

# 7

## Generating Functions

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THE MOST POWERFUL WAY to deal with sequences of numbers, as far as anybody knows, is to manipulate infinite series that “generate” those sequences. We’ve learned a lot of sequences and we’ve seen a few generating functions; now we’re ready to explore generating functions in depth, and to see how remarkably useful they are.

### 7.1 DOMINO THEORY AND CHANGE

Generating functions are important enough, and for many of us new enough, to justify a relaxed approach as we begin to look at them more closely. So let’s start this chapter with some fun and games as we try to develop our intuitions about generating functions. We will study two applications of the ideas, one involving dominoes and the other involving coins.

How many ways  $T_n$  are there to completely cover a  $2 \times n$  rectangle with  $2 \times 1$  dominoes? We assume that the dominoes are identical (either because they’re face down, or because someone has rendered them indistinguishable, say by painting them all red); thus only their orientations — vertical or horizontal — matter, and we can imagine that we’re working with domino-shaped tiles. For example, there are three tilings of a  $2 \times 3$  rectangle, namely  $\begin{smallmatrix} \square & \square & \square \\ \square & \square & \square \end{smallmatrix}$ ,  $\begin{smallmatrix} \square & \square & \square \\ \square & \square & \square \end{smallmatrix}$ , and  $\begin{smallmatrix} \square & \square & \square \\ \square & \square & \square \end{smallmatrix}$ ; so  $T_3 = 3$ .

To find a closed form for general  $T_n$  we do our usual first thing, look at small cases. When  $n = 1$  there’s obviously just one tiling,  $\begin{smallmatrix} \square \\ \square \end{smallmatrix}$ ; and when  $n = 2$  there are two,  $\begin{smallmatrix} \square & \square \\ \square & \square \end{smallmatrix}$  and  $\begin{smallmatrix} \square & \square \\ \square & \square \end{smallmatrix}$ .

How about when  $n = 0$ ; how many tilings of a  $2 \times 0$  rectangle are there? It’s not immediately clear what this question means, but we’ve seen similar situations before: There is one permutation of zero objects (namely the empty permutation), so  $0! = 1$ . There is one way to choose zero things from  $n$  things (namely to choose nothing), so  $\binom{n}{0} = 1$ . There is one way to partition the empty set into zero nonempty subsets, but there are no such ways to partition a nonempty set; so  $\{n\} = [n=0]$ . By such reasoning we can conclude that

there's just one way to tile a  $2 \times 0$  rectangle with dominoes, namely to use no dominoes; therefore  $T_0 = 1$ . (This spoils the simple pattern  $T_n = n$  that holds when  $n = 1, 2$ , and  $3$ ; but that pattern was probably doomed anyway, since  $T_0$  wants to be  $1$  according to the logic of the situation.) A proper understanding of the null case turns out to be useful whenever we want to solve an enumeration problem.

Let's look at one more small case,  $n = 4$ . There are two possibilities for tiling the left edge of the rectangle — we put either a vertical domino or two horizontal dominoes there. If we choose a vertical one, the partial solution is  $\square$  and the remaining  $2 \times 3$  rectangle can be covered in  $T_3$  ways. If we choose two horizontals, the partial solution  $\square\square$  can be completed in  $T_2$  ways. Thus  $T_4 = T_3 + T_2 = 5$ . (The five tilings are  $\square\square\square$ ,  $\square\square\square$ ,  $\square\square\square$ ,  $\square\square\square$ , and  $\square\square\square$ .)

We now know the first five values of  $T_n$ :

| $n$   | 0 | 1 | 2 | 3 | 4 |
|-------|---|---|---|---|---|
| $T_n$ | 1 | 1 | 2 | 3 | 5 |

These look suspiciously like the Fibonacci numbers, and it's not hard to see why: The reasoning we used to establish  $T_4 = T_3 + T_2$  easily generalizes to  $T_n = T_{n-1} + T_{n-2}$ , for  $n \geq 2$ . Thus we have the same recurrence here as for the Fibonacci numbers, except that the initial values  $T_0 = 1$  and  $T_1 = 1$  are a little different. But these initial values are the consecutive Fibonacci numbers  $F_1$  and  $F_2$ , so the  $T$ 's are just Fibonacci numbers shifted up one place:

$$T_n = F_{n+1}, \quad \text{for } n \geq 0.$$

(We consider this to be a closed form for  $T_n$ , because the Fibonacci numbers are important enough to be considered "known." Also,  $F_n$  itself has a closed form (6.123) in terms of algebraic operations.) Notice that this equation confirms the wisdom of setting  $T_0 = 1$ .

But what does all this have to do with generating functions? Well, we're about to get to that — there's another way to figure out what  $T_n$  is. This new way is based on a bold idea. Let's consider the "sum" of all possible  $2 \times n$  tilings, for all  $n \geq 0$ , and call it  $T$ :

*To boldly go  
where no tiling has  
gone before.*

$$T = 1 + \square + \square\square + \square\square\square + \square\square\square\square + \square\square\square\square\square + \dots \quad (7.1)$$

(The first term '1' on the right stands for the null tiling of a  $2 \times 0$  rectangle.) This sum  $T$  represents lots of information. It's useful because it lets us prove things about  $T$  as a whole rather than forcing us to prove them (by induction) about its individual terms.

The terms of this sum stand for tilings, which are combinatorial objects. We won't be fussy about what's considered legal when infinitely many tilings

are added together; everything can be made rigorous, but our goal right now is to expand our consciousness beyond conventional algebraic formulas.

We've added the patterns together, and we can also multiply them—by juxtaposition. For example, we can multiply the tilings  $\square$  and  $\square$  to get the new tiling  $\square\square$ . But notice that multiplication is not commutative; that is, the order of multiplication counts:  $\square\square$  is different from  $\square\square$ .

Using this notion of multiplication it's not hard to see that the null tiling plays a special role—it is the multiplicative identity. For instance,  $1 \times \square = \square \times 1 = \square$ .

Now we can use domino arithmetic to manipulate the infinite sum  $T$ :

$$\begin{aligned} T &= 1 + \square + \square\square + \square\square\square + \square\square\square\square + \square\square\square\square\square + \dots \\ &= 1 + \square(1 + \square + \square\square + \dots) + \square(1 + \square + \square\square + \dots) \\ &= 1 + \square T + \square T. \end{aligned} \tag{7.2}$$

Every valid tiling occurs exactly once in each right side, so what we've done is reasonable even though we're ignoring the cautions in Chapter 2 about “absolute convergence.” The bottom line of this equation tells us that everything in  $T$  is either the null tiling, or is a vertical tile followed by something else in  $T$ , or is two horizontal tiles followed by something else in  $T$ .

So now let's try to solve the equation for  $T$ . Replacing the  $T$  on the left by  $1T$  and subtracting the last two terms on the right from both sides of the equation, we get

$$(1 - \square - \square)T = 1. \tag{7.3}$$

For a consistency check, here's an expanded version:

$$\begin{array}{r} 1 + \square + \square\square + \square\square\square + \square\square\square\square + \square\square\square\square\square + \dots \\ - \square - \square\square - \square\square\square - \square\square\square\square - \square\square\square\square\square - \dots \\ - \square - \square\square - \square\square\square - \square\square\square\square - \square\square\square\square\square - \dots \\ \hline 1 \end{array}$$

Every term in the top row, except the first, is cancelled by a term in either the second or third row, so our equation is correct.

So far it's been fairly easy to make combinatorial sense of the equations we've been working with. Now, however, to get a compact expression for  $T$  we cross a combinatorial divide. With a leap of algebraic faith we divide both sides of equation (7.3) by  $1 - \square - \square$  to get

$$T = \frac{1}{1 - \square - \square}. \tag{7.4}$$

ave a gut feel-  
that these  
ns must con-  
ge, as long as  
dominoes are  
all enough.

(Multiplication isn't commutative, so we're on the verge of cheating, by not distinguishing between left and right division. In our application it doesn't matter, because  $|$  commutes with everything. But let's not be picky, unless our wild ideas lead to paradoxes.)

The next step is to expand this fraction as a power series, using the rule

$$\frac{1}{1-z} = 1 + z + z^2 + z^3 + \dots$$

The null tiling  $|$ , which is the multiplicative identity for our combinatorial arithmetic, plays the part of 1, the usual multiplicative identity; and  $\square + \boxminus$  plays  $z$ . So we get the expansion

$$\begin{aligned} \frac{|}{1-\square-\boxminus} &= | + (\square + \boxminus) + (\square + \boxminus)^2 + (\square + \boxminus)^3 + \dots \\ &= | + (\square + \boxminus) + (\square + \square\boxminus + \boxminus\square + \boxminus\boxminus) \\ &\quad + (\square\square + \square\boxminus + \boxminus\square + \boxminus\boxminus + \square\boxminus\square + \boxminus\square\boxminus + \boxminus\boxminus\square + \boxminus\boxminus\boxminus) + \dots \end{aligned}$$

This is  $T$ , but the tilings are arranged in a different order than we had before. Every tiling appears exactly once in this sum; for example,  $\square\boxminus\square\boxminus$  appears in the expansion of  $(\square + \boxminus)^7$ .


We can get useful information from this infinite sum by compressing it down, ignoring details that are not of interest. For example, we can imagine that the patterns become unglued and that the individual dominoes commute with each other; then a term like  $\square\boxminus\square\boxminus$  becomes  $\square^4\boxminus^6$ , because it contains four verticals and six horizontals. Collecting like terms gives us the series

$$T = | + \square + \boxminus + \square^2 + \boxminus^2 + \square^3 + 2\square\boxminus^2 + \boxminus^4 + 3\square^2\boxminus^2 + \boxminus^6 + \dots$$

The  $2\square\boxminus^2$  here represents the two terms of the old expansion,  $\square\boxminus$  and  $\boxminus\square$ , that have one vertical and two horizontal dominoes; similarly  $3\square^2\boxminus^2$  represents the three terms  $\square\square$ ,  $\square\boxminus$ , and  $\boxminus\square$ . We're essentially treating  $\square$  and  $\boxminus$  as ordinary (commutative) variables.

We can find a closed form for the coefficients in the commutative version of  $T$  by using the binomial theorem:

$$\begin{aligned} \frac{|}{1-(\square + \boxminus^2)} &= | + (\square + \boxminus^2) + (\square + \boxminus^2)^2 + (\square + \boxminus^2)^3 + \dots \\ &= \sum_{k \geq 0} (\square + \boxminus^2)^k \\ &= \sum_{j, k \geq 0} \binom{k}{j} \square^j \boxminus^{2k-2j} \\ &= \sum_{j, m \geq 0} \binom{j+m}{j} \square^j \boxminus^{2m}. \end{aligned} \tag{7.5}$$

(The last step replaces  $k-j$  by  $m$ ; this is legal because we have  $\binom{k}{j} = 0$  when  $0 \leq k < j$ .) We conclude that  $\binom{j+m}{j}$  is the number of ways to tile a  $2 \times (j+2m)$  rectangle with  $j$  vertical dominoes and  $2m$  horizontal dominoes. For example, we recently looked at the  $2 \times 10$  tiling , which involves four verticals and six horizontals; there are  $\binom{4+3}{4} = 35$  such tilings in all, so one of the terms in the commutative version of  $T$  is  $35 \square^4 \square^6$ .

We can suppress even more detail by ignoring the orientation of the dominoes. Suppose we don't care about the horizontal/vertical breakdown; we only want to know about the total number of  $2 \times n$  tilings. (This, in fact, is the number  $T_n$  we started out trying to discover.) We can collect the necessary information by simply substituting a single quantity,  $z$ , for  $\square$  and  $\square$ . And we might as well also replace  $1$  by  $1$ , getting


$$T = \frac{1}{1 - z - z^2}. \quad (7.6)$$

This is the generating function (6.117) for Fibonacci numbers, except for a missing factor of  $z$  in the numerator; so we conclude that the coefficient of  $z^n$  in  $T$  is  $F_{n+1}$ .

The compact representations  $1/(1-\square-\square)$ ,  $1/(1-\square-\square^2)$ , and  $1/(1-z-z^2)$  that we have deduced for  $T$  are called *generating functions*, because they generate the coefficients of interest.

Incidentally, our derivation implies that the number of  $2 \times n$  domino tilings with exactly  $m$  pairs of horizontal dominoes is  $\binom{n-m}{m}$ . (This follows because there are  $j = n - 2m$  vertical dominoes, hence there are

$$\binom{j+m}{j} = \binom{j+m}{m} = \binom{n-m}{m}$$

ways to do the tiling according to our formula.) We observed in Chapter 6 that  $\binom{n-m}{m}$  is the number of Morse code sequences of length  $n$  that contain  $m$  dashes; in fact, it's easy to see that  $2 \times n$  domino tilings correspond directly to Morse code sequences. (The tiling  corresponds to '•- -•-•'.) Thus domino tilings are closely related to the continuant polynomials we studied in Chapter 6. It's a small world.

We have solved the  $T_n$  problem in two ways. The first way, guessing the answer and proving it by induction, was easier; the second way, using infinite sums of domino patterns and distilling out the coefficients of interest, was fancier. But did we use the second method only because it was amusing to play with dominoes as if they were algebraic variables? No; the real reason for introducing the second way was that the infinite-sum approach is a lot more powerful. The second method applies to many more problems, because it doesn't require us to make magic guesses.

Let's generalize up a notch, to a problem where guesswork will be beyond us. How many ways  $U_n$  are there to tile a  $3 \times n$  rectangle with dominoes?

The first few cases of this problem tell us a little: The null tiling gives  $U_0 = 1$ . There is no valid tiling when  $n = 1$ , since a  $2 \times 1$  domino doesn't fill a  $3 \times 1$  rectangle, and since there isn't room for two. The next case,  $n = 2$ , can easily be done by hand; there are three tilings,  $\begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix}$ ,  $\begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix}$ , and  $\begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix}$ , so  $U_2 = 3$ . (Come to think of it we already knew this, because the previous problem told us that  $T_3 = 3$ ; the number of ways to tile a  $3 \times 2$  rectangle is the same as the number to tile a  $2 \times 3$ .) When  $n = 3$ , as when  $n = 1$ , there are no tilings. We can convince ourselves of this either by making a quick exhaustive search or by looking at the problem from a higher level: The area of a  $3 \times 3$  rectangle is odd, so we can't possibly tile it with dominoes whose area is even. (The same argument obviously applies to any odd  $n$ .) Finally, when  $n = 4$  there seem to be about a dozen tilings; it's difficult to be sure about the exact number without spending a lot of time to guarantee that the list is complete.

So let's try the infinite-sum approach that worked last time:

$$U = | + \begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix} + \begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix} + \begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix} + \begin{smallmatrix} \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \end{smallmatrix} + \begin{smallmatrix} \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \end{smallmatrix} + \begin{smallmatrix} \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \end{smallmatrix} + \begin{smallmatrix} \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \end{smallmatrix} + \begin{smallmatrix} \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \end{smallmatrix} + \cdots. \quad (7.7)$$

Every non-null tiling begins with either  $\begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix}$  or  $\begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix}$  or  $\begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix}$ ; but unfortunately the first two of these three possibilities don't simply factor out and leave us with  $U$  again. The sum of all terms in  $U$  that begin with  $\begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix}$  can, however, be written as  $\begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix} V$ , where

$$V = \begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix} + \begin{smallmatrix} \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \end{smallmatrix} + \begin{smallmatrix} \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \end{smallmatrix} + \begin{smallmatrix} \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \end{smallmatrix} + \begin{smallmatrix} \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \end{smallmatrix} + \cdots$$

is the sum of all domino tilings of a mutilated  $3 \times n$  rectangle that has its lower left corner missing. Similarly, the terms of  $U$  that begin with  $\begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix}$  can be written  $\begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix} \Lambda$ , where

$$\Lambda = \begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix} + \begin{smallmatrix} \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \end{smallmatrix} + \begin{smallmatrix} \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \end{smallmatrix} + \begin{smallmatrix} \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \end{smallmatrix} + \begin{smallmatrix} \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \\ \blacksquare & \blacksquare & \blacksquare \end{smallmatrix} + \cdots$$

consists of all rectangular tilings lacking their upper left corner. The series  $\Lambda$  is a mirror image of  $V$ . These factorizations allow us to write

$$U = | + \begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix} V + \begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix} \Lambda + \begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix} U.$$

And we can factor  $V$  and  $\Lambda$  as well, because such tilings can begin in only two ways:

$$\begin{aligned} V &= \begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix} U + \begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix} V, \\ \Lambda &= \begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix} U + \begin{smallmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{smallmatrix} \Lambda. \end{aligned}$$

Now we have three equations in three unknowns ( $U$ ,  $V$ , and  $\Lambda$ ). We can solve them by first solving for  $V$  and  $\Lambda$  in terms of  $U$ , then plugging the results into the equation for  $U$ :

$$\begin{aligned} V &= (I - \square) \square U, & \Lambda &= (I - \square) \square U; \\ U &= I + \square(I - \square) \square U + \square(I - \square) \square U + \square U. \end{aligned}$$

And the final equation can be solved for  $U$ , giving the compact formula

$$U = \frac{I}{I - \square(I - \square) \square - \square(I - \square) \square - \square}. \quad (7.8)$$

earned in another  
ss about “regular  
pressions.” If I’m  
mistaken, we  
write

$$= (\square \square^* \square + \square \square^* \square + \square)^*$$

the language of  
ular expressions;  
here must be  
ne connection  
ween regular  
pressions and gen-  
ting functions.

This expression defines the infinite sum  $U$ , just as (7.4) defines  $T$ .

The next step is to go commutative. Everything simplifies beautifully when we detach all the dominoes and use only powers of  $\square$  and  $\square$ :

$$\begin{aligned} U &= \frac{1}{1 - \square^2 \square (1 - \square^3)^{-1} - \square^2 \square (1 - \square^3)^{-1} - \square^3} \\ &= \frac{1 - \square^3}{(1 - \square^3)^2 - 2\square^2 \square} \\ &= \frac{(1 - \square^3)^{-1}}{1 - 2\square^2 \square (1 - \square^3)^{-2}} \\ &= \frac{1}{1 - \square^3} + \frac{2\square^2 \square}{(1 - \square^3)^3} + \frac{4\square^4 \square^2}{(1 - \square^3)^5} + \frac{8\square^6 \square^3}{(1 - \square^3)^7} + \cdots \\ &= \sum_{k \geq 0} \frac{2^k \square^{2k} \square^k}{(1 - \square^3)^{2k+1}} \\ &= \sum_{k, m \geq 0} \binom{m+2k}{m} 2^k \square^{2k} \square^{k+3m}. \end{aligned}$$

(This derivation deserves careful scrutiny. The last step uses the formula  $(1 - w)^{-2k-1} = \sum_m \binom{m+2k}{m} w^m$ , identity (5.56).) Let’s take a good look at the bottom line to see what it tells us. First, it says that every  $3 \times n$  tiling uses an even number of vertical dominoes. Moreover, if there are  $2k$  verticals, there must be at least  $k$  horizontals, and the total number of horizontals must be  $k + 3m$  for some  $m \geq 0$ . Finally, the number of possible tilings with  $2k$  verticals and  $k + 3m$  horizontals is exactly  $\binom{m+2k}{m} 2^k$ .

We now are able to analyze the  $3 \times 4$  tilings that left us doubtful when we began looking at the  $3 \times n$  problem. When  $n = 4$  the total area is 12, so we need six dominoes altogether. There are  $2k$  verticals and  $k + 3m$  horizontals,

for some  $k$  and  $m$ ; hence  $2k + k + 3m = 6$ . In other words,  $k + m = 2$ . If we use no verticals, then  $k = 0$  and  $m = 2$ ; the number of possibilities is  $\binom{2+0}{2}2^0 = 1$ . (This accounts for the tiling  $\begin{smallmatrix} \square & \square \\ \square & \square \end{smallmatrix}$ .) If we use two verticals, then  $k = 1$  and  $m = 1$ ; there are  $\binom{1+2}{1}2^1 = 6$  such tilings. And if we use four verticals, then  $k = 2$  and  $m = 0$ ; there are  $\binom{0+4}{0}2^2 = 4$  such tilings, making a total of  $U_4 = 11$ . In general if  $n$  is even, this reasoning shows that  $k + m = \frac{1}{2}n$ , hence  $\binom{m+2k}{m} = \binom{n/2+k}{n/2-k}$  and the total number of  $3 \times n$  tilings is

$$U_n = \sum_k \binom{n/2+k}{n/2-k} 2^k = \sum_m \binom{n-m}{m} 2^{n/2-m}. \quad (7.9)$$

As before, we can also substitute  $z$  for both  $\square$  and  $\sqcap$ , getting a generating function that doesn't discriminate between dominoes of particular persuasions. The result is

$$U = \frac{1}{1 - z^3(1 - z^3)^{-1} - z^3(1 - z^3)^{-1} - z^3} = \frac{1 - z^3}{1 - 4z^3 + z^6}. \quad (7.10)$$

If we expand this quotient into a power series, we get

$$U = 1 + U_2 z^3 + U_4 z^6 + U_6 z^9 + U_8 z^{12} + \dots,$$

a generating function for the numbers  $U_n$ . (There's a curious mismatch between subscripts and exponents in this formula, but it is easily explained. The coefficient of  $z^9$ , for example, is  $U_6$ , which counts the tilings of a  $3 \times 6$  rectangle. This is what we want, because every such tiling contains nine dominoes.)

We could proceed to analyze (7.10) and get a closed form for the coefficients, but it's better to save that for later in the chapter after we've gotten more experience. So let's divest ourselves of dominoes for the moment and proceed to the next advertised problem, "change."

How many ways are there to pay 50 cents? We assume that the payment must be made with pennies  $\textcircled{1}$ , nickels  $\textcircled{5}$ , dimes  $\textcircled{10}$ , quarters  $\textcircled{25}$ , and half-dollars  $\textcircled{50}$ . George Pólya [298] popularized this problem by showing that it can be solved with generating functions in an instructive way.

*Ah yes, I remember when we had half-dollars.*

Let's set up infinite sums that represent all possible ways to give change, just as we tackled the domino problems by working with infinite sums that represent all possible domino patterns. It's simplest to start by working with fewer varieties of coins, so let's suppose first that we have nothing but pennies. The sum of all ways to leave some number of pennies (but just pennies) in change can be written

$$\begin{aligned} P &= \not{x} + \textcircled{1} + \textcircled{1}\textcircled{1} + \textcircled{1}\textcircled{1}\textcircled{1} + \textcircled{1}\textcircled{1}\textcircled{1}\textcircled{1} + \dots \\ &= \not{x} + \textcircled{1} + \textcircled{1}^2 + \textcircled{1}^3 + \textcircled{1}^4 + \dots \end{aligned}$$



The first term stands for the way to leave no pennies, the second term stands for one penny, then two pennies, three pennies, and so on. Now if we're allowed to use both pennies and nickels, the sum of all possible ways is

$$\begin{aligned} N &= P + \textcircled{5} P + \textcircled{5}\textcircled{5} P + \textcircled{5}\textcircled{5}\textcircled{5} P + \textcircled{5}\textcircled{5}\textcircled{5}\textcircled{5} P + \cdots \\ &= (\not{\$} + \textcircled{5} + \textcircled{5}^2 + \textcircled{5}^3 + \textcircled{5}^4 + \cdots) P, \end{aligned}$$

since each payment has a certain number of nickels chosen from the first factor and a certain number of pennies chosen from  $P$ . (Notice that  $N$  is *not* the sum  $\not{\$} + \textcircled{1} + \textcircled{5} + (\textcircled{1} + \textcircled{5})^2 + (\textcircled{1} + \textcircled{5})^3 + \cdots$ , because such a sum includes many types of payment more than once. For example, the term  $(\textcircled{1} + \textcircled{5})^2 = \textcircled{1}\textcircled{1} + \textcircled{1}\textcircled{5} + \textcircled{5}\textcircled{1} + \textcircled{5}\textcircled{5}$  treats  $\textcircled{1}\textcircled{5}$  and  $\textcircled{5}\textcircled{1}$  as if they were different, but we want to list each set of coins only once without respect to order.)

Similarly, if dimes are permitted as well, we get the infinite sum

$$D = (\not{\$} + \textcircled{10} + \textcircled{10}^2 + \textcircled{10}^3 + \textcircled{10}^4 + \cdots) N,$$

which includes terms like  $\textcircled{10}^3 \textcircled{5}^3 \textcircled{1}^5 = \textcircled{10}\textcircled{10}\textcircled{10}\textcircled{5}\textcircled{5}\textcircled{5}\textcircled{1}\textcircled{1}\textcircled{1}\textcircled{1}\textcircled{1}$  when it is expanded in full. Each of these terms is a different way to make change.

ns of the realm.

Adding quarters and then half-dollars to the realm of possibilities gives

$$\begin{aligned} Q &= (\not{\$} + \textcircled{25} + \textcircled{25}^2 + \textcircled{25}^3 + \textcircled{25}^4 + \cdots) D; \\ C &= (\not{\$} + \textcircled{50} + \textcircled{50}^2 + \textcircled{50}^3 + \textcircled{50}^4 + \cdots) Q. \end{aligned}$$

Our problem is to find the number of terms in  $C$  worth exactly 50¢.

A simple trick solves this problem nicely: We can replace  $\textcircled{1}$  by  $z$ ,  $\textcircled{5}$  by  $z^5$ ,  $\textcircled{10}$  by  $z^{10}$ ,  $\textcircled{25}$  by  $z^{25}$ , and  $\textcircled{50}$  by  $z^{50}$ . Then each term is replaced by  $z^n$ , where  $n$  is the monetary value of the original term. For example, the term  $\textcircled{50}\textcircled{10}\textcircled{5}\textcircled{5}\textcircled{1}$  becomes  $z^{50+10+5+5+1} = z^{71}$ . The four ways of paying 13 cents, namely  $\textcircled{10}\textcircled{1}^3$ ,  $\textcircled{5}\textcircled{1}^8$ ,  $\textcircled{5}^2\textcircled{1}^3$ , and  $\textcircled{1}^{13}$ , each reduce to  $z^{13}$ ; hence the coefficient of  $z^{13}$  will be 4 after the  $z$ -substitutions are made.

Let  $P_n$ ,  $N_n$ ,  $D_n$ ,  $Q_n$ , and  $C_n$  be the numbers of ways to pay  $n$  cents when we're allowed to use coins that are worth at most 1, 5, 10, 25, and 50 cents, respectively. Our analysis tells us that these are the coefficients of  $z^n$  in the respective power series

$$\begin{aligned} P &= 1 + z + z^2 + z^3 + z^4 + \cdots, \\ N &= (1 + z^5 + z^{10} + z^{15} + z^{20} + \cdots) P, \\ D &= (1 + z^{10} + z^{20} + z^{30} + z^{40} + \cdots) N, \\ Q &= (1 + z^{25} + z^{50} + z^{75} + z^{100} + \cdots) D, \\ C &= (1 + z^{50} + z^{100} + z^{150} + z^{200} + \cdots) Q. \end{aligned}$$

Obviously  $P_n = 1$  for all  $n \geq 0$ . And a little thought proves that we have  $N_n = \lfloor n/5 \rfloor + 1$ : To make  $n$  cents out of pennies and nickels, we must choose either 0 or 1 or ... or  $\lfloor n/5 \rfloor$  nickels, after which there's only one way to supply the requisite number of pennies. Thus  $P_n$  and  $N_n$  are simple; but the values of  $D_n$ ,  $Q_n$ , and  $C_n$  are increasingly more complicated.

*How many pennies are there, really? If  $n$  is greater than, say,  $10^{10}$ , I bet that  $P_n = 0$  in the "real world."*

One way to deal with these formulas is to realize that  $1 + z^m + z^{2m} + \dots$  is just  $1/(1 - z^m)$ . Thus we can write

$$\begin{aligned} P &= 1/(1 - z), \\ N &= P/(1 - z^5), \\ D &= N/(1 - z^{10}), \\ Q &= D/(1 - z^{25}), \\ C &= Q/(1 - z^{50}). \end{aligned}$$

Multiplying by the denominators, we have

$$\begin{aligned} (1 - z)P &= 1, \\ (1 - z^5)N &= P, \\ (1 - z^{10})D &= N, \\ (1 - z^{25})Q &= D, \\ (1 - z^{50})C &= Q. \end{aligned}$$

Now we can equate coefficients of  $z^n$  in these equations, getting recurrence relations from which the desired coefficients can quickly be computed:

$$\begin{aligned} P_n &= P_{n-1} + [n=0], \\ N_n &= N_{n-5} + P_n, \\ D_n &= D_{n-10} + N_n, \\ Q_n &= Q_{n-25} + D_n, \\ C_n &= C_{n-50} + Q_n. \end{aligned}$$

For example, the coefficient of  $z^n$  in  $D = (1 - z^{25})Q$  is equal to  $Q_n - Q_{n-25}$ ; so we must have  $Q_n - Q_{n-25} = D_n$ , as claimed.

We could unfold these recurrences and find, for example, that  $Q_n = D_n + D_{n-25} + D_{n-50} + D_{n-75} + \dots$ , stopping when the subscripts get negative. But the non-iterated form is convenient because each coefficient is computed with just one addition, as in Pascal's triangle.

Let's use the recurrences to find  $C_{50}$ . First,  $C_{50} = C_0 + Q_{50}$ ; so we want to know  $Q_{50}$ . Then  $Q_{50} = Q_{25} + D_{50}$ , and  $Q_{25} = Q_0 + D_{25}$ ; so we also want to know  $D_{50}$  and  $D_{25}$ . These  $D_n$  depend in turn on  $D_{40}$ ,  $D_{30}$ ,  $D_{20}$ ,  $D_{15}$ ,  $D_{10}$ ,  $D_5$ , and on  $N_{50}$ ,  $N_{45}$ , ...,  $N_5$ . A simple calculation therefore suffices

to determine all the necessary coefficients:

| n     | 0 | 5 | 10 | 15 | 20 | 25 | 30 | 35 | 40 | 45 | 50 |
|-------|---|---|----|----|----|----|----|----|----|----|----|
| $P_n$ | 1 | 1 | 1  | 1  | 1  | 1  | 1  | 1  | 1  | 1  | 1  |
| $N_n$ | 1 | 2 | 3  | 4  | 5  | 6  | 7  | 8  | 9  | 10 | 11 |
| $D_n$ | 1 | 2 | 4  | 6  | 9  | 12 | 16 |    | 25 |    | 36 |
| $Q_n$ | 1 |   |    |    |    | 13 |    |    |    |    | 49 |
| $C_n$ | 1 |   |    |    |    |    |    |    |    |    | 50 |

The final value in the table gives us our answer,  $C_{50}$ : There are exactly 50 ways to leave a 50-cent tip.

How about a closed form for  $C_n$ ? Multiplying the equations together gives us the compact expression

$$C = \frac{1}{1-z} \frac{1}{1-z^5} \frac{1}{1-z^{10}} \frac{1}{1-z^{25}} \frac{1}{1-z^{50}}, \quad (7.11)$$

but it's not obvious how to get from here to the coefficient of  $z^n$ . Fortunately there is a way; we'll return to this problem later in the chapter.

More elegant formulas arise if we consider the problem of giving change when we live in a land that mints coins of every positive integer denomination ( $\textcircled{1}$ ,  $\textcircled{2}$ ,  $\textcircled{3}$ , ...) instead of just the five we allowed before. The corresponding generating function is an infinite product of fractions,

$$\frac{1}{(1-z)(1-z^2)(1-z^3)\dots},$$

and the coefficient of  $z^n$  when these factors are fully multiplied out is called  $p(n)$ , the number of *partitions* of  $n$ . A partition of  $n$  is a representation of  $n$  as a sum of positive integers, disregarding order. For example, there are seven different partitions of 5, namely

$$5 = 4+1 = 3+2 = 3+1+1 = 2+2+1 = 2+1+1+1 = 1+1+1+1+1;$$

hence  $p(5) = 7$ . (Also  $p(2) = 2$ ,  $p(3) = 3$ ,  $p(4) = 5$ , and  $p(6) = 11$ ; it begins to look as if  $p(n)$  is always a prime number. But  $p(7) = 15$ , spoiling the pattern.) There is no closed form for  $p(n)$ , but the theory of partitions is a fascinating branch of mathematics in which many remarkable discoveries have been made. For example, Ramanujan proved that  $p(5n+4) \equiv 0 \pmod{5}$ ,  $p(7n+5) \equiv 0 \pmod{7}$ , and  $p(11n+6) \equiv 0 \pmod{11}$ , by making ingenious transformations of generating functions (see Andrews [11, Chapter 10]).

## 7.2 BASIC MANEUVERS

Now let's look more closely at some of the techniques that make power series powerful.

First a few words about terminology and notation. Our generic generating function has the form

$$G(z) = g_0 + g_1z + g_2z^2 + \cdots = \sum_{n \geq 0} g_n z^n, \quad (7.12)$$

and we say that  $G(z)$ , or  $G$  for short, is the generating function for the sequence  $\langle g_0, g_1, g_2, \dots \rangle$ , which we also call  $\langle g_n \rangle$ . The coefficient  $g_n$  of  $z^n$  in  $G(z)$  is often denoted  $[z^n]G(z)$ , as in Section 5.4.

The sum in (7.12) runs over all  $n \geq 0$ , but we often find it more convenient to extend the sum over all integers  $n$ . We can do this by simply regarding  $g_{-1} = g_{-2} = \cdots = 0$ . In such cases we might still talk about the sequence  $\langle g_0, g_1, g_2, \dots \rangle$ , as if the  $g_n$ 's didn't exist for negative  $n$ .

Two kinds of "closed forms" come up when we work with generating functions. We might have a closed form for  $G(z)$ , expressed in terms of  $z$ ; or we might have a closed form for  $g_n$ , expressed in terms of  $n$ . For example, the generating function for Fibonacci numbers has the closed form  $z/(1 - z - z^2)$ ; the Fibonacci numbers themselves have the closed form  $(\phi^n - \hat{\phi}^n)/\sqrt{5}$ . The context will explain what kind of closed form is meant.

Now a few words about perspective. The generating function  $G(z)$  appears to be two different entities, depending on how we view it. Sometimes it is a function of a complex variable  $z$ , satisfying all the standard properties proved in calculus books. And sometimes it is simply a formal power series, with  $z$  acting as a placeholder. In the previous section, for example, we used the second interpretation; we saw several examples in which  $z$  was substituted for some feature of a combinatorial object in a "sum" of such objects. The coefficient of  $z^n$  was then the number of combinatorial objects having  $n$  occurrences of that feature.

When we view  $G(z)$  as a function of a complex variable, its convergence becomes an issue. We said in Chapter 2 that the infinite series  $\sum_{n \geq 0} g_n z^n$  converges (absolutely) if and only if there's a bounding constant  $A$  such that the finite sums  $\sum_{0 \leq n \leq N} |g_n z^n|$  never exceed  $A$ , for any  $N$ . Therefore it's easy to see that if  $\sum_{n \geq 0} g_n z^n$  converges for some value  $z = z_0$ , it also converges for all  $z$  with  $|z| < |z_0|$ . Furthermore, we must have  $\lim_{n \rightarrow \infty} |g_n z_0^n| = 0$ ; hence, in the notation of Chapter 9,  $g_n = O(|1/z_0|^n)$  if there is convergence at  $z_0$ . And conversely if  $g_n = O(M^n)$ , the series  $\sum_{n \geq 0} g_n z^n$  converges for all  $|z| < 1/M$ . These are the basic facts about convergence of power series.

But for our purposes convergence is usually a red herring, unless we're trying to study the asymptotic behavior of the coefficients. Nearly every

*If physicists can get away with viewing light sometimes as a wave and sometimes as a particle, mathematicians should be able to view generating functions in two different ways.*

operation we perform on generating functions can be justified rigorously as an operation on formal power series, and such operations are legal even when the series don't converge. (The relevant theory can be found, for example, in Bell [23], Niven [282], and Henrici [182, Chapter 1].)

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Furthermore, even if we throw all caution to the winds and derive formulas without any rigorous justification, we generally can take the results of our derivation and prove them by induction. For example, the generating function for the Fibonacci numbers converges only when  $|z| < 1/\phi \approx 0.618$ , but we didn't need to know that when we proved the formula  $F_n = (\phi^n - \bar{\phi}^n)/\sqrt{5}$ . The latter formula, once discovered, can be verified directly, if we don't trust the theory of formal power series. Therefore we'll ignore questions of convergence in this chapter; it's more a hindrance than a help.

So much for perspective. Next we look at our main tools for reshaping generating functions—adding, shifting, changing variables, differentiating, integrating, and multiplying. In what follows we assume that, unless stated otherwise,  $F(z)$  and  $G(z)$  are the generating functions for the sequences  $\langle f_n \rangle$  and  $\langle g_n \rangle$ . We also assume that the  $f_n$ 's and  $g_n$ 's are zero for negative  $n$ , since this saves us some bickering with the limits of summation.

It's pretty obvious what happens when we add constant multiples of  $F$  and  $G$  together:

$$\begin{aligned}\alpha F(z) + \beta G(z) &= \alpha \sum_n f_n z^n + \beta \sum_n g_n z^n \\ &= \sum_n (\alpha f_n + \beta g_n) z^n.\end{aligned}\tag{7.13}$$

This gives us the generating function for the sequence  $\langle \alpha f_n + \beta g_n \rangle$ .

Shifting a generating function isn't much harder. To shift  $G(z)$  right by  $m$  places, that is, to form the generating function for the sequence  $\langle 0, \dots, 0, g_0, g_1, \dots \rangle = \langle g_{n-m} \rangle$  with  $m$  leading 0's, we simply multiply by  $z^m$ :

$$z^m G(z) = \sum_n g_n z^{n+m} = \sum_n g_{n-m} z^n, \quad \text{integer } m \geq 0.\tag{7.14}$$

This is the operation we used (twice), along with addition, to deduce the equation  $(1 - z - z^2)F(z) = z$  on our way to finding a closed form for the Fibonacci numbers in Chapter 6.

And to shift  $G(z)$  left  $m$  places—that is, to form the generating function for the sequence  $\langle g_m, g_{m+1}, g_{m+2}, \dots \rangle = \langle g_{n+m} \rangle$  with the first  $m$  elements discarded—we subtract off the first  $m$  terms and then divide by  $z^m$ :

$$\frac{G(z) - g_0 - g_1 z - \dots - g_{m-1} z^{m-1}}{z^m} = \sum_{n \geq m} g_n z^{n-m} = \sum_{n \geq 0} g_{n+m} z^n.\tag{7.15}$$

(We can't extend this last sum over all  $n$  unless  $g_0 = \dots = g_{m-1} = 0$ .)

Replacing the  $z$  by a constant multiple is another of our tricks:

$$G(cz) = \sum_n g_n (cz)^n = \sum_n c^n g_n z^n; \quad (7.16)$$

this yields the generating function for the sequence  $\langle c^n g_n \rangle$ . The special case  $c = -1$  is particularly useful.

Often we want to bring down a factor of  $n$  into the coefficient. Differentiation is what lets us do that:

*I fear differentiating  $dz$ 's.*

$$G'(z) = g_1 + 2g_2 z + 3g_3 z^2 + \cdots = \sum_n (n+1)g_{n+1} z^n. \quad (7.17)$$

Shifting this right one place gives us a form that's sometimes more useful,

$$zG'(z) = \sum_n n g_n z^n. \quad (7.18)$$

This is the generating function for the sequence  $\langle n g_n \rangle$ . Repeated differentiation would allow us to multiply  $g_n$  by any desired polynomial in  $n$ .

Integration, the inverse operation, lets us divide the terms by  $n$ :

$$\int_0^z G(t) dt = g_0 z + \frac{1}{2} g_1 z^2 + \frac{1}{3} g_2 z^3 + \cdots = \sum_{n \geq 1} \frac{1}{n} g_{n-1} z^n. \quad (7.19)$$

(Notice that the constant term is zero.) If we want the generating function for  $\langle g_n/n \rangle$  instead of  $\langle g_{n-1}/n \rangle$ , we should first shift left one place, replacing  $G(t)$  by  $(G(t) - g_0)/t$  in the integral.

Finally, here's how we multiply generating functions together:

$$\begin{aligned} F(z)G(z) &= (f_0 + f_1 z + f_2 z^2 + \cdots)(g_0 + g_1 z + g_2 z^2 + \cdots) \\ &= (f_0 g_0) + (f_0 g_1 + f_1 g_0)z + (f_0 g_2 + f_1 g_1 + f_2 g_0)z^2 + \cdots \\ &= \sum_n \left( \sum_k f_k g_{n-k} \right) z^n. \end{aligned} \quad (7.20)$$

As we observed in Chapter 5, this gives the generating function for the sequence  $\langle h_n \rangle$ , the *convolution* of  $\langle f_n \rangle$  and  $\langle g_n \rangle$ . The sum  $h_n = \sum_k f_k g_{n-k}$  can also be written  $h_n = \sum_{k=0}^n f_k g_{n-k}$ , because  $f_k = 0$  when  $k < 0$  and  $g_{n-k} = 0$  when  $k > n$ . Multiplication/convolution is a little more complicated than the other operations, but it's very useful—so useful that we will spend all of Section 7.5 below looking at examples of it.

Multiplication has several special cases that are worth considering as operations in themselves. We've already seen one of these: When  $F(z) = z^m$  we get the shifting operation (7.14). In that case the sum  $h_n$  becomes the single term  $g_{n-m}$ , because all  $f_k$ 's are 0 except for  $f_m = 1$ .

**Table 335** Generating function manipulations.

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|   |
|---|
| $\alpha F(z) + \beta G(z) = \sum_n (\alpha f_n + \beta g_n) z^n$  |
| $z^m G(z) = \sum_n g_{n-m} z^n, \quad \text{integer } m \geq 0$   |
| $\frac{G(z) - g_0 - g_1 z - \cdots - g_{m-1} z^{m-1}}{z^m} = \sum_{n \geq 0} g_{n+m} z^n, \quad \text{integer } m \geq 0$ |
| $G(cz) = \sum_n c^n g_n z^n$  |
| $G'(z) = \sum_n (n+1) g_{n+1} z^n$  |
| $zG'(z) = \sum_n n g_n z^n$   |
| $\int_0^z G(t) dt = \sum_{n \geq 1} \frac{1}{n} g_{n-1} z^n$  |
| $F(z)G(z) = \sum_n \left( \sum_k f_k g_{n-k} \right) z^n$   |
| $\frac{1}{1-z} G(z) = \sum_n \left( \sum_{k \leq n} g_k \right) z^n$  |

---

Another useful special case arises when  $F(z)$  is the familiar function  $1/(1-z) = 1 + z + z^2 + \cdots$ ; then all  $f_k$ 's (for  $k \geq 0$ ) are 1 and we have the important formula

$$\frac{1}{1-z} G(z) = \sum_n \left( \sum_{k \geq 0} g_{n-k} \right) z^n = \sum_n \left( \sum_{k \leq n} g_k \right) z^n. \quad (7.21)$$

Multiplying a generating function by  $1/(1-z)$  gives us the generating function for the cumulative sums of the original sequence.

Table 335 summarizes the operations we've discussed so far. To use all these manipulations effectively it helps to have a healthy repertoire of generating functions in stock. Table 336 lists the simplest ones; we can use those to get started and to solve quite a few problems.

Each of the generating functions in Table 336 is important enough to be memorized. Many of them are special cases of the others, and many of

**Table 336** Simple sequences and their generating functions.

| sequence   | generating function                        | closed form             |
|--|--|-------------------------|
| $\langle 1, 0, 0, 0, 0, \dots \rangle$   | $\sum_{n \geq 0} [n=0] z^n$                | 1                       |
| $\langle 0, \dots, 0, 1, 0, 0, \dots \rangle$  | $\sum_{n \geq 0} [n=m] z^n$                | $z^m$                   |
| $\langle 1, 1, 1, 1, 1, \dots \rangle$   | $\sum_{n \geq 0} z^n$                      | $\frac{1}{1-z}$         |
| $\langle 1, -1, 1, -1, 1, \dots \rangle$   | $\sum_{n \geq 0} (-1)^n z^n$               | $\frac{1}{1+z}$         |
| $\langle 1, 0, 1, 0, 1, 0, \dots \rangle$  | $\sum_{n \geq 0} [2 \nmid n] z^n$          | $\frac{1}{1-z^2}$       |
| $\langle 1, 0, \dots, 0, 1, 0, \dots, 0, 1, 0, \dots \rangle$                        | $\sum_{n \geq 0} [m \nmid n] z^n$          | $\frac{1}{1-z^m}$       |
| $\langle 1, 2, 3, 4, 5, 6, \dots \rangle$  | $\sum_{n \geq 0} (n+1) z^n$                | $\frac{1}{(1-z)^2}$     |
| $\langle 1, 2, 4, 8, 16, 32, \dots \rangle$  | $\sum_{n \geq 0} 2^n z^n$                  | $\frac{1}{1-2z}$        |
| $\langle 1, 4, 6, 4, 1, 0, 0, \dots \rangle$   | $\sum_{n \geq 0} \binom{4}{n} z^n$         | $(1+z)^4$               |
| $\langle 1, c, \binom{c}{2}, \binom{c}{3}, \dots \rangle$                            | $\sum_{n \geq 0} \binom{c}{n} z^n$         | $(1+z)^c$               |
| $\langle 1, c, \binom{c+1}{2}, \binom{c+2}{3}, \dots \rangle$                        | $\sum_{n \geq 0} \binom{c+n-1}{n} z^n$     | $\frac{1}{(1-z)^c}$     |
| $\langle 1, c, c^2, c^3, \dots \rangle$  | $\sum_{n \geq 0} c^n z^n$                  | $\frac{1}{1-cz}$        |
| $\langle 1, \binom{m+1}{m}, \binom{m+2}{m}, \binom{m+3}{m}, \dots \rangle$           | $\sum_{n \geq 0} \binom{m+n}{m} z^n$       | $\frac{1}{(1-z)^{m+1}}$ |
| $\langle 0, 1, \frac{1}{2}, \frac{1}{3}, \frac{1}{4}, \dots \rangle$                 | $\sum_{n \geq 1} \frac{1}{n} z^n$          | $\ln \frac{1}{1-z}$     |
| $\langle 0, 1, -\frac{1}{2}, \frac{1}{3}, -\frac{1}{4}, \dots \rangle$               | $\sum_{n \geq 1} \frac{(-1)^{n+1}}{n} z^n$ | $\ln(1+z)$              |
| $\langle 1, 1, \frac{1}{2}, \frac{1}{6}, \frac{1}{24}, \frac{1}{120}, \dots \rangle$ | $\sum_{n \geq 0} \frac{1}{n!} z^n$         | $e^z$                   |

them can be derived quickly from the others by using the basic operations of Table 335; therefore the memory work isn't very hard.

For example, let's consider the sequence  $\langle 1, 2, 3, 4, \dots \rangle$ , whose generating function  $1/(1-z)^2$  is often useful. This generating function appears near the

*Hint: If the sequence consists of binomial coefficients, its generating function usually involves a binomial,  $1 \pm z$ .*



middle of Table 336, and it's also the special case  $m = 1$  of  $\langle 1, \binom{m+1}{m}, \binom{m+2}{m}, \binom{m+3}{m}, \dots \rangle$ , which appears further down; it's also the special case  $c = 2$  of the closely related sequence  $\langle 1, c, \binom{c+1}{2}, \binom{c+2}{3}, \dots \rangle$ . We can derive it from the generating function for  $\langle 1, 1, 1, 1, \dots \rangle$  by taking cumulative sums as in (7.21); that is, by dividing  $1/(1-z)$  by  $(1-z)$ . Or we can derive it from  $\langle 1, 1, 1, 1, \dots \rangle$  by differentiation, using (7.17).

The sequence  $\langle 1, 0, 1, 0, \dots \rangle$  is another one whose generating function can be obtained in many ways. We can obviously derive the formula  $\sum_n z^{2n} = 1/(1-z^2)$  by substituting  $z^2$  for  $z$  in the identity  $\sum_n z^n = 1/(1-z)$ ; we can also apply cumulative summation to the sequence  $\langle 1, -1, 1, -1, \dots \rangle$ , whose generating function is  $1/(1+z)$ , getting  $1/(1+z)(1-z) = 1/(1-z^2)$ . And there's also a third way, which is based on a general method for extracting the even-numbered terms  $\langle g_0, 0, g_2, 0, g_4, 0, \dots \rangle$  of *any* given sequence: If we add  $G(-z)$  to  $G(+z)$  we get

$$G(z) + G(-z) = \sum_n g_n (1 + (-1)^n) z^n = 2 \sum_n g_n [n \text{ even}] z^n;$$

therefore

$$\frac{G(z) + G(-z)}{2} = \sum_n g_{2n} z^{2n}. \quad (7.22)$$

The odd-numbered terms can be extracted in a similar way,

$$\frac{G(z) - G(-z)}{2} = \sum_n g_{2n+1} z^{2n+1}. \quad (7.23)$$

In the special case where  $g_n = 1$  and  $G(z) = 1/(1-z)$ , the generating function for  $\langle 1, 0, 1, 0, \dots \rangle$  is  $\frac{1}{2}(G(z) + G(-z)) = \frac{1}{2}\left(\frac{1}{1-z} + \frac{1}{1+z}\right) = \frac{1}{1-z^2}$ .

Let's try this extraction trick on the generating function for Fibonacci numbers. We know that  $\sum_n F_n z^n = z/(1-z-z^2)$ ; hence

$$\begin{aligned} \sum_n F_{2n} z^{2n} &= \frac{1}{2} \left( \frac{z}{1-z-z^2} + \frac{-z}{1+z-z^2} \right) \\ &= \frac{1}{2} \left( \frac{z+z^2-z^3-z+z^2+z^3}{(1-z^2)^2-z^2} \right) = \frac{z^2}{1-3z^2+z^4}. \end{aligned}$$

This generates the sequence  $\langle F_0, 0, F_2, 0, F_4, \dots \rangle$ ; hence the sequence of alternate  $F$ 's,  $\langle F_0, F_2, F_4, F_6, \dots \rangle = \langle 0, 1, 3, 8, \dots \rangle$ , has a simple generating function:

$$\sum_n F_{2n} z^n = \frac{z}{1-3z+z^2}. \quad (7.24)$$

## 7.3 SOLVING RECURRENCES

Now let's focus our attention on one of the most important uses of generating functions: the solution of recurrence relations.

Given a sequence  $\langle g_n \rangle$  that satisfies a given recurrence, we seek a closed form for  $g_n$  in terms of  $n$ . A solution to this problem via generating functions proceeds in four steps that are almost mechanical enough to be programmed on a computer:

- 1 Write down a single equation that expresses  $g_n$  in terms of other elements of the sequence. This equation should be valid for all integers  $n$ , assuming that  $g_{-1} = g_{-2} = \cdots = 0$ .
- 2 Multiply both sides of the equation by  $z^n$  and sum over all  $n$ . This gives, on the left, the sum  $\sum_n g_n z^n$ , which is the generating function  $G(z)$ . The right-hand side should be manipulated so that it becomes some other expression involving  $G(z)$ .
- 3 Solve the resulting equation, getting a closed form for  $G(z)$ .
- 4 Expand  $G(z)$  into a power series and read off the coefficient of  $z^n$ ; this is a closed form for  $g_n$ .

This method works because the single function  $G(z)$  represents the entire sequence  $\langle g_n \rangle$  in such a way that many manipulations are possible.

### *Example 1: Fibonacci numbers revisited.*

For example, let's rerun the derivation of Fibonacci numbers from Chapter 6. In that chapter we were feeling our way, learning a new method; now we can be more systematic. The given recurrence is

$$\begin{aligned} g_0 &= 0; & g_1 &= 1; \\ g_n &= g_{n-1} + g_{n-2}, & \text{for } n \geq 2. \end{aligned}$$

We will find a closed form for  $g_n$  by using the four steps above.

Step 1 tells us to write the recurrence as a "single equation" for  $g_n$ . We could say

$$g_n = \begin{cases} 0, & \text{if } n \leq 0; \\ 1, & \text{if } n = 1; \\ g_{n-1} + g_{n-2}, & \text{if } n > 1; \end{cases}$$

but this is cheating. Step 1 really asks for a formula that doesn't involve a case-by-case construction. The single equation

$$g_n = g_{n-1} + g_{n-2}$$

works for  $n \geq 2$ , and it also holds when  $n \leq 0$  (because we have  $g_0 = 0$  and  $g_{\text{negative}} = 0$ ). But when  $n = 1$  we get 1 on the left and 0 on the right.

Fortunately the problem is easy to fix, since we can add  $[n=1]$  to the right; this adds 1 when  $n = 1$ , and it makes no change when  $n \neq 1$ . So, we have

$$g_n = g_{n-1} + g_{n-2} + [n=1];$$

this is the equation called for in Step 1.

Step 2 now asks us to transform the equation for  $\langle g_n \rangle$  into an equation for  $G(z) = \sum_n g_n z^n$ . The task is not difficult:

$$\begin{aligned} G(z) &= \sum_n g_n z^n = \sum_n g_{n-1} z^n + \sum_n g_{n-2} z^n + \sum_n [n=1] z^n \\ &= \sum_n g_n z^{n+1} + \sum_n g_n z^{n+2} + z \\ &= zG(z) + z^2G(z) + z. \end{aligned}$$

Step 3 is also simple in this case; we have

$$G(z) = \frac{z}{1 - z - z^2},$$

which of course comes as no surprise.

Step 4 is the clincher. We carried it out in Chapter 6 by having a sudden flash of inspiration; let's go more slowly now, so that we can get through Step 4 safely later, when we meet problems that are more difficult. What is

$$[z^n] \frac{z}{1 - z - z^2},$$

the coefficient of  $z^n$  when  $z/(1 - z - z^2)$  is expanded in a power series? More generally, if we are given any rational function

$$R(z) = \frac{P(z)}{Q(z)},$$

where  $P$  and  $Q$  are polynomials, what is the coefficient  $[z^n] R(z)$ ?

There's one kind of rational function whose coefficients are particularly nice, namely

$$\frac{a}{(1 - \rho z)^{m+1}} = \sum_{n \geq 0} \binom{m+n}{n} a \rho^n z^n. \quad (7.25)$$

(The case  $\rho = 1$  appears in Table 336, and we can get the general formula shown here by substituting  $\rho z$  for  $z$ .) A finite sum of functions like (7.25),

$$S(z) = \frac{a_1}{(1 - \rho_1 z)^{m_1+1}} + \frac{a_2}{(1 - \rho_2 z)^{m_2+1}} + \cdots + \frac{a_l}{(1 - \rho_l z)^{m_l+1}}, \quad (7.26)$$

also has nice coefficients,

$$[z^n]S(z) = a_1 \binom{m_1+n}{m_1} \rho_1^n + a_2 \binom{m_2+n}{m_2} \rho_2^n + \cdots + a_l \binom{m_l+n}{m_l} \rho_l^n. \quad (7.27)$$

We will show that every rational function  $R(z)$  such that  $R(0) \neq \infty$  can be expressed in the form

$$R(z) = S(z) + T(z), \quad (7.28)$$

where  $S(z)$  has the form (7.26) and  $T(z)$  is a polynomial. Therefore there is a closed form for the coefficients  $[z^n]R(z)$ . Finding  $S(z)$  and  $T(z)$  is equivalent to finding the “partial fraction expansion” of  $R(z)$ .

Notice that  $S(z) = \infty$  when  $z$  has the values  $1/\rho_1, \dots, 1/\rho_l$ . Therefore the numbers  $\rho_k$  that we need to find, if we’re going to succeed in expressing  $R(z)$  in the desired form  $S(z) + T(z)$ , must be the reciprocals of the numbers  $\alpha_k$  where  $Q(\alpha_k) = 0$ . (Recall that  $R(z) = P(z)/Q(z)$ , where  $P$  and  $Q$  are polynomials; we have  $R(z) = \infty$  only if  $Q(z) = 0$ .)

Suppose  $Q(z)$  has the form

$$Q(z) = q_0 + q_1 z + \cdots + q_m z^m, \quad \text{where } q_0 \neq 0 \text{ and } q_m \neq 0.$$

The “reflected” polynomial

$$Q^R(z) = q_0 z^m + q_1 z^{m-1} + \cdots + q_m$$

has an important relation to  $Q(z)$ :

$$\begin{aligned} Q^R(z) &= q_0(z - \rho_1) \cdots (z - \rho_m) \\ \iff Q(z) &= q_0(1 - \rho_1 z) \cdots (1 - \rho_m z). \end{aligned}$$

Thus, the roots of  $Q^R$  are the reciprocals of the roots of  $Q$ , and vice versa. We can therefore find the numbers  $\rho_k$  we seek by factoring the reflected polynomial  $Q^R(z)$ .

For example, in the Fibonacci case we have

$$Q(z) = 1 - z - z^2; \quad Q^R(z) = z^2 - z - 1.$$

The roots of  $Q^R$  can be found by setting  $(a, b, c) = (1, -1, -1)$  in the quadratic formula  $(-b \pm \sqrt{b^2 - 4ac})/2a$ ; we find that they are

$$\phi = \frac{1 + \sqrt{5}}{2} \quad \text{and} \quad \hat{\phi} = \frac{1 - \sqrt{5}}{2}.$$

Therefore  $Q^R(z) = (z - \phi)(z - \hat{\phi})$  and  $Q(z) = (1 - \phi z)(1 - \hat{\phi} z)$ .

Once we've found the  $\rho$ 's, we can proceed to find the partial fraction expansion. It's simplest if all the roots are distinct, so let's consider that special case first. We might as well state and prove the general result formally:

**Rational Expansion Theorem for Distinct Roots.**

If  $R(z) = P(z)/Q(z)$ , where  $Q(z) = q_0(1 - \rho_1 z) \dots (1 - \rho_l z)$  and the numbers  $(\rho_1, \dots, \rho_l)$  are distinct, and if  $P(z)$  is a polynomial of degree less than  $l$ , then

$$[z^n] R(z) = a_1 \rho_1^n + \dots + a_l \rho_l^n, \quad \text{where } a_k = \frac{-\rho_k P(1/\rho_k)}{Q'(1/\rho_k)}. \quad (7.29)$$

Proof: Let  $a_1, \dots, a_l$  be the stated constants. Formula (7.29) holds if  $R(z) = P(z)/Q(z)$  is equal to

$$S(z) = \frac{a_1}{1 - \rho_1 z} + \dots + \frac{a_l}{1 - \rho_l z}.$$

And we can prove that  $R(z) = S(z)$  by showing that the function  $T(z) = R(z) - S(z)$  is not infinite as  $z \rightarrow 1/\rho_k$ . For this will show that the rational function  $T(z)$  is never infinite; hence  $T(z)$  must be a polynomial. We also can show that  $T(z) \rightarrow 0$  as  $z \rightarrow \infty$ ; hence  $T(z)$  must be zero.

Let  $\alpha_k = 1/\rho_k$ . To prove that  $\lim_{z \rightarrow \alpha_k} T(z) \neq \infty$ , it suffices to show that  $\lim_{z \rightarrow \alpha_k} (z - \alpha_k)T(z) = 0$ , because  $T(z)$  is a rational function of  $z$ . Thus we want to show that

$$\lim_{z \rightarrow \alpha_k} (z - \alpha_k)R(z) = \lim_{z \rightarrow \alpha_k} (z - \alpha_k)S(z).$$

The right-hand limit equals  $\lim_{z \rightarrow \alpha_k} a_k(z - \alpha_k)/(1 - \rho_k z) = -a_k/\rho_k$ , because  $(1 - \rho_k z) = -\rho_k(z - \alpha_k)$  and  $(z - \alpha_k)/(1 - \rho_j z) \rightarrow 0$  for  $j \neq k$ . The left-hand limit is

$$\lim_{z \rightarrow \alpha_k} (z - \alpha_k) \frac{P(z)}{Q(z)} = P(\alpha_k) \lim_{z \rightarrow \alpha_k} \frac{z - \alpha_k}{Q(z)} = \frac{P(\alpha_k)}{Q'(\alpha_k)},$$

by L'Hospital's rule. Thus the theorem is proved.

Returning to the Fibonacci example, we have  $P(z) = z$  and  $Q(z) = 1 - z - z^2 = (1 - \phi z)(1 - \hat{\phi} z)$ ; hence  $Q'(z) = -1 - 2z$ , and

$$\frac{-\rho P(1/\rho)}{Q'(1/\rho)} = \frac{-1}{-1 - 2/\rho} = \frac{\rho}{\rho + 2}.$$

According to (7.29), the coefficient of  $\phi^n$  in  $[z^n] R(z)$  is therefore  $\phi/(\phi + 2) = 1/\sqrt{5}$ ; the coefficient of  $\hat{\phi}^n$  is  $\hat{\phi}/(\hat{\phi} + 2) = -1/\sqrt{5}$ . So the theorem tells us that  $F_n = (\phi^n - \hat{\phi}^n)/\sqrt{5}$ , as in (6.123).

When  $Q(z)$  has repeated roots, the calculations become more difficult, but we can beef up the proof of the theorem and prove the following more general result:

**General Expansion Theorem for Rational Generating Functions.**

If  $R(z) = P(z)/Q(z)$ , where  $Q(z) = q_0(1 - \rho_1 z)^{d_1} \dots (1 - \rho_l z)^{d_l}$  and the numbers  $(\rho_1, \dots, \rho_l)$  are distinct, and if  $P(z)$  is a polynomial of degree less than  $d_1 + \dots + d_l$ , then

$$[z^n] R(z) = f_1(n)\rho_1^n + \dots + f_l(n)\rho_l^n \quad \text{for all } n \geq 0, \quad (7.30)$$

where each  $f_k(n)$  is a polynomial of degree  $d_k - 1$  with leading coefficient

$$\begin{aligned} a_k &= \frac{(-\rho_k)^{d_k} P(1/\rho_k) d_k}{Q^{(d_k)}(1/\rho_k)} \\ &= \frac{P(1/\rho_k)}{(d_k - 1)! q_0 \prod_{j \neq k} (1 - \rho_j/\rho_k)^{d_j}}. \end{aligned} \quad (7.31)$$

This can be proved by induction on  $\max(d_1, \dots, d_l)$ , using the fact that

$$R(z) - \frac{a_1(d_1 - 1)!}{(1 - \rho_1 z)^{d_1}} - \dots - \frac{a_l(d_l - 1)!}{(1 - \rho_l z)^{d_l}}$$

is a rational function whose denominator polynomial is not divisible by  $(1 - \rho_k z)^{d_k}$  for any  $k$ .

**Example 2: A more-or-less random recurrence.**

Now that we've seen some general methods, we're ready to tackle new problems. Let's try to find a closed form for the recurrence

$$\begin{aligned} g_0 &= g_1 = 1; \\ g_n &= g_{n-1} + 2g_{n-2} + (-1)^n, \quad \text{for } n \geq 2. \end{aligned} \quad (7.32)$$

It's always a good idea to make a table of small cases first, and the recurrence lets us do that easily:

| $n$      | 0 | 1  | 2 | 3  | 4  | 5  | 6  | 7  |
|----------|---|----|---|----|----|----|----|----|
| $(-1)^n$ | 1 | -1 | 1 | -1 | 1  | -1 | 1  | -1 |
| $g_n$    | 1 | 1  | 4 | 5  | 14 | 23 | 52 | 97 |

No closed form is evident, and this sequence isn't even listed in Sloane's *Handbook* [330]; so we need to go through the four-step process if we want to discover the solution.

Step 1 is easy, since we merely need to insert fudge factors to fix things when  $n < 2$ : The equation

$$g_n = g_{n-1} + 2g_{n-2} + (-1)^n[n \geq 0] + [n=1]$$

holds for all integers  $n$ . Now we can carry out Step 2:

$$\begin{aligned} G(z) &= \sum_n g_n z^n = \sum_n g_{n-1} z^n + 2 \sum_n g_{n-2} z^n + \sum_{n \geq 0} (-1)^n z^n + \sum_{n=1} z^n \\ &= zG(z) + 2z^2 G(z) + \frac{1}{1+z} + z. \end{aligned}$$

3.: The upper  
ex on  $\sum_{n=1} z^n$   
ot missing!

(Incidentally, we could also have used  $\binom{-1}{n}$  instead of  $(-1)^n[n \geq 0]$ , thereby getting  $\sum_n \binom{-1}{n} z^n = (1+z)^{-1}$  by the binomial theorem.) Step 3 is elementary algebra, which yields

$$G(z) = \frac{1+z(1+z)}{(1+z)(1-z-2z^2)} = \frac{1+z+z^2}{(1-2z)(1+z)^2}.$$

And that leaves us with Step 4.

The squared factor in the denominator is a bit troublesome, since we know that repeated roots are more complicated than distinct roots; but there it is. We have two roots,  $\rho_1 = 2$  and  $\rho_2 = -1$ ; the general expansion theorem (7.30) tells us that

$$g_n = a_1 2^n + (a_2 n + c)(-1)^n$$

for some constant  $c$ , where

$$a_1 = \frac{1+1/2+1/4}{(1+1/2)^2} = \frac{7}{9}; \quad a_2 = \frac{1-1+1}{1-2/(-1)} = \frac{1}{3}.$$

(The second formula for  $a_k$  in (7.31) is easier to use than the first one when the denominator has nice factors. We simply substitute  $z = 1/\rho_k$  everywhere in  $R(z)$ , except in the factor where this gives zero, and divide by  $(d_k - 1)!$ ; this gives the coefficient of  $n^{d_k-1} \rho_k^n$ .) Plugging in  $n = 0$  tells us that the value of the remaining constant  $c$  had better be  $\frac{2}{9}$ ; hence our answer is

$$g_n = \frac{7}{9} 2^n + \left(\frac{1}{3}n + \frac{2}{9}\right)(-1)^n. \quad (7.33)$$

It doesn't hurt to check the cases  $n = 1$  and  $2$ , just to be sure that we didn't foul up. Maybe we should even try  $n = 3$ , since this formula looks weird. But it's correct, all right.

Could we have discovered (7.33) by guesswork? Perhaps after tabulating a few more values we may have observed that  $g_{n+1} \approx 2g_n$  when  $n$  is large.

And with chutzpah and luck we might even have been able to smoke out the constant  $\frac{7}{9}$ . But it sure is simpler and more reliable to have generating functions as a tool.

**Example 3: Mutually recursive sequences.**

Sometimes we have two or more recurrences that depend on each other. Then we can form generating functions for both of them, and solve both by a simple extension of our four-step method.

For example, let's return to the problem of  $3 \times n$  domino tilings that we explored earlier this chapter. If we want to know only the total number of ways,  $U_n$ , to cover a  $3 \times n$  rectangle with dominoes, without breaking this number down into vertical dominoes versus horizontal dominoes, we needn't go into as much detail as we did before. We can merely set up the recurrences

$$\begin{aligned} U_0 &= 1, & U_1 &= 0; & V_0 &= 0, & V_1 &= 1; \\ U_n &= 2V_{n-1} + U_{n-2}, & V_n &= U_{n-1} + V_{n-2}, & & & & \text{for } n \geq 2. \end{aligned}$$

Here  $V_n$  is the number of ways to cover a  $3 \times n$  rectangle-minus-corner, using  $(3n-1)/2$  dominoes. These recurrences are easy to discover, if we consider the possible domino configurations at the rectangle's left edge, as before. Here are the values of  $U_n$  and  $V_n$  for small  $n$ :

| $n$   | 0 | 1 | 2 | 3 | 4  | 5  | 6  | 7  |
|-------|---|---|---|---|----|----|----|----|
| $U_n$ | 1 | 0 | 3 | 0 | 11 | 0  | 41 | 0  |
| $V_n$ | 0 | 1 | 0 | 4 | 0  | 15 | 0  | 56 |

(7.34)

Let's find closed forms, in four steps. First (Step 1), we have

$$U_n = 2V_{n-1} + U_{n-2} + [n=0], \quad V_n = U_{n-1} + V_{n-2},$$

for all  $n$ . Hence (Step 2),

$$U(z) = 2zV(z) + z^2U(z) + 1, \quad V(z) = zU(z) + z^2V(z).$$

Now (Step 3) we must solve two equations in two unknowns; but these are easy, since the second equation yields  $V(z) = zU(z)/(1-z^2)$ ; we find

$$U(z) = \frac{1-z^2}{1-4z^2+z^4}; \quad V(z) = \frac{z}{1-4z^2+z^4}. \quad (7.35)$$

(We had this formula for  $U(z)$  in (7.10), but with  $z^3$  instead of  $z^2$ . In that derivation,  $n$  was the number of dominoes; now it's the width of the rectangle.)

The denominator  $1-4z^2+z^4$  is a function of  $z^2$ ; this is what makes  $U_{2n+1} = 0$  and  $V_{2n} = 0$ , as they should be. We can take advantage of this



nice property of  $z^2$  by retaining  $z^2$  when we factor the denominator: We need not take  $1 - 4z^2 + z^4$  all the way to a product of four factors  $(1 - \rho_k z)$ , since two factors of the form  $(1 - \rho_k z^2)$  will be enough to tell us the coefficients. In other words if we consider the generating function

$$W(z) = \frac{1}{1 - 4z + z^2} = W_0 + W_1 z + W_2 z^2 + \cdots, \quad (7.36)$$

we will have  $V(z) = zW(z^2)$  and  $U(z) = (1 - z^2)W(z^2)$ ; hence  $V_{2n+1} = W_n$  and  $U_{2n} = W_n - W_{n-1}$ . We save time and energy by working with the simpler function  $W(z)$ .

The factors of  $1 - 4z + z^2$  are  $(z - 2 - \sqrt{3})$  and  $(z - 2 + \sqrt{3})$ , and they can also be written  $(1 - (2 + \sqrt{3})z)$  and  $(1 - (2 - \sqrt{3})z)$  because this polynomial is its own reflection. Thus it turns out that we have

$$\begin{aligned} V_{2n+1} = W_n &= \frac{3+2\sqrt{3}}{6}(2+\sqrt{3})^n + \frac{3-2\sqrt{3}}{6}(2-\sqrt{3})^n; \\ U_{2n} = W_n - W_{n-1} &= \frac{3+\sqrt{3}}{6}(2+\sqrt{3})^n + \frac{3-\sqrt{3}}{6}(2-\sqrt{3})^n \\ &= \frac{(2+\sqrt{3})^n}{3-\sqrt{3}} + \frac{(2-\sqrt{3})^n}{3+\sqrt{3}}. \end{aligned} \quad (7.37)$$

This is the desired closed form for the number of  $3 \times n$  domino tilings.

Incidentally, we can simplify the formula for  $U_{2n}$  by realizing that the second term always lies between 0 and 1. The number  $U_{2n}$  is an integer, so we have

$$U_{2n} = \left\lceil \frac{(2+\sqrt{3})^n}{3-\sqrt{3}} \right\rceil, \quad \text{for } n \geq 0. \quad (7.38)$$

In fact, the other term  $(2 - \sqrt{3})^n / (3 + \sqrt{3})$  is extremely small when  $n$  is large, because  $2 - \sqrt{3} \approx 0.268$ . This needs to be taken into account if we try to use formula (7.38) in numerical calculations. For example, a fairly expensive name-brand hand calculator comes up with 413403.0005 when asked to compute  $(2 + \sqrt{3})^{10} / (3 - \sqrt{3})$ . This is correct to nine significant figures; but the true value is slightly *less* than 413403, not slightly greater. Therefore it would be a mistake to take the ceiling of 413403.0005; the correct answer,  $U_{20} = 413403$ , is obtained by *rounding* to the nearest integer. Ceilings can be hazardous.

**Example 4: A closed form for change.**

When we left the problem of making change, we had just calculated the number of ways to pay 50¢. Let's try now to count the number of ways there are to change a dollar, or a million dollars — still using only pennies, nickels, dimes, quarters, and halves.

known slippery  
ers too.

The generating function derived earlier is

$$C(z) = \frac{1}{1-z} \frac{1}{1-z^5} \frac{1}{1-z^{10}} \frac{1}{1-z^{25}} \frac{1}{1-z^{50}};$$

this is a rational function of  $z$  with a denominator of degree 91. Therefore we can decompose the denominator into 91 factors and come up with a 91-term “closed form” for  $C_n$ , the number of ways to give  $n$  cents in change. But that’s too horrible to contemplate. Can’t we do better than the general method suggests, in this particular case?

One ray of hope suggests itself immediately, when we notice that the denominator is almost a function of  $z^5$ . The trick we just used to simplify the calculations by noting that  $1 - 4z^2 + z^4$  is a function of  $z^2$  can be applied to  $C(z)$ , if we replace  $1/(1-z)$  by  $(1+z+z^2+z^3+z^4)/(1-z^5)$ :

$$\begin{aligned} C(z) &= \frac{1+z+z^2+z^3+z^4}{1-z^5} \frac{1}{1-z^5} \frac{1}{1-z^{10}} \frac{1}{1-z^{25}} \frac{1}{1-z^{50}} \\ &= (1+z+z^2+z^3+z^4)\check{C}(z^5), \\ \check{C}(z) &= \frac{1}{1-z} \frac{1}{1-z} \frac{1}{1-z^2} \frac{1}{1-z^5} \frac{1}{1-z^{10}}. \end{aligned}$$

The compressed function  $\check{C}(z)$  has a denominator whose degree is only 19, so it’s much more tractable than the original. This new expression for  $C(z)$  shows us, incidentally, that  $C_{5n} = C_{5n+1} = C_{5n+2} = C_{5n+3} = C_{5n+4}$ ; and indeed, this set of equations is obvious in retrospect: The number of ways to leave a 53¢ tip is the same as the number of ways to leave a 50¢ tip, because the number of pennies is predetermined modulo 5.

But  $\check{C}(z)$  still doesn’t have a really simple closed form based on the roots of the denominator. The easiest way to compute the coefficients of  $\check{C}(z)$  is probably to recognize that each of the denominator factors is a divisor of  $1 - z^{10}$ . Hence we can write

*Now we’re also getting compressed reasoning.*

$$\check{C}(z) = \frac{A(z)}{(1-z^{10})^5}, \quad \text{where } A(z) = A_0 + A_1z + \cdots + A_{31}z^{31}. \quad (7.39)$$

The actual value of  $A(z)$ , for the curious, is

$$\begin{aligned} &(1+z+\cdots+z^9)^2(1+z^2+\cdots+z^8)(1+z^5) \\ &= 1 + 2z + 4z^2 + 6z^3 + 9z^4 + 13z^5 + 18z^6 + 24z^7 \\ &\quad + 31z^8 + 39z^9 + 45z^{10} + 52z^{11} + 57z^{12} + 63z^{13} + 67z^{14} + 69z^{15} \\ &\quad + 69z^{16} + 67z^{17} + 63z^{18} + 57z^{19} + 52z^{20} + 45z^{21} + 39z^{22} + 31z^{23} \\ &\quad + 24z^{24} + 18z^{25} + 13z^{26} + 9z^{27} + 6z^{28} + 4z^{29} + 2z^{30} + z^{31}. \end{aligned}$$

Finally, since  $1/(1-z^{10})^5 = \sum_{k \geq 0} \binom{k+4}{4} z^{10k}$ , we can determine the coefficient  $\check{C}_n = [z^n] \check{C}(z)$  as follows, when  $n = 10q + r$  and  $0 \leq r < 10$ :

$$\begin{aligned} \check{C}_{10q+r} &= \sum_{j,k} A_j \binom{k+4}{4} [10q + r = 10k + j] \\ &= A_r \binom{q+4}{4} + A_{r+10} \binom{q+3}{4} + A_{r+20} \binom{q+2}{4} + A_{r+30} \binom{q+1}{4}. \end{aligned} \quad (7.40)$$

This gives ten cases, one for each value of  $r$ ; but it's a pretty good closed form, compared with alternatives that involve powers of complex numbers.

For example, we can use this expression to deduce the value of  $C_{50q} = \check{C}_{10q}$ . Then  $r = 0$  and we have

$$C_{50q} = \binom{q+4}{4} + 45 \binom{q+3}{4} + 52 \binom{q+2}{4} + 2 \binom{q+1}{4}.$$

The number of ways to change 50¢ is  $\binom{5}{4} + 45 \binom{4}{4} = 50$ ; the number of ways to change \$1 is  $\binom{6}{4} + 45 \binom{5}{4} + 52 \binom{4}{4} = 292$ ; and the number of ways to change \$1,000,000 is

$$\begin{aligned} \binom{2000004}{4} + 45 \binom{2000003}{4} + 52 \binom{2000002}{4} + 2 \binom{2000001}{4} \\ = 66666793333412666685000001. \end{aligned}$$

**Example 5: A divergent series.**

Now let's try to get a closed form for the numbers  $g_n$  defined by

$$\begin{aligned} g_0 &= 1; \\ g_n &= ng_{n-1}, \quad \text{for } n > 0. \end{aligned}$$

After staring at this for a few nanoseconds we realize that  $g_n$  is just  $n!$ ; in fact, the method of summation factors described in Chapter 2 suggests this answer immediately. But let's try to solve the recurrence with generating functions, just to see what happens. (A powerful technique should be able to handle easy recurrences like this, as well as others that have answers we can't guess so easily.)

The equation

$$g_n = ng_{n-1} + [n=0]$$

holds for all  $n$ , and it leads to

$$G(z) = \sum_n g_n z^n = \sum_n ng_{n-1} z^n + \sum_{n=0} z^n.$$

To complete Step 2, we want to express  $\sum_n ng_{n-1} z^n$  in terms of  $G(z)$ , and the basic maneuvers in Table 335 suggest that the derivative  $G'(z) =$

everyday people  
are talking  
into seconds.

$\sum_n n g_n z^{n-1}$  is somehow involved. So we steer toward that kind of sum:

$$\begin{aligned} G(z) &= 1 + \sum_n (n+1) g_n z^{n+1} \\ &= 1 + \sum_n n g_n z^{n+1} + \sum_n g_n z^{n+1} \\ &= 1 + z^2 G'(z) + zG(z). \end{aligned}$$

Let's check this equation, using the values of  $g_n$  for small  $n$ . Since

$$\begin{aligned} G &= 1 + z + 2z^2 + 6z^3 + 24z^4 + \cdots, \\ G' &= 1 + 4z + 18z^2 + 96z^3 + \cdots, \end{aligned}$$

we have

$$\begin{aligned} z^2 G' &= z^2 + 4z^3 + 18z^4 + 96z^5 + \cdots, \\ zG &= z + z^2 + 2z^3 + 6z^4 + 24z^5 + \cdots, \\ 1 &= 1. \end{aligned}$$

These three lines add up to  $G$ , so we're fine so far. Incidentally, we often find it convenient to write ' $G$ ' instead of ' $G(z)$ '; the extra ' $(z)$ ' just clutters up the formula when we aren't changing  $z$ .

Step 3 is next, and it's different from what we've done before because we have a differential equation to solve. But this is a differential equation that we can handle with the hypergeometric series techniques of Section 5.6; those techniques aren't too bad. (Readers who are unfamiliar with hypergeometrics needn't worry—this will be quick.)

First we must get rid of the constant '1', so we take the derivative of both sides:

$$\begin{aligned} G' &= (z^2 G' + zG + 1)' = (2zG' + z^2 G'') + (G + zG') \\ &= z^2 G'' + 3zG' + G. \end{aligned}$$

The theory in Chapter 5 tells us to rewrite this using the  $\partial$  operator, and we know from exercise 6.13 that

$$\partial G = zG', \quad \partial^2 G = z^2 G'' + zG'.$$

Therefore the desired form of the differential equation is

$$\partial G = z\partial^2 G + 2z\partial G + zG = z(\partial + 1)^2 G.$$

According to (5.109), the solution with  $g_0 = 1$  is the hypergeometric series  $F(1, 1; ; z)$ .

*"This will be quick."  
That's what the  
doctor said just  
before he stuck me  
with that needle.  
Come to think of it,  
"hypergeometric"  
sounds a lot like  
"hypodermic."*

Step 3 was more than we bargained for; but now that we know what the function  $G$  is, Step 4 is easy—the hypergeometric definition (5.76) gives us the power series expansion:

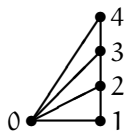
$$G(z) = F\left(\begin{matrix} 1, 1 \\ \end{matrix} \middle| z\right) = \sum_{n \geq 0} \frac{1^{\overline{n}} 1^{\overline{n}} z^n}{n!} = \sum_{n \geq 0} n! z^n.$$

We've confirmed the closed form we knew all along,  $g_n = n!$ .

Notice that the technique gave the right answer even though  $G(z)$  diverges for all nonzero  $z$ . The sequence  $n!$  grows so fast, the terms  $|n! z^n|$  approach  $\infty$  as  $n \rightarrow \infty$ , unless  $z = 0$ . This shows that formal power series can be manipulated algebraically without worrying about convergence.

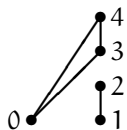
**Example 6: A recurrence that goes all the way back.**

Let's close this section by applying generating functions to a problem in graph theory. A *fan* of order  $n$  is a graph on the vertices  $\{0, 1, \dots, n\}$  with  $2n - 1$  edges defined as follows: Vertex 0 is connected by an edge to each of the other  $n$  vertices, and vertex  $k$  is connected by an edge to vertex  $k + 1$ , for  $1 \leq k < n$ . Here, for example, is the fan of order 4, which has five vertices and seven edges.



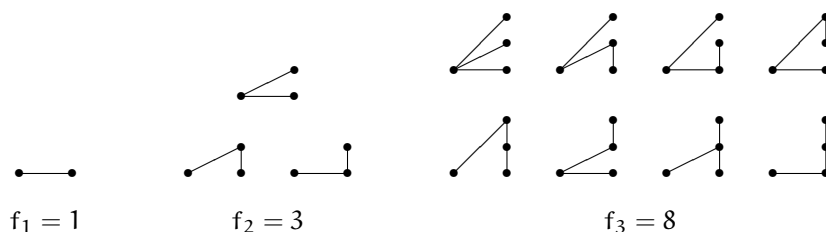
The problem of interest: How many spanning trees  $f_n$  are in such a graph? A *spanning tree* is a subgraph containing all the vertices, and containing enough edges to make the subgraph connected yet not so many that it has a cycle. It turns out that every spanning tree of a graph on  $n + 1$  vertices has exactly  $n$  edges. With fewer than  $n$  edges the subgraph wouldn't be connected, and with more than  $n$  it would have a cycle; graph theory books prove this.

There are  $\binom{2n-1}{n}$  ways to choose  $n$  edges from among the  $2n - 1$  present in a fan of order  $n$ , but these choices don't always yield a spanning tree. For instance the subgraph



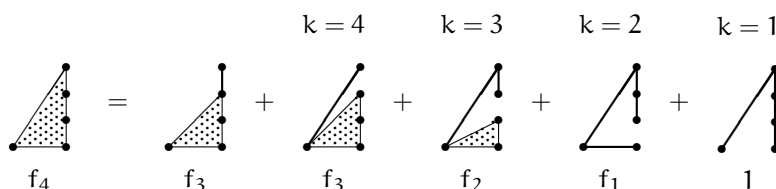
has four edges but is not a spanning tree; it has a cycle from 0 to 4 to 3 to 0, and it has no connection between  $\{1, 2\}$  and the other vertices. We want to count how many of the  $\binom{2n-1}{n}$  choices actually do yield spanning trees.

Let's look at some small cases. It's pretty easy to enumerate the spanning trees for  $n = 1, 2$ , and 3:



(We need not show the labels on the vertices, if we always draw vertex 0 at the left.) What about the case  $n = 0$ ? At first it seems reasonable to set  $f_0 = 1$ ; but we'll take  $f_0 = 0$ , because the existence of a fan of order 0 (which should have  $2n - 1 = -1$  edges) is dubious.

Our four-step procedure tells us to find a recurrence for  $f_n$  that holds for all  $n$ . We can get a recurrence by observing how the topmost vertex (vertex  $n$ ) is connected to the rest of the spanning tree. If it's not connected to vertex 0, it must be connected to vertex  $n - 1$ , since it must be connected to the rest of the graph. In this case, any of the  $f_{n-1}$  spanning trees for the remaining fan (on the vertices 0 through  $n - 1$ ) will complete a spanning tree for the whole graph. Otherwise vertex  $n$  is connected to 0, and there's some number  $k \leq n$  such that vertices  $n, n - 1, \dots, k$  are connected directly but the edge between  $k$  and  $k - 1$  is not present. Then there can't be any edges between 0 and  $\{n - 1, \dots, k\}$ , or there would be a cycle. If  $k = 1$ , the spanning tree is therefore determined completely. And if  $k > 1$ , any of the  $f_{k-1}$  ways to produce a spanning tree on  $\{0, 1, \dots, k - 1\}$  will yield a spanning tree on the whole graph. For example, here's what this analysis produces when  $n = 4$ :



The general equation, valid for  $n \geq 1$ , is

$$f_n = f_{n-1} + f_{n-1} + f_{n-2} + f_{n-3} + \cdots + f_1 + 1.$$

(It almost seems as though the '1' on the end is  $f_0$  and we should have chosen  $f_0 = 1$ ; but we will doggedly stick with our choice.) A few changes suffice to make the equation valid for all integers  $n$ :

$$f_n = f_{n-1} + \sum_{k < n} f_k + [n > 0]. \quad (7.41)$$

This is a recurrence that “goes all the way back” from  $f_{n-1}$  through all previous values, so it’s different from the other recurrences we’ve seen so far in this chapter. We used a special method to get rid of a similar right-side sum in Chapter 2, when we solved the quicksort recurrence (2.12); namely, we subtracted one instance of the recurrence from another ( $f_{n+1} - f_n$ ). This trick would get rid of the  $\sum$  now, as it did then; but we’ll see that generating functions allow us to work directly with such sums. (And it’s a good thing that they do, because we will be seeing much more complicated recurrences before long.)

Step 1 is finished; Step 2 is where we need to do a new thing:

$$\begin{aligned} F(z) &= \sum_n f_n z^n = \sum_n f_{n-1} z^n + \sum_{k,n} f_k z^n [k < n] + \sum_n [n > 0] z^n \\ &= zF(z) + \sum_k f_k z^k \sum_n [n > k] z^{n-k} + \frac{z}{1-z} \\ &= zF(z) + F(z) \sum_{m>0} z^m + \frac{z}{1-z} \\ &= zF(z) + F(z) \frac{z}{1-z} + \frac{z}{1-z}. \end{aligned}$$

The key trick here was to change  $z^n$  to  $z^k z^{n-k}$ ; this made it possible to express the value of the double sum in terms of  $F(z)$ , as required in Step 2.

Now Step 3 is simple algebra, and we find

$$F(z) = \frac{z}{1-3z+z^2}.$$

Those of us with a zest for memorization will recognize this as the generating function (7.24) for the even-numbered Fibonacci numbers. So, we needn’t go through Step 4; we have found a somewhat surprising answer to the spans-of-fans problem:

$$f_n = F_{2n}, \quad \text{for } n \geq 0. \quad (7.42)$$

## 7.4 SPECIAL GENERATING FUNCTIONS

Step 4 of the four-step procedure becomes much easier if we know the coefficients of lots of different power series. The expansions in Table 336 are quite useful, as far as they go, but many other types of closed forms are possible. Therefore we ought to supplement that table with another one, which lists power series that correspond to the “special numbers” considered in Chapter 6.

**Table 352** Generating functions for special numbers.

$$\frac{1}{(1-z)^{m+1}} \ln \frac{1}{1-z} = \sum_{n \geq 0} (H_{m+n} - H_m) \binom{m+n}{n} z^n \quad (7.43)$$

$$\frac{z}{e^z - 1} = \sum_{n \geq 0} B_n \frac{z^n}{n!} \quad (7.44)$$

$$\frac{F_m z}{1 - (F_{m-1} + F_{m+1})z + (-1)^m z^2} = \sum_{n \geq 0} F_{mn} z^n \quad (7.45)$$

$$\sum_k \left\{ \begin{matrix} m \\ k \end{matrix} \right\} \frac{k! z^k}{(1-z)^{k+1}} = \sum_{n \geq 0} n^m z^n \quad (7.46)$$

$$(z^{-1})^{\overline{-m}} = \frac{z^m}{(1-z)(1-2z)\dots(1-mz)} = \sum_{n \geq 0} \left\{ \begin{matrix} n \\ m \end{matrix} \right\} z^n \quad (7.47)$$

$$z^{\overline{m}} = z(z+1)\dots(z+m-1) = \sum_{n \geq 0} \left[ \begin{matrix} m \\ n \end{matrix} \right] z^n \quad (7.48)$$

$$(e^z - 1)^m = m! \sum_{n \geq 0} \left\{ \begin{matrix} n \\ m \end{matrix} \right\} \frac{z^n}{n!} \quad (7.49)$$

$$\left( \ln \frac{1}{1-z} \right)^m = m! \sum_{n \geq 0} \left[ \begin{matrix} n \\ m \end{matrix} \right] \frac{z^n}{n!} \quad (7.50)$$

$$\left( \frac{z}{\ln(1+z)} \right)^m = \sum_{n \geq 0} \frac{z^n}{n!} \left\{ \begin{matrix} m \\ m-n \end{matrix} \right\} / \binom{m-1}{n} \quad (7.51)$$

$$\left( \frac{z}{1-e^{-z}} \right)^m = \sum_{n \geq 0} \frac{z^n}{n!} \left[ \begin{matrix} m \\ m-n \end{matrix} \right] / \binom{m-1}{n} \quad (7.52)$$

$$e^{z+wz} = \sum_{m,n \geq 0} \binom{n}{m} w^m \frac{z^n}{n!} \quad (7.53)$$

$$e^{w(e^z-1)} = \sum_{m,n \geq 0} \left\{ \begin{matrix} n \\ m \end{matrix} \right\} w^m \frac{z^n}{n!} \quad (7.54)$$

$$\frac{1}{(1-z)^w} = \sum_{m,n \geq 0} \left[ \begin{matrix} n \\ m \end{matrix} \right] w^m \frac{z^n}{n!} \quad (7.55)$$

$$\frac{1-w}{e^{(w-1)z} - w} = \sum_{m,n \geq 0} \left\langle \begin{matrix} n \\ m \end{matrix} \right\rangle w^m \frac{z^n}{n!} \quad (7.56)$$



Table 352 is the database we need. The identities in this table are not difficult to prove, so we needn't dwell on them; this table is primarily for reference when we meet a new problem. But there's a nice proof of the first formula, (7.43), that deserves mention: We start with the identity

$$\frac{1}{(1-z)^{x+1}} = \sum_n \binom{x+n}{n} z^n$$

and differentiate it with respect to  $x$ . On the left,  $(1-z)^{-x-1}$  is equal to  $e^{(x+1)\ln(1/(1-z))}$ , so  $d/dx$  contributes a factor of  $\ln(1/(1-z))$ . On the right, the numerator of  $\binom{x+n}{n}$  is  $(x+n)\dots(x+1)$ , and  $d/dx$  splits this into  $n$  terms whose sum is equivalent to multiplying  $\binom{x+n}{n}$  by

$$\frac{1}{x+n} + \dots + \frac{1}{x+1} = H_{x+n} - H_x.$$

Replacing  $x$  by  $m$  gives (7.43). Notice that  $H_{x+n} - H_x$  is meaningful even when  $x$  is not an integer.

By the way, this method of differentiating a complicated product — leaving it as a product — is usually better than expressing the derivative as a sum. For example the right side of

$$\frac{d}{dx}((x+n)^n \dots (x+1)^1) = (x+n)^n \dots (x+1)^1 \left( \frac{n}{x+n} + \dots + \frac{1}{x+1} \right)$$

would be a lot messier written out as a sum.

The general identities in Table 352 include many important special cases. For example, (7.43) simplifies to the generating function for  $H_n$  when  $m = 0$ :

$$\frac{1}{1-z} \ln \frac{1}{1-z} = \sum_n H_n z^n. \quad (7.57)$$

This equation can also be derived in other ways; for example, we can take the power series for  $\ln(1/(1-z))$  and divide it by  $1-z$  to get cumulative sums.

Identities (7.51) and (7.52) involve the respective ratios  $\left\{ \begin{smallmatrix} m \\ m-n \end{smallmatrix} \right\} / \binom{m-1}{n}$  and  $\left[ \begin{smallmatrix} m \\ m-n \end{smallmatrix} \right] / \binom{m-1}{n}$ , which have the undefined form  $0/0$  when  $n \geq m$ . However, there is a way to give them a proper meaning using the Stirling polynomials of (6.45), because we have

$$\left\{ \begin{smallmatrix} m \\ m-n \end{smallmatrix} \right\} / \binom{m-1}{n} = (-1)^{n+1} n! m \sigma_n(n-m); \quad (7.58)$$

$$\left[ \begin{smallmatrix} m \\ m-n \end{smallmatrix} \right] / \binom{m-1}{n} = n! m \sigma_n(m). \quad (7.59)$$

Thus, for example, the case  $m = 1$  of (7.51) should not be regarded as the power series  $\sum_{n \geq 0} (z^n/n!) \left\{ \begin{smallmatrix} 1 \\ 1-n \end{smallmatrix} \right\} / \binom{0}{n}$ , but rather as

$$\frac{z}{\ln(1+z)} = - \sum_{n \geq 0} (-z)^n \sigma_n(n-1) = 1 + \frac{1}{2}z - \frac{1}{12}z^2 + \dots.$$

Identities (7.53), (7.55), (7.54), and (7.56) are “double generating functions” or “super generating functions” because they have the form  $G(w, z) = \sum_{m,n} g_{m,n} w^m z^n$ . The coefficient of  $w^m$  is a generating function in the variable  $z$ ; the coefficient of  $z^n$  is a generating function in the variable  $w$ . Equation (7.56) can be put into the more symmetrical form

$$\frac{e^w - e^z}{we^z - ze^w} = \sum_{m,n} \left\langle \begin{matrix} m+n+1 \\ m \end{matrix} \right\rangle \frac{w^m z^n}{(m+n+1)!}. \quad (7.60)$$

## 7.5 CONVOLUTIONS

The *convolution* of two given sequences  $\langle f_0, f_1, \dots \rangle = \langle f_n \rangle$  and  $\langle g_0, g_1, \dots \rangle = \langle g_n \rangle$  is the sequence  $\langle f_0 g_0, f_0 g_1 + f_1 g_0, \dots \rangle = \langle \sum_k f_k g_{n-k} \rangle$ . We have observed in Sections 5.4 and 7.2 that convolution of sequences corresponds to multiplication of their generating functions. This fact makes it easy to evaluate many sums that would otherwise be difficult to handle.

*I always thought convolution was what happens to my brain when I try to do a proof.*

### Example 1: A Fibonacci convolution.

For example, let's try to evaluate  $\sum_{k=0}^n F_k F_{n-k}$  in closed form. This is the convolution of  $\langle F_n \rangle$  with itself, so the sum must be the coefficient of  $z^n$  in  $F(z)^2$ , where  $F(z)$  is the generating function for  $\langle F_n \rangle$ . All we have to do is figure out the value of this coefficient.

The generating function  $F(z)$  is  $z/(1-z-z^2)$ , a quotient of polynomials; so the general expansion theorem for rational functions tells us that the answer can be obtained from a partial fraction representation. We can use the general expansion theorem (7.30) and grind away; or we can use the fact that

$$\begin{aligned} F(z)^2 &= \left( \frac{1}{\sqrt{5}} \left( \frac{1}{1-\phi z} - \frac{1}{1-\hat{\phi} z} \right) \right)^2 \\ &= \frac{1}{5} \left( \frac{1}{(1-\phi z)^2} - \frac{2}{(1-\phi z)(1-\hat{\phi} z)} + \frac{1}{(1-\hat{\phi} z)^2} \right) \\ &= \frac{1}{5} \sum_{n \geq 0} (n+1) \phi^n z^n - \frac{2}{5} \sum_{n \geq 0} F_{n+1} z^n + \frac{1}{5} \sum_{n \geq 0} (n+1) \hat{\phi}^n z^n. \end{aligned}$$

Instead of expressing the answer in terms of  $\phi$  and  $\hat{\phi}$ , let's try for a closed form in terms of Fibonacci numbers. Recalling that  $\phi + \hat{\phi} = 1$ , we have

$$\begin{aligned} \phi^n + \hat{\phi}^n &= [z^n] \left( \frac{1}{1-\phi z} + \frac{1}{1-\hat{\phi} z} \right) \\ &= [z^n] \frac{2 - (\phi + \hat{\phi})z}{(1-\phi z)(1-\hat{\phi} z)} = [z^n] \frac{2-z}{1-z-z^2} = 2F_{n+1} - F_n. \end{aligned}$$

Hence

$$F(z)^2 = \frac{1}{5} \sum_{n \geq 0} (n+1)(2F_{n+1} - F_n)z^n - \frac{2}{5} \sum_{n \geq 0} F_{n+1} z^n,$$

and we have the answer we seek:

$$\sum_{k=0}^n F_k F_{n-k} = \frac{2nF_{n+1} - (n+1)F_n}{5}. \quad (7.61)$$

For example, when  $n = 3$  this formula gives  $F_0F_3 + F_1F_2 + F_2F_1 + F_3F_0 = 0 + 1 + 1 + 0 = 2$  on the left and  $(6F_4 - 4F_3)/5 = (18 - 8)/5 = 2$  on the right.

**Example 2: Harmonic convolutions.**

The efficiency of a certain computer method called “samplesort” depends on the value of the sum

$$T_{m,n} = \sum_{0 \leq k < n} \binom{k}{m} \frac{1}{n-k}, \quad \text{integers } m, n \geq 0.$$

Exercise 5.58 obtains the value of this sum by a somewhat intricate double induction, using summation factors. It's much easier to realize that  $T_{m,n}$  is just the  $n$ th term in the convolution of  $\langle \binom{0}{m}, \binom{1}{m}, \binom{2}{m}, \dots \rangle$  with  $\langle 0, \frac{1}{1}, \frac{1}{2}, \dots \rangle$ . Both sequences have simple generating functions in Table 336:

$$\sum_{n \geq 0} \binom{n}{m} z^n = \frac{z^m}{(1-z)^{m+1}}; \quad \sum_{n > 0} \frac{z^n}{n} = \ln \frac{1}{1-z}.$$

Therefore, by (7.43),

$$\begin{aligned} T_{m,n} &= [z^n] \frac{z^m}{(1-z)^{m+1}} \ln \frac{1}{1-z} = [z^{n-m}] \frac{1}{(1-z)^{m+1}} \ln \frac{1}{1-z} \\ &= (H_n - H_m) \binom{n}{n-m}. \end{aligned}$$

In fact, there are many more sums that boil down to this same sort of convolution, because we have

$$\frac{1}{(1-z)^{r+1}} \ln \frac{1}{1-z} \cdot \frac{1}{(1-z)^{s+1}} = \frac{1}{(1-z)^{r+s+2}} \ln \frac{1}{1-z}$$

for all  $r$  and  $s$ . Equating coefficients of  $z^n$  gives the general identity

$$\begin{aligned} \sum_k \binom{r+k}{k} \binom{s+n-k}{n-k} (H_{r+k} - H_r) \\ = \binom{r+s+n+1}{n} (H_{r+s+n+1} - H_{r+s+1}). \end{aligned} \quad (7.62)$$

This seems almost too good to be true. But it checks, at least when  $n = 2$ :

*Because it's so harmonic.*

$$\begin{aligned} \binom{r+1}{1} \binom{s+1}{1} \frac{1}{r+1} + \binom{r+2}{2} \binom{s+0}{0} \left( \frac{1}{r+2} + \frac{1}{r+1} \right) \\ = \binom{r+s+3}{2} \left( \frac{1}{r+s+3} + \frac{1}{r+s+2} \right). \end{aligned}$$

Special cases like  $s = 0$  are as remarkable as the general case.

And there's more. We can use the convolution identity

$$\sum_k \binom{r+k}{k} \binom{s+n-k}{n-k} = \binom{r+s+n+1}{n}$$

to transpose  $H_r$  to the other side, since  $H_r$  is independent of  $k$ :

$$\begin{aligned} \sum_k \binom{r+k}{k} \binom{s+n-k}{n-k} H_{r+k} \\ = \binom{r+s+n+1}{n} (H_{r+s+n+1} - H_{r+s+1} + H_r). \end{aligned} \quad (7.63)$$

There's still more: If  $r$  and  $s$  are nonnegative integers  $l$  and  $m$ , we can replace  $\binom{r+k}{k}$  by  $\binom{l+k}{l}$  and  $\binom{s+n-k}{n-k}$  by  $\binom{m+n-k}{m}$ ; then we can change  $k$  to  $k-l$  and  $n$  to  $n-m-l$ , getting

$$\sum_{k=0}^n \binom{k}{l} \binom{n-k}{m} H_k = \binom{n+1}{l+m+1} (H_{n+1} - H_{l+m+1} + H_l),$$

integers  $l, m, n \geq 0$ . (7.64)

Even the special case  $l = m = 0$  of this identity was difficult for us to handle in Chapter 2! (See (2.36).) We've come a long way.

**Example 3: Convolutions of convolutions.**

If we form the convolution of  $\langle f_n \rangle$  and  $\langle g_n \rangle$ , then convolve this with a third sequence  $\langle h_n \rangle$ , we get a sequence whose  $n$ th term is

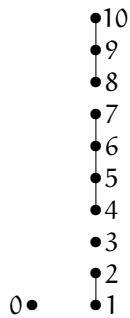
$$\sum_{j+k+l=n} f_j g_k h_l.$$

The generating function of this three-fold convolution is, of course, the three-fold product  $F(z)G(z)H(z)$ . In a similar way, the  $m$ -fold convolution of a sequence  $\langle g_n \rangle$  with itself has  $n$ th term equal to

$$\sum_{k_1+k_2+\dots+k_m=n} g_{k_1} g_{k_2} \dots g_{k_m}$$

and its generating function is  $G(z)^m$ .

We can apply these observations to the spans-of-fans problem considered earlier (Example 6 in Section 7.3). It turns out that there's another way to compute  $f_n$ , the number of spanning trees of an  $n$ -fan, based on the configurations of tree edges between the vertices  $\{1, 2, \dots, n\}$ : The edge between vertex  $k$  and vertex  $k+1$  may or may not be selected for the tree; and each of the ways to select these edges connects up certain blocks of adjacent vertices. For example, when  $n = 10$  we might connect vertices  $\{1, 2\}$ ,  $\{3\}$ ,  $\{4, 5, 6, 7\}$ , and  $\{8, 9, 10\}$ :



How many spanning trees can we make, by adding additional edges to vertex 0? We need to connect 0 to each of the four blocks; and there are two ways to join 0 with  $\{1, 2\}$ , one way to join it with  $\{3\}$ , four ways with  $\{4, 5, 6, 7\}$ , and three ways with  $\{8, 9, 10\}$ , or  $2 \cdot 1 \cdot 4 \cdot 3 = 24$  ways altogether. Summing over all possible ways to make blocks gives us the following expression for the total number of spanning trees:

$$f_n = \sum_{m \geq 0} \sum_{\substack{k_1 + k_2 + \dots + k_m = n \\ k_1, k_2, \dots, k_m \geq 0}} k_1 k_2 \dots k_m. \quad (7.65)$$

For example,  $f_4 = 4 + 3 \cdot 1 + 2 \cdot 2 + 1 \cdot 3 + 2 \cdot 1 \cdot 1 + 1 \cdot 2 \cdot 1 + 1 \cdot 1 \cdot 2 + 1 \cdot 1 \cdot 1 \cdot 1 = 21$ .

This is the sum of  $m$ -fold convolutions of the sequence  $\langle 0, 1, 2, 3, \dots \rangle$ , for  $m = 1, 2, 3, \dots$ ; hence the generating function for  $\langle f_n \rangle$  is

$$F(z) = G(z) + G(z)^2 + G(z)^3 + \dots = \frac{G(z)}{1 - G(z)}$$

where  $G(z)$  is the generating function for  $\langle 0, 1, 2, 3, \dots \rangle$ , namely  $z/(1 - z)^2$ . Consequently we have

$$F(z) = \frac{z}{(1 - z)^2 - z} = \frac{z}{1 - 3z + z^2},$$

as before. This approach to  $\langle f_n \rangle$  is more symmetrical and appealing than the complicated recurrence we had earlier.

**Example 4: A convoluted recurrence.**

Our next example is especially important. In fact, it's the "classic example" of why generating functions are useful in the solution of recurrences.

Suppose we have  $n + 1$  variables  $x_0, x_1, \dots, x_n$  whose product is to be computed by doing  $n$  multiplications. How many ways  $C_n$  are there to insert parentheses into the product  $x_0 \cdot x_1 \cdot \dots \cdot x_n$  so that the order of multiplication is completely specified? For example, when  $n = 2$  there are two ways,  $x_0 \cdot (x_1 \cdot x_2)$  and  $(x_0 \cdot x_1) \cdot x_2$ . And when  $n = 3$  there are five ways,

$$\begin{aligned} &x_0 \cdot (x_1 \cdot (x_2 \cdot x_3)), \quad x_0 \cdot ((x_1 \cdot x_2) \cdot x_3), \quad (x_0 \cdot x_1) \cdot (x_2 \cdot x_3), \\ &\quad (x_0 \cdot (x_1 \cdot x_2)) \cdot x_3, \quad ((x_0 \cdot x_1) \cdot x_2) \cdot x_3. \end{aligned}$$

Thus  $C_2 = 2$ ,  $C_3 = 5$ ; we also have  $C_1 = 1$  and  $C_0 = 1$ .

Let's use the four-step procedure of Section 7.3. What is a recurrence for the  $C$ 's? The key observation is that there's exactly one ' $\cdot$ ' operation outside all of the parentheses, when  $n > 0$ ; this is the final multiplication that ties everything together. If this ' $\cdot$ ' occurs between  $x_k$  and  $x_{k+1}$ , there are  $C_k$  ways to fully parenthesize  $x_0 \cdot \dots \cdot x_k$ , and there are  $C_{n-k-1}$  ways to fully parenthesize  $x_{k+1} \cdot \dots \cdot x_n$ ; hence

$$C_n = C_0 C_{n-1} + C_1 C_{n-2} + \dots + C_{n-1} C_0, \quad \text{if } n > 0.$$

By now we recognize this expression as a convolution, and we know how to patch the formula so that it holds for all integers  $n$ :

$$C_n = \sum_k C_k C_{n-1-k} + [n=0]. \quad (7.66)$$

Step 1 is now complete. Step 2 tells us to multiply by  $z^n$  and sum:

$$\begin{aligned} C(z) &= \sum_n C_n z^n \\ &= \sum_{k,n} C_k C_{n-1-k} z^n + \sum_{n=0} z^n \\ &= \sum_k C_k z^k \sum_n C_{n-1-k} z^{n-k} + 1 \\ &= C(z) \cdot zC(z) + 1. \end{aligned}$$

Lo and behold, the convolution has become a product, in the generating-function world. Life is full of surprises.

*The authors jest.*

Step 3 is also easy. We solve for  $C(z)$  by the quadratic formula:

$$C(z) = \frac{1 \pm \sqrt{1-4z}}{2z}.$$

But should we choose the  $+$  sign or the  $-$  sign? Both choices yield a function that satisfies  $C(z) = zC(z)^2 + 1$ , but only one of the choices is suitable for our problem. We might choose the  $+$  sign on the grounds that positive thinking is best; but we soon discover that this choice gives  $C(0) = \infty$ , contrary to the facts. (The correct function  $C(z)$  is supposed to have  $C(0) = C_0 = 1$ .) Therefore we conclude that

$$C(z) = \frac{1 - \sqrt{1-4z}}{2z}.$$

Finally, Step 4. What is  $[z^n]C(z)$ ? The binomial theorem tells us that

$$\sqrt{1-4z} = \sum_{k \geq 0} \binom{1/2}{k} (-4z)^k = 1 + \sum_{k \geq 1} \frac{1}{2k} \binom{-1/2}{k-1} (-4z)^k;$$

hence, using (5.37),

$$\begin{aligned} \frac{1 - \sqrt{1-4z}}{2z} &= \sum_{k \geq 1} \frac{1}{k} \binom{-1/2}{k-1} (-4z)^{k-1} \\ &= \sum_{n \geq 0} \binom{-1/2}{n} \frac{(-4z)^n}{n+1} = \sum_{n \geq 0} \binom{2n}{n} \frac{z^n}{n+1}. \end{aligned}$$

The number of ways to parenthesize,  $C_n$ , is  $\binom{2n}{n} \frac{1}{n+1}$ .

We anticipated this result in Chapter 5, when we introduced the sequence of *Catalan numbers*  $\langle 1, 1, 2, 5, 14, \dots \rangle = \langle C_n \rangle$ . This sequence arises in dozens of problems that seem at first to be unrelated to each other [46], because many situations have a recursive structure that corresponds to the convolution recurrence (7.66).

For example, let's consider the following problem: How many sequences  $\langle a_1, a_2, \dots, a_{2n} \rangle$  of  $+1$ 's and  $-1$ 's have the property that

$$a_1 + a_2 + \dots + a_{2n} = 0$$

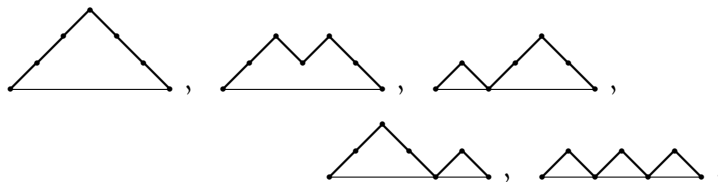
and have all their partial sums

$$a_1, \quad a_1 + a_2, \quad \dots, \quad a_1 + a_2 + \dots + a_{2n}$$

nonnegative? There must be  $n$  occurrences of  $+1$  and  $n$  occurrences of  $-1$ . We can represent this problem graphically by plotting the sequence of partial

the convo-  
ed recurrence  
led us to an  
recurring con-  
vulsion.

sums  $s_n = \sum_{k=1}^n a_k$  as a function of  $n$ : The five solutions for  $n = 3$  are



These are “mountain ranges” of width  $2n$  that can be drawn with line segments of the forms  $\nearrow$  and  $\searrow$ . It turns out that there are exactly  $C_n$  ways to do this, and the sequences can be related to the parenthesis problem in the following way: Put an extra pair of parentheses around the entire formula, so that there are  $n$  pairs of parentheses corresponding to the  $n$  multiplications. Now replace each ‘ $\cdot$ ’ by  $+1$  and each ‘ $)$ ’ by  $-1$  and erase everything else. For example, the formula  $x_0 \cdot ((x_1 \cdot x_2) \cdot (x_3 \cdot x_4))$  corresponds to the sequence  $\langle +1, +1, -1, +1, +1, -1, -1, -1 \rangle$  by this rule. The five ways to parenthesize  $x_0 \cdot x_1 \cdot x_2 \cdot x_3$  correspond to the five mountain ranges for  $n = 3$  shown above.

Moreover, a slight reformulation of our sequence-counting problem leads to a surprisingly simple combinatorial solution that avoids the use of generating functions: How many sequences  $\langle a_0, a_1, a_2, \dots, a_{2n} \rangle$  of  $+1$ ’s and  $-1$ ’s have the property that

$$a_0 + a_1 + a_2 + \dots + a_{2n} = 1,$$

when all the partial sums

$$a_0, \quad a_0 + a_1, \quad a_0 + a_1 + a_2, \quad \dots, \quad a_0 + a_1 + \dots + a_{2n}$$

are required to be *positive*? Clearly these are just the sequences of the previous problem, with the additional element  $a_0 = +1$  placed in front. But the sequences in the new problem can be enumerated by a simple counting argument, using a remarkable fact discovered by George Raney [302] in 1959: *If  $\langle x_1, x_2, \dots, x_m \rangle$  is any sequence of integers whose sum is  $+1$ , exactly one of the cyclic shifts*

$$\langle x_1, x_2, \dots, x_m \rangle, \quad \langle x_2, \dots, x_m, x_1 \rangle, \quad \dots, \quad \langle x_m, x_1, \dots, x_{m-1} \rangle$$

*has all of its partial sums positive.* For example, consider the sequence  $\langle 3, -5, 2, -2, 3, 0 \rangle$ . Its cyclic shifts are

$$\begin{array}{ll} \langle 3, -5, 2, -2, 3, 0 \rangle & \langle -2, 3, 0, 3, -5, 2 \rangle \\ \langle -5, 2, -2, 3, 0, 3 \rangle & \langle 3, 0, 3, -5, 2, -2 \rangle \checkmark \\ \langle 2, -2, 3, 0, 3, -5 \rangle & \langle 0, 3, -5, 2, -2, 3 \rangle \end{array}$$

and only the one that’s checked has entirely positive partial sums.





whose total sum is  $+1$ . Such sequences can be called *m-Raney sequences*. If there are  $k$  occurrences of  $(1-m)$  and  $mn+1-k$  occurrences of  $+1$ , we have

$$k(1-m) + (mn+1-k) = 1,$$

hence  $k = n$ . There are  $\binom{mn+1}{n}$  sequences with  $n$  occurrences of  $(1-m)$  and  $mn+1-n$  occurrences of  $+1$ , and Raney's lemma tells us that the number of such sequences with all partial sums positive is exactly

$$\binom{mn+1}{n} \frac{1}{mn+1} = \binom{mn}{n} \frac{1}{(m-1)n+1}. \quad (7.67)$$

So this is the number of *m-Raney sequences*. Let's call this a Fuss-Catalan number  $C_n^{(m)}$ , because the sequence  $\langle C_n^{(m)} \rangle$  was first investigated by N. I. Fuss [135] in 1791 (many years before Catalan himself got into the act). The ordinary Catalan numbers are  $C_n = C_n^{(2)}$ .

Now that we know the answer, (7.67), let's play "Jeopardy" and figure out a question that leads to it. In the case  $m = 2$  the question was: "What numbers  $C_n$  satisfy the recurrence  $C_n = \sum_k C_k C_{n-1-k} + [n=0]$ ?" We will try to find a similar question (a similar recurrence) in the general case.

The trivial sequence  $\langle +1 \rangle$  of length 1 is clearly an *m-Raney sequence*. If we put the number  $(1-m)$  at the right of any *m-Raney* sequence, we get an *m-Raney sequence*; the partial sums stay positive as they increase to  $+2$ , then  $+3$ ,  $\dots$ ,  $+m$ , and  $+1$ . Conversely, we can show that all *m-Raney sequences*  $\langle a_0, \dots, a_{mn} \rangle$  arise in this way, if  $n > 0$ : The last term  $a_{mn}$  must be  $(1-m)$ . The partial sums  $s_j = a_0 + \dots + a_{j-1}$  are positive for  $1 \leq j \leq mn$ , and  $s_{mn} = m$  because  $s_{mn} + a_{mn} = 1$ . Let  $k_1$  be the largest index  $\leq mn$  such that  $s_{k_1} = 1$ ; let  $k_2$  be largest such that  $s_{k_2} = 2$ ; and so on. Thus  $s_{k_j} = j$  and  $s_k > j$ , for  $k_j < k \leq mn$  and  $1 \leq j \leq m$ . It follows that  $k_m = mn$ , and we can verify without difficulty that each of the subsequences  $\langle a_0, \dots, a_{k_1-1} \rangle$ ,  $\langle a_{k_1}, \dots, a_{k_2-1} \rangle$ ,  $\dots$ ,  $\langle a_{k_{m-1}}, \dots, a_{k_m-1} \rangle$  is an *m-Raney sequence*. We must have  $k_1 = mn_1 + 1$ ,  $k_2 - k_1 = mn_2 + 1$ ,  $\dots$ ,  $k_m - k_{m-1} = mn_m + 1$ , for some nonnegative integers  $n_1, n_2, \dots, n_m$ .

Therefore  $\binom{mn+1}{n} \frac{1}{mn+1}$  is the answer to the following two interesting questions: "What are the numbers  $C_n^{(m)}$  defined by the recurrence

$$C_n^{(m)} = \left( \sum_{n_1+n_2+\dots+n_m=n-1} C_{n_1}^{(m)} C_{n_2}^{(m)} \dots C_{n_m}^{(m)} \right) + [n=0] \quad (7.68)$$

for all integers  $n$ ?" "If  $G(z)$  is a power series that satisfies

$$G(z) = z G(z)^m + 1, \quad (7.69)$$

what is  $[z^n] G(z)$ ?"

*(Attention, computer scientists: The stack interpretation now applies with respect to an m-ary operation, instead of the binary multiplication considered earlier.)*

Notice that these are not easy questions. In the ordinary Catalan case ( $m = 2$ ), we solved (7.6g) for  $G(z)$  and its coefficients by using the quadratic formula and the binomial theorem; but when  $m = 3$ , none of the standard techniques gives any clue about how to solve the cubic equation  $G = zG^3 + 1$ . So it has turned out to be easier to answer this question before asking it.

Now, however, we know enough to ask even harder questions and deduce their answers. How about this one: "What is  $[z^n] G(z)^l$ , if  $l$  is a positive integer and if  $G(z)$  is the power series defined by (7.6g)?" The argument we just gave can be used to show that  $[z^n] G(z)^l$  is the number of sequences of length  $mn + l$  with the following three properties:

- Each element is either  $+1$  or  $(1 - m)$ .
- The partial sums are all positive.
- The total sum is  $l$ .

For we get all such sequences in a unique way by putting together  $l$  sequences that have the  $m$ -Raney property. The number of ways to do this is

$$\sum_{n_1 + n_2 + \dots + n_l = n} C_{n_1}^{(m)} C_{n_2}^{(m)} \dots C_{n_l}^{(m)} = [z^n] G(z)^l.$$

Raney proved a generalization of his lemma that tells us how to count such sequences: *If  $\langle x_1, x_2, \dots, x_m \rangle$  is any sequence of integers with  $x_j \leq 1$  for all  $j$ , and with  $x_1 + x_2 + \dots + x_m = l > 0$ , then exactly  $l$  of the cyclic shifts*

$$\langle x_1, x_2, \dots, x_m \rangle, \langle x_2, \dots, x_m, x_1 \rangle, \dots, \langle x_m, x_1, \dots, x_{m-1} \rangle$$

*have all positive partial sums.*

For example, we can check this statement on the sequence  $\langle -2, 1, -1, 0, 1, 1, -1, 1, 1 \rangle$ . The cyclic shifts are

$$\begin{array}{ll} \langle -2, 1, -1, 0, 1, 1, -1, 1, 1 \rangle & \langle 1, -1, 1, 1, 1, -2, 1, -1, 0, 1 \rangle \\ \langle 1, -1, 0, 1, 1, -1, 1, 1, -2 \rangle & \langle -1, 1, 1, 1, -2, 1, -1, 0, 1, 1 \rangle \\ \langle -1, 0, 1, 1, -1, 1, 1, -2, 1 \rangle & \langle 1, 1, 1, -2, 1, -1, 0, 1, 1, -1 \rangle \checkmark \\ \langle 0, 1, 1, -1, 1, 1, -2, 1, -1 \rangle & \langle 1, 1, -2, 1, -1, 0, 1, 1, -1, 1 \rangle \\ \langle 1, 1, -1, 1, 1, -2, 1, -1, 0 \rangle \checkmark & \langle 1, -2, 1, -1, 0, 1, 1, -1, 1, 1 \rangle \end{array}$$

and only the two examples marked ' $\checkmark$ ' have all partial sums positive. This generalized lemma is proved in exercise 13.

A sequence of  $+1$ 's and  $(1 - m)$ 's that has length  $mn + l$  and total sum  $l$  must have exactly  $n$  occurrences of  $(1 - m)$ . The generalized lemma tells us that  $l/(mn + l)$  of these  $\binom{mn+l}{n}$  sequences have all partial sums positive;

hence our tough question has a surprisingly simple answer:

$$[z^n] G(z)^l = \binom{mn+l}{n} \frac{l}{mn+l}, \quad (7.70)$$

for all integers  $l > 0$ .

Readers who haven't forgotten Chapter 5 might well be experiencing *déjà vu*: "That formula looks familiar; haven't we seen it before?" Yes, indeed; Lambert's equation (5.60) says that

$$[z^n] \mathcal{B}_t(z)^r = \binom{tn+r}{n} \frac{r}{tn+r}.$$

Therefore the generating function  $G(z)$  in (7.69) must actually be the generalized binomial series  $\mathcal{B}_m(z)$ . Sure enough, equation (5.59) says

$$\mathcal{B}_m(z)^{1-m} - \mathcal{B}_m(z)^{-m} = z,$$

which is the same as

$$\mathcal{B}_m(z) - 1 = z\mathcal{B}_m(z)^m.$$

Let's switch to the notation of Chapter 5, now that we know we're dealing with generalized binomials. Chapter 5 stated a bunch of identities without proof. We have now closed part of the gap by proving that the power series  $\mathcal{B}_t(z)$  defined by

$$\mathcal{B}_t(z) = \sum_n \binom{tn+1}{n} \frac{z^n}{tn+1}$$

has the remarkable property that

$$\mathcal{B}_t(z)^r = \sum_n \binom{tn+r}{n} \frac{r z^n}{tn+r},$$

whenever  $t$  and  $r$  are positive integers.

Can we extend these results to *arbitrary* values of  $t$  and  $r$ ? Yes; because the coefficients  $\binom{tn+r}{n} \frac{r}{tn+r}$  are polynomials in  $t$  and  $r$ . The general  $r$ th power defined by

$$\mathcal{B}_t(z)^r = e^{r \ln \mathcal{B}_t(z)} = \sum_{n \geq 0} \frac{(r \ln \mathcal{B}_t(z))^n}{n!} = \sum_{n \geq 0} \frac{r^n}{n!} \left( - \sum_{m \geq 1} \frac{(1 - \mathcal{B}_t(z))^m}{m} \right)^n$$

has coefficients that are polynomials in  $t$  and  $r$ ; and those polynomials are equal to  $\binom{tn+r}{n} \frac{r}{tn+r}$  for infinitely many values of  $t$  and  $r$ . So the two sequences of polynomials must be identically equal.

Chapter 5 also mentions the generalized exponential series

$$\mathcal{E}_t(z) = \sum_{n \geq 0} \frac{(tn+1)^{n-1}}{n!} z^n,$$

which is said in (5.60) to have an equally remarkable property:

$$[z^n] \mathcal{E}_t(z)^r = \frac{r(tn+r)^{n-1}}{n!}. \quad (7.71)$$

We can prove this as a limiting case of the formulas for  $\mathcal{B}_t(z)$ , because it is not difficult to show that

$$\mathcal{E}_t(z)^r = \lim_{x \rightarrow \infty} \mathcal{B}_{xt}(z/x)^{xr}.$$

## 7.6 EXPONENTIAL GF'S

Sometimes a sequence  $\langle g_n \rangle$  has a generating function whose properties are quite complicated, while the related sequence  $\langle g_n/n! \rangle$  has a generating function that's quite simple. In such cases we naturally prefer to work with  $\langle g_n/n! \rangle$  and then multiply by  $n!$  at the end. This trick works sufficiently often that we have a special name for it: We call the power series

$$\widehat{G}(z) = \sum_{n \geq 0} g_n \frac{z^n}{n!} \quad (7.72)$$

the *exponential generating function* or “egf” of the sequence  $\langle g_0, g_1, g_2, \dots \rangle$ . This name arises because the exponential function  $e^z$  is the egf of  $\langle 1, 1, 1, \dots \rangle$ .

Many of the generating functions in Table 352 are actually egf's. For example, equation (7.50) says that  $(\ln \frac{1}{1-z})^m/m!$  is the egf for the sequence  $\langle \begin{bmatrix} 0 \\ m \end{bmatrix}, \begin{bmatrix} 1 \\ m \end{bmatrix}, \begin{bmatrix} 2 \\ m \end{bmatrix}, \dots \rangle$ . The ordinary generating function for this sequence is much more complicated (and also divergent).

Exponential generating functions have their own basic maneuvers, analogous to the operations we learned in Section 7.2. For example, if we multiply the egf of  $\langle g_n \rangle$  by  $z$ , we get

$$\sum_{n \geq 0} g_n \frac{z^{n+1}}{n!} = \sum_{n \geq 1} g_{n-1} \frac{z^n}{(n-1)!} = \sum_{n \geq 0} n g_{n-1} \frac{z^n}{n!};$$

this is the egf of  $\langle 0, g_0, 2g_1, \dots \rangle = \langle n g_{n-1} \rangle$ .

Differentiating the egf of  $\langle g_0, g_1, g_2, \dots \rangle$  with respect to  $z$  gives

$$\sum_{n \geq 0} n g_n \frac{z^{n-1}}{n!} = \sum_{n \geq 1} g_n \frac{z^{n-1}}{(n-1)!} = \sum_{n \geq 0} g_{n+1} \frac{z^n}{n!}; \quad (7.73)$$

we having  
yet?

this is the egf of  $\langle g_1, g_2, \dots \rangle$ . Thus differentiation on egf's corresponds to the left-shift operation  $(G(z) - g_0)/z$  on ordinary gf's. (We used this left-shift property of egf's when we studied hypergeometric series, (5.106).) Integration of an egf gives

$$\int_0^z \sum_{n \geq 0} g_n \frac{t^n}{n!} dt = \sum_{n \geq 0} g_n \frac{z^{n+1}}{(n+1)!} = \sum_{n \geq 1} g_{n-1} \frac{z^n}{n!}; \quad (7.74)$$

this is a right shift, the egf of  $\langle 0, g_0, g_1, \dots \rangle$ .

The most interesting operation on egf's, as on ordinary gf's, is multiplication. If  $\hat{F}(z)$  and  $\hat{G}(z)$  are egf's for  $\langle f_n \rangle$  and  $\langle g_n \rangle$ , then  $\hat{F}(z)\hat{G}(z) = \hat{H}(z)$  is the egf for a sequence  $\langle h_n \rangle$  called the *binomial convolution* of  $\langle f_n \rangle$  and  $\langle g_n \rangle$ :

$$h_n = \sum_k \binom{n}{k} f_k g_{n-k}. \quad (7.75)$$

Binomial coefficients appear here because  $\binom{n}{k} = n!/k!(n-k)!$ , hence

$$\frac{h_n}{n!} = \sum_{k=0}^n \frac{f_k}{k!} \frac{g_{n-k}}{(n-k)!};$$

in other words,  $\langle h_n/n! \rangle$  is the ordinary convolution of  $\langle f_n/n! \rangle$  and  $\langle g_n/n! \rangle$ .

Binomial convolutions occur frequently in applications. For example, we defined the Bernoulli numbers in (6.79) by the implicit recurrence

$$\sum_{j=0}^m \binom{m+1}{j} B_j = [m=0], \quad \text{for all } m \geq 0;$$

this can be rewritten as a binomial convolution, if we substitute  $n$  for  $m+1$  and add the term  $B_n$  to both sides:

$$\sum_k \binom{n}{k} B_k = B_n + [n=1], \quad \text{for all } n \geq 0. \quad (7.76)$$

We can now relate this recurrence to power series (as promised in Chapter 6) by introducing the egf for Bernoulli numbers,  $\hat{B}(z) = \sum_{n \geq 0} B_n z^n/n!$ . The left-hand side of (7.76) is the binomial convolution of  $\langle B_n \rangle$  with the constant sequence  $\langle 1, 1, 1, \dots \rangle$ ; hence the egf of the left-hand side is  $\hat{B}(z)e^z$ . The egf of the right-hand side is  $\sum_{n \geq 0} (B_n + [n=1])z^n/n! = \hat{B}(z) + z$ . Therefore we must have  $\hat{B}(z) = z/(e^z - 1)$ ; we have proved equation (6.81), which appears also in Table 352 as equation (7.44).

Now let's look again at a sum that has been popping up frequently in this book,

$$S_m(n) = 0^m + 1^m + 2^m + \cdots + (n-1)^m = \sum_{0 \leq k < n} k^m.$$

This time we will try to analyze the problem with generating functions, in hopes that it will suddenly become simpler. We will consider  $n$  to be fixed and  $m$  variable; thus our goal is to understand the coefficients of the power series

$$S(z) = S_0(n) + S_1(n)z + S_2(n)z^2 + \cdots = \sum_{m \geq 0} S_m(n)z^m.$$

We know that the generating function for  $\langle 1, k, k^2, \dots \rangle$  is

$$\frac{1}{1-kz} = \sum_{m \geq 0} k^m z^m,$$

hence

$$S(z) = \sum_{m \geq 0} \sum_{0 \leq k < n} k^m z^m = \sum_{0 \leq k < n} \frac{1}{1-kz}$$

by interchanging the order of summation. We can put this sum in closed form,

$$\begin{aligned} S(z) &= \frac{1}{z} \left( \frac{1}{z^{-1}-0} + \frac{1}{z^{-1}-1} + \cdots + \frac{1}{z^{-1}-n+1} \right) \\ &= \frac{1}{z} (H_{z^{-1}} - H_{z^{-1}-n}); \end{aligned} \quad (7.77)$$

but we know nothing about expanding such a closed form in powers of  $z$ .

Exponential generating functions come to the rescue. The egf of our sequence  $\langle S_0(n), S_1(n), S_2(n), \dots \rangle$  is

$$\widehat{S}(z, n) = S_0(n) + S_1(n) \frac{z}{1!} + S_2(n) \frac{z^2}{2!} + \cdots = \sum_{m \geq 0} S_m(n) \frac{z^m}{m!}.$$

To get these coefficients  $S_m(n)$  we can use the egf for  $\langle 1, k, k^2, \dots \rangle$ , namely

$$e^{kz} = \sum_{m \geq 0} k^m \frac{z^m}{m!},$$

and we have

$$\widehat{S}(z, n) = \sum_{m \geq 0} \sum_{0 \leq k < n} k^m \frac{z^m}{m!} = \sum_{0 \leq k < n} e^{kz}.$$

And the latter sum is a geometric progression, so there's a closed form

$$\widehat{S}(z, n) = \frac{e^{nz} - 1}{e^z - 1}. \quad (7.78)$$

Eureka! All we need to do is figure out the coefficients of this relatively simple function, and we'll know  $S_m(n)$ , because  $S_m(n) = m! [z^m] \widehat{S}(z, n)$ .

Here's where Bernoulli numbers come into the picture. We observed a moment ago that the egf for Bernoulli numbers is

$$\widehat{B}(z) = \sum_{k \geq 0} B_k \frac{z^k}{k!} = \frac{z}{e^z - 1};$$

hence we can write

$$\begin{aligned} \widehat{S}(z, n) &= \widehat{B}(z) \frac{e^{nz} - 1}{z} \\ &= \left( B_0 \frac{z^0}{0!} + B_1 \frac{z^1}{1!} + B_2 \frac{z^2}{2!} + \cdots \right) \left( n \frac{z^0}{1!} + n^2 \frac{z^1}{2!} + n^3 \frac{z^2}{3!} + \cdots \right). \end{aligned}$$

The sum  $S_m(n)$  is  $m!$  times the coefficient of  $z^m$  in this product. For example,

$$\begin{aligned} S_0(n) &= 0! \left( B_0 \frac{n}{1! 0!} \right) = n; \\ S_1(n) &= 1! \left( B_0 \frac{n^2}{2! 0!} + B_1 \frac{n}{1! 1!} \right) = \frac{1}{2}n^2 - \frac{1}{2}n; \\ S_2(n) &= 2! \left( B_0 \frac{n^3}{3! 0!} + B_1 \frac{n^2}{2! 1!} + B_2 \frac{n}{1! 2!} \right) = \frac{1}{3}n^3 - \frac{1}{2}n^2 + \frac{1}{6}n. \end{aligned}$$

We have therefore derived the formula  $\square_n = S_2(n) = \frac{1}{3}n(n - \frac{1}{2})(n - 1)$  for the umpteenth time, and this was the simplest derivation of all: In a few lines we have found the general behavior of  $S_m(n)$  for all  $m$ .

The general formula can be written

$$S_{m-1}(n) = \frac{1}{m} (B_m(n) - B_m(0)), \quad (7.79)$$

where  $B_m(x)$  is the *Bernoulli polynomial* defined by

$$B_m(x) = \sum_k \binom{m}{k} B_k x^{m-k}. \quad (7.80)$$

Here's why: The Bernoulli polynomial is the binomial convolution of the sequence  $\langle B_0, B_1, B_2, \dots \rangle$  with  $\langle 1, x, x^2, \dots \rangle$ ; hence the exponential generating



function for  $\langle B_0(x), B_1(x), B_2(x), \dots \rangle$  is the product of their egf's,

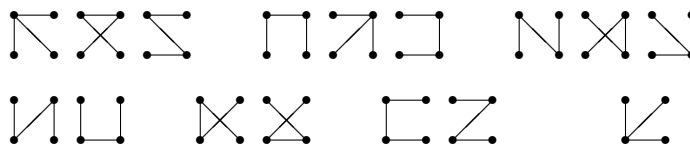
$$\widehat{B}(z, x) = \sum_{m \geq 0} B_m(x) \frac{z^m}{m!} = \frac{z}{e^z - 1} \sum_{m \geq 0} x^m \frac{z^m}{m!} = \frac{ze^{xz}}{e^z - 1}. \quad (7.81)$$

Equation (7.79) follows because the egf for  $\langle 0, S_0(n), 2S_1(n), \dots \rangle$  is, by (7.78),


$$z \frac{e^{nz} - 1}{e^z - 1} = \widehat{B}(z, n) - \widehat{B}(z, 0).$$

Let's turn now to another problem for which egf's are just the thing: How many spanning trees are possible in the *complete graph* on  $n$  vertices  $\{1, 2, \dots, n\}$ ? Let's call this number  $t_n$ . The complete graph has  $\frac{1}{2}n(n-1)$  edges, one edge joining each pair of distinct vertices; so we're essentially looking for the total number of ways to connect up  $n$  given things by drawing  $n-1$  lines between them.

We have  $t_1 = t_2 = 1$ . Also  $t_3 = 3$ , because a complete graph on three vertices is a fan of order 2; we know that  $f_2 = 3$ . And there are sixteen spanning trees when  $n = 4$ :


(7.82)

Hence  $t_4 = 16$ .

Our experience with the analogous problem for fans suggests that the best way to tackle this problem is to single out one vertex, and to look at the blocks or components that the spanning tree joins together when we ignore all edges that touch the special vertex. If the non-special vertices form  $m$  components of sizes  $k_1, k_2, \dots, k_m$ , then we can connect them to the special vertex in  $k_1 k_2 \dots k_m$  ways. For example, in the case  $n = 4$ , we can consider the lower left vertex to be special. The top row of (7.82) shows  $3t_3$  cases where the other three vertices are joined among themselves in  $t_3$  ways and then connected to the lower left in 3 ways. The bottom row shows  $2 \cdot 1 \times t_2 t_1 \times \binom{3}{2}$  solutions where the other three vertices are divided into components of sizes 2 and 1 in  $\binom{3}{2}$  ways; there's also the case  where the other three vertices are completely unconnected among themselves.

This line of reasoning leads to the recurrence

$$t_n = \sum_{m \geq 0} \frac{1}{m!} \sum_{k_1 + \dots + k_m = n-1} \binom{n-1}{k_1, k_2, \dots, k_m} k_1 k_2 \dots k_m t_{k_1} t_{k_2} \dots t_{k_m}$$

for all  $n > 1$ . Here's why: There are  $\binom{n-1}{k_1, k_2, \dots, k_m}$  ways to assign  $n-1$  elements to a sequence of  $m$  components of respective sizes  $k_1, k_2, \dots, k_m$ ; there are  $t_{k_1} t_{k_2} \dots t_{k_m}$  ways to connect up those individual components with spanning trees; there are  $k_1 k_2 \dots k_m$  ways to connect vertex  $n$  to those components; and we divide by  $m!$  because we want to disregard the order of the components. For example, when  $n = 4$  the recurrence says that

$$t_4 = 3t_3 + \frac{1}{2} \left( \binom{3}{1,2} 2t_1 t_2 + \binom{3}{2,1} 2t_2 t_1 \right) + \frac{1}{6} \left( \binom{3}{1,1,1} t_1^3 \right) = 3t_3 + 6t_2 t_1 + t_1^3.$$

The recurrence for  $t_n$  looks formidable at first, possibly even frightening; but it really isn't bad, only convoluted. We can define

$$u_n = n t_n$$

and then everything simplifies considerably:

$$\frac{u_n}{n!} = \sum_{m \geq 0} \frac{1}{m!} \sum_{k_1 + k_2 + \dots + k_m = n-1} \frac{u_{k_1}}{k_1!} \frac{u_{k_2}}{k_2!} \dots \frac{u_{k_m}}{k_m!}, \quad \text{if } n > 1. \quad (7.83)$$

The inner sum is the coefficient of  $z^{n-1}$  in the egf  $\hat{U}(z)$ , raised to the  $m$ th power; and we obtain the correct formula also when  $n = 1$ , if we add in the term  $\hat{U}(z)^0$  that corresponds to the case  $m = 0$ . So

$$\frac{u_n}{n!} = [z^{n-1}] \sum_{m \geq 0} \frac{1}{m!} \hat{U}(z)^m = [z^{n-1}] e^{\hat{U}(z)} = [z^n] z e^{\hat{U}(z)}$$

for all  $n > 0$ , and we have the equation

$$\hat{U}(z) = z e^{\hat{U}(z)}. \quad (7.84)$$

Progress! Equation (7.84) is almost like

$$\mathcal{E}(z) = e^{z \mathcal{E}(z)},$$

which defines the generalized exponential series  $\mathcal{E}(z) = \mathcal{E}_1(z)$  in (5.59) and (7.71); indeed, we have

$$\hat{U}(z) = z \mathcal{E}(z).$$

So we can read off the answer to our problem:

$$t_n = \frac{u_n}{n} = \frac{n!}{n} [z^n] \hat{U}(z) = (n-1)! [z^{n-1}] \mathcal{E}(z) = n^{n-2}. \quad (7.85)$$

The complete graph on  $\{1, 2, \dots, n\}$  has exactly  $n^{n-2}$  spanning trees, for all  $n > 0$ .

## 7.7 DIRICHLET GENERATING FUNCTIONS

There are many other possible ways to generate a sequence from a series; any system of “kernel” functions  $K_n(z)$  such that

$$\sum_n g_n K_n(z) = 0 \implies g_n = 0 \text{ for all } n$$

can be used, at least in principle. Ordinary generating functions use  $K_n(z) = z^n$ , and exponential generating functions use  $K_n(z) = z^n/n!$ ; we could also try falling factorial powers  $z^{\underline{n}}$ , or binomial coefficients  $z^{\underline{n}}/n! = \binom{z}{n}$ .

The most important alternative to gf’s and egf’s uses the kernel functions  $1/n^z$ ; it is intended for sequences  $\langle g_1, g_2, \dots \rangle$  that begin with  $n = 1$  instead of  $n = 0$ :

$$\tilde{G}(z) = \sum_{n \geq 1} \frac{g_n}{n^z}. \quad (7.86)$$

This is called a *Dirichlet generating function* (dgf), because the German mathematician Gustav Lejeune Dirichlet (1805–1859) made much of it.

For example, the dgf of the constant sequence  $\langle 1, 1, 1, \dots \rangle$  is

$$\sum_{n \geq 1} \frac{1}{n^z} = \zeta(z). \quad (7.87)$$

This is Riemann’s *zeta function*, which we have also called the generalized harmonic number  $H_\infty^{(z)}$  when  $z > 1$ .

The product of Dirichlet generating functions corresponds to a special kind of convolution:

$$\tilde{F}(z)\tilde{G}(z) = \sum_{l, m \geq 1} \frac{f_l}{l^z} \frac{g_m}{m^z} = \sum_{n \geq 1} \frac{1}{n^z} \sum_{l, m \geq 1} f_l g_m [l \cdot m = n].$$

Thus  $\tilde{F}(z)\tilde{G}(z) = \tilde{H}(z)$  is the dgf of the sequence

$$h_n = \sum_{d \mid n} f_d g_{n/d}. \quad (7.88)$$

For example, we know from (4.55) that  $\sum_{d \mid n} \mu(d) = [n=1]$ ; this is the Dirichlet convolution of the Möbius sequence  $\langle \mu(1), \mu(2), \mu(3), \dots \rangle$  with  $\langle 1, 1, 1, \dots \rangle$ , hence

$$\tilde{M}(z)\zeta(z) = \sum_{n \geq 1} \frac{[n=1]}{n^z} = 1. \quad (7.89)$$

In other words, the dgf of  $\langle \mu(1), \mu(2), \mu(3), \dots \rangle$  is  $\zeta(z)^{-1}$ .

Dirichlet generating functions are particularly valuable when the sequence  $\langle g_1, g_2, \dots \rangle$  is a *multiplicative function*, namely when

$$g_{mn} = g_m g_n \quad \text{for } m \perp n.$$

In such cases the values of  $g_n$  for all  $n$  are determined by the values of  $g_p$  when  $n$  is a power of a prime, and we can factor the dgf into a product over primes:

$$\tilde{G}(z) = \prod_{p \text{ prime}} \left( 1 + \frac{g_p}{p^z} + \frac{g_{p^2}}{p^{2z}} + \frac{g_{p^3}}{p^{3z}} + \dots \right). \quad (7.90)$$

If, for instance, we set  $g_n = 1$  for all  $n$ , we obtain a product representation of Riemann's zeta function:

$$\zeta(z) = \prod_{p \text{ prime}} \left( \frac{1}{1 - p^{-z}} \right). \quad (7.91)$$

The Möbius function has  $\mu(p) = -1$  and  $\mu(p^k) = 0$  for  $k > 1$ , hence its dgf is

$$\tilde{M}(z) = \prod_{p \text{ prime}} (1 - p^{-z}); \quad (7.92)$$

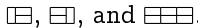
this agrees, of course, with (7.89) and (7.91). Euler's  $\varphi$  function has  $\varphi(p^k) = p^k - p^{k-1}$ , hence its dgf has the factored form

$$\tilde{\Phi}(z) = \prod_{p \text{ prime}} \left( 1 + \frac{p-1}{p^z - p} \right) = \prod_{p \text{ prime}} \frac{1 - p^{-z}}{1 - p^{1-z}}. \quad (7.93)$$

We conclude that  $\tilde{\Phi}(z) = \zeta(z-1)/\zeta(z)$ .

## Exercises

### Warmups

- 1 An eccentric collector of  $2 \times n$  domino tilings pays \$4 for each vertical domino and \$1 for each horizontal domino. How many tilings are worth exactly \$ $m$  by this criterion? For example, when  $m = 6$  there are three solutions: .
- 2 Give the generating function and the exponential generating function for the sequence  $\langle 2, 5, 13, 35, \dots \rangle = \langle 2^n + 3^n \rangle$  in closed form.
- 3 What is  $\sum_{n \geq 0} H_n / 10^n$ ?
- 4 The general expansion theorem for rational functions  $P(z)/Q(z)$  is not completely general, because it restricts the degree of  $P$  to be less than the degree of  $Q$ . What happens if  $P$  has a larger degree than this?

- 5 Find a generating function  $S(z)$  such that

$$[z^n] S(z) = \sum_k \binom{r}{k} \binom{r}{n-2k}.$$

### Basics

- 6 Show that the recurrence (7.32) can be solved by the repertoire method, without using generating functions.
- 7 Solve the recurrence

$$\begin{aligned} g_0 &= 1; \\ g_n &= g_{n-1} + 2g_{n-2} + \cdots + ng_0, \quad \text{for } n > 0. \end{aligned}$$

- 8 What is  $[z^n] (\ln(1-z))^2 / (1-z)^{m+1}$ ?
- 9 Use the result of the previous exercise to evaluate  $\sum_{k=0}^n H_k H_{n-k}$ .
- 10 Set  $r = s = -1/2$  in identity (7.62) and then remove all occurrences of  $1/2$  by using tricks like (5.36). What amazing identity do you deduce?
- 11 This problem, whose three parts are independent, gives practice in the manipulation of generating functions. We assume that  $A(z) = \sum_n a_n z^n$ ,  $B(z) = \sum_n b_n z^n$ ,  $C(z) = \sum_n c_n z^n$ , and that the coefficients are zero for negative  $n$ .
- a If  $c_n = \sum_{j+2k \leq n} a_j b_k$ , express  $C$  in terms of  $A$  and  $B$ .
- b If  $nb_n = \sum_{k=0}^n 2^k a_k / (n-k)!$ , express  $A$  in terms of  $B$ .
- c If  $r$  is a real number and if  $a_n = \sum_{k=0}^n \binom{r+k}{k} b_{n-k}$ , express  $A$  in terms of  $B$ ; then use your formula to find coefficients  $f_k(r)$  such that  $b_n = \sum_{k=0}^n f_k(r) a_{n-k}$ .
- 12 How many ways are there to put the numbers  $\{1, 2, \dots, 2n\}$  into a  $2 \times n$  array so that rows and columns are in increasing order from left to right and from top to bottom? For example, one solution when  $n = 5$  is

$$\begin{pmatrix} 1 & 2 & 4 & 5 & 8 \\ 3 & 6 & 7 & 9 & 10 \end{pmatrix}.$$

- 13 Prove Raney's generalized lemma, which is stated just before (7.70).
- 14 Solve the recurrence

$$\begin{aligned} g_0 &= 0, & g_1 &= 1, \\ g_n &= -2ng_{n-1} + \sum_k \binom{n}{k} g_k g_{n-k}, & \text{for } n > 1, \end{aligned}$$

by using an exponential generating function.

deduce that Clark  
is really  
German.

- 15 The *Bell number*  $\omega_n$  is the number of ways to partition  $n$  things into subsets. For example,  $\omega_3 = 5$  because we can partition  $\{1, 2, 3\}$  in the following ways:

$$\{1, 2, 3\}; \quad \{1, 2\} \cup \{3\}; \quad \{1, 3\} \cup \{2\}; \quad \{1\} \cup \{2, 3\}; \quad \{1\} \cup \{2\} \cup \{3\}.$$

Prove that  $\omega_{n+1} = \sum_k \binom{n}{k} \omega_{n-k}$ , and use this recurrence to find a closed form for the exponential generating function  $P(z) = \sum_n \omega_n z^n / n!$ .

- 16 Two sequences  $\langle a_n \rangle$  and  $\langle b_n \rangle$  are related by the convolution formula

$$b_n = \sum_{k_1 + 2k_2 + \cdots + nk_n = n} \binom{a_1 + k_1 - 1}{k_1} \binom{a_2 + k_2 - 1}{k_2} \cdots \binom{a_n + k_n - 1}{k_n};$$

also  $a_0 = 0$  and  $b_0 = 1$ . Prove that the corresponding generating functions satisfy  $\ln B(z) = A(z) + \frac{1}{2}A(z^2) + \frac{1}{3}A(z^3) + \cdots$ .

- 17 Show that the exponential generating function  $\widehat{G}(z)$  of a sequence is related to the ordinary generating function  $G(z)$  by the formula

$$\int_0^\infty \widehat{G}(zt) e^{-t} dt = G(z),$$

if the integral exists.

- 18 Find the Dirichlet generating functions for the sequences

- a  $g_n = \sqrt{n}$ ;
- b  $g_n = \ln n$ ;
- c  $g_n = [n \text{ is squarefree}]$ .

Express your answers in terms of the zeta function. (Squarefreeness is defined in exercise 4.13.)

- 19 Every power series  $F(z) = \sum_{n \geq 0} f_n z^n$  with  $f_0 = 1$  defines a sequence of polynomials  $f_n(x)$  by the rule

$$F(z)^x = \sum_{n \geq 0} f_n(x) z^n,$$

where  $f_n(1) = f_n$  and  $f_n(0) = [n=0]$ . In general,  $f_n(x)$  has degree  $n$ . Show that such polynomials always satisfy the convolution formulas

$$\sum_{k=0}^n f_k(x) f_{n-k}(y) = f_n(x+y);$$

$$(x+y) \sum_{k=0}^n k f_k(x) f_{n-k}(y) = x n f_n(x+y).$$

(The identities in Tables 202 and 272 are special cases of this trick.)

*What do you mean, "in general"? If  $f_1 = f_2 = \cdots = f_{m-1} = 0$ , the degree of  $f_n(x)$  is at most  $\lfloor n/m \rfloor$ .*

- 20** A power series  $G(z)$  is called *differentiably finite* if there exist finitely many polynomials  $P_0(z), \dots, P_m(z)$ , not all zero, such that

$$P_0(z)G(z) + P_1(z)G'(z) + \cdots + P_m(z)G^{(m)}(z) = 0.$$

A sequence of numbers  $\langle g_0, g_1, g_2, \dots \rangle$  is called *polynomially recursive* if there exist finitely many polynomials  $p_0(z), \dots, p_m(z)$ , not all zero, such that

$$p_0(n)g_n + p_1(n)g_{n+1} + \cdots + p_m(n)g_{n+m} = 0$$

for all integers  $n \geq 0$ . Prove that a generating function is differentiably finite if and only if its sequence of coefficients is polynomially recursive.

### Homework exercises

- 21** A robber holds up a bank and demands \$500 in tens and twenties. He also demands to know the number of ways in which the cashier can give him the money. Find a generating function  $G(z)$  for which this number is  $[z^{500}]G(z)$ , and a more compact generating function  $\tilde{G}(z)$  for which this number is  $[z^{50}]\tilde{G}(z)$ . Determine the required number of ways by (a) using partial fractions; (b) using a method like (7.39).
- 22** Let  $P$  be the sum of all ways to “triangulate” polygons:

$$P = \text{---} + \triangle + \begin{array}{|c|} \hline \diagup \\ \hline \end{array} + \begin{array}{|c|} \hline \diagdown \\ \hline \end{array} + \begin{array}{|c|} \hline \diagup \quad \diagdown \\ \hline \end{array} + \begin{array}{|c|} \hline \diagup \quad \diagdown \quad \diagup \\ \hline \end{array} + \begin{array}{|c|} \hline \diagup \quad \diagdown \quad \diagup \quad \diagdown \\ \hline \end{array} + \cdots.$$

(The first term represents a degenerate polygon with only two vertices; every other term shows a polygon that has been divided into triangles. For example, a pentagon can be triangulated in five ways.) Define a “multiplication” operation  $A \triangle B$  on triangulated polygons  $A$  and  $B$  so that the equation

$$P = \text{---} + P \triangle P$$

is valid. Then replace each triangle by ‘ $z$ ’; what does this tell you about the number of ways to decompose an  $n$ -gon into triangles?

- 23** In how many ways can a  $2 \times 2 \times n$  pillar be built out of  $2 \times 1 \times 1$  bricks?
- 24** How many spanning trees are in an  $n$ -wheel (a graph with  $n$  “outer” vertices in a cycle, each connected to an  $(n+1)$ st “hub” vertex), when  $n \geq 3$ ?

- 25 Let  $m \geq 2$  be an integer. What is a closed form for the generating function of the sequence  $\langle n \bmod m \rangle$ , as a function of  $z$  and  $m$ ? Use this generating function to express ' $n \bmod m$ ' in terms of the complex number  $\omega = e^{2\pi i/m}$ . (For example, when  $m = 2$  we have  $\omega = -1$  and  $n \bmod 2 = \frac{1}{2} - \frac{1}{2}(-1)^n$ .)
- 26 The second-order Fibonacci numbers  $\langle \mathfrak{F}_n \rangle$  are defined by the recurrence

$$\begin{aligned}\mathfrak{F}_0 &= 0; & \mathfrak{F}_1 &= 1; \\ \mathfrak{F}_n &= \mathfrak{F}_{n-1} + \mathfrak{F}_{n-2} + F_n, & \text{for } n > 1.\end{aligned}$$

Express  $\mathfrak{F}_n$  in terms of the usual Fibonacci numbers  $F_n$  and  $F_{n+1}$ .

- 27 A  $2 \times n$  domino tiling can also be regarded as a way to draw  $n$  disjoint lines in a  $2 \times n$  array of points:

$$\begin{array}{ccccccc} | & \text{---} & \text{---} & | & \text{---} & | & | \\ | & \text{---} & \text{---} & | & \text{---} & | & | \end{array}$$

If we superimpose two such patterns, we get a set of cycles, since every point is touched by two lines. For example, if the lines above are combined with the lines

$$\begin{array}{ccccccc} | & | & \text{---} & \text{---} & \text{---} & \text{---} & \\ | & | & \text{---} & \text{---} & \text{---} & \text{---} & \end{array},$$

the result is

$$\begin{array}{ccccccc} | & | & \text{---} & \text{---} & \text{---} & \text{---} & \\ | & | & \text{---} & \text{---} & \text{---} & \text{---} & \end{array}.$$

The same set of cycles is also obtained by combining

$$\begin{array}{ccccccc} | & | & \text{---} & \text{---} & \text{---} & | & | \\ | & | & \text{---} & \text{---} & \text{---} & | & | \end{array} \quad \text{with} \quad \begin{array}{ccccccc} | & \text{---} & \text{---} & | & \text{---} & \text{---} & \\ | & \text{---} & \text{---} & | & \text{---} & \text{---} & \end{array}.$$

But we get a unique way to reconstruct the original patterns from the superimposed ones if we assign orientations to the vertical lines by using arrows that go alternately up/down/up/down/... in the first pattern and alternately down/up/down/up/... in the second. For example,

$$\begin{array}{ccccccc} | & \text{---} & \text{---} & | & \text{---} & | & | \\ | & \text{---} & \text{---} & | & \text{---} & | & | \end{array} + \begin{array}{ccccccc} | & | & \text{---} & \text{---} & \text{---} & \text{---} & \\ | & | & \text{---} & \text{---} & \text{---} & \text{---} & \end{array} = \begin{array}{ccccccc} | & | & \text{---} & \text{---} & \text{---} & \text{---} & \\ | & | & \text{---} & \text{---} & \text{---} & \text{---} & \end{array}.$$

The number of such oriented cycle patterns must therefore be  $T_n^2 = F_{n+1}^2$ , and we should be able to prove this via algebra. Let  $Q_n$  be the number of oriented  $2 \times n$  cycle patterns. Find a recurrence for  $Q_n$ , solve it with generating functions, and deduce algebraically that  $Q_n = F_{n+1}^2$ .

- 28 The coefficients of  $A(z)$  in (7.39) satisfy  $A_r + A_{r+10} + A_{r+20} + A_{r+30} = 100$  for  $0 \leq r < 10$ . Find a "simple" explanation for this.



29 What is the sum of Fibonacci products

$$\sum_{m>0} \sum_{\substack{k_1+k_2+\dots+k_m=n \\ k_1, k_2, \dots, k_m > 0}} F_{k_1} F_{k_2} \dots F_{k_m} ?$$

30 If the generating function  $G(z) = 1/(1 - \alpha z)(1 - \beta z)$  has the partial fraction decomposition  $a/(1 - \alpha z) + b/(1 - \beta z)$ , what is the partial fraction decomposition of  $G(z)^n$ ?

31 What function  $g(n)$  of the positive integer  $n$  satisfies the recurrence

$$\sum_{d \mid n} g(d) \varphi(n/d) = 1,$$

where  $\varphi$  is Euler's totient function?

32 An *arithmetic progression* is an infinite set of integers

$$\{an + b\} = \{b, a + b, 2a + b, 3a + b, \dots\}.$$

A set of arithmetic progressions  $\{a_1 n + b_1\}, \dots, \{a_m n + b_m\}$  is called an *exact cover* if every nonnegative integer occurs in one and only one of the progressions. For example, the three progressions  $\{2n\}, \{4n + 1\}, \{4n + 3\}$  constitute an exact cover. Show that if  $\{a_1 n + b_1\}, \dots, \{a_m n + b_m\}$  is an exact cover such that  $2 \leq a_1 \leq \dots \leq a_m$ , then  $a_{m-1} = a_m$ . *Hint:* Use generating functions.

### Exam problems

33 What is  $[w^m z^n] (\ln(1 + z))/(1 - wz)$ ?

34 Find a closed form for the generating function  $\sum_{n \geq 0} G_n(z) w^n$ , if

$$G_n(z) = \sum_{k \leq n/m} \binom{n - mk}{k} z^{mk}.$$

(Here  $m$  is a fixed positive integer.)

35 Evaluate the sum  $\sum_{0 < k < n} 1/k(n - k)$  in two ways:

- a Expand the summand in partial fractions.
- b Treat the sum as a convolution and use generating functions.

36 Let  $A(z)$  be the generating function for  $\langle a_0, a_1, a_2, a_3, \dots \rangle$ . Express  $\sum_n a_{\lfloor n/m \rfloor} z^n$  in terms of  $A$ ,  $z$ , and  $m$ .

- 37 Let  $a_n$  be the number of ways to write the positive integer  $n$  as a sum of powers of 2, disregarding order. For example,  $a_4 = 4$ , since  $4 = 2 + 2 = 2 + 1 + 1 = 1 + 1 + 1 + 1$ . By convention we let  $a_0 = 1$ . Let  $b_n = \sum_{k=0}^n a_k$  be the cumulative sum of the first  $a$ 's.
- Make a table of the  $a$ 's and  $b$ 's up through  $n = 10$ . What amazing relation do you observe in your table? (Don't prove it yet.)
  - Express the generating function  $A(z)$  as an infinite product.
  - Use the expression from part (b) to prove the result of part (a).
- 38 Find a closed form for the double generating function

$$M(w, z) = \sum_{m, n \geq 0} \min(m, n) w^m z^n.$$

Generalize your answer to obtain, for fixed  $m \geq 2$ , a closed form for

$$M(z_1, \dots, z_m) = \sum_{n_1, \dots, n_m \geq 0} \min(n_1, \dots, n_m) z_1^{n_1} \dots z_m^{n_m}.$$

- 39 Given positive integers  $m$  and  $n$ , find closed forms for

$$\sum_{1 \leq k_1 < k_2 < \dots < k_m \leq n} k_1 k_2 \dots k_m \quad \text{and} \quad \sum_{1 \leq k_1 \leq k_2 \leq \dots \leq k_m \leq n} k_1 k_2 \dots k_m.$$

(For example, when  $m = 2$  and  $n = 3$  the sums are  $1 \cdot 2 + 1 \cdot 3 + 2 \cdot 3$  and  $1 \cdot 1 + 1 \cdot 2 + 1 \cdot 3 + 2 \cdot 2 + 2 \cdot 3 + 3 \cdot 3$ .) *Hint:* What are the coefficients of  $z^m$  in the generating functions  $(1 + a_1 z) \dots (1 + a_n z)$  and  $1/(1 - a_1 z) \dots (1 - a_n z)$ ?

- 40 Express  $\sum_k \binom{n}{k} (kF_{k-1} - F_k)(n - k)_i$  in closed form.
- 41 An *up-down permutation* of order  $n$  is an arrangement  $a_1 a_2 \dots a_n$  of the integers  $\{1, 2, \dots, n\}$  that goes alternately up and down:

$$a_1 < a_2 > a_3 < a_4 > \dots.$$

For example, 35142 is an up-down permutation of order 5. If  $A_n$  denotes the number of up-down permutations of order  $n$ , show that the exponential generating function of  $\langle A_n \rangle$  is  $(1 + \sin z)/\cos z$ .

- 42 A space probe has discovered that organic material on Mars has DNA composed of five symbols, denoted by  $(a, b, c, d, e)$ , instead of the four components in earthling DNA. The four pairs  $cd$ ,  $ce$ ,  $ed$ , and  $ee$  never occur consecutively in a string of Martian DNA, but any string without forbidden pairs is possible. (Thus  $bbcda$  is forbidden but  $bbdca$  is OK.) How many Martian DNA strings of length  $n$  are possible? (When  $n = 2$  the answer is 21, because the left and right ends of a string are distinguishable.)

- 43 The *Newtonian generating function* of a sequence  $\langle g_n \rangle$  is defined to be

$$\dot{G}(z) = \sum_n g_n \binom{z}{n}.$$

Find a convolution formula that defines the relation between sequences  $\langle f_n \rangle$ ,  $\langle g_n \rangle$ , and  $\langle h_n \rangle$  whose Newtonian generating functions are related by the equation  $\dot{F}(z)\dot{G}(z) = \dot{H}(z)$ . Try to make your formula as simple and symmetric as possible.

- 44 Let  $q_n$  be the number of possible outcomes when  $n$  numbers  $\{x_1, \dots, x_n\}$  are compared with each other. For example,  $q_3 = 13$  because the possibilities are

$$\begin{aligned} x_1 < x_2 < x_3; & \quad x_1 < x_2 = x_3; & \quad x_1 < x_3 < x_2; & \quad x_1 = x_2 < x_3; \\ x_1 = x_2 = x_3; & \quad x_1 = x_3 < x_2; & \quad x_2 < x_1 < x_3; \\ x_2 < x_1 = x_3; & \quad x_2 < x_3 < x_1; & \quad x_2 = x_3 < x_1; \\ x_3 < x_1 < x_2; & \quad x_3 < x_1 = x_2; & \quad x_3 < x_2 < x_1. \end{aligned}$$

Find a closed form for the egf  $\hat{Q}(z) = \sum_n q_n z^n/n!$ . Also find sequences  $\langle a_n \rangle$ ,  $\langle b_n \rangle$ ,  $\langle c_n \rangle$  such that

$$q_n = \sum_{k \geq 0} k^n a_k = \sum_k \left\{ \begin{matrix} n \\ k \end{matrix} \right\} b_k = \sum_k \left\langle \begin{matrix} n \\ k \end{matrix} \right\rangle c_k, \quad \text{for all } n > 0.$$

- 45 Evaluate  $\sum_{m,n>0} [m \perp n]/m^2 n^2$ .

- 46 Evaluate

$$\sum_{0 \leq k \leq n/2} \binom{n-2k}{k} \left( \frac{-4}{27} \right)^k$$

in closed form. *Hint:*  $z^3 - z^2 + \frac{4}{27} = (z + \frac{1}{3})(z - \frac{2}{3})^2$ .

- 47 Show that the numbers  $U_n$  and  $V_n$  of  $3 \times n$  domino tilings, as given in (7.34), are closely related to the fractions in the Stern–Brocot tree that converge to  $\sqrt{3}$ .
- 48 A certain sequence  $\langle g_n \rangle$  satisfies the recurrence

$$ag_n + bg_{n+1} + cg_{n+2} + d = 0, \quad \text{integer } n \geq 0,$$

for some integers  $(a, b, c, d)$  with  $\gcd(a, b, c, d) = 1$ . It also has the closed form

$$g_n = \lfloor \alpha(1 + \sqrt{2})^n \rfloor, \quad \text{integer } n \geq 0,$$

for some real number  $\alpha$  between 0 and 1. Find  $a, b, c, d$ , and  $\alpha$ .

- 49 This is a problem about powers and parity.

*Kissinger, take note.*

- a Consider the sequence  $\langle a_0, a_1, a_2, \dots \rangle = \langle 2, 2, 6, \dots \rangle$  defined by the formula

$$a_n = (1 + \sqrt{2})^n + (1 - \sqrt{2})^n.$$

Find a simple recurrence relation that is satisfied by this sequence.

- b Prove that  $\lceil (1 + \sqrt{2})^n \rceil \equiv n \pmod{2}$  for all integers  $n > 0$ .  
 c Find a number  $\alpha$  of the form  $(p + \sqrt{q})/2$ , where  $p$  and  $q$  are positive integers, such that  $\lfloor \alpha^n \rfloor \equiv n \pmod{2}$  for all integers  $n > 0$ .

### Bonus problems

- 50 Continuing exercise 22, consider the sum of all ways to decompose polygons into polygons:

$$Q = \text{---} + \triangle + \square + \begin{array}{|c|} \hline \diagup \\ \hline \end{array} + \begin{array}{|c|} \hline \diagdown \\ \hline \end{array} \\ + \text{pentagon} + \text{pentagon with 1 diagonal} + \text{pentagon with 2 diagonals} + \text{pentagon with 3 diagonals} + \text{pentagon with 4 diagonals} + \text{pentagon with 5 diagonals} + \dots$$

Find a symbolic equation for  $Q$  and use it to find a generating function for the number of ways to draw nonintersecting diagonals inside a convex  $n$ -gon. (Give a closed form for the generating function as a function of  $z$ ; you need not find a closed form for the coefficients.)

- 51 Prove that the product

$$2^{mn/2} \prod_{\substack{1 \leq j \leq m \\ 1 \leq k \leq n}} \left( \left( \cos^2 \frac{j\pi}{m+1} \right) \square^2 + \left( \cos^2 \frac{k\pi}{n+1} \right) \square^2 \right)^{1/4}$$

is the generating function for tilings of an  $m \times n$  rectangle with dominoes. (There are  $mn$  factors, which we can imagine are written in the  $mn$  cells of the rectangle. If  $mn$  is odd, the middle factor is zero. The coefficient of  $\square^j \square^k$  is the number of ways to do the tiling with  $j$  vertical and  $k$  horizontal dominoes.) *Hint:* This is a difficult problem, really beyond the scope of this book. You may wish to simply verify the formula in the case  $m = 3$ ,  $n = 4$ .

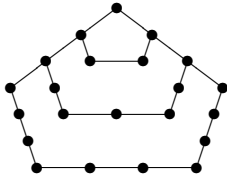
*Is this a hint or a warning?*

- 52 Prove that the polynomials defined by the recurrence

$$p_n(y) = \left(y - \frac{1}{4}\right)^n - \sum_{k=0}^{n-1} \binom{2n}{2k} \left(\frac{-1}{4}\right)^{n-k} p_k(y), \quad \text{integer } n \geq 0,$$

have the form  $p_n(y) = \sum_{m=0}^n |n_m| y^m$ , where  $|n_m|$  is a positive integer for  $1 \leq m \leq n$ . *Hint:* This exercise is very instructive but not very easy.

- 53 The sequence of *pentagonal numbers*  $\langle 1, 5, 12, 22, \dots \rangle$  generalizes the triangular and square numbers in an obvious way:



Let the  $n$ th triangular number be  $T_n = n(n+1)/2$ ; let the  $n$ th pentagonal number be  $P_n = n(3n-1)/2$ ; and let  $U_n$  be the  $3 \times n$  domino-tiling number defined in (7.38). Prove that the triangular number  $T_{(U_{4n+2}-1)/2}$  is also a pentagonal number. *Hint:*  $3U_{2n}^2 = (V_{2n-1} + V_{2n+1})^2 + 2$ .

- 54 Consider the following curious construction:

|   |   |    |    |   |    |    |    |    |    |     |     |     |    |    |      |     |
|---|---|----|----|---|----|----|----|----|----|-----|-----|-----|----|----|------|-----|
| 1 | 2 | 3  | 4  | 5 | 6  | 7  | 8  | 9  | 10 | 11  | 12  | 13  | 14 | 15 | 16   | ... |
| 1 | 2 | 3  | 4  |   | 6  | 7  | 8  | 9  |    | 11  | 12  | 13  | 14 |    | 16   | ... |
| 1 | 3 | 6  | 10 |   | 16 | 23 | 31 | 40 |    | 51  | 63  | 76  | 90 |    | 106  | ... |
| 1 | 3 | 6  |    |   | 16 | 23 | 31 |    |    | 51  | 63  | 76  |    |    | 106  | ... |
| 1 | 4 | 10 |    |   | 26 | 49 | 80 |    |    | 131 | 194 | 270 |    |    | 376  | ... |
| 1 | 4 |    |    |   | 26 | 49 |    |    |    | 131 | 194 |     |    |    | 376  | ... |
| 1 | 5 |    |    |   | 31 | 80 |    |    |    | 211 | 405 |     |    |    | 781  | ... |
| 1 |   |    |    |   | 31 |    |    |    |    | 211 |     |     |    |    | 781  | ... |
| 1 |   |    |    |   | 32 |    |    |    |    | 243 |     |     |    |    | 1024 | ... |

(Start with a row containing all the positive integers. Then delete every  $m$ th column; here  $m = 5$ . Then replace the remaining entries by partial sums. Then delete every  $(m-1)$ st column. Then replace with partial sums again, and so on.) Use generating functions to show that the final result is the sequence of  $m$ th powers. For example, when  $m = 5$  we get  $\langle 1^5, 2^5, 3^5, 4^5, \dots \rangle$  as shown.

- 55 Prove that if the power series  $F(z)$  and  $G(z)$  are differentiably finite (as defined in exercise 20), then so are  $F(z) + G(z)$  and  $F(z)G(z)$ .

### Research problems

- 56 Prove that there is no “simple closed form” for the coefficient of  $z^n$  in  $(1 + z + z^2)^n$ , as a function of  $n$ , in some large class of “simple closed forms.”
- 57 Prove or disprove: If all the coefficients of  $G(z)$  are either 0 or 1, and if all the coefficients of  $G(z)^2$  are less than some constant  $M$ , then infinitely many of the coefficients of  $G(z)^2$  are zero.