-1} + by_{-1}$ | | | | | | | | |
| $r_0 = b$ | | $x_0 = 0; y_0 = 1$ | $b = ax_0 + by_0$ | | | | | | | | |
| $ \begin{aligned} r_1 &= a \bmod b \\ q_1 &= \lfloor a/b \rfloor \end{aligned} $ | $a = q_1 b + r_1$ | $\begin{vmatrix} x_1 = x_{-1} - q_1 x_0 = 1 \\ y_1 = y_{-1} - q_1 y_0 = -q_1 \end{vmatrix}$ | $r_1 = ax_1 + by_1$ | | | | | | | | |
| $ \begin{aligned} r_2 &= b \bmod r_1 \\ q_2 &= \lfloor b/r_1 \rfloor \end{aligned} $ | $b = q_2 r_1 + r_2$ | $\begin{vmatrix} x_2 = x_0 - q_2 x_1 \\ y_2 = y_0 - q_2 y_1 \end{vmatrix}$ | $r_2 = ax_2 + by_2$ | | | | | | | | |
| $ \begin{aligned} r_3 &= r_1 \bmod r_2 \\ q_3 &= \lfloor r_1/r_2 \rfloor \end{aligned} $ | $r_1 = q_3 r_2 + r_3$ | $ \begin{aligned} x_3 &= x_1 - q_3 x_2 \\ y_3 &= y_1 - q_3 y_2 \end{aligned} $ | $r_3 = ax_3 + by_3$ | | | | | | | | |
| • | • | • | • | | | | | | | | |
| • | • | • | • | | | | | | | | |
| • | • | • | • | | | | | | | | |
| $\begin{vmatrix} r_n = r_{n-2} \bmod r_{n-1} \\ q_n = \lfloor r_{n-2}/r_{n-1} \rfloor \end{vmatrix}$ | $r_{n-2} = q_n r_{n-1} + r_n$ | $\begin{vmatrix} x_n = x_{n-2} - q_n x_{n-1} \\ y_n = y_{n-2} - q_n y_{n-1} \end{vmatrix}$ | $r_n = ax_n + by_n$ | | | | | | | | |
| $r_{n+1} = r_{n-1} \bmod r_n = 0$ $q_{n+1} = \lfloor r_{n-1}/r_n \rfloor$ | $r_{n-1} = q_{n+1}r_n + 0$ | | $d = \gcd(a, b) = r_n$ $x = x_n; y = y_n$ | | | | | | | | |

$$a = 1759 \qquad b = 550 \qquad r_1 = a \times 1 + b \times 1$$

$$r_1 = r_{1-2} - r_{1-1} \cdot q_1 \qquad q_1 = \frac{r_{1-2}}{r_{1-1}} \qquad x_1 = x_{1-2} - q_1 \times 1 \cdot 1 \qquad q_1 = q_1 \cdot q_1$$

Extended Euclidean Algorithm Example

Used to find x and y such that ax + by = d = gcd(a,b)

| i | r_i | q_i | x_i | Y_i |
|----|-------|-------|-------|-------|
| -1 | 1759 | | 1 | 0 |
| 0 | 550 | | 0 | 1 |
| 1 | 109 | 3 | 1 | -3 |
| 2 | 5 | 5 | -5 | 16 |
| 3 | 4 | 21 | 106 | -339 |
| 4 | 1 | 1 | -111 | 355 |
| 5 | 0 | 4 | | |

Result: d = 1; x = -111; y = 355

(This table can be found on page 39 in the textbook)

Prime Numbers

- Prime numbers only have divisors of 1 and itself
 - They cannot be written as a product of other numbers
- Prime numbers are central to number theory
- Any integer a > 1 can be factored in a unique way as

$$a = p_1^{a1} p_2^{a2} \dots p_t^{at}$$

where $p_1 < p_2 < ... < p_t$ are prime numbers and where each a_i is a positive integer

This is known as the fundamental theorem of arithmetic

Table 2.5 Primes Under 2000

| 2 | 101 | 211 | 307 | 401 | 503 | 601 | 701 | 809 | 907 | 1009 | 1103 | 1201 | 1301 | 1409 | 1511 | 1601 | 1709 | 1801 | 1901 |
|----|-----|-----|-----|-----|-----|-----|-----|-----|-----|------|------|------|------|------|------|------|------|------|------|
| 3 | 103 | 223 | 311 | 409 | 509 | 607 | 709 | 811 | 911 | 1013 | 1109 | 1213 | 1303 | 1423 | 1523 | 1607 | 1721 | 1811 | 1907 |
| 5 | 107 | 227 | 313 | 419 | 521 | 613 | 719 | 821 | 919 | 1019 | 1117 | 1217 | 1307 | 1427 | 1531 | 1609 | 1723 | 1823 | 1913 |
| 7 | 109 | 229 | 317 | 421 | 523 | 617 | 727 | 823 | 929 | 1021 | 1123 | 1223 | 1319 | 1429 | 1543 | 1613 | 1733 | 1831 | 1931 |
| 11 | 113 | 233 | 331 | 431 | 541 | 619 | 733 | 827 | 937 | 1031 | 1129 | 1229 | 1321 | 1433 | 1549 | 1619 | 1741 | 1847 | 1933 |
| 13 | 127 | 239 | 337 | 433 | 547 | 631 | 739 | 829 | 941 | 1033 | 1151 | 1231 | 1327 | 1439 | 1553 | 1621 | 1747 | 1861 | 1949 |
| 17 | 131 | 241 | 347 | 439 | 557 | 641 | 743 | 839 | 947 | 1039 | 1153 | 1237 | 1361 | 1447 | 1559 | 1627 | 1753 | 1867 | 1951 |
| 19 | 137 | 251 | 349 | 443 | 563 | 643 | 751 | 853 | 953 | 1049 | 1163 | 1249 | 1367 | 1451 | 1567 | 1637 | 1759 | 1871 | 1973 |
| 23 | 139 | 257 | 353 | 449 | 569 | 647 | 757 | 857 | 967 | 1051 | 1171 | 1259 | 1373 | 1453 | 1571 | 1657 | 1777 | 1873 | 1979 |
| 29 | 149 | 263 | 359 | 457 | 571 | 653 | 761 | 859 | 971 | 1061 | 1181 | 1277 | 1381 | 1459 | 1579 | 1663 | 1783 | 1877 | 1987 |
| 31 | 151 | 269 | 367 | 461 | 577 | 659 | 769 | 863 | 977 | 1063 | 1187 | 1279 | 1399 | 1471 | 1583 | 1667 | 1787 | 1879 | 1993 |
| 37 | 157 | 271 | 373 | 463 | 587 | 661 | 773 | 877 | 983 | 1069 | 1193 | 1283 | | 1481 | 1597 | 1669 | 1789 | 1889 | 1997 |
| 41 | 163 | 277 | 379 | 467 | 593 | 673 | 787 | 881 | 991 | 1087 | , | 1289 | | 1483 | | 1693 | | | 1999 |
| 43 | 167 | 281 | 383 | 479 | 599 | 677 | 797 | 883 | 997 | 1091 | | 1291 | | 1487 | | 1697 | | | |
| 47 | 173 | 283 | 389 | 487 | | 683 | | 887 | | 1093 | | 1297 | | 1489 | | 1699 | | | |
| 53 | 179 | 293 | 397 | 491 | | 691 | | | | 1097 | | | | 1493 | | | | | |
| 59 | 181 | | | 499 | | | | | | | | | | 1499 | | | | | |
| 61 | 191 | | | | | | | | | | | | | | | | | | |
| 67 | 193 | | | | | | | | | | | | | | | | | | |
| 71 | 197 | | | | | | | | | | | | | | | | | | |
| 73 | 199 | | | | | | | | | | | | | | | | | | |
| 79 | | | | | | | | | | | | | | | | | | | |
| 83 | | | | | | | | | | | | | | | | | | | |
| 89 | | | | | | | | | | | | | | | | | | | |
| 97 | | | | | | | | | | | | | | | | | | | |
| | | | | | | | | | | | | | | | | | | | |

Fermat's Theorem

How to solve modulus operation for big numbers?

$$7^{180}$$
=? (mod 19)

Fermat's Theorem

- States the following:
 - If p is prime and a is a positive integer not divisible by p then

$$a^{p-1} \equiv 1 \pmod{p}$$

- An alternate form is:
 - If p is prime and a is a positive integer, then

$$a^p \equiv a \pmod{p}$$

Fermat's Theorem Examples

$$a = 7, p = 19$$

 $7^2 = 49 \equiv 11 \pmod{19}$
 $7^4 \equiv 121 \equiv 7 \pmod{19}$
 $7^8 \equiv 49 \equiv 11 \pmod{19}$
 $7^{16} \equiv 121 \equiv 7 \pmod{19}$
 $a^{p-1} = 7^{18} = 7^{16} \times 7^2 \equiv 7 \times 11 \equiv 1 \pmod{19}$

$$p = 5, a = 3$$
 $a^p = 3^5 = 243 \equiv 3 \pmod{5} = a \pmod{p}$
 $p = 5, a = 10$ $a^p = 10^5 = 100000 \equiv 10 \pmod{5} \equiv 0 \pmod{5} = a \pmod{p}$

Fermat's Theorem

How to solve modulus operation for big numbers?

•
$$7^{180}$$
=? (mod 19)

$$7^{180} = (7^{18})^{10} = (1)^{10} = 1 \pmod{19}$$

Euler's Totient Function $\emptyset(n)$

- $\phi(n)$ = The number of positive integers less than n and relatively prime to n.
- $\phi(1) = 1$
- $\emptyset(p) = p 1$ (where p and q are prime)
- $\emptyset(pq) = \emptyset(p)\emptyset(q) = (p-1)(q-1)$

Table 2.6

Some Values of Euler's Totient Function $\emptyset(n)$

| n | φ(<i>n</i>) |
|----|---------------|
| 1 | 1 |
| 2 | 1 |
| 3 | 2 |
| 4 | 2 |
| 5 | 4 |
| 6 | 2 |
| 7 | 6 |
| 8 | 4 |
| 9 | 6 |
| 10 | 4 |

| n | $\phi(n)$ |
|----|-----------|
| 11 | 10 |
| 12 | 4 |
| 13 | 12 |
| 14 | 6 |
| 15 | 8 |
| 16 | 8 |
| 17 | 16 |
| 18 | 6 |
| 19 | 18 |
| 20 | 8 |

| n | φ(<i>n</i>) |
|----|---------------|
| 21 | 12 |
| 22 | 10 |
| 23 | 22 |
| 24 | 8 |
| 25 | 20 |
| 26 | 12 |
| 27 | 18 |
| 28 | 12 |
| 29 | 28 |
| 30 | 8 |

Euler's Theorem

 States that for every a and n that are relatively prime:

$$a^{\varnothing(n)} \equiv 1 \pmod{n}$$

An alternative form is:

$$a^{\emptyset(n)+1} \equiv a \pmod{n}$$

Note: The second form does not require a and n to be relatively prime.

Euler's Theorem Examples

```
a = 3; n = 10; \phi(10) = 4; a^{\phi(n)} = 3^4 = 81 \equiv 1 \pmod{10} \equiv 1 \pmod{n}
a = 2; n = 11; \phi(11) = 10; a^{\phi(n)} = 2^{10} = 1024 \equiv 1 \pmod{11} \equiv 1 \pmod{n}
```

Miller-Rabin Algorithm

Miller-Rabin Algorithm

- Typically used to test a large number for primality
- Algorithm is:

```
TEST (n)
```

- Find integers k, q, with k > 0, q odd, so that $(n 1) = 2^k q$;
- Select a random integer a, 1 < a < n 1;
- if a^q mod n = 1 then return ("inconclusive");
- for j = 0 to k 1 do
- if $(a^{2^{n}} \otimes n + n 1)$ then return ("inconclusive");
- return ("composite");

Miller-Rabin Algorithm

- The pseudocode below can be used to calculate k and q in the first step of the Miller-Rabin Algorithm
- Algorithm is:

```
Input (n)
```

```
1. • k \leftarrow 0;

• q \leftarrow (n-1);

• while ((q mod 2) ==0) /*While q is even*/

• \{k \leftarrow k+1;

• q \leftarrow q/2;\}

• return (k,q);
```

Deterministic Primality Algorithm

- Prior to 2002 there was no known method of efficiently proving the primality of very large numbers
- All of the algorithms in use produced a probabilistic result
- In 2002 Agrawal, Kayal, and Saxena developed an algorithm that efficiently determines whether a given large number is prime
 - Known as the AKS algorithm
 - Does not appear to be as efficient as the Miller-Rabin algorithm

Chinese Remainder Theorem (CRT)

- Believed to have been discovered by the Chinese mathematician Sun-Tsu in around 100 A.D.
- One of the most useful results of number theory
- Says it is possible to reconstruct integers in a certain range from their residues modulo a set of pairwise relatively prime moduli
- Can be stated in several ways

Provides a way to manipulate (potentially very large) numbers mod *M* in terms of tuples of smaller numbers

- This can be useful when *M* is 150 digits or more
- However, it is necessary to know beforehand the factorization of M

Chinese Remainder Theorem

• Given pairwise coprime positive integers n_1, n_2, \dots, n_k and arbitrary integers a_1, a_2, \dots, a_k the system of equations

$$x \equiv a_1 \ (mod \ n_1)$$

$$x \equiv a_2 \ (mod \ n_2)$$

$$x \equiv a_k \ (mod \ n_k)$$

has a unique solution for x.

Chinese Remainder Theorem

Algorithm to find a solution for x using the CRT:

- 1. Compute $N=n_1 \times n_2 \times \cdots \times n_k$.
- 2. For each $i=1,2,\ldots,k$, compute

$$y_i = rac{N}{n_i} = n_1 n_2 \cdots n_{i-1} n_{i+1} \cdots n_k.$$

- 3. For each $i=1,2,\ldots,k$, compute $z_i\equiv y_i^{-1} mod n_i$ using Euclid's extended algorithm (z_i exists since n_1,n_2,\ldots,n_k are pairwise coprime).
- 4. The integer $x=\sum_{i=1}^k a_iy_iz_i$ is a solution to the system of congruences, and $x \mod N$ is the unique solution modulo N.

$$x = 1 \quad \text{mod} 3$$

$$x = 3$$

$$x = 6 \quad \text{mod} 3$$

$$x = 3$$

$$X = \sum_{i=1}^{3} a_i y_i z_i = (1)(35)(2) + (4)(21)(1) + (6)(15)(1) =$$

$$70 + 84 + 90 = 244 = 34 \pmod{105}$$

Table 2.7
Powers of Integers, Modulo 19

| a | a^2 | a^3 | a^4 | a^5 | a^6 | a^7 | a^8 | a^9 | a ¹⁰ | a ¹¹ | a ¹² | a ¹³ | a^{14} | a ¹⁵ | a ¹⁶ | a ¹⁷ | a ¹⁸ |
|----|-------|-------|-------|-------|-------|-------|-------|-------|-----------------|-----------------|-----------------|-----------------|----------|-----------------|-----------------|-----------------|-----------------|
| 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| 2 | 4 | 8 | 16 | 13 | 7 | 14 | 9 | 18 | 17 | 15 | 11 | 3 | 6 | 12 | 5 | 10 | 1 |
| 3 | 9 | 8 | 5 | 15 | 7 | 2 | 6 | 18 | 16 | 10 | 11 | 14 | 4 | 12 | 17 | 13 | 1 |
| 4 | 16 | 7 | 9 | 17 | 11 | 6 | 5 | 1 | 4 | 16 | 7 | 9 | 17 | 11 | 6 | 5 | 1 |
| 5 | 6 | 11 | 17 | 9 | 7 | 16 | 4 | 1 | 5 | 6 | 11 | 17 | 9 | 7 | 16 | 4 | 1 |
| 6 | 17 | 7 | 4 | 5 | 11 | 9 | 16 | 1 | 6 | 17 | 7 | 4 | 5 | 11 | 9 | 16 | 1 |
| 7 | 11 | 1 | 7 | 11 | 1 | 7 | 11 | 1 | 7 | 11 | 1 | 7 | 11 | 1 | 7 | 11 | 1 |
| 8 | 7 | 18 | 11 | 12 | 1 | 8 | 7 | 18 | 11 | 12 | 1 | 8 | 7 | 18 | 11 | 12 | 1 |
| 9 | 5 | 7 | 6 | 16 | 11 | 4 | 17 | 1 | 9 | 5 | 7 | 6 | 16 | 11 | 4 | 17 | 1 |
| 10 | 5 | 12 | 6 | 3 | 11 | 15 | 17 | 18 | 9 | 14 | 7 | 13 | 16 | 8 | 4 | 2 | 1 |
| 11 | 7 | 1 | 11 | 7 | 1 | 11 | 7 | 1 | 11 | 7 | 1 | 11 | 7 | 1 | 11 | 7 | 1 |
| 12 | 11 | 18 | 7 | 8 | 1 | 12 | 11 | 18 | 7 | 8 | 1 | 12 | 11 | 18 | 7 | 8 | 1 |
| 13 | 17 | 12 | 4 | 14 | 11 | 10 | 16 | 18 | 6 | 2 | 7 | 15 | 5 | 8 | 9 | 3 | 1 |
| 14 | 6 | 8 | 17 | 10 | 7 | 3 | 4 | 18 | 5 | 13 | 11 | 2 | 9 | 12 | 16 | 15 | 1 |
| 15 | 16 | 12 | 9 | 2 | 11 | 13 | 5 | 18 | 4 | 3 | 7 | 10 | 17 | 8 | 6 | 14 | 1 |
| 16 | 9 | 11 | 5 | 4 | 7 | 17 | 6 | 1 | 16 | 9 | 11 | 5 | 4 | 7 | 17 | 6 | 1 |
| 17 | 4 | 11 | 16 | 6 | 7 | 5 | 9 | 1 | 17 | 4 | 11 | 16 | 6 | 7 | 5 | 9 | 1 |
| 18 | 1 | 18 | 1 | 18 | 1 | 18 | 1 | 18 | 1 | 18 | 1 | 18 | 1 | 18 | 1 | 18 | 1 |
| | | | | | | | | | | | | | | | | | |

The possible primitive roots of prime number 19 are
 2, 3, 10, 13, 14 and 15

Discrete Logarithms

 Let a be a primitive root of prime number p then for any integer b, we have

$$b \equiv a^i \pmod{p}$$
 where $0 \le i \le p-1$

- This exponent i is referred to as the discrete logarithm of b for the base a(mod p).
- We denote this value as $dlog_{a,p}(b)$.

Discrete Logarithms

- The discrete logarithm (dlog) i for an integer b to the base (a,p) is written as follows $dlog_{a,p}b = i$
- This implies that aⁱ mod p = b
- The integer p is a prime number and a is a primitive root of p
- Choosing a to be a primitive root of p ensures that a discrete logarithm value exists for values of b
- Discrete Logarithms are used in several cryptographic algorithms

Table 2.8

Tables of Discrete Logarithms, Modulo 19

(a) Discrete logarithms to the base 2, modulo 19

| а | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 |
|------------------|----|---|----|---|----|----|---|---|---|----|----|----|----|----|----|----|----|----|
| $\log_{2,19}(a)$ | 18 | 1 | 13 | 2 | 16 | 14 | 6 | 3 | 8 | 17 | 12 | 15 | 5 | 7 | 11 | 4 | 10 | 9 |

(b) Discrete logarithms to the base 3, modulo 19

| a | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 |
|------------------|----|---|---|----|---|---|---|---|---|----|----|----|----|----|----|----|----|----|
| $\log_{3,19}(a)$ | 18 | 7 | 1 | 14 | 4 | 8 | 6 | 3 | 2 | 11 | 12 | 15 | 17 | 13 | 5 | 10 | 16 | 9 |

(c) Discrete logarithms to the base 10, modulo 19

| a | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 |
|-------------------|----|----|---|----|---|---|----|----|----|----|----|----|----|----|----|----|----|----|
| $\log_{10,19}(a)$ | 18 | 17 | 5 | 16 | 2 | 4 | 12 | 15 | 10 | 1 | 6 | 3 | 13 | 11 | 7 | 14 | 8 | 9 |

(d) Discrete logarithms to the base 13, modulo 19

| | 1 | 2 | 2 | 4 | ~ | - | | 0 | 0 | 10 | 1.1 | 10 | 10 | 1.4 | 1.7 | 1.0 | 1.7 | 10 |
|-------------------|----|----|----|---|----|----|----|----|----|----|-----|----|----|-----|-----|-----|-----|----|
| a | I | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 |
| $\log_{13,19}(a)$ | 18 | 11 | 17 | 4 | 14 | 10 | 12 | 15 | 16 | 7 | 6 | 3 | 1 | 5 | 13 | 8 | 2 | 9 |

(e) Discrete logarithms to the base 14, modulo 19

| а | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 |
|-------------------|----|----|---|---|----|---|---|---|----|----|----|----|----|----|----|----|----|----|
| $\log_{14,19}(a)$ | 18 | 13 | 7 | 8 | 10 | 2 | 6 | 3 | 14 | 5 | 12 | 15 | 11 | 1 | 17 | 16 | 4 | 9 |

(f) Discrete logarithms to the base 15, modulo 19

| a | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 |
|-------------------|----|---|----|----|---|----|----|----|---|----|----|----|----|----|----|----|----|----|
| $\log_{15,19}(a)$ | 18 | 5 | 11 | 10 | 8 | 16 | 12 | 15 | 4 | 13 | 6 | 3 | 7 | 17 | 1 | 2 | 14 | 9 |

Summary

- Understand the concept of divisibility and the division algorithm
- Understand how to use the Euclidean algorithm to find the greatest common divisor
- Present an overview of the concepts of modular arithmetic
- Explain the operation of the extended Euclidean algorithm
- Discuss key concepts relating to prime numbers



- Understand Fermat's theorem
- Understand Euler's theorem
- Define Euler's totient function
- Make a presentation on the topic of testing for primality
- Explain the Chinese remainder theorem
- Define discrete logarithms