

Divisibility

- We say that a nonzero b **divides** a if a = mb for some m, where a, b, and m are integers
- b divides a if there is no remainder on division
- The notation b | a is commonly used to mean b divides a
- If b | a we say that b is a divisor of a

The positive divisors of 24 are 1, 2, 3, 4, 6, 8, 12, and 24 13 | 182; - 5 | 30; 17 | 289; - 3 | 33; 17 | 0

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Properties of Divisibility

- If $a \mid 1$, then $a = \pm 1$
- If $a \mid b$ and $b \mid a$, then $a = \pm b$
- Any b ≠ o divides o
- If a | b and b | c, then a | c

11 | 66 and 66 | 198 = 11 | 198

• If $b \mid g$ and $b \mid h$, then $b \mid (mg + nh)$ for arbitrary integers m and n

Properties of Divisibility

- To see this last point, note that:
 - If $b \mid g$, then g is of the form $g = b * g_t$ for some integer g_t
- If $b \mid h$, then h is of the form h = b * h, for some integer h,
- · So:
- mg + nh = mbg, + nbh, = b * (mg, + nh,) and therefore, b divides mg + nh

 $\begin{array}{l} b=7;\ g=14;\ h=63;\ m=3;\ n=2\\ 7\left|14\ \text{and}\ 7\right|63.\\ \text{To show}\ 7\left(3^*14+2^*63\right),\\ \text{we have}\ (3^*14+2^*63)=7(3^*2+2^*9),\\ \text{and it is obvious that}\ 7\left|\left(7(3^*2+2^*9)\right)\right|. \end{array}$

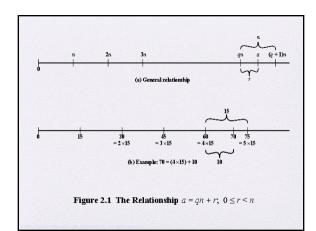
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Division Algorithm

 Given any positive integer n and any nonnegative integer a, if we divide a by n we get an integer quotient q and an integer remainder r that obey the following relationship:

a = qn + r $0 \le r < n; q = \lfloor a/n \rfloor$

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Euclidean Algorithm



- One of the basic techniques of number theory
- Procedure for determining the greatest common divisor of two positive integers
- Two integers are relatively prime if their only common positive integer factor is 1

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Greatest Common Divisor

- The greatest common divisor of a and b is the largest integer that divides both a and b
- We can use the notation gcd(a,b) to mean the greatest common divisor of a and b
- We also define gcd(0,0) = 0
- Positive integer c is said to be the gcd of a and b if:
 - c is a divisor of a and b
 - Any divisor of a and b is a divisor of c
- An equivalent definition is:

gcd(a,b) = max[k, such that k | a and k | b]

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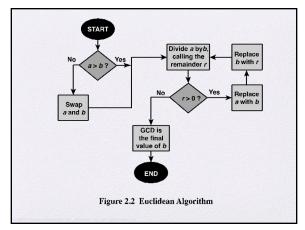
GCD

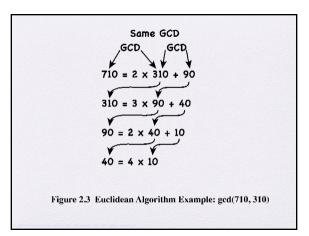
- Because we require that the greatest common divisor be positive, gcd(a,b) = gcd(a,b) = gcd(a,b) = gcd(-a,b) = gcd(-a,b)
- In general, gcd(a,b) = gcd(|a|, |b|)

gcd(60, 24) = gcd(60, -24) = 12

- Also, because all nonzero integers divide o, we have gcd(a,o) = | a |
- We stated that two integers a and b are relatively prime if their
 only common positive integer factor is 1; this is equivalent to
 saying that a and b are relatively prime if gcd(a,b) = 1

8 and 15 are relatively prime because the positive divisors of 8 are 1, 2, 4, and 8, and the positive divisors of 15 are 1, 3, 5, and 15. So 1 is the only integer on both lists.





	Tabl	e 2.1	
Euclid Dividend	ean Algo	rithm Ex	kample Remainder
a = 1160718174	b = 316258250	q ₁ = 3	r ₁ = 211943424
b = 316258250	r ₁ = 211943424	q ₂ = 1	r ₂ = 104314826
r ₁ = 211943424	r ₂ = 104314826	q ₃ = 2	r ₃ = 3313772
r ₂ = 104314826	r ₃ = 3313772	q ₄ = 31	r ₄ = 1587894
r ₃ = 3313772	r ₄ = 1587894	q ₅ = 2	r ₅ = 137984
r ₄ = 1587894	r ₅ = 137984	q ₆ = 11	r ₆ = 70070
r ₅ = 137984	r ₆ = 70070	q ₇ = 1	r ₇ = 67914
r ₆ = 70070	r ₇ = 67914	q ₈ = 1	r ₈ = 2156
r ₇ = 67914	r ₈ = 2156	q ₉ = 31	r ₉ = 1078
r ₈ = 2156	r ₉ = 1078	q ₁₀ = 2	r ₁₀ = 0

Modular Arithmetic

- The modulus
 - If a is an integer and n is a positive integer, we define a mod n to be the remainder when a is divided by n; the integer n is called the modulus
 - Thus, for any integer a:

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a = qn + r 0 \le r < n; q = [a/n]
a = [a/n] * n + (a \mod n)
```

11 mod 7 = 4; - 11 mod 7 = 3

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Modular Arithmetic

- Congruent modulo n
 - Two integers a and b are said to be congruent modulo n if (a mod n) = (b mod n)
 - This is written as $a \equiv b \pmod{n}$
 - Note that if $a \equiv o \pmod{n}$, then $n \mid a$

73 ≡ 4 (mod 23); 21 ≡ - 9 (mod 10)

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Properties of Congruences

- Congruences have the following properties:
 - 1. $a \equiv b \pmod{n}$ if $n \mid (a-b)$
 - 2. $a \equiv b \pmod{n}$ implies $b \equiv a \pmod{n}$
 - 3. $a \equiv b \pmod{n}$ and $b \equiv c \pmod{n}$ imply $a \equiv c \pmod{n}$
- To demonstrate the first point, if n (a b), then (a b) = kn for some k
 - So, we can write a = b + kn
 - Therefore, (a mod n) = (remainder when b + kn is divided by n) = (remainder when b is divided by n) = (b mod n)

23 = 8 (mod 5) because 23 - 8 = 15 = 5 * 3 -11 = 5 (mod 8) because -11 - 5 = -16 = 8 * (-2) 81 = 0 (mod 27) because 81 - 0 = 81 = 27 * 3

Modular Arithmetic

- Modular arithmetic exhibits the following properties:
 - 1. $[(a \mod n) + (b \mod n)] \mod n = (a + b) \mod n$
 - 2. $[(a \mod n) (b \mod n)] \mod n = (a b) \mod n$
 - 3. $[(a \mod n) * (b \mod n)] \mod n = (a * b) \mod n$
- We demonstrate the first property:
 - Define (a mod n) = r_a and (b mod n) = r_b. Then we can write a = r_a + jn for some integer j and b = r_b + kn for some integer k
 - Then

 $(a + b) \mod n = (ra + jn + rb + kn) \mod n$

- $= (ra + rb + (k + j)n) \mod n$
- = (ra + rb) mod n
- $= [(a \bmod n) + (b \bmod n)] \bmod n$

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Remaining Properties:

• Examples of the three remaining properties:

11 mod 8 = 3; 15 mod 8 = 7

[(11 mod 8) + (15 mod 8)] mod 8 = 10 mod 8 = 2

(11 + 15) mod 8 = 26 mod 8 = 2

[(11 mod 8) - (15 mod 8)] mod 8 = -4 mod 8 = 4

(11 - 15) mod 8 = -4 mod 8 = 4

[(11 mod 8) * (15 mod 8)] mod 8 = 21 mod 8 = 5

(11 * 15) mod 8 = 165 mod 8 = 5

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Modular Arithmetic Inverses

- Modular Additive Inverse
 - If (a + b) mod n = 0, then a and b are modular additive inverses of each other mod n.
- Modular Multiplicative Inverses
 - If (a * b) mod n = 1, then a and b are modular multiplicative inverses of each other mod n.

Modular Arithmetic Inverses

- Modular additive inverses can be used to carry out subtraction operations mod n
 - $(a b) \mod n = (a + (-b)) \mod n$, where -b is the modular additive inverse of b.
- Modular multiplicative inverses can be used to carry out division operations mod n
 - (a/b)mod $n = (a * (b^{-1}))$ mod n, where b^{-1} is the modular multiplicative inverse of b.

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			Гab	le 2	2.2(a)			
	Ar	rith	me	tic	Mc	du	10 8	3	
+	0	1	2	3	4	5	6	7	
0	0	1	2	3	4	5	6	7	
1	1	2	3	4	5	6	7	0	
2	2	3	4	5	6	7	0	1	
3	3	4	5	6	7	0	1	2	
4	4	5	6	7	0	1	2	3	
5	5	6	7	0	1	2	3	4	
6	6	7	0	1	2	3	4	5	
7	7	0	1	2	3	4	5	6	
						(This take			

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		Т	abl	e 2	.2(t)			
	Mu	ltip	lica	tior	Mo	odu	lo 8		
×	0	1	2	3	4	5	6	7	
0	0	0	0	0	0	0	0	0	
1	0	1	2	3	4	5	6	7	
2	0	2	4	6	0	2	4	6	
3	0	3	6	1	4	7	2	5	
4	0	4	0	4	0	4	0	4	
5	0	5	2	7	4	1	6	3	
6	0	6	4	2	0	6	4	2	
7	0	7	6	5	4	3	2	1	
a gibbi Branson Edge	ation ships then	Shirt Market In	gnie roservot.	13 1 3 1 1 1 1 1 1 1 1		This table can	be found on p	age 33 in the t	extbook)

	w	-w	w ⁻¹	
Table 2.2(c)	0	0	=	
	1	7	1	
Additive	2	6	_	
and	3	5	3	
Multiplicative	4	4		
Inverse	5	3	5	
Modulo 8	6	2	I	
	7	1	7	
© 2020 Pearson Education, Inc., Hoboken, NJ. All rights reserved.	(This ta	ble can be found	on page 33 in the	textbr

roperties of M	Table 2.3 Iodular Arithmetic for Integers in 2
roperties or ivi	loddidi 74 ft iiilede for iiilegers iii 2
Property	Expression
Commutative Laws	$(w+x) \bmod n = (x+w) \bmod n$ $(w\times x) \bmod n = (x\times w) \bmod n$
Associative Laws	$[(w+x)+y] \bmod n = [w+(x+y)] \bmod n$ $[(w\times x)\times y] \bmod n = [w\times (x\times y)] \bmod n$
Distributive Law	$[w \times (x + y)] \operatorname{mod} n = [(w \times x) + (w \times y)] \operatorname{mod} n$
Identities	$(0+w) \bmod n = w \bmod n$ $(1 \times w) \bmod n = w \bmod n$
	For each $w \in \mathbb{Z}_n$, there exists a z such that $w + z \equiv 0 \mod n$

Extend	ed Euclic	lean Alg	orithm E	xample
	and y such that ax			
i	r_i	q_i	x_i	Y _i
-1	1759		1	0
0	550		0	1
1	109	3	1	-3
2	5	5	-5	16
3	4	21	106	-339
4	1	1	-111	355
5	0	4		
Danultu d = 4	= -111; y = 355			

Extended Euclidean Algorithm

- The extended Euclidean algorithm can be used to calculate the modular multiplicative inverse
 - 1. Assume a > b
 - 2. If gcd(a,b) = 1 i.e. they are mutually prime
 - 3. Then the modular multiplicative inverse of $b \mod a = y$; y derived from solving for ax + by = d using the extended Euclidean algorithm
- From the previous example the modular multiplicative inverse of 550 mod 1759 is 355 since gcd(1759,550) = 1

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Prime Numbers

- Prime numbers only have divisors of 1 and itself
 - They cannot be written as a product of other numbers
- Prime numbers are central to number theory
- Any integer a > 1 can be factored in a unique way as

$$a = p_1^{a1} * p_2^{a2} * ... * p_t^{at}$$

where $p_1 \! < \! p_2 \! < \! \ldots < \! p_t$ are prime numbers and where each a_i is a positive integer

 This is known as the fundamental theorem of arithmetic

										ole	-								
							Pri	me	es I	Jno	der	20	00						
2	101	211	307	401	503	601	701	809	907	1009	1103	1201	1301	1409	1511	1601	1709	1801	190
3	103	223	311	409	509	607	709	811	911	1013	1109	1213	1303	1423	1523	1607	1721	1811	190
5	107	227	313	419	521	613	719	821	919	1019	1117	1217	1307	1427	1531	1609	1723	1823	191
7	109	229	317	421	523	617	727	823	929	1021	1123	1223	1319	1429	1543	1613	1733	1831	193
-11	113	233	331	431	541	619	733	827	937	1031	1129	1229	1321	1433	1549	1619	1741	1847	193.
13	127	239	337	433	547	631	739	829	941	1033	1151	1231	1327	1439	1553	1621	1747	1861	1949
17	131	241	347	439	557	641	743	839	947	1039	1153	1237	1361	1447	1559	1627	1753	1867	195
19	137	251	349	443	563	643	751	853	953	1049	1163	1249	1367	1451	1567	1637	1759	1871	197
23	139	257	353	449	569	647	757	857	967	1051	1171	1259	1373	1453	1571	1657	1777	1873	197
29	149	263	359	457	571	653	761	859	971	1061	1181	1277	1381	1459	1579	1663	1783	1877	198
31	151	269	367	461	577	659	769	863	977	1063	1187	1279	1399	1471	1583	1667	1787	1879	199
37	157	271	373	463	587	661	773	877	983	1069	1193	1283		1481	1597	1669	1789	1889	199
41	163	277	379	467	593	673	787	881	991	1087		1289		1483		1693			199
43	167	281	383	479	599	677	797	883	997	1091		1291		1487		1697			
47	173	283	389	487		683		887		1093		1297		1489		1699			
53	179	293	397	491		691				1097				1493					
59	181			499										1499					
61	191																		
67	193																		
71	197																		
73	199																		
79																			
83																			
89																			
97																			

Fermat's Theorem

- States the following:
 - If *p* is prime and *a* is a positive integer not divisible by *p* then

$$a^{p-1} \equiv 1 \pmod{p}$$

- An alternate form is:
 - If p is prime and a is a positive integer, then

$$a^p \equiv a \pmod{p}$$

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Table 2.6 Some Values of Euler's Totient Function g(n)

n	φ(<i>n</i>)	n	φ(<i>n</i>)
1	1	11	10
2	1	12	4
3	2	13	12
4	2	14	6
5	4	15	8
6	2	16	8
7	6	17	16
8	4	18	6
9	6	19	18
10	4	20	8

n	φ(<i>n</i>)
21	12
22	10
23	22
24	8
25	20
26	12
27	18
28	12
29	28
30	8

(This table can be found on page 44 in the textbook

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Euler's Theorem

• States that for every *a* and *n* that are relatively prime:

$$a^{\emptyset(n)} \equiv 1 \pmod{n}$$

An alternative form is:

$$a^{\varnothing(n)+1} \equiv a \pmod{n}$$

Miller-Rabin Algorithm Typically used to test a large number for primality Algorithm is: TEST (n) Test (n) Select a random integer a, 1 < a < n - 1; Test ("inconclusive"); Test (a^{2-max} mod n = 1 then return ("inconclusive"); Test (a^{2-max} mod n = n - 1) then return ("inconclusive"); Test (a^{2-max} mod n = n - 1) then return ("inconclusive");

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$\label{eq:matrix} \begin{tabular}{ll} \textbf{Miller-Rabin Algorithm} \\ \hline & \textbf{ The pseudocode below can be used to calculate k and q in the first step of the Miller-Rabin Algorithm} \\ \hline & \textbf{ Algorithm is:} \\ \hline & \textbf{ Input (n)} \\ \hline & \textbf{ 1. } & \textbf{ k} \leftarrow \textbf{ 0;} \\ \hline & \textbf{ 2. } & \textbf{ q} \leftarrow (\textbf{n-1}); \\ \hline & \textbf{ 3. } & \textbf{ while ((q mod 2) == 0) /*While q is even*/} \\ \hline & \textbf{ 4. } & \textbf{ k} \leftarrow \textbf{ k+1;} \\ \hline & \textbf{ 5. } & \textbf{ q} \leftarrow \textbf{ q/2;} \\ \hline & \textbf{ 6. } & \textbf{ return (k,q);} \\ \hline \end{tabular}$

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Deterministic Primality Algorithm

- Prior to 2002 there was no known method of efficiently proving the primality of very large numbers
- All of the algorithms in use produced a probabilistic result
- In 2002 Agrawal, Kayal, and Saxena developed an algorithm that efficiently determines whether a given large number is prime
- Known as the AKS algorithm
- Does not appear to be as efficient as the Miller-Rabin algorithm



Chinese Remainder Theorem (CRT)

- Believed to have been discovered by the Chinese mathematician Sun-Tsu in around 100 A.D.
- · One of the most useful results of number theory
- Says it is possible to reconstruct integers in a certain range from their residues modulo a set of pairwise relatively prime moduli
- Can be stated in several ways

Provides a way to manipulate (potentially very large) numbers mod *M* in terms of tuples of smaller

This can be useful when M is 150 digits or more
However, it is necessary to know beforehand the factorization of M

This can be useful when M is 150 digits or more
factorization of M

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Discrete Logarithms

- The discrete logarithm (dlog) i for or an integer b to the base (a,p) is written as follows dlog a,p b = i
- This implies that $a^i \mod p = b$
- The integer p is a prime number and a is a primitive root of p
- Choosing **a** to be a primitive root of **p** ensures that a discrete logarithm value exists for values of **b**
- Discrete Logarithms are used in several cryptographic algorithms

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(This table can be found on page 53 in the te

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2, 3, 10, 13, 14 and 15

							Т	ab	le 2	8.9								
			Т	able	es of	f Dis	cre	te L	oga	rithr	ns,	Мос	dulc	19				
				(a) E	iscre	te log	arith	ms to	the b	ase 2	, mod	lulo 1	9					
a	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18
$\log_{2.19}(a)$	18	1	13	2	16	14	6	3	8	17	12	15	5	7	11	4	10	9
				(b) I	iscre	te log	arith	ms to	the b	ase 3	, mod	lulo 1	9					
a	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18
$\log_{3.19}(a)$	18	7	1	14	4	8	6	3	2	11	12	15	17	13	5	10	16	9
a log _{10,19} (a)	1 18	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18
10210,19(4)				(d) D		e loga			the b	ase 13	3, mo	dulo 1	19					
		2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18
a	1																	
	18	11	17	4	14	10	12	15	16	7	6	3	1	5	13	8	2	9
	18	11		(e) D			3 19				6	17.5	9	5	13	8	2	9
	1	2	17	(e) D	iscret	e logs	3 19	ns to	the b	ase 14	6 1, mod	dulo 1	13	5	15	16	2	18
log _{13,19} (a)	1 18 18	11 2 13	17	(e) D	iscret	e loga	3 19	ns to	the b	ase 14	6 I, mo	dulo 1						
log _{13,19} (a)	1	2	17	(e) D	iscret	e loga	7 6	ns to	9 14	10 5	6 1, mod	12 15	13		15	16	17	18
log _{13,19} (a)	1	2 13	17 3 7	(e) D	iscret	e loga	7 6	ns to	9 14	10 5	6 1, mod	12 15	13		15	16 16	17 4	18

Understand the concept of divisibility and the division algorithm Understand how to use the Euclidean algorithm to find the greatest common divisor Present an overview of the concepts of modular arithmetic Explain the operation of the extended Euclidean algorithm Discuss key concepts relating to prime numbers Understand Fermat's theorem Understand Euler's totient function Make a presentation on the topic of testing for primality Explain the Chinese remainder theorem Define discrete logarithms