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Combinatorial optimization:
Sparse max-flow (C++)
Min-cost max-flow (C++)
Push-relabel max-flow (C++)
Min-cost matching (C++)
Max bipartite matching (C++)
Max Bipartite Matching(Hopcroft-Karp)(C++)
Global min cut (C++)
Geometry:
Convex hull(Monotone chain) (C++)
Convex hull(Graham scan) (C++)
Points in ConvexPolygon(C++)
Polygon Area(C++)
ClosestPair(C++)
CircleIntersection Area(C++)
SegmentIntersection(C++)
Miscellaneous geometry (C++)
Slow Delaunay triangulation (C++)
Numerical algorithms:
Number theoretic algorithms (modular, Chinese remainder, linear Diophantine) (C++)
Systems of linear equations, matrix inverse, determinant (C++)
Reduced row echelon form, matrix rank (C++)
Fast Fourier transform (C++)
Simplex algorithm (C++)
Graph algorithms:
_____
Fast Dijkstra's algorithm (C++)
Strongly connected components (C)
Data structures:
_____
Suffix arrays (C++)
Binary Indexed Tree
Union-Find Set (C/C++)
KD-tree (C++)
Miscellaneous:
______
Longest increasing subsequence (C++)
Dates (C++)
Regular expressions (Java)
Prime numbers (C++)
Knuth-Morris-Pratt (C++)
LCA-array(C++)
LCA-dfs(C++)
SegmentTree(Range Update)(C++)
SegmentTree(2-D) (C++)
Trie-2D(C++)
Heavy-Light-Decomposition(C++)
BipartiteMatching(C++)
2-SAT(C++)
MatrixExpo(C++)
Bitmask(subset of a mask) (C++)
Base 4 Use in DP(C++)
Histrogram Problem(C++)
Matching Pattern(Kmp+DP)(C++)
Maximal Sub-Rectangle(C++)
RectangleUnion(C++)
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Dinic.cc 1/27
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```
// Adjacency list implementation of Dinic's blocking flow algorithm.
// This is very fast in practice, and only loses to push-relabel flow.
      - source
      - sink
      - maximum flow value
      - To obtain the actual flow values, look at all edges with
#include <cmath>
#include <vector>
#include <iostream>
#include <queue>
using namespace std;
const int INF = 2000000000;
struct Edge
    int from, to, cap, flow, index;
    Edge(int from, int to, int cap, int flow, int index) :
    from(from), to(to), cap(cap), flow(flow), index(index) {}
};
struct Dinic
    int N;
    vector<vector<Edge> > G;
    vector<Edge *> dad;
    vector<int> Q;
    Dinic(int N) : N(N), G(N), dad(N), Q(N) {}
    void AddEdge(int from, int to, int cap)
        // For Directed Graph....!!
        G[from].push_back(Edge(from, to, cap, 0, G[to].size()));
        if (from == to) G[from].back().index++;
        G[to].push_back(Edge(to, from, 0, 0, G[from].size() - 1));
    }
    long long BlockingFlow(int s, int t)
        fill(dad.begin(), dad.end(), (Edge *) NULL);
        dad[s] = &G[0][0] - 1;
        int head = 0, tail = 0;
        Q[tail++] = s;
        while (head < tail)</pre>
            int x = Q[head++];
            for (int i = 0; i < G[x].size(); i++)</pre>
                Edge &e = G[x][i];
                if (!dad[e.to] && e.cap - e.flow > 0)
                    dad[e.to] = &G[x][i];
                    Q[tail++] = e.to;
```

```
if (!dad[t]) return 0;
        long long totflow = 0;
        for (int i = 0; i < G[t].size(); i++)</pre>
           Edge *start = &G[G[t][i].to][G[t][i].index];
           int amt = INF;
           for (Edge *e = start; amt && e != dad[s]; e = dad[e->from])
               if (!e)
                {
                   amt = 0;
                   break;
               amt = min(amt, e->cap - e->flow);
           if (amt == 0) continue;
           for (Edge *e = start; amt && e != dad[s]; e = dad[e->from])
               e->flow += amt;
               G[e->to][e->index].flow -= amt;
           totflow += amt;
       return totflow;
    }
   long long GetMaxFlow(int s, int t)
       long long totflow = 0;
       while (long long flow = BlockingFlow(s, t))
           totflow += flow;
       return totflow;
};
MinCostMaxFlow.cc 2/27
// Implementation of min cost max flow algorithm using adjacency
// matrix (Edmonds and Karp 1972). This implementation keeps track of
// cap[j][i]). For a regular max flow, set all edge costs to 0.
// Running time, O(|V|^2) cost per augmentation
      min cost max flow: O(|V|^4 * MAX_EDGE_COST) augmentations
      - sink
      - (maximum flow value, minimum cost value)
      - To obtain the actual flow, look at positive values only.
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
```

```
typedef vector<VI> VVI;
typedef long long L;
typedef vector<L> VL;
typedef vector<VL> VVL;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
const L INF = numeric_limits<L>::max() / 4;
struct MinCostMaxFlow
   int N;
   VVL cap, flow, cost;
   VI found;
   VL dist, pi, width;
   VPII dad;
   MinCostMaxFlow(int N) :
        N(N), cap(N, VL(N)), flow(N, VL(N)), cost(N, VL(N)),
        found(N), dist(N), pi(N), width(N), dad(N) {}
   void AddEdge(int from, int to, L cap, L cost)
       this->cap[from][to] = cap;
       this->cost[from][to] = cost;
   void Relax(int s, int k, L cap, L cost, int dir)
       L val = dist[s] + pi[s] - pi[k] + cost;
       if (cap && val < dist[k])</pre>
           dist[k] = val;
           dad[k] = make_pair(s, dir);
           width[k] = min(cap, width[s]);
    }
   L Dijkstra(int s, int t)
       fill(found.begin(), found.end(), false);
       fill(dist.begin(), dist.end(), INF);
       fill(width.begin(), width.end(), 0);
       dist[s] = 0;
        width[s] = INF;
        while (s != -1)
            int best = -1;
            found[s] = true;
            for (int k = 0; k < N; k++)
                if (found[k]) continue;
                Relax(s, k, cap[s][k] - flow[s][k], cost[s][k], 1);
                Relax(s, k, flow[k][s], -cost[k][s], -1);
                if (best == -1 || dist[k] < dist[best]) best = k;</pre>
            s = best;
        }
        for (int k = 0; k < N; k++)
           pi[k] = min(pi[k] + dist[k], INF);
        return width[t];
    }
   pair<L, L> GetMaxFlow(int s, int t)
```

```
L totflow = 0, totcost = 0;
        while (L amt = Dijkstra(s, t))
            totflow += amt;
            for (int x = t; x != s; x = dad[x].first)
                if (dad[x].second == 1)
                    flow[dad[x].first][x] += amt;
                    totcost += amt * cost[dad[x].first][x];
                else
                {
                    flow[x][dad[x].first] -= amt;
                    totcost -= amt * cost[x][dad[x].first];
            }
       return make_pair(totflow, totcost);
    }
};
PushRelabel.cc 3/27
===========
// Adjacency list implementation of FIFO push relabel maximum flow
// achieve the worst-case.
       - source
       - To obtain the actual flow values, look at all edges with
#include <cmath>
#include <vector>
#include <iostream>
#include <queue>
using namespace std;
typedef long long LL;
struct Edge
    int from, to, cap, flow, index;
    Edge(int from, int to, int cap, int flow, int index) :
    from(from), to(to), cap(cap), flow(flow), index(index) {}
};
struct PushRelabel
    int N;
    vector<vector<Edge> > G;
```

```
vector<LL> excess;
vector<int> dist, active, count;
queue<int> Q;
PushRelabel(int N) : N(N), G(N), excess(N), dist(N), active(N), count(2*N) {}
void AddEdge(int from, int to, int cap)
    G[from].push_back(Edge(from, to, cap, 0, G[to].size()));
    if (from == to) G[from].back().index++;
    G[to].push_back(Edge(to, from, 0, 0, G[from].size() - 1));
}
void Enqueue(int v)
    if (!active[v] && excess[v] > 0)
       active[v] = true;
       Q.push(v);
}
void Push(Edge &e)
    int amt = int(min(excess[e.from], LL(e.cap - e.flow)));
   if (dist[e.from] <= dist[e.to] | amt == 0) return;</pre>
    e.flow += amt;
   G[e.to][e.index].flow -= amt;
   excess[e.to] += amt;
   excess[e.from] -= amt;
    Enqueue(e.to);
}
void Gap(int k)
    for (int v = 0; v < N; v++)
        if (dist[v] < k) continue;</pre>
        count[dist[v]]--;
        dist[v] = max(dist[v], N+1);
        count[dist[v]]++;
        Enqueue(v);
}
void Relabel(int v)
    count[dist[v]]--;
    dist[v] = 2*N;
    for (int i = 0; i < G[v].size(); i++)</pre>
        if (G[v][i].cap - G[v][i].flow > 0)
            dist[v] = min(dist[v], dist[G[v][i].to] + 1);
    count[dist[v]]++;
    Enqueue(v);
}
void Discharge(int v)
    for (int i = 0; excess[v] > 0 && i < G[v].size(); i++) Push(G[v][i]);
    if (excess[v] > 0)
        if (count[dist[v]] == 1)
            Gap(dist[v]);
        else
            Relabel(v);
```

```
}
    LL GetMaxFlow(int s, int t)
        count[0] = N-1;
        count[N] = 1;
        dist[s] = N;
        active[s] = active[t] = true;
        for (int i = 0; i < G[s].size(); i++)</pre>
            excess[s] += G[s][i].cap;
            Push(G[s][i]);
        while (!Q.empty())
            int v = Q.front();
            Q.pop();
            active[v] = false;
            Discharge(v);
        LL totflow = 0;
        for (int i = 0; i < G[s].size(); i++) totflow += G[s][i].flow;</pre>
        return totflow;
    }
};
MinCostMatching.cc 4/27
// Min cost bipartite matching via shortest augmenting paths
// This is an O(n^3) implementation of a shortest augmenting path
\//\ algorithm for finding min cost perfect matchings in dense
// graphs. In practice, it solves 1000x1000 problems in around 1
    cost[i][j] = cost for pairing left node i with right node j
    Lmate[i] = index of right node that left node i pairs with
    Rmate[j] = index of left node that right node j pairs with
// The values in cost[i][j] may be positive or negative. To perform
// maximization, simply negate the cost[][] matrix.
#include <algorithm>
#include <cstdio>
#include <cmath>
#include <vector>
using namespace std;
typedef vector<double> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
double MinCostMatching(const VVD &cost, VI &Lmate, VI &Rmate)
    int n = int(cost.size());
    // construct dual feasible solution
    VD u(n);
    VD v(n);
    for (int i = 0; i < n; i++)</pre>
```

```
u[i] = cost[i][0];
    for (int j = 1; j < n; j++) u[i] = min(u[i], cost[i][j]);</pre>
for (int j = 0; j < n; j++)</pre>
   v[j] = cost[0][j] - u[0];
    for (int i = 1; i < n; i++) v[j] = min(v[j], cost[i][j] - u[i]);</pre>
//\ {\tt construct\ primal\ solution\ satisfying\ complementary\ slackness}
Lmate = VI(n, -1);
Rmate = VI(n, -1);
int mated = 0;
for (int i = 0; i < n; i++)</pre>
    for (int j = 0; j < n; j++)
        if (Rmate[j] != -1) continue;
        if (fabs(cost[i][j] - u[i] - v[j]) < 1e-10)</pre>
            Lmate[i] = j;
            Rmate[j] = i;
            mated++;
            break;
   }
}
VD dist(n);
VI dad(n);
VI seen(n);
// repeat until primal solution is feasible
while (mated < n)</pre>
{
    int s = 0;
    while (Lmate[s] != -1) s++;
    // initialize Dijkstra
    fill(dad.begin(), dad.end(), -1);
    fill(seen.begin(), seen.end(), 0);
    for (int k = 0; k < n; k++)</pre>
        dist[k] = cost[s][k] - u[s] - v[k];
    int j = 0;
    while (true)
        j = -1;
        for (int k = 0; k < n; k++)
            if (seen[k]) continue;
            if (j == -1 || dist[k] < dist[j]) j = k;</pre>
        seen[j] = 1;
        if (Rmate[j] == -1) break;
        // relax neighbors
        const int i = Rmate[j];
        for (int k = 0; k < n; k++)
            if (seen[k]) continue;
```

```
if (dist[k] > new_dist)
                   dist[k] = new_dist;
                   dad[k] = j;
           }
        // update dual variables
        for (int k = 0; k < n; k++)
           if (k == j || !seen[k]) continue;
           const int i = Rmate[k];
           v[k] += dist[k] - dist[j];
           u[i] -= dist[k] - dist[j];
        u[s] += dist[j];
        while (dad[j] >= 0)
           const int d = dad[j];
           Rmate[j] = Rmate[d];
           Lmate[Rmate[j]] = j;
           j = d;
        Rmate[j] = s;
       Lmate[s] = j;
       mated++;
    }
    double value = 0;
    for (int i = 0; i < n; i++)</pre>
        value += cost[i][Lmate[i]];
    return value;
MaxBipartiteMatching.cc 5/27
_____
    INPUT: w[i][j] = edge between row node i and column node j
    OUTPUT: mr[i] = assignment for row node i, -1 if unassigned
#include <vector>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool FindMatch(int i, const VVI &w, VI &mr, VI &mc, VI &seen)
    for (int j = 0; j < w[i].size(); j++)</pre>
        if (w[i][j] && !seen[j])
```

const double new_dist = dist[j] + cost[i][k] - u[i] - v[k];

```
seen[j] = true;
            if (mc[j] < 0 | FindMatch(mc[j], w, mr, mc, seen))</pre>
                mr[i] = j;
                mc[j] = i;
                return true;
            }
       }
   return false;
int BipartiteMatching(const VVI &w, VI &mr, VI &mc)
    mr = VI(w.size(), -1);
    mc = VI(w[0].size(), -1);
    int ct = 0;
    for (int i = 0; i < w.size(); i++)</pre>
        VI seen(w[0].size());
       if (FindMatch(i, w, mr, mc, seen)) ct++;
    return ct;
}
MaximumBipartite Matching(Hopcroft-Karp):
_____
n: number of nodes on left side, nodes are numbered 1 to n
m: number of nodes on right side, nodes are numbered n+1 to n+m
G = NIL[0] ? G1[G[1--n]] ? G2[G[n+1--n+m]]
bool bfs() {
   int i, u, v, len;
    queue< int > Q;
    for(i=1; i<=n; i++) {</pre>
        if(match[i]==NIL) {
           dist[i] = 0;
           Q.push(i);
        else dist[i] = INF;
    dist[NIL] = INF;
    while(!Q.empty()) {
        u = Q.front(); Q.pop();
        if(u!=NIL) {
            len = G[u].size();
            for(i=0; i<len; i++) {</pre>
                v = G[u][i];
                if(dist[match[v]] == INF) {
                    dist[match[v]] = dist[u] + 1;
                    Q.push(match[v]);
           }
    return (dist[NIL]!=INF);
bool dfs(int u) {
    int i, v, len;
    if(u!=NIL) {
        len = G[u].size();
        for(i=0; i<len; i++) {</pre>
```

```
v = G[u][i];
            if(dist[match[v]]==dist[u]+1) {
                if(dfs(match[v])) {
                   match[v] = u;
                    match[u] = v;
                   return true;
                }
            }
        dist[u] = INF;
       return false;
    return true;
}
int hopcroft_karp() {
   int matching = 0, i;
   CLR(match);
    while(bfs())
       for(i=1; i<=n; i++)</pre>
           if(match[i] == NIL && dfs(i))
               matching++;
    return matching;
}
MinCut.cc 6/27
-----
// Adjacency matrix implementation of Stoer-Wagner min cut algorithm.
// - (min cut value, nodes in half of min cut)
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
const int INF = 1000000000;
pair<int, VI> GetMinCut(VVI &weights)
    int N = weights.size();
   VI used(N), cut, best_cut;
    int best_weight = -1;
    for (int phase = N-1; phase >= 0; phase--)
       VI w = weights[0];
       VI added = used;
       int prev, last = 0;
        for (int i = 0; i < phase; i++)</pre>
           prev = last;
           last = -1;
            for (int j = 1; j < N; j++)</pre>
```

```
if (!added[j] && (last == -1 | | w[j] > w[last])) last = j;
            if (i == phase-1)
                for (int j = 0; j < N; j++) weights[prev][j] += weights[last][j];</pre>
                for (int j = 0; j < N; j++) weights[j][prev] = weights[prev][j];</pre>
                used[last] = true;
                cut.push_back(last);
                if (best_weight == -1 || w[last] < best_weight)</pre>
                    best_cut = cut;
                    best_weight = w[last];
            }
            else
                for (int j = 0; j < N; j++)</pre>
                    w[j] += weights[last][j];
                added[last] = true;
    return make_pair(best_weight, best_cut);
ConvexHull(Monotone chain).cc 7/27
// Compute the 2D convex hull of a set of points using the monotone chain
// #defined.
   INPUT: a vector of input points, unordered.
   OUTPUT: a vector of points in the convex hull, counterclockwise, starting
              with bottommost/leftmost point
#include <cstdio>
#include <cassert>
#include <vector>
#include <algorithm>
#include <cmath>
using namespace std;
#define REMOVE_REDUNDANT
typedef double T;
const T EPS = 1e-7;
struct PT
    T x, y;
    PT() {}
    PT(T x, T y) : x(x), y(y) {}
    bool operator<(const PT &rhs) const</pre>
        return make_pair(y,x) < make_pair(rhs.y,rhs.x);</pre>
    bool operator==(const PT &rhs) const
        return make_pair(y,x) == make_pair(rhs.y,rhs.x);
};
T cross(PT p, PT q)
```

```
return p.x*q.y-p.y*q.x;
}
T area2(PT a, PT b, PT c)
   return cross(a,b) + cross(b,c) + cross(c,a);
#ifdef REMOVE_REDUNDANT
bool between (const PT &a, const PT &b, const PT &c)
   #endif
void ConvexHull(vector<PT> &pts)
   sort(pts.begin(), pts.end());
   pts.erase(unique(pts.begin(), pts.end()), pts.end());
   vector<PT> up, dn;
   for (int i = 0; i < pts.size(); i++)</pre>
       while (up.size() > 1 && area2(up[up.size()-2], up.back(), pts[i]) >= 0) up.pop_back();
       while (dn.size() > 1 && area2(dn[dn.size()-2], dn.back(), pts[i]) <= 0) dn.pop_back();</pre>
       up.push_back(pts[i]);
       dn.push_back(pts[i]);
   pts = dn;
   for (int i = (int) up.size() - 2; i >= 1; i--) pts.push_back(up[i]);
#ifdef REMOVE_REDUNDANT
   if (pts.size() <= 2) return;</pre>
   dn.clear();
   dn.push_back(pts[0]);
   dn.push_back(pts[1]);
   for (int i = 2; i < pts.size(); i++)</pre>
       if (between(dn[dn.size()-2], dn[dn.size()-1], pts[i])) dn.pop_back();
       dn.push_back(pts[i]);
   if (dn.size() >= 3 \&\& between(dn.back(), dn[0], dn[1]))
       dn[0] = dn.back();
       dn.pop_back();
   pts = dn;
#endif
ConvexHull(Graham scan):
ConvexHull : Graham's Scan O(n lg n), integer implementation
P[]: holds all the points, C[]: holds points on the hull
np: number of points in P[], nc: number of points in C[]
to handle duplicate, call makeUnique() before calling convexHull()
call convexHull() if you have np >= 3
to remove co-linear points on hull, call compress() after convexHull()
* /
point P[MAX], C[MAX], P0;
inline int triArea2(const point &a, const point &b, const point &c) {
   return (a.x*(b.y-c.y) + b.x*(c.y-a.y) + c.x*(a.y-b.y));
inline int sqDist(const point &a, const point &b) {
```

```
return ((a.x-b.x)*(a.x-b.x) + (a.y-b.y)*(a.y-b.y));
inline bool comp(const point &a, const point &b) {
    int d = triArea2(P0, a, b);
    if(d < 0) return false;</pre>
    if(!d && sqDist(P0, a) > sqDist(P0, b)) return false;
    return true;
inline bool normal(const point &a, const point &b) {
    return ((a.x==b.x) ? a.y < b.y : a.x < b.x);</pre>
inline bool issame(const point &a, const point &b) {
    return (a.x == b.x && a.y == b.y);
inline void makeUnique(int &np) {
    sort(&P[0], &P[np], normal);
    np = unique(&P[0], &P[np], issame) - P;
void convexHull(int &np, int &nc) {
    int i, j, pos = 0;
    for(i = 1; i < np; i++)</pre>
        if(P[i].y<P[pos].y | | (P[i].y==P[pos].y && P[i].x<P[pos].x))</pre>
            pos = i;
    swap(P[0], P[pos]);
    P0 = P[0];
    sort(&P[1], &P[np], comp);
    for(i = 0; i < 3; i++) C[i] = P[i];</pre>
    for(i = j = 3; i < np; i++) {</pre>
        while(triArea2(C[j-2], C[j-1], P[i]) < 0) j--;
        C[j++] = P[i];
    nc = j;
}
void compress(int &nc) {
    int i, j, d;
    C[nc] = C[0];
    for(i=j=1; i < nc; i++) {</pre>
        d = triArea2(C[j-1], C[i], C[i+1]);
        if(d \mid | (!d \&\& issame(C[j-1], C[i+1]))) C[j++] = C[i];
    nc = j;
}
Points In convexPolygon:
C[] array of points of convex polygon in ccw order,
nc number of points in C, p target points.
returns true if p is inside C (including edge) or false otherwise.
complexity O(lg n)
* /
inline bool inConvexPoly(point *C, int nc, const point &p) {
    int st = 1, en = nc - 1, mid;
    while(en - st > 1) {
       mid = (st + en) >> 1;
        if(triArea2(C[0], C[mid], p) < 0) en = mid;</pre>
        else st = mid;
    if(triArea2(C[0], C[st], p) < 0) return false;</pre>
```

```
if(triArea2(C[st], C[en], p) < 0) return false;</pre>
    if(triArea2(C[en], C[0], p) < 0) return false;</pre>
    return true;
Ploygon Area:
===========
P[\,] holds the points, must be either in cw or ccw
function returns double of the area.
inline int dArea(int np) {
   int area = 0;
    for(int i = 0; i < np; i++) {</pre>
        area += p[i].x*p[i+1].y - p[i].y*p[i+1].x;
    return abs(area);
ClosestPair:
closestPair(Point *X, Point *Y, int n);
X contains the points sorted by x co-ordinate,
Y contains the points sorted by y co-ordinate,
One additional item in Point structure is needed, the original index.
* /
typedef long long i64;
typedef struct { int x, y, i; } Point;
int flag[MAX];
inline i64 sq(const i64 &x) {
    return x*x;
inline i64 sqdist(const Point &a, const Point &b) {
    return sq(a.x-b.x) + sq(a.y-b.y);
inline i64 closestPair(Point *X, Point *Y, int n) {
    if(n == 1) return INF;
    if(n == 2) return sqdist(X[0], X[1]);
    int i, j, k, n1, n2, ns, m = n >> 1;
    Point Xm = X[m-1], *XL, *XR, *YL, *YR, *YS;
    i64 lt, rt, dd, tmp;
    XL = new Point[m], YL = new Point[m];
    XR = new Point[m+1], YR = new Point[m+1];
    YS = new Point[n];
    for(i = 0; i < m; i++) XL[i] = X[i], flag[X[i].i] = 0;</pre>
    for(; i < n; i++) XR[i - m] = X[i], flag[X[i].i] = 1;</pre>
    for(i = n2 = n1 = 0; i < n; i++) {</pre>
        if(!flag[Y[i].i]) YL[n1++] = Y[i];
        else YR[n2++] = Y[i];
    }
    lt = closestPair(XL, YL, n1);
    rt = closestPair(XR, YR, n2);
    dd = min(lt, rt);
    for(i = ns = 0; i < n; i++)
```

```
if(sq(Y[i].x - Xm.x) < dd)
           YS[ns++] = Y[i];
    for(j = 0; j < ns; j++)
        for(k = j + 1; k < ns && sq(YS[k].y - YS[j].y) < dd; k++)
           dd = min(dd, sqdist(YS[j], YS[k]));
    delete[] XL; delete[] XR;
    delete[] YL; delete[] YR;
    delete[] YS;
    return dd;
}
CircleIntersection Area:
This code assumes the circle center and radius to be integer.
Change this when necessary.
inline double commonArea(const Circle &a, const Circle &b) {
   int dsq = sqDist(a.c, b.c);
    double d = sqrt((double)dsq);
    if(sq(a.r + b.r) <= dsq) return 0;</pre>
    if(a.r >= b.r \&\& sq(a.r-b.r) >= dsq) return pi * b.r * b.r;
    if(a.r <= b.r && sq(b.r-a.r) >= dsq) return pi * a.r * a.r;
    double angleA = 2.0 * acos((a.r * a.r + dsq - b.r * b.r) / (2.0 * a.r * d));
    double angleB = 2.0 * acos((b.r * b.r + dsq - a.r * a.r) / (2.0 * b.r * d));
    return 0.5 * (a.r * a.r * (angleA - sin(angleA)) + b.r * b.r * (angleB - sin(angleB)));
}
Segment intersection:
Segment intersection in 2D integer space.
P1, p2 makes first segment, p3, p4 makes the second segment
inline bool intersect(const Point &p1, const Point &p2, const Point &p3, const Point &p4) {
   i64 d1, d2, d3, d4;
    d1 = direction(p3, p4, p1);
    d2 = direction(p3, p4, p2);
    d3 = direction(p1, p2, p3);
    d4 = direction(p1, p2, p4);
    if(((d1 < 0 \&\& d2 > 0)) | (d1 > 0 \&\& d2 < 0))
    && ((d3 < 0 && d4 > 0) | (d3 > 0 && d4 < 0))) return true;
    if(!d3 && onsegment(p1, p2, p3)) return true;
    if(!d4 && onsegment(p1, p2, p4)) return true;
    if(!d1 && onsegment(p3, p4, p1)) return true;
    if(!d2 && onsegment(p3, p4, p2)) return true;
    return false;
Geometry.cc 8/27
==========
#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>
using namespace std;
double INF = 1e100;
```

```
struct PT
   double x, y;
   PT() {}
    PT(double x, double y) : x(x), y(y) {}
    PT(const PT &p) : x(p.x), y(p.y)
    PT operator + (const PT &p) const
       return PT(x+p.x, y+p.y);
    }
    PT operator - (const PT &p) const
       return PT(x-p.x, y-p.y);
    PT operator * (double c) const
       return PT(x*c, y*c );
    PT operator / (double c) const
       return PT(x/c, y/c );
};
double dot(PT p, PT q)
   return p.x*q.x+p.y*q.y;
double dist2(PT p, PT q)
    return dot(p-q,p-q);
double cross(PT p, PT q)
    return p.x*q.y-p.y*q.x;
ostream &operator<<(ostream &os, const PT &p)
    os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90(PT p)
   return PT(-p.y,p.x);
PT RotateCW90(PT p)
{
   return PT(p.y,-p.x);
PT RotateCCW(PT p, double t)
    return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c)
    return a + (b-a)*dot(c-a, b-a)/dot(b-a, b-a);
}
// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c)
```

double EPS = 1e-12;

```
double r = dot(b-a,b-a);
    if (fabs(r) < EPS) return a;</pre>
    r = dot(c-a, b-a)/r;
    if (r < 0) return a;</pre>
    if (r > 1) return b;
    return a + (b-a)*r;
// compute distance from \ensuremath{\text{c}} to segment between a and \ensuremath{\text{b}}
double DistancePointSegment(PT a, PT b, PT c)
    return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z, double a, double b, double c, double d)
    return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
// determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d)
    return fabs(cross(b-a, c-d)) < EPS;</pre>
}
bool LinesCollinear(PT a, PT b, PT c, PT d)
    // determine if line segment from a to b intersects with
bool SegmentsIntersect(PT a, PT b, PT c, PT d)
    if (LinesCollinear(a, b, c, d))
        if (dist2(a, c) < EPS || dist2(a, d) < EPS ||</pre>
                dist2(b, c) < EPS | dist2(b, d) < EPS) return true;</pre>
        if (dot(c-a, c-b) > 0 & dot(d-a, d-b) > 0 & dot(c-b, d-b) > 0)
            return false;
        return true;
    if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
    if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
    return true;
// compute intersection of line passing through a and b
PT ComputeLineIntersection(PT a, PT b, PT c, PT d)
    b=b-a;
    d=c-d;
    c=c-a;
    assert(dot(b, b) > EPS && dot(d, d) > EPS);
    return a + b*cross(c, d)/cross(b, d);
}
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c)
    b=(a+b)/2;
```

```
c = (a+c)/2;
    return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
\ensuremath{//} strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an {\rm *exact*} test using
\//\ integer arithmetic by taking care of the division appropriately
\ensuremath{//} (making sure to deal with signs properly) and then by writing exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q)
    bool c = 0;
    for (int i = 0; i < p.size(); i++)</pre>
        int j = (i+1)%p.size();
        if ((p[i].y <= q.y && q.y < p[j].y ||</pre>
                p[j].y \le q.y && q.y < p[i].y) &&
                q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y))
    return c;
}
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q)
    for (int i = 0; i < p.size(); i++)</pre>
        if (dist2(ProjectPointSegment(p[i], p[(i+1)*p.size()], q), q) < EPS)</pre>
            return true;
    return false;
}
// compute intersection of line through points a and b with
// circle centered at c with radius r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r)
    vector<PT> ret;
    b = b-a;
    a = a-c;
    double A = dot(b, b);
    double B = dot(a, b);
    double C = dot(a, a) - r*r;
    double D = B*B - A*C;
    if (D < -EPS) return ret;</pre>
    ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
    if (D > EPS)
        ret.push_back(c+a+b*(-B-sqrt(D))/A);
    return ret;
}
// compute intersection of circle centered at a with radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R)
{
    vector<PT> ret;
    double d = sqrt(dist2(a, b));
    if (d > r+R \mid | d+min(r, R) < max(r, R)) return ret;
    double x = (d*d-R*R+r*r)/(2*d);
    double y = sqrt(r*r-x*x);
    PT v = (b-a)/d;
    ret.push_back(a+v*x + RotateCCW90(v)*y);
    if (y > 0)
        ret.push_back(a+v*x - RotateCCW90(v)*y);
    return ret;
```

```
// This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or
\ensuremath{//} counterclockwise fashion. Note that the centroid is often known as
// the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p)
    double area = 0;
    for(int i = 0; i < p.size(); i++)</pre>
        int j = (i+1) % p.size();
        area += p[i].x*p[j].y - p[j].x*p[i].y;
    return area / 2.0;
}
double ComputeArea(const vector<PT> &p)
    return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p)
    PT c(0,0);
    double scale = 6.0 * ComputeSignedArea(p);
    for (int i = 0; i < p.size(); i++)</pre>
        int j = (i+1) % p.size();
        c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
    return c / scale;
\//\ tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p)
    for (int i = 0; i < p.size(); i++)</pre>
        for (int k = i+1; k < p.size(); k++)</pre>
            int j = (i+1) % p.size();
            int 1 = (k+1) % p.size();
            if (i == 1 || j == k) continue;
            if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
                 return false;
    return true;
int main()
    cerr << RotateCCW90(PT(2,5)) << endl;</pre>
    cerr << RotateCW90(PT(2,5)) << endl;</pre>
    cerr << RotateCCW(PT(2,5),M_PI/2) << endl;</pre>
    cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;</pre>
```

```
cerr << ProjectPointSegment(PT(-5,-2), PT(10,4), PT(3,7)) << " "
     << ProjectPointSegment(PT(7.5,3), PT(10,4), PT(3,7)) << " "</pre>
     << ProjectPointSegment(PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;</pre>
// expected: 6.78903
cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;</pre>
cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "</pre>
     << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "</pre>
     << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;</pre>
cerr << LinesCollinear(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "</pre>
     << LinesCollinear(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "</pre>
     << LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;</pre>
cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << " "
     << SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT(0,5)) << " "</pre>
     << SegmentsIntersect(PT(0,0), PT(2,4), PT(2,-1), PT(-2,1)) << " "
     << SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT(1,7)) << endl;
cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << endl;</pre>
cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;</pre>
vector<PT> v;
v.push back(PT(0,0));
v.push back(PT(5,0));
v.push back(PT(5,5));
v.push_back(PT(0,5));
cerr << PointInPolygon(v, PT(2,2)) << " "</pre>
     << PointInPolygon(v, PT(2,0)) << " "</pre>
     << PointInPolygon(v, PT(0,2)) << " "</pre>
     << PointInPolygon(v, PT(5,2)) << " "</pre>
     << PointInPolygon(v, PT(2,5)) << endl;</pre>
cerr << PointOnPolygon(v, PT(2,2)) << " "</pre>
     << PointOnPolygon(v, PT(2,0)) << " "</pre>
     << PointOnPolygon(v, PT(0,2)) << " "</pre>
     << PointOnPolygon(v, PT(5,2)) << " "</pre>
     << PointOnPolygon(v, PT(2,5)) << endl;</pre>
             (5,4) (4,5)
             (4,5) (5,4)
             (4,5) (5,4)
vector<PT> u = CircleLineIntersection(PT(0,6), PT(2,6), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " ";</pre>
cerr << endl;</pre>
u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " ";</pre>
cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " ";</pre>
cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5);
```

```
cerr << endl;</pre>
    u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 10, sqrt(2.0)/2.0);
    for (int i = 0; i < u.size(); i++) cerr << u[i] << " ";</pre>
    cerr << endl;</pre>
    u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt(2.0)/2.0);
    for (int i = 0; i < u.size(); i++) cerr << u[i] << " ";</pre>
    cerr << endl;</pre>
    // centroid should be (1.1666666, 1.166666)
    PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
    vector<PT> p(pa, pa+4);
    PT c = ComputeCentroid(p);
    cerr << "Area: " << ComputeArea(p) << endl;</pre>
    cerr << "Centroid: " << c << endl;</pre>
    return 0;
}
JavaGeometry.java 9/27
// In this example, we read an input file containing three lines, each
// contains a list of points, p[1], p[2], ...
// Our goal is to determine:
// (1) whether B - A is a single closed shape (as opposed to multiple shapes)
  (2) the area of B - A
    (3) whether each p[i] is in the interior of B - A
  8 6
   The area is singular.
import java.util.*;
import java.awt.geom.*;
import java.io.*;
public class JavaGeometry
    // make an array of doubles from a string
    static double[] readPoints(String s)
        String[] arr = s.trim().split("\s++");
        double[] ret = new double[arr.length];
        for (int i = 0; i < arr.length; i++) ret[i] = Double.parseDouble(arr[i]);</pre>
        return ret;
    }
    // make an Area object from the coordinates of a polygon
    static Area makeArea(double[] pts)
```

for (int i = 0; i < u.size(); i++) cerr << u[i] << " ";</pre>

```
Path2D.Double p = new Path2D.Double();
    p.moveTo(pts[0], pts[1]);
    for (int i = 2; i < pts.length; i += 2) p.lineTo(pts[i], pts[i+1]);</pre>
   p.closePath();
   return new Area(p);
// compute area of polygon
static double computePolygonArea(ArrayList<Point2D.Double> points)
   Point2D.Double[] pts = points.toArray(new Point2D.Double[points.size()]);
    double area = 0;
   for (int i = 0; i < pts.length; i++)</pre>
        int j = (i+1) % pts.length;
       area += pts[i].x * pts[j].y - pts[j].x * pts[i].y;
   return Math.abs(area)/2;
// compute the area of an Area object containing several disjoint polygons
static double computeArea(Area area)
   double totArea = 0;
    PathIterator iter = area.getPathIterator(null);
   ArrayList<Point2D.Double> points = new ArrayList<Point2D.Double>();
    while (!iter.isDone())
       double[] buffer = new double[6];
        switch (iter.currentSegment(buffer))
        case PathIterator.SEG MOVETO:
        case PathIterator.SEG_LINETO:
            points.add(new Point2D.Double(buffer[0], buffer[1]));
            break;
        case PathIterator.SEG_CLOSE:
            totArea += computePolygonArea(points);
            points.clear();
            break;
        iter.next();
    return totArea;
// avoid wrapping the Scanner object for file reading in a
public static void main(String args[]) throws Exception
    Scanner scanner = new Scanner(new File("input.txt"));
    // Scanner scanner = new Scanner (System.in);
    double[] pointsA = readPoints(scanner.nextLine());
    double[] pointsB = readPoints(scanner.nextLine());
    Area areaA = makeArea(pointsA);
   Area areaB = makeArea(pointsB);
    areaB.subtract(areaA);
    // areaB.exclusiveOr (areaA);
    // areaB.add (areaA);
    // areaB.intersect (areaA);
```

```
boolean isSingle = areaB.isSingular();
        // areaB.isEmpty();
        if (isSingle)
           System.out.println("The area is singular.");
        else
           System.out.println("The area is not singular.");
        // (2) compute the area of B - A
        System.out.println("The area is " + computeArea(areaB) + ".");
        // (3) determine whether each p[i] is in the interior of B - A
        while (scanner.hasNextDouble())
           double x = scanner.nextDouble();
            assert(scanner.hasNextDouble());
           double y = scanner.nextDouble();
            if (areaB.contains(x,y))
               System.out.println ("Point belongs to the area.");
            else
               System.out.println ("Point does not belong to the area.");
        }
        // Finally, some useful things we didn't use in this example:
            Ellipse2D.Double ellipse = new Ellipse2D.Double (double x, double y,
                                                             double w. double h);
             creates an ellipse inscribed in box with bottom-left corner (x,y)
              and upper-right corner (x+y,w+h)
           Rectangle2D.Double rect = new Rectangle2D.Double (double x, double y,
                                                               double w, double h);
              creates a box with bottom-left corner (x,y) and upper-right
        // Each of these can be embedded in an Area object (e.g., new Area (rect)).
    }
Geom3D. java 10/27
==========
public class Geom3D
    // distance from point (x, y, z) to plane aX + bY + cZ + d = 0
   public static double ptPlaneDist(double x, double y, double z,
                                    double a, double b, double c, double d)
    {
        return Math.abs(a*x + b*y + c*z + d) / Math.sqrt(a*a + b*b + c*c);
    }
   // distance between parallel planes aX + bY + cZ + d1 = 0 and
    // aX + bY + cZ + d2 = 0
   public static double planePlaneDist(double a, double b, double c,
```

// (1) determine whether B - A is a single closed shape (as

```
{
        return Math.abs(d1 - d2) / Math.sqrt(a*a + b*b + c*c);
    }
    // distance from point (px, py, pz) to line (x1, y1, z1)-(x2, y2, z2)  
    \//\ (or ray, or segment; in the case of the ray, the endpoint is the
    public static final int LINE = 0;
    public static final int SEGMENT = 1;
    public static final int RAY = 2;
    public static double ptLineDistSq(double x1, double y1, double z1,
                                      double x2, double y2, double z2, double px, double py, double pz,
                                      int type)
    {
        double pd2 = (x1-x2)*(x1-x2) + (y1-y2)*(y1-y2) + (z1-z2)*(z1-z2);
        double x, y, z;
        if (pd2 == 0)
            x = x1;
            y = y1;
            z = z1;
        else
            double u = ((px-x1)*(x2-x1) + (py-y1)*(y2-y1) + (pz-z1)*(z2-z1)) / pd2;
            x = x1 + u * (x2 - x1);
            y = y1 + u * (y2 - y1);
            z = z1 + u * (z2 - z1);
            if (type != LINE && u < 0)</pre>
                x = x1;
                y = y1;
                z = z1;
            if (type == SEGMENT && u > 1.0)
                x = x2;
                y = y2;
                z = z2;
        return (x-px)*(x-px) + (y-py)*(y-py) + (z-pz)*(z-pz);
    public static double ptLineDist(double x1, double y1, double z1,
                                    double x2, double y2, double z2, double px, double py, double pz,
                                    int type)
        return Math.sqrt(ptLineDistSq(x1, y1, z1, x2, y2, z2, px, py, pz, type));
Delaunay.cc 11/27
===========
// Slow but simple Delaunay triangulation. Does not handle
// degenerate cases (from O'Rourke, Computational Geometry in C)
// Running time: O(n^4)
// INPUT: x[] = x-coordinates
```

}

double d1, double d2)

```
// OUTPUT: triples = a vector containing m triples of indices
#include<vector>
using namespace std;
typedef double T;
struct triple
    int i, j, k;
    triple() {}
    triple(int i, int j, int k) : i(i), j(j), k(k) {}
};
vector<triple> delaunayTriangulation(vector<T>& x, vector<T>& y)
    int n = x.size();
    vector<T> z(n);
    vector<triple> ret;
    for (int i = 0; i < n; i++)</pre>
        z[i] = x[i] * x[i] + y[i] * y[i];
    for (int i = 0; i < n-2; i++)
        for (int j = i+1; j < n; j++)
            for (int k = i+1; k < n; k++)</pre>
                if (j == k) continue;
                double xn = (y[j]-y[i])*(z[k]-z[i]) - (y[k]-y[i])*(z[j]-z[i]);
                double yn = (x[k]-x[i])*(z[j]-z[i]) - (x[j]-x[i])*(z[k]-z[i]);
                double zn = (x[j]-x[i])*(y[k]-y[i]) - (x[k]-x[i])*(y[j]-y[i]);
                bool flag = zn < 0;</pre>
                for (int m = 0; flag && m < n; m++)</pre>
                    flag = flag && ((x[m]-x[i])*xn +
                                     (y[m]-y[i])*yn +
                                     (z[m]-z[i])*zn <= 0);
                if (flag) ret.push_back(triple(i, j, k));
    return ret;
int main()
    T xs[] = \{0, 0, 1, 0.9\};
    T ys[] = \{0, 1, 0, 0.9\};
    vector<T> x(&xs[0], &xs[4]), y(&ys[0], &ys[4]);
    vector<triple> tri = delaunayTriangulation(x, y);
    int i;
    for(i = 0; i < tri.size(); i++)</pre>
       printf("%d %d %d\n", tri[i].i, tri[i].j, tri[i].k);
    return 0;
}
```

```
// This is a collection of useful code for solving problems that
// involve modular linear equations. Note that all of the
// algorithms described here work on nonnegative integers.
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int,int> PII;
// return a % b (positive value)
int mod(int a, int b)
   return ((a%b)+b)%b;
int gcd(int a, int b)
   int tmp;
   while(b)
       a%=b;
       tmp=a;
       a=b;
       b=tmp;
   return a;
}
// computes lcm(a,b)
int lcm(int a, int b)
   return a/gcd(a,b)*b;
// returns d = gcd(a,b); finds x,y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y)
   int xx = y = 0;
   int yy = x = 1;
   while (b)
       int q = a/b;
       int t = b;
       b = a b;
       a = t;
       t = xx;
       xx = x-q*xx;
       x = t;
       t = yy;
       yy = y-q*yy;
       y = t;
    }
   return a;
}
// finds all solutions to ax = b (mod n)
VI modular_linear_equation_solver(int a, int b, int n)
{
   int x, y;
   VI solutions;
    int d = extended_euclid(a, n, x, y);
```

```
if (!(b%d))
        x = mod (x*(b/d), n);
        for (int i = 0; i < d; i++)</pre>
            solutions.push_back(mod(x + i*(n/d), n));
    return solutions;
}
// computes b such that ab = 1 \pmod{n}, returns -1 on failure
int mod_inverse(int a, int n)
    int x, y;
    int d = extended_euclid(a, n, x, y);
    if (d > 1) return -1;
    return mod(x,n);
// z % x = a, z % y = b. Here, z is unique modulo M = lcm(x,y).
// Return (z,M). On failure, M = -1.
PII chinese_remainder_theorem(int x, int a, int y, int b)
    int s, t;
    int d = extended_euclid(x, y, s, t);
    if (a%d != b%d) return make_pair(0, -1);
    return make_pair(mod(s*b*x+t*a*y,x*y)/d, x*y/d);
}
// z % x[i] = a[i] for all i. Note that the solution is
// unique modulo M = lcm_i (x[i]). Return (z,M). On
// failure, M = -1. Note that we do not require the a[i]'s
PII chinese_remainder_theorem(const VI &x, const VI &a)
    PII ret = make_pair(a[0], x[0]);
    for (int i = 1; i < x.size(); i++)</pre>
        ret = chinese_remainder_theorem(ret.second, ret.first, x[i], a[i]);
        if (ret.second == -1) break;
    return ret;
// computes x and y such that ax + by = c; on failure, x = y = -1
void linear_diophantine(int a, int b, int c, int &x, int &y)
    int d = gcd(a,b);
    if (c%d)
       x = y = -1;
    }
    else
        x = c/d * mod_inverse(a/d, b/d);
        y = (c-a*x)/b;
}
int main()
    cout << gcd(14, 30) << endl;</pre>
```

```
int x, y;
    int d = extended_euclid(14, 30, x, y);
    cout << d << " " << x << " " << y << endl;</pre>
    // expected: 95 45
    VI sols = modular_linear_equation_solver(14, 30, 100);
    for (int i = 0; i < (int) sols.size(); i++) cout << sols[i] << " ";</pre>
    cout << endl;</pre>
    // expected: 8
    cout << mod_inverse(8, 9) << endl;</pre>
    int xs[] = {3, 5, 7, 4, 6};
    int as[] = \{2, 3, 2, 3, 5\};
    PII ret = chinese_remainder_theorem(VI (xs, xs+3), VI(as, as+3));
    cout << ret.first << " " << ret.second << endl;</pre>
    ret = chinese_remainder_theorem (VI(xs+3, xs+5), VI(as+3, as+5));
    cout << ret.first << " " << ret.second << endl;</pre>
    linear_diophantine(7, 2, 5, x, y);
    cout << x << " " << y << endl;</pre>
}
GaussJordan.cc 13/27
============
// Gauss-Jordan elimination with full pivoting.
// (1) solving systems of linear equations (AX=B)
// (2) inverting matrices (AX=I)
   (3) computing determinants of square matrices
// INPUT: a[][] = an nxn matrix
// OUTPUT: X = an nxm matrix (stored in b[][])
             A^{-1} = an nxn matrix (stored in a[][])
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan(VVT &a, VVT &b)
    const int n = a.size();
    const int m = b[0].size();
    VI irow(n), icol(n), ipiv(n);
```

```
T \det = 1;
    for (int i = 0; i < n; i++)</pre>
        int pj = -1, pk = -1;
        for (int j = 0; j < n; j++) if (!ipiv[j])</pre>
                 for (int k = 0; k < n; k++) if (!ipiv[k])
                         if (pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk]))
                              pj = j;
                              pk = k;
        if (fabs(a[pj][pk]) < EPS)</pre>
             cerr << "Matrix is singular." << endl;</pre>
             exit(0);
        ipiv[pk]++;
        swap(a[pj], a[pk]);
        swap(b[pj], b[pk]);
        if (pj != pk) det *= -1;
        irow[i] = pj;
        icol[i] = pk;
        T c = 1.0 / a[pk][pk];
        det *= a[pk][pk];
        a[pk][pk] = 1.0;
        for (int p = 0; p < n; p++) a[pk][p] *= c;</pre>
        for (int p = 0; p < m; p++) b[pk][p] *= c;</pre>
        for (int p = 0; p < n; p++) if (p != pk)
                 c = a[p][pk];
                 a[p][pk] = 0;
                 for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;</pre>
                 for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;</pre>
            }
    }
    for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p])
            for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);</pre>
    return det;
int main()
    const int n = 4;
    const int m = 2;
    double A[n][n] = \{ \{1,2,3,4\}, \{1,0,1,0\}, \{5,3,2,4\}, \{6,1,4,6\} \};
    double B[n][m] = \{ \{1,2\}, \{4,3\}, \{5,6\}, \{8,7\} \};
    VVT a(n), b(n);
    for (int i = 0; i < n; i++)</pre>
        a[i] = VT(A[i], A[i] + n);
        b[i] = VT(B[i], B[i] + m);
    double det = GaussJordan(a, b);
    // expected: 60
    cout << "Determinant: " << det << endl;</pre>
```

```
cout << "Inverse: " << endl;</pre>
    for (int i = 0; i < n; i++)</pre>
        for (int j = 0; j < n; j++)
          cout << a[i][j] << ' ';
        cout << endl;</pre>
    }
    // expected: 1.63333 1.3
    cout << "Solution: " << endl;</pre>
    for (int i = 0; i < n; i++)</pre>
        for (int j = 0; j < m; j++)
          cout << b[i][j] << ' ';
       cout << endl;</pre>
    }
}
ReducedRowEchelonForm.cc 14/27
// INPUT: a[][] = an nxn matrix
// OUTPUT: rref[][] = an nxm matrix (stored in a[][])
            returns rank of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPSILON = 1e-10;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
int rref(VVT &a)
    int n = a.size();
    int m = a[0].size();
    int r = 0;
    for (int c = 0; c < m; c++)</pre>
        int j = r;
        for (int i = r+1; i < n; i++)</pre>
            if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
        if (fabs(a[j][c]) < EPSILON) continue;</pre>
        swap(a[j], a[r]);
        T s = 1.0 / a[r][c];
        for (int j = 0; j < m; j++) a[r][j] *= s;</pre>
        for (int i = 0; i < n; i++) if (i != r)</pre>
            {
```

```
T t = a[i][c];
                 for (int j = 0; j < m; j++) a[i][j] -= t * a[r][j];</pre>
        r++;
    }
    return r;
int main()
    const int n = 5;
    const int m = 4;
     \textbf{double} \ A[n][m] \ = \ \left\{ \ \{16,2,3,13\}, \{5,11,10,8\}, \{9,7,6,12\}, \{4,14,15,1\}, \{13,21,21,13\} \ \right\}; 
    VVT a(n);
    for (int i = 0; i < n; i++)</pre>
        a[i] = VT(A[i], A[i] + n);
    int rank = rref (a);
    // expected: 4
    cout << "Rank: " << rank << endl;</pre>
    cout << "rref: " << endl;</pre>
    for (int i = 0; i < 5; i++)</pre>
        for (int j = 0; j < 4; j++)
           cout << a[i][j] << ' ';</pre>
        cout << endl;</pre>
    }
}
FFT_new.cpp 15/27
==========
#include <cassert>
#include <cstdio>
#include <cmath>
struct cpx
    cpx() {}
    cpx(double aa):a(aa) {}
    cpx(double aa, double bb):a(aa),b(bb) {}
    double a;
    double b;
    double modsq(void) const
        return a * a + b * b;
    }
    cpx bar(void) const
        return cpx(a, -b);
    }
};
cpx operator +(cpx a, cpx b)
    return cpx(a.a + b.a, a.b + b.b);
```

```
cpx operator *(cpx a, cpx b)
   return cpx(a.a * b.a - a.b * b.b, a.a * b.b + a.b * b.a);
cpx operator /(cpx a, cpx b)
   cpx r = a * b.bar();
   return cpx(r.a / b.modsq(), r.b / b.modsq());
cpx EXP(double theta)
   return cpx(cos(theta),sin(theta));
const double two_pi = 4 * acos(0);
// size: length of the input/output {MUST BE A POWER OF 2}
// RESULT: out[k] = \sum_{j=0}^{s-1} in[j] * exp(dir * 2pi * i * j * k / size)
void FFT(cpx *in, cpx *out, int step, int size, int dir)
   if(size < 1) return;</pre>
   if(size == 1)
       out[0] = in[0];
       return;
   FFT(in, out, step * 2, size / 2, dir);
   FFT(in + step, out + size / 2, step * 2, size / 2, dir);
   for(int i = 0; i < size / 2; i++)
        cpx even = out[i];
        cpx odd = out[i + size / 2];
        out[i] = even + EXP(dir * two_pi * i / size) * odd;
        out[i + size / 2] = even + EXP(dir * two_pi * (i + size / 2) / size) * odd;
}
// f[0...N-1] and g[0..N-1] are numbers
  Want to compute the convolution h, defined by
// h[n] = sum of f[k]g[n-k] (k = 0, ..., N-1).
// Let F[0...N-1] be FFT(f), and similarly, define G and H.
// The convolution theorem says H[n] = F[n]G[n] (element-wise product).
// To compute h[\ ] in O(N\ log\ N) time, do the following:
// 1. Compute F and G (pass dir = 1 as the argument).
    2. Get H by element-wise multiplying F and G.
       and *dividing by N*. DO NOT FORGET THIS SCALING FACTOR.
int main(void)
   printf("If rows come in identical pairs, then everything works.\n");
   cpx a[8] = \{0, 1, cpx(1,3), cpx(0,5), 1, 0, 2, 0\};
   cpx b[8] = \{1, cpx(0,-2), cpx(0,1), 3, -1, -3, 1, -2\};
   cpx A[8];
   cpx B[8];
   FFT(a, A, 1, 8, 1);
   FFT(b, B, 1, 8, 1);
```

```
for(int i = 0 ; i < 8 ; i++)</pre>
       printf("%7.21f%7.21f", A[i].a, A[i].b);
    }
   printf("\n");
   for(int i = 0 ; i < 8 ; i++)</pre>
       cpx Ai(0,0);
       for(int j = 0 ; j < 8 ; j++)
           Ai = Ai + a[j] * EXP(j * i * two_pi / 8);
       printf("%7.21f%7.21f", Ai.a, Ai.b);
    }
   printf("\n");
   cpx AB[8];
   for(int i = 0 ; i < 8 ; i++)
       AB[i] = A[i] * B[i];
   cpx aconvb[8];
   FFT(AB, aconvb, 1, 8, -1);
   for(int i = 0 ; i < 8 ; i++)</pre>
       aconvb[i] = aconvb[i] / 8;
   for(int i = 0 ; i < 8 ; i++)</pre>
       printf("%7.21f%7.21f", aconvb[i].a, aconvb[i].b);
   }
   printf("\n");
   for(int i = 0 ; i < 8 ; i++)
       cpx aconvbi(0,0);
       for(int j = 0 ; j < 8 ; j++)
           aconvbi = aconvbi + a[j] * b[(8 + i - j) % 8];
       printf("%7.21f%7.21f", aconvbi.a, aconvbi.b);
   printf("\n");
   return 0;
Simplex.cc 16/27
==========
// INPUT: A -- an m x n matrix
// b -- an m-dimensional vector
        c -- an n-dimensional vector
        x -- a vector where the optimal solution will be stored
// above, nan if infeasible)
//\ \mbox{To} use this code, create an LPSolver object with A, b, and c as
// arguments. Then, call Solve(x).
#include <iostream>
#include <iomanip>
#include <vector>
#include <cmath>
```

```
using namespace std;
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver
   int m, n;
   VI B, N;
   VVD D;
   LPSolver(const VVD &A, const VD &b, const VD &c):
       m(b.size()), n(c.size()), N(n+1), B(m), D(m+2, VD(n+2))
       for (int i = 0; i < m; i++) for (int j = 0; j < n; j++) D[i][j] = A[i][j];</pre>
       for (int i = 0; i < m; i++)</pre>
           B[i] = n+i;
           D[i][n] = -1;
           D[i][n+1] = b[i];
        for (int j = 0; j < n; j++)
           N[j] = j;
           D[m][j] = -c[j];
       N[n] = -1;
       D[m+1][n] = 1;
   void Pivot(int r, int s)
        for (int i = 0; i < m+2; i++) if (i != r)</pre>
                for (int j = 0; j < n+2; j++) if (j != s)
                       D[i][j] = D[r][j] * D[i][s] / D[r][s];
        for (int j = 0; j < n+2; j++) if (j != s) D[r][j] /= D[r][s];</pre>
        for (int i = 0; i < m+2; i++) if (i != r) D[i][s] /= -D[r][s];</pre>
        D[r][s] = 1.0 / D[r][s];
        swap(B[r], N[s]);
    }
   bool Simplex(int phase)
        int x = phase == 1 ? m+1 : m;
        while (true)
            int s = -1;
            for (int j = 0; j <= n; j++)</pre>
                if (phase == 2 && N[j] == -1) continue;
                if (D[x][s] >= -EPS) return true;
            int r = -1;
            for (int i = 0; i < m; i++)</pre>
                if (D[i][s] <= 0) continue;</pre>
                if (r == -1 | D[i][n+1] / D[i][s] < D[r][n+1] / D[r][s] | |</pre>
                        D[i][n+1] \ / \ D[i][s] \ == \ D[r][n+1] \ / \ D[r][s] \ \&\& \ B[i] \ < \ B[r]) \ r \ = \ i;
            }
```

```
if (r == -1) return false;
            Pivot(r, s);
    }
    DOUBLE Solve(VD &x)
        int r = 0;
        for (int i = 1; i < m; i++) if (D[i][n+1] < D[r][n+1]) r = i;
        if (D[r][n+1] <= -EPS)</pre>
            Pivot(r, n);
            if (!Simplex(1) | D[m+1][n+1] < -EPS) return -numeric_limits<DOUBLE>::infinity();
            for (int i = 0; i < m; i++) if (B[i] == -1)</pre>
                     int s = -1;
                     for (int j = 0; j <= n; j++)</pre>
                        if (s == -1 | | D[i][j] < D[i][s] | | D[i][j] == D[i][s] && N[j] < N[s]) s = j;
        if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity();
        for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][n+1];</pre>
        return D[m][n+1];
};
int main()
{
    const int m = 4;
    const int n = 3i
    DOUBLE _A[m][n] =
        { 6, -1, 0 },
        { -1, -5, 0 },
        { 1, 5, 1 },
        \{-1, -5, -1\}
    };
    DOUBLE _b[m] = \{ 10, -4, 5, -5 \};
    DOUBLE _{c[n]} = \{ 1, -1, 0 \};
    VVD A(m);
    VD b(_b, _b + m);
    VD c(_c, _c + n);
    for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n);</pre>
    LPSolver solver(A, b, c);
    VD x;
    DOUBLE value = solver.Solve(x);
    cerr << "VALUE: "<< value << endl;</pre>
    cerr << "SOLUTION:";</pre>
    for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];</pre>
    cerr << endl;</pre>
    return 0;
FastDijkstra.cc 17/27
-----
// Implementation of Dijkstra's algorithm using adjacency lists
// and priority queue for efficiency.
```

```
#include <queue>
#include <stdio.h>
using namespace std;
const int INF = 2000000000;
typedef pair<int,int> PII;
int main()
   int N, s, t;
   scanf ("%d%d%d", &N, &s, &t);
   vector<vector<PII> > edges(N);
   for (int i = 0; i < N; i++)</pre>
       int M;
       scanf ("%d", &M);
       for (int j = 0; j < M; j++)
           int vertex, dist;
           scanf ("%d%d", &vertex, &dist);
           edges[i].push_back (make_pair (dist, vertex)); // note order of arguments here
    }
    // use priority queue in which top element has the "smallest" priority
   priority_queue<PII, vector<PII>, greater<PII> > Q;
   vector<int> dist(N, INF), dad(N, -1);
   Q.push (make_pair (0, s));
   dist[s] = 0;
   while (!Q.empty())
       PII p = Q.top();
       if (p.second == t) break;
        Q.pop();
        int here = p.second;
        for (vector<PII>::iterator it=edges[here].begin(); it!=edges[here].end(); it++)
            if (dist[here] + it->first < dist[it->second])
                dist[it->second] = dist[here] + it->first;
                dad[it->second] = here;
                Q.push (make_pair (dist[it->second], it->second));
       }
   printf ("%d\n", dist[t]);
    if (dist[t] < INF)</pre>
        for(int i=t; i!=-1; i=dad[i])
           printf ("%d%c", i, (i==s?'\n':' '));
   return 0;
}
SCC.cc 18/27
-----
#include<memory.h>
struct edge
   int e, nxt;
};
```

```
int V, E;
edge e[MAXE], er[MAXE];
int sp[MAXV], spr[MAXV];
int group_cnt, group_num[MAXV];
bool v[MAXV];
int stk[MAXV];
void fill_forward(int x)
    int i;
   v[x]=true;
    for(i=sp[x]; i; i=e[i].nxt) if(!v[e[i].e]) fill_forward(e[i].e);
   stk[++stk[0]]=x;
void fill_backward(int x)
    int i;
    v[x]=false;
    group_num[x]=group_cnt;
    for(i=spr[x]; i; i=er[i].nxt) if(v[er[i].e]) fill_backward(er[i].e);
void add_edge(int v1, int v2) //add edge v1->v2
    e [++E].e=v2;
    e [E].nxt=sp [v1];
    sp[v1]=E;
    er[ E].e=v1;
    er[E].nxt=spr[v2];
    spr[v2]=E;
}
void SCC()
    int i;
    stk[0]=0;
    memset(v, false, sizeof(v));
    for(i=1; i<=V; i++) if(!v[i]) fill_forward(i);</pre>
    group_cnt=0;
   for(i=stk[0]; i>=1; i--) if(v[stk[i]])
        group_cnt++;
       fill_backward(stk[i]);
SuffixArray.cc 19/27
// INPUT: string s
           of substring s[i...L-1] in the list of sorted suffixes.
           That is, if we take the inverse of the permutation suffix[],
           we get the actual suffix array.
#include <vector>
#include <iostream>
#include <string>
using namespace std;
struct SuffixArray
    const int L;
```

```
string s;
    vector<vector<int> > P;
    vector<pair<int,int>,int> > M;
    SuffixArray(const string &s) : L(s.length()), s(s), P(1, vector<int>(L, 0)), M(L)
        for (int i = 0; i < L; i++) P[0][i] = int(s[i]);</pre>
        for (int skip = 1, level = 1; skip < L; skip *= 2, level++)</pre>
            P.push_back(vector<int>(L, 0));
            for (int i = 0; i < L; i++)</pre>
                M[i] = make_pair(make_pair(P[level-1][i], i + skip < L ? P[level-1][i + skip] : -1000), i);</pre>
            sort(M.begin(), M.end());
            for (int i = 0; i < L; i++)</pre>
                P[level][M[i].second] = (i > 0 && M[i].first == M[i-1].first) ? P[level][M[i-1].second] : i;
    }
    vector<int> GetSuffixArray()
        return P.back();
    int LongestCommonPrefix(int i, int j)
        int len = 0;
        if (i == j) return L - i;
        for (int k = P.size() - 1; k >= 0 && i < L && j < L; k--)
            if (P[k][i] == P[k][j])
                i += 1 << k;
                j += 1 << k;
                len += 1 << k;
        return len;
int main()
         el is the 3'rd suffix
           l is the 4'th suffix
    SuffixArray suffix("bobocel");
    vector<int> v = suffix.GetSuffixArray();
    // Expected output: 0 5 1 6 2 3 4
    for (int i = 0; i < v.size(); i++) cout << v[i] << " ";</pre>
    cout << endl;</pre>
    cout << suffix.LongestCommonPrefix(0, 2) << endl;</pre>
BIT.cc 20/27
=========
```

};

}

```
using namespace std;
#define LOGSZ 17
int tree[(1<<LOGSZ)+1];</pre>
int N = (1 << LOGSZ);
// add v to value at x
void set(int x, int v)
   while(x <= N)</pre>
    {
       tree[x] += v;
       x += (x & -x);
}
// get cumulative sum up to and including \boldsymbol{x}
int get(int x)
   int res = 0;
   while(x)
       res += tree[x];
       x = (x \& -x);
   return res;
}
// get largest value with cumulative sum less than or equal to x;
// for smallest, pass x-1 and add 1 to result
int getind(int x)
   int idx = 0, mask = N;
   while(mask && idx < N)</pre>
       int t = idx + mask;
       if(x >= tree[t])
           idx = t;
           x -= tree[t];
       mask >>= 1;
   return idx;
UnionFind.cc 21/27
===========
int find(vector <int>& C, int x)
   return (C[x]==x) ? x : C[x]=find(C, C[x]); //C++
int find(int x)
   return (C[x]==x)?x:C[x]=find(C[x]); //C
}
KDTree.cc 22/27
==========
// probably good enough for most things (current it's a 2D-tree)
```

```
// - constructs from n points in O(n lg^2 n) time
// - handles nearest-neighbor query in O(\lg\,n) if points are well distributed
// - worst case for nearest-neighbor may be linear in pathological case
// Sonny Chan, Stanford University, April 2009
#include <iostream>
#include <vector>
#include <limits>
#include <cstdlib>
using namespace std;
typedef long long ntype;
const ntype sentry = numeric_limits<ntype>::max();
// point structure for 2D-tree, can be extended to 3D
struct point
{
   ntype x, y;
    point(ntype xx = 0, ntype yy = 0) : x(xx), y(yy) {}
};
bool operator==(const point &a, const point &b)
    return a.x == b.x && a.y == b.y;
}
// sorts points on x-coordinate
bool on_x(const point &a, const point &b)
    return a.x < b.x;</pre>
// sorts points on y-coordinate
bool on_y(const point &a, const point &b)
    return a.y < b.y;</pre>
ntype pdist2(const point &a, const point &b)
    ntype dx = a.x-b.x, dy = a.y-b.y;
    return dx*dx + dy*dy;
struct bbox
    ntype x0, x1, y0, y1;
    bbox() : x0(sentry), x1(-sentry), y0(sentry), y1(-sentry) {}
    // computes bounding box from a bunch of points
    void compute(const vector<point> &v)
    {
        for (int i = 0; i < v.size(); ++i)</pre>
            x0 = min(x0, v[i].x);
            x1 = max(x1, v[i].x);
            y0 = min(y0, v[i].y);
            y1 = max(y1, v[i].y);
```

```
// squared distance between a point and this bbox, 0 if inside
   ntype distance(const point &p)
       if (p.x < x0)
                               return pdist2(point(x0, y0), p);
           if (p.y < y0)
            else if (p.y > y1) return pdist2(point(x0, y1), p);
                               return pdist2(point(x0, p.y), p);
        else if (p.x > x1)
            if (p.y < y0)
                               return pdist2(point(x1, y0), p);
            else if (p.y > y1) return pdist2(point(x1, y1), p);
                               return pdist2(point(x1, p.y), p);
       else
            if (p.y < y0)
                               return pdist2(point(p.x, y0), p);
           else if (p.y > y1) return pdist2(point(p.x, y1), p);
                               return 0;
    }
};
struct kdnode
   bool leaf;
                   // true if this is a leaf node (has one point)
                   // the single point of this is a leaf
   point pt;
   bbox bound;
   kdnode *first, *second; // two children of this kd-node
   kdnode() : leaf(false), first(0), second(0) {}
   ~kdnode()
       if (first) delete first;
       if (second) delete second;
    // intersect a point with this node (returns squared distance)
   ntype intersect(const point &p)
    {
       return bound.distance(p);
    // recursively builds a kd-tree from a given cloud of points
   void construct(vector<point> &vp)
       bound.compute(vp);
        // if we're down to one point, then we're a leaf node
       if (vp.size() == 1)
           leaf = true;
           pt = vp[0];
        }
        else
            // split on x if the bbox is wider than high (not best heuristic...)
            if (bound.x1-bound.x0 >= bound.y1-bound.y0)
                sort(vp.begin(), vp.end(), on_x);
```

```
else
                sort(vp.begin(), vp.end(), on_y);
            // divide by taking half the array for each child
            // (not best performance if many duplicates in the middle)
            int half = vp.size()/2;
            vector<point> vl(vp.begin(), vp.begin()+half);
            vector<point> vr(vp.begin()+half, vp.end());
            first = new kdnode();
            first->construct(vl);
            second = new kdnode();
           second->construct(vr);
    }
};
struct kdtree
   kdnode *root;
    // constructs a kd-tree from a points (copied here, as it sorts them)
   kdtree(const vector<point> &vp)
       vector<point> v(vp.begin(), vp.end());
       root = new kdnode();
       root->construct(v);
    }
    ~kdtree()
    {
        delete root;
    // recursive search method returns squared distance to nearest point
   ntype search(kdnode *node, const point &p)
        if (node->leaf)
            // commented special case tells a point not to find itself
            return pdist2(p, node->pt);
        ntype bfirst = node->first->intersect(p);
        ntype bsecond = node->second->intersect(p);
        // choose the side with the closest bounding box to search first
        if (bfirst < bsecond)</pre>
           ntype best = search(node->first, p);
            if (bsecond < best)</pre>
                best = min(best, search(node->second, p));
            return best;
        }
        else
            ntype best = search(node->second, p);
            if (bfirst < best)</pre>
                best = min(best, search(node->first, p));
            return best;
    }
```

```
ntype nearest(const point &p)
       return search(root, p);
   }
};
// some basic test code here
int main()
   // generate some random points for a kd-tree
   vector<point> vp;
   for (int i = 0; i < 100000; ++i)</pre>
       vp.push_back(point(rand()%100000, rand()%100000));
   kdtree tree(vp);
   for (int i = 0; i < 10; ++i)</pre>
       point q(rand()%100000, rand()%100000);
        \texttt{cout} << "Closest squared distance to (" << q.x << ", " << q.y << ")"
            << " is " << tree.nearest(q) << endl;</pre>
    }
   return 0;
}
LongestIncreasingSubsequence.cc 23/27
_____
// Given a list of numbers of length n, this routine extracts a
   INPUT: a vector of integers
    OUTPUT: a vector containing the longest increasing subsequence
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int,int> PII;
typedef vector<PII> VPII;
#define STRICTLY_INCREASNG
VI LongestIncreasingSubsequence(VI v)
   VPII best;
   VI dad(v.size(), -1);
   for (int i = 0; i < v.size(); i++)</pre>
#ifdef STRICTLY_INCREASNG
       PII item = make_pair(v[i], 0);
       VPII::iterator it = lower_bound(best.begin(), best.end(), item);
        item.second = i;
```

```
#else
       VPII::iterator it = upper_bound(best.begin(), best.end(), item);
#endif
        if (it == best.end())
           dad[i] = (best.size() == 0 ? -1 : best.back().second);
           best.push_back(item);
        else
           dad[i] = dad[it->second];
           *it = item;
    }
   VI ret;
   for (int i = best.back().second; i >= 0; i = dad[i])
       ret.push_back(v[i]);
   reverse(ret.begin(), ret.end());
   return ret;
}
Dates.cc 24/27
==========
// as integers from 1 to 31, and years are expressed as 4-digit
#include <iostream>
#include <string>
using namespace std;
string dayOfWeek[] = {"Mon", "Tue", "Wed", "Thu", "Fri", "Sat", "Sun"};
// converts Gregorian date to integer (Julian day number)
int dateToInt (int m, int d, int y)
   return
       1461 * (y + 4800 + (m - 14) / 12) / 4 +
        367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
        3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
       d - 32075;
void intToDate (int jd, int &m, int &d, int &y)
{
   int x, n, i, j;
   x = jd + 68569;
   n = 4 * x / 146097;
   x = (146097 * n + 3) / 4;
   i = (4000 * (x + 1)) / 1461001;
   x = 1461 * i / 4 - 31;
   j = 80 * x / 2447;
   d = x - 2447 * j / 80;
   x = j / 11;
   m = j + 2 - 12 * x;
   y = 100 * (n - 49) + i + x;
```

```
// converts integer (Julian day number) to day of week
    string intToDay (int jd)
       return dayOfWeek[jd % 7];
    int main (int argc, char **argv)
       int jd = dateToInt (3, 24, 2004);
       int m, d, y;
       intToDate (jd, m, d, y);
       string day = intToDay (jd);
            Wed
       cout << jd << endl << m << "/" << d << "/" << y << endl << day << endl;</pre>
   LogLan.java 25/27
    ==========
    // This is a solution for
    // Loglan: a logical language
    // http://acm.uva.es/p/v1/134.html
    \ensuremath{//} In this problem, we are given a regular language, whose rules can be
    // inferred directly from the code. For each sentence in the input, we must
    // determine whether the sentence matches the regular expression or not. The
    // code consists of (1) building the regular expression (which is fairly
    // complex) and (2) using the regex to match sentences.
    import java.util.*;
    import java.util.regex.*;
   public class LogLan
       public static String BuildRegex ()
            String space = " +";
           String A = "([aeiou])";
            String C = "([a-z&&[^aeiou]])";
           String MOD = "(g" + A + ")";
           String BA = "(b" + A + ")";
            String DA = (d' + A + )';
            String LA = "(1" + A + ")";
            String NAM = "([a-z]*" + C + ")";
            String PREDA = "(" + C + C + A + C + A + " | " + C + A + C + C + A + ")";
            String predstring = "(" + PREDA + "(" + space + PREDA + ")*)";
            String predname = "(" + LA + space + predstring + " | " + NAM + ")";
            String preds = "(" + predstring + "(" + space + A + space + predstring + ")*)";
            String predclaim = "(" + predname + space + BA + space + preds + "|" + DA + space + preds + ")";
            String verbpred = "(" + MOD + space + predstring + ")";
            String statement = "(" + predname + space + verbpred + space + predname + " | " +predname + space +
verbpred + ")";
            String sentence = "(" + statement + " | " + predclaim + ")";
            return "^" + sentence + "$";
        }
```

```
String regex = BuildRegex();
        Pattern pattern = Pattern.compile (regex);
        Scanner s = new Scanner(System.in);
        while (true)
            // In this problem, each sentence consists of multiple lines, where the last
            \ensuremath{//} line is terminated by a period. The code below reads lines until
            // encountering a line whose final character is a '.'. Note the use of
                 s.charAt() to extract characters from a Java string
                 s.compareTo(t) < 0 if s < t, lexicographically</pre>
                 s.indexOf("apple") returns index of first occurrence of "apple" in s
                 s.replace(c,d) replaces occurrences of character c with d
                  s.toLowerCase() / s.toUpperCase() returns a new lower/uppercased string
                  Double.parseDouble(s) converts s to a double
            String sentence = "";
            while (true)
                sentence = (sentence + " " + s.nextLine()).trim();
                if (sentence.equals("#")) return;
                if (sentence.charAt(sentence.length()-1) == '.') break;
            // now, we remove the period, and match the regular expression
            String removed_period = sentence.substring(0, sentence.length()-1).trim();
            if (pattern.matcher (removed_period).find())
                System.out.println ("Good");
            else
                System.out.println ("Bad!");
    }
Primes.cc 26/27
==========
// O(sqrt(x)) Exhaustive Primality Test
#include <cmath>
#define EPS 1e-7
typedef long long LL;
bool IsPrimeSlow (LL x)
    if(x<=1) return false;</pre>
    if(x<=3) return true;</pre>
```

public static void main (String args[])

```
if (!(x%2) | !(x%3)) return false;
   LL s=(LL)(sqrt((double)(x))+EPS);
   for(LL i=5; i<=s; i+=6)</pre>
      if (!(x%i) || !(x%(i+2))) return false;
   }
   return true;
}
                                  17 19
               47
     41
          43
     97 101 103
    157 163 167
                         389 397 401 409
                                            419
                                                  421
                                                        431
                                                             433
    439 443 449
                   457
                         461 463 467 479
                                                  491
    The largest prime smaller than 1000000000 is 999999937.
    The largest prime smaller than 10000000000 is 9999999967.
    The largest prime smaller than 10000000000 is 99999999977.
    The largest prime smaller than 1000000000000 is 9999999999973
    The largest prime smaller than 1000000000000000 is 99999999999999937
KMP.cpp 27/27
_____
/*
Searches for the string w in the string s (of length k). Returns the
0-based index of the first match (k if no match is found). Algorithm
runs in O(k) time.
* /
#include <iostream>
#include <string>
#include <vector>
using namespace std;
typedef vector<int> VI;
void buildTable(string& w, VI& t)
   t = VI(w.length());
```

```
int i = 2, j = 0;
        t[0] = -1;
        t[1] = 0;
        while(i < w.length())</pre>
            if(w[i-1] == w[j])
                t[i] = j+1;
                i++;
                j++;
            else if(j > 0) j = t[j];
            else
               t[i] = 0;
                i++;
        }
    int KMP(string& s, string& w)
        int m = 0, i = 0;
        VI t;
        buildTable(w, t);
        while(m+i < s.length())</pre>
            if(w[i] == s[m+i])
                i++;
                if(i == w.length()) return m;
            else
               m += i-t[i];
                if(i > 0) i = t[i];
        return s.length();
    int main()
        string a = (string) "The example above illustrates the general technique";
        string b = "table";
        int p = KMP(a, b);
        cout << p << ": " << a.substr(p, b.length()) << " " << b << endl;</pre>
    LCA_by array:
    -----
        You are given a rooted tree, where each node contains an integer value.
        And the value of a node is strictly greater than the value of its parent.
       Now you are given a node and an integer query. You have to find the greatest possible parent of this
node
       (may include the node itself), whose value if greater than or equal to the given query integer.
    * /
    #include <cstdio>
    #include <cstring>
    const int MAX = 100000;
```

```
const int LOG = 20;
int T[MAX], V[MAX];
int P[MAX][LOG];
int main()
    int test, cs, n, u, m, i, j, q, lg;
    V[0] = 1; T[0] = -1;
    scanf("%d", &test);
    for(cs = 1; cs <= test; cs++)</pre>
        scanf("%d %d", &n, &q);
        for(i = 1; i < n; i++)</pre>
            scanf("%d %d", &u, &m); // ith node's parent and ith node's value.
            T[i] = u, V[i] = m;
        for(i = 0; i < n; i++) P[i][0] = T[i];</pre>
        for(j = 1; (1 << j) < n; j++)
            for(i = 0; i < n; i++)</pre>
                if(P[i][j-1] != -1) P[i][j] = P[P[i][j-1]][j-1];
                else P[i][j] = -1;
        for(lg = 1; (1 << lg) < n; lg++);</pre>
        lg--;
        printf("Case %d:\n", cs);
        while(q--)
            scanf("%d %d", &u, &m);
            for(j = lg; j >= 0; j--)
                if(P[u][j] != -1 \&\& V[P[u][j]] >= m) u = P[u][j];
            printf("%d\n", u);
    return 0;
LCA_By -dfs :
_____
    You are given a rooted tree, where each node contains an integer value.
    And the value of a node is strictly greater than the value of its parent.
    Now you are given a node and an integer query. You have to find the greatest possible parent of this
    (may include the node itself), whose value if greater than or equal to the given query integer.
#include <cstdio>
#include <cstring>
#define MAXN 100100
#define LOGMAXN 18
int color[MAXN],parent[MAXN][LOGMAXN],maxValue[MAXN],store[MAXN];
int start[MAXN],finish[MAXN],T,step,nodes;
vector<int> adj[MAXN];
void dfs(int u,int par)
```

```
int i,v;color[u]=1;start[u]=T++;
    parent[u][0]=par;
    for(i=1;i<=step;i++) parent[u][i]=parent[parent[u][i-1]][i-1];</pre>
    maxValue[u]=store[u];
    maxValue[u]=max(maxValue[u],maxValue[par]);
    REP(i,SZ(adj[u]))
        v=adj[u][i];if(color[v]) continue;
        dfs(v,u);
   finish[u]=T++;
bool IsAnchestor(int u,int v)
    if(start[u]<=start[v] && finish[u]>=finish[v]) return true;
    return false;
int lca_query(int v,int val)
    int i,j,p;
    for(i=step;i>=0;i--)
        for(j=step;j>=0;j--)
           p=parent[v][j];
           if(maxValue[p]>=val){v=p;break;}
    return v;
int main()
    int test,Case=1,i,j,Q,u,v,val,ans;scanf("%d",&test);
    while(test--)
        scanf("%d %d",&nodes,&Q);
       for(int i=0;i<=nodes;i++) adj[i].clear();</pre>
        store[0]=1;
        for(int i=0;i<nodes-1;i++)</pre>
            u=i+1;scanf("%d %d",&v,&val);
            store[u]=val;
            adj[u].pb(v);
            adj[v].pb(u);
        for(step=0;two(step)<=nodes;step++);</pre>
        memset(color,0,sizeof(color));
        T=0;dfs(0,0);
        printf("Case %d:\n",Case++);
        while(Q--)
            scanf("%d %d",&v,&val);
            ans=lca_query(v,val);
            printf("%d\n",ans);
    return 0;
}
SegmentTree (Range Update):
-----
#include<stdio.h>
#include<string.h>
#define mem(a,b) memset(a,b,sizeof(a))
```

```
#define MAX 100002
using namespace std;
typedef long long 11;
#define type int
type N,Q;
type tree[4*MAX],col[4*MAX];
type arr[MAX];
type build(int node,int left,int right)
    if(left==right)
        tree[node]=arr[left];
        return tree[node];
    int mid=(right+left)/2;
    type p=build(node*2,left,mid);
    type q=build(node*2+1,mid+1,right);
    tree[node]=p+q;
    return tree[node];
void refresh(int node,int left,int right)
    /// some change here.
    if(left==right) tree[node]+=col[node];
    else
        col[node*2]+=col[node];
        col[node*2+1]+=col[node];
        tree[node] = tree[node] + (right - left + 1) * col[node];
    col[node]=0;
    return;
type quary(int node,int left,int right,int i,int j)
    if(col[node]) refresh(node,left,right);
    if(right<i or left>j) return 0;
    if(left>=i and right<=j) return tree[node];</pre>
    int mid=(left+right)/2;
    type p=quary(node*2,left,mid,i,j);
    type q=quary(node*2+1,mid+1,right,i,j);
    return p+q;
void update(int node,int left,int right,int i,int j,type val)
    if(right<i or left>j) return;
    if(left>=i and right<=j)</pre>
        col[node]+=val;
        refresh(node, left, right);
        return;
    int mid=(left+right)/2;
    update(node*2,left,mid,i,j,val);
    update(node*2+1,mid+1,right,i,j,val);
    if(col[node*2]) refresh(node*2,left,mid);
    if(col[node*2+1]) refresh(node*2+1,mid+1,right);
    tree[node]=tree[node*2]+tree[node*2+1];
    return;
```

```
int main()
    scanf("%lld %lld",&N,&Q);
    memset(tree, 0, sizeof(tree));
    memset(col,0,sizeof(col));
    /// scan element.
    build(1,0,N-1);
    int x,y,z;type val;
    11 ans=quary(1,0,N-1,y,z);
    update(1,0,N-1,y,z,val);
    return 0;
SegmentTree(2D):
==========
A rectangular matrix a[N][M], and enter the amount of search queries (or minimum / maximum)
at some small matrix [x_1, x_2, y_1, y_2] and requests modification of individual elements
of the matrix (ie, queries of the form a [x][y] = p).
So, we will build a two-dimensional tree segments:
the first tree segments in the first coordinate ( x), then - on the second ( y).
#include<stdio.h>
#include<algorithm>
using namespace std;
typedef long long 11;
#define MAX 505
#define type int
type tree[4*MAX][4*MAX],
int N,grid[MAX][MAX];
void build_y(int vx,int lx,int rx,int vy,int ly,int ry)
    if(ly==ry)
        if(lx==rx) tree[vx][vy]=grid[lx][ly];
        else tree[vx][vy]=max(tree[vx*2][vy],tree[vx*2+1][vy]);
        return;
    int my=(ly+ry)/2;
    build_y(vx,lx,rx,vy*2,ly,my);
    build_y(vx,lx,rx,vy*2+1,my+1,ry);
    tree[vx][vy]=max(tree[vx][vy*2],tree[vx][vy*2+1]);
void build_x(int vx,int lx,int rx)
    if(lx!=rx)
        int mx=(lx+rx)/2;
        build_x(vx*2,lx,mx);
       build_x(vx*2+1,mx+1,rx);
    build_y(vx,lx,rx,1,1,N);
type max_y(int vx,int vy,int tly,int _try,int ly,int ry)
```

```
if(ry<tly or ly>_try) return 0;
    if(tly>=ly and _try<=ry)</pre>
        int t=tree[vx][vy];
        return t;
    int tmy=(tly+_try)/2;
    int ml=max_y(vx,vy*2,tly,tmy,ly,ry);
    int m2=max_y(vx,vy*2+1,tmy+1,_try,ly,ry);
    return max(m1,m2);
type max_x(int vx,int tlx,int trx,int lx,int rx,int ly,int ry)
    if(rx<tlx or lx>trx) return 0;
    if(tlx>=lx and trx<=rx);</pre>
       int t=max_y(vx,1,1,N,ly,ry);
       return t;
    int tmx=(tlx+trx)/2;
    int m1=max_x(vx*2,tlx,tmx,lx,rx,ly,ry);
    int m2=max_x(vx*2+1,tmx+1,trx,lx,rx,ly,ry);
    return max(m1,m2);
}
int main()
    int N,s1,s2,s,cnt;
   cin>>N>>s>>s1>>s2;
   build_x(1,1,N);
   cnt=max_x(1,1,N,s1,s1+s-1,s2,s2+s-1);
    return 0;
}
Trie-2D(array):
==========
    Trie implementation using array, faster and takes less memory.
    Each node can contain arbitrary data as needed for solving the problem.
    The ALPHABET, MAX and scale() may need tuning as necessary.
const int ALPHABET = 26;
const int MAX = 100000;
                           // Number of string * number of Alphabet
#define scale(x) (x-'a')
int next[MAX][ALPHABET];
int data[MAX];
struct Trie
    int n, root;
    void init()
       root = 0, n = 1;
        data[root] = 0;
        memset(next[root], -1, sizeof(next[root]));
    void insert(char *s)
        int curr = root, i, k;
        for(i = 0; s[i]; i++)
```

```
k = scale(s[i]);
           if(next[curr][k] == -1)
               next[curr][k] = n;
               memset(next[n], -1, sizeof(next[n]));
            }
            curr = next[curr][k];
           data[curr] = s[i]; // optional
        data[curr] = 0; // sentinel, optional
    bool find(char *s)
       int curr = root, i, k;
        for(i = 0; s[i]; i++)
           k = scale(s[i]);
           if(next[curr][k] == -1) return false;
           curr = next[curr][k];
       return (data[curr] == 0);
};
int main()
   Trie T;T.init();
   return 0;
Heavy-Light-Decomposition:
-----
/*
You are given a tree (an acyclic undirected connected graph) with {\tt N} nodes,
and edges numbered 1, 2, 3...N-1.
We will ask you to perfrom some instructions of the following form:
   CHANGE i ti : change the cost of the i-th edge to ti
    QUERY a b: ask for the maximum edge cost on the path from node a to node b.
#include <iostream>
#include <cstring>
#include <cstdio>
#include <algorithm>
#include <cmath>
#include <vector>
#define MAX 10010
using namespace std;
struct Edge
    int u, v, w;
    Edge() {};
    Edge(int a, int b, int c) \{u = a, v = b, w = c;\}
};
vector<Edge> edge[MAX];
Edge alledge[MAX];
int t, n, idx;
int dep[MAX], size[MAX];
int son[MAX], tid[MAX];
int pre[MAX], top[MAX];
```

```
bool vis[MAX];
/// heavy light Decomposition....
void find_heavy_edge(int x, int father, int depth)
    vis[x] = true;pre[x] = father;dep[x] = depth;
    size[x] = 1; son[x] = -1; int maxsize = 0;
    int sz = edge[x].size();
    for (int i = 0; i < sz; ++i)</pre>
        int child = edge[x][i].v;
        if (!vis[child])
            find_heavy_edge(child, x, depth + 1);
            size[x] += size[child];
            if (size[child] > maxsize)
                maxsize = size[child];
                son[x] = child;
    return;
}
void connect_heavy_edge(int x, int ance)
    vis[x] = true; tid[x] = ++idx; top[x] = ance;
    if (son[x] != -1) {connect_heavy_edge(son[x], ance);}
    int sz = edge[x].size();
    for (int i = 0; i < sz; ++i)</pre>
        int child = edge[x][i].v;
        if (!vis[child]){connect_heavy_edge(child, child);}
    return;
/// segment tree
int Max[MAX<<2], val[MAX];</pre>
void build(int rt,int l, int r)
    if (1 == r){Max[rt] = val[1];return ;}
   int m = (1 + r) >> 1;
    build(rt << 1,1,m);</pre>
    build((rt << 1) | 1,m + 1, r);
    Max[rt] = max(Max[rt<<1], Max[rt<<1|1]);
void update(int rt,int 1, int r, int p, int c)
    if (1 == r){Max[rt] = c;return ;}
    int m = (1 + r) >> 1;
    if (p <= m) {update( rt << 1,1, m, p, c); }</pre>
    else {update((rt << 1) | 1,m + 1, r, p, c);}
    Max[rt] = max(Max[rt << 1], Max[rt << 1|1]);
}
int query(int rt,int 1, int r, int L, int R)
    if (L <= 1 && R >= r){return Max[rt];}
    int m = (1 + r) >> 1;
    int tmp = -(1 << 30);
    if (L <= m) {tmp = max(tmp, query(rt << 1,1, m, L, R));}</pre>
    if (R > m) \{ tmp = max(tmp, query((rt << 1) | 1, m + 1, r, L, R)); \}
    return tmp;
```

```
void CHANGE(int x, int val)
    if (dep[alledge[x].u] > dep[alledge[x].v]){
       update(1,2, n,tid[alledge[x].u], val);
    else
       update(1,2, n,tid[alledge[x].v], val);
    }
}
int QUERY(int a, int b)
    int ans = -(1 << 30);
    while (top[a] != top[b])
        if (dep[top[a]] < dep[top[b]]) swap(a, b);</pre>
        ans = \max(ans, query(1,2, n,tid[top[a]), tid[a]));
        a = pre[top[a]];
    if (dep[a] > dep[b]) swap(a, b);
    if (a != b) ans = max(ans, query(1,2, n,tid[a] + 1, tid[b]));
    return ans;
int main()
    scanf("%d", &t);
    while (t--)
       scanf("%d", &n);
        for (int i = 0; i <= n; ++i){edge[i].clear();}</pre>
        int a, b, c;
        for (int i = 1; i <= n - 1; ++i)</pre>
            scanf("%d%d%d", &a, &b, &c);
           alledge[i].u = a;
           alledge[i].v = b;
            alledge[i].w = c;
            edge[a].push_back(Edge(a, b, c));
            edge[b].push_back(Edge(b, a, c));
        memset(vis, false, sizeof(vis));
        find_heavy_edge(1, 1, 1);
        idx = 0;
        memset(vis, false, sizeof(vis));
        connect_heavy_edge(1, 1);
        for (int i = 1; i <= n - 1; ++i)</pre>
            if (dep[alledge[i].u] > dep[alledge[i].v])
                val[tid[alledge[i].u]] = alledge[i].w;
            else
                val[tid[alledge[i].v]] = alledge[i].w;
        build(1,2, n);char op[10];
        while (scanf("%s", op))
            if (op[0] == 'D') break;
            scanf("%d%d", &a, &b);
            if (op[0] == 'Q')
                printf("%d\n", QUERY(a, b));
            else
                CHANGE(a, b);
    }
```

```
return 0;
BCC():
======
#include <cstdio>
#include <iostream>
#include <vector>
#include <stack>
using namespace std;
#define mp make_pair
#define SZ(c) (int)c.size()
#define MAX 10000
const double EPS = 1e-9;
const int INF = 0x7f7f7f7f;
vector< int > G[MAX];
stack<pair<int,int> > S;
int dfstime;
int low[MAX], vis[MAX], used[MAX];
int N,M;
void dfs(int u, int par)
    int v, i, sz = G[u].size();pair<int,int> e, curr;
    used[u] = 1;vis[u] = low[u] = ++dfstime;
    for(i = 0; i < sz; i++)</pre>
       v = G[u][i];
       if(v == par) continue;
        if(!used[v])
            S.push(mp(u, v));
            dfs(v, u);
            if(low[v] >= vis[u])
                // new component
                curr = mp(u, v);int cnt=0;
                do
                    e = S.top();
                    S.pop();
                     // e is an edge in current bcc
                     //if(e==curr and cnt==1) break; /// single edge cheak.
                while(e != curr);
            low[u] = min(low[u], low[v]);
        else if(vis[v] < vis[u])</pre>
            S.push(mp(u, v));
            low[u] = min(low[u], vis[v]);
    }
    return;
void BCC()
    for(int i=0;i<=N;i++){low[i]=vis[i]=used[i]=0;}</pre>
    while(!S.empty()) S.pop();
    dfstime=0;
    for(int i=0;i<N;i++)</pre>
```

```
if(used[i]==0) dfs(i,-1);
    /// if low is differt between two nodes. Then we can say that...
    /// they are in two different bcc.
    return;
int main()
    while(cin>>N and N)
        cin>>M;int a,b;
        for(int i=0;i<=N;i++) G[i].clear();</pre>
        for(int i=0;i<M;i++)</pre>
            cin>>a>>b; // 0 base
           G[a].push_back(b);
           G[b].push_back(a);
       BCC();
    return 0;
BipartiteMatching():
===========
#include <cstdio>
#include <iostream>
#include <vector>
#include <stack>
#include<string.h>
using namespace std;
#define MAX 1200
#define FOR(i,a,b) for(int i=a;i<b;i++)</pre>
#define SZ(a) (int)a.size()
#define mem(a,b) memset(a,b,sizeof(a))
int Right[MAX],used[MAX];
vector<int>G[MAX];
int N,M;
bool dfs(int u)
    if(used[u]) return false;
    used[u]=true;
    FOR(i,0,SZ(G[u]))
        int v=G[u][i];
        if(Right[v]==-1 || dfs(Right[v]))
            Right[v]=u;
            return true;
    }
    return false;
int BPM()
    mem(Right,-1);int cnt=0;
    for(int i=0;i<N;i++)</pre>
        mem(used,0);
        if(dfs(i)) cnt++;
```

```
return cnt;
int main()
    int i,j,u,v;cin>>N>>M;
    for(int i=0;i<=N;i++) G[i].clear();</pre>
    for(int i=0;i<M;i++)</pre>
       cin>>u>>v;
       G[u].push_back(v);
    cout<<BPM()<<endl;</pre>
    return 0;
}
2-SAT:
==========
/// Properties
// (a or b) make two edge: !a-->b , !b-->a
// If there's an edge (a,b) then there's also an edge (!b,!a).
// If there were a path from x to both y and !y,
// then there would have been a path from x to !y and from !y to !x.
// a 2-CNF formula A is unsatisfiable iff
// there exists a variable x, such that:
// there is a path from x to !x in the graph
// there is a path from !x to x in the graph.
#include<stdio.h>
#include<string.h>
#include<iostream>
#include<vector>
using namespace std;
#define SZ(a) push_back
                  (int)a.size()
#define INF (1<<28)
#define MAX 16100
vector<int>adj[MAX],trans[MAX],graph[MAX];
int group_no[MAX],N,M,YES,t_value[MAX];
bool color[16100];
vector<int>S;
void addedge(int u,int v)
    if (u<0) u=(u*-1)+M;
    if (v<0) v=(v*-1)+M;
    adj[u].pb(v);
    trans[v].pb(u);
void dfs(int u)
    color[u]=1;
    for(int i=0;i<adj[u].size();i++)</pre>
       int v=adj[u][i];
       if(!color[v]) dfs(v);
    S.push_back(u);
}
void dfs_2(int u,int grp_no)
    color[u]=1;
```

```
group_no[u]=grp_no;
    graph[grp_no].pb(u);
    for(int i=0;i<trans[u].size();i++)</pre>
        int v=trans[u][i];
        if(!color[v]) dfs_2(v,grp_no);
    return;
}
void dfs_3(int num,int val)
    color[num] = 1;
    for(int i=0;i<graph[num].size();i++)</pre>
        int node = graph[num][i];
        t_value[node] = val;
    for(int i=0;i<graph[num].size();i++)</pre>
        int node = graph[num][i];
        node = (node+M)%N;
        if(!node) node = N;
        if(!color[group_no[node]]) dfs_3(group_no[node],!val);
    return;
}
void SCC()
    int i,cnt,u; S.clear();
    memset(color,0,sizeof(color));
    for(i=1;i<=N;i++) if(!color[i]) dfs(i);</pre>
    memset(color,0,sizeof(color)); cnt=1;
    for(i=S.size()-1;i>=0;i--)
        u=S[i];
        if(!color[u]) dfs_2(u,cnt++);
    YES=1;
    for(i=1;i<=N/2;i++)</pre>
        if(group_no[i]==group_no[i+M]){YES=0;return;}
    mem(t_value,-1);
    mem(color,0);
    for(int i=1;i<cnt;i++) if(!color[i]) dfs_3(i,0);</pre>
    return;
int main()
    int a,b;
    scanf("%d %d",&N,&M);
    for(int i=0;i<=2*M;i++)</pre>
        adj[i].clear();
        trans[i].clear();
        graph[i].clear();
    for(int i=0;i<N;i++)</pre>
        scanf("%d %d",&a,&b);
```

```
addedge(-1*a,b);
        addedge(-1*b,a);
    }
    N = 2*M;
    SCC();
    if(!YES) puts("NO");
    else
        vector<int>ans;ans.clear();
        puts("YES");
        for(int i=1;i<=M;i++)</pre>
             if(t_value[i]==1)
                 ans.pb(i);
        printf("%d",ans.size());
        for(int i=0;i<ans.size();i++) printf(" %d",ans[i]);</pre>
        printf("\n");
    return 0;
}
MatrixExpo:
==========
const int ROW=5, CUL=5;
typedef int type;
struct Matrix
    type mat[ROW][CUL];
    Matrix(){;}
    void view(int dimension) {for(int i=1;i<=dimension;i++)</pre>
    {for(int j=1;j<=dimension;j++)printf("%4d",mat[i][j]);puts("");}}</pre>
};
Matrix MatrixMul(Matrix &A, Matrix &B, int dimension)
    Matrix Res;
    for(int i=1;i<=dimension;i++)</pre>
    for(int j=1;j<=dimension;j++)</pre>
        Res.mat[i][j]=0;
        for(int k=1;k<=dimension;k++)</pre>
            Res.mat[i][j]=(Res.mat[i][j]+(A.mat[i][k]*B.mat[k][j]));
    return Res;
}
Matrix MatrixExpo(Matrix &A, ll n,int dimension)
    Matrix Res, X=A;
    for(int i=1;i<=dimension;i++)</pre>
        for(int j=1;j<=dimension;j++)</pre>
            Res.mat[i][j]=(i==j)?1:0;
             /// Identical Matrix
    while(n)
        if((n&1)==1) Res=MatrixMul(Res,X,dimension);
        X=MatrixMul(X,X,dimension);
        n/=2;
```

```
return Res;
Matrix MakeBase(ll p,ll q)
   Matrix base;
   //make base matrix.
   return base;
BitMask:
=========
int nmask=mask;
while(nmask>=0)
   int tmp=mask&~nmask;
   if(val[nmask]) ret=min(ret,rec(tmp)+1);
   nmask--;
   if(nmask>=0) nmask=nmask&mask;
Base 4 Use in DP:
_____
int POW[11];
void PowerTake()
   POW[0]=1;
   for(int i=1;i<=10;i++)</pre>
       POW[i]=POW[i-1]*4;
   return;
inline int get_bit(int mask,int pos)
   int bit=(mask/POW[pos])%4;
   return bit;
inline int set_bit(int mask,int pos,int val)
   int bit=get_bit(mask,pos);
   mask-=(POW[pos]*bit);
   mask+=(POW[pos]*val);
   return mask;
HistroGram Problem:
Finds largest rectangular area in a histogram in O(n)
typedef long long 11;
long long calc(int *ht, int n)
   ll ret = 0;
   int top = 1, st[MAX], i;
   ht[0] = st[0] = ht[++n] = 0;
   for(i = 1; i <= n; i++)</pre>
```

```
while(top > 1 && ht[st[top-1]] >= ht[i])
           ret = _max(ret, (i64)ht[st[top-1]]*(i64)(i - st[top-2]-1));
           top--;
       st[top++] = i;
    return ret;
}
Matching Pattern(Kmp+DP):
_____
M patarn can not present in N-digits Numbers.
* /
11 dp[MAX][MAX],N,len,M=10000007;
11 b[MAX]; /// use for kmp.
string str;
void KMP()
    mem(b,0);int i,k=0;b[k]=0;
    len=SZ(str);
    for(i=1;i<len;i++)</pre>
       while(k>0 and str[i]!=str[k]) k=b[k-1];
       if(str[k]==str[i]) k++;b[i]=k;
    return;
}
11 rec(int cur,int match)
    if(match>=len) return 0;
   if(cur==N) return 1;
    11 &ret=dp[cur][match];
    if(ret!=-1) return ret;
   ret=0;
    for(ll i=0;i<=9;i++)</pre>
       11 k=match;
       while(k>0 and (str[k]-'0')!=i) k=b[k-1];
       if((str[k]-'0')==i) k++;
       ret=(rec(cur+1,k)+ret)%M;
    return ret%M;
}
Maximal Sub-Rectangle:
// 11951 - Area(UVA).
// Maximum area with minimum cost
#define lim 110
int n,m,K, arr[lim][lim];
int cumarr[lim][lim];
int maxarea,mincost;
int update(int area,int cost)
    if(area>maxarea){
       maxarea=area;
```

```
mincost=cost;
    else if(area==maxarea) mincost=min(mincost,cost);
int main()
    cin>>n>>m>>K;maxarea=0;mincost=0;
    mem(cumarr,0);int i,j,k,l;
    for(i=1;i<=n;i++)</pre>
        for(j=1;j<=m;j++)</pre>
            cin>>arr[i][j];
    for(i=1;i<=n;i++)</pre>
        for(j=1;j<=m;j++)</pre>
            cumarr[i][j]=cumarr[i][j-1]+arr[i][j];
    for(i=1;i<=m;i++){</pre>
        for(j=i;j<=m;j++)</pre>
            1=1;
            int area=0,cost=0;
            //optimal row chosing by 2 pointer
            for(k=1;k<=n;k++)
                if(k!=1) area-=(j-i+1); //cancelling previous
                cost-=(cumarr[k-1][j]-cumarr[k-1][i-1]);
                while(l<=n)</pre>
                     if(cost+cumarr[1][j]-cumarr[1][i-1]>K)break;
                     area+=(j-i+1);
                     cost+=(cumarr[1][j]-cumarr[1][i-1]);
                     1++;
                update(area,cost);
            }
    cout<<maxarea<<" "<<mincost<<endl;</pre>
    return 0;
RectangleUnion:
===========
/*
Given x1,x2,y1,y2.A person will advertise in all the houses whose x co-ordinate
is between x1 and x2 (inclusive) and y co-ordinate is between y1 and y2 (inclusive).
Output the total number of houses that are advertised by at least k people.
//complexity: k*nlog(n)
* /
struct cord
    int x, y1, y2, val; //val for starting or ending
    cord(int _x=0, int _y1=0, int _y2=0, int _val=0)
        x=_x, y1=_y1, y2=_y2, val=_val;
    }
};
cord pnt[MAX];
int ans[4*MAX], upd[4*MAX];
vector<int>y;
```

```
bool cmp(cord a, cord b)
    if(a.x==b.x) return (a.val > b.val);
    return (a.x < b.x);</pre>
int k;
int update(int node, int st, int end, int i, int j, int val)
    if(j<=y[st] || i>=y[end]) return ans[node];
    if(y[st]>=i && y[end]<=j)</pre>
        upd[node]+=val;
        if(upd[node]>=k) /// at least k person.
            return ans[node] = y[end]-y[st];
        else
            if(end-st==1)return ans[node]=0;
            int mid=(st+end)>>1, ret1, ret2;
            if(upd[node] == k-1&&val == -1) upd[2*node] ++;
            if(upd[node] == k-1&&val == -1) upd[2*node+1]++;
            ret1 = update(2*node, st, mid, i, j, val);
            ret2 = update(2*node+1, mid, end, i, j, val);
            return ans[node] = ret1+ret2;
    }
    int mid=(st+end)>>1, ret1, ret2;
    ret1 = update(2*node, st, mid, i, j, val);
    ret2 = update(2*node+1, mid, end, i, j, val);
        //special attention to mid
    if(upd[node] < k) ans[node] = ret1+ret2;</pre>
    else ans[node]=y[end]-y[st];
    return ans[node];
int main()
    int i, j, n, x1, y1, x2, y2, cnt=0, m;
    scanf("%d %d", &n,&k);y.clear();
    for(i=0;i<n;i++)</pre>
        scanf("%d%d%d%d", &x1, &y1, &x2, &y2);
        pnt[cnt++] = cord(x1, y1, y2, 1);
        pnt[cnt++] = cord(x2, y1, y2, -1);
        y.pb(y1), y.pb(y2);
    sort(y.begin(), y.end());
    y.resize(unique(y.begin(), y.end())-y.begin());
    n = SZ(y);sort(&pnt[0], &pnt[0]+cnt, cmp);
    memset(ans, 0, sizeof ans); memset(upd, 0, sizeof upd);
    ll sum=0, now; x1=-1; //any value
    for(i=0;i<cnt;i++)</pre>
        x2 = pnt[i].x;
        now = x2-x1;
        sum+=now*ans[1];
        update(1, 0, n-1, pnt[i].y1, pnt[i].y2, pnt[i].val);
        x1 = x2;
    cout<<sum<<endl;</pre>
```

```
return 0;
```