

Generalized Reed-Solomon Codes

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Overview

- Data transmission has become vital in today's society.
- Transmission methods are not perfect.
- Error correcting codes help to alleviate the burden of transmission.

Outline

- 1 Defining our Field
- 2 Codes
- 3 Generalized Reed-Solomon Codes
- 4 Conclusion

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Introducing \mathbb{F}_{64}

The finite field of 64 elements.

[Insert picture of ASCII table showing the elements]

Representation of our Elements

- Each element is encoded as a 6-bit binary string.
- We can represent that string as a polynomial in $\mathbb{F}_2[x]$.

For example:

- $\# \rightarrow 000011 \rightarrow x + 1$

Addition in \mathbb{F}_{64}

- Regular polynomial addition, but the coefficients are added modulo 2.
- $(x^4 + x^2 + x + 1) + (x^3 + x + 1) = x^4 + x^3 + x^2.$

Multiplication in \mathbb{F}_{64}

- Similar to regular multiplication, except the result cannot be larger than degree 5.
- We have to use an irreducible polynomial to mod out terms that are of too high a degree.
- The polynomial: $x^6 + x + 1 \rightarrow x^6 = x + 1$.

Multiplication Example

- $(x^5 + x^2 + x + 1)(x^4 + x)$
- $= x^9 + x^6 + x^6 + x^5 + x^4 + x^3 + x^2 + x$
- $= x^9 + x^5 + x^4 + x^3 + x^2 + x$
- $= (x^6)(x^3) + x^5 + x^4 + x^3 + x^2 + x$
- $= (x + 1)(x^3) + x^5 + x^4 + x^3 + x^2 + x$
- $= x^4 + x^3 + x^5 + x^4 + x^3 + x^2 + x$
- $= x^5 + x^2 + x$

Multiplication Table

[Insert Multiplication Table Here]

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 - Cyclic Redundancy Check
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Introduction to Codes

A *code* is a rule for converting information from one representation into another.

- $A \rightarrow 0$
- $B \rightarrow 1$
- \vdots
- $Z \rightarrow 25$

Encoding

Encoding is the act of conversion, following the rules of the code.

- HELLO \rightarrow 8 5 12 12 15

Errors

Random errors are a type of error that corrupts individual symbols during transmission.

- 8 5 **12** 12 15 \rightarrow 8 5 **19** 12 15

Burst errors are errors that corrupt a large chunk of symbols.

- 8 5 **12 12** 15 \rightarrow 8 5 **19 24** 5

Error Correction

Error detecting codes are a type of code that can detect when these errors occur.

Error correcting codes are a type of code that can correct these errors.

Introduction to CRC

- Used to detect accidental changes in data.
- Appends a check value to the message prior to transmission.
- After transmission the check value is recomputed and compared to the original value.

Application and Integrity

- CRC is good at detecting random errors and burst errors.
- It is not suitable for detecting intentional modifications to the data.

Example

To compute a binary CRC with a 3-bit check value:

- Start with our message encoded in binary: 11010011101100
- Find a irreducible polynomial of degree 3: $x^3 + x + 1 \rightarrow 1011$

Coming soon:

The remainder of the CRC example.

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 - Reed-Solomon Codes
 - Generalized Reed-Solomon Codes
 - Encoding GRS Codes
 - Decoding GRS Codes
- 4 Conclusion

History

- Introduced in 1960 by I.S. Reed and G. Solomon.
- Useful in practical applications and mathematically interesting.
- Used CD players and deep-space communications.

The Code

- Reed-Solomon (RS) codes are a linear, error correcting code.
- A RS-code is capable of correcting $\frac{n-k}{2}$ symbol errors. Where n is the total number of symbols, and k is the number of data symbols.
- Great at correcting burst errors.

Representation

A message encoded using an RS-code is a polynomial with the message symbols embedded in the coefficients.

For our example, the polynomial would be in $\mathbb{F}_{64}[x]$ and is our codeword.

- $H + Ex + Lx^2 + Lx^3 + Ox^4.$

Generalized Reed-Solomon Codes

- There is an alternate representation, known as Generalized Reed-Solomon (GRS) Codes.
- In this representation the message polynomial is evaluated at n distinct points and multiplied by a scalar.
- The result is a n -dimensional vector, which is used as the codeword.
- This is the representation that will be used for the remainder of the presentation.

The Code

Let F be a field. Choose nonzero elements $\hat{v} = v_1, v_2, \dots, v_n \in F$ and distinct elements $\hat{\alpha} = \alpha_1, \alpha_2, \dots, \alpha_n \in F$.

$$C = \text{GRS}_{n,k}(\hat{\alpha}, \hat{v}) = \{(v_1 f(\alpha_1), v_2 f(\alpha_2), \dots, v_n f(\alpha_n)) \mid f(x) \in F[x]_k\}$$

- $F = \mathbb{F}_{64}$.
- $\hat{v} = (1, 1, \dots, 1)$.
- $\hat{\alpha} = (A, B, \dots, W)$.
- $f(x)$ is our *message polynomial*.

Message Polynomial

- The message: "THIS IS MAJOR TOM".
- $f(x) = T + Hx + Ix^2 + Sx^3 + x^4 + Ix^5 + Sx^6 + x^7 + Mx^8 + Ax^9 + Jx^{10} + Ox^{11} + Rx^{12} + x^{13} + Tx^{14} + Ox^{15} + Mx^{16}$
- $\hat{c} = \hat{p} + \hat{e}$

The Codeword

The codeword, $\hat{c} = (v_1 f(\alpha_1), v_2 f(\alpha_2), \dots, v_n f(\alpha_n))$

- $\hat{c} = (!f(A), !f(B), \dots, !f(W))$
- $\hat{c} = \{T, (, U, 8, P, K, G, N, W, P, 4, K, ', -, N, (, M, H, K, 1, ", <, K\}$

The Received Message

- $\hat{c} = \{T, (, U, 8, P, K, \mathbf{G}, N, W, P, \mathbf{4}, K, ', -, N, (, M, H, K, 1, ", <, K\}$
- $\hat{p} = \{T, (, U, 8, P, K, \mathbf{P}, N, W, P, \mathbf{\&}, K, ', -, N, (, M, H, K, 1, ", \mathbf{R}, K\}$

Finding the Dual Code

- Want to find \hat{u} such that $C^\perp = \text{GRS}_{n,k}(\hat{\alpha}, \hat{u})$
- $L(x) = (x - \alpha_1)(x - \alpha_2)\dots(x - \alpha_n)$
- $L_i(x) = \frac{L(x)}{(x - \alpha_i)}$

Finding the Dual Code

- $\hat{u} = (L_1(\alpha_1), L_2(\alpha_2), \dots, L_n(\alpha_n),)$
- $\hat{u} = (L_1(A), L_2(B), \dots, L_{23}(W))$
- $\hat{u} = \{2, D, V, +, 9, O,], G, *, \wedge, 3, 5, X, ,, A, A, >, <, C, 8, G, E, : \}$

Finding the Syndrome Polynomial



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Final Thoughts

Questions