

10/2/2020

## UNIT-1

### MATHEMATICAL LOGIC

$$(p \vee q) \equiv (q \vee p)$$

$$\sim(p \vee q) \equiv \sim p \wedge \sim q$$

$$p \vee (q \vee r) \equiv (p \vee q) \vee r$$

$$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r).$$

$$\sim(p \rightarrow q) \equiv p \wedge \sim q \text{ or } \sim p \vee q$$

Suppose  $p \rightarrow q$  is given proposition

i.  $q \rightarrow p$  → converse

ii.  $\sim p \rightarrow \sim q$  → inverse

iii.  $\sim q \rightarrow \sim p$  → contrapositive.

Conjunctive Simplification:

$$(p \wedge q) \Rightarrow p$$

Disjunctive Amplification

$$p \Rightarrow (p \vee q)$$

$$\begin{array}{c} p, q, r \\ \frac{p \rightarrow q \quad q \rightarrow r}{p \rightarrow r} \end{array}$$

$$p \wedge (p \rightarrow q) \Rightarrow q$$

$$(p \rightarrow q) \wedge \sim q \Rightarrow \sim p$$

1. Let  $p$  and  $q$  be primitive stmts for which  $p \rightarrow q$  is false. Determine the truth values of the foll. compound proposition.

- i.  $p \wedge q$
- ii.  $\sim p \vee q$
- iii.  $\sim q \rightarrow \sim p$

$$\begin{array}{ccccc} p & q & p \rightarrow q & p \wedge q & \sim p \vee q & \sim q \rightarrow \sim p \\ \hline \text{F} & \text{T} & \text{F} & \text{F} & \text{T} & \text{T} \\ \text{T} & \text{F} & \text{T} & \text{F} & \text{F} & \text{F} \\ \text{F} & \text{T} & \text{T} & \text{F} & \text{F} & \text{T} \\ \text{F} & \text{F} & \text{T} & \text{F} & \text{F} & \text{F} \end{array}$$

2. PT:  $[(p \rightarrow q) \wedge (q \rightarrow r)] \rightarrow (p \rightarrow r)$  is a tautology

P	q	r	$p \rightarrow q$	$q \rightarrow r$	$(p \rightarrow q) \wedge (q \rightarrow r)$	$p \rightarrow r$	Given P.
T	T	T	T	T	T	T	T
T	T	F	T	F	F	F	T
T	F	T	F	T	F	T	T
T	F	F	F	F	F	F	T
F	T	T	T	T	T	T	T
F	T	F	T	F	F	T	T
F	F	T	T	T	T	T	T
F	F	F	T	T	T	T	T

3.  $[P \rightarrow (q \wedge r)] \Leftrightarrow [(P \rightarrow q) \wedge (P \rightarrow r)]$  using truth table  
 sol<sup>n</sup>:

P	q	r	$q \wedge r$	$P \rightarrow q \wedge r$	$P \rightarrow q$	$P \rightarrow r$	$(P \rightarrow q) \wedge (P \rightarrow r)$
T	T	T	T	T	T	T	T
T	T	F	F	F	F	F	F
T	F	T	F	F	F	T	F
T	F	F	F	F	F	F	F
F	T	T	T	T	T	T	T
F	T	F	F	T	T	T	T
F	F	T	F	T	T	T	T
F	F	F	F	T	T	T	T

Hence verified

11/2/2020

1. " If  $x$  is not a real no then it is not a rational no. and not an irrational no.

$$\sim [\sim p \rightarrow (\sim q \wedge \sim r)] \Rightarrow \sim p \wedge \sim (\sim q \wedge \sim r)$$

Sol<sup>n</sup>: P :  $x$  is a real no.

q :  $x$  is a rational no.

r :  $x$  is a irrational no.

$$\Rightarrow \sim p \wedge (q \vee r)$$

i.e.  $x$  is not a real no and it is a rational no. or it is an irrational no.

2. Consider  $(p \vee q) \wedge \sim(\sim p \wedge q) \Leftrightarrow p$

Sol<sup>n</sup>: Consider  $(p \vee q) \wedge \sim(p \wedge q)$

$$(p \vee q) \wedge (p \vee \sim q) \quad \text{De Morgan's.}$$

$$p \vee (q \wedge \sim q)$$

$$p \vee F_0$$

∴ Implication (Modus Ponens)

3. PT  $[(p \rightarrow q) \wedge \sim q] \Rightarrow \sim p$

Sol<sup>n</sup>: By Modus Tollens i.e.

if  $p \rightarrow q$  is true and  $q$  is false then  
 $p$  is false

$$p \rightarrow q$$

$$\frac{\sim q}{\sim p}$$

$$\therefore \sim p$$

4. If I study then I don't fail an exam  
 If I don't fail in exam, my father gifts me 2 wheels

Sol<sup>n</sup>:

Rule of syllogism

$$\begin{array}{c} p \rightarrow q \\ q \rightarrow r \\ \hline p \rightarrow r \end{array}$$

5. If you invest in the stock market then you will get rich

If you get rich then you will be happy

6. Prove the following argument is valid

$$\begin{array}{c} p \rightarrow r \\ q \rightarrow r \\ \hline (p \vee q) \rightarrow r \end{array}$$

Sol<sup>n</sup>:  $(p \rightarrow r) \wedge (q \rightarrow r) \Rightarrow (p \vee q) \rightarrow r$

Consider  $(p \rightarrow r) \wedge (q \rightarrow r)$

$$p \wedge (\neg p \vee q) \vee \\ (p \wedge p \rightarrow q) \Rightarrow q$$

$$q \wedge (\neg p \wedge q)$$

$$(q \wedge \neg p) \wedge (q \wedge \neg q)$$

$$\wedge F$$

$$(\neg p \wedge p) \vee (p \wedge q)$$

$$F_0 \vee (p \wedge q)$$

$$p \wedge (\neg p \vee q) \vee (q \wedge \neg(p \wedge q)) \Leftrightarrow q$$

$$\vee F_0$$

$$p \wedge \neg(p \rightarrow q)$$

19-2-2020:

Prove that the following arguments are valid

1. Smoking is healthy

If smoking is healthy then cigarettes are prescribed by doctors

∴ Cigarettes are prescribed by doctors

Soln:  $p \rightarrow$  Smoking is healthy

$q \rightarrow$  Cigarettes are prescribed by doctors.

$p$

$\frac{p \rightarrow q}{}$

$\therefore q$

By Modus Ponens this argument is valid

5. Continued.

$p$ : invest in stock market

$q$ : you will be rich

$r$ : you will be happy.

$p \rightarrow q$

Rule

$q \rightarrow r$

By laws of Syllogism.

$\therefore p \rightarrow r$

2. Prove that  $(p \vee q) \wedge \neg(\neg p \wedge q) \Leftrightarrow p$

Sol<sup>n</sup>:  $(p \vee q) \wedge \neg(\neg p \wedge q)$

$$\Leftrightarrow p \vee (q \wedge \neg q)$$

Distribution Law

$$\Leftrightarrow p \vee F_0$$

Identity Law

$$\Leftrightarrow p$$

Domination.

3. Prove that  $[p \vee (q \vee r) \wedge \neg q] \Rightarrow p \vee r$ .

Sol<sup>n</sup>: Consider  $p \vee (q \vee r) \wedge \neg q$

$$\Leftrightarrow \cancel{p \vee (q \vee r)} \wedge \neg q$$

$$\cancel{p \vee r}$$

$$\Rightarrow p \vee \underline{(q \vee r) \wedge \neg q}$$

$$\Rightarrow p \vee \cancel{r}$$

Rule of disjunction  
syllogism.

4. Simplify:  $(p \vee q) \wedge (\neg p \wedge q)$

Sol<sup>n</sup>: Consider  $(p \vee q) \wedge (\neg p \wedge q)$

$$p \wedge (\neg p \wedge q)$$

Disjunctive

$$(p \wedge \neg p) \wedge q$$

Amplification

Associative Law.

$$= F \wedge q = F_0.$$

$$5. p \vee \sim(p \wedge q)$$

Sol<sup>n</sup>: Consider

$$p \vee \sim(p \wedge q)$$

$$\Rightarrow p \vee \sim p$$

$$\Rightarrow T_0$$

$$6. \text{ PT } (p \rightarrow q) \wedge [\sim q \wedge (\tau \vee \sim q)] \Leftrightarrow \sim(q \vee p)$$

Sol<sup>n</sup>.

Consider

$$(p \rightarrow q) \wedge [\underline{\sim q \wedge (\tau \vee \sim q)}]$$

$$(p \rightarrow q) \wedge \sim q$$

Absorption

7. PT the following stmts are valid

$$\begin{array}{c} i. \sim p \rightarrow q \\ \sim q \\ \hline p \end{array}$$

By Modus Tollen's

$$\begin{array}{c} p \rightarrow q \\ \sim q \\ \hline \sim p. \end{array}$$

According to Modus Tollen's the given Stmt is valid.

8.

$$\frac{\begin{array}{c} p \rightarrow r \\ q \rightarrow r \end{array}}{(p \vee q) \rightarrow r}.$$

consider:  $(p \rightarrow r) \wedge (q \rightarrow r)$

$$\Leftrightarrow (\neg p \vee r) \wedge (\neg q \vee r)$$

$$\Leftrightarrow (r \vee \neg p) \wedge (r \vee \neg q)$$

$$\Leftrightarrow r \vee (\neg p \wedge \neg q)$$

$$\Leftrightarrow \neg(p \vee q) \vee r.$$

$$\Leftrightarrow (p \vee q) \rightarrow r.$$

9.  $p \rightarrow (q \rightarrow r)$

$$\frac{\begin{array}{c} \neg q \rightarrow \neg p \\ p \end{array}}{r}$$

Sol<sup>n</sup>:  $p \rightarrow (q \rightarrow r) \wedge (\neg q \rightarrow \neg p) \wedge p.$

$$p \rightarrow (q \rightarrow r) \wedge p' \wedge (\neg q \rightarrow \neg p)$$

$$p \wedge (p \rightarrow (q \rightarrow r)) \wedge (\neg q \rightarrow \neg p)$$

$$(q \rightarrow r) \wedge (p \rightarrow q)$$

$$(p \rightarrow q) \wedge (q \rightarrow r) \Leftrightarrow p \rightarrow r \Leftrightarrow np \vee r \Rightarrow r \vee np \Rightarrow r.$$

3 3-2020

2. If all triangles are right angled, then no triangle is equiangular.

Soln:  $p(x)$ :  $x$  is right angled

$q(x)$ :  $x$  is equiangular

$$\{\forall x \in T, p(x) \rightarrow \forall x \in T, \neg q(x)\}$$

negation of conditional is

$$\sim \{\forall x \in T, p(x) \rightarrow \forall x \in T, \neg q(x)\}$$

$$\sim(p \rightarrow q) \Rightarrow \neg p \rightarrow \neg q \quad p \wedge \neg q$$

$$\{\forall x \in T, p(x)\} \wedge \sim \{\forall x \in T, \neg q(x)\}$$

$$\sim \text{of } A \rightarrow F$$

$$\therefore \exists x \in T, p(x) \wedge \exists x \in T, q(x)$$

"All triangles are right angled and some are equiangular."

Prove the following

$$1. \forall x, p(x) \rightarrow \exists x, p(x)$$

$$2. \forall x [p(x) \vee q(x)] \rightarrow \forall x p(x) \vee \exists x q(x)$$

1. Consider

$$\forall x, p(x) \Rightarrow p(x) \text{ is true for every } x \in S$$

$$\Rightarrow p(a) \text{ is true for some } a \in S$$

$$\Rightarrow p(a) \text{ is true for some}$$

$$x = a \in S$$

$$\Rightarrow p(x) \text{ is true for some } x$$

$$\Rightarrow \exists x, p(x)$$

2. Consider

$$\forall x [p(x) \vee q(x)]$$

$$\Rightarrow p(x) \text{ or } q(x) \text{ is true for every } x \in S$$

$$\Rightarrow \{p(x) \text{ is true for every } x \in S\} \vee$$

$$\{q(x) \text{ is true for every } x \in S\}$$

$$\Rightarrow \{p(x) \text{ is true for every } x \in S\} \vee$$

$$\{q(a) \text{ is true for some } a \in S\}$$

$$\Rightarrow \{ p(x) \text{ is true for every } x \in S \} \vee \\ \{ q(x) \text{ is true for some } x \in S \} \\ \Rightarrow \text{pclose} \{ \forall x, p(x) \} \vee \{ \exists x, q(x) \}$$

Find whether the following argument is valid.

1. No engineering student of I <sup>or</sup> II Sem studies logic.

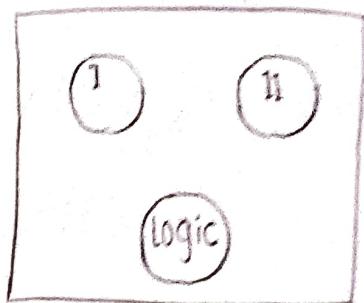
Anil is an engg. student who studies logic

$\therefore$  Anil is not in II Semester.

Soln:  $p(x)$  :  $x$  is in I Sem

$q(x)$  :  $x$  is in II Sem

$r(x)$  :  $x$  studies logic.



$$\forall x, \{ p(x) \vee q(x) \} \rightarrow \neg r(x)$$

$$r(a)$$

$$\therefore \neg q(a)$$

Consider

$$\forall x, [\{ p(x) \vee q(x) \} \rightarrow \neg r(x)] \wedge r(a) \Rightarrow \neg q(a)$$

Consider

$$\forall x \{ [p(x) \vee q(x)] \rightarrow \sim r(x) \} \wedge r(a)$$

$$\Rightarrow \{ [p(a) \vee q(a)] \rightarrow \sim r(a) \} \wedge r(a)$$

$$\Rightarrow \sim [\sim r(a)] \rightarrow \sim [p(a) \vee q(a)] \wedge r(a) \quad \text{Contraposition}$$

$$\Rightarrow [r(a) \rightarrow \sim \{ p(a) \vee q(a) \}] \wedge \sim r(a)$$

$$\Rightarrow r(a) \wedge \{ r(a) \rightarrow \sim [p(a) \vee q(a)] \} \quad \text{commutative}$$

per  $p \wedge (p \rightarrow q) \Rightarrow q$  Modus Ponens.

$$\Rightarrow \sim (p(a) \vee q(a))$$

$$\Rightarrow \sim p(a) \wedge \sim q(a)$$

$$\Rightarrow \sim q(a) \wedge \sim p(a)$$

$$\Rightarrow \sim q(a) \quad \text{Conjunctive simplification.}$$

Prove that the following argument is valid.

$$\forall x [p(x) \rightarrow q(x)]$$

$$\forall x [q(x) \rightarrow r(x)]$$

---

$$\forall x [p(x) \rightarrow r(x)]$$

sopn: According to Rule of Syllogism  
above argument is valid.

Prove the following argument is valid

$$\forall x [p(x) \rightarrow \{q(x) \wedge r(x)\}]$$

$$\forall x [p(x) \wedge s(x)]$$

$$\therefore \forall x [r(x) \wedge s(x)]$$

sopn:

$$\Rightarrow \{p(x) \wedge [p(x) \rightarrow \{q(x) \wedge r(x)\}]\} \wedge s(x)$$
$$\{s(x) \wedge [\quad] \}$$

$$\Rightarrow \{q(x) \wedge r(x)\} \wedge s(x) \quad \text{— Modus Ponens}$$

$$q(x) \wedge [r(x) \wedge s(x)] \quad \text{— Associative}$$

$$\Rightarrow [r(x) \wedge s(x)] \wedge q(x) \quad \text{— Commutative}$$

$$\Rightarrow r(x) \wedge s(x) \quad \text{— Conjunctive Simplification}$$

## Method of Proof :

For any given condition of the form  $p \rightarrow q$  the process of establishing that the condition is true by using the rules or laws of logic constitutes a proof a condition

### \* Direct proof :

The direct method of proving the condition  $p \rightarrow q$  is true is as below.

1. Hypothesis : (Assumption)  $p$  is true
2. Analysis :: Starting with hypothesis and applying rules of logic infer that  $q$  is true
3. Conclusion :  $p \rightarrow q$  is true.

This is known as method of direct proof.

### Indirect proof :

For the given condition  $p \rightarrow q$  show that the contrapositive  $\sim q \rightarrow \sim p$  is true

## \* Proof by contradiction.

The steps are as below.

1. Hypothesis : Assume that  $p \rightarrow q$  is false  
i.e. Assume  $p$  is true and  $q$  is false
2. Analysis : Starting with the hypothesis that  $q$  is false and by applying rules of logic and other known facts prove that  $p$  is also false.

This contradicts the assumption that  $p$  is true.

3. Conclusion : Because of the contradiction of the analysis we infer that  $p \rightarrow q$  is true.

4-3-2020

## Problems on Direct proof

1. The square of an odd integer is an odd integer

Sol<sup>n</sup>: If  $n$  is an odd integer then  $n^2$  is an odd integer

p : n is odd integer

q :  $n^2$  is odd integer

Direct proof:

Assume p is true

i.e. n is an odd integer

$\Rightarrow n = 2k+1$ , k is an integer

$$n^2 = (2k+1)^2$$

$$= 4k^2 + 4k + 1$$

$$= 4(k^2 + k) + 1$$

$\therefore n^2$  is odd integer

$\therefore p \rightarrow q$  is true

2. Prove that for all integers K and l , if K and l are both odd then  $K+l$  is even and  $KL$  is odd

SOL:

p : K and l are odd

q :  $K+l$  is even and  $KL$  is odd

r :  $KL$  is odd.

To prove :  $p \rightarrow (q \wedge r)$

Assume p is true

i.e. k and l are odd

$$k = 2m+1, m \text{ is integer}$$

$$l = 2n+1, n \text{ is integer}$$

$$\text{Now, } k+l = (2m+1) + (2n+1)$$

$$= 2m+2n+2$$

$$= 2(m+n+1)$$

$\therefore k+l$  is even

Hence q is true.

$$\text{Now, } kl = (2m+1)(2n+1)$$

$$= 4mn + 2m + 2n + 1$$

$$= 2(mn+m+n) + 1$$

$\therefore kl$  is odd

Hence r is true.

Hence  $p \rightarrow (q \wedge r)$  is true.

## Problems on Method of Contradiction:

1. If  $n$  is an odd integer then  $n+11$  is an even integer.

Soln:  $p$ :  $n$  is an odd integer

$q$ :  $n+11$  is an even integer

Contradiction:

Assume  $p \rightarrow q$  is false.

i.e.  $p$  is true and  $q$  is false

$q$  is false i.e.  $n+11$  is an odd integer

$n+11 = 2k+1$ ,  $k$  is integer

$$n = 2k-10$$

$$n = 2(k-5)$$

$\therefore n$  is an even integer

$p$  is false which contradicts the assumption that  $p$  is true

$\therefore p \rightarrow q$  is true.

2. For every integer  $n$ , if  $n^2$  is odd then  $n$  is odd

Sol<sup>n</sup>:  $P: n^2$  is odd

$q: n$  is odd

contradiction

Assume  $P \rightarrow q$  is false

i.e.  $P$  is true and  $q$  is false

$q$  is false i.e.  $\neg q$   $n$  is even.

$$n = 2k \quad k \text{ is integer}$$

$$\text{Now } n^2 = (2k)^2$$

$$n^2 = 4k^2$$

$\therefore n^2$  is even

$P$  is false which contradicts the assumption that  $P$  is true

$\therefore P \rightarrow q$  is true.

## Problems on Indirect Proof

1. The product of 2 even integers is an even integer.

Soln: If a and b are even integers then ab is an even integer.

P: a and b are even integers.

q: ab is even integer

Indirect proof:  $\neg p \vee \neg q \rightarrow \neg p$

Assume  $\neg q$  is true

i.e. ab is an odd integer

ab is not divisible by 2

then a is not divisible by 2, also b is not divisible by 2

$\Rightarrow$  a is odd integer and b is odd integer

i.e. p is false

$\neg p$  is true

Hence  $\neg q \rightarrow \neg p$  is true.

$p \rightarrow q$  is true.

6-3-2020

2 PT for all real nos.  $x$  and  $y$ , if  $x+y \geq 100$   
then  $x \geq 50$  and or  $y \geq 50$ .

Sol<sup>n</sup>:  $P : x+y \geq 100$

$$q : x \geq 50 \quad r : y \geq 50$$

Indirect Proof  $P \rightarrow (q \vee r)$

$$\sim(q \vee r) \rightarrow \sim P.$$

$$\text{i.e. } (\sim q \wedge \sim r) \rightarrow \sim P.$$

Assume  $(\sim q \wedge \sim r)$  is true.

i.e.  $x \leq 50$ , and  $y \leq 50$

$$x+y \leq 100.$$

i.e.  $P$  is false

$\sim P$  is true.

Hence  $\sim(q \vee r) \rightarrow \sim P$

$\therefore P \rightarrow (q \vee r)$  is true.

## Stirling Number:

If  $A$  and  $B$  are finite set with cardinality of  $|A|=m$  and  $|B|=n$  where  $m \geq n$  then the number of onto functions from  $A$  to  $B$  is given by

$$P(m,n) = \sum_{k=0}^n (-1)^k \cdot \left( {}^n C_{n-k} \right) (n-k)^m$$

with and  $\frac{P(m,n)}{n!}$  is called the stirling number

of second kind and is denoted as  $S(m,n)$

$$S(m,n) = \frac{1}{n!} \sum_{k=0}^n (-1)^k \left( {}^n C_{n-k} \right) (n-k)^m$$

- The Stirling number represents the no. of ways in which it is possible to assign  $m$  distinct objects into  $n$  identical places with no place left empty.

\* Note :  $S(m,1) = 1$

$$S(m,m) = 1 \quad \text{for } m \geq 1.$$

The no. of possible ways to assign  $m$  distinct objects to  $n$  identical places with empty places allowed is given by

$$P(m) = \sum_{i=1}^n S(m, i), \quad m \geq 1$$

1. Evaluate  $S(5, 4)$ ,  $S(8, 6)$ .

Sol<sup>n</sup>: i.  $m = 5$ ,  $n = 4$

$$\begin{aligned} S(5, 4) &= \frac{1}{n!} \sum_{k=0}^n (-1)^k \left( {}^n C_{n-k} \right) (n-k)^m \\ &= \frac{1}{4!} \left[ (-1)^0 {}^4 C_4 * (4-0)^5 + \right. \\ &\quad (-1)^1 {}^4 C_3 * (4-1)^5 + \\ &\quad (-1)^2 {}^4 C_2 * (4-2)^5 + \\ &\quad (-1)^3 {}^4 C_1 * (4-3)^5 + \\ &\quad \left. (-1)^4 {}^4 C_0 * (4-4)^5 \right] \end{aligned}$$

$$\begin{aligned} &= \frac{1}{24} \left[ -1 \times 1 \times 1024 + (-1) \times 4 \times 243 \right. \\ &\quad \left. + 1 \times 6 \times 32 + (-1) \times 4 \times 1 + 0 \right] \end{aligned}$$

$$= \frac{1}{24} \left\{ -1024 - 972 + 32 + 4 \right\} = \frac{-1020}{24} = -\frac{255}{6}$$

$$2. A = \{1, 2, 3, 4, 5, 6\}$$

$$B = \{w, x, y, z\}$$

$$\text{soln. } |A| = 6 \quad |B| = 4$$

$$P(6, 4) = ?$$

$$P(m, n) = \sum_{k=0}^n (-1)^k {}^n C_{n-k} (n-k)^m$$

$$= \sum_{k=0}^4 (-1)^k$$

$$= [(-1)^0 {}^4 C_4 (4-0)^6 + \\ (-1)^1 {}^4 C_3 (4-1)^6 + \\ (-1)^2 {}^4 C_2 (4-2)^6 + \\ (-1)^3 {}^4 C_1 (4-3)^6 + \\ (-1)^4 {}^4 C_0 (4-4)^6]$$

$$= [-1 \times 1 \times 4096 - 4 \times 729 + 6 \times 64 \\ - 4 \times 1 + 0]$$

$$= -6632$$

3. There are 6 programmers who can assist 8 executives in how many ways can the executives be assisted so that each programmer assist at least one executive.

$$\text{Soln: } P(m, n) = n! \cdot S(m, n)$$

$$|m| = 8 \quad |n| = 6$$

$$P(m, n) = \frac{1}{n!} \sum_{k=0}^n (-1)^k \binom{n}{k} (n-k)^m$$

$$= 6! \left[ (-1)^0 {}^6C_0 (6-0)^8 + \right.$$

$$(-1)^1 {}^6C_1 (6-1)^8 +$$

$$(-1)^2 {}^6C_2 (6-2)^8 +$$

$$(-1)^3 {}^6C_3 (6-3)^8 +$$

$$(-1)^4 {}^6C_4 (6-4)^8 +$$

$$(-1)^5 {}^6C_5 (6-5)^8 +$$

$$(-1)^6 {}^6C_6 (6-6)^8 \Big]$$

$$= 6! \left[ -1 \times \right.$$

$$= 191520$$

3. Let  $A = \{1, 2, 3, 4\}$   $B = \{1, 2, 3, 4, 5, 6\}$

a. find how many functions are there from A to B

b. How ~~may~~ many of these are one to one

c. How many are onto.

i. find how many functions are there from B to A.

ii. how many of these are onto

Soln:  $|A| = m = 4$ .

$$|B| = n = 6$$

a. The no. of possible functions from A to B

$$n^m = 6^4 = 1296.$$

b. There will be  $\frac{n!}{(n-m)!}$  no. of one to one.

functions from A to B.

$$\frac{6!}{(6-4)!} = 360$$

11 - 3 - 2020.

Prove that if 30 dictionaries in a library contain a total of 61327 pages, then atleast one of the dictionaries must have atleast 2045 pages.

Sol:

### Pigeon Hole

If  $m$  pigeons occupy  $n$  pigeon holes, then atleast 1 pigeon hole must contain  $p+1$  or more pigeons, where  $p = \left\lfloor \frac{m-1}{n} \right\rfloor$ .

$$p = \left\lfloor \frac{61327-1}{30} \right\rfloor = [2044.2] \\ = 2044.$$

Treating pages as pigeons and dictionaries as pigeon holes by pigeon hole principle atleast one of the dictionary <sup>must contain</sup>  $p+1$  or more pages where

$$P = 2044.$$

ST if any 5 nos from 1 to 8 are chosen  
then 2 of them will have their sum = 9

Soln: Let us consider the following sets

$$A_1 = \{1, 8\} \quad A_2 = \{2, 7\} \quad A_3 = \{3, 6\}$$

$$A_4 = \{4, 5\}$$

These are the only sets containing the nos from 1 to 8 whose sum is 9

Since every no. from 1 to 8 belongs to one of these sets each of the 5 no. chosen must belong to one of the sets. also there are only 4 sets, 2 of the 5 chosen nos have to belong to the same set.

And these 2 nos have their sum = 9

## Principle of Inclusion and Exclusion:

If  $S$  is a finite set  $A$  and  $B$  are subsets of  $S$  then  $|A \cup B| = |A| + |B| - |A \cap B|$

$$A \cap \bar{B} = |A \cup B| = S - |A \cup B|$$

$$= S - \{ |A| + |B| - |A \cap B| \}$$

$$= S - |A| - |B| + |A \cap B| \quad \text{--- 2}$$

One 1 is known as addition principle and 2 is known as principle and exclusion principle

$$\text{For } n \text{ sets } |\bar{A}_1 \cap \bar{A}_2 \cap \dots \cap \bar{A}_n| = S - \sum |A_i| + \sum |A_i \cap A_j| - \sum |A_i \cap A_j \cap A_k| + (-1)^n (A_1 \cap A_2 \cap A_3 \cap \dots \cap A_n)$$

1 A computer company requires 30 programmers to handle system programming jobs and 40 programmers for application programming jobs.

If the company appoints 55 programmers to carry out these jobs [how many of them performs job of both type]

Soln:

- ii. how many handle only system programming job
- ii. how many handle only application programming job.

$$\text{Soln: } |S| = 55 : |A \cup B|$$

$$|A| = 30 : |B| = 40$$

$$\begin{aligned} \text{i. } |A \cap B| &= |A \cup B| - |A| - |B| \\ &= 55 - 30 - 40 \\ &= 15 \end{aligned}$$

$$\text{ii. } |A| - |A \cap B| = 30 - 15 = 15$$

$$\text{iii. } |B| - |A \cap B| = 40 - 15 = 25$$

2. 30 cars are assembled in a factory, the options available are a music system, AC and power windows. It is known that 15 of the cars have music system, 8 AC, and 6 have power window further 3 have all the options

do not

determine how many cars have any options at all

12-3-2020

$$\text{soln: } |S| = 30$$

$$|A_1| = 15$$

$$|B_2| = 8$$

$$|A_3| = 6.$$

$$|A_1 \cap B_2 \cap A_3| = 3.$$

To find:  $|\bar{A}_1 \cap \bar{B}_2 \cap \bar{A}_3|$

$$|\bar{A} \cap \bar{B} \cap \bar{C}| = |S| - [\sum |A_i|] + [\sum |A_i \cap A_j|] - |A_1 \cap B_2 \cap A_3|$$

$$= 30 - 29 + \sum |A_i \cap A_j| - 3$$

$$= \sum |A_i \cap A_j| - 2$$

We know

$|A_i \cap A_1 \cap A_2 \cap A_3|$  is a subset of  $A_i \cap A_j$  for

$$i, j = 1, 2, 3$$

i.e. Each of  $|A_i \cap A_j|$ , which are 3 in number is greater than or equal to  $|A_1 \cap A_2 \cap A_3|$

$$\sum |A_i \cap A_j| \geq 3|A_1 \cap A_2 \cap A_3|$$

$$\geq 3 \times 3$$

$$\sum |A_i \cap A_j| \geq 9 - 2.$$

using 2 in 1

$$|\bar{A}_1 \cap \bar{A}_2 \cap \bar{A}_3| \geq 9 - 2$$

$$|\bar{A}_1 \cap \bar{A}_2 \cap \bar{A}_3| \geq 7$$

Hence there will be greater than or equal to 7 cars which do not have any options.

3. In how many ways can the 26 letters of the English alphabet be permuted so that none of the patterns CAR, DOG, PUN, and BYTE occurs

SOLN: Let  $S$  denote the set of all the permutations of 26 letters i.e.

$$|S| = 26!$$

Let  $A_1$  be the set of all the permutations in which CAR appears.,  $A_2$  is the set of all permutations in which DOG is in  $A_3$  and  $A_4$  has all the permutations in which PUN and BYTE appears

The set  $A_1$  has the cardinality  $|A_1| = 24!$

The word CAR consists of 3 letters which form a single block.

set  $A_1$  consists of all permutations which contain this single block and 23 remaining letters.

$$\text{Hence } |A_2| = 24! \quad \text{as}$$

$$|A_3| = 24!$$

$$|A_4| = 23!$$

$$|A_1 \cap A_2| = (26 - 6 + 2)!$$

$$= 22!$$

$$|A_1 \cap A_3| = |A_2 \cap A_3| = 22!$$

Also we can find that  $|A_1 \cap A_2| = |A_1 \cap A_3| = |A_2 \cap A_3| = 22!$

$$|A_1 \cap A_4| = (26 - 7 + 2)! = 21!$$

$$\text{Hence } |A_2 \cap A_4| = |A_3 \cap A_4| = 21!$$

$$|A_1 \cap A_2 \cap A_3| = (26 - 9 + 3)! = 20!$$

$$|A_2 \cap A_3 \cap A_4| = (26 - 10 + 3)! = 19!$$

$$|A_1 \cap A_3 \cap A_4| = 19!$$

$$|A_1 \cap A_2 \cap A_4| = 19!$$

$$|A_1 \cap A_2 \cap A_3 \cap A_4| = (26 - 13 + 4)! = 17!$$

We need to find

$$\begin{aligned} |\bar{A}_1 \cap \bar{A}_2 \cap \bar{A}_3 \cap \bar{A}_4| &= |S| - \sum |A_i| + \sum |A_i \cap A_j| \\ &\quad - \sum |A_i \cap A_j \cap A_k| \\ &\quad + |A_1 \cap A_2 \cap A_3 \cap A_4| \\ &= \cancel{26!} - \cancel{7!} + \cancel{108!} * \cancel{- 77!} + \cancel{17}. \end{aligned}$$

4. Find the no. of non-negative integer solution of the equation  $x_1 + x_2 + x_3 + x_4 = 18$  under the condition  $x_i \leq 7$  for  $i = 1, 2, 3, 4$

Note :

$$C(n+r-1, r) = \frac{(n+r-1)!}{r!(n-1)!}$$

represents the no. of non negative integer solu<sup>n</sup>  
of the equn  $x_1 + x_2 + x_3 + \dots + x_n = 7$ .

Soln. Let  $S$  denote the set of all non-negative  
integer solution of the given equn.

The number of such solution is

$$C(4+18-1, 18) = C(21, 18)$$

$$\therefore |S| = C(21, 18)$$

Let  $A_1$  be the subset of  $S$  containing  
non-negative integer soln of the given equn  
with the condition  $x_1 > 7$   $x_2, x_3, x_4 \geq 0$

$$A_1 = \{(x_1, x_2, x_3, x_4) \in S \mid x_1 > 7\}$$

$$A_2 = \{(x_1, x_2, x_3, x_4) \in S \mid x_2 > 7\}$$

$$A_3 = \{(x_1, x_2, x_3, x_4) \in S \mid x_3 > 7\}$$

$$A_4 = \{ \quad \quad \quad \in S \mid x_4 > 7 \}$$

We now require to find cardinality of

$$|\bar{A}_1 \cap \bar{A}_2 \cap \bar{A}_3 \cap \bar{A}_4|$$

Let us set  $y_1 = x_1 - 8$  but  $x_1 \geq 7$   
 $\Rightarrow x_1 \geq 8$   
 $\therefore y \geq 0$

i.e. We shall have

$$x_1 + x_2 + x_3 + x_4 = 18$$

$$y_1 + x_2 + x_3 + x_4 = 10$$

The no. of such non-negative integer solutions  
 will be

$$C(4+10-1, 10) = C(13, 10)$$

$$|A_1| = C(13, 10)$$

111<sup>4</sup> because of symmetry

$$|A_2| = |A_3| = |A_4| = C(13, 10)$$

Now let us take

$$y_1 = x_1 - 8 \quad y_2 = x_2 - 8$$

$$y_1 \geq 0 \quad y_2 \geq 0$$

$$y_1 + y_2 + x_3 + x_4 = 2$$

$\therefore$  The no. of such non-negative integer sol's  
 will be

$$C(4+2-1, 2) = C(5, 2)$$

$$|A_1 \cap A_2| = C(5, 2)$$

$$\text{By symmetry } |A_1 \cap A_3| = |A_1 \cap A_4| = |A_2 \cap A_3| = |A_2 \cap A_4| \\ = |A_3 \cap A_4| = C(5,2)$$

In the given equat<sup>n</sup> more than 2  $x_i$ 's cannot be greater than 7 simultaneously  
hence

$$|A_1 \cap A_2 \cap A_3| = |A_2 \cap A_3 \cap A_4| = |A_1 \cap A_3 \cap A_4| = \\ |A_1 \cap A_2 \cap A_4| = |A_1 \cap A_2 \cap A_3 \cap A_4| \\ = 0$$

$$|\bar{A}_1 \cap \bar{A}_2 \cap \bar{A}_3 \cap \bar{A}_4| = |S| - \sum |A_{ii}| + \sum |A_i \cap A_j| - \sum |A_i \cap A_j \cap A_k| \\ + |A_1 \cap A_2 \cap A_3 \cap A_4| \\ = C(22, 18) + 6(C(5, 2)) \\ - 4C(13, 10)$$