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Q3. (20 marks)

Question A

Answer:

```
Function find_most_frequent(k):
    most_frequency = 0                # 0(n)
    for i in k:                       # 0(1)
        count_for_now = 0            # 0(n)
        for j in k:                  # 0(1)
            if j == i:                # 0(1)
                times += 1            # 0(1)
            if times larger than most_frequency: # 0(1)
                most_frequency = times # 0(1)
    return most_frequency
```

Use two for loop. Suppose k = an array of integers and $\text{most_frequency} = 0$.
for i in k , using $\text{count_for_now} = 0$ to count the number i occurring in the array of integers. And then for j in k , if the number i is equal to the number j , then $\text{count_for_now} += 1$. After finish the loop of k , comparing the number count_for_now with the number most_frequency , and chose the number most_frequency is equal to the biggest one.
Return the most_frequency

Time complexity:

Two for loop, the first one: for i in k , is equal to $O(n)$, the second one: for j in k , is equal to $O(n)$. Therefore, the algorithm has $O(n^2)$ time complexity

Space complexity:

only need count_for_now and most_frequency has fixed length 2, so it is $O(1)$

Question B

```
Function find_most_frequent(k):
    Set an empty Dictionary, called frequent_one    # 0(1)
    For i in k:                                     # 0(n)
        If i in frequent_one:                       # 0(1)
            Value of key i in frequent_one + 1      # 0(1)
        else:
            value of key i in frequent_one = 1      # 0(1)
    sort the largest integer in frequent_one values # 0(1)
```

Create a dictionary, called $\text{frequent_one} = \{\}$. Suppose k = an array of

integers, for i in k if i exist in the frequent_one dictionary, then the value of that key plus one. Else, create an new key with starting value 1. Sort the largest integer in frequent_one dictionary values and find the biggest value and return the keys of that value

Time complexity:

for i in range(length(k)) is equal to $O(n)$. Therefore, $T(n) = O(n)$

Space complexity:

The frequent_one dictionary is fixed, therefore, so it is $O(1)$

Q4. (10 Marks)

Consider the following Pseudo-code that uses the Credit ADT from Q1.

```
function has_desc(accounts, desc):  
    for account in accounts:  
        for transaction in transactions(account):  
            if transaction matches desc:  
                return true  
    return false
```

Let n be the number of cards in the Credit object, and let m be the number of cards in the accounts list.

Give the worst-case time complexity for has_desc, under the following assumptions. You must justify

your answer.

a) Assuming m is in $O(1)$, and the cost of transactions(cc) is in $O(\log n)$

Answer:

$$T = (O(\log n) + O(1) * 10) * O(1) = O(\log n)$$

The cost of transactions(cc) is in $O(\log n)$, meaning it doesn't run in a constant time. Thus the inner loop is in $O(n)$

b) Assuming m is in $O(n)$ and the cost of transactions(cc) is in $O(n)$

Answer:

$$T = (O(n) + O(1) * 10) * O(n) = O(n^2)$$