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1 Startup templates

1.1 template

```
1 #include <vector>
2 #include < list >
3 #include <map>
4 #include <set>
5 #include <deque>
6 #include <stack>
  #include <bitset>
  #include <algorithm>
9 #include <functional>
10 #include <numeric>
11 #include <utility>
12 #include <sstream>
13 #include <iostream>
14 #include <cstdio>
15 #include <cmath>
16 #include <cstdlib>
  #include <cstring>
17
18
  #include <cassert>
19
20 using namespace std;
21
22
  typedef long long ll;
23
  typedef pair<int, int> pii;
  template<typename T> int size(T& a) { return (int) a.size(); }
26
  template<typename T> T sqr(T a) { return a * a; }
28 #define _(a, b) memset((a), (b), sizeof(a))
29 #define fs first
30 #define sc second
31 #define pb push_back
  #define mp make_pair
33 #define all(a) a.begin(), a.end()
34 #define REP(i, a, b) for(int i = (a); i < (b); ++i)
35 #define REPD(i, a, b) for (int i = (b) - 1; i >= a; --i)
36 #define ve vector
```

1.2 gvimrc

```
set autoread
set autoindent
set autochdir
set cindent
set number
syntax on
set shiftwidth =4
set tabstop =4
colorscheme desert
set gfn =Monospace\ 12
```

2 Graph Algorithms

2.1 Kuhn Max Matching

```
// need: graph( head, to, nxt ), used, mate
   bool kuhn(int v) {
 3
       if (used[v]) return false;
 4
        used[v] = true;
 5
        for(int e = head[v]; e != -1; e = nxt[e]) {
 6
            if (mate[to[e]] = -1) {
 7
                 mate\,[\,to\,[\,e\,]\,]\ =\ v\,;
 8
                 return true;
 9
            }
10
11
       for(int e = head[v]; e != -1; e = nxt[e]) {
12
            if (kuhn(mate[to[e]])) {
13
                 mate[to[e]] = v;
14
                 return true;
15
16
17
       return false;
18
19
20
   int max_matching() {
21
       int res = 0;
22
       REP(\,i\;,\;\;0\,,\;\;n\,)\;\;\{
23
            _(used, false);
24
            if (kuhn(i)) res++;
25
26
       return res;
27
```

2.2 Dinic Max-Flow

```
// need: graph( head, nxt, to, capa, flow ), dist, q
   bool bfs(int src, int dest) {
         \begin{array}{l} -(\operatorname{dist}, -1); \\ \operatorname{dist}[\operatorname{src}] = 0; \end{array}
 3
 4
         int \dot{H} = 0;
 5
 6
         q[H ++] = src;
 7
         REP(i, 0, H) {
 8
              int cur = q[i];
 9
              for (int e = head[cur]; e != -1; e = nxt[e]) {
10
                    if (capa[e] > flow[e] && dist[to[e]] = -1) {
                         dist[to[e]] = dist[cur] + 1;
11
                         q[H ++] = to[e];
12
13
              }
14
15
         return dist[dest] >= 0;
16
17
18
19
   int dfs(int cur, int curflow) {
20
         if (cur == dest) return curflow;
21
         int d;
         \mathbf{for}\,(\,\mathbf{int}\&\,\,e\,\,=\,\,\mathrm{work}\,[\,\mathrm{cur}\,]\,;\ e\,\,!{=}\,\,-1;\,\,e\,\,=\,\,\mathrm{nxt}\,[\,e\,]\,)\ \ \{
22
23
              if (capa[e] > flow[e] && (dist[to[e]] = dist[cur] + 1) &&
24
                         (d = dfs(to[e], min(curflow, capa[e] - flow[e])))) {
                    flow [e] += d;
flow [e ^ 1] -= d;
25
26
27
                    return d;
28
29
30
         return 0;
31
32
33
   int dinic() {
34
         int res = 0;
35
         while (bfs(src, des)) {
36
              int d:
37
              memcpy(work, head, sizeof(head));
```

2.3 Hungary Algo

```
// need: a[n][m], all indices start with 1
   vector < int > u (n+1), v (m+1), p (m+1), way (m+1);
3
   for (int i=1; i<=n; ++i) {
     p[0] = i;
     int j0 = 0;
6
     vector < int > minv (m+1, INF);
     vector < char > used (m+1, false);
7
8
9
       used[j0] = true;
10
       int i0 = p[j0], delta = INF, j1;
11
       for (int j=1; j < m; ++j)
          if (!used[j]) {
12
            int cur = a[i0][j]-u[i0]-v[j];
13
            if \ (\operatorname{cur} < \operatorname{minv}[j])
14
15
              minv[j] = cur, way[j] = j0;
16
            if (minv[j] < delta)
17
              delta = minv[j], j1 = j;
18
19
       for (int j=0; j < m; ++j)
20
          if (used[j])
21
           u[p[j]] += delta, v[j] -= delta;
22
23
           minv[j] -= delta;
24
       j0 = j1;
25
     } while (p[j0] != 0);
26
     do {
27
       int j1 = way[j0];
28
       p[j0] = p[j1];
29
       j0 = j1;
30
     } while (j0);
31
32
   // restore ans[] -- selected column for each row
33 for (int j=1; j < m; ++j)
    ans[p[j]] = j;
35
   // cost
36|\inf \cos t = -v[0];
```

2.4 Min Cut

```
pair<int, ve<int> > GetMinCut(ve< ve<int> > &weights) {
     int N = weights.size();
     ve<int> used(N), cut, best_cut;
3
4
     int best_weight = -1;
5
6
     REPD(phase, 0, N) {
7
        ve < int > w = weights [0];
8
        ve<int> added = used;
9
        int prev, last = 0;
10
        REP(i, 0, phase) {
11
          prev = last;
12
          \begin{array}{l} last = -1; \\ REP(j \; , \; 1 \; , \; N) \end{array}
13
             if (!added[j] && (last == -1 || w[j] > w[last])) last = j;
14
15
          if (i = phase -1) {
16
            REP(j, 0, N) weights[prev][j] += weights[last][j];
            REP(j\;,\;\;0\;,\;\;N)\;\;weights[\,j\,][\,prev\,]\;=\;weights[\,prev\,][\,j\,];
17
18
             used[last] = true;
19
             cut.pb(last);
20
             if (best\_weight == -1 \mid \mid w[last] < best\_weight) {
```

```
21
              best_cut = cut;
22
              best\_weight = w[last];
23
            }
24
         } else {
25
           REP(j, 0, N)
26
             w[j] += weights[last][j];
27
           added[last] = true;
28
29
30
     }
31
     return mp(best_weight, best_cut);
32
```

2.5 Bridges

```
//\ need:\ graph\left(\ head\,,\ to\,,\ nxt\ \right),\ used\,,\ tin\,,\ fup\,,\ timer
   void dfs (int v, int par = -1) {
 3
     used[v] = true;
     tin[v] = fup[v] = timer++;
 4
     for(int e = head[v]; e != -1; e = nxt[e]) {
 6
       int u = to[e];
 7
        \quad if \ (u == par) \quad continue;
 8
        if (used[u])
 9
          fup[v] = min (fup[v], tin[u]);
10
        else {
11
          dfs(u, v);
12
          fup[v] = min (fup[v], fup[u]);
          if (fup[u] > tin[v])
13
            IS_BRIDGE(v, u);
14
15
16
17
18
19
   void find_bridges() {
20
     timer = 0;
21
      _{-}(used, 0)
22
     REP(i, 0, N)
       if (!used[i]) dfs (i);
23
24
```

2.6 Cut Vertices

```
//\ need:\ graph\ (head,\ to,\ nxt)\,,\ tin\,,\ fup\,,\ used\,,\ timer
   void dfs (int v, int par = -1) {
      used[v] = true;
 3
 4
       tin[v] = fup[v] = timer++;
 5
      int children = 0;
      \mbox{ for (int } e = head [v]; \ e != -1; \ e = nxt[e]) \ \{
 6
 7
         int u = to[e];
 8
         if (u == par) continue;
         if (used[u])
 9
10
            fup[v] = min (fup[v], tin[u]);
11
         else {
12
            dfs(u, v);
            \texttt{fup}\left[\overset{.}{v}\right] \; = \; \min \; \; \left(\; \texttt{fup}\left[\,v\,\right] \;, \; \; \texttt{fup}\left[\,u\,\right] \right) \;;
13
14
            if (\sup [u] > = \lim [v] \&\& p ! = -1)
              IS_CUTPOINT(v);
15
           ++children;
16
17
18
19
       if (p = -1 \&\& children > 1)
         IS_CUTPOINT(v);
20
21
22
23
   int main() {
24
      timer = 0;
       _(used, 0);
25
26
       dfs (0);
27
```

2.7 Min-Cost Max-Flow

```
// need: graph (head, nxt, to, from, capa, cost, flow)
   // pi, dist, prve
 3 void updatePotentials() {
     memcpy(pi, dist, sizeof(int) * N);
 4
 5
 6
 7
   bool fordBellman(int src, int dst) {
     REP(i, 0, N) dist[i] = INF;
 9
      dist[src] = 0;
10
     bool changed;
11
12
     REP(phase, 0, N) {
13
        changed = false;
14
       REP(v, 0, N) {
15
          if(dist[v] == INF) continue;
16
          for(int e = head[v]; e != -1; e = nxt[e]) {
17
            \quad \textbf{int} \ u \, = \, to \, [\, e \, ] \, ; \quad
             if(capa[e] > flow[e] \&\& dist[u] > dist[v] + cost[e]) {
18
19
               dist[u] = dist[v] + cost[e];
20
               prve[u] = e;
21
               changed = true;
22
23
          }
24
25
        if (!changed) break;
26
27
     return ! changed;
28
29
30
   set < pii > q;
31
   bool dijkstra (int src, int dst) {
     REP(i, 0, N) dist[i] = INF;
33
     dist[src] = 0;
34
     q.insert(mp(0, 0));
35
36
     \mathbf{while}\,(\,\mathrm{size}\,(\,\mathrm{q}\,)\,)\  \, \{\,
37
        pii tmp = (*q.begin());
38
        \mathbf{int} \ v = \mathrm{tmp.sc} \, , \ d = \mathrm{tmp.fs} \, ;
39
        q.erase(q.begin());
40
        if (d != dist[v]) continue;
41
42
        for(int e = head[v]; e != -1; e = nxt[e]) {
43
          \mathbf{int}\ u\,=\,\mathrm{to}\,[\,e\,]\,;
44
          if(capa[e] > flow[e] \&\& dist[u] > dist[v] + cost[e] - pi[v] + pi[u]) {
45
             dist[u] = dist[v] + cost[e] - pi[v] + pi[u];
            prve[u] = e;
46
47
            q.insert(mp(dist[u], u));
48
49
50
51
     return dist[dst] != INF;
52
53
54
   pii minCostMaxFlow(int src, int dst) {
55
      if(!fordBellman(src, dst)) return mp(0, 0);
     int sumFlow = 0, sumCost = 0;
56
57
58
     do {
59
        int curFlow = INF, curCost = 0;
        int cur = dst;
60
61
        while(cur != src) {
62
          int e = prve[cur];
63
          curFlow = min(curFlow, capa[e] - flow[e]);
64
          curCost += cost[e];
65
          cur = from [e];
66
67
        cur = dst;
68
        while(cur != src) {
69
          int e = prve[cur];
          flow[e] += curFlow;
flow[e ^ 1] -= curFlow;
70
71
          cur = from [e];
72
```

```
73      }
74      sumCost += curFlow * curCost;
75      updatePotentials();
76      } while(dijkstra(src, dst));
77
78      return mp(sumFlow, sumCost);
79    }
```

2.8 Strongly Connected Components

```
// need: graph (head, to, nxt), reverse graph (rhead, rto, rnxt)
   // used, compID, order
   void dfs1(int v) {
 3
     used[v] = true;
 5
     for(int e = head[v]; e != -1; e = nxt[e]) {
 6
       if (!used[to[e]]) dfs1(to[e]);
 7
 8
     order.pb(v);
 9
10
  void dfs2(int v, int id) {
11
     used[v] = true;
12
     compID[v] = id;
13
     for(int e = rhead[v]; e != -1; e = rnxt[e]) {
14
       if (!used[ rto[e] ]) dfs2(rto[e]);
15
16 }
17
   void main() {
18
     _(used, false);
     REP(i, 0, N)
19
20
       if (!used[i]) dfs1(i);
21
     _(used, false);
     \dot{\mathbf{n}}\dot{\mathbf{t}} id = 0;
22
23
     REPD(i, 0, N) {
24
       int v = order[i];
25
       if(!used[v]) dfs2(v, id++);
26
27
```

2.9 2-SAT

```
Problem: (a \lor c) \& (a \lor !b) \& ...
Edges: (a \lor b) is equivalent to (!a \Rightarrow b) \lor (!b \Rightarrow a)
Solution: there is no solution iff for some x \ compID[x] = compID[!x], else see code below
```

```
// need: graph, scc
   int main() {
 2
      _(used, false);
     REP(i, 0, N)
 5
        if (!used[i]) dfs1 (i);
 6
 7
      _{-}(\text{compID}, -1);
 8
     int id = 0;
 9
     REPD(i, 0, N) {
10
        int v = order[i];
11
        \mathbf{if} \pmod{[v]} = -1 \operatorname{dfs2}(v, id++);
12
     }
13
14
     REP(i, 0, N)
15
        if (compID[i] = compID[i^1]) 
          puts ("NO SOLUTION");
16
17
          return 0;
18
     \overrightarrow{REP}(i, 0, N) {
19
20
        int ans = comp[i] > comp[i^1] ? i : i^1;
21
        printf ("%d ", ans);
22
23
```

3 Linear Algebra

3.1 Gauss Elimination

```
//Ax = B. RETURN: determinant, A \rightarrow A^{(-1)}, B \rightarrow solution
   typedef double T;
    typedef vector<T> VT;
    typedef vector <VT> VVT;
 5
 6
   T Gauss Jordan (VVT &a, VT &b) {
 7
       \mathbf{const} int n = a.size();
 8
       ve < int > irow(n), icol(n), ipiv(n);
 9
      T \det = 1;
10
11
      REP(i, 0, n) {
12
          int pj = -1, pk = -1;
13
          REP(j, 0, n) if (!ipiv[j])
             REP(k, 0, n) if (!ipiv[k])
14
15
                if (pj = -1 || fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j; pk = k; }
16
           if \ (fabs(a[pj][pk]) < EPS) \ \{ \ cerr << "Matrix is singular." << endl; \ exit(0); \ \} 
          ipiv [pk]++;
17
          swap(a[pj], a[pk]);
swap(b[pj], b[pk]);
18
19
20
          if (pj != pk) det *= -1;
21
          irow[i] = pj;
22
          icol[i] = pk;
23
         T\ c\ =\ 1.0\ /\ a\,[\,pk\,]\,[\,pk\,]\,;
24
25
          \det \ *= \ a [pk][pk];
26
          a[pk][pk] = 1.0;
          \begin{array}{lll} REP(\,p\,,\;\;0\,,\;\;n) & a\,[\,pk\,]\,[\,p\,] \;\; *=\; c\,; \\ b\,[\,pk\,] \;\; *=\; c\,; \end{array} 
27
28
29
         REP(p, 0, n) if (p != pk) {
30
             c = a[p][pk];
31
             a\,[\,p\,]\,[\,pk\,] \ = \ 0\,;
32
             REP(\,q\,,\ 0\,,\ n\,)\ a\,[\,p\,]\,[\,q\,]\ -\!=\ a\,[\,pk\,]\,[\,q\,]\ *\ c\,;
33
             b[p] = b[pk] * c;
34
35
36
      \begin{array}{lll} \text{REPD}(p, \ 0, \ n) & \textbf{if} \ (irow\,[p] \ != \ icol\,[p]) \ \{ \\ \text{REP}(k, \ 0, \ n) \ swap(a\,[k][\,irow\,[p]] \, , \ a\,[k][\,icol\,[p]]) \, ; \end{array}
37
38
39
40
41
       return det;
42
```

3.2 Fast-Fourier Transform

```
typedef complex<double> base;
 3
   void fft (vector<base> & a, bool invert) {
 4
      int n = (int) a.size();
 6
      for (int i = 1, j = 0; i < n; i++) {
 7
         \quad \textbf{int} \quad \text{bit} \ = \ n \ >> \ 1;
 8
         for (; j \ge bit; bit \ge 1)
          j -= bit;
 9
10
         j += bit;
11
         if (i < j) swap (a[i], a[j]);
12
13
14
      \mathbf{for} \ (\mathbf{int} \ \operatorname{len} \ = \ 2; \ \operatorname{len} \ <= \ \mathrm{n}; \ \operatorname{len} \ <<= \ 1) \ \{
15
         double ang = 2 * PI/len * (invert ? -1 : 1);
16
         base \ wlen \ (\cos(ang)\,, \ \sin(ang))\,;
17
         for (int i = 0; i < n; i += len) {
           base w (1);
18
19
           for (int j = 0; j < len/2; j++) {
20
              base u = a[i+j], v = a[i+j+len/2] * w;
21
              a[i+j] = u + v;
22
              a[i+j+len/2] = u - v;
```

3.3 Simplex

```
// maximize c^T x
   //Ax \ll b
   // x >= 0
 3
 5
   struct LPSolver {
      int m, n;
 7
      ve < int > B, N;
      ve < ve < double > > D;
 8
 9
10
       LPSolver(\textbf{const} \ ve < ve < \textbf{double} > > \&A, \ \textbf{const} \ ve < \textbf{double} > \&b \,, \ \textbf{const} \ ve < \textbf{double} > \&c) \ :
         m(\texttt{b.size())}\,,\,\,n(\texttt{c.size())}\,,\,\,N(\texttt{n+1})\,,\,\,B(\texttt{m})\,,\,\,D(\texttt{m+2},\,\,\texttt{ve}<\!\!\textbf{double}>\!\!(\texttt{n+2}))\  \, \{
11
12
           REP(i, 0, m) REP(j, 0, n) D[i][j] = A[i][j];
13
            REP(i, 0, m) \{ B[i] = n+i; D[i][n] = -1; D[i][n+1] = b[i]; \}
14
           REP(j, 0, n) \{ N[j] = j; D[m][j] = -c[j]; \}
15
           N[n] = -1; D[m+1][n] = 1;
16
17
      void Pivot(int r, int s)
18
19
         REP(i, 0, m + 2) if (i != r)
           \overrightarrow{REP}(j, 0, n + 2) if (j != s)
20
21
              D[i][j] = D[r][j] * D[i][s] / D[r][s];
22
         REP(j\;,\;\;0\;,\;\;n\;+\;2)\quad i\;f\quad (\;j\;\;!=\;s\;)\;\;D[\;r\;][\;j\;]\;\;/=\;D[\;r\;][\;s\;]\;;
23
         REP(i, 0, n + 2) \quad \textbf{if} \quad (i != r) \quad D[i][s] /= -D[r][s];
         D[r][s] = 1.0 / D[r][s];
24
25
         swap(B[r], N[s]);
26
27
28
      bool Simplex(int phase) {
29
         int x = phase == 1 ? m+1 : m;
         while (true) {
30
31
            int s = -1;
32
            REP(j, 0, n + 1) {
33
               if (phase = 2 && N[j] = -1) continue;
34
                \mbox{if } (s = -1 \ || \ D[x][j] < D[x][s] \ || \ D[x][j] = D[x][s] \ \&\& \ N[j] < N[s]) \ s = j; 
35
36
            if (D[x][s] >= -EPS) return true;
37
            \mathbf{int} \quad r \ = \ -1;
38
           REP(i, 0, m) {
39
               if (D[i][s] \le 0) continue;
                if \ (r = -1 \ || \ D[i][n+1] \ / \ D[i][s] < D[r][n+1] \ / \ D[r][s] \ || 
40
41
                    D[i][n+1] / D[i][s] == D[r][n+1] / D[r][s] && B[i] < B[r]) r = i;
42
43
            if (r = -1) return false;
44
            Pivot(r, s);
45
         }
46
      }
47
48
      double Solve (ve<double> &x) {
49
         int r = 0;
50
         REP(i, 1, m) if (D[i][n+1] < D[r][n+1]) r = i;
51
         if (D[r][n+1] \le -EPS) {
52
            Pivot(r, n);
            \textbf{if} \hspace{0.2cm} (!\, Simplex\, (1) \hspace{0.2cm} || \hspace{0.2cm} D[m+1][n+1] \hspace{0.2cm} < -EPS) \hspace{0.2cm} \textbf{return} \hspace{0.2cm} - numeric\_limits \\ < \textbf{double} > :: infinity\, () \hspace{0.2cm} ; \\
53
           REP(i, 0, m) if (B[i] = -1) {
54
55
               int s = -1;
56
              REP(\,j\;,\;\;0\,,\;\;n\;+\;1)
57
                  if \ (s = -1 \ || \ D[i][j] < D[i][s] \ || \ D[i][j] = D[i][s] \ \&\& \ N[j] < N[s]) \ s = j; \\ 
58
               Pivot(i, s);
59
            }
60
61
         if (!Simplex(2)) return numeric_limits < double > :: infinity();
62
         x = ve < double > (n);
```

```
63 REP(i, 0, m) if (B[i] < n) x[B[i]] = D[i][n+1];
64 return D[m][n+1];
65 }
66 };
```

4 String Algorithms

4.1 Suffix Array

```
struct entry {
     int nr[2], p;
   } L[MAXN], tmp[MAXN];
   #define eq(a, b) ((a).nr[0] = (b).nr[0] && (a).nr[1] = (b).nr[1])
 4
 5 int cnt [MAXN], p[2][MAXN];
 7
   \mathbf{void} \ \ \mathrm{radixPass} \, (\, \mathrm{entry} \ * \ \mathrm{a} \, , \ \ \mathbf{int} \ \ \mathrm{N}, \ \ \mathbf{int} \ \ \mathrm{pass} \, , \ \ \mathbf{int} \ \ \mathrm{K}, \ \ \mathrm{entry} \ * \ \mathrm{b}) \ \ \{
 8
     memset(cnt, 0, (K + 1) * sizeof(int));
     REP(i, 0, N) cnt[a[i].nr[pass]]++;
 9
10
     int sum = 0;
11
     REP(i, 0, K+1) {
12
       sum += cnt[i];
13
        cnt[i] = sum - cnt[i];
14
     REP(i, 0, N) b[cnt[a[i].nr[pass]]++] = a[i];
15
16|}
17
18
   void makeSA(int * s , int N, int * suftab , int * isuftab) {
19
     REP(i, 0, N) p[0][i] = s[i];
20
     int k = 200;
     21
22
23
24
25
          L[i].p = i;
26
27
        radixPass(L, N, 1, k, tmp);
28
        radixPass(tmp, N, 0, k, L);
29
        k = 1;
30
       REP(i, 0, N)
          \label{eq:pstep} p \, [\, step \, ] \, [\, L \, [\, i \, \, ] \, . \, p \, ] \, = \, i \, > \, 0 \, \, \&\& \, \, eq \, (L \, [\, i \, ] \, , \, \, L \, [\, i \, - \, 1 \, ] \, ) \, \, \, ?
31
32
             p[step][L[i-1].p] : k++;
33
34
        if(k > N) break;
35
36
     REP(i, 0, N) suftab[i] = L[i].p;
37
     REP(i, 0, N) isuftab [suftab [i]] = i;
38
39
40 void makeLCP(int * s, int * suftab, int * isuftab, int N, int * lcptab) {
41
     int cur = 0;
42
     REP(i, 0, N)
43
        if (isuftab [i] == 0) continue;
        44
45
        while(ii < N && jj < N && s[ii] == s[jj]) ii++, jj++, cur++;
46
        lcptab[isuftab[i]] = cur --;
47
        if(cur < 0) cur = 0;
48
49
```

4.2 Suffix Tree from Suffix Array

```
1
   struct Seg {
 2
      \mathbf{int}\ \mathrm{lb}\ ,\ \mathrm{rb}\ ,\ \mathrm{lcp}\ ;
      vector < Seg*> childList;
 4
      void init(int l, int i, int j) {
 5
         lb = i; rb = j; lcp = l;
 6
 7
      void add(Seg * son) {
 8
         childList.pb(son);
 9
10
11 typedef Seg* pSeg;
12
13 struct Stack {
      pSeg\ segs\left[ M\!A\!X\!N<<\ 1\,\right];
14
15
      int size;
```

```
void push(pSeg seg) {
17
        segs[size++] = seg;
18
19
     pSeg pop() {
20
       return segs[--size];
21
22
     pSeg top() {
23
       return segs[size - 1];
24
25
   } stack;
26
27
   pSeg top() { return stack.top(); }
28
   void push(pSeg seg) { stack.push(seg); }
29
   pSeg pop() { return stack.pop(); }
31
   pSeg init(int lcp, int lb, int rb) {
32
     pSeg ret = new Seg;
33
     ret->init(lcp, lb, rb);
34
     \textbf{return} \hspace{0.1in} \texttt{ret} \; ;
35
36
   pSeg makeTree() {
37
38
     stack.size = 0;
     pSeg lastInterval = NULL;
39
     int lastSingleton = 0;
41
     stack.push(init(0, 0, -1));
42
     REP(i\ ,\ 1\ ,\ N)\ \{
43
        int lb = i - 1;
        pSeg singleton = init(N - suftab[i - 1] - 1, i - 1, i - 1);
44
45
        //process\left(singleton\right);
46
        \mathbf{while}(\,\mathrm{lcptab}\,[\,\mathrm{i}\,]\,<\,\mathrm{top}\,(\,)-\!\!>\!\!\mathrm{lcp}\,)\  \, \{
47
          if(singleton != NULL) {
48
49
             top()->add(singleton);
50
             singleton = NULL;
51
          top()->rb = i - 1;
52
          lastInterval = pop();
53
54
          //process(lastInterval);
55
          lb = lastInterval->lb;
56
          if(lcptab[i] <= top()->lcp) {
57
             top()->add(lastInterval);
58
             lastInterval = NULL;
59
60
61
        if(lcptab[i] > top()->lcp)
62
          if(lastInterval != NULL)
63
            pSeg seg = init(lcptab[i], lb, -1);
64
            seg->add(lastInterval);
65
            push (seg);
66
             lastInterval = NULL;
67
          } else push(init(lcptab[i], lb, -1));
68
69
        if(singleton != NULL) {
70
          top()->add(singleton);
71
72
73
     assert (stack.size == 1);
74
     //process(top());
75
     return top();
76
```

4.3 Z-function

```
9 | if (i+z[i]-1 > r)

10 | l = i, r = i+z[i]-1;

11 | }

12 | return z;

13 | }
```

4.4 Suffix Automata

```
struct state {
     int len, link;
3
     map<char, int> next;
4
   };
   state st [MAXLEN * 2];
  int sz, last;
8
9
   void sa_init() {
10
     sz = last = 0;
     st[0].len = 0;
11
12
     st[0].link = -1;
13
    ++sz;
14
15
16
   {f void} sa_extend ({f char} c) {
17
     int cur = sz++;
18
     st[cur].len = st[last].len + 1;
19
     int p;
     for (p=last; p!=-1 && !st[p].next.count(c); p=st[p].link)
20
21
       st[p].next[c] = cur;
22
     if (p = -1)
       st[cur].link = 0;
23
24
     else {
25
       int q = st[p].next[c];
       if (st[p].len + 1 = st[q].len)
26
27
          st[cur].link = q;
28
       \mathbf{else} \ \{
29
         int clone = sz++;
30
         st[clone].len = st[p].len + 1;
         st[clone].next = st[q].next;
31
32
          st[clone].link = st[q].link;
33
         for (; p!=-1 \&\& st[p].next[c]==q; p=st[p].link)
34
            st[p].next[c] = clone;
35
         st[q]. link = st[cur]. link = clone;
36
37
38
     last = cur;
39
```

4.5 Palindromes

```
for (i = 0; i < n; i++)
 2
      if(i > r) k = 1;
 3
      else k = \min(d1[l + r - i], r - i);
 4
 5
      while (0 \le i-k \&\& i+k \le n \&\& s[i-k] == s[i+k]) k++;
 6
      d1[i] = k;
 7
      \mathbf{if}(\mathbf{i} + \mathbf{k} - 1 > \mathbf{r})
 8
         r = i + k - 1, l = i - k + 1;
9
10
11
   \mbox{for}\,(\,i \ = \ 0\,; \ i \ < \ n\,; \ i++)\{
12
      if(i > r) k = 0;
13
      else k = \min(d2[1 + r - i + 1], r - i + 1);
14
15
      while (i + k < n \& i - k - 1 >= 0 \& s[i+k] == s[i - k - 1]) k++;
16
      d2[i] = k;
17
18
      \mathbf{if}(i + k - 1 > r)
19
         l \; = \; i \; - \; k \, , \; \; r \; = \; i \; + \; k \; - \; 1 \, ;
20
```

4.6 Lyndon decomposition & Duval

```
// Lyndon decomposition
   for(int i = 0; i < n;) {
     int j=i+1, k=i;
      while (j < n \&\& s[k] <= s[j]) {
 5
        if (s[k] < s[j])
 6
          k = i;
        _{
m else}
          ++k;
9
        ++j;
10
11
      while (i \le k) {
12
        cout << s.substr (i, j-k) << ' ';
13
        i += j - k;
14
15 }
16
17
   string min_cyclic_shift (string s) {
18
      s += s;
19
     int n = (int) s.length();
20
      int i=0, ans=0;
21
      while (i < n/2) {
        ans = i;
23
        int j=i+1, k=i;
        while (j < n \&\& s[k] <= s[j]) { if (s[k] < s[j])
24
25
26
             k \; = \; i \; ;
27
           _{
m else}
28
             +\!\!+\!\!k;
29
          ++j;
30
31
        \mathbf{while} \ (\mathtt{i} <= \mathtt{k}) \quad \mathtt{i} \ +\!\!\!= \mathtt{j} \ - \mathtt{k};
32
33
     return s.substr (ans, n/2);
34 }
```

5 Modular

```
All algorithms described here work on nonnegative integers.
 3
   // return a % b (positive value)
  int mod(int a, int b) {
     return ((a%b)+b)%b;
 6
 8
   // computes gcd(a,b)
 9 int gcd(int a, int b) {
10
     int tmp;
11
     while(b){a%=b; tmp=a; a=b; b=tmp;}
12
13 }
14
15
   // computes lcm(a,b)
16 int lcm(int a, int b) {
17
     return a/\gcd(a,b)*b;
18
19
   // returns d = gcd(a,b); finds x,y such that d = ax + by
20
21
   int extended_euclid(int a, int b, int &x, int &y) {
22
     int xx = y = 0;
23
     int yy = x = 1;
24
     while (b) {
25
        int q = a/b;
26
        int t = b; b = a\%b; a = t;
27
        t = xx; xx = x-q*xx; x = t;
28
        t = yy; yy = y-q*yy; y = t;
29
30
     return a;
31
32
33 // finds all solutions to ax = b \pmod{n}
34 ve<int> modular_linear_equation_solver(int a, int b, int n) {
35
     \mathbf{int}\ \mathbf{x}\,,\ \mathbf{y}\,;
36
     ve<int> solutions;
37
     int d = extended_euclid(a, n, x, y);
38
     if (!(b%d)) {
39
        x = mod (x*(b/d), n);
40
        \  \  \, \textbf{for}\  \, (\,\textbf{int}\  \  \, \textbf{i}\,\,=\,\,0\,;\  \  \, \textbf{i}\,\,<\,\, \textbf{d}\,;\  \  \, \textbf{i}\,++)
41
          solutions.push_back(mod(x + i*(n/d), n));
42
43
     return solutions;
44
45
46
   // computes b such that ab = 1 \pmod{n}, returns -1 on failure
47 int mod_inverse(int a, int n) {
     int x, y;
     \mathbf{int} \ d = \mathtt{extended\_euclid} \left( \mathtt{a} \,, \ \mathtt{n} \,, \ \mathtt{x} \,, \ \mathtt{y} \right);
49
50
     if (d > 1) return -1;
51
     return \mod(x,n);
52|}
54
   // find z such that z % x = a, z % y = b. Here, z is unique modulo M = lcm(x, y).
55
   // Return (z,M). On failure, M=-1.
   pii chinese_remainder_theorem(int x, int a, int y, int b) {
56
57
     int s, t;
58
     int d = extended_euclid(x, y, s, t);
59
     if (a\%d != b\%d) return make_pair(0, -1);
60
     return make_pair(mod(s*b*x+t*a*y,x*y)/d, x*y/d);
61
62
63 // Chinese remainder theorem: find z such that
   /\!/\;z\;\%\;x[\,i\,]\;=\;a[\,i\,]\;\;for\;\;all\;\;i\,.\quad Note\;\;that\;\;the\;\;solution\;\;is
64
   // unique modulo M = lcm_i(x[i]). Return (z,M). On failure, M = -1.
66 pii chinese_remainder_theorem (const ve<int> &x, const ve<int> &a) {
     pii ret = make_pair(a[0], x[0]);
68
     for (int i = 1; i < x.size(); i++) {
69
        ret = chinese_remainder_theorem(ret.second, ret.first, x[i], a[i]);
70
        if (ret.second = -1) break;
71
     }
     return ret;
```