

# Public Key Encryption

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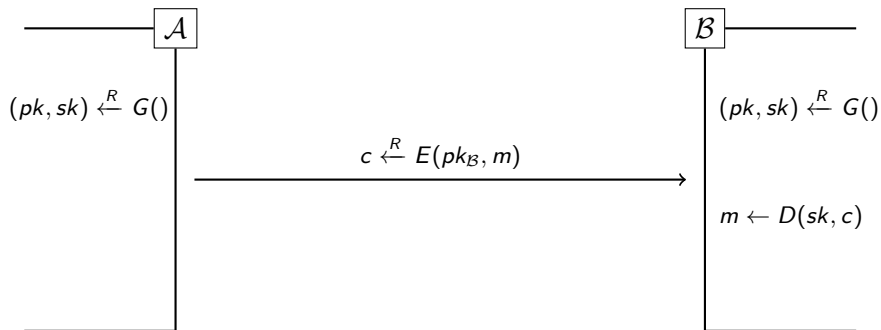
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# 1 Overview

# Recap

- Last lecture we discussed an example of how to exchange secret keys "in the clear"
- This allows us to re-use our symmetric key cryptography protocols to exchange information in the public!
- We also discussed two real-world examples of key exchange (Diffie-Hellman and RSA)
- In the lecture prior to that we introduced the notion of asymmetric encryption, outlined its basics. Today we will dive into its security!

# Public Key Encryption Overview



# Benefits

- 1 Public key encryption is often referred to as *assymmetric encryption* because the encryptor and decryptor use different keys (unlike *symmetric encryption* which uses the same keys)
- 2 After the public key is securely obtained, there is only one interaction to send a message!
- 3 We can re-use the public key many times
- 4 Anyone can post their public key for everyone else to see (no key exchange required), this means that the secret keys must not be derivable from public keys

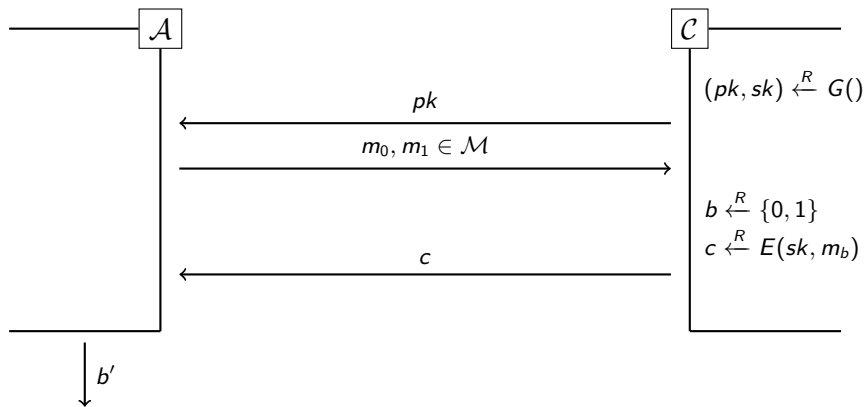
## 2 Semantic Security

# Review of Semantic Security

- The intuition of semantic security is that the probability a computationally bounded adversary can learn anything about a message from its ciphertext is negligible
- Semantic security guarantees that a message cannot be recovered from a ciphertext



# Public Key Semantic Security Attack Game



if  $b' = b$ , then the adversary wins

# Semantic Security Randomization

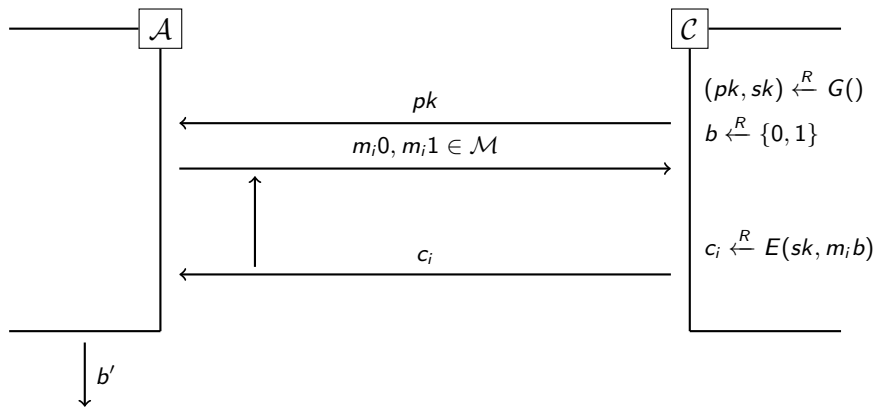
- For public key semantic security, the encryption function must be random. Can we think of an attack???
- Future Homework Assignment!

### 3 Chosen Plaintext Attack Security

# Public Key CPA Security

- Semantic security does not imply CPA security in symmetric key cryptography schemes
- The intuition behind this is that in a symmetric key security setting, the attacker cannot encrypt their own messages into their own cipher texts (because they don't have access to the key)
- In a public key setting, the adversary doesn't even need to interact with the challenger to get cipher texts

# Public Key CPA Attack Game



if  $b' = b$ , then the adversary wins

## 4 Chosen Ciphertext Attack Security

# CCA Motivation

- Imagine I collected homework by email and required that y'all encrypt your homework with my public key  $pk$
- What if another student intercepted that stream of data and strategically flipped bits until the submission read as "from: student1" to "from: student2"
- Therefore, we need our systems to be secure against adversaries generating valid cipher texts

# CCA Construction

$$E(pk, m) = [x \xleftarrow{R} \mathcal{X}, y \leftarrow F(pk, x), k \leftarrow H(x), c \xleftarrow{R} E_s(k, m)] \rightarrow (c, y)$$

$$D(pk, (y, c)) = [x \leftarrow I(sk, y), k \leftarrow H(x), m \leftarrow D_s(k, c)] \rightarrow m$$



## 5 Shamir Secret Sharing

# Motivation

- suppose you have a secret that you want to share with your friends
- but you only want your friends to unlock this secret if they are all together
- wouldn't it be cool if there was a mathematical way to ensure they can only unlock the secret if they are all together?

# Overview

- Enter polynomial interpolation
- if you have two points, you can find a unique line that intersects them, if you have 3 points, you can identify a unique parabola and so on
- you can encode your secret somewhere along a function, and then tell all of your friends a different point along the curve expressed by the function
- if enough of your friends get together, they can figure out the exact curve and find the secret!