The UniDyn Mathematica package: Self-derived unitary operator rotations in quantum mechanics

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Abstract

We have developed a *Mathematica* algorithm for symbolically calculating the unitary transformations of quantum-mechanical operators. The algorithm obtains closed-form analytical results, does not rely on a matrix representation of operators, and is applicable to both bounded systems like coupled spins and unbounded systems like harmonic oscillators. The rotations are "self derived" from the operators' underlying commutation relations. Example calculations are presented involving magnetic resonance and quantum optics.

1 Introduction

Given a time-independent Hamiltonian \mathcal{H} and an initial (density) operator $\rho(0)$, the Evolver algorithm described below implements the following unitary rotation:

$$\rho(t) = e^{-i\mathcal{H}t} \ \rho(0) \ e^{+i\mathcal{H}t} \equiv \text{Evolver}[\mathcal{H}, t, \rho(0)]$$
 (1)

Knowing the commutation relations among the operators comprising \mathcal{H} and $\rho(0)$, it is often possible to obtain a closed-form algebraic solution to Eq. 1. The Evolver algorithm obtains a closed-form solution to Eq. 1 by extending an approach introduced by Slichter in Section 2.3 of his *Principles of Magnetic Resonance* text [1]. The only input to the algorithm is the commutation relations for the relevant operators.

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Slichter was interested in calculating the unitary evolution of angular momentum operators arising in spin-physics problems. Below we introduce the Slichter procedure by considering a few representative examples. We show that his procedure works well for harmonic oscillator problems as well, including problems involving the coupling of a harmonic oscillator with a two-level system, a minimum model for describing electron transfer in chemical reactions. It will become apparent that automating Slichter's procedure in a computer algebra program like *Mathematica* is possible but challenging. We introduce a generalization of the Slichter algorithm that is well suited for automation by a computer algebra program. We have implemented this algorithm as a *Mathematica* function, Evolver. This function evaluates Eq. 1 given an \mathcal{H} , a $\rho(0)$, and the commutation relations between the operators comprising \mathcal{H} and $\rho(0)$.

2 The algorithm

To understand the Evolver algorithm it is essential to understand Slichter's approach to calculating the unitary evolution of an operator. To help the reader understand and appreciate Slichter's approach, let us briefly review two methods commonly used for solving Eq. 1. As an example, consider the case where $\mathcal{H} = \omega I_z$ and $\rho(0) = I_+ = I_x + i I_y$. In this case we want to compute

$$\rho(t) = e^{-i\omega t I_z} I_+ e^{+i\omega t I_z}. \tag{2}$$

Method 1 — One approach to computing $\rho(t)$ is to expand Eq. 2 in a Taylor series,

$$\rho(t) = \rho(0) + \rho^{(1)}(0) t + \frac{1}{2!} \rho^{(2)}(0) t^2 + \frac{1}{3!} \rho^{(3)}(0) t^3 + \cdots$$
 (3)

The coefficients are obtained by differentiating ρ and setting $t \to 0$. Taking the first derivative,

$$\rho^{(1)}(t) = e^{-i\omega t I_z} (-i\omega I_z) I_+ e^{+i\omega t I_z} + e^{-i\omega t I_z} I_+ (+i\omega I_z) e^{+i\omega t I_z}$$
(4a)

$$= e^{-i\omega t I_z} \left(-i\omega [I_z, I_+] \right) e^{+i\omega t I_z}$$
(4b)

$$\rho^{(1)}(0) = -i \,\omega \,I_{+} \tag{4c}$$

where $\rho^{(n)}$ represents the n^{th} derivative with respect to time and where we have used $[I_z, I_+] = I_+$ to simplify the commutator. Taking the second derivative,

$$\rho^{(2)}(t) = e^{-i\omega t I_z} \left((-i\omega)^2 [I_z, [I_z, I_+]] \right) e^{+i\omega t I_z}$$
(5a)

$$\rho^{(2)}(0) = (-i\,\omega)^2 I_+ \tag{5b}$$

By induction, we see that

$$\rho^{(n)}(0) = (-i\,\omega)^n \,I_+ \tag{6}$$

Substituting this finding into the Taylor expansion gives

$$\rho(t) = I_{+} \left(1 + (-i \omega t) + \frac{1}{2!} (-i \omega t)^{2} + \frac{1}{3!} (-i \omega t)^{3} + \cdots \right)$$
 (7)

We are now supposed to recognize the term in parenthesis as the Taylor series of $e^{-i\omega t}$. This insight enables us to resum the infinite series in the Taylor expansion to obtain the closed-form result

$$\rho(t) = I_+ e^{-i\omega t}. \tag{8}$$

Method 2 — A second approach to evaluating Eq. 2 is to expand the exponential using the Löwdin projection-operator theorem [2]. This theorem allows us to expand a function of an operator — I_z here — in terms involving the function evaluated at the operator's eigenvalues times an operator that project's onto the eigenvalue's subspace. Taking the total spin angular momentum to be I=1/2 for simplicity, the relevant eigenvalues are +1/2 and -1/2 and the relevant projection operators are $\mathcal{P}_{1/2}=|\alpha\rangle\langle\alpha|$ and $\mathcal{P}_{-1/2}=|\beta\rangle\langle\beta|$. Applying Löwdin's theorem,

$$e^{-i\,\omega t\,I_z} = e^{-i\,\omega t/2} \,|\alpha\rangle\,\langle\alpha| + e^{+i\,\omega t/2} \,|\beta\rangle\,\langle\beta|\,. \tag{9}$$

Substituting this result into Eq. 2 yields

$$\rho(t) = \left(e^{-i\,\omega t/2}\,|\alpha\rangle\,\langle\alpha| + e^{+i\,\omega t/2}\,|\beta\rangle\,\langle\beta|\right)I_{+}\left(e^{+i\,\omega t/2}\,|\alpha\rangle\,\langle\alpha| + e^{-i\,\omega t/2}\,|\beta\rangle\,\langle\beta|\right)$$

Applying the relations $I_+ |\alpha\rangle = 0$, $I_+ |\beta\rangle = |\alpha\rangle$, $\langle \alpha | \alpha \rangle = 1$, $\langle \beta | \alpha \rangle = 0$, and $|\alpha\rangle\langle\beta| = I_+$, this expression simplifies to

$$\rho(t) = I_+ e^{-i\omega t}. \tag{10}$$

While both these approaches yield closed-form solutions, each is hardly extensible. The first method requires the resumming of a Taylor series; this step would be difficult or impossible to automate for any evolution more complicated that the one above. The second method requires obtaining the eigenvalues of the Hamiltonian, usually by reducing it to matrix form and diagonalizing it. It is hard to see how to apply this diagonalization procedure in an infinite-level system like the idealized harmonic oscillator. Even for a spin problem where the number of levels is finite, diagonalizing $\mathcal H$ would force us to write down the matrix representation of $\mathcal H$ which would in turn commit us to specifying the total angular momentum I of the spin we are interested in. In many problems, we would like to obtain a solution valid for a spin of *any* I.

Method 3, Example 1 — Now consider Slichter's procedure. Let us take another look at the derivative of ρ :

$$\dot{\rho}(t) = e^{-i\,\omega t\,I_z} \left(-i\,\omega[I_z, I_+] \right) e^{+i\,\omega t\,I_z} \tag{11a}$$

$$= -i \omega \left(e^{-i \omega t I_z} I_+ e^{+i \omega t I_z} \right)$$
 (11b)

As before we have used $[I_z, I_+] = I_+$ to reduce the commutator in Eq. 11a. The key insight in the Slichter procedure is that the term in parenthesis in Eq. 11b is nothing more than the original time dependent density operator, $\rho(t)$. This insight allows us to write Equation 11b as

$$\dot{\rho}(t) = -i \,\omega \,\rho(t) \tag{12}$$

In this differential equation, ω is a *number*, while $\rho(t)$ is an *operator*. The solution to this differential equation is

$$\rho(t) = \rho(0) e^{-i \omega t} = I_{+} e^{-i \omega t}, \tag{13}$$

which is easily verified by back substitution.

Example 2 — An analogous calculation arises in a harmonic oscillator problem where the Hamiltonian is $\mathcal{H} = \omega(a^{\dagger}a + 1/2)$ and the initial operator is $\rho(0) = a^{\dagger}$:

$$\rho(t) = e^{-i\omega t(a^{\dagger}a + 1/2)} a^{\dagger} e^{+i\omega t(a^{\dagger}a + 1/2)}$$
(14)

Using the same procedure and the commutation relation $[a^{\dagger}a, a^{\dagger}] = a^{\dagger}[a, a^{\dagger}] + [a^{\dagger}, a^{\dagger}]a = a^{\dagger}$, this equation reduces to

$$\rho(t) = a^{\dagger} e^{-i \,\omega t}.\tag{15}$$

These first two Slichter-procedure example cases have in common that the commutator of the Hamiltonian with the operator of interest is simply proportional to the operator. As a result of this underlying commutation relation, the problem of calculating Eq. 1 in both cases has been reduced to the problem of solving a first-order differential equation, Eq. 12. In light of the two previous methods, the Slichter procedure is rather remarkable. It allows us to obtain a closed-form solution for Eq. 2 without resorting to Taylor series and without even requiring knowledge of the Hamiltonian's eigenvalues.

Example 3 — The Slichter procedure is readily applied to more complicated unitary-evolution problems. Consider the case where $\mathcal{H} = \omega I_z$ and $\rho(0) = I_x$. Then

$$\rho(t) = e^{-i\omega t I_z} I_x e^{+i\omega t I_z}. \tag{16}$$

Taking the time derivative we obtain

$$\dot{\rho}(t) = e^{-i\,\omega t\,I_z} \left(-i\,\omega[I_z, I_x]\right) e^{+i\,\omega t\,I_z} \tag{17a}$$

$$=\omega\left(e^{-i\,\omega t\,I_z}\,I_y\,e^{+i\,\omega t\,I_z}\right)\tag{17b}$$

where we have used $[I_z, I_x] = i I_y$ to reduce the commutator in Eq. 17a. In contrast to the previous case, $\dot{\rho}(t)$ is not proportional to $\dot{\rho}$. Taking another time derivative we obtain

$$\ddot{\rho}(t) = \omega e^{-i \omega t I_z} \left(-i \omega [I_z, I_y] \right) e^{+i \omega t I_z}$$
(18a)

$$= -\omega^2 \left(e^{-i\,\omega t\,I_z} \,I_x \,e^{+i\,\omega t\,I_z} \right) \tag{18b}$$

where we have used $[I_z, I_y] = -i I_x$ to reduce the commutator in Eq. 17a. The term in parenthesis in Eq. 18b is proportional to $\rho(t)$ and consequently

$$\ddot{\rho}(t) = -\omega^2 \, \rho(t). \tag{19}$$

The solution to this second-order differential equation is

$$\rho(t) = \rho(0)\sin(\omega t) + \frac{\dot{\rho}(0)}{\omega}\cos(\omega t). \tag{20}$$

We are given that $\rho(0) = I_x$ and we see from Eq. 17b that $\dot{\rho}(0) = \omega I_y$. Plugging these initial conditions into the above equation we obtain

$$\rho(t) = I_x \sin(\omega t) + I_y \cos(\omega t) \tag{21}$$

as the solution to Eq. 16.

We can imagine automating the steps leading to Eqs. 13, 14, and 21. What is required is the ability to

- 1. perform non-commutative algebra,
- 2. evaluate commutators, and
- 3. determine which differential equation an evolved operator satisfies.

Mathematica has a native function for carrying out non-commutative multiplication, but this function has essentially no simplification rules associated with it; for example,

This limitation was resolved beautifully by Helton and co-workers, whose NCALgebra package [3] gives *Mathematica* the ability to manipulate noncommuting algebraic expressions. This package allows the user to define variables as either commuting (e.g., ω) or non-commuting (e.g., I_+ , I_x , a^{\dagger} , and so on). This package forms an excellent starting point for implementing noncommutative multiplication, and expanding and simplifying commutators. We find, however, that using the package adds significant computational overhead, increasing computation time over ten-fold. In the code below, we instead introduce our own non-commutative multiplication function and accompanying simplification rules.

Based on our experience so far, developing an algorithm capable of self-deriving Eqs. 13, 14, and 21 would now seem to be a matter of calculating

$$\rho^{(1)} = U^{\dagger} [-i\mathcal{H}, \rho(0)] U \tag{22}$$

$$\rho^{(2)} = U^{\dagger} [-i\mathcal{H}, [-i\mathcal{H}, \rho(0)]] U = U^{\dagger} [-i\mathcal{H}, \rho^{(1)}] U$$
 (23)

and continuing until we obtain an expression that is proportional to $\rho = U^{\dagger} \rho(0) U$. If n iterations are required, then $\rho(t)$ must satisfy an n^{th} order differential equation. If the equation's coefficients and initial conditions can be extracted from the available non-commutative expressions correctly, then the differential equation can be fed to Mathematica to solve. This is roughly the procedure we will use. Because it would have to handle solving a large number of possible differential equations, however, significant effort would be required to implement the procedure just outlined in an automated way. The next example highlights why this is so. This example will serve as our launching point for introducing a revised, simple, and general algorithm for implementing the Slichter procedure for evaluating unitary evolution in an automated way.

Example 4 — Consider a unitary evolution with $\mathcal{H} = \Delta \omega I_z + \omega_1 I_x$ and $\rho(0) = I_z$. This rotation is involved in calculating the evolution of the difference in populations of a two level system when near-resonant irradiation is applied. The two-level system could be, for example, a spin I = 1/2 particle like a proton spin in a magnetic resonance experiment or it could be the lowest two electronic energy levels of an atom in a quantum optics experiment. The unitary evolution we are interested in computing is

$$\rho(t) = e^{-i(\Delta\omega I_z + \omega_1 I_x)t} I_z e^{+i(\Delta\omega I_z + \omega_1 I_x)t} \equiv U^{\dagger} I_x U$$
(24)

Computing the first few derivatives, we find

$$\dot{\rho} = U^{\dagger} \left(-i \left[\Delta \omega I_z + \omega_1 I_x, I_z \right] \right) U \tag{25a}$$

$$= U^{\dagger} (-\omega_1 I_y) U \tag{25b}$$

$$\ddot{\rho} = U^{\dagger} \left(-i \left[\Delta \omega I_z + \omega_1 I_x, -\omega_1, I_y \right] \right) U \tag{25c}$$

$$= U^{\dagger} (\Delta \omega \, \omega_1 I_x - \omega_1^2 I_z) U \tag{25d}$$

$$\ddot{\rho} = U^{\dagger} \left(-i \left[\Delta \omega I_z + \omega_1 I_x, -\omega_1, \Delta \omega \, \omega_1 I_x - \omega_1^2 I_z \right] \right) U \tag{25e}$$

$$= U^{\dagger} \left(\omega_1 \left(\omega_1^2 + \Delta \omega^2\right) I_y\right) U \tag{25f}$$

Neither $\dot{\rho}$ nor $\ddot{\rho}$ is proportional to ρ . We see, however, that $\ddot{\rho}$ is proportional to $\dot{\rho}$. Defining $\sigma \equiv \dot{\rho}$, we see that σ satisfies the following differential equation

$$\ddot{\sigma} = -(\Delta\omega^2 + \omega_1^2)\,\sigma\tag{26}$$

with initial conditions

$$\sigma(0) = \dot{\rho}(0) = -\omega_1 I_y \tag{27a}$$

$$\dot{\sigma}(0) = \ddot{\rho}(0) = \Delta\omega \,\omega_1 I_x - \omega_1^2 I_z \tag{27b}$$

The solution to Eq. 26 is

$$\sigma(t) = \sigma(0)\cos(\omega_{\text{eff}}t) + \frac{\dot{\sigma}(0)}{\omega_{\text{eff}}}\sin(\omega_{\text{eff}}t)$$
 (28)

with

$$\omega_{\text{eff}} \equiv \sqrt{\Delta \omega^2 + \omega_1^2} \tag{29}$$

an effective evolution frequency. Plugging in initial conditions,

$$\sigma(t) = \frac{\Delta\omega \,\omega_1}{\omega_{\text{eff}}} I_x \sin(\omega_{\text{eff}} t) - \omega_1 I_y \cos(\omega_{\text{eff}} t) - \frac{\omega_1^2}{\omega_{\text{eff}}} I_z \sin(\omega_{\text{eff}} t)$$
 (30)

To obtain an equation for $\rho(t)$ we need to integrate this equation for $\sigma(t)$, taking care to include a constant of integration. Noting that

$$\int \cos(\omega_{\rm eff} t) dt = \frac{\sin(\omega_{\rm eff} t)}{\omega_{\rm eff}} \text{ and } \int \sin(\omega_{\rm eff} t) dt = -\frac{\cos(\omega_{\rm eff} t)}{\omega_{\rm eff}},$$

the solution for $\rho(t)$ becomes

$$\rho(t) = -\frac{\Delta\omega \,\omega_1}{\omega_{\text{eff}}^2} I_x \sin(\omega_{\text{eff}} t) - \frac{\omega_1}{\omega_{\text{eff}}} I_y \sin(\omega_{\text{eff}} t) + \frac{\omega_1^2}{\omega_{\text{eff}}^2} I_z \cos(\omega_{\text{eff}} t) + c \quad (31)$$

The integration constant c is determined by requiring the above equation to satisfy $\rho(0) = I_z$, the initial condition. Solving for the integration constant,

$$c = \frac{\Delta\omega \,\omega_1}{\Delta\omega^2 + \omega_1^2} I_x + \frac{\Delta\omega^2}{\Delta\omega^2 + \omega_1^2} I_z \tag{32}$$

Plugging this integration constant into Eq. 31 we obtain the following solution for the density operator in Eq. 24:

$$\rho(t) = \left(\frac{\Delta\omega^2}{\Delta\omega^2 + \omega_1^2} I_z + \frac{\Delta\omega\omega_1}{\Delta\omega^2 + \omega_1^2} I_x\right) + \frac{\omega_1}{\sqrt{\Delta\omega^2 + \omega_1^2}} \sin\left(\sqrt{\Delta\omega^2 + \omega_1^2} t\right) + \left(\frac{\omega_1^2}{\Delta\omega^2 + \omega_1^2} I_z - \frac{\Delta\omega\omega_1}{\Delta\omega^2 + \omega_1^2} I_x\right) \cos\left(\sqrt{\Delta\omega^2 + \omega_1^2} t\right).$$
(33)

We see in this example an answer for $\rho(t)$ that is markedly more complicated that in the previous examples, the solution to (essentially) a third order differential equation. To uncover this differential equation, is was necessary to compare subsequent derivatives of $\rho(t)$ not to $\rho(0)$ but to $\dot{\rho}(0)$ instead.

Considering in total the four examples of unitary evolution considered so far, it would seem that any algorithm we develop needs to consider the possibility that the density operator satisfies a first, second, or even thirdorder differential equation.

Method 4 — We now show by example that the single differential equation in each of these four cases can be reduced to a set of *coupled first-order* differential equations — a significant step towards making a simple automated, unitary evolution algorithm. We would expect the four cases to each still require a decision on the number of coupled of equations required to solve for ρ in each case. Surprisingly, no decision is needed. We will show that all the cases we have discussed so far can be handled by a set of four coupled differential equations. This is the essential new insight implemented here in the Evolver algorithm.

To understand the new method, consider transforming the second-order differential equation in Eq. 19 into two coupled first-order equations. Let the two new variables be

$$x_1 = \rho \text{ and } x_2 = \dot{\rho}. \tag{34}$$

Taking the time derivative of each of these variables we obtain

$$\dot{x}_1 = \dot{\rho} = x_2, \text{ and} \tag{35a}$$

$$\dot{x}_1 = \ddot{\rho} = -\omega^2 \rho = -\omega^2 x_1.$$
 (35b)

It is apparent that these two variables satisfy the following set of coupled first order equations

$$\frac{d}{dt} \begin{bmatrix} x_2 \\ x_1 \end{bmatrix} = \begin{bmatrix} 0 & -\omega^2 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} x_2 \\ x_1 \end{bmatrix} \text{ with } \begin{bmatrix} x_2(0) \\ x_1(0) \end{bmatrix} = \begin{bmatrix} \omega I_y \\ I_x \end{bmatrix}. \tag{36}$$

Solving this set of coupled equations, we find $x_1(t) = \rho(t) = I_x \sin(\omega t) + I_y \cos(\omega t)$, the expected answer.

What if we did not know how many coupled equations were necessary to solve the problem? Let's consider what would happen if we guessed that three coupled equations were needed instead of two. Defining

$$x_1 = \rho \text{ and } x_2 = \dot{\rho} \text{ and } x_3 = \ddot{\rho}, \tag{37}$$

and taking the time derivative gives

$$\dot{x}_1 = \dot{\rho} = x_2,\tag{38a}$$

$$\dot{x}_2 = \ddot{\rho} = x_3, \text{ and} \tag{38b}$$

$$\dot{x}_3 = \ddot{\rho} \tag{38c}$$

$$= \omega^2 U^{\dagger} \left(-i \,\omega[I_z, I_x] \right) U \tag{38d}$$

$$= -\omega^3 U^{\dagger} I_y U \tag{38e}$$

$$= -\omega^2 \,\dot{\rho} \tag{38f}$$

$$= -\omega^2 x_2. \tag{38g}$$

In writing \dot{x}_3 we have used the shorthand $U(t) \equiv \exp[i \omega t I_z]$ and employed Eq. 17b to simplify the result. The three variables satisfy the following set of coupled first-order differential equations

$$\frac{d}{dt} \begin{bmatrix} x_3 \\ x_2 \\ x_1 \end{bmatrix} = \begin{bmatrix} 0 & -\omega^2 & 0 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix} \begin{bmatrix} x_3 \\ x_2 \\ x_1 \end{bmatrix} \text{ with } \begin{bmatrix} x_3(0) \\ x_2(0) \\ x_1(0) \end{bmatrix} = \begin{bmatrix} -\omega^2 I_x \\ \omega I_y \\ I_x \end{bmatrix}$$
(39)

Solving this new set of three coupled equations gives $x_1(t) = \rho(t) = I_x \sin(\omega t) + I_y \cos(\omega t)$ — the *same answer* that was obtained by solving the set of two coupled equations, Eqs. 36. This finding suggests that we are at liberty to "overguess" the number of equations required to solve the problem.

Let us see how the new method would be applied to solve the Example 4 problem with $\mathcal{H} = \Delta \omega I_z + \omega_1 I_x$ and $\rho(0) = I_z$. The time derivatives up to third order may be summarized as

$$\begin{bmatrix}
\rho^{(3)} \\
\rho^{(2)} \\
\rho^{(1)} \\
\rho^{(0)}
\end{bmatrix} = U^{\dagger} \begin{bmatrix}
-\Delta\omega(\Delta\omega^{2} + \omega_{1}^{2}) I_{y} \\
-\Delta\omega(\Delta\omega I_{x} - \omega_{1}I_{z}) \\
\Delta\omega I_{y} \\
I_{x}
\end{bmatrix} U$$
(40)

with $\rho^{(n)}$ the n^{th} derivative of ρ with respect to time and

$$U \equiv \exp\left[i\left(\Delta\omega I_z + \omega_1 I_x\right)t\right]. \tag{41}$$

We see that $\rho^{(3)} = -(\Delta \omega^2 + \omega_1^2) \, \rho^{(1)}$, as we showed before. Defining $x_n = \rho^{(n-1)}$ for n=1 through 4, we obtain the following set of four coupled equations describing the time evolution of the density operator:

$$\frac{d}{dt} \begin{bmatrix} x_4 \\ x_3 \\ x_2 \\ x_1 \end{bmatrix} = \underbrace{\begin{bmatrix} 0 & -(\Delta\omega^2 + \omega_1^2) & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}}_{\Omega} \underbrace{\begin{bmatrix} x_4 \\ x_3 \\ x_2 \\ x_1 \end{bmatrix}}_{\lambda} \tag{42}$$

with

$$\underbrace{\begin{bmatrix} x_4(0) \\ x_3(0) \\ x_2(0) \\ x_1(0) \end{bmatrix}}_{\lambda_0} = \begin{bmatrix} -\Delta\omega(\Delta\omega^2 + \omega_1^2) I_y \\ -\Delta\omega(\Delta\omega I_x - \omega_1 I_z) \\ \Delta\omega I_y \\ I_x \end{bmatrix}$$
(43)

Solving these coupled equations, we obtain Eq. 33 for $x_1 = \rho(t)$.

In the above equations we have defined λ as the vector of x_n 's we are interested in and Ω as the matrix of coefficients coupling λ to $\dot{\lambda}$. The coupled equations are summarized as

$$\frac{d\lambda}{dt} - \mathbf{\Omega} \cdot \lambda = 0 \tag{44}$$

with λ_0 given. The Evolver algorithm proceeds as follows.

1. Evaluate the derivatives $\rho^{(1)}$ through $\rho^{(3)}$ by repeatedly applying computing $\rho^{(n+1)} = [-iH, \rho^{(n)}]$, starting from $\rho^{(0)} = \rho(0)$. Assign the initial-condition vector $\lambda_0 = (\rho^{(3)}, \dots, \rho^{(0)})$.

- 2. Fill in the lower 3 rows of Ω by placing a 1 in entries one below the diagonal and 0 in all the other positions.
- 3. Look for an entry in the list $(\rho^{(2)}, \rho^{(1)}, \rho^{(0)})$ that is proportional to $\rho^{(3)}$, starting with the $\rho^{(2)}$ (entry n=3) and working towards $\rho^{(0)}$ (entry n=1). Call the entry where the match occurs n_{match} . Call the ratio between $\rho^{(3)}$ and the matching entry r.
- 4. Fill in the upper row of Ω with zeros. Assign the value r to the entry in the n_{match} column of the upper row of Ω , counting from the right to the left.
- 5. Feed Eq. 44, including the initial condition, to *Mathematica* to solve. The time-dependent density operator is $\rho(t) = \lambda_1(t)$.

3 The code

The code consists of .m files and *Mathematica* notebooks. The code was created in the literate-programming style, with LaTeX-style text, equations, and tables interspersed with code in the form of commented-out text. We use LaTeX's lstinputlisting package to import the code and typeset it nicely. This package also takes care of importing and un-commenting the embedded comments so they can be processed in LaTeX. The package moreover implements a long list of character substitutions, allowing the code listings below to have variables appear with subscripts and superscripts.

The code, at the time of writing, includes over 120 *unit tests* covering all the functions defined below. More detailed applications of the code are given in separate *Mathematica* notebooks and will be highlighted separately below.

3.1 Operators and scalars

Our first task is to define functions that enable *Mathematica* to distinguish between (non-commutative) operators and (commutative) scalars. This is done in the OpQ.m package, whose listing appears below. A variable is either a commutative *scalar* or non-commutative *operator*. Operators have a *phylum* and an *order*. Operators of different phyla commute. The operator order is used for sorting.

Listing 1: unidyn/OpQ.m

A *simple operator* is a symbol whose upvalue for SimpleOperatorQ[] is defined to be **True**. The default return for SimpleOperatorQ is **False**, with the result that unless defined otherwise, all arguments to the query SimpleOpertorQ[] return **False**. By default then, all quantities are scalars.

```
Clear[SimpleOperatorQ]
SimpleOperatorQ[x ] := False;
```

An *operator* is any expression one of whose atoms is a simple operator. The function call to **Level** in OperatorQ[] remarkably pulls out a list of all symbols used in the expression x (all its *atoms*.) The function ScalarQ[] tests if its argument is a *scalar*. By scalar we mean *not an operator*. As the default return for OperatorQ[] is **False**, the default return for ScalarQ[] is **True**; all quantities default to scalars.

You "create" an operator by defining an upvalue for it, e.g., arranging so that SimpleOperator[the operator] returns **True**. By using upvalues, you associate this definition not with the function SimpleOperator[], but with the operator itself. This makes for faster computations. The added definitions make it easy to create many operators at once by passing multiple variables or a list of variables to the function CreateOperator[].

```
Clear[CreateOperator];

CreateOperator[a_0_Symbol] := (Clear[a_0];

SimpleOperatorQ[a_0] ^:= True)

CreateOperator[a_0_,b_0__] := (CreateOperator[a_0];

CreateOperator[a_0];)

CreateOperator[a_0_?VectorQ] := (CreateOperator /@ a_0;)
```

You don't really need to "create" a scalar, since this is the default category for any symbol, given the above definitions. Nevertheless, by defining the upvalue of ScalarQ to be **True**, you can speed up computations which involve testing to see whether or not an object is a scalar.

```
 \begin{split} \textbf{Clear}[\text{CreateScalar}]; \\ \text{CreateScalar}[a_0\_\text{Symbol}] &:= \\ & (\textbf{Clear}[a_0]; \\ \text{Scalar}[a_0] \land := \textbf{True}) \\ \text{CreateScalar}[a_0\_\text{List}] &:= \\ & (\text{CreateScalar} \ / @ \ a_0;) \\ \text{CreateScalar}[a_0\_, b_0\_] &:= \\ & (\text{CreateScalar}[a_0]; \\ \text{CreateScalar}[b_0];) \\ \end{split}
```

Passing a matrix to SimpleOperator invokes the following function call. This function assigns the operators in the matrix a *phylum* and an *order* which is determined by the operators location in the matrix.

```
CreateOperator[a_0?ListQ] := 

Module[{val, m, n},

(Clear[#]; SimpleOperatorQ[#] ^= True) &
```

The idea of an operator having a *phylum* and an *order* can be understood best with an example. Consider, for example, the function call

```
CreateOperator[\{\{I_x, I_y, I_z\}, \{S_x, S_y, S_z\}\}]
```

This call will create six operators which, when passed to the functions phylum and order, will return the following values:

Ор	phylum[Op]	order[Op]
I_{χ}	1	1
I_{ν}	1	2
I_z	1	3
S_x S_y S_z	2	1
S_y	2	2
S_z	2	3

Having a phylum and an order assigned to each operator will be used below to sort operators and to decide whether or not two perators commute.

Unit tests for the functions in the OpQ.m package are organized into a separate package, OpQ-tests.m, whose listing follows. Unit testing is built on *Mathematica*'s VerificationTest function. When the return value of the VerificationTest function is passed to the TestReport in a *Mathematica* notebook, the notebook displays a nice graphical report card.

Listing 2: unidyn/OpQ-tests.m

Create a shorthand function for creating unit tests.

Any variable will test **True** when queried by ScalarQ[], whether it has been defined as a scalar or not.

```
vtest [" 1", ScalarQ[a_0] == True]
vtest [" 2", CreateScalar[b_0];
ScalarQ[b_0] == True]
```

For OperatorQ[] to return **True**, in contrast, the variable must have been created as an operator first

```
vtest [" 3", OperatorQ[c_0] == False]
vtest [" 4", CreateOperator[d_0];
OperatorQ[d_0] == True]
```

Any expression containing an operator will be queried **True** by OperatorQ[]. This is so when the expression contains only an operator or when the expression contains any combination of scalars and operators. The only case in which OperatorQ[] returns **False** is is it is passed an expression containing only scalars.

```
vtest [" 5", OperatorQ[Times[\mathbf{Exp}[d_0], d_0]] == \mathbf{True}] vtest [" 6", OperatorQ[Times[b_0, d_0]] == \mathbf{True}] vtest [" 7", OperatorQ[Times[a_0, b_0]] == \mathbf{False}]
```

Operators created in a batch using

CreateOperator[matrix]

are assigned a phylum and an order which is determined by their position in the matrix.

```
CreateOperator[\{\{I_x, I_y, I_z\}, \{S_x, S_y, S_z\}\}\}];
```

```
vtest [" 8", phylum[I_x] == 1]
vtest [" 9", order[I_x] == 1]
vtest [" 10", phylum[I_y] == 1]
vtest [" 11", order[I_y] == 2]
vtest [" 12", phylum[S_x] == 2]
vtest [" 13", order[S_x] == 1]
vtest [" 14", phylum[I_x]! = phylum[S_x]]
vtest [" 15", order[I_x] == order[S_x]]
```

Another, more stringent test of the *phylum* and *order* system:

```
CreateOperator[\{\{I_x, I_y, I_z\}, \{a, a^{\dagger}\}\}\}];
```

```
vtest ["16", Not[Mult[a, a^{\dagger}] === aa^{\dagger}]] vtest ["17", phylum /@\{I_x, I_y, I_z\} === \{1, 1, 1\}] vtest ["18", phylum /@\{a, a^{\dagger}\} === \{2, 2\}] vtest ["19", order /@\{I_x, I_y, I_z\} === \{1, 2, 3\}] vtest ["20", order /@\{a, a^{\dagger}\} === \{1, 2\}]
```

Very that the operators created in a big list are in fact recognized as operators.

```
vtest ["21", OperatorQ /@\{a_0, I_x, a\} == \{False, True, True\}]
```

4 Non-commutative multiplication

We next define our own non-commutative multiplication function, Mult. This is done in the Mult.m package. We provide a sorting function which rearranges operators being multiplied so that, if possible, operators of a higher *phylum* appear in front of operators of a lower *phylum*. During this rearrangement, non-commuting operators are not allowed to pass each other.

Listing 3: unidyn/Mult.m

First define bottom-out and empty cases:

```
Clear[Mult];

Mult[a_0] := a_0

Mult[] := 1
```

Mult distributes over Plus and is associative:

```
 \begin{split} & \mathsf{Mult}[a_0\_\_,b_0\_\mathsf{Plus},c_0\_\_] := \\ & \mathsf{Plus} @@ (\mathsf{Mult}[a_0,\,\#,\,c_0]\,\&\,) \ /@ \ \mathbf{List} @@ \ b_0 \\ & \mathsf{Mult}[a_0\_\_,\!\mathsf{Mult}[b_0\_\_,c_0\_\_],\!d_0\_\_] := \\ & \mathsf{Mult}[a_0,\!b_0,\!c_0,\!d_0] \end{aligned}
```

Define rules for factoring out scalars:

```
\begin{aligned} & \mathsf{Mult}[a_0\_\_, b_0\_?\mathsf{ScalarQ}, c_0\_\_] := \\ & \mathsf{Times}[b_0, \mathsf{Mult}[a_0, c_0]] \\ & \mathsf{Mult}[a_0\_\_, b_0\_?\mathsf{ScalarQ} \ c_0\_\_, d_0\_\_] := \\ & \mathsf{Times}[b_0, \mathsf{Mult}[a_0, c_0, d_0]] \end{aligned}
```

The function NCSort will sort the operators in a list into canonical order, being careful not to allow non-commuting operators to "pass" each other. The list that you pass to NCSort must contain *only* operators. How this functions works is probably easiest seen with an example. Suppose the following operators were initialized as follows:

CreateOperator[
$$\{\{I_x, I_y, I_z\}, \{S_x, S_y, S_z\}\}$$
]

Asked to sort the list of operators

$$a = \{S_x, I_y, I_x, S_y\},\$$

we would report the sorted list

$$a_{\text{sorted}} = \{I_y, I_x, S_x, S_y\}.$$

The S operator, being of a higher phylum that an I operator, should be passed through the I operators. We do not pass I_x to the left of I_y because I_x and I_y do not commute.

In the code below, the statement Map[phylum[#]&, a] will create a list populated by the operators' phyla, in this case $\{2,1,1,2\}$. To create p we multiply by the number of operators in the list plus one, 5 in this example, and add to this $\{1,2,3,4\}$. This gives $p=\{11,7,8,14\}$. An operator's ranking in this list depends on both its phylum and on it's location in the original list. Sorting p will tell us how to order the operators in the original list a. The variable p_{new} keeps track of the locations of the elements of p in the *sorted* version of p; in this example $p_{\text{new}} = \{3,1,2,4\}$. We recognize p_{new} as the order in which the input operators should appear in p_{new} as the order in which the input operators should appear in p_{new} as the order in which the input operators should appear in p_{new} as the order in which the input operators should appear in p_{new} as the order in which the input operators should appear in p_{new} as the order in which the input operators should appear in p_{new} and p_{new} as the order in which the input operators should appear in p_{new} as the order in which the input operators should appear in p_{new} as the order in which the input operators should appear in p_{new} as the order in which the input operators should appear in p_{new} as the order in which the input operators should appear in p_{new} as the order in which the input operators should be the order in p_{new} and p_{new} as the order in which the input operators should be the order in p_{new} as the order in p_{new} and p_{new} and p_{new} and p_{new} and p_{new} and p_{new} and p_{new} as the order in p_{new} and p_{new} an

```
\begin{split} & \mathsf{NCSort}[a_0\_\mathsf{List}] := \\ & \mathsf{Module}[\{n, p, a_{\mathrm{new}}, p_{\mathrm{new}}\}, \\ & n = \mathsf{Length}[a_0]; \\ & p = (n+1) \; \mathsf{Map}[\mathsf{phylum}[\#]\&, \, a_0] + \mathsf{Table}[i, \{i, 1, n\}]; \\ & p_{\mathrm{new}} = (\mathsf{Position}[\mathsf{Sort}[p], \#] \; [[1, 1]])\& \; \; /@ \; p; \\ & a_{\mathrm{new}} = \; \mathsf{Table}[0, \{i, 1, n\}]; \\ & \mathsf{Do}[a_{\mathrm{new}}[[p_{\mathrm{new}}[[i]]]] = a_0[[i]], \{i, 1, n\}]; \\ & \mathsf{Return}[a_{\mathrm{new}}]] \end{split}
```

The function SortedMult is the same as Mult except that it reorders the operators in the call list by applying NCSort before passing the result to Mult.

```
SortedMult[a_0__] := Mult[Sequence @@ NCSort[List[a_0]]]
```

The function MultSort reorders all the operators in a NonCommutativeMultiply call.

```
MultSort[a_0] := a_0 /. Mult \rightarrow SortedMult
```

Extensive unit testing shows that Mult distributes over addition, is associative, and pulls scalars out in front as desired. Testing is also carried out on the sorting function.

Listing 4: unidyn/Mult-tests.m

First define some operators and scalars.

```
CreateOperator[{{I_x, I_y, I_z},{S_x, S_y, S_z}];
CreateScalar[{a_0, b_0, c_0, d_0}];
```

In the following test, the left hand side should resolve to $\operatorname{Mult}[I_x,I_y]$ while the right-hand side should resolve to $\operatorname{Times}[I_x,I_y]$. These are *not* the same. If Mult does not properly recognize that I_x and I_y are operators and instead treats them as scalars, then the tests below will inadvertently pass.

```
vtest ["01", Not[Mult[I_x,I_y] === I_xI_y]]
```

Test that Mult distributes over addition and is associative. Note how Mult handles products of scalars and products of operators differently. Below we test for sameness (===) instead of equality (==). This is because if the two sides are *not* the same, then the equality test (using ==) is undefined — neither true nor false — and the unit test does not fail as we wish.

```
vtest ["02a", Mult[a_0, b_0+ c_0] === a_0b_0+ a_0c_0]
```

```
vtest ["02b", Mult[a_0 + b_0, c_0] === a_0c_0 + b_0c_0]
```

Test that Mult distributes over addition and is associative. Note how Mult handles products of scalars and products of operators differently. Below we test for sameness (===) instead of equality (==). This is because if the two sides are *not* the same, then the equality test (using ==) is undefined — neither true nor false — and the unit test does not fail as we wish.

```
vtest ["03", \operatorname{Mult}[I_y, I_x] \ a_0b_0 === a_0b_0\operatorname{Mult}[I_y, I_x]] vtest ["04a", \operatorname{Mult}[a_0, b_0 + c_0] === a_0b_0 + a_0c_0] vtest ["04b", \operatorname{Mult}[a_0 + b_0, c_0] === a_0c_0 + b_0c_0] vtest ["04c", \operatorname{Mult}[I_x, I_y + I_z] === \operatorname{Mult}[I_x, I_y] + \operatorname{Mult}[I_x, I_z]] vtest ["04d", \operatorname{Mult}[I_x + I_y, I_z] === \operatorname{Mult}[I_x, I_z] + \operatorname{Mult}[I_y, I_z]
```

Here is an example where we apparently do not have to call NCExpand[].

```
vtest ["05a", Mult[I_x, Mult[S_x, S_y], I_z] === Mult[I_x, S_x, S_y, I_z]]
```

Test that scalars get factored out properly.

```
vtest ["06a", Mult[I_x, 2 \mid a_0, S_x] === a_0 2 \mid Mult[I_x, S_x]]
vtest ["06b", Mult[I_x, Mult[a_0, I_y], Mult[b_0, S_x]] === a_0 b_0 Mult[I_x, I_y, S_x]]
```

Test our sorting function:

```
vtest ["07a", NCSort[\{S_x, I_x, I_y, S_z\}] === \{I_x, I_y, S_x, S_z\}]
```

The function MultSort[] is used to order the operators in a standing Mult[] function.

```
vtest ["08a", SortedMult[I_y, S_x, I_x] === Mult[I_y, I_x, S_x]] vtest ["08b", MultSort[Mult[I_y, S_x, I_x]] === Mult[I_y, I_x, S_x]]
```

If we now define the operators to have a different natural order, then the above tests fail. This confirms that the operators are being sorted according to our rules.

```
CreateOperator[{{S_x, S_y, S_z},{I_x, I_y, I_z}]; vtest ["08b", Not[SortedMult[I_y, S_x, I_x] === Mult[I_y, I_x, S_x]]] vtest ["08d", Not[MultSort[Mult[I_y, S_x, I_x]] === Mult[I_y, I_x, S_x]]]
```

Check that MultSort[] works as expected when scalars are peppered into the list of operators being multiplied.

```
vtest ["09a", MultSort[Mult[a_0 I_y, S_x, I_x]]

=== a_0Mult[S_x, I_y, I_x]]

expr1 = Mult[S_x, a_0S_z] + Mult[I_x, b_0S_x];

expr2 = MultSort[Mult[expr1, I_x]];

expr3 = a_0Mult[S_x, S_z, I_x] + b_0Mult[S_x, I_x, I_x];

vtest ["10a", expr2 === expr3]
```

5 Commutator

Define a commutator function which factors out scalars and has rules for simplifying more complex commutators, like and [AB, C] and [A, BC], involving three non-commuting operators.

Listing 5: unidyn/Comm.m

```
Comm[a_0_?ScalarQ, b_0_] := 0
    Comm[a_0\_, b_0\_?ScalarQ] := 0
    Comm[a_{0}, a_{0}] := 0
The commutator distributes over Plus:
    Comm[a_0___, b_0_Plus] := Plus @@ (Comm[a_0, #] & ) /@ List @@ b_0
    Comm[a_0_Plus, b_0_] := Plus @@ (Comm[#, b_0] & ) /@ List @@ a_0
Rules for factoring out scalars:
    Comm[a_0__ b_0_?ScalarQ, c_0__] := b_0Comm[a_0, c_0]
    Comm[a_0__, b_0_?ScalarQ c_0__] := b_0Comm[a_0, c_0]
Two rules for simplifying commutators involving products:
    Comm[a_0\_Mult, C_] :=
       Module[{A,B}],
          A = (List @@ a_0)[[1]];
          B = \text{Mult @@ Rest}[\text{List @@ } a_0];
          Mult[A, Comm[B, C]] + Mult[Comm[A, C], B]
    Comm[A_{-}, b_{0}_{-}Mult] :=
       Module [\{B,C\},
```

Finally, a commutator between different operators of a different phyla is zero:

```
Comm[a_0_,b_0_] := 0 /; phylum[a_0] != phylum[b_0];
```

 $C = \text{Mult @@ Rest}[\text{List @@ } b_0];$

Mult[Comm[A, B], C] + Mult[B, Comm[A, C]]

B = (List @@ $b_0)[[1]];$

1

Extensive testing shows that the commutator function factors out scalars correctly and correctly simplifies over a dozen commutation identities, including the Jacobi identity.

Listing 6: unidyn/Comm-tests.m

```
CreateScalar[\{a_0, b_0, c_0, d_0\}];
CreateOperator[\{\{I_x, I_y, I_z\}, \{S_x, S_y, S_z\}\}];
```

The commutator is distributive. The commutator of a scalar with an operator is zero. The commutator of an operator with iteslf is zero. The commutator between two operators of different phyla is zero. The commutator of two operators of the same phyla is left unevaluated.

```
vtest ["01a", Comm[I_x, I_y + I_z] === Comm[I_x, I_y] + Comm[I_x, I_z]] vtest ["01b", Comm[a_0, I_x] === 0] vtest ["01c", Comm[I_x, I_x] === 0] vtest ["01d", Comm[I_x, S_y] === 0] vtest ["01e", Comm[S_x, S_y] === Comm[S_x, S_y]]
```

Scalars are properly factored out of commutators.

```
vtest ["02a", Comm[a_0 S_x, S_y] === a_0Comm[S_x, S_y]] vtest ["02b", Comm[S_x, b_0 S_y] === b_0Comm[S_x, S_y]] vtest ["02c", Comm[a_0 S_x, b_0 S_y] === a_0 b_0Comm[a_0 S_x, b_0 S_y] === a_0 b_0Comm[a_0 S_x, b_0 S_y]
```

Test the [AB,C] and [A,BC] expansion rules, adding in scalars to both positions. In test 03b, we see that the S_x operator in the second place on the commutator can be pulled out front since it commutes with the operator in the first place of commutator.

```
 \begin{array}{l} \mathsf{vtest} \ ["03a", \ \mathsf{Comm}[a_0 \ \mathsf{Mult}[S_x, S_y], \ b_0 S_z] \\ === a_0 b_0 (\mathsf{Mult}[S_x, \ \mathsf{Comm}[S_y, S_z]] + \mathsf{Mult}[\mathsf{Comm}[S_x, S_z], \ S_y])] \\ \mathsf{vtest} \ ["03b", \ \mathsf{Comm}[a_0 \ S_x, b_0 \mathsf{Mult}[S_y, S_z]] \\ === a_0 b_0 (\mathsf{Mult}[S_y, \ \mathsf{Comm}[S_x, S_z]] + \mathsf{Mult}[\mathsf{Comm}[S_x, S_y], \ S_z])] \\ \mathsf{vtest} \ ["03c", \ \mathsf{Comm}[S_x, \ \mathsf{Mult}[S_x, S_y]] \\ === \mathsf{Mult}[S_x, \ \mathsf{Comm}[S_x, S_y]]] \\ \end{array}
```

Test the Jacobi identity. For this identity to resolve to zero, we must tell *Mathematica* that [A,B] equals A**B-B**A. First, check that this substitution does what we expect, then test the Jacobi identity.

```
 \begin{split} & \text{vtest} \ [\text{``04a''}, \ (\text{Comm}[I_x, I_y] \ //. \ \text{Comm}[S_{x\_}, S_{y\_}] \ \rightarrow \ \text{Mult}[S_x, S_y] \ - \ \text{Mult}[S_y, S_x]) \\ & === \ \text{Mult}[I_x, I_y] \ - \ \text{Mult}[I_y, I_x]] \\ & \text{vtest} \ [\text{``04b''}, \ ((\text{Comm}[I_x, \text{Comm}[I_y, I_z]] \ + \ \text{Comm}[I_y, \text{Comm}[I_z, I_x]] \ + \ \text{Comm}[I_z, \text{Comm}[I_z, I_y]]) \\ & //. \ \text{Comm}[S_{x\_}, S_{y\_}] \ \rightarrow \ \text{Mult}[S_x, S_y] \ - \ \text{Mult}[S_y, S_x]) \\ & === 0] \end{aligned}
```

Test the table of ten identities under "Advanced Identities" at the Wikipedia page for "Commutator", https://en.wikipedia.org/wiki/Commutator. I find that the first equality number 9 in "Advanced Identities" does not actually evaluate to True.

```
CreateOperator[\{\{A, B, C, D, E\$sym\}\}\}];
vtest ["05.01",
Comm[A, Mult[B, C]]
=== Mult[Comm[A, B], C]
+ Mult[B, Comm[A, C]]
```

```
vtest ["05.02",
Comm[A, Mult[B, C, D]]
 === Mult[Comm[A, B], C, D]
 + Mult[B, Comm[A, C], D]
 + Mult[B, C, Comm[A, D]]
]
vtest ["05.03",
Comm[A, Mult[B, C, D, E$sym]]
 === Mult[Comm[A, B], C, D, E$sym]
 + Mult[B, Comm[A, C], D, E$sym]
 + Mult[B, C, Comm[A, D], E$sym]
 + Mult[B, C, D, Comm[A, E$sym]]
]
vtest ["05.04",
Comm[Mult[A, B], C]
 === Mult[A, Comm[B, C]]
 + Mult[Comm[A, C], B]
vtest ["05.05",
Comm[Mult[A, B, C], D]
 === Mult[A, B, Comm[C, D]]
 + Mult[A, Comm[B, D], C]
 + Mult[Comm[A, D], B, C]
]
vtest ["05.06",
Comm[Mult[A, B, C, D], E$sym]
 === Mult[A, B, C, Comm[D, E\$sym]]
 + Mult[A, B, Comm[C, E$sym], D]
 + Mult[A, Comm[B, E$sym], C, D]
 + Mult[Comm[A, E$sym], B, C, D]
]
vtest ["05.07",
Comm[A, B+ C]
 === Comm[A, B]
 + Comm[A, C]
]
vtest ["05.08",
Comm[A + B, C + D]
 === Comm[A, C]
 + Comm[A, D]
 + Comm[B, C]
 + Comm[B, D]
]
vtest ["05.09",
```

```
Comm[Mult[A, B], Mult[C, D]]
 === Mult[A, Comm[B, C], D]
 + Mult[A, C, Comm[B, D]]
 + Mult[Comm[A, C], D, B]
 + Mult[C, Comm[A, D], B]
1
vtest ["05.10",
Module[{left, right },
  left = Comm[Comm[A, C], Comm[B, D]]
  //. Comm[A_?OperatorQ, B_?OperatorQ] \rightarrow Mult[A,B] – Mult[B,A];
  right = Comm[Comm[A, B], C], D]
 + Comm[Comm[Comm[B, C], D], A]
 + Comm[Comm[C, D], A], B]
 + Comm[Comm[D, A], B], C]
  //. Comm[A_?OperatorQ, B_?OperatorQ] \rightarrow Mult[A,B] – Mult[B,A];
  left === right]
1
```

6 Spins

Declare angular momentum operators for one spin I_x , I_y , and I_z and specify the relevant commutation relations. To speed up computation, each operator's commutation relations are stored as an *upvalue* of the operator instead of as a *downvalue* of the commutator function. Additional product-operator simplification rules are defined if the spin is declared to have an angular momentum of $\ell = 1/2$.

Listing 7: unidyn/Spins.m

```
SpinSingle$CreateOperators[I_{x}_,I_{y}_,I_{z}_,L_:Null] := Module[{nonexistent},
```

Test if the operators exist; if they do not already exist, then create them. If an operator Op has been created already, then CommutativeQ[Op] will return **True**. Unless I_x , I_y , and I_z all already exist as operators, then create all three operators afresh. In the code below it is important that we call Mult'CommutativeQ and not just CommutativeQ.

```
nonexistent = \mathbf{Not}[\mathsf{OperatorQ}[I_x]] \parallel \mathbf{Not}[\mathsf{OperatorQ}[I_y]] \parallel \mathbf{Not}[\mathsf{OperatorQ}[I_z]];

If [\mathsf{nonexistent} == \mathbf{True}, \mathbf{Clear}[I_x, I_y, I_z];
```

```
CreateOperator[\{\{I_x, I_y, I_z\}\}\}];
Message[SpinSingle$CreateOperators::create],
Message[SpinSingle$CreateOperators::nocreate];];
```

The commutation relations are defined as upvalues of the various spin angular momentum operators. That is, the commutation relations are associated with the operators and not with the Comm function. The following commutation relations hold for any L.

```
\begin{array}{ll} I_x \text{ /: } & \mathsf{Comm}[I_x,I_y] = \mathsf{I}\ I_z; \\ I_x \text{ /: } & \mathsf{Comm}[I_x,I_z] = -\mathsf{I}\ I_y; \\ I_y \text{ /: } & \mathsf{Comm}[I_y,I_x] = -\mathsf{I}\ I_z; \\ I_y \text{ /: } & \mathsf{Comm}[I_y,I_z] = \mathsf{I}\ I_x; \\ I_z \text{ /: } & \mathsf{Comm}[I_z,I_y] = -\mathsf{I}\ I_x; \\ I_z \text{ /: } & \mathsf{Comm}[I_z,I_x] = \mathsf{I}\ I_y; \end{array}
```

Message[SpinSingle\$CreateOperators::comm]

Switch[L,

When the total angular momentum L=1/2, additional rules are defined to simplify products of the angular momentum operators.

```
\begin{split} &I_{z} \: / : \quad \text{Mult}[a\_\_,I_{z},I_{z},b\_\_] := \text{Mult}[a,b]/4; \\ &I_{y} \: / : \quad \text{Mult}[a\_\_,I_{y},I_{y},b\_\_] := \text{Mult}[a,b]/4; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{x},b\_\_] := \text{Mult}[a,b]/4; \\ &I_{z} \: / : \quad \text{Mult}[a\_\_,I_{z},I_{x},b\_\_] := \text{I} \quad \text{Mult}[a,I_{y},b]/2; \\ &I_{z} \: / : \quad \text{Mult}[a\_\_,I_{z},I_{y},b\_\_] := -\text{I} \quad \text{Mult}[a,I_{x},b]/2; \\ &I_{y} \: / : \quad \text{Mult}[a\_\_,I_{y},I_{z},b\_\_] := \text{I} \quad \text{Mult}[a,I_{z},b]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{z},b\_\_] := -\text{I} \quad \text{Mult}[a,I_{y},b]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{y},b\_\_] := \text{I} \quad \text{Mult}[a,I_{z},b]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{y},b\_\_] := \text{I} \quad \text{Mult}[a,I_{z},b]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{y},b\_\_] := \text{I} \quad \text{Mult}[a,I_{z},b]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{y},b\_\_] := \text{I} \quad \text{Mult}[a,I_{z},b]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{y},b\_\_] := \text{I} \quad \text{Mult}[a,I_{z},b]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{y},b\_\_] := \text{I} \quad \text{Mult}[a,I_{z},b]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{y},b\_\_] := \text{I} \quad \text{Mult}[a,I_{z},b]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{y},b\_\_] := \text{I} \quad \text{Mult}[a,I_{z},b]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{y},b\_\_] := \text{I} \quad \text{Mult}[a,I_{z},b]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{y},b\_\_] := \text{I} \quad \text{Mult}[a,I_{z},b]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{y},b\_\_] := \text{I} \quad \text{Mult}[a,I_{z},b]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{y},b\_\_] := \text{I} \quad \text{Mult}[a,I_{x},b]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{y},b\_\_] := \text{I} \quad \text{Mult}[a,I_{x},b]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{y},b\_\_] := \text{I} \quad \text{Mult}[a,I_{x},b]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{y},b\_\_] := \text{I} \quad \text{Mult}[a,I_{x},b]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{y},b\_\_] := \text{I} \quad \text{Mult}[a,I_{x},b]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{x},I_{x},b\_\_] := \text{I} \quad \text{Mult}[a,I_{x},B]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{x},I_{x},b\_\_] := \text{I} \quad \text{Mult}[a\_\_,I_{x},B]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{x},I_{x},B]/2; \\ &I_{x} \: / : \quad \text{Mult}[a\_\_,I_{x},I_{x},I_{x},B]/2; \\ &I_{x} \:
```

Message[SpinSingle\$CreateOperators::simplify],

When the total angular momentum L is unspecified, no such simplification rules are defined.

```
Null, \label{lem:message} $\operatorname{SpinSingle\$CreateOperators::nosimplify}$ ]; $\operatorname{Return}[\{I_x,\,I_y,\,I_z\}]$ ]
```

Test the commutation relations for one spin and for a system of two spins, one of them $\ell = 1/2$ and the other with ℓ unspecified.

Listing 8: unidyn/Spins-tests.m

If there is one operator in the requested list of three spin operators that is not defined, then create all three spin operators afresh. Here we check that the test does what we want. If one of the operators is undefined, the test should come out false.

```
Clear[I_x, I_y, I_z];
CreateOperator[{{I_x, I_y}}];
tests = OperatorQ /@{I_x, I_y, I_z};
vtest ["00a", And @@ tests == False]
```

If, on the other hand, all three spin operators have been defined already, then the test should come out true.

```
Clear[I_x, I_y, I_z];
CreateOperator[{{I_x, I_y, I_z}}];
tests = OperatorQ /@{I_x, I_y, I_z};
vtest ["00b", And @@ tests == True]
```

We should also check the limiting case that *none* of the operators have been defined yet.

```
Clear[I_x, I_y, I_z];
tests = OperatorQ /@ {I_x, I_y, I_z};
vtest ["00c", And @@ Not /@ tests == True]
```

Create spin angular momentum operators with the total angular momentum unspecified. Test that the canonical angular momentum commutation relation holds true. With the total angular momentum unspecified, the product I_xI_y cannot be simplified further.

```
Clear[I_x, I_y, I_z]

SpinSingle$CreateOperators[I_x, I_y, I_z];

vtest ["01a", Comm[I_x, I_y] === | I_z]

vtest ["01b", Not[Mult[I_x, I_y] === Mult[I_y, I_x]]]

vtest ["01c", Not[Mult[I_x, I_y] === | I_z/2]]

vtest ["01d", Not[Mult[I_z, I_z] === 1/4]]
```

Create two sets of operators, one set for I spins and one set for S spins. Assign the I-spin operators the properties of I=1/2 angular momentum operators.

```
Clear[I_x, I_y, I_z, S_x, S_y, S_z];
CreateOperator[{{I_x, I_y, I_z},{S_x, S_y, S_z}}];
SpinSingle$CreateOperators[I_x, I_y, I_z, 1/2];
```

We have not defined the S operators to be spin operators yet. Nevertheless, the commutator of one S operator with another should be non-zero. On the other hand, the commutator of an S operator with an I operator should be zero.

```
vtest ["02a", Not[Comm[S_x, S_z] === 0]] vtest ["02b", Comm[I_z, S_z] === 0]
```

The canonical commutation relations hold true for the *I*-spin operators. In addition, the product of two spin angular momentum operators can be further simplified.

```
vtest ["03a", Comm[I_x, I_y] === | I_z]
vtest ["03b", Mult[I_x, I_y] === | I_z/2]
vtest ["03c", Mult[I_z, I_z] === 1/4]
```

Test that the double commutator of I_x with the free-evolution Hamiltonian returns the expected reult.

```
Clear[h, w, rho0, rho2];
CreateScalar[w];
h = w I_z;
rho0 = I_x;
rho2 = Comm[-I h, Comm[-I h,rho0]];
vtest ["04a", rho2 == - Mult[w, w, I_x]]
```

Try this test again with another spin operator on the backend.

```
Clear[I_x, I_y, I_z, S_x, S_y, S_z]
CreateOperator[{{I_x, I_y, I_z},{S_x, S_y, S_z}];
SpinSingle$CreateOperators[I_x, I_y, I_z];
SpinSingle$CreateOperators[S_x, S_y, S_z];

Clear[h, w, rho0, rho2];
CreateScalar[w];
h = w Mult[I_z, S_z];
rho0 = I_x;
rho2 = Comm[-I h, Comm[-I h,rho0]];
vtest ["04b", rho2 == - Mult[w, w, I_x, S_z, S_z]]
```

7 Harmonic oscillator

Declare the annihilation (or lowering) operator a and the creation (or raising) operator a^{\dagger} for a harmonic oscillator, and specify the relevant commutation relations.

Listing 9: unidyn/Osc.m

The commutation relations are defined as *upvalues* of the lowering and raising operators. Here aR = a, the lowering operator, and $aL = a^{\dagger}$, the raising operator.

```
OscSingle$CreateOperators[a_{-}, a^{\dagger}_{-}] := Module[{nonexistent},
```

Test if the operators exist; if they do not already exist, then create them.

```
nonexistent = 

Not[OperatorQ[a]] ||

Not[OperatorQ[a<sup>†</sup>]];
```

Check commutation relations among the harmonic-oscillator operators: a, a^{\dagger} , the number operator $N=a^{\dagger}a$, the unitless position operator $Q=(a^{\dagger}+a)/\sqrt{2}$, and the unitless momentum operator $P=i(a^{\dagger}-a)/\sqrt{2}$. Confirm that harmonic oscillator can exist "on top of" spin operators and are sorted correctly.

Listing 10: unidyn/Osc-tests.m

Create raising and lowering operators for a single harmonic oscillator. Check that the ooperators are indeed created. Check that they have the expected commutation relations.

```
Clear[a, a^{\dagger}, a_0, Nop];

CreateScalar[\{a_0\}];

OscSingle$CreateOperators[a, a^{\dagger}];

vtest ["01a", OperatorQ /@ \{a, a^{\dagger}\} == {True, True}]

vtest ["01b", Comm[a, a^{\dagger}] === 1]

vtest ["01c", Comm[a^{\dagger}, a] === -1]
```

Check that scalars are factored out of the commutations relations involving the harmonic oscillator raising and lower operators.

```
vtest ["01d", Comm[a_0 \ a, a^{\dagger}] === a_0]
vtest ["01e", Comm[a, a_0 a^{\dagger}] === a_0]
```

Test the commutation relations by defining the number operator, Nop = $N = a^{\dagger}a$, and checking the commutation relations $[N, a^{\dagger}] = a^{\dagger}$ and [N, a] = -a.

```
Nop = Mult[a^{\dagger}, a];

vtest ["01d", Comm[Nop, a^{\dagger}] === a^{\dagger}]

vtest ["01e", Comm[Nop, a] === -a]

Clear[a, a^{\dagger}, a_0, Nop];
```

Check that we can define harmonic oscillator operators "on top of" an existing operator that commutes with the harmonic-oscillator operators.

```
Clear[a, a^{\dagger}, a, Q];

CreateScalar[{a}];

CreateOperator[{{Q},{a^{\dagger}, a}}]

OscSingle$CreateOperators[a, a^{\dagger}];

vtest ["02a", Comm[a, a^{\dagger}] === 1]

vtest ["02b", Comm[a^{\dagger}, a] === -1]

vtest ["02c", Comm[a, Q] == 0]

vtest ["02d", Comm[a, a] == 0]

vtest ["02e", Comm[a, a] === 0]

vtest ["02e", MultSort[Mult[a, Q]] === Mult[a, a]]

Clear[a, a^{\dagger}, a, Q];
```

Test the commutations for the position and momentum operators.

```
Clear[a, a^{\dagger}, Q, P];

OscSingle$CreateOperators[a, a^{\dagger}];

Q = (a^{\dagger} + a)/\sqrt{2};

P = I (a^{\dagger} - a)/\sqrt{2};

vtest ["03a", Comm[Q, P] === I]

vtest ["03b", Comm[P, Q] === -I]

Clear[a, a^{\dagger}, Q, P];
```

Abother check that we can define harmonic oscillator operators "on top of" existing operators. In the following tests we prove that an operator like I_x (not defined as a spin operator, just an operator) commutes with one of the harmonic oscillator operators while the harmonic oscillator commutation relations are retained.

```
Clear[I_x, I_y, I_z, a, a^{\dagger}, Nop];
CreateOperator[{{I_x, I_y, I_z},{a^{\dagger}, a}}]
OscSingle$CreateOperators[a, a^{\dagger}];
Nop = Mult[a^{\dagger}, a];
vtest ["04a", Not[Comm[I_x, I_y] === 0]]
vtest ["04b", Comm[I_x, a] == 0]
vtest ["04c", Comm[Mult[I_x, I_y], a] == 0]
vtest ["04d", Comm[a, a^{\dagger}] === 1]
vtest ["04e", Comm[Nop, a^{\dagger}] === a^{\dagger}]
vtest ["04f", Comm[Nop, a] === -a]
```

We defined the $\it I$ operators to have higher precedence than the harmonic oscillator operators. Check that MultSort pulls the $\it I$ operators out front as expected.

```
vtest ["05", MultSort[Mult[a, I_x, a^{\dagger}, a, I_y]]
```

```
=== Mult[I_x, I_y, a, a^{\dagger}, a]]

Clear[I_x, I_y, I_z, a, a^{\dagger}, Nop];
```

8 Unitary Evolution

Implement the Evolver algorithm described above. The function can be passed an option to print out intermediate variables, which is helpful for debugging.

Listing 11: unidyn/Evolve.m

The evolution operator distribute over Plus and Mult.

```
 \begin{split} \textbf{Clear}[\textbf{Evolve}]; \\ \textbf{Evolve}[\mathcal{H}\_\_,\,t\_\_,\,\rho\_\textbf{Plus}] &:= \textbf{Plus} \ @@ \ (\textbf{Evolve}[\mathcal{H},\,t,\,\#]\&) \ \ /@ \ \textbf{List} \ @@ \ \rho \end{split}   \begin{split} \textbf{Evolve}[\mathcal{H}\_\_,\,t\_\_,\,\rho\_\textbf{Mult}] &:= \textbf{Mult} \ @@ \ (\textbf{Evolve}[\mathcal{H},\,t,\,\#]\&) \ \ /@ \ \textbf{List} \ @@ \ \rho \end{split}
```

Scalars in front of the density operator should be pulled out front.

```
Evolve[\mathcal{H}_{\underline{\phantom{M}}}, t_{\underline{\phantom{M}}}, Times[a_?ScalarQ, \rho_{\underline{\phantom{M}}}] := a Evolve[\mathcal{H}, t, \rho]
```

A test to see if all the terms in a sum commute with each other. It is important to include an AllCommutingQ::usage statement at the top of this package so that this function is available in the General' context of the notebook.

```
AllCommutingQ[\mathcal{H}_] := Module[{H$list, Comm$matrix}, If [Head[\mathcal{H}] === Plus, H$list = List @@ \mathcal{H}; Comm$matrix = Outer[Comm, H$list, H$list]; Return[And @@ ((# === 0)& /@ Flatten[Comm$matrix])], Return[False] ]
```

If all the terms in the Hamiltonian commute, then we may distribute the Evolve operator over the terms in the Hamiltonian.

```
Evolve[\mathcal{H}_?AllCommutingQ, t_, \rho_] := Mult @@ (Evolve[#, t, \rho]&) /@ List @@ \mathcal{H}
```

A function to coerce *Mathematica* into writing simpler looking expressions, from http://mathematica.stackexchange.com/questions/5403/how-to-get-fullsimplify-to-fully-simplify-my-expression-with-custom-complexity-f

```
VisualComplexity:=(Count[ToBoxes[#],Except[" "|"("|")", String], Infinity ]&)
```

The Evolver function follows. By default, the function will not print out intermediate results during the computation.

```
Options[Evolver] = {quiet \rightarrow True};

Evolver[\mathcal{H}_{-}, t_{-}, \rho_{0}_{-}, opts : OptionsPattern[]] :=

Module[{k, a$vect, q, r, r$value, X, x4, x3, x2, x1, time, system, sol},

Clear[\rho];

\rho[0] = \rho_{0};
```

Evalute the derivates of the density operator by repeatedly applying the commutator. Simplify them as much as possible. Getting the simplication right is tricky. A special function has to be fed to FullSimplify to get it to return useful results.

```
\begin{array}{l} \textbf{Do}[\\ \rho[\text{k+1}] = (-\text{I } \text{Comm}[\mathcal{H}, \rho[\text{k}]] \ /. \ \text{Mult} \rightarrow \text{SortedMult}) \\ \text{ } \text{ } \text{ } \text{FullSimplify } [\text{\#, ComplexityFunction} \rightarrow \text{VisualComplexity}]\&, \\ \text{ } \{\text{k}, 0, 4\}\\ ]; \end{array}
```

Print out the vector of density-operator derivatives if asked.

```
If [OptionValue[quiet] == False, Print[" \[Rho] matrix = ", \rho[#]& /@ {0,1,2,3,4} // MatrixForm]];
```

Look for an entry in the $(\rho^{(2)}, \rho^{(1)}, \rho^{(0)})$ list that is proportional to $\rho^{(3)}$. Stop when you find it. Determining *proportional to* is tricky. Here we use the ability of the NCALgebra package to compute a symbolic inverse of an operator. When $(\rho^{(n)})^{-1} * \rho^{(3)}$ is a scalar, then we have found a match. We are implicity assuming that the entries $(\rho^{(2)}, \rho^{(1)}, \rho^{(0)})$ have an inverse. This will be true of then entries involve Hermitian operators. If the entries involve *non-Hermitian operators*, however, like I_+ or I_- then the entries might not have a proper inverse. We do not, at present, test whether the operators in the list are Hermitian or not.

```
\label{eq:asymptotic_series} \begin{split} &a\$\text{vect}=\{0,0,0,0\};\\ &r = \text{Null};\\ &r = \text{Catch}[\\ &\textbf{Do}[\\ &q = \text{Mult}[\\ &&\text{Divide}[\rho[3], \text{Mult}[\rho[k]]]]\\ &&] \text{//} &\text{FullSimplify}[\#, \text{ComplexityFunction} \rightarrow \text{VisualComplexity}]\&\ ;\\ &&\text{If}\left[\text{ScalarQ}[q]==&\textbf{True}, \text{Throw}[\{3-k, q\}]],\\ &&\{k,0,2\}\\ &]\\ &]\\ &]; \end{split}
```

If we have not found a match by now, then throw up our hands and exit the function. Spit out the full $(\rho^{(3)}, \rho^{(2)}, \rho^{(1)}, \rho^{(0)})$ vector, in case the user can spot the *proportional to* condition.

```
If [r=== Null,
Message[Evolver::unsolvable];
Return[\rho[#]& /@ {0,1,2,3,4}],
```

```
r$value=r
];
```

Set up the coupling matrix Ω based on the matching condition.

```
a\$vect[[r\$value [[1]]]] = r\$value [[2]];
A = \{a\$vect,\{1, 0, 0, 0\},\{0, 1, 0, 0\},\{0, 0, 1, 0\}\};
```

If asked, spit out the coupling matrix for inspection.

```
If [OptionValue[quiet] == False, Print[" \[CapitalOmega] = ", A // MatrixForm]];
```

Set up the four coupled equations and solve them. Respect *Mathematica* version 8 standards here and feed DSolve[] a list where each element of the list is an equation. First set up the equations.

```
 \begin{split} & \text{X[time\_]} = \{\text{x4[time]}, \ \text{x3[time]}, \ \text{x2[time]}, \ \text{x1[time]}\}; \\ & \text{lhs\$list} = \text{D[X[time]}, \text{time]}; \\ & \text{rhs\$list} = \text{A} . \ \text{X[time]}; \\ & \text{eqns} = (\text{lhs\$list} \ [[\#]] \ == \text{rhs\$list} \ [[\#]])\& \ /@ \{1,2,3,4\}; \\ & \text{system} = \{\text{eqns}, \\ & \text{x4[0]} == \rho \text{[3]}, \ \text{x3[0]} == \rho \text{[2]}, \ \text{x2[0]} == \rho \text{[1]}, \ \text{x1[0]} == \rho \text{[0]}; \\ & \text{If [OptionValue[quiet]} == \textbf{False}, \ \text{Print[" system of equations} = ", \ \text{system} \ // \ \text{MatrixForm]}; \\ \end{aligned}
```

Now solve them and print out the solution with the private variables replaced by public ones.

```
sol = DSolve[system,\{x1,x2,x3,x4\},time]; 
 If [OptionValue[quiet] == False, Print[" 1st solution = ", sol [[1]][[1]] ]]; 
 If [OptionValue[quiet] == False, Print[" 1st solution w/ substitution = ", x1[time] /. 
 sol [[1]][[1]] /. time \rightarrow t]];
```

Return only the first element of the solution, $\lambda_1(t)$.

```
Return[FullSimplify[x1[time] /. sol[[1]] /. time \rightarrow t]];
```

Unit testing includes verifying that Evolver distributes over sums and factors our scalars correctly. Unit testing also verifies that *Mathematica's* differential equation solver is behaving as expected. We are not worried about bugs in the solver. We are instead concerned about changes in function syntax. The solver function's syntax changed between *Mathematica* version 8 and 10 and could change again. We added an equation-solver unit test to flag a future syntax change.

We test the following cases of unitary evolution in the rotating frame:

```
1. e^{-i\omega t I_z}I_x e^{+i\omega t I_z}
```

free evolution of transverse magnetization in a magnetic field

- 2. $e^{-i\omega t I_x}I_z e^{+i\omega t I_x}$ on-resonance nutation of longitudinal (i.e., thermal-equilibrium) magnetization in the rotating frame
- 3. $e^{-i\omega t I_z}I_+e^{+i\omega t I_z}$ free evolution of complex transverse magnetization in a magnetic field
- 4. $e^{-idtI_zS_z}I_xe^{+idtI_zS_z}$ evolution of transverse magnetization under a scalar coupling
- 5. $e^{-i(\Delta I_z + \omega I_x)t}I_z e^{+i(\Delta I_z + \omega I_x)t}$ off-resonance nutation of longitudinal magnetization
- 6. $e^{-\omega t \frac{1}{2}(a^{\dagger} a + a a^{\dagger})} a e^{+\omega t \frac{1}{2}(a^{\dagger} a + a a^{\dagger})}$ evolution of the lowering (and raising) operator under the harmonic oscillator Hamiltonian
- 7. $e^{-\omega t \frac{1}{2}(a^{\dagger} a + a a^{\dagger})} Q e^{+\omega t \frac{1}{2}(a^{\dagger} a + a a^{\dagger})}$ evolution of the position (and momentum) operator under the harmonic oscillator Hamiltonian
- 8. $e^{-i\delta q P} Q e^{+i\delta q P}$ a position kick
- 9. $e^{-i\delta q Q} P e^{+i\delta q Q}$ a momentum kick

The spin off-resonance nutation, case 5, is an especially demanding test because it generates terms involving I_x , I_y , and I_z operators from solutions to a third-order differential equation.

Listing 12: unidyn/Evolve-tests.m

Global debugging flag to force Evolver to be noisy instead of quiet. Set to False for noisy output and to True for quiet output.

quiet\$query = **True**;

Show that the Evolve operator is distributive over both sums and non-commutative products of operators. Show that scalars are pulled out front, even if buried in a oomplicated product of operators and scalars. For this test, we'll make three sets of two commuting operators representing, for example, three independent harmonic oscillators.

$$\label{eq:createOperator} \begin{split} & CreateOperator[\{\{Q,R\},\{S,U\},\{V,W\}\}]\\ & CreateScalar[\{q,r,s,u,v,w\}] \end{split}$$

```
 \begin{array}{l} \text{vtest} \, [\text{``01a} > \text{distribute addition''}, \\ \text{Evolve}[H,\,t,\,\,Q+R+S] \\ \text{===} \, \text{Evolve}[H,\,t,\,\,Q] + \text{Evolve}[H,\,t,\,\,R] + \text{Evolve}[H,\,t,\,\,S]] \\ \text{vtest} \, [\text{``01b} > \text{distribute multiplication''}, \\ \text{Evolve}[H,\,t,\,\,\text{Mult}[Q,\,R,\,S]] \\ \text{===} \, \text{Mult}[\text{Evolve}[H,\,t,\,\,Q],\,\,\text{Evolve}[H,\,t,\,\,R],\,\,\text{Evolve}[H,\,t,\,\,S]]] \\ \text{vtest} \, [\text{``01c} > \text{distribute complicated expression''}, \\ \text{Evolve}[H,\,t,\,\,\text{Mult}[(Q\,q),\,\,(r\,R),\,\,(s\,S)] + u\,\,U] \\ \text{===} \, u\,\,\text{Evolve}[H,\,t,\,\,U] + q\,\,r\,\,s\,\,\text{Mult}[\text{Evolve}[H,\,t,\,\,Q],\,\,\text{Evolve}[H,\,t,\,\,R],\,\,\text{Evolve}[H,\,t,\,\,S]]] \\ \end{array}
```

Make up some Hamiltonians and see if they pass the all-terms-commuting test. The first test is particularly important. In test 03a and 03b below, nothing happens; the Hamiltonian is either so simple that it can be broken into pieces, 03a, or contains terms which do not commut, 03b. In test 03c we have a Hamiltonian whose three terms commute, and in this case the Evolve operator can be expanded.

```
\begin{split} &H0=Q;\\ &H1=q\;\text{Mult}[Q,\,R]+s\;\text{Mult}[S,\,U]+\text{Mult}[U,\,S]+\text{Mult}[V,\,V,\,W];\\ &H2=q\;Q+s\;S+v\;\text{Mult}[Q,\,S];\\ &\text{vtest}\; [\text{``02a}>\text{commuting}\;\text{test}\;1\text{''},\,\text{AllCommuting}Q[H0]===&\,\textbf{False}]\\ &\text{vtest}\; [\text{``02b}>\text{commuting}\;\text{test}\;2\text{''},\,\text{AllCommuting}Q[H1]===&\,\textbf{False}]\\ &\text{vtest}\; [\text{``02c}>\text{commuting}\;\text{test}\;3\text{''},\,\text{AllCommuting}Q[H2]===&\,\textbf{True}]\\ &\text{vtest}\; [\text{``03a}>\text{Evolve}\;\text{expand}\;\text{test}\;1\text{''},\,\text{Evolve}[H0,\,t,\,Q]===&\,\text{Evolve}[Q,\,t,\,Q]]\\ &\text{vtest}\; [\text{``03b}>\text{Evolve}\;\text{expand}\;\text{test}\;2\text{''},\,\text{Evolve}[H1,\,t,\,Q]\\ &==&\,\text{Evolve}[q\;\text{Mult}[Q,R]+s\;\text{Mult}[S,U]+\text{Mult}[U,\,S]+\text{Mult}[V,\,V,\,W],\,t,\,Q]]\\ &\text{vtest}\; [\text{``03c}>\text{Evolve}\;\text{expand}\;\text{test}\;3\text{''},\,\text{Evolve}[H2,\,t,\,Q]\\ &===&\,\text{Mult}[\text{Evolve}[q\;Q,\,t,\,Q],\,\text{Evolve}[s\;S,\,t,\,Q],\,\text{Evolve}[v\;\text{Mult}[Q,S],\,t,\,Q]]]\\ &\text{Clear}[H0,\,H1,\,H2]\\ &\text{Clear}[Q,\,R,\,S,\,U,\,V,\,W,\,q,\,r,\,s,\,u,\,v,\,w] \end{split}
```

To test the unitary evolution operator on spins, let us set up a model two-spin system. This system is comprised of an L = 1/2 I spin and an unspecified-L S spin.

```
Clear[I_x, I_y, I_z, S_x, S_y, S_z, \omega, d_0, \Delta, \rho, t]
Clear[A, r, lhs$list, rhs$list, time, eqns, system, \rho_{\text{calc}}, \rho_{\text{known}}, X]
CreateScalar[\omega, d_0, \Delta];
CreateOperator[{{I_x, I_y, I_z},{S_x, S_y, S_z}}]
SpinSingle$CreateOperators[I_x, I_y, I_z, L=1/2];
SpinSingle$CreateOperators[S_x, S_y, S_z, L=1/2];
```

Test the differential equation solver first. In *Mathematica* version 10 gives more leeway in how the equations are set up – you can set one list equal to another, for example. In *Mathematica* version 8, in contrast the syntax is not so forgiving. Let us mock-up an equation by hand and feed it to the solver.

```
A = {{0,-\triangle ^2,0,0},{1,0,0,0},{0,1,0,0},{0,0,1,0}};

r = {I_x, I_y\triangle, -I_x \triangle^2, -I_y \triangle^3, I_x\triangle^4};
```

```
(\rho[\#-1] = r[[\#]])& /@\{1,2,3,4,5\};
      X[time] = \{x4[time], x3[time], x2[time], x1[time]\};
      lhs = D[X[time], time];
      rhs$list = A . X[time];
      eqns = ( lhs\{list [[#]] == rhs\{list [[#]])\} /@ \{1,2,3,4\};
      system = {eqns,
        x4[0]== \rho[3], x3[0]== \rho[2], x2[0]== \rho[1], x1[0]== \rho[0];
        sol = DSolve[system, \{x1, x2, x3, x4\}, time];
      \rho_{\text{calc}}= (x1[time] /. sol[[1]] /. time \rightarrow t) // Expand // ExpToTrig // FullSimplify ;
      \rho_{\mathsf{known}} = I_x \mathsf{Cos}[t \ \Delta] + I_y \mathsf{Sin}[t \ \Delta];
      vtest ["04a > DSolve test", \rho_{calc} === \rho_{known}]
Free evolution of I_x:
      vtest ["05a > free evolution of Ix",
         FullSimplify [ExpToTrig[Expand[
           Evolver[\omega I_z, t, I_x, quiet \rightarrow quiet$query]]]]
        ===I_x \mathbf{Cos}[\omega \ t] + I_y \mathbf{Sin}[\omega \ t]]
On-resonance nutation of I_z:
      vtest ["05b > on-resonance nutation of Iz",
         FullSimplify [ExpToTrig[Expand[
           Evolver[\omega I_x, t, I_z, quiet \rightarrow quiet$query]]]]
        ===I_z \mathbf{Cos}[\omega \ t] - I_y \mathbf{Sin}[\omega \ t]]
Free evolution of I_+:
      vtest ["05c > free evolution of I+",
         FullSimplify [Evolver[\omega I_z, t, I_x, quiet \rightarrow quiet$query]
           + I Evolver[\omega I_z, t, I_y, quiet \rightarrow quiet$query]]
        === \mathbf{Exp}[-\mathsf{I}\ \omega t](I_x + \mathsf{I}\ I_y)]
Evolution under a scalar coupling:
      vtest ["05d > scalar-coupling evolution of Ix",
        Evolver[d_0 Mult[I_z, S_z], t, I_x, quiet \rightarrow quiet$query]
        ===I_x \mathbf{Cos}[d_0 \ t/2] + 2 \ \mathbf{Mult}[I_y, S_z] \ \mathbf{Sin}[d_0 \ t/2]]
Off-resonance nutation of I_z. It is important to carefully set the assumptions used by Simplify.
      Assumptions = \{Element[\Delta, Reals], \Delta > 0, Element[\omega, Reals], \omega > = 0\};
      c_1 = (\Delta^2)/(\Delta^2 + \omega^2);
      c_2 = (\Delta \omega)/(\Delta^2 + \omega^2);
      c_3 = (\omega^2)/(\Delta^2 + \omega^2);
      c_4 = \omega/\text{Sqrt}[\Delta^2 + \omega^2];
```

```
\begin{aligned} &\omega_{\text{eff}} = \text{Sqrt}[\Delta^{\wedge}2 + \omega^{\wedge}2]; \\ &\rho_{\text{known}} = \text{Collect}[ \\ &c_1I_z + c_2I_x + (c_3\ I_z - c_2I_x)\ \mathbf{Cos}[\omega_{\text{eff}}\ t] - c_4I_y\mathbf{Sin}[\omega_{\text{eff}}\ t] \ // \\ & \text{Expand, } \{I_x,\ I_y,\ I_z\}]; \\ &\rho_{\text{calc}} = \text{Collect}[ \\ & \text{Evolver}[\Delta\ I_z + \omega I_x,\ t,\ I_z,\ \text{quiet}\ \to \text{quiet}\$\text{query}] \\ &// & \text{FullSimplify}\ ,\ \{I_x,\ I_y,\ I_z\},\ \text{Expand}]; \\ &\text{vtest}\ [\text{``05e}\ > \text{Off-resonance nutation of of Iz''},\ \rho_{\text{calc}} === \rho_{\text{known}}] \end{aligned}
```

Clear[A, r, lhs\$list, rhs\$list, time, eqns, system, ρ_{calc} , ρ_{known} , X]

Harmonic oscillator evolution. First, create the harmonic oscillator Hamiltonian in symmetric form. Evolve the lowering operator and confirm that it picks up the expected phase factor. Evolve the raising operator and confirm that it picks up the expected (conjugate) phase factor.

```
Clear[a, a^{\dagger}, \omega, Q, P, \mathcal{H}, Q, P, \delta q, \delta p];

CreateScalar[\omega, delta$x$sym, \delta p];

OscSingle$CreateOperators[a, a^{\dagger}];

\mathcal{H} = \omega(\mathsf{Mult}[a, a^{\dagger}] + \mathsf{Mult}[a^{\dagger}, a])/2;

vtest ["06a1 > free evolution of lowering operator",

Simplify [TrigToExp[Expand[Evolver[\mathcal{H}, t, a, quiet \rightarrow quietquery]]]]

=== a \text{ Exp}[l \ \omega t \ ]]

vtest ["06a2 > free evolution of raising operator",

Simplify [TrigToExp[Expand[Evolver[\mathcal{H}, t, a^{\dagger}, quiet \rightarrow quietquery]]]]

=== a^{\dagger} \text{ Exp}[-l \ \omega t \ ]
```

Evolve the position operator. To do this, write the position operator in terms of the raising and lowering operators, calculate the evolution, and rewite the answer in terms of the position and momentum operator. For the re-write step, create a *new* set of non-commuting operators, to avoid chasing our tail.

```
CreateOperator[{{Q,P}}];

{Q, P} = {(a^{\dagger} + a)/\sqrt{2}, | (a^{\dagger} - a)/\sqrt{2}};

QP$rules = {a^{\dagger} \rightarrow (Q - | P)/\sqrt{2}, a \rightarrow (Q + | P)/\sqrt{2}};
```

Write Q and P in terms of raising and lower operators, write the raising and lowering operators in terms of position and momentum, and you should get the original Q and P back again.

```
vtest ["06b > test Q definition ", Simplify [Q /. QP$rules] === Q] vtest ["06c > test P definition ", Simplify [P /. QP$rules] === P]
```

Now we are ready to evolve position and momentum.

```
vtest ["06d > free evolution of Q", 
Simplify [ExpToTrig[Expand[Evolver[\mathcal{H}, t, Q, quiet \rightarrow quiet$query] /. QP$rules]]] 
== Q \mathbf{Cos}[\omega\ t] - P \mathbf{Sin}[\omega\ t]] 
vtest ["06e > free evolution of P", 
Simplify [ExpToTrig[Expand[Evolver[\mathcal{H}, t, P, quiet \rightarrow quiet$query] /. QP$rules]]] 
== P \mathbf{Cos}[\omega\ t] + Q \mathbf{Sin}[\omega\ t]]
```

Check that the operator $e^{-i\delta q\,P}$ delivers a position kick and that the operator $e^{-i\delta p\,X}$ delivers a momentum kick. These are examples of evolution where the commuting series terminates.

```
 \begin{array}{l} {\rm vtest} \, [\text{``06f'} > {\rm position'} \, \, {\rm kick''}, \\ {\rm Simplify} \, [{\rm Evolver}[\delta q \, P, \, {\rm t}, \, Q] \, \, /. \, {\rm QP\$rules'} \sim {\rm Join} \sim \{{\rm t} \, \longrightarrow \, 1\}] \\ {\rm == } \, {\rm Q} \, - \, \delta q] \\ \\ {\rm vtest} \, [\text{``06g'} > {\rm momentum'} \, {\rm kick''}, \\ {\rm Simplify} \, [{\rm Evolver}[\delta p \, Q, \, {\rm t}, \, P] \, \, /. \, {\rm QP\$rules'} \sim {\rm Join} \sim \{{\rm t} \, \longrightarrow \, 1\}] \\ {\rm == } \, {\rm P} \, + \, \delta p] \\ \end{array}
```

Clean up:

Clear[$a, a^{\dagger}, \omega, Q, P, \mathcal{H}, Q, P, \delta q, \delta p$];

References

- [1] C. P. Slichter, Principles of magnetic resonance, 3rd edition, 1990.
- [2] P.-O. Löwdin, Quantum theory of many-particle systems. III. Extension of the Hartree-Fock scheme to include degenerate systems and correlation effects, *Phys. Rev.*, **1955**, *97*, 1509–1520.
- [3] J. W. Helton, M. de Oliveira, M. Stankus, and R. L. Miller, Ncalgebra version 5.0.6, 2015.